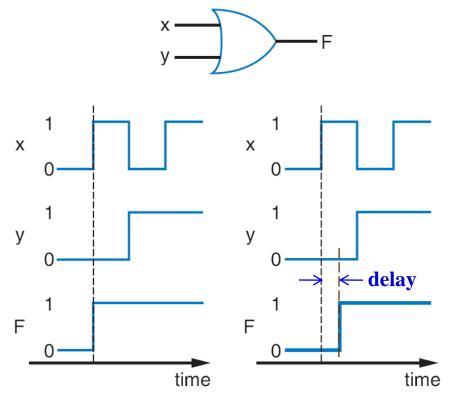
Topic 6 Latches and Flip Flops

Introduction

- Beginning from this lecture, we will:
 - Design a new type of building blocks, latch & flip-flop, that store value of a bit, a sequential circuit
 - Combine the blocks to build multi-bit storage a register
 - Describe the sequential behavior of a digital circuit using a finite state machine (FSM)
 - Convert a finite state machine to a controller a sequential circuit having a register and combinational logic

Reality of Combinational Circuit

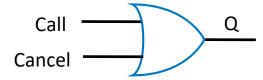
Non-Ideal Gate Behavior -- Delay



- Real gates have some delay
 - Outputs don't change immediately after inputs change

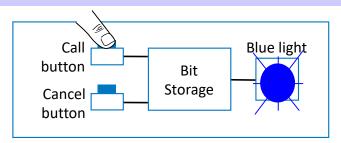
Example of Needing Bit Storage

- Flight attendant call button
 - Button pressed: provides a "1"
 - Press call: light turns on
 - Stays on after button released
 - Press cancel: light turns off
 - Logic gate circuit to implement this?

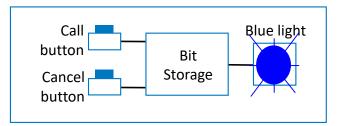


Doesn't work. Q=1 when Call=1, but doesn't stay 1 when Call returns to 0

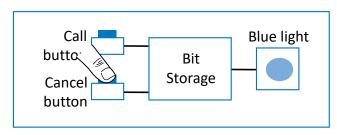
Need some form of "storage" in the circuit



1. Call button pressed – light turns on



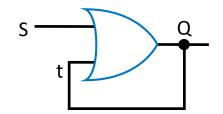
2. Call button released – light stays on

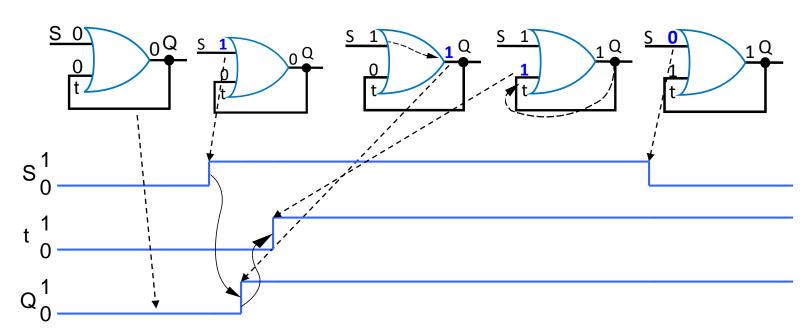


3. Cancel button pressed – light turns off

First Attempt at Implementation of Bit Storage

- We need some sort of feedback
 - Does circuit on the right do what we want?
 - Partially: Once Q becomes 1 (when S=1), Q stays 1 forever – no value of S can bring Q back to 0

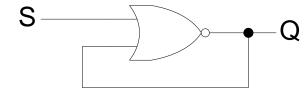




Concepts of Sequential Circuit

Sequential circuit

Combinational circuit with feedbacks



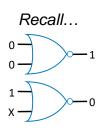
- Due to the feedback, output of a sequential circuit is decided by
 - Present inputs, and
 - Past input sequence and
 - Past outputs sequence

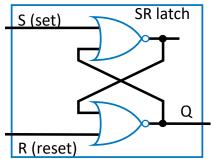
Timing concepts

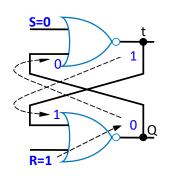
- input-output propagation delay
- clock
- Other timing issues

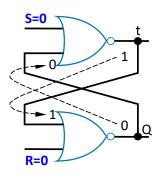
Second Attempt at Bit Storage – SR Latch

