

Discrete Loop Soup.

(1) Definitions:

Def: i) $\tilde{I}_n : \{0, 1, \dots, n\} \rightarrow E$. $\tilde{I} = \bigcup \tilde{I}_n$

ii) $\tilde{I}_{r,n} = \tilde{I}_n \cap \{l(0) = l(n), l(i) \neq l(i+1)\}$

space of rooted (based) loops of

length n . $\tilde{I}_r = \bigcup_{n \in \mathbb{Z}^+} \tilde{I}_{r,n}$.

iii) $\tilde{I}^* = \tilde{I}_r / \sim$, where $l \sim l'$ if
 $\exists n \in \mathbb{Z}$ s.t. $\theta_n(l) = l'$.

Rmk: We endow them with σ -algebra
 generated by finite loops.

Def: (Measure)

i) $\tilde{m}_r(l) = \frac{1}{n} P_{x_0, x_1} \dots P_{x_n, x_0}$ for $\forall l =$
 $(x_0, x_1, \dots, x_n) \in \tilde{I}_r$.

ii) $\tilde{m}^* = \tilde{\pi}^* \circ \tilde{m}_r$ $\tilde{\pi}^* : \tilde{I}_r \rightarrow \tilde{I}^*$ canonical.

iii) Set $j_l(\gamma) = \#\{i \in \{1, \dots, \tilde{\pi}(l)\} \mid x_i = \gamma\}$ for
 $l = (x_0, \dots, x_{\tilde{\pi}(l)})$. $\tilde{m}_x(l) =: j_{x_0}^{-1} \cdot P_{x_0, x_1} \dots P_{x_n, x_0}$

Rmk: For $\mathcal{M}_{0,x}$ is set of unrooted loops visit
 point x in D . Then:

$\tilde{m}^*|_{\mathcal{M}_{0,x}} = L \circ \tilde{m}_x$ where L is loop-erasure.

Def: i) $\tilde{\nu} = \sum_i \delta_{x_i}^*$ point measure on $(\tilde{L}^*, \tilde{L}^*, |$
 ν is σ -finite. $|Z| \leq |S|$.

ii) $\tilde{\mu}_\alpha \in \mathcal{M}$ on $(\tilde{L}^*, \tilde{L}^*)$ is law of PPP with
 intensity $\leq \tilde{\nu}^*$ for $\alpha > 0$.

iii) Set \mathcal{L}_α is loop-soup $\sim \tilde{\mu}_\alpha$.

(2) Properties of measure:

prop. $F \subset \subset E$. E is transient graph. Then:

$$\tilde{\nu}^*(\mathcal{L}^* \cap F \neq \emptyset) = \log \det (h|_{F \times F})$$

Pf: i) $\tilde{\nu}^*(\mathcal{L}^* \subset R) = \tilde{\mu}_r(\mathcal{L} \subset R)$

$$= \sum_{k \geq 2} \tilde{\mu}_r(\mathcal{L} \subset R, \tilde{f}(\mathcal{L}) = k)$$

$$= \sum_{k \geq 2} \frac{\text{Tr}(P|_{R \times R})^k}{k}$$

2) LHS = $\lim_{n \rightarrow \infty} \tilde{\mu}_r(\mathcal{L} \cap F \neq \emptyset, \mathcal{L} \subset B_n)$

$$= \lim_{n \rightarrow \infty} \tilde{\mu}_r(\mathcal{L} \subset B_n) - \mu_r(\mathcal{L} \subset B_n/F)$$

$$= \lim_{n \rightarrow \infty} \log \frac{\det(I - P|_{B_n/F \times B_n/F})}{\det(I - P|_{B_n \times B_n})}$$

where $B_n \subset \subset E$. $B_n \uparrow E$.

