Sı	Stormtroop3rs[C] ICPC Team Notebook 2017 umário		7.5 Seg Tree com Lazy 7.6 Union-Find 7.7 Treap 7.8 Seg Tree 2D 7.9 Polyce 7.10 KD2	29 29 31 31
	Template 1.1 Template C++	1 8 1	8.1 KMP	33 35
	2.1 Triângulo de Pascal 2.2 GCD-LCM 2.3 Bezout Theorem 2.4 Teorema Chinês dos Restos 2.5 Crivo de Eratóstenes 2.6 Divisores de N 2.7 Funções com Números Primos (Crivo, Fatoração, PHI, etc)	2 2 2 3 3 3 3	8.4 Suffix Array (Gugu) 8.5 Rolling Hash 8.6 Longest Commom Prefix with Hash 8.7 Minimum Lexicographic Rotation 8.8 Longest Palindrome (Manacher algorithm) 8.9 Autômato de Sufixos	37 37 37 37
	2.8 Exponenciação Modular Rápida 2.9 Exponenciação de Matriz 2.10 Floyd Cycle Detection 2.11 Brent Cycle Detection 2.12 Romberg's method - Calcula Integral (UFS2010) 2.13 Pollard's rho algorithm (UFS2010) 2.14 Miller-Rabin's algorithm (UFS2010) 2.15 Quantidade de dígitos de N! na base B 2.16 Quantiade de zeros a direita de N! na base B 2.17 Baby Step Giant Step	4 9 4 4 5 5 5 5 5 5	PD 9.1 Soma acumulada 2D 9.2 Knuth Optimization 9.3 Convex Hull Trick 9.4 Longest Increasing Subsequence 9.5 Kadane 1D 9.6 Kadane 2D 9.7 Merge Sort com num de Inversoes 9.8 Quick Sort	39 40 41 41 41 41
	2.18 Primos num intervalo	6 10	0 Miscelânia 10.1 Calendário	$\begin{array}{c} 42 \\ 42 \end{array}$
3	Geometria 2D 3.1 Geometria 2D Library	$\overset{\scriptscriptstyle{6}}{\overset{\scriptscriptstyle{6}}{}}}$	Template	
4	Polígonos 2D4.1Polígono 2D Library4.2Convex Hull4.3Minimum Enclosing Circle	9	.1 Template C++	
5	Geometria 3D 5.1 Geometria 3D Library	12 12	<pre>#include <bits stdc++.h=""> using namespace std;</bits></pre>	
6	Grafos 6.1 Topological Sort 6.2 Dijkstra 6.3 Floyd-Warshall 6.4 Bellman-Ford 6.5 Vértices de Articulação e Pontes 6.6 Tarjan 6.7 Kosaraju 6.8 2-Sat 6.9 LCA 6.10 Maximum Bipartite Matching 6.11 Hopcroft Karp - Maximum Bipartite Matching 6.12 Network Flow (lento) 6.13 Network Flow - Dinic 6.14 Min Cost Max Flow 6.15 Min Cost Max Flow (Stefano) 6.16 Tree Isomorphism 6.17 Stoer Wagner - Minimum Cut 6.18 Erdos Gallai 6.19 Stable Marriage 6.20 Hungarian Max Bipartite Matching with Cost	17 17 18 18 18 19 20 20 21 21 21 22 23 24	<pre>#define rep(i, a, b) for(int i=(a); i<(b); i++) #define pb push_back #define mk make_pair #define debug(x) cout<<line<<": "<<#x<<"="" "<<x<<endl;="" #define="" (c).begin(),="" (c).end()="" (c).resize(unique(all(c))-c.="" 3.1415926535897932384626433832795028841971="" all(c)="" begin());="" f="" first="" int="" ll;="" long="" pair<int,="" pi="" s="" second="" sort(all(c));="" typedef="" unique(c)=""> ii; const int INF = 0x3f3f3f3f3f; const double EPS = 1e-9; //(-1: x<y) (0:="" (1:="" x="">y) inline int cmp(double x, double y = 0, double tol = EPS){</y)></line<<":></pre>	
7	Estruturas de Dados 7.1 BIT		return ((x <= y+tol) ? (x+tol < y) ? -1:0:1); } // STL Utils * //Leitura de String:	

```
#include <sst.ream>
istringstream ins; // ins.str(s); ins >> a >> b; //IN
ostringstream outs; // outs << a << b; s = outs.str(); <math>//OUT
upper_bound(v.begin(), v.end(), key);//justafter the last
   element found
lower_bound(v.begin(), v.end(), key);//first element found, or,
    or case
//the element is not found, first position it could be inserted
    without
//violating the order (first element greater or equal key)
binary_search(v.begin(), v.end(), key);
next_permutation(v.begin(), v.end());
//Heap without priority change
priority_queue<ii, vector<ii>, greater<ii> > Q;
Q.push<ii(priority, key)>; //insert element
int u = Q.top().second; Q.pop(); //remove element
//Heap with priority change ( priorities in v[] )
struct comp{
  inline bool operator() (const int i, const int j) {
    return (v[i] != v[i]) ? (v[i] < v[i]) : (i < i);
set<int, comp> s;
s.erase(x); v[x]=new_priority; s.insert(x); //priority change
v[x] = priority; s.insert(x); //insert element
int u = *s.begin(); s.erase(*s.begin()); //remove element
remove(all(v), val); //removes all elements with val key
v.resize( remove(all(v), val) - v.begin()); // remove and
   resize the vector
reverse(all(v));
find(all(v), key); //returns first iterator to key
find_if(all(v), pred);//where pred is bool pred(T key)
//String Functions ( s is a string )
s.find(string pattern, int pos=0); // find a pattern in string
starting from position zero. returns index of first element if
pattern is found and a large number if not found
// substring de s comecando em pos e com tamanho tam
s.substr(int pos, int tam);
int main(){
  return 0;
```

2 Numerical algorithms

2.1 Triângulo de Pascal

```
/*Calcula os numeros binomiais (N,K) = N!/(K!(N-K)!)
```

```
(N,K) representa o numero de maneiras de criar um subconjunto
    de tamanho
K dado um conjunto de tamanho N. A ordem dos elementos nao
   importa.*/
const int MAXN = 50;
long long C[MAXN][MAXN];
void calc pascal() {
  memset(C, 0, sizeof(C));
  for (int i = 0; i < MAXN; ++i) {</pre>
    C[i][0] = C[i][i] = 1;
    for (int j = 1; j < i; ++j)
      C[i][j] = C[i-1][j-1] + C[i-1][j];
//Pascal triangle elements:
C(33, 16) = 1.166.803.110 [int limit]
C(34, 17) = 2.333.606.220 [unsigned int limit]
C(66, 33) = 7.219.428.434.016.265.740 [int64_t limit]
C(67, 33) = 14.226.520.737.620.288.370 [uint64_t limit]
//Fatorial
12! = 479.001.600 [(unsigned) int limit]
20! = 2.432.902.008.176.640.000 [(unsigned) int64 t limit ]
```

2.2 GCD-LCM

```
11 gcd(l1 a, l1 b){return b==0 ? a:gcd(b, a%b);}
11 lcm(l1 a, l1 b){
   if (a and b) return abs(a)/gcd(a, b)*abs(y);
   else return abs(a | y);
}
```

2.3 Bezout Theorem

```
/* Aplicacoes:
Determinar a solucao da equacao a*x+b*y = gcd(a, b), onde a e b
    sao dois
numeros inteiros naturais.
Como chamar: egcd(a, b), Retorno: a tupla (gcd(a, b), x, y)
Determina tambem o Inverso Modular */
struct Triple{
 int d, x, y; // long long ?
 Triple (int q, int w, int e):d(q), x(w), y(e) {}
Triple egcd (int a, int b) {
 if (!b) return Triple(a, 1, 0);
 Triple q = egcd(b, a%b);
 return Triple(q.d, q.y, q.x - a/b * q.y);
int invMod (int a, int n) {
 Triple t = egcd (a, n);
 if (t.d > 1) return 0;
```

```
return (t.x%n+n)%n;
}
```

