

beepSeek-Prover-V1.5: Harnessing Proof Assistant Feedback for Reinforcement Learning and Monte-Carlo Tree Search

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DeepSeek-AI

https://github.com/deepseek-ai/DeepSeek-Prover-V1.5

Abstract

proving in Lean 4, which enhances DeepSeek-Prover-V1 by optimizing both training and inference processes. trained on DeepSeekMath-Base with specialization in formal mathematical languages, the model undergoes supervised fine-tuning using an enhanced formal theorem proving dataset derived from DeepSeek-Prover-V1. ther refinement is achieved through reinforcement learning from proof assistant feedback (RLPAF) ond the single-pass whole-proof generation approach of DeepSeek-Prover-V1, we propose RMaxTS, a variant of Monte-Carlo tree search that employs an intrinsic-reward-driven exploration strategy to generate diverse proof paths. PSeek-Prover-V1.5 demonstrates significant improvements over DeepSeek-Prover-V1, achieving new state-of-the-art results on the test set of the high school level miniF2F benchmark (63.5%) and the undergraduate level ProofNet benchmark (25.3%).

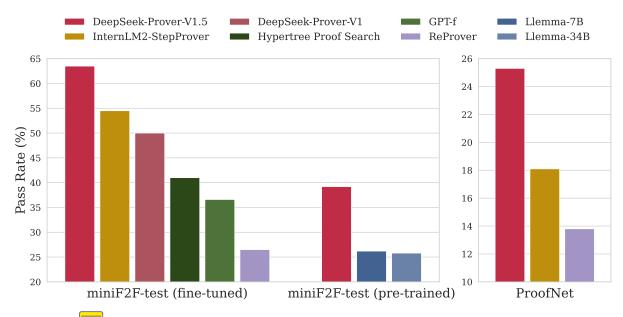


Figure 1 | strates of models on formal theorem proving benchmarks in Lean 4: the high school level miniF2F-test benchmark (Zheng et al., 2022) and the undergraduate level ProofNet benchmark (Azerbayev et al., 2023). compare both the pre-trained and fine-tuned versions of DeepSeek-Prover-V1.5 with strong baselines.

1. Introduction

reasoning and theorem proving in artificial intelligence. spite notable progress in natural language domains, language models still encounter substantial challenges in formal theorem proving, e.g. using Lean (Moura and Ullrich, 2021) and Isabelle (Paulson, 1994), which requires rigorous derivations satisfying formal specifications of the verification system. In advanced models like GPT-4 (OpenAI, 2023) struggle with complex formal proofs, underscoring the intricate nature of both the coding and the mathematics involved. Formal theorem proving model must not only grasp the syntax and semantics of formal systems like the Lean theorem prover but also align abstract mathematical reasoning with precise formal representation.

guage models in formal theorem proving typically employ two strategies: proof-step generation (Polu and Sutskever, 2020; Jiang et al., 2022; Lample et al., 2022; Yang et al., 2023; Wu et al., 2024) and whole-proof generation (Jiang et al., 2022; Zhao et al., 2023; Wang et al., 2023). of-step generation predicts each subsequent tactic and verifies it using the formal verifier to obtain updated information about the current tactic state, often utilizing tree search techniques to construct valid proofs. ontrast, whole-proof generation is computationally efficient, which produces an entire proof code based on the theorem statement, requiring less communication budget to coordinate between the prover model and the formal theorem verifier. **DeepSeek-Prover-V1** (Xin et al., 2024) has achieved state-of-the-art results in Lean 4 with whole-proof generation, this paradigm presents its unique challenges. equires long-horizon sequence prediction without access to intermediate tactic states, and future tactics depend on these hidden results. = lean's tactic mode, proofs are constructed through a sequence of tactics that transform the proof state. sequential nature introduces the risk of compounding errors (Ross et al., 2011), where a single misinterpretation can lead to significant deviations from a valid proof path. re specifically, the auto-regressive model may have incorrect believes on intermediate tactic states when generating long proofs.

eamlessly integrate intermediate tactic states in proof-step generation while maintaining the simplicity and computational efficiency of whole-proof generation, we have developed a unified approach in DeepSeek-Prover-V1.5. proof-step and whole-proof generation techniques through a truncate-and-resume mechanism. process begins with standard whole-proof generation, where the language model completes the proof code following the theorem statement prefix. Lean prover then verifies this code. = e proof is correct and complete, the procedure terminates. = n error is detected, the code is truncated at the first error message, and any subsequent code is discarded. generated proof code is then used as a prompt for the generation of next proof segment. enhance the accuracy of the model's new completions, we append the latest state from the Lean 4 prover as a comment at the end of the prompt. tably, our method is not restricted to resuming from the last successfully applied tactic. = integrate the truncate-and-resume mechanism into Monte-Carlo tree search (MCTS; Coulom, 2006) in which the truncation points are scheduled by the tree search policy. Eddition, we propose a novel reward-free exploration algorithm for MCTS to address the reward sparsity issue of proof search. = assign the tree search agent intrinsic motivation, a.k.a. curiosity (Schmidhuber, 2010), to extensively explore the tactic state space. se algorithmic modules extend the functionality of our whole-proof generation model to become a flexible tool for interactive theorem proving, which can effectively utilize the proof assistant feedback and generate diverse solution candidates.

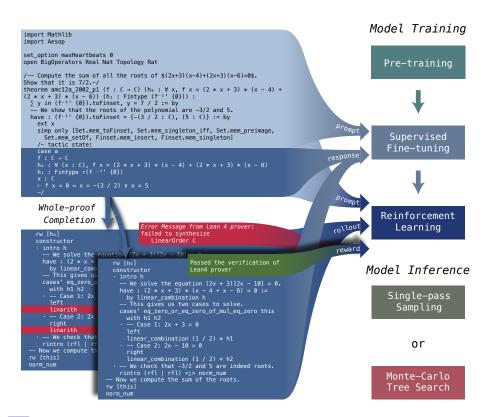


Figure 2 | Framework. DeepSeek-Prover-V1.5 is trained through pre-training, supervised fine-tuning, and reinforcement learning. Fring supervised fine-tuning, the pre-trained model receives an incomplete theorem proof ending with a tactic state comment keyword. From model is trained to predict the content of this tactic state (auxiliary objective) and complete the subsequent proof steps (main objective). The reinforcement learning stage, given an incomplete theorem proof and ground-truth tactic state from the Lean prover, we roll out the fine-tuned model to generate multiple proof candidates, which are then verified by the Lean prover. Verification results for these candidates are used as binary (0-1) rewards to further optimize the model and enhance its alignment with the formal specifications of the verification system. From model inference, we offer two alternatives: single-pass sampling and Monte-Carlo tree search.

Contributions

present a comprehensive framework for developing a language model-based formal mathematics prover, integrating several key components: large-scale mathematical pre-training, formal mathematics corpus construction and augmentation, online reinforcement learning from proof assistant feedback, and a tree search methodology for long-term planning in theorem proving. pre-trained model, supervised fine-tuned model, and reinforcement learning model, along with the code for the Monte-Carlo tree search algorithm, are publicly available for further research and application.

re-Training: We enhance our base model's capabilities in formal theorem proving and mathematical reasoning by further pre-training on high-quality mathematics and code data, with a focus on formal languages such as Lean, Isabelle, and Metamath.

upervised Fine-Tuning: We improve the Lean 4 code completion dataset by implement-

ing two data augmentation techniques. t, we use DeepSeek-Coder V2 236B (Zhu et al.,

2024) to annotate natural language chain-of-thought comments alongside Lean 4 code, aligning formal theorem proving with natural language reasoning. ond, we insert intermediate tactic state information within the Lean 4 proof code, enabling our model to leverage compiler feedback effectively. resulting dataset is then used to fine-tune the pre-trained model.

reinforcement Learning: We employ the GRPO algorithm (Shao et al., 2024) to perform reinforcement learning from proof assistant feedback (RLPAF) on the supervised fine-tuned model. fication results from the Lean prover serve as reward supervision, enhancing the model's alignment with the formal specifications of the verification system. onte-Carlo Tree Search: We advance the tree search method in formal theorem proving by introducing a novel abstraction and a corresponding search algorithm. retruncate-and-resume mechanism acts as a state-action abstraction, seamlessly integrating the tree search process into the whole-proof generation framework. present RMaxTS, an innovative Monte-Carlo tree search algorithm that leverages the RMax (Brafman and Tennenholtz, 2002) strategy to tackle exploration challenges in sparse-reward proof search problems. assigning intrinsic rewards, this algorithm encourages the prover agent to generate diverse planning paths, thereby fostering extensive exploration of the proof space.