3) With Jacobi's equality:

$$\frac{\det(I - P|_{(B_n/F)^2})}{\det(I - P|_{B_n^2})} = \det(I - P|_{B_n \times B_n})^{-1} |_{F \times F}$$

$$\xrightarrow{n \rightarrow \infty} \det(h|_{F \times F})$$

Cor. For $\{x_i\}_n \in E$. we have:

$$\tilde{m}^* \subset x_i \in \mathcal{L}^*. \forall (1 \leq i \leq n) =$$

$$\sum_{\substack{A = \{x_i\}_n \\ A \neq \emptyset}} (-1)^{|A|+1} \log \det (h|_{A \times A})$$

Def: i) For $\mathcal{L} \in \tilde{\mathcal{L}}_r$. multiplicity of \mathcal{L} is:

$$mult(\mathcal{L}) = \max \{k \mid \mathcal{L} = (\mathcal{L}_1, \mathcal{L}_2, \dots, \mathcal{L}_k), \mathcal{L}_i \in \tilde{\mathcal{L}}_r, \\ \text{and } \mathcal{L}_1 = \mathcal{L}_2 = \dots = \mathcal{L}_k\}.$$

Rmk: If $\tilde{f}(\mathcal{L}^*) = n$. Then: we have.

$$\tilde{m}^*(\mathcal{L}^*) = \frac{n}{mult(\mathcal{L})} \cdot \tilde{m}_r(\mathcal{L}) \text{ for the}$$

$$\text{loop } \mathcal{L} = (\mathcal{Z}_1, \dots, \mathcal{Z}_n)$$

ii) For $S_1, S_2 \subset E$. $S_1 \cap S_2 = \emptyset$. Set map:

$$\begin{array}{ccc} \tilde{\mathcal{L}}^* \cap \{\mathcal{L}^* \text{ visits } S_1, S_2\} & \longrightarrow & P(\tilde{\mathcal{L}}_r) \\ \mathcal{L}(S_1, S_2) : \mathcal{L}^* & \longmapsto & R \end{array}$$

R is subset of $\tilde{\mathcal{L}}_r$ containing \mathcal{L}' satisfies:

a) $\mathcal{L} \sim \mathcal{L}'$ in $\tilde{\mathcal{L}}^*$

b) $\mathcal{L}' = (\mathcal{Z}_1, \dots, \mathcal{Z}_n)$. st. $\mathcal{Z}_1 \in S_1$.

c) $\exists i$. st. $\mathcal{Z}_i \in S_2$ and $\mathcal{Z}_j \notin S_1 \cup S_2, \forall j > i$.

Rmk: $\mathcal{L}(S_1, S_2) \cap \mathcal{L}^* \neq \emptyset \Leftrightarrow \mathcal{L} \text{ visits } S_1, S_2$.

iii) For $\mathcal{L}' = (\mathcal{Z}_1, \dots, \mathcal{Z}_n) \in \mathcal{L}(S_1, S_2) \cap \mathcal{L}^*$. Set $\mathcal{Z}_0 = 1$

$$\mathcal{Z}_{2k+1} = \inf \{j > \mathcal{Z}_{2k} \mid \mathcal{Z}_j \in S_2\}, \mathcal{Z}_{2k+2} = \inf \{j > \mathcal{Z}_{2k+1} \mid \mathcal{Z}_j \in S_1\}$$

and set $\inf \{Q\} = n+1$.

iv) Set $k(u)$ satisfies $\sum_{k(u)-1}^n \mathbb{1}_{k(u)} = n$. $\sum_{k(u)}^n \mathbb{1}_{k(u)} = n+1$

prop. $\forall \ell^* \in \tilde{L}^*$ visits S_1, S_2 . $\tilde{f}(\ell^*) = n$. discrete loop

$$\Rightarrow \tilde{m}^*(\ell^*) = \frac{n}{k(u)} \cdot \sum_{\ell' \in L(S_1, S_2)(u)} \tilde{m}_r(\ell')$$

Pf: It follows $k(u) = m(u) \cdot \# L(S_1, S_2)(u)$

(It's guaranteed by c) in def ii))

prop. For $E = \mathbb{Z}^d$, $d \geq 3$, $k_x = 0$.