2.4 Teorema Chinês dos Restos

```
#define MAXN 1000
/* Aplicacoes:
Determinar x tal que x = a[i] (mod m[i]).
Exemplo: Para a[] = \{1, 2, 3\} \in m[] = \{5, 6, 7\}
x = 206
Como chamar a funcao:
  1) Chamar a funcao crt()
Depende do Inverso Modular (Bezout ou Fermat)
Testado: URI 1831
*/
int n;
int a[MAXN], m[MAXN];
int crt(){
 int M = 1, x = 0;
  for (int i = 0; i < n; ++i) M *= m[i];</pre>
  for(int i = 0; i < n; ++i)
    x += a[i] * invMod(M/m[i], m[i]) * (M/m[i]);
  return (((x%M)+M)%M);
```

2.5 Crivo de Eratóstenes

```
bitset<10000005> bs; vector<int> primos;
void crivo(ll limite = 10000000LL){ // calcula primos ate
    limite
  primos.clear();
 bs.set();
  bs[0] = bs[1] = 0;
  for(ll i = 2; i <= limite; i++)</pre>
    if (bs[i]) {
      for(ll j = i*i; j <= limite; j += i)</pre>
       bs[j] = 0;
      primos.push_back(i);
    }
bool isPrime(ll N, ll limite) {
  if (N <= limite) return bs[N];</pre>
  for(int i = 0; i < (int) primos.size(); i++)</pre>
    if (N%primos[i] == 0) return false;
  return true; }
```

2.6 Divisores de N

```
// Retorna todos os divisores naturais de N em O(sqrt(N)) vector<ll> divisores(ll N) {
```

```
vector<ll> divisor; divisor.clear();
for(ll div = 1, k; div*div <= N; div++) {
   if (N%div == 0) {
      divisor.pb(div);
      k = N/div;
      if (k != div) divisor.pb(k);
   }
}
// sort(divisor.begin(), divisor.end()); // caso precise ordenado
return divisor;
}</pre>
```

2.7 Funções com Números Primos (Crivo, Fatoração, PHI, etc)

```
map<int, int> factors;
//Encontra os fatores primos de N .: N = p1^e1 * ... *pi^ei
// factors armazena em first o fator primo e em segundo seu
   expoente
void primeFactors(ll N) {
  factors.clear();
  while (N\%2==0) ++factors[2], N>>=1;
  for(11 PF=3; PF*PF <= N; PF+=2) {</pre>
    while (N%PF == 0)
      N /= PF, factors[PF]++;
  if (N > 1) factors [N] = 1;
//Funcoess derivadas dos numeros primos
void NumberTheory(ll N) {
  primeFactors(N);
  map<int, int> :: iterator f;//<fator primo, expoente>
  11 Totient=N, numDiv=1, sumDiv=1, numDiffPF, sumPF=0;
  numDiffPF = factors.size();
  for(f = factors.begin(); f != factors.end(); f++){
   ll PF = f \rightarrow fst, power = f \rightarrow snd;
   Totient -= Totient/PF;
    numDiv \star = (power+1);
    sumDiv \star = ((11) pow((double) PF, power+1.0)-1)/(PF-1);
    sumPF += PF;
  printf ("Totiente/Euler-Phi de N = %lld\n", Totient);
  printf ("qt de divisores de N = %lld\n", numDiv);
  printf ("soma dos divisores de N = %lld\n", sumDiv);
  printf ("qt de fatores primos distintos = %lld\n", numDiffPF)
  printf ("soma dos fatores primos = %lld\n", sumPF);
//Calcula Euler Phi para cada valor do intervalo [1, N]
#define MM 1000010
int phi[MM];
```

```
void crivo_euler_phi(int N) {
  for (int i = 1; i <= N; i++)</pre>
    phi[i] = i;
  for(int i = 2; i <= N; i++)
    if (phi[i] == i) {
      for(int k = i; k <= N; k+=i)
        phi[k] = (phi[k]/i) * (i-1);
}
//Qtde de fatores primos distintos de cada valor do range [2,
    MAX N1
#define MAX_N 10000000
int NDPF[MAX_N]; //
void NumDiffPrimeFactors() {
 memset(NDPF, 0, sizeof NDPF);
  for(int i = 2; i < MAX_N; i++)</pre>
    if (NDPF[i] == 0)
      for (int j = i; j < MAX_N; j += i)</pre>
        NDPF[j]++;
```

2.8 Exponenciação Modular Rápida

```
/* Calcula (B^P) %MOD em O(log(P))
Calcula o inverso modular de b(modulo mod) se mod for primo.
    basta
fazer invB = fastpow(b, mod-2, mod); */
int fastpow(int b, int p, int mod) {
    int ret = 1;
    for(ll pow = b; p > 0; p >>= 1, pow = (pow*pow) %mod)
        if(p & 1) ret = (ret*pow) %mod;
    return ret;
}
```

2.9 Exponenciação de Matriz

```
const int M = 2; ll mod = le9+7;
int sz = 2;
ll mat[M][M], ans[M][M], tmp[M][M];

void mult(ll a[][M], ll b[][M]) {
  rep(i, 0, sz)
    rep(j, 0, sz) {
     tmp[i][j] = 0;
     rep(k, 0, sz)
        tmp[i][j] += a[i][k]*b[k][j];
     tmp[i][j] %= mod;
    }
  memcpy(a, tmp, sizeof tmp);
}
// ans = mat^n
```

```
void fastExp(ll ans[][M], ll n) {
    // inicializar mat
    mat[0][0] = mat[0][1] = mat[1][0] = 1;
    mat[1][1] = 0;
    // matriz identidade
    rep(i,0,sz) rep(j,0,sz) ans[i][j] = (i==j);
    while(n) {
        if (n & 1) mult(ans, mat);
        n >>= 1;
        mult(mat, mat);
    }
        //n-\'esino termo de fibonacci
    //cout << ans[1][0]*fib(1) + ans[1][1] * fib(0) << "\n";
}</pre>
```

2.10 Floyd Cycle Detection

```
// faz o mesmo que brent-cycle
ii floyd_cycle(int x) {
  int t = f(x), h = f(f(x)), start = 0, lenght = 1;
  while (t != h) {t = f(t); h = f(f(h)); }

h = x;
  while (t != h) { t = f(t); h = f(h); start++; }

t = start; h = f(t);
  while (t != h) { h = f(h); lenght++; }
  return ii(start, lenght);
}
```

2.11 Brent Cycle Detection

```
/*Dado uma sequencia f(x0), f(f(x0)), ..., f(f(...f(x0))), ela
ciclica. A funcao retorna o tamanho do ciclo e o valor que o
   inicia */
ii brent_cycle(int x) {
    int p = 1, length = 1, t = x, start = 0;
    int h = f(x);
    while (t != h) {
       if (p == length) { t = h; p \star= 2; length = 0; }
        h = f(h);
        ++length;
    t = h = x;
    for (int i = length; i != 0; --i) h = f(h);
    while (t != h) {
       t = f(t); h = f(h);
        ++start;
    return ii(start, length);
```

2.12 Romberg's method - Calcula Integral (UFS2010)

```
//Calcula a integral de f[a, b]
typedef long double ld;

ld f(double x) {
    // return f(x)
}

ld romberg(ld a, ld b) {
    ld R[16][16], div = (b-a)/2;
    R[0][0] = div*(f(a)+f(b));
    for(int n = 1; n <= 15; n++, div /= 2) {
        R[n][0] = R[n-1][0]/2;
        for(ld sample = a+div; sample<b; sample += 2*div)
        R[n][0] += div*f(a+sample);
    }
    for(int m = 1; m <= 15; m++)
        for(int n = m; n <= 15; n++)
        R[n][m] = R[n][m-1]+1/(pow(4, m)-1)*(R[n][m-1]-R[n-1][m-1])
        ;
    return R[15][15];
}</pre>
```

2.13 Pollard's rho algorithm (UFS2010)

```
//Retorna um fator primo de N

ll f(ll val) {return (val*val+pollard_r) %pollard_n; }

ll myabs(ll a) {return a >= 0 ? a:-a; }

ll pollard(ll n) {
    srand(unsigned(time(0)));
    pollard_n = n;
    long long d = 1;
    do {
        d = 1;
        pollard_r = rand() %n;
        long long x = 2, y = 2;
        while(d == 1)
            x = f(x), y = f(f(y)), d = __gcd(myabs(x-y), n);
    } while(d == n);
    return d;
}
```