Summary of Evaluations and Metrics

iniF2F: In the single-pass whole-proof generation setting, DeepSeek-Prover-V1.5 achieved a pass rate of 60.2% on the test set of miniF2F, marking a significant improvement of absolute 10.2 percentage points over DeepSeek-Prover-V1's 50.0%. orporating tree search techniques further elevated the pass rate to a new state-of-the-art 63.5%.

whole-proof generation setting for ProofNet, with pass rates of 21.6% on the validation set and 23.7% on the test set. integration of tree search techniques further enhanced these results, achieving new state-of-the-art pass rates of 25.4% on the validation set and 25.3% on the test set.

Model Training

Pre-training

mathematical language, we further pre-train our base model (Shao et al., 2024). s refinement involved training on high-quality datasets that include both code and natural language mathematical content. specifically focused on formal languages widely used in proof assistants, such as Lean, Isabelle, and Metamath. designate this improved model as DeepSeek-Prover-V1.5-Base.

Supervised Fine-tuning

is section, we explore the methodology and processes involved in the supervised fine-tuning (SFT) of DeepSeek-Prover-V1.5. cifically, we augment the proof dataset from DeepSeek-Prover-V1 by adding detailed explanatory comments. is enhancement aims to improve the alignment between natural language descriptions and Lean 4 code, thereby facilitating better formal mathematical reasoning. ditionally, we incorporate intermediate tactic state

information as an auxiliary prediction task to support the truncate-and-resume mechanism used in the Monte-Carlo Tree Search process. refer to the resulting model as DeepSeek-Prover-V1.5-SFT.

Curation. We develop a comprehensive Lean 4 code completion dataset for the supervised fine-tuning. is dataset includes synthetic proof code derived from a wide range of formal theorems. Ese theorems are sourced from various projects, such as the standard Lean 4 math library Mathlib4 (Mathlib Community, 2020), synthetic theorems from DeepSeek-Prover-V1 (Xin et al., 2024) and Lean Workbook (Ying et al., 2024), and validation sets from the miniF2F (Zheng et al., 2022) and ProofNet (Azerbayev et al., 2023) benchmarks. Laugment the formal proof data, we employed an expert iteration process (Polu and Sutskever, 2020). involves generating proofs using the language model, verifying the generated proof data, retraining the model with the verified data, and then using the optimized model to generate additional proof data. ween each iteration, we use DeepSeek-Coder V2 236B (Zhu et al., 2024) to annotate the thought process before the proof code as comments. ally, we tailor these data for the truncate-and-resume mechanism for Monte-Carlo Tree Search (details in Section 3.1). resulting proof dataset consists of 9,645k sequences.

Emught-augmented Proof Generation. DeepSeek-Prover-V1, we identified a significant gap between problem-solving strategies in natural language and theorem proving in Lean. natural language, models generate detailed deduction steps to construct proofs, whereas in Lean, they often rely on a sequence of high-level tactic calls to brute-force solutions. high-level tactics, while effective, obscure their internal workings and outcomes, hindering the model's ability to resolve complex proof goals with structured mathematical reasoning. address this issue, we develop an approach that incorporates natural language reasoning before generating theorem proof code. illar to Lean-STaR (Lin et al., 2024), which performs isolated chain-of-thought reasoning (Wei et al., 2022; Feng et al., 2023) before each proof step, our method integrates this reasoning directly as comments within the proof code. : use the DeepSeek-Coder V2 236B (Zhu et al., 2024) to enhance existing data in DeepSeek-Prover-V1 in two ways: 🢳, by inserting a complete natural language solution at the beginning of the proof block, and second, by alternately inserting specific natural language steps for corresponding Lean tactics ining the model with this data format enforces it to propose complete mathematical reasoning at the beginning of the proof block and detailed step planning before each tactic **=** is approach successfully develops new behaviors, employing delicate mathematical thinking to guide the generation of tactics. he training data, two distinct guiding prompts are used to differentiate between the CoT (Chain of Thought) mode and the non-CoT mode for proof code completion. =mples of input and output in both modes can be found in Appendix A.

mpt Augmentation with Tactic State Information. To implement the truncate-and-resume mechanism for Monte-Carlo Tree Search, we needed to extract tactic information from the code generated by the model. enhanced the Lean REPL (Read-Eval-Print Loop; Leanprover Community, 2023) with data extraction tools from the LeanDojo (Yang et al., 2023) project. allowed us to extract tactic information in triples, which include the position of each tactic, as well as the tactic states before and after its application. is information helps us identify the specific tactic code that triggers verification errors (used in the expansion step for tree search, see Section 3.2). each tactic in a generated valid formal proof, we insert the tactic state returned by the verifier as a comment "/- tactic state: ... -/". ring training, we use all tokens following "/- tactic state: " as responses to calculate the supervised fine-tuning loss, while the tokens before this comment is used as prompts and do not contribute to the training loss calculation.

rain for 9B tokens, using a batch size of 2,048 and a constant learning rate of 1e-4 e training process begins with 100 warm-up steps to stabilize the learning dynamics. ining examples are randomly concatenated to form sequences, with a maximum context length of 4,096 tokens.

Reinforcement Learning from Proof Assistant Feedback

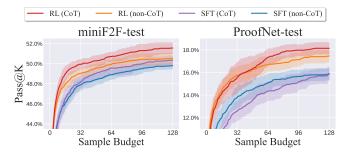
forcement learning (RL) has been proven effective in enhancing the mathematical reasoning capabilities of supervised fine-tuned language models (Shao et al., 2024). further advance DeepSeek-Prover-V1.5-SFT, we incorporate a reinforcement learning phase, resulting in the model DeepSeek-Prover-V1.5-RL. s phase leverages RL to enhance performance based on verification feedback from the Lean 4 prover. specifics of this RL process are detailed below.

In the reinforcement learning stage, we use a subset of theorem statements from the supervised fine-tuning dataset as training prompts. select theorems for which DeepSeek-Prover-V1.5-SFT has a moderate success rate in generating correct proofs upon multiple attempts. ensures that the model has room for improvement while still being able to receive positive feedback. er filtering, we retain approximately 4.5k unique theorem statements. ch theorem is prefixed with both CoT and non-CoT guiding prompts to enhance the model's proof generation capabilities in both modes.

wards. When training LLMs via RL, a trained reward model typically provides feedback signals. Ontrast, formal theorem proving benefits from the rigorous verification of generated proofs by proof assistants, offering a significant advantage. Cifically, each generated proof receives a reward of 1 if verified as correct, and 0 otherwise. Cifically, each generated proof receives a reward of 1 if verified as correct, and 0 otherwise. Cifically, each generated proof receives a reward of 1 if verified as correct, and 0 otherwise. Cifically, each generated proof receives a reward of 1 if verified as correct, and 0 otherwise. Cifically, each generated proof receives a reward of 1 if verified as correct, and 0 otherwise. Cifically, each generated proof receives a reward of 1 if verified as correct, and 0 otherwise. Cifically, each generated proof receives a reward of 1 if verified as correct, and 0 otherwise. Cifically, each generated proof receives a reward of 1 if verified as correct, and 0 otherwise. Cifically, each generated proof receives a reward of 1 if verified as correct, and 0 otherwise. Cifically, each generated proof receives a reward of 1 if verified as correct, and 0 otherwise. Cifically, each generated proof receives a reward of 1 if verified as correct, and 0 otherwise. Cifically, each generated proof receives a reward of 1 if verified as correct, and 0 otherwise. Cifically, each generated proof receives a reward of 1 if verified as correct, and 0 otherwise. Cifically, each generated proof receives a reward of 1 if verified as correct, and 0 otherwise. Cifically, each generated proof receives a reward of 1 if verified as correct, and 0 otherwise. Cifically, each generated proof receives a reward of 1 if verified as correct, and 0 otherwise. Cifically, each generated proof receives a reward of 1 if verified as correct, and 0 otherwise. Cifically, each generated proof receives a reward of 1 if verified as correct, and 0 otherwise. Cifically, each generated proof receives a reward of 1 if verified as correct

nforcement Learning Algorithm. We employ the Group Relative Policy Optimization (GRPO; Shao et al., 2024) as our RL algorithm, which has demonstrated superior effectiveness and efficiency compared to PPO (Schulman et al., 2017), primarily because it eliminates the necessity of training an additional critic model. cifically, GRPO samples a group of candidate proofs for each theorem prompt and optimizes the model based on the relative rewards of the outputs within the group. r prompt selection strategy is designed to likely include both correct and incorrect proofs among the candidates, aligning well with the group-relative nature of GRPO and thereby enhancing the training process.

ning Setting. We conduct RL training based on the SFT model, which serves as both the initial model and the reference model for imposing the Kullback-Leibler (KL) divergence penalty. use a constant learning rate of 5e-6, and the KL penalty coefficient is set to 0.02. each theorem, we sample a group of 32 candidate proofs, with maximum length set to 2,048. etraining batch size is configured to 512.