i) For $\lambda > 1$, $\exists C = C(d, \lambda) < \infty$, s.t. $\forall N \geq 1$, $M \geq \lambda N$
and $k \in B(0, N)$, $\tilde{m}^*(k \leftrightarrow \partial B(0, M)) \leq C \cdot \text{cap}(k) \cdot M^{2-d}$

ii) Set $F(d, N) = I_{\{d=3\}} \cdot N + I_{\{d=4\}} \cdot \frac{N^2}{\log N} + I_{\{d \geq 5\}} \cdot N^2$

For $\lambda > 1$, $\exists C = C(d, \lambda) < \infty$, s.t. $\forall N \geq 1$, $M \geq \lambda N$

and $k \in B(0, N)$. We have:

$$\tilde{m}^*(k \leftrightarrow \partial B(0, M), \text{cap}(k) > F(d, N)) \geq C \cdot \text{cap}(k) \cdot M^{2-d}$$

Pf: i) Set $z_0 = M_k$, $z_{2k+1} = \inf \{n > z_{2k} \mid z_n \in \partial B(0, M)\}$

$$z_{2k+2} = \inf \{n > z_{2k+1} \mid z_n \in k\}$$

$$\Rightarrow \tilde{m}^*(k \leftrightarrow \partial B(0, M)) = \sum_{n \geq 1} \frac{1}{n} \sum_{x \in \partial k} \mathbb{P}^x(X_{z_{2n}} = x)$$

(count $m(u) = n$, $n \geq 1$, with $\tilde{m}_r(\cdot) = 0$)

By Markov prop. and Markov ineqn.:

$$\begin{aligned} \mathbb{P}^x(X_{z_{2n+2}} = x) &\leq \mathbb{P}^0(X_{z_2} = x) \cdot \left(\max_{\partial k} \mathbb{P}^x(z_2 < \infty) \right)^n \\ &\leq C \cdot \mathbb{P}^0(X_{z_2} = x) \cdot C^n, \quad C < 1. \end{aligned}$$

$$S_1 = \tilde{M}^* \subset K \leftrightarrow \partial B(0, m) \leq C \cdot P^0(Z_2 < \infty)$$

By last exit id. we have:

$$P^z(M_k < \infty) \leq M^{2-k} \text{cap}(K), \quad \forall z \in \partial B(0, m)$$

$$\text{ii) Check: } P^x(\text{cap}(\{Z_0, \dots, Z_{H_k}\}) > C F(\lambda, n)) \\ \geq C z_c(\lambda, n) > 0.$$

(3) Decomposition of Loops

in Excursion:

Next, we set $E = \mathbb{Z}^d$, $d \geq 3$, $X_0 = 0$, X_n is
a SRW on \mathbb{Z}^d , $R = A \cap B$, $A, B \subset \mathbb{Z}^d$.

Def: $\mathcal{L}_{A,B} = \{ \ell^* \in \tilde{\mathcal{L}}^* \mid \ell^* \cap A, \ell^* \cap B \neq \emptyset \}$.

Lemma. There $\exists C = C(d)$, $C = C(d) < \infty$, st.

$\forall n \geq 1, m \geq 2n, A \subset B(0, n), X \in B(0, m)$.

$$C \tilde{e}_{A \cap B} \leq P_x(X_{H_A} = \eta \mid H_A < \infty) \leq C \tilde{e}_{A \cap B}$$

Lemma. $\mathcal{I}_{A,B}^q$ is a PPP on $\mathcal{L}_{A,B} \stackrel{A}{=} \mathcal{L}(A, B) \cup \mathcal{L}(A, B)$

with intensity $\leq \tilde{m}_{A,B}$, where $\tilde{m}_{A,B}(\ell)$

$$\stackrel{n=|\ell|}{=} \frac{1}{k(\ell)} P_x(\ell(X_0, \dots, X_n) = \ell, X_n \in X_1) \text{ for } \ell$$

$\in \mathcal{L}(A, B)$. Then $\pi^* \in \mathcal{I}_{A,B}^q$ is PPP on

$\mathcal{L}_{A,B}$ with intensity $\leq \tilde{I}_{A,B}^q$.