2.14 Miller-Rabin's algorithm (UFS2010)

```
// Teste de primalidade
bool miller_rabin(int n, int base) {
  if(n <= 1) return false;
  if(n%2 == 0) return n == 2;
  int s = 0, d = n-1;</pre>
```

```
while(d%2 == 0) d /= 2, ++s;
int base_d = fastpow(base, d, n);
if(base_d == 1) return true;
int base_2r = base_d;
for(int i = 0; i < s; ++i) {
   if(base_2r == 1) return false;
   if(base_2r == n-1) return true;
   base_2r = (long long)base_2r*base_2r%n;
}
return false;
}
bool isprime(int n) {
   if(n == 2 || n == 7 || n == 61) return true;
   return miller_rabin(n, 2) && miller_rabin(n, 7) &&
        miller_rabin(n, 61);
}</pre>
```

2.15 Quantidade de dígitos de N! na base B

```
int NumOfDigitsInFactorial(int N, int B) {
   double logFatN = 0;
   for (int i = 1; i <= N; i++)
     logFatN += log((double)i);
   int nd = floor(logFatN / log((double)B)) + 1;
   return nd;
}</pre>
```

2.16 Quantiade de zeros a direita de N! na base B

```
int NumOfTrailingZeros(int N, int B) {
  int nfact = fatora(B);
  int zeros = INF;
    // para cada fator de B, aux representa qtas vezes
    // fator[i]^expoente[i] aparece na representacao de N!
  for (int i=0; i < nfact; i++) {
    int soma = 0;
    int NN = N;
    while (NN) {
        soma += NN/fator[i];
        NN /= fator[i];
    }
    int aux = soma / expoente[i];
    zeros = min(zeros, aux);
  }
  return zeros;
}</pre>
```

2.17 Baby Step Giant Step

```
/*Determinar o menor E tal que B^E = N \pmod{P}, -1 se for impossivel
```

```
Requer: Bezout Theorem*/
11 bsgs(ll b, ll n, ll p) {
    if (n == 1) return 0;
    map <ll, int> table;
    ll m = sqrt(p) + 1, pot = 1, pot2 = 1;
    for (int j = 0; j < m; ++j) {
        if (pot == n) return j;
        table[(n * invMod(pot, p)) % p] = j;
        pot = (pot * b) % p;
    }
    for (int i = 0; i < m; ++i) {
        if (table.find(pot2) != table.end())
            return i * m + table[pot2];
        pot2 = (pot * pot2) % p;
    }
    return -1;
}</pre>
```

2.18 Primos num intervalo

```
//Encontra os primos no intervalo [n,m]
vector<int> ret;
void primesBetween(int n, int m) {
  ret.clear();
  vector<int> primes(m-n+1);
  for (int i = 0; i < m - n + 1; ++i)
    primes[i] = 0;
  for (int p = 2; p*p <= m; ++p) {
    int less = (n / p)*p;
   for (int j = less; j <= m; j += p)
      if (j != p \&\& j >= n)
        primes[j - n] = 1;
  for (int i = 0; i < m - n + 1; ++i) {
    if (primes[i] == 0 && n+i != 1) {
      ret.push_back(n+i);
  }
```

2.19 FFT

```
typedef complex<double> comp;
const double PI = 2*acos(0);
const int MAX_N = 1 << 20;
int rev[MAX_N];
comp roots[MAX_N];
void preCalc(int N, int BASE){</pre>
```

```
for (int i = 1; i < N; ++i) rev[i] = (rev[i >> 1] >> 1) + ((i)
      \& 1) << (BASE - 1);
 int NN = N \gg 1;
 roots[NN] = comp(1, 0);
 roots[NN + 1] = comp(cos(2 * PI / N), sin(2 * PI / N));
 for (int i = 2; i < NN; ++i) roots[NN + i] = roots[NN + i - 1
     | * roots[NN + 1];
 for (int i = NN - 1; i > 0; --i) roots[i] = roots[2 * i];
void fft( vector<comp> &a, bool invert ) {
 int N = a.size();
 if (invert) rep(i, 0, N) a[i] = conj(a[i]);
  rep(i, 0, N) if (i < rev[i]) swap(a[i], a[rev[i]]);
 for (int k = 1; k < N; k *= 2) {
   for (int i = 0; i < N; i += 2 * k) {
      rep(j, 0, k)
        comp B = a[i + j + k] * roots[k + j];
        a[i + j + k] = a[i + j] - B;
       a[i + j] = a[i + j] + B;
 if (invert) rep(i, 0, a.size()) a[i] /= N;
vector<comp> multiply_real( vector<comp> a, vector<comp> b,
   vector<comp> c ) {
  int n = a.size();
 int m = b.size();
 int base = 0, N = 1;
 while (N < n+m-1) base++, N <<= 1;
 preCalc(N, base);
 a.resize(N, comp(0, 0));
 c.resize(N);
  rep(i, 0, b.size()) a[i] = comp(real(a[i]), real(b[i]));
  fft(a, 0);
 rep(i, 0, N) {
   int j = (N - i) & (N - 1);
   c[i] = (a[i] * a[i] - conj(a[j] * a[j])) * comp(0, -0.25);
 fft(c, 1);
  return c;
```