Model	Pass@128		
Model	miniF2F-test	ProofNet-test	
Base (3-shot)	$29.7\% \pm 0.5\%$	$9.7\% \pm 0.7\%$	
SFT (non-CoT) SFT (CoT)	$49.8\% \pm 0.3\%$ $50.4\% \pm 0.4\%$	$15.9\% \pm 0.5\%$ $15.9\% \pm 0.6\%$	
RL (non-CoT) RL (CoT)	$50.5\% \pm 0.6\%$ $51.6\% \pm 0.5\%$	$17.5\% \pm 0.5\%$ $18.2\% \pm 0.5\%$	

re 3 | Comparison of model capabilities at different training stages. T' and "non-CoT" refer to evaluations using two guiding prompts. shaded region represents the range of standard deviations around the mean values. notation $\mu \pm \sigma$ indicates the average accuracy μ and the standard deviation σ .

Evaluation

lchmarks. We evaluate theorem-proving performance on the following benchmarks to compare model capabilities after each training stage:

iniF2F (Zheng et al., 2022) focuses on formal problem-solving skills for high-school level exercises and competitions, such as AMC, AIME, and IMO, with an emphasis on algebra and number theory. benchmark includes 244 validation and 244 test problems, originally in Lean 3 and manually converted to Lean 4.9.0, based on the version provided by Yang (2023).

ProofNet (Azerbayev et al., 2023) evaluates formal theorem-proving capabilities at the undergraduate level in mathematics. omprises 185 validation and 186 test problems from widely-used undergraduate textbooks, covering real and complex analysis, linear algebra, abstract algebra, and topology. ese problems were initially in Lean 3 and manually converted to Lean 4.9.0.

pendently sample three proof demonstrations from the validation set to construct the few-shot prompts or the miniF2F benchmark, we use human-written proofs from Yang (2023), while for the ProofNet benchmark, we use correct proofs generated by DeepSeek-Prover-V1.5-RL as few-shot demonstrations. DeepSeek-Prover-V1.5-SFT and DeepSeek-Prover-V1.5-RL, we employ two types of guiding prompts: one that encourages chain-of-thought (CoT) reasoning before each proof step, and one that does not (non-CoT).

measures the model's success in generating a correct proof within *K* attempts. ch model is deployed on a single A100-40G GPU, utilizing the vLLM framework (Kwon et al., 2023) for sample generation. sampling parameters are set with a temperature of 1, a top-p value of 0.95, and a maximum token limit of 2,048. generated proofs are then verified using the Lean 4 theorem prover. this verification, we import Mathlib4 (Mathlib Community, 2020) and Aesop (Limperg and From, 2023) to access predefined premises and tactics. e verification process is subject to a time limit of 300 seconds.

parison across Training Stages. Figure 3 presents a comparative analysis of each training stage on the miniF2F and ProofNet datasets. It base model, DeepSeek-Prover-V1.5-Base, achieves a notable pass rate, solving nearly one-third of the problems on the test set of the miniF2F benchmark using 3-shot prompting. It supervised fine-tuning stage, resulting in DeepSeek-Prover-V1.5-SFT, significantly outperforms the base model, with Pass@128 accuracy increasing by approximately two-thirds on miniF2F and doubling on ProofNet. It subsequent reinforcement learning stage further enhances the model's performance, improving Pass@K accuracy across all values of K. It ontrast to findings in natural language mathematics, such as those reported in DeepSeekMath (Shao et al., 2024), where reinforcement learning primarily boosts the correct response from TopK, we observe a genuine enhancement of fundamental capabilities in formal theorem proving. It is improvement is evident not only with a small sample budget but also remains stable as the sample budget increases. It is conclusion is further supported by later Monte-Carlo Tree Search experiments with larger sample budgets, as discussed in Section 4.2.

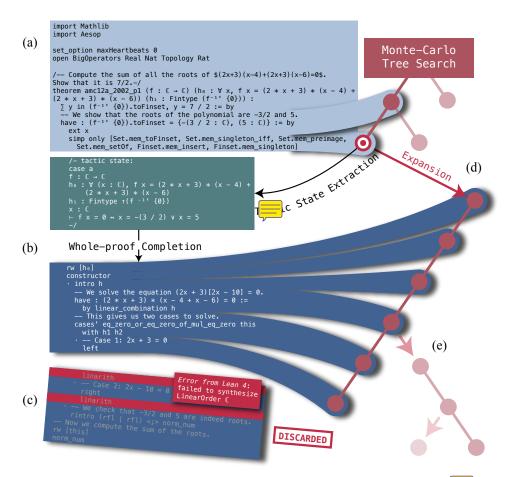
CoT generation modes for both DeepSeek-Prover-V1.5-SFT and DeepSeek-Prover-V1.5-RL. results in Figure 3 demonstrate that the CoT mode consistently outperforms the non-CoT mode across most settings. ecifically, DeepSeek-Prover-V1.5-RL, leveraging these enhanced theorem-proving patterns, achieves superior performance on both benchmarks, with an average accuracy of 51.6% on miniF2F and 18.2% on ProofNet. integration of natural language reasoning in CoT mode significantly enhances the planning and execution of formal proof writing. a detailed comparison of proof strategies with and without the use of natural language chain-of-thought, refer to the examples provided in Appendix A.

Exploration-oriented Monte-Carlo Tree Search

Tactic-level Tree Abstraction

mplement the tree search method in the whole-proof generation setting, we introduce a proof tree abstraction to define the tailored state and action space, leveraging a truncate-and-resume mechanism. Ighly following the paradigm of Yao et al. (2023), we begin by decomposing an incomplete proof into a sequence of tree nodes that correspond to individual proof steps, and then we utilize the partial content stored in these tree nodes to continue the proof generation process. If ure 4 illustrates the process of constructing a proof search tree from whole-proof generation.

then truncate the proof at the earliest verification error, ensuring that all subsequent tactic codes can be successfully applied to advance the proof towards the desired theorem. tactic codes are segmented into several code fractions, each containing a valid tactic code and its associated chain-of-thought comments, corresponding to a single tree edge that represents a tactic state transition. rough this abstraction, each tactic code is converted into a series of tree nodes, forming a path from the root to a specific node.



a node, we trace its corresponding incomplete proof code prefix, which includes the file header, initial statement, and successfully applied tactics from the ancestor nodes. The language model then generates the subsequent proof based on this prefix along with a comment block containing the current tactic state. The combined proof code (prefix and newly generated code) is verified by the Lean 4 prover. Performs are found, the tree-search procedure terminates. For are detected, we truncate the newly generated code at the first error message, discard the subsequent code, and parse the successful portion into tactics. Each tactic is added as a new node in the search tree, extending a chain of descendants beneath the selected node. Once the tree updates are complete, the next iteration of expansion begins by selecting an alternative candidate node, which is not limited to leaf nodes. s process repeats until a correct proof is found or the sample budget is exhausted.

tactic state, meaning each node in our proof tree can correspond to various tactic codes that achieve the same outcome. handle this, we store a set of these equivalent tactic codes at each node. he tree search agent expands a node, it randomly selects one tactic to use as a prompt for the language model. sprompt includes the incomplete proof code ending with the chosen tactic and the tactic state information from the Lean prover as a comment block. fine-tuned model (see Section 2.2) has been trained to recognize and utilize this format, using the incomplete code augmented with tactic state comments to guide subsequent proof generation.

Interactive Theorem Proving via Monte-Carlo Tree Search

proof search tree is developed using the standard Monte-Carlo Tree Search (MCTS) paradigm (MCTS; Coulom, 2006; Browne et al., 2012), which iteratively applies four steps: *Selection, Expansion, Simulation*, and *Backpropagation*. integrate the *Simulation* step into *Expansion* because our whole-proof generation model inherently performs a rollout from the expanded node. detailed design of the algorithm workflow is as follows.

ction. The selection step, *a.k.a.*the tree policy, starts from the root node and traverses downward to identify a promising node for expansion. objective of this algorithmic step is to trade off between exploration and exploitation (Kocsis and Szepesvári, 2006). tree policy at a tree node *s* is computed by selecting the action that maximizes the value from the set of valid operations:

$$TreePolicy(s) = \underset{a \in Children(s) \cup \{\emptyset\}}{\arg \max} Q_{UCB}(s, a), \tag{1}$$

the current node s, denoted by a special token $a = \emptyset$. sapproach uses a technique called *virtual node* (Wang et al., 2023), which assigns each node an imaginary child to represent the selection of the current node s for expansion. nables the tree search agent to continually expand non-leaf nodes, as the action space is supported by a generative model whose output scope cannot be determined by a fixed number of trails. value estimation $Q_{UCB}(s, a)$ of performing action a on node s is composed by two components:

$$\forall a \in Children(s) \cup \{\emptyset\}, \quad Q_{UCB}(s, a) = \underbrace{Q(s, a)}_{\text{litation}} + \underbrace{UCB(s, a)}_{\text{pration}},$$
 (2)

re Q(s,a) denotes a sample-based estimation of action values derived from the selection history, functioning as the exploitation component that retrieves high-value candidates from previous trials. B(s,a) denotes the exploration bonus computed by upper confidence bounds (UCB; Auer, 2002), which diminishes with the repeated execution of the state-action pair (s,a) re specifically, $Q_{UCB}(s,a)$ stands for an optimistic estimation of Q(s,a) and can serve as an upper bound with high probability. defer the discussion of detailed settings of node values and UCB bonus to Section 3.3.