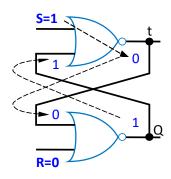
- Cross-coupled NOR gates
 - S: set (or preset) to 1
 - R: reset (or clear) to 0

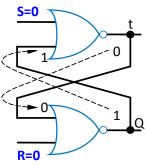


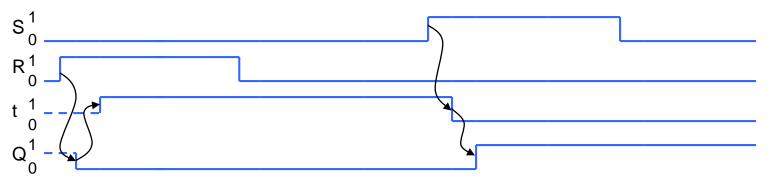






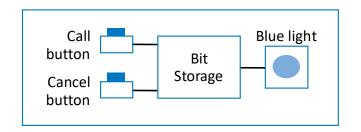


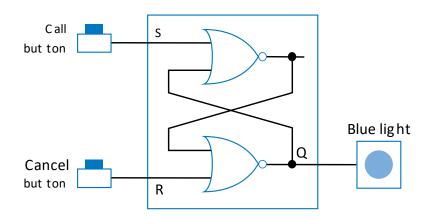




Example Using SR Latch for Bit Storage

- SR latch can serve as a bit storage, for example:
 - Call=1 : sets Q to 1
 - Q stays 1 even after Call=0
 - Cancel=1: resets Q to 0
- But, there's a problem...

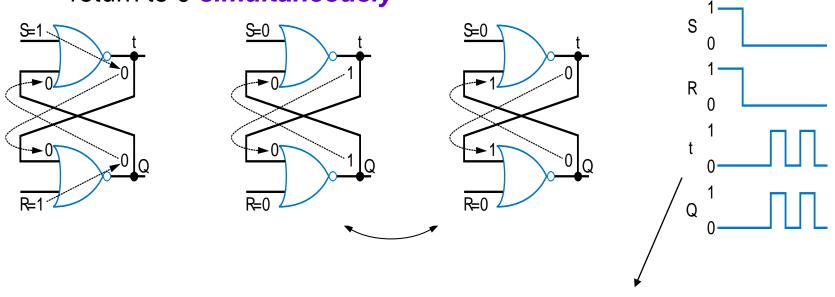




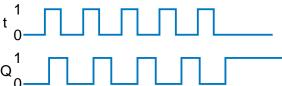
Problem with SR Latch

Problem

 If S=1 and R=1, we don't know what value Q will take when they both return to 0 simultaneously



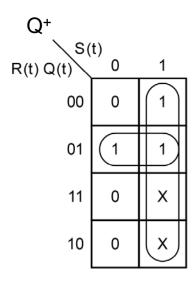
Q may oscillate. Then, because one path will be t slightly longer than the other, Q will eventually settle to 1 or 0 – but we don't know which.



Representation of SR Latch

- When discussing latches and flip-flops, we use
 - present state to represent current value of the Q output
 - next state to represent the new value of Q output responding to the current inputs and feedback of current output
- Characteristic table

S(t)	R(t)	Q(t)	Q(t	<u>+</u> Δ) — Q+	
0	0	0	0	hold	
0	0	1	1	Hold	
0	1	0	0	reset	
0	1	1	0	10301	
1	0	0	1	set	
1	0	1	1	301	
1	1	0	Х	not allowed	
1	1	1	Х	not anowed	

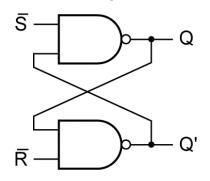


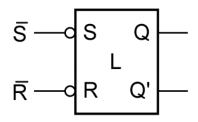
Characteristic equation

$$Q^+ = S + R'Q$$

Alternative Implementation of SR Latch

The cross-coupled RS latch can be implemented using NAND gates



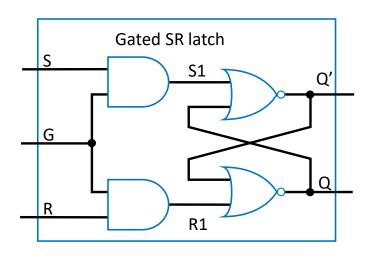


Characteristic table

	S	R	Q	Q ⁺	
•	0	0	0	Х	not allowed
	0	0	1	X	not allowed
	0	1	0	1	set
	0	1	1	11	301
	1	0	0	0	reset
	1	0	1	0	10001
	1	1	0	0	hold
	1	1	1	1	