3 Geometria 2D

3.1 Geometria 2D Library

```
const double EPS = 1e-9;
```

```
inline int cmp( double x, double y = 0, double tol = EPS) {
  return ( (x \le y + tol) ? (x + tol < y) ? -1 : 0 : 1);
int sgn(11 x) \{ return x < 0 ? -1 : x > 0; \}
struct point{
  double x, v;
  point (double x=0, double y=0): x(x), y(y) {}
  point operator + (point q) { return point (x+q.x, y+q.y);}
  point operator - (point q) { return point (x-q.x, y-q.y); }
  point operator *(double t) { return point(x*t, y*t); }
  point operator / (double t) { return point(x/t, y/t); }
  point operator / (point q) { // divisao de num complexos
    return point (x*q.x+y*q.y, y*q.x-x*q.y)/q.norma2();
  double norma2() { return (*this) * (*this); } // |u|^2
  double operator *(point q) { return x*q.x+y*q.y; }
  double operator %(point q) { return x*q.y-y*q.x; }
  int cmp(point q) const{
    if(int t = ::cmp(x, q.x)) return t;
    return ::cmp(y, q.y);
  bool operator == (point q) const{return cmp(q) == 0;};
  bool operator != (point q) const{return cmp(q) != 0;};
  bool operator < (point q) const{return cmp(q) < 0;};</pre>
ostream & operator << (ostream & os, const point &p) {
  os << "(" << p.x << "," << p.y << ")";
#define vec(a, b) (b-a)
typedef vector<point> polygon;
// Rotaciona um ponto em relacao a origem
point RotateCCW90 (point p) { return point (-p.y, p.x); }
point RotateCW90(point p) { return point(p.y, -p.x); }
// rotaciona um ponto P em A graus no sentido anti-horario.
// para rotacionar no sentido horario, basta A ser negativo
point RotateCCW(point p, double A) {
 A = (A/180.0) *acos(-1.0); // convertendo para radianos
  return point (p.x*cos(A)-p.y*sin(A), p.x*sin(A)+p.y*cos(A));
// tamanho de um vetor
double abs(point u) {
  return sqrt(u*u);
/*Decide se P->Q->R forma um giro a esquerda, direita ou se sao
colineares. Retorna 1 se R estiver a esquerda da reta PO. -1 se
estiver a direta e 0 se forem colineares.*/
int ccw(point p, point q, point r) {
```

```
return cmp(vec(p, q)%(vec(p, r)));
// Projeta o vetor A sobre a direcao do vetor B
point project(point a, point b) {
 return b*((a*b)/(b*b));
// Retorna a projecao do ponto P sobre reta definida por [A,B]
point projectPointLine(point p, point a, point b) {
 return p + project(p-a, b-a);
// Retorna o angulo entre os vetores AB e AC
double arg(point B, point A, point C) {
 point u = vec(A, B), v = vec(A, C);
 return atan2(u%v, u*v);
// verifica se x esta entre [l,r] *Tente evitar usar double!
bool is_in(double l, double r, double x) {
 if (1 > r) swap(1, r);
 return (1 <= x + EPS) && (x - EPS <= r);</pre>
// Determina se P esta entre o segmento fechado [A,B],
   inclusive
bool between(point p, point a, point b) {
 if (ccw(p,a,b) != 0) return false;
 else return is_in(a.x, b.x, p.x) && is_in(a.y, b.y, p.x);
// (Opcao2) Determina se P esta entre o segmento fechado [A,B],
    inclusive
// Se vc nao entendeu a anterior ou se acha que ela pode ser o
// use essa. Se nao der certo, tente mudar a precisao usada em
   cmp() para 1e-15
bool between2 (point p, point a, point b) {
 return cmp (abs (p-a) + abs (p-b), abs (a-b)) == 0;
// Distancia de ponto P para reta que passa por [A,B]
// armazena em C (por ref) o ponto projecao de P na reta
double distancePointLine(point p, point a, point b, point& c) {
 c = projectPointLine(p,a,b);
 return fabs((b-a)%(p-a))/abs(a-b); // or abs(p-c);
// Distancia de ponto P ao segmento [A,B]
// armazena em C (por ref) o ponto de projecao de P em [A,B]
// se este ponto sair do segmento, eh retornado o mais proximo
double distancePointSeg(point p, point a, point b, point& c) {
```

```
if ((b-a) * (p-a) <= 0) { c = a; return abs(a-p); }
  if ((a-b) * (p-b) <= 0) { c = b; return abs(b-p); }
 c = projectPointLine(p,a,b);
  return fabs((b-a)%(p-a))/abs(a-b); // or abs(p-c);
// Distancia entre (x,y,z) e plano ax+by+cz=d
double distancePointPlane (double x, double y, double z, double
   a, double b, double c, double d) {
  return fabs(a*x+b*y+c*z-d)/sqrt(a*a+b*b+c*c);
// Determina se os segmentos [A, B] e [C, D] se tocam
bool seg intersect(point a, point b, point c, point d) {
  int d1, d2, d3, d4;
 d1 = ccw(c, a, d); d2 = ccw(c, b, d);
  d3 = ccw(a, c, b); d4 = ccw(a, d, b);
 if (d1*d2 == -1 && d3*d4 == -1) return true;
 if (d1 == 0 && between(c, a, d)) return true;
 if (d2 == 0 && between(c, b, d)) return true;
 if (d3 == 0 && between(a, c, b)) return true;
  if (d4 == 0 && between(a, d, b)) return true;
  return false:
// Encontra a interseccao das retas (p-q) e (r-s) assumindo que
    existe
// apenas 1 intereccao
// Se for entre segmentos, verificar se interseptam primeiro
point line_intersect(point p, point q, point r, point s) {
 point a = q - p, b = s - r, c = point(p % q, r % s);
  return point(point(a.x,b.x)%c, point(a.y,b.y)%c)/(a%b);
// determine if lines from a to b and c to d are parallel or
   collinear
bool LinesParallel (point a, point b, point c, point d) { // Nao
  return fabs((b-a) % (c-d)) < EPS;</pre>
bool LinesCollinear(point a, point b, point c, point d) { // Nao
    testado
  return LinesParallel(a, b, c, d)
      && fabs((a-b) % (a-c)) < EPS
      && fabs((c-d) % (c-a)) < EPS;
// Retas, equacao na forma: ax+by+c = 0; com (b=1)
struct line{
```

```
double a, b, c;
 line(){}
 line (double m, point p) { // coeficiente e ponto sobre reta
   if (cmp(m) == 0) a = 0;
   else a = -m;
   b = 1;
    c = m*p.x-p.y;
 line(point p1, point p2) { // dois pontos que estao na reta
   if (fabs(p1.x-p2.x) < EPS){
      a = 1; b = 0; c = -p1.x;
     a = -(double) (p1.y-p2.y) / (p1.x-p2.x);
      c = -(double)(a*p1.x)-p1.v;
 bool parallelTo(line s){//verdadeiro se forem paralelas
   return cmp (a, s.a) == 0 \&\& cmp(b, s.b) == 0;
 bool sameAs(line s){//verdadeiro se forem retas coincidentes
   return this->parallelTo(s) && cmp(c, s.c) == 0;
 point intersectPoint(line s){
     point p;
     p.x = (s.b*c-b*s.c) / (s.a*b-a*s.b);
       if (fabs(b) > EPS) p.v = -(a*p.x+c)/b;
                         p.y = -(s.a*p.x+s.c)/s.b;
      else
      return p;
};
// Menor distancia entre dois pontos numa esfera de raio r
// lat = [-90,90]; long = [-180,180]
double spherical_distance(double lt1, double lo1, double lt2,
   double lo2, double r) {
  double pi = acos(-1);
  double a = pi*(lt1/180.0), b = pi*(lt2/180.0);
  double c = pi*((lo2-lo1)/180.0);
  return r*acos(sin(a)*sin(b) + cos(a)*cos(b)*cos(c));
// Heron's formula - area do triangulo(a,b,c) -1 se nao existe
double area heron(double a, double b, double c) {
 if (a < b) swap(a, b);
 if (a < c) swap(a, c);
 if (b < c) swap(b, c);
 if (a > b+c) return -1;
 return sqrt((a+(b+c))*(c-(a-b))*(c+(a-b))*(a+(b-c))/16.0);
/*Dado dois pontos (A, B) de uma circunferencia e seu raio R,
possivel obter seus possiveis centros (C1 e C2). Para obter o
```

```
outro centro, basta inverter os paramentros */
bool circle2PtsRad(point a, point b, double r, point &c) {
  point aux = a - b;
  double d = aux * aux;
  double det = r * r/d - 0.25;
  if (det < 0.0) return false;</pre>
  double h = sqrt(det);
  c.x = (a.x + b.x) * 0.5 + (a.y - b.y) * h;
  c.y = (a.y + b.y) * 0.5 + (b.x - a.x) * h;
  return true;
//***[Inicio] Funcoes que usam numeros complexos para pontos***
typedef complex<double> cxpt;
struct circle {
  cxpt c; double r;
  circle(cxpt c, double r) : c(c),r(r){}
  circle(){}
double cross(const cxpt &a, const cxpt &b) {
  return imag(conj(a)*b);
double dot(const cxpt &a, const cxpt &b) {
  return real(conj(a)*b);
// Area da interseccao de dois circulos
double circ_inter_area(circle &a, circle &b) {
  double d = abs(b.c-a.c);
  if (d <= (b.r - a.r)) return a.r*a.r*M_PI;</pre>
  if (d <= (a.r - b.r)) return b.r*b.r*M_PI;</pre>
  if (d >= a.r + b.r) return 0;
  double A = acos((a.r*a.r+d*d-b.r*b.r)/(2*a.r*d));
  double B = acos((b.r*b.r+d*d-a.r*a.r)/(2*b.r*d));
  return a.r*a.r*(A-0.5*sin(2*A))+b.r*b.r*(B-0.5*sin(2*B));
// Pontos de interseccao de dois circulos
// Intersects two circles and intersection points are in 'inter
// -1-> outside, 0-> inside, 1-> tangent, 2-> 2 intersections
int circ_circ_inter(circle &a, circle &b, vector<cxpt> &inter)
  double d2 = norm(b.c-a.c), rS = a.r+b.r, rD = a.r-b.r;
  if (d2 > rS*rS) return -1;
  if (d2 < rD*rD) return 0;</pre>
  double ca = 0.5*(1 + rS*rD/d2);
  cxpt z = cxpt(ca, sqrt((a.r*a.r/d2)-ca*ca));
  inter.push_back(a.c + (b.c-a.c)*z);
  if(abs(z.imag())>EPS)
    inter.push_back(a.c + (b.c-a.c)*conj(z));
  return inter.size();
```

```
// Line-circle intersection
// Intersects (infinite) line a-b with circle c
// Intersection points are in 'inter'
// 0 -> no intersection, 1 -> tangent, 2 -> two intersections
int line_circ_inter(cxpt a, cxpt b, circle c, vector<cxpt> &
   inter) {
   c.c -= a; b -= a;
   cxpt m = b*real(c.c/b);
    double d2 = norm(m-c.c);
    if (d2 > c.r*c.r) return 0;
    double l = sqrt((c.r*c.r-d2)/norm(b));
    inter.push_back(a + m + 1*b);
    if (abs(1)>EPS)
        inter.push_back(a + m - 1*b);
    return inter.size();
//***[FIM] Funcoes que usam numeros complexos para pontos***
```