The next step is invoking the proof generation model to expand the node nominated by the selection phase. Suming the incomplete proof codes stored on the node designated for expansion, we perform whole-proof generation to propose a series of subsequent tactics and submit the generated proof to Lean prover for verification. The a trial of proof completion is equivalent to conducting a single rollout of simulation within the standard MCTS framework. The next having found a new proof of the desired theorem. The rewise, we parse the verification feedback and truncate the generated proof to the assertion of the earliest verification error. The remaining tactics are transformed into a path of nodes to be merged into the search tree (see Figure 4). The important to note that, because we use the whole-proof generation setting—where the output is an entire proof consisting of a sequence of tactics, rather than just the next tactic—our expansion procedure may insert a path of tree nodes into the search tree during each iteration. So differs from the conventional MCTS designed for competitive games, which typically expands only one layer of children nodes per iteration (Silver et al., 2016, 2018; Schrittwieser et al., 2020).

Lackpropagation. The final phase of each tree search iteration is to update value statistics along the selection trajectory from the root to the expanded node, *i.e.*, updating the values associated with the tree policy stated in Eq. (1). $\tau = \{(root, s^{(1)}), (s^{(1)}, s^{(2)}), (s^{(2)}, s^{(3)}), \dots, (s^{(|\tau|-1)} = s_t, \emptyset)\}$ denote the selection trajectory of *t*-th iteration that ends with s_t as the expanding node update $Q_{UCB}(s,a)$ for all $(s,a) \in \tau$ by taking the most recent trajectory reward $R(\tau)$ into account (details refer to Eq. (7)). τ extrinsic source of rewards comes from the compiler feedback, specifically assigning a reward of τ extrinsic τ and τ for completed proofs and τ extrinsic τ of or unsolved ones. τ bection 3.3, we will introduce an intrinsic reward mechanism to augment the reward assignment that enhances the agent's incentive for exploration.

Intrinsic Rewards for Monte-Carlo Tree Search

he search problem of formal theorem proving, the extrinsic rewards are extremely sparse, *i.e.*, the search agent only obtains non-zero rewards when the proof is completely solved. For specifically, the proof search process forms a tree structure with only a narrow set of leaves delivering non-zero rewards, which matches a famous hard-exploration case (Krishnamurthy et al., 2016) in the literature of statistical reinforcement learning. Formote exploration in sparse-reward sequential decision making, one classical paradigm is constructing intrinsic rewards (Schmidhuber, 2010) that encourage the agent to not only optimize extrinsic rewards but also acquire general information about the interactive environment (Bellemare et al., 2016; Houthooft et al., 2016; Pathak et al., 2017; Burda et al., 2019). This section, we present our intrinsic-reward-driven exploration algorithm, *RMax applied to Tree Search* (RMaxTS), to incorporate reward-free exploration in the proof search problem.

ax applied to MCTS. We adopt RMax (Brafman and Tennenholtz, 2002), a classical exploration mechanism, to construct intrinsic rewards for Monte-Carlo tree search. core idea of RMax is to explore a broad coverage of the state space. agent awards itself a maximal amount of reward upon reaching an unseen state. the context of proof search, where no extrinsic rewards are provided until the proof is completed, our algorithmic procedure resembles ZeroRMax (Jin et al., 2020), in which the agent's exploration is driven solely by intrinsic rewards, i.e., setting $R(\tau) = R_{\text{intrinsic}}(\tau)$. in intrinsic reward of a tree expansion step is determined by whether a new node is added to the search tree,

$$R_{\text{intrinsic}}(\tau) = \mathbb{I} \left[\text{at least one new node is added to the search tree} \right],$$
 (3)

re τ denotes the most recent selection trajectory that requires a reward assignment for backpropagation. Es exploration strategy prioritizes the expansion of nodes where the prover model generates tactics that lead to a diverse range of tactic states. multiple Lean codes can result in the same transition of intermediate states, this heuristics can potentially reduce redundant generation and improve sample efficiency.

B for Non-stationary Rewards. The common setting of UCB exploration bonus for Monte-Carlo tree search is using UCB1 (Auer et al., 2002):

$$Q_{UCB1}(s,a) = \frac{W(s,a)}{N(s,a)} + \sqrt{\frac{2\ln\sum_{a'}N(s,a')}{N(s,a)}},$$
(4)

$$W(s,a) = \sum_{\tau \in \Gamma(s,a)} R(\tau), \tag{5}$$

$$N(s,a) = |\Gamma(s,a)|, \tag{6}$$

ere $\Gamma(s,a) = \{\tau \mid (s,a) \in \tau\}$ denotes the list of tree-policy trajectory τ containing (s,a) as an intermediate selection step. facilitate discussions, we organize the list $\Gamma(s, a) = \{\tau_1, \tau_2, \cdots\}$ such that newly collected trajectories have larger subscript indices. Ehis work, we propose to use an alternative variant of UCB method. = te that the derived intrinsic reward in Eq. (3) is a non-stationary reward signal whose expected value decays with the progress of exploration. t is because it becomes definitely harder to discover new nodes with unseen tactic states as the search tree expands through sophisticated exploration. tackle the non-stationarity, we consider discounted upper confidence bounds (DUCB; Garivier and Moulines, 2011), which uses a discount factor $\gamma \in (0,1)$ to smoothly drop those outdated feedback records:

$$Q_{DUCB}(s,a) = \frac{W_{\gamma}(s,a)}{N_{\gamma}(s,a)} + \sqrt{\frac{2\ln\sum_{a'}N_{\gamma}(s,a')}{N_{\gamma}(s,a)}},$$
(7)

$$W_{\gamma}(s,a) = \sum_{t=1}^{N(s,a)} \gamma^{N(s,a)-t} R(\tau_t),$$

$$N_{\gamma}(s,a) = \sum_{t=0}^{N(s,a)-1} \gamma^t,$$
(8)

$$N_{\gamma}(s,a) = \sum_{t=0}^{N(s,a)-1} \gamma^{t}, \tag{9}$$

==ere newly received feedback would be assigned a larger weight in the value estimation. practice, we set y = 0.99. Note that the role of discount factor y in DUCB differs from its role in value iteration for infinite-horizon MDPs. discounting is applied to tree search iterations rather than to the action-step horizon within a single trajectory.

Parallelization of Monte-Carlo Tree Search

🚃 nhance the efficiency of Monte-Carlo Tree Search (MCTS), we implement several established parallelization techniques as described by Chaslot et al. (2008).

cot Parallelization: We deploy 256 MCTS runners per node, with one language model per GPU and a batch size of 512 for proof generation. ee Lean prover is invoked through REPL and executed on a cluster with thousands of CPU cores, where each proof verification task is handled by an individual process, created and terminated in a sandbox. proof generation by language models and verification by Lean provers are handled asynchronously. setup allows MCTS runners to perform concurrent tree search operations, significantly accelerating the process. \equiv

ree Parallelization: We manage each search tree with 32 thread workers to parallelize the tree iteration steps. s method effectively schedules and balances the tasks of proof generation and Lean verification. h thread worker iteratively performs the tree search loop by selecting a candidate node for expansion, invoking the language model to generate the proof, verifying the generated proof with the Lean prover, and performing backpropagation.

tual Loss: To encourage diverse node selection among concurrent thread workers, we assign a virtual reward $R(\tau) = 0$ for ongoing iterations. \equiv s involves backpropagating a reward of 0 temporarily and updating it to the true reward upon completion. strategy promotes exploration of different nodes for expansion, thereby enhancing the overall search efficiency.

Comparison with Existing Methods

his section, we compare our proposed proof tree search method, which introduces a novel truncate-and-resume mechanism for whole-proof generation, with existing approaches. = rent methods for using language models in formal mathematics proof search generally fall into two main strategies:

Tulti-pass proof-step generation: This strategy breaks down the proving process into multiple episodes of tactic generation and verification, typically following a tree search pattern. Involves generating and verifying one tactic at a time, repeating the process for the next tactic until no proof goals remain. In able examples include GPT-f (Polu and Sutskever, 2020; Polu et al., 2022), Thor (Jiang et al., 2022), ReProver (Yang et al., 2023), Hypertree Proof Search (Lample et al., 2022), and InternLM2-StepProver (Wu et al., 2024). In gle-pass whole-proof generation: This approach generates and verify an entire proof in one attempt. In the proof is incorrect, the model generates a new proof in the next attempt. It hods in this category include DSP (Jiang et al., 2022), Subgoal-Prover Zhao et al. (2023), LEGO-Prover (Wang et al., 2023), Lyra (Zheng et al., 2023), and miniCTX (Hu et al., 2024).

proof tree search method uniquely bridges these two strategies, offering a novel hybrid approach. Parts with whole-proof generation, similar to the single-pass approach, but extends this by implementing a sophisticated truncate-and-resume mechanism. Is sprocess involves truncating the generated proof to its successful initial segment, parsing this segment into individual tactics, and resuming the tree search from this point. Is iterative process effectively implements a Monte-Carlo Tree Search, seamlessly integrating single-pass whole-proof generation with multi-pass proof-step generation. In sequently, we can train a single model with nearly identical objectives to support both strategies simultaneously. It experimental results demonstrate that this unified approach achieves superior performance in both settings. In combining the strengths of existing methods and introducing innovative techniques, our method offers a more versatile and effective solution for formal mathematics proof search, potentially paving the way for future advancements in this field.