Pf: By Transf prop. of PPPs.

Def: i) $\dot{I}_{A,B}^{\tau,i}$ is restriction of $\dot{I}_{A,B}^{\tau}$ on $\dot{L}_{A,B}^i$

$$=: \{ \alpha \in \dot{L}_{A,B} \mid k(\alpha) = j \}.$$

Rmk: i) $(\dot{I}_{A,B}^{\tau,i})_{j \geq 1}$ are indep PPPs with

$$\text{intensity } (\dot{I}_{A,B}^i)_{i \geq 1}$$

$$\text{ii) } \dot{I}_{A,B}^{\tau} = \sum_{j \geq 1} \dot{I}_{A,B}^{\tau,j}.$$

ii) For $\alpha^* \in \dot{L}_{A,B}$, $\alpha' \in \dot{L}_{A,B}(\alpha^*)$.

$$\text{set } \phi_1(\alpha') = 1, \quad \psi_1(\alpha') = \inf \{ j > 1 \mid X_j \in B \}.$$

$$\phi_k(\alpha') = \inf \{ j > \psi_{k-1}(\alpha') \mid X_j \in A \}.$$

$$\psi_k(\alpha') = \inf \{ j > \phi_k(\alpha') \mid X_j \in B \}.$$

iii) For $\alpha' \in \dot{L}_{A,B}^j$, $\alpha' = (X_1, \dots, X_n)$.

$$\text{set } \tilde{\phi}_i(\alpha') = X_{\phi_i(\alpha')}, \quad \tilde{\psi}_i(\alpha') = X_{\psi_i(\alpha')}, \quad 1 \leq i \leq j.$$

excursion from A to B:

$$\vec{W}_i(\alpha') = (X_{\phi_i(\alpha')}, \dots, X_{\psi_i(\alpha')}), \quad 1 \leq i \leq j$$

and excursion from B to A:

$$\overleftarrow{W}_i(\alpha') = (X_{\psi_i(\alpha')}, \dots, X_{\phi_{i+1}(\alpha')}), \quad 1 \leq i \leq j-1.$$

$$\overleftarrow{W}_j(\alpha') = (X_{\psi_j(\alpha')}, \dots, X_n, X_1)$$

$$\text{iv) } \dot{I}_{A,B}^{\tau,j} = \dot{I}((\tilde{\phi}_1(\alpha'), \tilde{\psi}_1(\alpha')), \dots, (\tilde{\phi}_j(\alpha'), \tilde{\psi}_j(\alpha'))) \\ \text{for } \alpha' \in \dot{L}_{A,B}^j$$

$$\vec{\sum}_{A \rightarrow B}^{\tau, i} = \sum \{ \vec{W}_1(c_1), \dots, \vec{W}_j(c_j) \} \quad i' \in \vec{\mathcal{I}}_{A \rightarrow B}^{\tau, i}$$

$$\overleftarrow{\sum}_{A \rightarrow B}^{\tau, i} = \sum \{ \overleftarrow{W}_1(c_1), \dots, \overleftarrow{W}_j(c_j) \} \quad i' \in \overleftarrow{\mathcal{I}}_{A \rightarrow B}^{\tau, i}$$

Prop. For path $w = (x_0, x_1, \dots)$ set $z_0(w) = 0$.

$$z_{2j+1}(w) = \inf \{ k > z_{2j}(w) \mid x_k \in B \}$$

$$z_{2j+2}(w) = \inf \{ k > z_{2j+1}(w) \mid x_k \in A \}$$

Then, for $\forall \tau > 0, j \geq 1$.

