

# Datenbanken

## Blatt 4

Gruppe 26

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# 1 Aufgabe

## 1.1

### 1.1.1

```

1  {[a.Firstname, a.Lastname, d.Firstname, d.Lastname] | a, d ∈ person
2      ∧ ∃ x ∈ directs (d.PersonId = x.PersonId
3          ∧ ∃ y ∈ acts (a.PersonId = y.PersonId
4              ∧ x.MovieId = y.MovieId
5          ) ∨ ∃ x2 ∈ acts (d.PersonId = x2.PersonId
6              ∧ ∃ y ∈ acts (a.PersonId = y.PersonId
7                  ∧ x2.MovieId = y.MovieId
8              )
9      ) ∧ ¬ (
10         ¬(d.Firstname = 'Lars' ∧ d.Lastname = 'von Trier')
11         ∧ ¬(d.Firstname = 'Michael' ∧ d.Lastname = 'Haneke')
12     )
13 }
```

### 1.1.2

```

1  {[c.name, c.location, s2.time, s2.price, m.title] | c ∈ cinema
2      ∧ ∃ r ∈ screeningRoom (r.inCinema = c.CinemaId
3          ∧ ∃ s1, s2 ∈ screening (s1.RoomId = r.RoomId
4              ∧ ∃ m ∈ movie (s1.MovieId = m.MovieId ∧ s2.MovieId = m.MovieId
5                  ∧ s1.date = Today ∧ s1.time = '9:00'
6              )
7          )
8      )
9  }
```

### 1.1.3

```

1  {[a.Firstname, a.Lastname] | a, d ∈ person
2      ∧ ∀ m ∈ movies (d.Firstname = 'Lars' ∧ d.Lastname = 'von Trier'
3          ∧ ∃ x ∈ directs (m.MovieId = x.MovieId
4              ∧ d.PersonId = x.PersonId
5          ) ⇒ (
6              ∃ y ∈ acts (m.MovieId = y.MovieId
7                  ∧ a.PersonId = y.PersonId
8              )
9          )
10     ) ∧ ∃ t ∈ acts (t.PersonId = a.PersonId)
11 }
```

## 1.2

### 1.2.1

```
1 { [afn, aln, dfn, dln] |
2   ∃ dPid ( [dPid, dfn, dln] ∈ Person
3     ∧ ∃ aPid ([aPid, afn, aln] ∈ Person
4       ∧ ∃ aMid, r ([aMid, aPid, r] ∈ acts
5         ∧ ∃ dMid ([dMid, dPid] ∈ directs
6           ∨ ∃ role ([dMid, dPid, role] ∈ acts
7             ∧ ¬(
8               ¬(dfn = '' ∧ dln='') )
9             ∧ ¬(dfn = '' ∧ dln='') )
10          )
11         )
12       )
13     )
14   )
15 }
16 }
```

### 1.2.2

```
1 { [n, l, time, p, t] |
2   ∃ cId ( [cId, n ,l] ∈ Cinema
3     ∧ ∃ rId, num ([rId, num, cId] ∈ screeningRoom
4       ∧ ∃ mId, d, len ([mId, rId, time, d, len, p] ∈ screening
5         ∧ ∃ time2, d2, len2, p2, rId2 ([mId, rId2, time2, d2, len2, p2] ∈ screening
6           ∧ ∃ lang, fsk, y ([mId, t, lang, fsk, y] ∈ Movie
7             ∧ time2 = (21*60)
8             ∧ d2 = today
9           )
10         )
11       )
12     )
13   )
14 }
```

### 1.2.3

```
1 { [afn, aln] |
2   ∃ aPid ( [aPid, afn, aln] ∈ Person
3     ∧ ∀ mId (
4       ∃ dPid, dfn, dln ( [dPid, dfn, dln] ∈ Person
5         ∧ [dPid, mId] ∈ directs
6         ∧ dfn = 'Vorname'
7         ∧ dln = 'Nachname'
8       ) ⇒ (
9         ∃ r (
10          [aPid, mId, r] ∈ acts
11        )
12      )
13    ) ∧ ∃ mId2, r2 (
14      [aPid, mId2, r2] ∈ acts
15    )
16  )
17 }
```