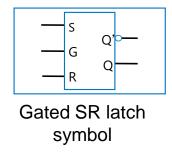
Gated SR Latch

• SR latch is enabled by a gate control signal G



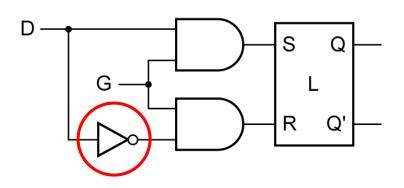
Characteristic Table

G	S	R	Q ⁺
0	X	X	Q; Latch locked
1	0	0	Q; Hold state
1	0	1	0; Reset state
1	1	0	1; Set state
1	1	1	not allowed



Solution to SR Latch Restriction - Gated D Latch

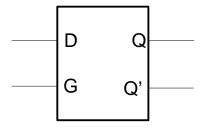
Solution to the unstable state problem caused by S = R = 1 in SR latch



Characteristic rable			
G D		Q ⁺	
1	0	0	
1	1	1	
Ο	Y	\circ	

Characteristic Table

 The input value is stored into the latch only when gate control G has high level – Level Sensitive

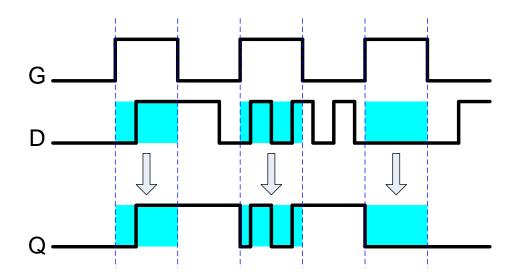


D latch symbol

Gated D Latch - Transparent Latch

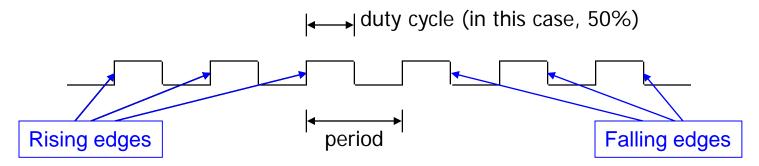
Properties of the D latch

- D latch is used as a temporary storage for a bit
- The binary information at the data input of the D latch is copied to the Q output when the control input G is high (or enabled)
- The output Q follows changes on the data input D as long as the control input G is enabled, so called a **transparent** latch



A Typical Control Input - Clock Signal

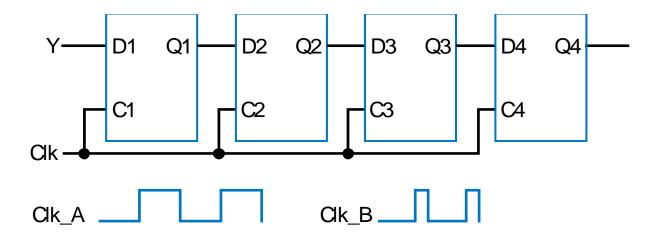
 Periodic pulse train used in sequential circuit to synchronize circuit behaviors



- Clock period: time interval between pulses
- Clock cycle: one such time interval
- Clock frequency: 1/period

Freq	Period		
100 GHz	0.01 ns		
10 GHz	0.1 ns		
1 GHz	1 ns		
100 MHz	10 ns		
10 MHz	100 ns		