Experimental Results

his section, we evaluate the theorem-proving capabilities of DeepSeek-Prover-V1.5 using two distinct benchmarks: miniF2F, which encompasses high-school level exercises and competition problems, and ProofNet, which pertains to undergraduate-level theorems. present the results for both complete proof generation and Monte-Carlo tree search methodologies, utilizing the same trained model and inference configuration as Section 2.4 to ensure consistency.

Main Results

resent a comparative analysis of DeepSeek-Prover-V1.5 against previous state-of-the-art language models, highlighting its performance and advancements.

eneral-purpose Models: GPT-3.5 and GPT-4 (OpenAI, 2023) are advanced generative AI models developed by OpenAI, known for their effectiveness across diverse tasks, including code generation. pite not being specifically designed for theorem proving, their extensive parameter scales provide significant capabilities. evaluation of these models in formal theorem proving is facilitated by COPRA (Thakur et al., 2023), an in-context learning agent that leverages these large language models to propose tactic applications. ditionally, we examine Llemma (Azerbayev et al., 2024), a series of

language models trained on extensive general mathematical corpora, commonly used as the base model for formal theorem proving.

et al., 2022) represents an initial effort to apply Transformers (Vaswani et al., 2017) to proof-step generation for theorem proving tasks, utilizing a best-first search module to construct complete proofs. sequent advancements include **ReProver** (Yang et al., 2023), **LLMStep** (Welleck and Saha, 2023), and **Lean-STaR** (Lin et al., 2024) pertree **Proof Search** (Lample et al., 2022) explores the use of Monte Carlo tree search in formal theorem proving using Lean. current works, **InternLM2-Math** (Ying et al., 2024) and **InternLM2-StepProver** (Wu et al., 2024), also demonstrate outstanding performance.

ric. We compare the performance of DeepSeek-Prover-V1.5 with state-of-the-art models using the pass@*K* accuracy metric, which evaluates the model's ability to generate a correct proof within *K* attempts. display the sample budget *K* according to the following rules to align the computation budget across different generation schemes.

of proofs generated, with large values of *K* factorized for the ease of comparison to tree search methods.

or best-first-search methods, following the notation of Azerbayev et al. (2024), we present $K = N \times S \times T$ where N denotes the number of best-first-search attempts, S denotes the number of tactics generated for each expansion, and T denotes the number of expansion iterations.

pr tree search methods, *e.g.*, RMaxTS and HTPS (Lample et al., 2022), we present $K = N \times T$ where N denotes the number of tree search attempts, and T denotes the number of model generations invoked in tree expansions.

lts on miniF2F. Table 1 provides a comparative analysis of various theorem-proving methods on the miniF2F-test dataset. he single-pass whole-proof generation setting, DeepSeek-Prover-V1.5-RL achieved the highest pass rate at 60.2%, marking a significant improvement of 10.2 percentage points over DeepSeek-Prover-V1's 50.0%. h a sampling budget limited to 128 attempts, DeepSeek-Prover-V1.5-RL proved 51.6% of the problems, significantly outperforming other whole-proof generation methods and is comparable to the leading tree search methods. the Tree Search Methods category, DeepSeek-Prover-V1.5-RL + RMaxTS leads with a pass rate of 62.7%, establishing a new state-of-the-art and creating a substantial gap with existing methods. btably, DeepSeek-Prover-V1.5-RL requires only 3200 whole-proof generation samplings to achieve a pass rate of 54.9%, surpassing the previous state-of-the-art result of InternLM2-StepProver, which performs 64 × 3200 tree searches to achieve 54.5%.

methods on the ProofNet dataset. epSeek-Prover-V1.5-RL achieved pass rates of 22.6% and 25.3% for the overall ProofNet dataset in the single-pass whole-proof generation setting and with the enhancement of RMaxTS, respectively. ese results surpass the existing state-of-the-art methods, ReProver (13.8%) and InternLM2-StepProver (18.1%). en the number of whole-proof generation attempts is restricted to 3200, DeepSeek-Prover-V1.5 also proved 21.7% of the theorems, demonstrating a 3.6% improvement over the previous state-of-the-art, InternLM2-StepProver.

Method	Sample budget	miniF2F-test
Single-pass Whole-Proof Generation Methods		
TheoremLlama [49]	128	33.6%
DaniCarla Burana VI [52]	128	$46.1\% \pm 0.5\%$
DeepSeek-Prover-V1 [53]	16×4096	50.0%
	128	$29.7\% \pm 0.5\%$
DeepSeek-Prover-V1.5-Base	3200	39.2%
•	6400	42.2%
	32	$48.2\% \pm 0.6\%$
	64	$49.6\% \pm 0.7\%$
Deer Cool, Proven V1 F CET	128	$50.4\% \pm 0.4\%$
DeepSeek-Prover-V1.5-SFT	3200	$53.3\% \pm 0.5\%$
	4×6400	$55.8\% \pm 0.7\%$
	16×6400	57.4%
	32	$50.0\% \pm 0.5\%$
	64	$50.7\% \pm 0.4\%$
D C I D MIEDI	128	$51.6\% \pm 0.5\%$
DeepSeek-Prover-V1.5-RL	3200	$54.9\% \pm 0.7\%$
	4×6400	$58.4\% \pm 0.6\%$
	16×6400	60.2%
Tree Search Methods		
COPRA (Code Llama) [45]	1×500	5.7%
COPRA (GPT-3.5) [45]	1×60	9.0%
COPRA (GPT-4) [45]	1×60	26.6%
Llemma-7B [5]	$1 \times 32 \times 100$	26.2%
Llemma-34B [5]	$1 \times 32 \times 100$	25.8%
ReProver [55]	-	26.5%
LLMStep [51]	$1 \times 32 \times 100$	27.9%
GPT-f [35]	$64 \times 8 \times 512$	36.6%
Hypertree Proof Search [23]	64×5000	41.0%
Lean-STaR [26]	$64 \times 1 \times 50$	46.3%
InternLM2-Math-7B [58]	$1 \times 32 \times 100$	30.3%
InternLM2-Math-Plus-7B [58]	$1 \times 32 \times 100$	43.4%
InternI M2 Ston Prover [52]	$1 \times 32 \times 100$	48.8%
InternLM2-StepProver [52]	$64 \times 32 \times 100$	54.5%
	1×3200	$53.5\% \pm 0.4\%$
Door Cook Drover V1 5 CET + DMTC	4×6400	$56.3\% \pm 0.3\%$
DeepSeek-Prover-V1.5-SFT + RMaxTS	16×6400	59.0%
	$32 \times 6400^{\dagger}$	60.2%
	1×3200	$55.0\% \pm 0.7\%$
Deep Cook Drover VI E DI + DMTC	4×6400	$59.6\% \pm 0.6\%$
DeepSeek-Prover-V1.5-RL + RMaxTS	16×6400	62.7 %
	$32 \times 6400^{\dagger}$	63.5%

le 1 | Comparison with state-of-the-art methods on the miniF2F-test dataset. In notation $\mu \pm \sigma$ denotes the average accuracy μ and the standard deviation σ . It less otherwise specified, DeepSeek-Prover-V1.5-Base results are based on 3-shot prompting, while DeepSeek-Prover-V1.5-SFT and RL employ CoT mode prompting. In symbol † indicates performance using a mixture strategy with two guiding prompts (see Section 4.2 for details).