4 Polígonos 2D

4.1 Polígono 2D Library

```
/*Poligono eh representado como um array de pontos T[i] sao os
  vertices do poligono. Existe uma aresta que conecta T[i] com
     T[i+1],
 e T[size-1] com T[0]. Logo assume-se que T[0] != T[size-1]
Poligono simples: Aquele em que as arestas nao se interceptam
Convexo: O angulo interno de T[i] com T[i-1] e T[i+1] <= 180
Concavo: Existe algum i que nao satisfaz a condicao anteriro
/* Retorna a area com sinal de um poligono T
   Se area > 0, T esta listado na ordem CCW */
double signedArea(polygon& T) {
 double area = 0;
 int n = T.size();
 if (n < 3) return 0;
 rep(i, 0, n)
   area += T[i] % T[(i+1)%n];
 return (area/2.0);
/* Retorna a area de um poligono T
   (pode ser concavo ou convexo) em O(N) */
double poly_area(polygon& T) {
 return fabs(signedArea(T));
/* Retorna a centroide de um poligono T em O(N) */
point centroide(polygon &T){
 int n = T.size();
```

```
double sgnArea = signedArea(T);
  point c = point(0,0);
  rep(i, 0, n){
   int k = (i+1) %n;
    c = c + (T[i] + T[k]) * (T[i] * T[k]);
  c = c / (sgnArea * 6.0);
  return c;
/* Retorna o perimetro do poligono T. (pode n funcionar como
   esperado se o poligono for uma linha reta (caso degenerado))
       */
double poly perimeter(polygon& T) {
  double perimeter = 0;
  int n = T.size();
  if (n < 3) return 0;
  rep(i, 0, n)
    perimeter += abs(T[i] - T[(i+1)%n]);
  return perimeter;
// tests whether or not a given polygon (in CW or CCW order) is
     simple
bool isSimple(const polygon &p) { // nao testado
  for (int i = 0; i < p.size(); i++) {</pre>
    for (int k = i+1; k < p.size(); k++) {
      int j = (i+1) % p.size();
      int 1 = (k+1) % p.size();
      if (i == 1 || j == k) continue;
      if (seg_intersect(p[i], p[j], p[k], p[l]))
        return false;
  return true;
//Retorna True se T for convexo. O(N)
bool isConvex(polygon& T) {
  int n = T.size();
  if (n < 3) return false;</pre>
  int giro = 0;
  rep(i, 0, n){ // encontra um giro valido
   int t = ccw(T[i],T[(i+1)%n], T[(i+2)%n]);
    if (t != 0) giro = t;
  if (giro == 0) return false; //todos pontos sao colineares
```

```
rep(i, 0, n){
    int t = ccw(T[i], T[(i+1)%n], T[(i+2)%n]);
    if (t != 0 && t != giro) return false;
  return true;
// Determina se P pertence a T, funciona para convexo ou
    concavo
// -1 borda, 0 fora, 1 dentro. O(N)
int in_poly(point p, polygon& T) {
  double a = 0; int N = T.size();
  rep(i, 0, N)
    if (between(p, T[i], T[(i+1)%N])) return -1;
    a += arg(T[i], p, T[(i+1)%N]);
  return cmp(a) != 0;
//determina se P pertence a B, funciona APENAS para convexo
bool PointInConvexPolygon(point P, polygon &B) {
  int ini = 1, fim = B.size()-2, mid, pos = -1;
  int giro = -1; // sentido horario
  while(ini<=fim) {</pre>
    mid = (ini+fim)/2;
    int aux = ccw(B[0], B[mid], P);
    if(aux == giro){
     pos = mid;
     ini = mid+1;
    }else{
      fim = mid-1;
  if(pos == -1) return false;
  if ( ccw(B[0], B[pos], P)!=qiro*-1 &&
      ccw(B[0], B[pos+1], P)!=qiro \&\&
      ccw(B[pos], B[pos+1], P) == qiro) // qiro // 0 na borda
    return true;
  return false;
// Determina o poligono interseccao de P e Q
// P e Q devem estar orientados anti-horario.
polygon poly_intersect(polygon& P, polygon& Q) {
  int m = Q.size(), n = P.size();
  int a = 0, b = 0, aa = 0, ba = 0, inflag = 0;
  polygon R;
  while ((aa<n || ba<m) && aa<2*n && ba<2*m) {</pre>
    point p1 = P[a], p2 = P[(a+1)%n], q1 = Q[b], q2 = Q[(b+1)%m]
    point A = p^2 - p^1, B = q^2 - q^1;
    int cross=cmp(A%B), ha=ccw(p2, q2, p1),
       hb=ccw(q^2, p^2, q^1);
    if (cross==0 \&\& ccw(p1, q1, p2)==0 \&\& cmp(A*B)<0) {
```

```
if (between(q1, p1, p2)) R.push_back(q1);
    if (between(q2, p1, p2)) R.push_back(q2);
    if (between (p1, q1, q2)) R.push_back (p1);
    if (between(p2, q1, q2)) R.push_back(p2);
    if (R.size() < 2) return polygon();</pre>
    inflag = 1; break;
  else if (cross!=0 && seg_intersect(p1, p2, q1, q2)){
    if (inflag == 0) aa = ba = 0;
    R.push_back(line_intersect(p1, p2, q1, q2));
    inflag = (hb > 0) ? 1:-1;
  if (cross==0 && hb<0 && ha<0) return R;</pre>
  bool t = cross==0 && hb==0 && ha==0;
  if (t?(inflag==1):(cross>=0)?(ha<=0):(hb>0)){
    if (inflag == -1) R.push_back(q2);
    ba++; b++; b %= m;
  else {
    if (inflag == 1) R.push_back(p2);
    aa++; a++; a %= n;
if (inflag == 0) {
  if (in_poly(P[0], Q)) return P;
  if (in_poly(Q[0], P)) return Q;
R.erase(unique(all(R)), R.end());
if (R.size() > 1 && R.front() == R.back()) R.pop_back();
return R;
```

4.2 Convex Hull

```
/*Encontra o convex hull de um conjunto de pontos;
pivot: Ponto base para a criacao do convex hull;
radial_lt(): Ordena os pontos em sentido anti-horario (ccw).
Input: Conjunto de pontos 2D;
Output: Conjunto de pontos do convex hull em sentido ccw;
(1) Se for preciso manter pontos colineares na borda do convex
hull, essa parte evita que eles sejam removidos;
*/
point pivot;
bool radial_lt(point a, point b) {
  int R = ccw(pivot, a, b); // verifica se A esta antes de B
  if (R == 0) // pivot, a, b sao colineares
    return (pivot-a) * (pivot-b) * (pivot-b);
  else
    return (R == 1); // 1 se A vem antes de B, sentido ccw
vector<point> convexhull(vector<point> &T) {
```

```
// Remove pontos duplicadados
sort(T.begin(), T.end()); //ordena por x e por y
T.resize( unique( T.begin(), T.end() ) - T.begin() );
int tam = 0, k, n = T.size();
vector<point> U; // convex hull
int P0 = min_element(T.begin(), T.end() ) - T.begin();
//nesse caso, pivot = ponto com menor x, depois menor y
pivot = T[P0];
swap(T[0], T[P0]);
sort(++T.begin(), T.end(), radial_lt);
/*(1)*/for(k=n-2; k>=0 && ccw(T[0],T[n-1],T[k])==0; k--);
reverse((k+1)+all(T));
// troque <= para < para manter pontos colineares</pre>
for (int i = 0; i < n; i++) {</pre>
  while (tam > 1 && ccw(U[tam-2], U[tam-1], T[i]) <= 0)</pre>
    U.pop_back(), tam--;
 U.pb(T[i]); tam++;
return U;
```