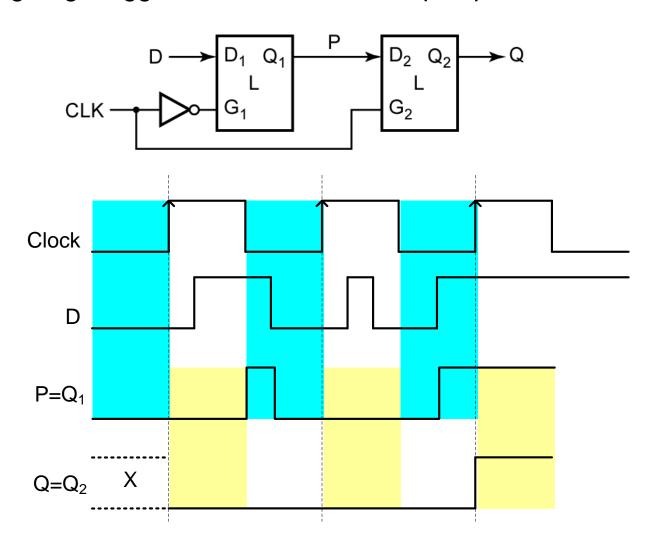
Problem with Level-Sensitive D Latch



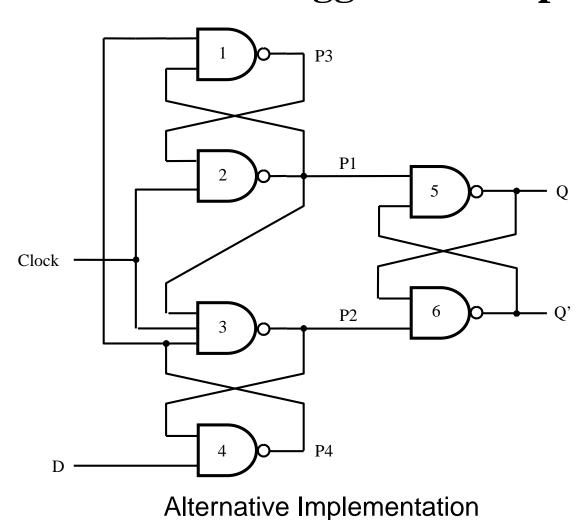
- D latch still has problem
 - When Clk=1, through how many latches will a signal travel?
 - Depends on for how long Clk=1
 - Clk_A -- signal may travel through multiple latches
 - Clk_B -- signal may travel through fewer latches
 - Hard to control the storage elements

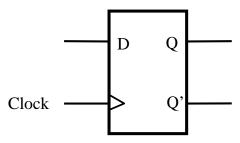
Rising-Edge Triggered D Flip Flop

Rising-edge triggered Master-Slave D flip flop



Alternative Implementation of Rising-Edge Triggered D Flip Flop



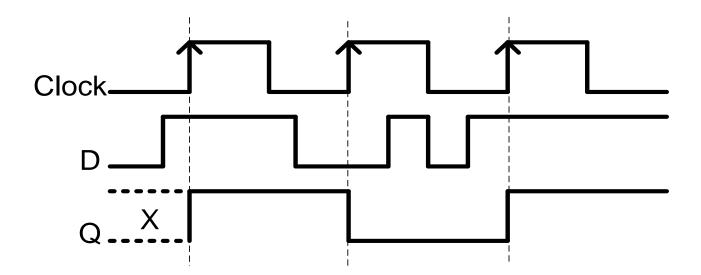


clock	D	Q ⁺
	0	0
	1	1
0	Χ	Q
1	Χ	Q

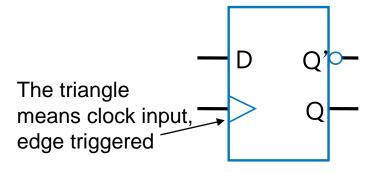
Characteristic equation: Q+ = D (at active clock edges)

Rising-Edge Triggered D Flip Flop

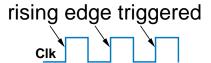
- Properties of the rising edge triggered D Flip Flop
 - The output changes only at the rising edges of the clock signal –
 Edge Sensitive
 - The output Q gets the value of input D at the time point of rising edge of clock

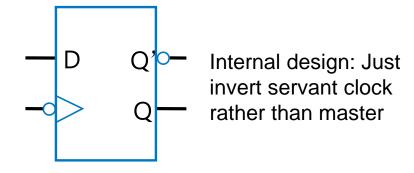


Symbols for D Flip-Flop

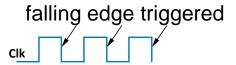


Symbol for rising-edge triggered D flip-flop



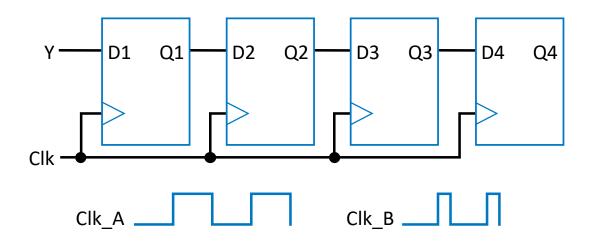


Symbol for falling-edge triggered D flip-flop



Application of D Flip-Flop

- Solves problem of concatenated D latches when G=1
 - In figure below, signal travels through exactly one flip-flop, for either Clk_A or Clk_B
 - On each rising edge of Clk, all four flip-flops are loaded simultaneously, doesn't matter how long Clk is 1.