Method	Sample budget	valid [‡]	ProofNet test	all
Single-pass Whole-Proof Generation Methods				
DeepSeek-Prover-V1.5-Base	128 3200	6.6% ± 0.9% 10.8%	9.7% ± 0.7% 15.6%	$7.5\% \pm 0.7\%$ 13.2%
DeepSeek-Prover-V1.5-SFT	128 3200 4×6400	19.9% ± 0.4% 20.7% ± 0.7% 22.2%	15.9% ± 0.6% 21.0% ± 0.9% 23.7%	17.9% ± 0.3% 20.9% ± 0.6% 22.9%
DeepSeek-Prover-V1.5-RL	128 3200 4×6400	20.1% ± 0.5% 21.4% ± 0.3% 21.6%	18.2% ± 0.5% 22.0% ± 0.5% 23.7%	19.1% ± 0.4% 21.7% ± 0.4% 22.6%
Tree Search Methods				
ReProver [55] InternLM2-StepProver [52]	1×32×100	-	-	13.8% 18.1%
DeepSeek-Prover-V1.5-SFT + RMaxTS	1 × 3200 4 × 6400	22.2% ± 0.7% 23.8%	21.6% ± 0.2% 25.8%	21.9% ± 0.4% 24.8%
DeepSeek-Prover-V1.5-RL + RMaxTS	1 × 3200 4 × 6400	22.0% ± 0.3% 25.4%	21.5% ± 0.8% 25.3%	21.8% ± 0.4% 25.3%

e 2 | Comparing with state-of-the-arts on the ProofNet dataset. ‡ = e that the validation set of ProofNet is used to perform expert iteration in supervised fine-tuning.

4.2. Re-Examining the Effectiveness of Training Strategies on Large-scale Sampling

We revisit the effects of several training modules in n a large-scale sampling setting, focusing on both single-pass whole-proof generation and Monte-Carlo tree search. The results demonstration that the observations and findings presented in Section 2.4 generalize to sampling scenarios with a large sample size.

General Enhancement of Reinforcement Learning. To support the claim that online reinforcement learning from verification feedback generally enhances the model capabilities, we compare our final model to the SFT-only version using a large sample budget. The comparison results are presented as two columns in Table 3. DeepSeek-Prover-V1.5-RL consistently outperforms the SFT model across all generation settings, regardless of whether the chain-of-thought strategy is applied. The results also indicate that the improvements gained from conducting online RL is orthogonal to those achieved through RMaxTS, which can be further combined to boost the performance. By integrating both CoT prompting and RMaxTS, DeepSeek-Prover-V1.5-RL achieves a pass rate of 62.7% on miniF2F-test. This performance shows a notable 3.7% improvement over the SFT model, highlighting the critical role of reinforcement learning in enhancing the overall effectiveness of the proof completion model.

CoT, non-CoT, and Mixture Strategy. We compare the performance of two generation modes, *i.e.*, non-CoT and CoT, on miniF2F-test dataset. The results, shown in Table 3, indicate that the advantage of CoT over the non-CoT mode is amplified as the sample budget increases. This suggests that the incorporation of natural language chain-of-thought can diversify the planning pathways of theorem proving, potentially leading to a broader range of reasoning strategies and more innovative solutions. Results also show that these two modes have complementary

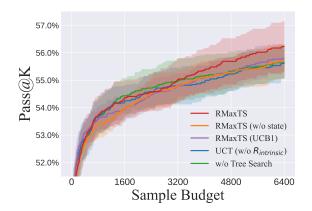
	Prompt mode	Sample budget	DeepSeek-Prover-V1.5	
			SFT	RL
Single-Pass Generation -	non-CoT	4×6400	$54.7\% \pm 0.4\%$	$56.5\% \pm 0.5\%$
		16×6400	56.1%	57.4%
	СоТ	4×6400	$55.8\% \pm 0.7\%$	$58.4\% \pm 0.5\%$
		16×6400	57.4%	60.2%
	non-CoT & CoT	$(2+2) \times 6400$	$56.1\% \pm 0.8\%$	$58.3\% \pm 0.6\%$
		$(8+8) \times 6400$	58.2%	60.7%
		$(16 + 16) \times 6400$	58.6%	61.1%
RMaxTS	non-CoT	4×6400	$55.7\% \pm 0.6\%$	$58.4\% \pm 0.6\%$
		16×6400	57.8%	59.4%
	СоТ	4×6400	$56.3\% \pm 0.3\%$	$59.6\% \pm 0.6\%$
		16×6400	59.0%	62.7%
	non-CoT & CoT	$(2+2) \times 6400$	$56.1\% \pm 0.8\%$	$60.0\% \pm 0.8\%$
		$(8+8)\times6400$	59.0%	63.1%
		$(16 + 16) \times 6400$	60.2%	63.5%

Table 3 | A large-scale ablation study to investigate the effectiveness of several algorithmic designs on model training. The results are evaluated on the miniF2F-test dataset.

advantages across different problems. The model's theorem proving strategy in the CoT mode is more systematic and proactive in mathematical thinking, while in the non-CoT mode, the model can efficiently use Lean high-level tactics to solve computational problems that can be addressed within Lean's automation mechanisms. To leverage these advantages, we consider a mixture strategy, denoted by non-CoT & CoT in Table 3, allocates half of sample budget to the CoT mode and the remains to the non-CoT mode. This simple combination of two guiding prompts shows great promise in further bootstrapping the performance of our proof completion model, achieving a pass rate of 63.5% on miniF2F-test. In Appendix B, we present example problems that illustrate the different advantages of the two generation modes.

4.3. Ablation Studies on RMaxTS

Intrinsic Rewards and Discounted UCB. We investigate the effectiveness of two core components of RMaxTS, i.e., the intrinsic rewards defined in Eq. (3) and the discounted upper confidence bound stated in Eq. (7). We start with a baseline implementing the standard UCT algorithm (Kocsis and Szepesvári, 2006) without intrinsic rewards, in which the exploration is driven exclusively by the UCB bonus. Note that, since no non-zero rewards are provided for this baseline, all variants of the UCB formula become equivalent, as node selection is determined solely by visitation counts. The experimental results in Figure 5 show that, in the absence of intrinsic rewards, the performance of UCT (without $R_{intrinsic}$) degenerates into a level comparable to that of non-search methods. Furthermore, we consider RMaxTS using the standard UCB1 (refer to Eq. (4)) instead of the discounted UCB, denoted by RMaxTS (DUCB \rightarrow UCB1). The results indicate that the performance of RMaxTS with UCB1 bonus is also moderate, comparable to that of UCT (without R_{intrinsic}). That is because UCB1 is designed to guarantee asymptotic performance through exhausted exploration (Auer et al., 2002) assuming the sample size to be sufficiently large. In contrast, the discounted UCB can accelerate the value propagation of non-stationary intrinsic rewards, preventing the guidance of R_{intrinsic} from being dominated by that of visitation counts. This demonstrates that the discounted UCB mechanism is a crucial complement to intrinsic-reward-driven exploration.



Sample budget	miniF2F-test
4×6400 16×6400	58.4% ± 0.5% 60.2%
$\begin{array}{c} 4 \times 6400 \\ 16 \times 6400 \end{array}$	$58.2\% \pm 0.3\%$ 61.1%
$\begin{array}{c} 4 \times 6400 \\ 16 \times 6400 \end{array}$	58.6% ± 0.3% 60.7%
4×6400 16×6400	58.4% ± 0.3% 61.1%
4×6400 16×6400	59.6% ± 0.6% 62.7%
	4 × 6400 16 × 6400 4 × 6400 16 × 6400 4 × 6400 4 × 6400 16 × 6400 4 × 6400 4 × 6400

Figure 5 | A modular ablation study examining the algorithmic design of RMaxTS. The experiments are conducted on the miniF2F-test dataset with DeepSeek-Prover-V1.5-RL using the CoT mode. The left panel presents the curves of Pass@K accuracy within 6400 generation samples. The results with a larger sample size are presented in the right panel.

Guidance of Tactic State Information. When expanding a tree node, we concatenate the intermediate tactic state information as a comment block to the incomplete code to guide the proof completion. With the provided auxiliary information, the proof completion model can enhance its internal representation of the tactic state, offering intermediate guidance for long-horizon planning. To demonstrate this advantage, we present experiments on RMaxTS that performs code completion directly from the raw incomplete code without accessing tactic state information, denoted by RMaxTS (without tactic state) in Figure 5. The results indicate that the performance gain from applying tree search becomes moderate in the absence of tactic state information, especially when tackling hard problems that require a large amount of samples. This highlights that the integration of compiler information is an essential component of the tree search algorithm, enhancing its overall effectiveness and sample efficiency.