4.3 Minimum Enclosing Circle

```
//Finds a circle of the minimum area enclosing a 2D point set.
typedef pair<point, double> circle; // {ponto, raio}
bool in_circle(circle C, point p) { // ponto dentro de circulo?
  return cmp(abs(p-C.first), C.second) <= 0;</pre>
// menor circulo que engloba o triangulo (P,Q,R)
point circumcenter(point p, point q, point r) {
 point a = p-r, b = q-r, c, ret;
 c = point(a*(p+r)/2, b*(q+r)/2);
 ret=point(c%point(a.y, b.y), point(a.x, b.x)%c)/(a%b);
 return ret;
circle spanning_circle(vector<point>& T) {
  int n = T.size();
  random_shuffle(all(T));
  circle C(point(), -INF);
  rep(i, 0, n) if(!in_circle(C, T[i])){
   C = circle(T[i], 0);
   rep(j, 0, i) if (!in_circle(C, T[j])){
     C = circle((T[i]+T[j])/2, abs(T[i]-T[j])/2);
     rep(k, 0, j) if (!in_circle(C, T[k])){
       point 0 = circumcenter(T[i], T[j], T[k]);
       C = circle(0, abs(0-T[k]));
```

```
return C;
}
```

5 Geometria 3D

5.1 Geometria 3D Library

```
// Tem algumas funcoes que possuem implementacoes distintas,
// mesmo objetivo. Maioria nao foi testada.
#define LINE 0
#define SEGMENT 1
#define RAY 2
int sqn(double x) {
  return (x > EPS) - (x < -EPS);
#define vec(ini, fim) (fim - ini)
struct PT{
  double x, y, z;
  PT() \{x = y = z = 0; \}
  PT (double x, double y, double z):x(x), y(y), z(z) {}
  PT operator + (PT q) {return PT(x+q.x,y+q.y,z+q.z);}
  PT operator - (PT q) {return PT(x-q.x,y-q.y,z-q.z);}
  PT operator * (double d) { return PT(x*d, y*d, z*d); }
  PT operator / (double d) { return PT(x/d, y/d, z/d); }
  double dist2() const {
    return x*x+y*y+z*z;
  double dist() const{
    return sqrt(dist2());
  bool operator == (const PT& a) const{
    return fabs(x - a.x) < EPS && fabs(y - a.y) < EPS && fabs(z
         - a.z) < EPS;
};
double dot (PT A, PT B) { // produto escalar
  return A.x*B.x + A.y*B.y + A.z*B.z;
PT cross(PT A, PT B) { // produto vetorial
  return PT (A.y*B.z-A.z*B.y, A.z*B.x-A.x*B.z, A.x*B.y-A.y*B.x
     );
inline double det(double a, double b, double c, double d) {
  return a*d - b*c;
```

```
inline double det (double a11, double a12, double a13, double a21,
    double a22, double a23, double a31, double a32, double a33) {
  return a11*det(a22,a23,a32,a33) - a12*det(a21,a23,a31,a33) +
      a13*det(a21,a22,a31,a32);
inline double det(const PT& a, const PT& b, const PT& c) {
  return det (a.x, a.y, a.z, b.x, b.y, b.z, c.x, c.y, c.z);
// tamanho do vetor A
double norma(PT A) {
  return sqrt(dot(A, A));
// distancia^2 de (a->b)
double distSq(PT a, PT b) {
  return dot(a-b, a-b);
// Projeta vetor A sobre o vetor B
PT project (PT A, PT B) { return B * dot(A, B) / dot(B, B); }
// Verifica se existe interseccao de segmentos
// (assumir que [A,B] e [C,D] sao coplanares)
bool seg_intersect(PT A, PT B, PT C, PT D) {
  return cmp(dot(cross(A-B, C-B), cross(A-B, D-B))) <= 0 &&
    cmp(dot(cross(C-D, A-D), cross(C-D, B-D))) \le 0;
// square distance between point and line, ray or segment
double ptLineDistSq(PT s1, PT s2, PT p, int type) {
  double pd2 = distSq(s1, s2);
  PT r;
  if(pd2 == 0)
    r = s1;
  else{
    double u = dot(p-s1, s2-s1) / pd2;
    r = s1 + (s2 - s1) *u;
    if(type != LINE && u < 0.0)
      r = s1;
    if (type == SEGMENT && u > 1.0)
      r = s2;
  return distSq(r, p);
// Distancia de ponto P ao segmento [A,B]
double dist_point_seg(PT P, PT A, PT B) {
```

```
PT PP = A + project(P-A, B-A);
 if (cmp(norma(A-PP) + norma(PP-B), norma(A-B)) == 0)
    return norma(P-PP);//distance point-line!
 else
   return min (norma (P-A), norma (P-B));
// Distance between lines ab and cd. TODO: Test this
double lineLineDistance(PT a, PT b, PT c, PT d) {
 PT v1 = b-a;
 PT v^2 = d-c;
 PT cr = cross(v1, v2);
 if (dot(cr, cr) < EPS) {</pre>
   PT proj = v1*(dot(v1, c-a)/dot(v1, v1));
   return sqrt (dot (c-a-proj, c-a-proj));
 } else {
   PT n = cr/sqrt(dot(cr, cr));
   PT p = dot(n, c - a);
   return sqrt(dot(p, p));
 }
// Menor distancia do segmento [A,B] ao segmento [C,D] (lento
#define dps dist_point_seq
double dist_seq_seq(PT A, PT B, PT C, PT D) {
 PT E = project (A-D, cross(B-A, D-C));
 // distance between lines!
 if (seq_intersect(A, B, C+E, D+E)) {
   return norma(E);
 }else {
   double dA = dps(A,C,D), dB = dps(B,C,D);
   double dC = dps(C, A, B), dD = dps(D, A, B);
   return min(min(dA, dB), min(dC, dD));
 }
// Menor distancia do segmento [A,B] ao segmento [C,D] (rapido
double dist_seq_seq2(PT A, PT B, PT C, PT D) {
 PT u(B-A), v(D-C), w(A-C);
 double a = dot(u, u), b = dot(u, v);
  double c = dot(v, v), d = dot(u, w), e = dot(v, w);
  double DD = a*c - b*b;
  double sc, sN, sD = DD;
 double tc, tN, tD = DD;
 if (DD < EPS) {
   sN = 0, sD = 1, tN = e, tD = c;
 }else{
   sN = (b*e - c*d);
```

```
tN = (a*e - b*d);
   if (sN < 0) {
    sN = 0, tN = e, tD = c;
   else if(sN > sD) 
     sN = sD, tN = e+b, tD = c;
 if (tN < 0) {
   tN = 0;
   if (-d < 0) sN = 0;
    else if (-d > a) sN = sD;
    else{
     sN = -d;
      sD = a;
 }else if(tN > tD) {
   tN = tD;
   if ((-d + b) < 0) sN = 0;
   else if (-d + b > a) sN = sD;
   else{
    sN = -d + b;
     sD = a;
  sc = fabs(sN) < EPS ? 0 : sN/sD;
  tc = fabs(tN) < EPS ? 0 : tN/tD;
  PT dP = w + (u*sc) - (v*tc);
  return norma(dP);
// Distancia de Ponto a Triangulo, dps = dist_point_seg
double dist_point_tri(PT P, PT A, PT B, PT C) {
 PT N = cross(B-A, C-A);
 PT PP = P - project (P-A, N);
PT R1, R2, R3;
 R1 = cross(B-A, PP-A);
 R2 = cross(C-B, PP-B);
 R3 = cross(A-C, PP-C);
  if (cmp(dot(R1,R2))) >= 0 && cmp(dot(R2,R3)) >= 0 &&
      cmp(dot(R3,R1)) >= 0) {
    return norma(P-PP);
    return min(dps(P,A,B), min(dps(P,B,C), dps(P,A,C)));
// compute a, b, c, d such that all points lie on ax + by + cz
   = d. TODO: test this
double planeFromPts(PT p1, PT p2, PT p3, double& a, double& b,
   double& c, double& d) {
 PT normal = cross(p2-p1, p3-p1);
```

```
a = normal.x; b = normal.y; c = normal.z;
                                                                         double v=0;
  d = -a*p1.x-b*p1.y-c*p1.z;
                                                                         for (i=0; i<poly.size(); i++)</pre>
                                                                           v+=fabs(signedTetrahedronVol(cent,poly[i][0],poly[i][1],
                                                                               poly[i][2]));
// project point onto plane. TODO: test this
                                                                         return v;
PT ptPlaneProj(PT p, double a, double b, double c, double d) {
  double l = (a*p.x+b*p.y+c*p.z+d)/(a*a+b*b+c*c);
  return PT(p.x-a*1, p.y-b*1, p.z-c*1);
// distance from point p to plane aX + bY + cZ + d = 0
                                                                       // Outras implementacoes [Usa struct PT]
double ptPlaneDist(PT p, double a, double b, double c, double d
                                                                       struct line{ // reta definida por um ponto p e direcao v
  return fabs(a*p.x + b*p.y + c*p.z + d) / sqrt(a*a + b*b + c*c
                                                                         PT p, v;
                                                                         line(){};
     );
                                                                         line (const PT& p, const PT& v):p(p), v(v) {
                                                                           assert(!(v == PT()));
// distance between parallel planes aX + bY + cZ + d1 = 0 and
                                                                         bool on(const PT& pt) const{
// aX + bY + cZ + d2 = 0
double planePlaneDist (double a, double b, double c, double d1,