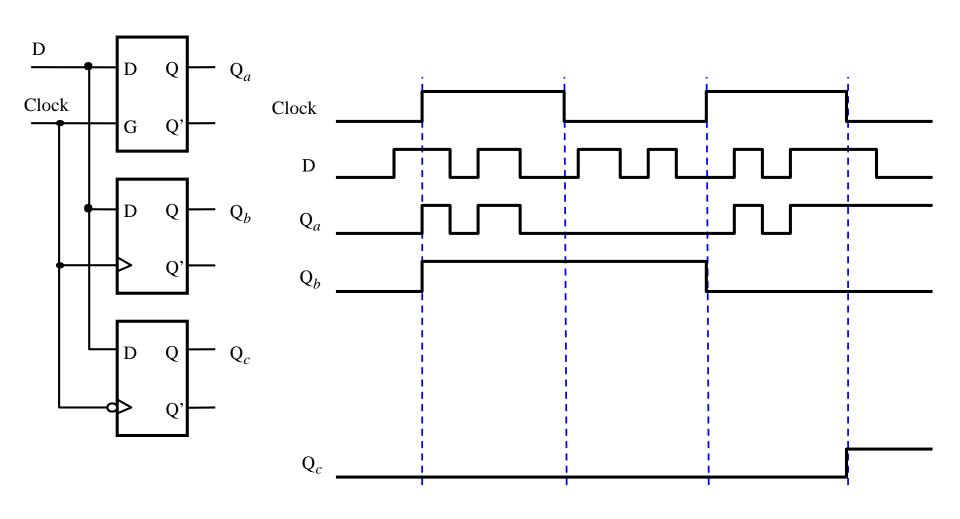
Flip Flop vs. Latch

- Both are storage elements in sequential circuits
- Flip flop
 - edge-sensitive, the input matters only at active edges (rising or falling)
 - behaviors are synchronous to the clock signal

Latch

- level-sensitive, the input matters whenever control has active level (high or low)
- behaviors are asynchronous to the clock signal

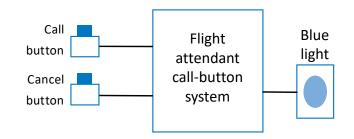
Flip Flop vs. Latch



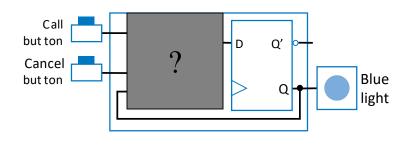
Flight-Attendant Call Button Using D Flip-Flop

- D flip-flop will store bit
- Inputs are Call, Cancel, and present output Q of D flip-flop
- Truth table

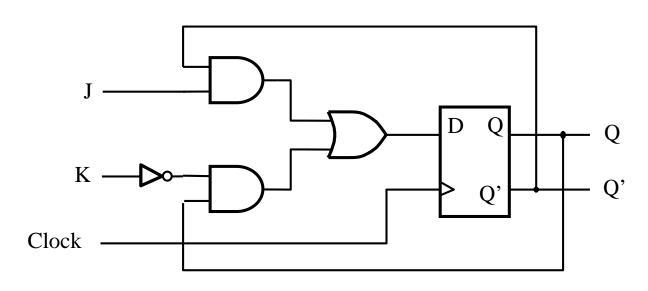
Call	Cancel	Q	Q+ = D
0	0	0	0
0	0	1	1
0	1	0	0
0	1	1	0
1	0	0	1
1	0	1	1
1	1	0	1
1	1	1	1

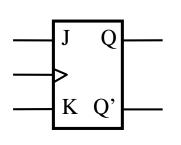


Circuit derived from truth table, using Chapter 2 combinational logic design process



Rising Edge-triggered J-K Flip Flop



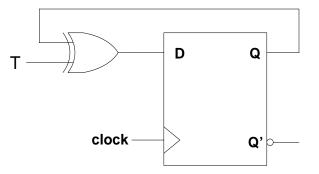


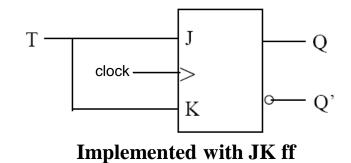
J	K	Q ⁺
0	0	Q
0	1	0
1	0	1
1	1	Q'