i) Intensity of $\sum_{A \rightarrow B}^{\tau, i}$ is

$$(c_1, b_1) \dots (c_j, b_j) \mapsto \frac{\tau}{j} p_{n_i} \left(\begin{array}{l} z_{2j} < \infty, \quad x_{z_{2j}} = a_i \\ x_{z_{2j+1}} = b_i, \quad x_{z_{2j+2}} = b_i \end{array} \right)$$

ii) Conditional on $\sum_{A \rightarrow B}^{\tau, i} = \sum \{ (c_{1i}, b_{1i}), \dots, (c_{ji}, b_{ji}) \}_{i=1}^n$

$\vec{\sum}_{A \rightarrow B}^{\tau, i}, \overleftarrow{\sum}_{A \rightarrow B}^{\tau, i}$ are indept and sampled as products of bridge measures $p_{n_{ik}, b_{ik}}^B$ resp.

$p_{b_{ik}, n_{ik+1}}^A$ where $p_{x, \eta}^m(\cdot) = p_x(\cdot \mid x_{n_m} = \eta)$

Remark: It means loops from $\mathcal{I}_{A \rightarrow B}^{\tau}$ can be sampled by:

i) First sample number and its starting and ending locations of all the loops in $\mathcal{I}_{A \rightarrow B}^{\tau}$

(It's achieved by sampling indeptly $(\sum_{A \rightarrow B}^{\tau, i})_j$)

iii) Sample indept RW bridges from $p_{\cdot, \cdot}^B$

and $p_{\cdot, \cdot}^A$

Pf: Note that present of M_t only depends on the nearest past and future.

$$\Rightarrow m_{A,B}(i) = \frac{1}{j} \mathbb{P}_{X_1} \left(\begin{array}{l} Z_{2j} < \infty, \quad X_{Z_{2j}} = x, \quad 1 \leq i \leq j-1 \\ X_{Z_{2(i-1)}} = \tilde{\varphi}_i(x'), \quad X_{Z_{2i-1}} = \tilde{\gamma}_i(x') \end{array} \right)$$

$$\cdot \prod_{i=1}^j \mathbb{P}_{\tilde{\varphi}_i(x'), \tilde{\gamma}_i(x')}^B (W_i(x')) \prod_{i=1}^j \mathbb{P}_{\tilde{\varphi}_i(x'), \tilde{\gamma}_i(x')}^A (V_i(x'))$$

prop. (Large Deviation)

For $Z_{A,B}^T$ total number of excursions from A to B of loops from $Z_{A,B}^T$. If $\sup_{B \in \mathcal{B}} \mathbb{P}_B (H_A < \infty) \leq \frac{1}{2e}$

Then: $\mathbb{P} (Z_{A,B}^T \geq k) \leq e^{T-k}$.

Pf: Set $Z_{A,B}^{T,j}$ is number of loops in $Z_{A,B}^{T,j}$

$$Z_{A,B}^T \sim \sum_{j=1}^T j Z_{A,B}^{T,j}$$

By prop 1), where $Z_{A,B}^{T,j} \sim \mathcal{P}(\lambda_j)$

$$\lambda_j = \frac{1}{j} \sum_{x \in A} \mathbb{P}_x (Z_{2j} < \infty, X_{Z_{2j}} = x)$$

$$\leq \frac{1}{j} \sup_A \mathbb{P}_x (Z_{2j} < \infty) \stackrel{\text{M.P.}}{\leq} \frac{1}{j} \left(\sup_A \mathbb{P}_x (Z_2 < \infty) \right)^j$$

$$\stackrel{\text{M.P.}}{\leq} \frac{1}{j} \left(\sup_B \mathbb{P}_B (H_A < \infty) \right)^j \leq \frac{\alpha}{j} (2e)^{-j}$$

$$\text{Combine: } \mathbb{E} (e^{Z_{A,B}^T}) = \prod_{j=1}^T \mathbb{E} (e^{j Z_{A,B}^{T,j}})$$

With exponential Chebyshev. inequality.