                                                                           return cross(pt - p, v) == PT();
    double d2) {
  return fabs (d1 - d2) / sqrt (a*a + b*b + c*c);
                                                                       };
                                                                       struct plane {
                                                                         PT n;
// Volume de Tetraedro
                                                                         double d;
double signedTetrahedronVol(PT A, PT B, PT C, PT D) {
                                                                         plane() : d(0) {}
  double A11 = A.x - B.x;
                                                                         plane (const PT &p1, const PT &p2,
  double A12 = A.x - C.x;
                                                                             const PT &p3) {
  double A13 = A.x - D.x;
                                                                           n = cross(p_2 - p_1, p_3 - p_1);
                                                                           d = -dot(n, p1);
  double A21 = A.y - B.y;
  double A22 = A.y - C.y;
                                                                           assert (side (p1) == 0);
  double A23 = A.y - D.y;
                                                                           assert (side (p_2) == 0);
  double A31 = A.z - B.z;
                                                                           assert (side (p3) == 0);
  double A32 = A.z - C.z;
                                                                         int side(const PT &p) const {
  double A33 = A.z - D.z;
  double det =
                                                                           return sqn(dot(n, p) + d);
    A11*A22*A33 + A12*A23*A31 +
    A13*A21*A32 - A11*A23*A32 -
                                                                       };
    A12*A21*A33 - A13*A22*A31;
  return det / 6;
                                                                       // interesecao de retas
                                                                       int intersec (const line& 11, const line& 12,
}
                                                                           PT& res){
                                                                         assert(!(11.v == PT()));
// Parameter is a vector of vectors of points - each interior
                                                                         assert(!(12.v == PT());
// represents the 3 points that make up 1 face, in any order.
                                                                         if (cross(11.v, 12.v) == PT()) {
// Note: The polyhedron must be convex, with all faces given as
                                                                           if (cross(11.v, 11.p - 12.p) == PT())
                                                                             return 2; // same
double polyhedronVol(vector<vector<PT> > poly) {
                                                                           return 0; // parallel
  int i, j;
  PT cent (0, 0, 0);
                                                                         PT n = cross(l1.v, l2.v);
  for (i=0; i<poly.size(); i++)</pre>
                                                                         PT p = 12.p - 11.p;
                                                                         if (sqn(dot(n,p)))
    for (j=0; j<3; j++)
      cent=cent+poly[i][j];
                                                                           return 0; // skew
                                                                         double t;
  cent=cent \star (1.0/(poly.size() \star3));
```

```
if (sqn(n.x))
    t = (p.y * 12.v.z - p.z * 12.v.y) / n.x;
  else if (sqn(n.y))
   t = (p.z * 12.v.x - p.x * 12.v.z) / n.y;
  else if (sqn(n.z))
    t = (p.x * 12.v.y - p.y * 12.v.x) / n.z;
  else
    assert (false);
  res = 11.p + 11.v * t;
  assert(l1.on(res)); assert(l2.on(res));
  return 1; // intersects
// distancia entre 2 retas
double dist(const line& l1,const line& l2){
  PT ret = 11.p - 12.p;
  ret = ret - 11.v * (dot(11.v, ret) / 11.v.dist2());
  PT tmp = 12.v - 11.v *
    (dot(l1.v, l2.v) / l1.v.dist2());
  if (sqn(tmp.dist2()))
    ret = ret - tmp * (dot(tmp,ret) / tmp.dist2());
  assert(fabs(dot(ret,l1.v)) < eps);</pre>
  assert(fabs(dot(ret,tmp)) < eps);</pre>
  assert(fabs(dot(ret, 12.v)) < eps);</pre>
  return ret.dist();
}
// Retorna os dois pontos mais proximos entre 11 e 12
void closest(const line& 11, const line& 12,
    PT& p1, PT& p2) {
  if (cross(11.v, 12.v) == PT())
    p1 = 11.p;
    p2 = 12.p - 11.v *
      (dot(11.v, 12.p - 11.p) / 11.v.dist2());
    return;
  PT p = 12.p - 11.p;
  double t1 = (
      dot(11.v,p) * 12.v.dist2() -
      dot(11.v, 12.v) * dot(12.v, p)
       ) / cross(l1.v, l2.v).dist2();
  double t^2 = (
      dot(12.v, 11.v) * dot(11.v, p) -
      dot(12.v,p) * 11.v.dist2()
       ) / cross(12.v,11.v).dist2();
  p1 = 11.p + 11.v * t1;
  p2 = 12.p + 12.v * t2;
  assert (11.on(p1));
  assert (12.on(p2));
//retorna a intersecao de reta com plano [retorna 1 se
    intersecao for ptl
int cross(const line &1, const plane &pl,
```

```
PT &res) {
  double d = dot(pl.n, l.v);
  if (sqn(d) == 0) {
    return (pl.side(l.p) == 0) ? 2 : 0;
  double t = (-dot(pl.n, l.p) - pl.d) / d;
  res = l.p + l.v * t;
#ifdef DEBUG
  assert(pl.side(res) == 0);
#endif
  return 1;
bool cross (const plane & plane & plane & p2,
    const plane& p3, PT& res) {
  double d = det(p1.n, p2.n, p3.n);
  if (sqn(d) == 0) {
    return false:
  PT px (p1.n.x, p2.n.x, p3.n.x);
  PT py (p1.n.y, p2.n.y, p3.n.y);
  PT pz(p1.n.z, p2.n.z, p3.n.z);
  PT p(-p1.d, -p2.d, -p3.d);
  res = PT(
      det(p,py,pz)/d,
      det(px,p,pz)/d
      det(px,py,p)/d
   );
#ifdef DEBUG
  assert(p1.side(res) == 0);
  assert (p_2.side (res) == 0);
  assert (p3.side (res) == 0);
#endif
  return true;
// retorna reta da intersecao de dois planos
int cross (const plane &p1, const plane &p2,
    line &res) {
  res.v = cross(p1.n, p2.n);
 if (res.v == PT()) {
    if ( (p1.n * (p1.d / p1.n.dist2())) ==
        (p2.n * (p2.d / p2.n.dist2())))
      return 2;
    else
      return 0:
  plane p3;
  p3.n = res.v;
  p3.d = 0;
  bool ret = cross(p1, p2, p3, res.p);
  assert (ret);
  assert (p1.side(res.p) == 0);
  assert(p2.side(res.p) == 0);
```

```
return 1;
// testes
int main(){
   line 1;
   1.p = PT(1, 1, 1);
   1.v = PT(1, 0, -1);
   plane p(PT(10, 11, 12), PT(9, 8, 7), PT(1, 3, 2));
   PT res;
    assert(cross(1, p, res) == 1);
   plane p1(PT(1, 2, 3), PT(4, 5, 6), PT(-1, 5, -4));
   plane p2(PT(3, 2, 1), PT(6, 5, 4), PT(239, 17, -42));
   line 1;
    assert (cross (p_1, p_2, 1) == 1);
   plane p1(PT(1, 2, 3), PT(4, 5, 6), PT(-1, 5, -4));
   plane p2(PT(1, 2, 3), PT(7, 8, 9), PT(3, -1, 10));
   line 1;
    assert (cross (p1, p2, 1) == 2);
   plane p1(PT(1, 2, 3), PT(4, 5, 6), PT(-1, 5, -4));
   plane p2(PT(1, 2, 4), PT(4, 5, 7), PT(-1, 5, -3));
    assert (cross (p_1, p_2, 1) == 0);
  line 11, 12;
  while (l1.p.load())
   11.v.load(); 11.v = 11.v - 11.p;
   12.p.load();
   12.v.load(); 12.v = 12.v - 12.p;
    if (l1.v == PT() || l2.v == PT()) continue;
    PT res;
    int cnt = intersec(11,12,res);
    double d = dist(11, 12);
    if (fabs(d) < eps)</pre>
      assert (cnt >= 1);
      assert (cnt == 0);
    PT p1, p2;
    closest (11, 12, p1, p2);
    assert (fabs ((p1-p2).dist() - d) < eps);
  plane a (PT(1,0,0),PT(0,1,0),PT(0,0,1));
 plane b(PT(-1,0,0),PT(0,-1,0),PT(0,0,-1));
 line 1:
  assert ((cross(a,b,l))==0);
```

```
return 0;
```