Conclusion, Limitation, and Future Work

present DeepSeek-Prover-V1.5, a language model with 7 billion parameters that outperforms all open-source models in formal theorem proving in Lean 4. ppSeek-Prover-V1.5 is initialized with DeepSeek-Prover-V1.5-Base, which extends the pre-training of DeepSeekMath-Base 7B using a specialized corpus for formal mathematical reasoning. pervised fine-tuning is conducted on a comprehensive Lean 4 code completion dataset, encompassing a wide range of formal theorems from various mathematical domains, beequently, we employ GRPO to enhance whole-proof generation through online reinforcement learning. In developing the DeepSeek-Prover-V1.5 model, we introduce RMaxTS, a variant of Monte-Carlo tree search, to improve problem-solving capabilities via large-scale search with extensive exploration. pse components form a comprehensive pipeline for training an LLM-based proof assistant, enabling DeepSeek-Prover-V1.5 to achieve significant improvements over DeepSeek-Prover-V1.

framework of DeepSeek-Prover-V1.5 is designed to establish an AlphaZero-like pipeline for formal theorem proving. — use of expert iteration and synthetic data mirrors the core trial-and-error loop of reinforcement learning, with the compiler oracle serving as the world model to provide environmental supervision. — thin the RL paradigm, the integrated tree search module has proven to be highly effective in advancing superhuman performance across various domains (Silver et al., 2016; Fawzi et al., 2022; Lutz et al., 2023). — eveloping DeepSeek-

Prover-V1.5, we focus on the exploration aspect of RL, introducing RMaxTS to diversify the generation of proof steps. wever, the exploitation aspect, specifically the problem of proof search, remains unexplored. romising future direction is training a critic model to assess incomplete proofs and prune search branches. ha partial-proof critic model would implicitly perform temporal credit assignment (Sutton, 1984), decomposing proof-level feedback into stepwise value differences (Arjona-Medina et al., 2019). reloping critic models for assessing long planning paths and providing guidance rewards presents a crucial and challenging problem (Ng and Russell, 2000; Sorg et al., 2010) that warrants further investigation.

real-world theory proving within complex, multi-theorem Lean files (Hu et al., 2024). shift is a natural extension of our whole-proof generation approach. robservations indicate that the current model already possesses some understanding of file-level context. ving forward, we will focus on enhancing this aspect to support cutting-edge Lean mathematical formalization developers with our language model advancements.

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Illustrative Examples of Non-CoT and CoT Prompting for Proof Completion

LaTeX expressions are utilized. Firstly, here is an example of whole-proof generation in non-CoT prompting mode:

Input:

```
Complete the following Lean 4 code: 
 '''lean4 
 import Mathlib 
 import Aesop 
 set_option maxHeartbeats 0 
 open BigOperators Real Nat Topology Rat 
 /-- Show that \frac{9x^2\sin^2x+4}{x\sin x} \geqslant 12 for 0 < x < \pi.-/ 
 theorem aime_1983_p9 (x : \mathbb{R}) (h<sub>0</sub> : 0 < x \land x < \text{Real.pi}) : 
 12 \leqslant (9 * (x ^2 * \text{Real.sin } x ^2) + 4) / (x * Real.sin x) := by
```

Output:

```
rw [le_div_iff (mul_pos h<sub>0</sub>.1 (Real.sin_pos_of_pos_of_lt_pi h<sub>0</sub>.1 h<sub>0</sub>.2))]
nlinarith [mul_self_nonneg (3 * x * Real.sin x - 2)]
'''
```

t, here is an example of whole-proof generation in CoT prompting mode. sproof was found using a Monte-Carlo tree search with a truncate-and-resume mechanism (see Section 3 for details). model receives a prompt including a proof code in progress, ending with a comment block containing the intermediate tactic state from the Lean 4 prover. subsequent code is generated with guidance from this annotated code prefix.

Input:

```
Plete the following Lean 4 code with explanatory comments preceding each line of code:

""lean4 import Mathlib import Aesop

set_option maxHeartbeats 0

open BigOperators Real Nat Topology Rat

|-- Show that \frac{9x^2 \sin^2 x + 4}{x \sin x} \geq 12 for 0 < x < \pi. -/
theorem aime_1983_p9 (x : \mathbb{R}) (h_0 : 0 < x \lambda x < Real.pi) :

12 \leq (9 * (x ^ 2 * Real.sin x ^ 2) + 4) / (x * Real.sin x) := by
|-- To find the minimum value of \frac{9x^2 \sin^2 x + 4}{x \sin x} for 0 < x < \pi, we need to show that it is at least 12. We start by noting that the expression can be rewritten using the division property of inequalities. We then use the fact that \sin x and x are positive in the given range to establish the necessary
```

```
inequalities. Finally, we apply these results to conclude that the minimum
 value is indeed 12.
-- We start by ensuring that the product x * \sin x is positive in the given
have h_1 : 0 < x * Real.sin x := by
  apply mul_pos
  -- x is positive in the range (0,\pi).
  exact h_0.1
  -- \sin x is positive in the range (0,\pi).
  exact Real.sin_pos_of_pos_of_lt_pi h<sub>0</sub>.1 h<sub>0</sub>.2
-- Using the division property of inequalities, we rewrite the expression.
rw [le_div_iff h<sub>1</sub>]
/- tactic state:
  x : \mathbb{R}
  h_0: 0 < x \wedge x < \pi
  h_1 : 0 < x * x.sin
  + 12 * (x * x.sin) \leq 9 * (x ^ 2 * x.sin ^ 2) + 4
```

Output:

```
-- This is equivalent to showing that 9x^2\sin^2x - 12x\sin x + 4 \ge 0, and the left-hand side can be rewritten as a perfect square (3x\sin x - 2)^2.

-- We use the fact that (3x\sin x - 2)^2 is non-negative to establish this. nlinarith [sq_nonneg (3 * x * Real.sin x - 2)]
```

In this case, the proof completion model first extracts the goal expression from the tactic state, *i.e.*, the inequality $12(x\sin x) \le 9(x^2\sin^2 x) + 4$ at the last line of the comment block, and then arranges it into the form $9x^2\sin^2 x - 12x\sin x + 4 \ge 0$ as part of the chain-of-thought process. Finally, the left-hand side is rewritten as a square term to establish non-negativity, closing the proof.

B. Example Solutions to MiniF2F-test Problems

To illustrate the different advantages and complementary of non-CoT mode and CoT mode of prompting, we provide examples where non-CoT mode succeeds while CoT mode fails, and vice versa.

B.1. Problems that non-CoT outperforms CoT

```
theorem induction_pord1p1on2powklt5on2 (n : \mathbb{N}) (h<sub>0</sub> : 0 < n) : (\Pi k in Finset.Icc 1 n, 1 + (1 : \mathbb{R}) / 2 \hat{\ } k) < 5 / 2 := by rw [Finset.Icc] simp_all [Nat.succ_le_iff, Nat.one_le_iff_ne_zero] have h<sub>1</sub> : 0 < 2 \hat{\ } k := by apply pow_pos <;> norm_num norm_num have h<sub>2</sub> : (2 : \mathbb{R}) \neq 0 := by norm_num field_simp rw [div_lt_div_iff] ring_nf norm_cast
```

```
nlinarith
all_goals norm_cast
all_goals linarith
```

```
theorem imo_1960_p2 (x : \mathbb{R}) (h<sub>0</sub> : 0 \leq 1 + 2 * x) (h<sub>1</sub> : (1 - Real.sqrt (1 + 2 *
               x)) ^2 = 0)
              (h_2 : 4 * x ^2 / (1 - Real.sqrt (1 + 2 * x)) ^2 < 2 * x + 9) : -(1 / 2)
             \leq x \wedge x < 45 / 8 := by
      norm_num at h<sub>0</sub> h<sub>1</sub> h<sub>2</sub>
      have h_3: 0 \le 1 + 2 * x := by linarith
      have h_4 : 0 < 1 + Real.sqrt (1 + 2 * x) := by
              nlinarith [Real.sqrt_nonneg (1 + 2 * x)]
      have h_5: 4 * x ^ 2 / (1 - Real.sqrt (1 + 2 * x)) ^ 2 < 2 * x + 9 := by
             linarith
      have h_6: 1 - Real.sqrt (1 + 2 * x) \neq 0 := by
              intro h
              apply h_1
              nlinarith
      have h_7: 4 * x ^ 2 / (1 - Real.sqrt (1 + 2 * x)) ^ 2 = (1 + Real.sqrt (1 + 2 * x)) ^ 2 = (1 + Real.sqrt (1 + 2 * x)) ^ 3 = (1 + Real.sqrt (1 + 2 * x)) ^ 3 = (1 + Real.sqrt (1 + 2 * x)) ^ 3 = (1 + Real.sqrt (1 + 2 * x)) ^ 3 = (1 + Real.sqrt (1 + 2 * x)) ^ 3 = (1 + Real.sqrt (1 + 2 * x)) ^ 3 = (1 + Real.sqrt (1 + 2 * x)) ^ 3 = (1 + Real.sqrt (1 + 2 * x)) ^ 3 = (1 + Real.sqrt (1 + 2 * x)) ^ 3 = (1 + Real.sqrt (1 + 2 * x)) ^ 3 = (1 + Real.sqrt (1 + 2 * x)) ^ 3 = (1 + Real.sqrt (1 + 2 * x)) ^ 3 = (1 + Real.sqrt (1 + 2 * x)) ^ 3 = (1 + Real.sqrt (1 + 2 * x)) ^ 3 = (1 + Real.sqrt (1 + 2 * x)) ^ 3 = (1 + Real.sqrt (1 + 2 * x)) ^ 3 = (1 + Real.sqrt (1 + 2 * x)) ^ 3 = (1 + Real.sqrt (1 + 2 * x)) ^ 3 = (1 + Real.sqrt (1 + 2 * x)) ^ 3 = (1 + Real.sqrt (1 + 2 * x)) ^ 3 = (1 + Real.sqrt (1 + 2 * x)) ^ 3 = (1 + Real.sqrt (1 + 2 * x)) ^ 3 = (1 + Real.sqrt (1 + 2 * x)) ^ 3 = (1 + Real.sqrt (1 + 2 * x)) ^ 3 = (1 + Real.sqrt (1 + 2 * x)) ^ 3 = (1 + Real.sqrt (1 + 2 * x)) ^ 3 = (1 + Real.sqrt (1 + 2 * x)) ^ 3 = (1 + Real.sqrt (1 + 2 * x)) ^ 3 = (1 + Real.sqrt (1 + 2 * x)) ^ 3 = (1 + Real.sqrt (1 + 2 * x)) ^ 3 = (1 + Real.sqrt (1 + 2 * x)) ^ 3 = (1 + Real.sqrt (1 + 2 * x)) ^ 3 = (1 + Real.sqrt (1 + 2 * x)) ^ 3 = (1 + Real.sqrt (1 + 2 * x)) ^ 3 = (1 + Real.sqrt (1 + 2 * x)) ^ 3 = (1 + Real.sqrt (1 + 2 * x)) ^ 3 = (1 + Real.sqrt (1 + 2 * x)) ^ 3 = (1 + Real.sqrt (1 + 2 * x)) ^ 3 = (1 + Real.sqrt (1 + 2 * x)) ^ 3 = (1 + Real.sqrt (1 + 2 * x)) ^ 3 = (1 + Real.sqrt (1 + 2 * x)) ^ 3 = (1 + Real.sqrt (1 + 2 * x)) ^ 3 = (1 + Real.sqrt (1 + 2 * x)) ^ 3 = (1 + Real.sqrt (1 + 2 * x)) ^ 3 = (1 + Real.sqrt (1 + 2 * x)) ^ 3 = (1 + Real.sqrt (1 + 2 * x)) ^ 3 = (1 + Real.sqrt (1 + 2 * x)) ^ 3 = (1 + Real.sqrt (1 + 2 * x)) ^ 3 = (1 + Real.sqrt (1 + 2 * x)) ^ 3 = (1 + Real.sqrt (1 + 2 * x)) ^ 3 = (1 + Real.sqrt (1 + 2 * x)) ^ 3 = (1 + Real.sqrt (1 + 2 * x)) ^ 3 = (1 + Real.sqrt (1 + 2 * x)) ^ 3 = (1 + Real.sqrt (1 + 2 * x)) ^ 3 = (1 + Real.sqrt (1 + 2 * x)) ^ 3 = (1 + Real.sqrt (1 + 2 * x)) ^ 3 = (1 + Real.sqrt (1 + 2 * x)) ^ 3 = (1 + Real.sqrt (
            2 * x)) ^ 2 := by
              field simp [h<sub>6</sub>]
              nlinarith [sq_sqrt (show 0 \le 1 + 2 * x by linarith)]
      rw [h_7] at h_5
      constructor <;> nlinarith [sq_sqrt (show 0 ≤ 1 + 2 * x by linarith)]
```