Characteristic equation:

$$Q^+ = JQ' + K'Q$$

Rising Edge-triggered T Flip Flop



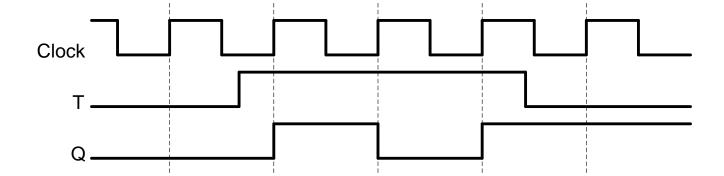


Implemented with D ff

	Т	Q	
Clock	 >	Q'	

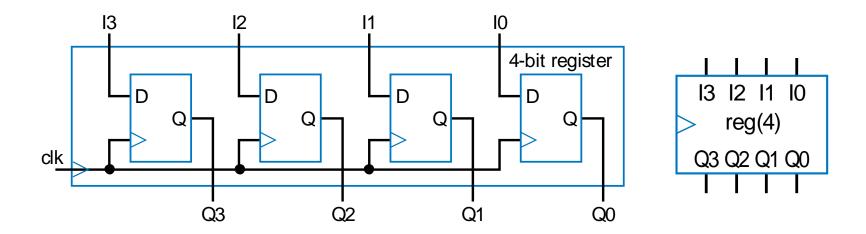
clock	Т	Q ⁺
	0	Q
	1	Q'

Characteristic equation: $Q^+ = T'Q + TQ' = T \oplus Q$



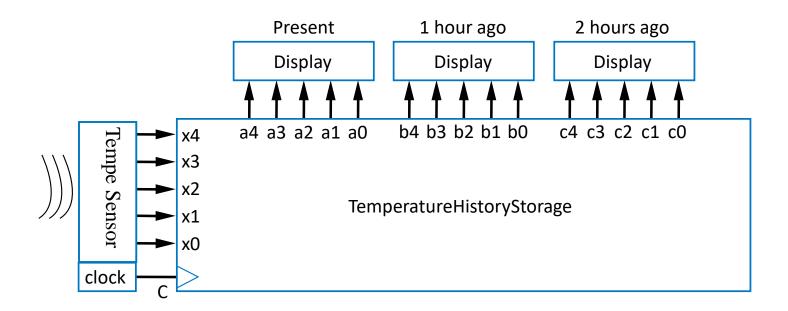
Basic Register

- Typically, we store multi-bit items
 - e.g., storing a 4-bit binary number
- Register: multiple flip-flops sharing clock signal



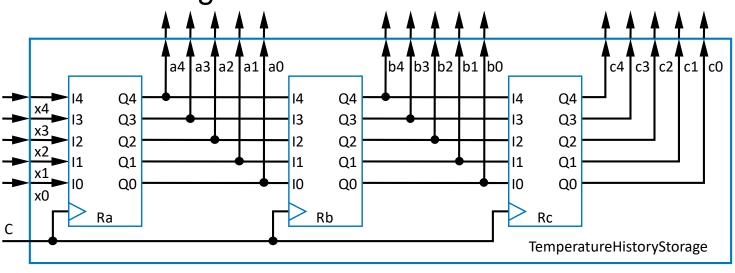
Example Using Registers: Temperature Display

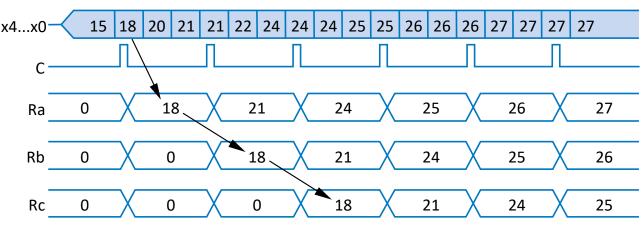
- Temperature history display
 - Sensor outputs temperature as 5-bit binary number
 - Timer pulses C every hour
 - Record temperature on each pulse, display last three recorded values