(*) Properties of Loop soups:

① Def: For $x \in E$. \mathcal{L}_D^x is set of loops in \mathcal{L} going through point x in D

prop. $\tilde{P}_1^x(\mathcal{L}_D^x = \emptyset) = e^{-\tilde{M}^*(M_D, x)} = 1/h_0(x, x)$

Pf: Define u is the sum of $\tilde{M}^*(M_D, x)$ over all rooted loops from x to x in D visiting x only once

$$\Rightarrow \tilde{M}^*(M_D, x) = \sum_{j \geq 1} u^j / j = -\ln(1-u)$$

$$\text{Besides: } h_0(x, x) = 1 + \sum_{j \geq 1} u^j = 1/(1-u)$$

Cor. If \mathcal{L}_D is set of loops in \mathcal{L} staying in D . Then: $\tilde{P}_1^x(\mathcal{L}_D = \emptyset) = 1/h_0$.

Pf: LHS = $\prod_{k \geq 0} \tilde{P}_1^x(\text{No loop in } D \setminus \{x_0 \dots x_k\} \text{ goes through } x_{k+1})$

$$= \prod_{k \geq 0} h_0 / \{x_0 \dots x_k\}(x_{k+1}, x_{k+1}) = 1/h_0$$

where define $x_0 = x$.

Cor. $\tilde{M}^*(M_D, x)$ is finite if $|D| < \infty$.

Pf: $\sum_D \tilde{M}^*(M_D, x) < \infty$ if $|D| < \infty$.

Cor. $\tilde{P}_x(\mathcal{I}_0^x = \mathcal{Q}) = (1/hoc(x,x))^r$

$\tilde{P}_x(\mathcal{I}_0 = \mathcal{Q}) = (tot(h_0))^{-r}$

Pf: Note for $\mathcal{I} \sim PPP(\alpha M)$

$\mathcal{I}' \sim PPP(\alpha' M)$. indep loop soups.

$\mathcal{I} + \mathcal{I}' \sim PPP((\alpha + \alpha') M)$.

Consider $s \in \mathbb{Z}^+ \Rightarrow \mathcal{Q}^+ \Rightarrow \mathcal{K}^+$.

prop. For $N = \# \mathcal{I}_0^x$. We choose uniformly at random

an order of these N rooted loops $\lambda_1, \dots, \lambda_N$

Let $\tilde{\lambda} = \lambda_1 \circ \lambda_2 \circ \dots \circ \lambda_N$. concatenation of them

If \mathcal{L} is a rooted loop from x to x in D .

Then $\tilde{P}_1(\tilde{\lambda} = \mathcal{L}) = (2d)^{-\tilde{S}(\mathcal{L})} / hoc(x,x)$

Pf: Lemma. $\forall k \in \mathbb{Z}^+$. $\sum_{r \geq 1} \sum_{\sum_{i=1}^r j_i = k} \frac{1}{r!} \prod_{i=1}^r j_i = 1$.

Note: Fix $(\mathcal{L}^i)_1^r$. $\mathcal{L}^1 \circ \dots \circ \mathcal{L}^r = \mathcal{L}$.

$\tilde{P}_1(\lambda_1 = \mathcal{L}^1, \dots, \lambda_N = \mathcal{L}^r, N(\mathcal{M}_{D,x}) = r) =$

$$= \frac{1}{r!} \cdot \prod_{i=1}^r \tilde{m}^*(\mathcal{L}^i) \cdot e^{-\tilde{m}^*(\mathcal{M}_{D,x})}$$

$$= (2d)^{-\tilde{S}(\mathcal{L})} \cdot \frac{1}{r!} \cdot \frac{1}{\prod_{i=1}^r j_i} \cdot 1/hoc(x,x)$$

Sum over $r \geq 1$ and $\sum_{i=1}^r j_i = j_{\mathcal{L}}(x)$

Rmk: $\tilde{\lambda} \stackrel{\mathcal{L}}{\sim} \lambda$. the loop from x to x that we erase before last visit to x .

Besides, the constructions are similar!

By the prop and remark above, we can recover the whole process of Wilson's algorithm in $D = \{X_1, \dots, X_n\}$ with root x_0 by sampling unrooted loop soups when $r=1$ and indept WST.