## 2 Aufgabe

### 2.1 Beispiel

$$\{n | \neg(n \in \text{Professoren})\}$$

### 2.2 Erklärung

Man definiert die Domäne der Formel des Ausdrucks. Dies ist eine Menge, welche folgende Elemente beinhaltet:

- alle in der Formel vorkommende Konstanten
- alle Attributwerte von Relationen, welche in der Formel referenziert werden

Das Ergebnis des Ausdrucks muss nun Teilmenge dieser Domäne sein, um sicher zu stellen, dass das Ergebnis der Anfrage endlich ist.

## 3 Aufgabe

$$\{n | \neg(n \in \text{Professoren})\}$$

Eine unsichere Anfrage wie in Aufgabe 2, kann in der relationalen Algebra nicht konstruiert werden. Wie der Name schon sagt, beschreibt die relationale Algebra Rechenoperationen auf Relationen, welche an sich bereits endlich sind. Alle Operatoren der relationalen Algebra ermöglichen höchstens das vervielfachen der Anzahl der betrachteten Elemente, aber nie das Wachstum ins unendliche.

## 4 Aufgabe

### 4.1

```
1 select m.Title
2 from movie m, directs d, person p
3 where m.MovieId = d.MovieId
4       and d.PersonId = p.PersonId
5       and p.Lastname = "Murnau"
6       and p.Firstname = "Friedrich Wilhelm"
```

### 4.2

```
1 select m.Title
2 from movie m, directs d, person p
3 where m.MovieId = d.MovieId
4       and d.PersonId = p.PersonId
5       and p.Lastname = 'Burton'
6       and p.Firstname = 'Tim'
7       and m.MovieId not in (
8         select a.MovieId
9         from acts a, person pjd
10        where a.PersonId = pjd.PersonId
11              and pjd.Lastname = 'Depp'
12              and pjd.Firstname = 'Johnny')
```

### 4.3

```
1 select m.Title, c.Name, c.Location, s.Time, s.Price
2 from movie m, cinema c, screening s, genre g, hasGenre hg, screeningRoom sr
3 where s.Date = Today()
4       and s.MovieId = m.MovieId
5       and m.MovieId = hg.MovieId
6       and hg.GenreId = g.GenreId
7       and g.Name = 'Documentary'
8       and s.RoomId = sr.RoomId
9       and sr.inCinema = c.CinemaId
```

### 4.4

```
1 Select a.Firstname, a.Lastname, d.Firstname, d.Lastname
2 from person a, person d
3 where
4     exists(
5         select *
6         from directs x
7         where d.PersonId = x.PersonId
8         and exists(
9             select *
10            from acts y
11            where a.PersonId = y.PersonId
12                  and x.MovieId = y.MovieId
13        )or exists(
14            select *
15            from acts x2
16            where d.PersonId = x2.PersonId
17            and exists(
18                select *
19                from acts y
20                where a.PersonId = y.PersonId
21                      and x2.MovieId = y.MovieId
```

```
22         ) and not (
23         not(d.Firstname = "Lars" and d.Lastname = "von Trier")
24         and not(d.Firstname = "Michael" and d.Lastname = "Haneke")
25     )
26 )
27 )
```

## 4.5

```
1  Select c.Name, c.Location, s.Price, m.Title
2  from cinema c, screening s, screeningRoom r, movie m
3  where m.MovieId = s.MovieId
4         and s.RoomId = r.RoomId
5         and c.CinemaId = r.inCinema
6         and exists (
7             Select m.MovieId = s.MovieId
8             and s.RoomId = r.RoomId
9             and c.CinemaId = r.inCinema
10            and s.Date = Today()
11            and s.Time = 2100
12 )
```