Optimization: Cost Estimation: Scans

R&G Chapter 15

(slides adapted from content by J.Gehrke, J.Shanmugasundaram, and/or C.Koch)

Reminder: Project 1 Deadline Extended to Tuesday Night

20 bonus points if you've already submitted

Project 2 will be assigned Wednesday

Homework 4 (due Sunday night) will be posted tonight (12:01 Monday Morning)

Enumerating Possible Plans

- There are many algorithms suitable for processing a specific data set:
 - I. How do we identify feasible algorithms.
 - 2. How do we choose between them

Single Relation Queries

```
SELECT O.Rank, COUNT(*),
FROM Officers O
WHERE O.Rank >= 2
AND O.Age > 20 AND O.Age < 30
GROUP BY O.Rank
HAVING COUNT(DISTINCT O.Ship) > 2
```

What is this query asking?

Single Relation Queries

```
SELECT O.Rank, COUNT(*),
```

FROM Officers O

WHERE O.Rank >= 2

AND O.Age > 20 AND O.Age < 30

GROUP BY O.Rank

HAVING COUNT(DISTINCT O.Ship) > 2

Group Exercise: Compose a (naive) RA expression for this query (using scan, select, project, aggregate)

Single Relation Queries

Group Exercise: Compose a (naive) RA expression for this query (using scan, select, project, aggregate)

Evaluation Strategies

- How do we read in the data (Which Access Path)?
 - No Index
 - Single Unsorted Index
 - Single Index Sorted on the GB Key
 - Multiple-Index
 - Index-Only Access
- How do we implement the Group-By?
 - Hash-Grouping
 - Sort/Group

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Friday, March 1, 13

Single Unsorted Index Scans have a Fixed Reduction factor. We'll cover this next week Sorted Index scans allow us to avoid sorting for the aggregate

Multiple-index scans are possible if we have indices that store record IDs. We can scan each index and compute the intersection of all record IDs.

If the index key includes ALL attributes referenced by the query, we can actually walk the index and reconstruct the relevant tuples from the key.



I) File Scan

Cost:

Officers
50 bytes/tuple
80 tuples/page
500 pages

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I) File Scan

Cost: 500 IOs

I) File Scan

Cost: 500 IOs

2) Select/Project/Write Cost: ?

I) File Scan Cost: 500 IOs

2) Select/Project/Write Cost: <500 IOs

Reduction Factor

Selection Eliminates Tuples

Projection Eliminates Columns

Selection/Projection Reduce The Size of a Relation!

Reduction Factor

(The fraction of Tuples/Bytes present in the output)

I) File Scan

Cost: 500 IOs

2) Select/Project/Write Cost: <500 IOs

Reduction Factor 0.75

I) File Scan

Cost: 500 IOs

2) Select/Project/Write Cost: 375 IOs

Reduction Factor 0.75

I) File Scan

Cost: 500 IOs

2) Select/Project/Write Cost: 375 IOs

Reduction Factor 0.75

3) Finish the Sort (50 frames) Cost: ?

I) File Scan

Cost: 500 IOs

2) Select/Project/Write Cost: 375 IOs

Reduction Factor 0.75

3) Finish the Sort (50 frames) Cost: 375 IOs

I) File Scan

Cost: 500 IOs

2) Select/Project/Write Cost: 375 IOs

Reduction Factor 0.75

3) Finish the Sort (50 frames) Cost: 375 IOs

4) Scan + Aggregate Cost: ?

Officers
50 bytes/tuple
80 tuples/page

500 pages

I) File Scan

Cost: 500 IOs

2) Select/Project/Write Cost: 375 IOs

Reduction Factor 0.75

3) Finish the Sort (50 frames) Cost: 375 IOs

4) Scan + Aggregate Cost: [included]

<u>Officers</u>

50 bytes/tuple

80 tuples/page

500 pages



Reduction Factors

- Reduction factors give us a way to estimate how many tuples are in a relation.
- This is crucial for estimating the performance of a query plan.
- Not just in terms of the IO cost of the plan, but also the compute cost.
 - IO: #pages ↔ Compute: #tuples

(Selection)

Simplify to Conjunctive Normal Form

$$\sigma_{a_1 \wedge \dots \wedge a_n}(R) = \sigma_{a_1}(\dots \sigma_{a_n}(R))$$

Split Conjunctions

(Selection)

Simplify to Conjunctive Normal Form

$$a_1 \lor \ldots \lor a_n = \neg(\neg a_1 \land \ldots \neg a_n)$$
(De Morgan's Law)

$$RF(a_1 \lor ... \lor a_n) = 1 - ((1 - RF(a_1)) \cdot ... \cdot (1 - RF(a_n)))$$

$$(RF(\neg a) = I-RF(a))$$



SELECT *
FROM R
WHERE A = 1

Assume a Uniform Distribution

$$\frac{1}{NKeys(A)}$$

How do we get NKeys?

```
SELECT *
FROM R, S
WHERE R.A IN (1, 2, ...)
```

Assume a Uniform Distribution

$$\frac{\#Values}{NKeys(A)}$$

SELECT *
FROM R
WHERE A > 1

Assume a Uniform Distribution

$$\frac{High(A) - value}{High(A) - Low(A)}$$

How do we get High/Low?

```
SELECT *
FROM R, S
WHERE R.A = S.B
```

Nested-Loop Index Scan on the Smaller Relation

 $\frac{1}{\max(NKeys(A), NKeys(B))}$



Index Scan

I) Index Scan

Cost: ? IOs

What Indexes are Available?

O.Age?

O.Rank?

Rank > 2

$$RF = ?$$

RF = ?

Ranks: 0-5 (in increments of 0.5)

Ages: 15-100 (in increments of 1)

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Assuming a uniform distribution, 85 Age keys, 11 Rank keys 9 legitimate ages, 9 / 85 = 0.1056 legitimate ranks, 6 / 11 = 0.545Use the Age index!

Index Scan-Sort/Group

I) Index Scan

Cost: ~53 IOs

2) Select/Project/Write Cost: ~23 IOs Assume Projection has an RF of 0.8

Ranks: 0-5 (in increments of 0.5)

Ages: 15-100 (in increments of 1)

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50 bytes/tuple
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500 pages

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We still get the 0.545 reduction factor on the 2nd select Assume another 0.8 on the projection Total = 0.4346

Index Scan-Sort/Group

I) Index Scan

Cost: ~53 IOs

2) Select/Project/Write Cost: 23 IOs Assume Projection has an RF of 0.8

Ranks: 0-5 (in increments of 0.5)

Ages: 15-100 (in increments of 1)

Officers
50 bytes/tuple
80 tuples/page
500 pages

23

Friday, March 1, 13

We still get the 0.545 reduction factor on the 2nd select Assume another 0.8 on the projection Total = 0.4346

Index Scan-Sort/Group

I) Index Scan

Cost: ~53 IOs

2) Select/Project/Write Cost: 23 IOs Assume Projection has an RF of 0.8

3) In-Mem Sort (50 frames) Cost: 0 IOs

4) Scan + Aggregate Cost: 0 IOs

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Ranks: 0-5 (in increments of 0.5)

Ages: 15-100 (in increments of 1)

<u>Officers</u>

50 bytes/tuple

80 tuples/page

500 pages

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We still get the 0.545 reduction factor on the 2nd select Assume another 0.8 on the projection Total = 0.4346