1) Sample N_1 loops in $\mathcal{L}_0^{x_0}$, order it uniformly at random and concatenate it. we obtain a loop η_1 .

2) Jump η_1 to η_2 .

3) Sample N_2 loops in $\mathcal{L}_{D/\eta_1}^{\eta_2}$ and repeat 1)

4) Proceed until reach the root x_0 .

\Rightarrow We recover a whole loop from $(\eta_1, \dots, \eta_{s+1})$

Prop. (Converse)

Start from Wilson's algorithm:

to read off rooted loops λ_i from η_i to η_i in $D / \{\eta_1, \dots, \eta_{i-1}\}$ (non-loop-erased)

i) For each i indeptly, if λ_i returns k times to η_i , choose to split it into r smaller loops, with j_1, \dots, j_r times return to η_i , resp.

s.t. $\sum_{i=1}^r j_i = k$. with prob: $\tilde{\mathbb{P}}_i(\lambda_i = \lambda'_1 \dots \lambda'_r) =$

$N(m_0, \eta_i) / \tilde{\mathbb{P}}_i(\tilde{\lambda} = \lambda) = 1/r! \prod_{i=1}^r j_i!$

ii) Then we obtain a point process of rooted loop in D from $(\lambda_i)_i^s$.

\Rightarrow This process is indpt of $\text{KST } \tau$ instead by Wilson's algorithm with law of loop-soaps in D with $\tau = 1$.

② Occupation Field:

i) Def: $V = (V(x))_{x \in D}$ is occupation field of \mathcal{L}_0 loop soaps with $\tau = 1$ in D . i.e. $V(x)$ is total number of visit to x by all loops in \mathcal{L}_0 .

prop. $V \sim \tilde{V} = V - 1$

Pf: By prop. above. we know it's identical dist. as Occupation field of Wilson's algorithm.

Cor. For $\tau > 0$. V_τ is occupation field of loop soap with intensity $\propto \tilde{m}^\tau$.

$$\Rightarrow \mathbb{E} \left[\frac{1}{n} \sum_{i=1}^n (1 + k(x_i))^{-V(x_i)} \right] = \left(\frac{1 - \bar{A}_1}{1 - \bar{A}_{0,k}} \right)^\tau$$

Pf: By infinite divisibility.

ii) Def: (Occupation Time on edge)

Set $T = (T(e))_{e \in E(G)}$ is occupation field on $E(G)$ i.e. $T(e)$ is total number of times visit to edge e by all loops in \mathcal{L}_0 .

Rank: Set $S(x) = \frac{1}{2} \sum_{e \in E(x)} T(e) \Rightarrow S = V^{1/2}$

$$E \sim \left(\prod_i \frac{1}{(1+kx_i)^{-S(x_i)}} \right) = \left(\frac{1-\bar{A}_0}{1-\bar{A}_0 k} \right)^{1/2}$$

Remark: i) $t = (t(e))_{e \in E(G)} \in \mathbb{R}^{E(G)}$ $|t| = \sum_{e \in E(G)} t(e)$

ii) $S = (S(x))_{x \in V} \in \mathbb{R}^V$ $S(x) = \frac{1}{2} \sum_{e \in E(x)} t(e)$

iii) Set $\mathcal{P}(2n) = (2n)! / 2^n \cdot n!$ is the number of possible pairs of $\{1, 2, \dots, 2n\}$.

prop: Law of occupation field T

$$\tilde{\mathbb{P}}_1(T=t) = |G|^{-\frac{1}{2}} \cdot (2d)^{-|t|} \cdot \prod_{x \in V} \mathcal{P}(2S(x)) \cdot \prod_{e \in E(G)} \frac{1}{t(e)!}$$

cor. Condition on $T=t$, law of loop soup is:

at each site x , choose indeptly a pair of $2S(x)$ adjacent pipes uniformly among all choices.

Rank: When knowing T , the missing info.

on loop-soup is how to connect all these jumps along edges to each other.