B.2. Problems that CoT outperforms non-CoT

```
theorem mathd_algebra_459 (a b c d : Q) (h_0 : 3 * a = b + c + d) (h_1 : 4 * b =
    a + c + d
    (h_2 : 2 * c = a + b + d) (h_3 : 8 * a + 10 * b + 6 * c = 24) : \uparrow d.den +
   d.num = 28 := by
 Given the system of equations involving the times that Anna (a), Bertram
   (b), Carli (c), and David (d) can hold their breaths:
 1. 3a = b + c + d
 2. 4b = a + c + d
 3. 2c = a + b + d
 4. 8a + 10b + 6c = 24
 We need to find the sum of the numerator and the denominator of the
   fraction representing David's time in minutes, simplified to its lowest
   terms, and show that it is 28.
 First, we solve the system of equations step by step:
 - From the first equation, we can express d in terms of a,b, and c:
    d = 3a - b - c
  - Substitute d into the second equation:
    4b = a + c + (3a - b - c)
    Simplifying, we get:
    4b = 4a - b
    5b = 4a
    b = \frac{4a}{5}
 - Substitute b into the third equation:
    2c = a + \frac{4a}{5} + (3a - b - c)
    Simplifying, we get:
    2c = a + \frac{4a}{5} + 3a - \frac{4a}{5} - c
```

```
2c = 4a - c
  3c = 4a
  c = \frac{4a}{3}
- Substitute b and c into the fourth equation:
  8a + 10\left(\frac{4a}{5}\right) + 6\left(\frac{4a}{3}\right) = 24
  Simplifying, we get:
  8a + 8a + 8a = 24
  24a = 24
  a = 1
Now, we find b and c:
-b = \frac{4a}{5} = \frac{4}{5}
-c = \frac{4a}{3} = \frac{4}{3}
Finally, we find d:
-d = 3a - b - c = 3 \cdot 1 - \frac{4}{5} - \frac{4}{3} = 3 - \frac{4}{5} - \frac{4}{3} = \frac{45}{15} - \frac{12}{15} - \frac{20}{15} = \frac{13}{15}
The fraction \frac{13}{15} is already in its simplest form, and the sum of the
 numerator and the denominator is:
13 + 15 = 28
-/
have h_4 : a = 1 := by
  have h_5: 8 * a + 10 * b + 6 * c = 24 := h_3
  have h_6: 24 * a = 24 := by
     -- Solve the system of equations to find the value of a
     linarith
  exact mul_right_cancel<sub>0</sub> (show (24 : \mathbb{Q}) \neq 0 by norm_num)
     (by linarith)
have h_5 : b = 4 / 5 := by
   -- Solve for b using the value of a
  linarith
have h_6: c = 4 / 3 := by
   -- Solve for c using the value of a
  linarith
have h_7: d = 13 / 15 := by
  -- Solve for d using the values of a, b, and c
-- Sum the numerator and the denominator of d
norm_num [h<sub>7</sub>]
<;> rfl
```

```
theorem numbertheory_x5neqy2p4 (x y : \mathbb{Z}) : x ^ 5 \neq y ^ 2 + 4 := by /-

To show that for any two integers x and y, x^5 \neq y^2 + 4, we can use modular arithmetic to demonstrate that the equation x^5 \equiv y^2 + 4 \pmod{11} has no solutions. Specifically, we will consider the possible values of x and y modulo 11 and show that none of these combinations satisfy the equation.

-/

-- Use modular arithmetic to show that the equation has no solutions modulo 11.

have : \neg(x ^ 5 % 11 = (y ^ 2 + 4) % 11) := by

-- Consider all possible values of x modulo 11.

have : x % 11 = 0 v x % 11 = 1 v x % 11 = 2 v x % 11 = 3 v x % 11 = 4 v x % 11 = 5 v

x % 11 = 6 v x % 11 = 7 v x % 11 = 8 v x % 11 = 9 v x % 11 = 10 := by omega

-- Consider all possible values of y modulo 11.

have : y % 11 = 0 v y % 11 = 1 v y % 11 = 2 v y % 11 = 3 v y % 11 = 4 v y
```

```
theorem amc12_2000_p12 (a m c : \mathbb{N}) (h<sub>0</sub> : a + m + c = 12) :
    a * m * c + a * m + m * c + a * c \leq 112 := by
  Given nonnegative integers A, M, and C such that A+M+C=12, we need to find
   the maximum value of A \cdot M \cdot C + A \cdot M + M \cdot C + A \cdot C.
  To solve this, we can use the fact that the expression A \cdot M \cdot C + A \cdot M + M \cdot C + A \cdot C
    is a polynomial in terms of A, M, and C. We can try different values of A, M,
    and C that sum to 12 and compute the value of the polynomial for each
   combination. The maximum value of the polynomial will then be the solution.
  -- We use a lemma that states if the sum of three numbers is less than or
   equal to a certain value, then the product of these numbers is also less
    than or equal to a certain value.
  have h_1: a \leq 12 \land m \leq 12 \land c \leq 12 := by
    constructor
    \cdot -- Since a + m + c = 12, a cannot be greater than 12.
      omega
    · constructor
      · -- Similarly, m cannot be greater than 12.
        omega
      · -- And c cannot be greater than 12.
  -- Using the lemma, we compute the value of the polynomial for different
   values of a, m, and c that sum to 12.
  cases' h1 with h1_left h1_right
  cases' h<sub>1</sub>_right with h<sub>1</sub>_middle h<sub>1</sub>_right
  -- We use interval_cases to try different values of a, m, and c.
  interval_cases a <;> interval_cases m <;> interval_cases c <;>
    -- For each combination, we compute the value of the polynomial and check
   if it is less than or equal to 112.
    simp_all (config := {decide := true})
```