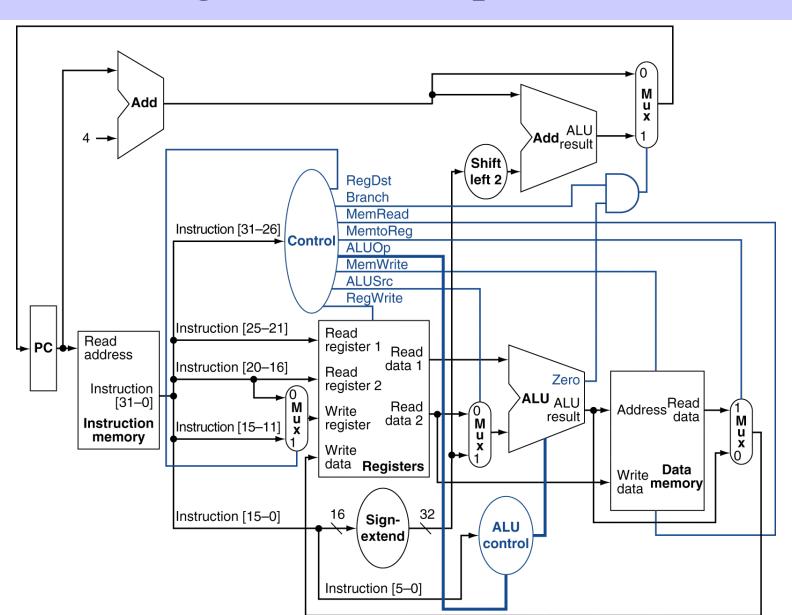
Example Using Registers: Temperature Display

• Use three 5-bit registers



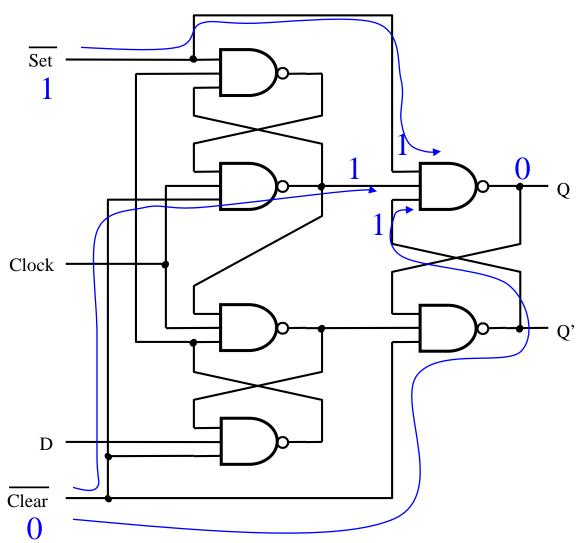


Big Picture – Simplified CPU



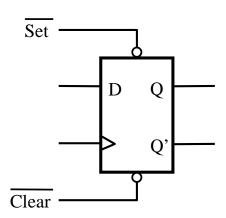
Implementation of Asynchronous Control Input

Control signals decides the output value directly independent of the clock signal



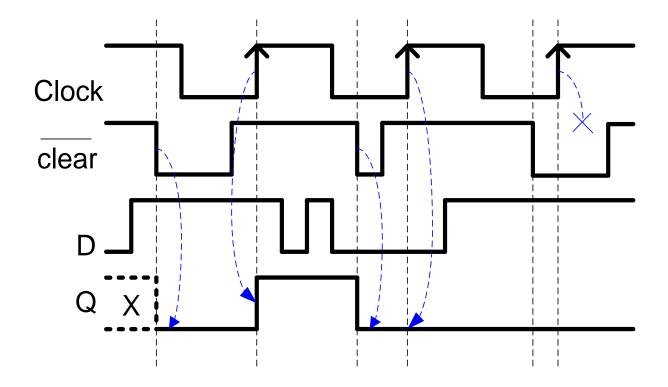
Control Inputs for Flip Flops

- Asynchronous:
 - control signals do not depend on the clock signal
- Synchronous:
 - control signals depend on the clock signal
- Active low:
 - It controls when it's low
- Active high:
 - It controls when it's high



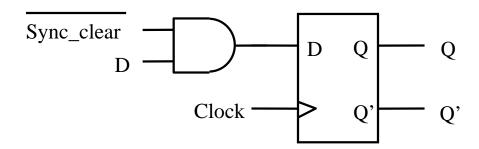
Asynchronous Control Input

D flip flop with active low asynchronous Clear



Implementation of Synchronous Control Input

- Synchronous Clear
 - control signal depends on the active edge (either rising or falling) of the clock signal



Flip-Flops with Control Inputs

D flip flop with active low synchronous Clear