6 Grafos

6.1 Topological Sort

```
vector<int> adj[MAXV];
void topo_sort(int N) {
    queue<int> q; // priority_queue?
    vector<int> in_degree(N,0);

    rep(i, 0, N) rep(j, 0, adj[i].size())
    in_degree[ adj[i][j] ] ++;

rep( i, 0, n ) if (in_degree[i] == 0) q.push(i);
while( !q.empty() ) {
    int u = q.front(); q.pop();
    rep(i, 0, adj[u].size()) {
        int v = adj[u][i];
        if ( --in_degree[v] == 0 ) q.push(v);
    }
}
```

6.2 Dijkstra

```
#define MAXV 100000
int dist[MAXV], pi[MAXV];// dist from s and pointer to parent
vector<ii> adj[MAXV]; //edge = {v, dist}
int dijkstra(int s, int t, int n){
  priority_queue<ii> pq;
  memset(pi, -1, sizeof pi);
  memset(dist, INF, sizeof dist); // INF eh suficiente?
  pq.push(ii(dist[s]=0, s));
  while (!pq.empty()) {
    ii top = pq.top(); pq.pop();
    int u = top.second, d = -top.first;
    if (d != dist[u]) continue;
    if (u == t) break; // terminou antes
    rep(i, 0, (int) adj[u].size()) {
      int v = adj[u][i].F;
      int cost = adj[u][i].S;
      if (dist[v] > dist[u] + cost){
        dist[v] = dist[u] + cost;
        pi[v] = u;
        pq.push(ii(-dist[v], v));
  return dist[t];
```

6.3 Floyd-Warshall

```
int adj[400][400], path[400][400];
int n, m; // #vertices, #arestas
// adi[u][v] = custo de {U->V}
// path[u][v] = k .: K vem logo apos U no caminho ate V
void read_graph() {
 memset(adj, INF, sizeof adj); // para menor caminho
  rep(i, 0, n) adj[i][i] = 0; // para menor caminho
  int u, v, w;
  rep(i, 0, m) {
   cin >> u >> v >> w;
   adi[u][v] = w;
   path[u][v] = v;
 }
void floyd() {
  rep(k, 0, n)
   rep(i, 0, n)
      rep(j, 0, n)
        if(adj[i][k] + adj[k][j] < adj[i][j]){
          adj[i][j] = adj[i][k] + adj[k][j];
          path[i][j] = path[i][k];
vector<int> findPath(int s, int d){
  vector<int> Path; Path.pb(s);
 while (s != d) \{ s = path[s][d]; Path.pb(s); \}
  return Path;
/*Aplicacoes:
1-Encontrar o fecho transitivo (saber se U conseque visitar V)
.: adj[u][v] = (adj[u][k] & adj[k][v]);
   (inicializar adj com 0)
2-Minimizar a maior aresta do caminho entre U e V
: adj[u][v] = min(adj[u][v], max(adj[u][k], adj[k][v]));
   (inicializar adj com INF)
3-Maximizar a menor aresta do caminho entre U e V
:: adj[u][v] = max(adj[u][v], min(adj[u][k], adj[k][u]));
   (inicializar adj com -INF) ∗/
```

6.4 Bellman-Ford

```
#define MAXV 400
//Vertices indexados em 0.
int V, E;//#vertices, #arestas
vector<ii> adj[MAXV];
```

```
11 dist[MAXV];
//Menor custo de S para todos vertices em O(V^3)
//bellman() retorna FALSE se o grafo tem ciclo com custo
//dist[v] contem o menor custo de s ate v.
bool bellman(int s) {
  rep(i, 0, V) dist[i] = INF;
  dist[s] = 0;
  rep(i, 0, V-1)
   rep(u, 0, V) {
      rep(j, 0, adj[u].size()){
        int v = adj[u][j].F, duv = adj[u][j].S;
        dist[v] = min(dist[v], dist[u] + duv);
  // verifica se tem ciclo com custo negativo
  rep(u, 0, V)
    rep(j, 0, adj[u].size()){
      int v = adj[u][j].F, duv = adj[u][j].S;
      if (dist[v] > dist[u] + duv) return false;
  return true;
```

6.5 Vértices de Articulação e Pontes

```
#define MAXV 100000
vector<int> adj[MAXV];
int dfs_num[MAXV], dfs_low[MAXV], dfs_parent[MAXV];
int dfscounter, V, dfsRoot, rootChildren, ans;
int articulation[MAXV], articulations;
vector<ii> bridges;
void articulationPointAndBridge(int u) {
 dfs low[u] = dfs num[u] = dfscounter++;
 rep(i, 0, adj[u].size()) {
   int v = adj[u][i];
   if (dfs_num[v] == -1) {
     dfs_parent[v] = u;
     if (u == dfsRoot) rootChildren++;
     articulationPointAndBridge(v);
     if(dfs_low[v] >= dfs_num[u]) articulation[u] = true;
     if (dfs_low[v] > dfs_num[u]) bridges.pb(mk(u,v));
      dfs_low[u] = min(dfs_low[u], dfs_low[v]);
   else if (v != dfs_parent[u])
      dfs_low[u] = min(dfs_low[u], dfs_num[v]);
int main(){
```

```
// read graph
dfscounter = 0;
rep(i, 0, V) {
  dfs_low[i] = dfs_parent[i] = articulation[i] = 0;
  dfs num[i] = -1;
articulations = 0;
bridges.clear();
rep(i, 0, V)
  if (dfs_num[i] == -1){
    dfsRoot = i; rootChildren = 0;
    articulationPointAndBridge(i);
    articulation[dfsRoot] = (rootChildren > 1);
printf("#articulations = %d\n", articulations);
rep(i, 0, V)
  if(articulation[i]) printf("Vertex %d\n", i);
printf("#bridges = %d\n",bridges.size());
rep(i, 0, bridges.size())
  printf("Bridge %d<->%d\n", bridges[i].F, bridges[i].S);
return 0;
```

6.6 Tarjan

```
#define MAXV 100000
vector<int> adj[MAXV]; int V;
int dfs_num[MAXV], dfs_low[MAXV], vis[MAXV], SCC[MAXV];
int dfsCounter, numSCC;
vector<int> S; //global variables
void tarjanSCC(int u) {
  dfs_low[u] = dfs_num[u] = dfsCounter++; //dfs_low[u] <=</pre>
     dfs_num[u]
  S.push back(u); //stores u in a vector based on order of
     visitation
  vis[u] = 1;
  rep(i, 0, adj[u].size()) {
   int v = adj[u][i];
   if (dfs_num[v] == -1)
      tarjanSCC(v);
   if (vis[v]) //condition for update
      dfs_low[u] = min(dfs_low[u], dfs_low[v]);
  if (dfs_low[u] == dfs_num[u]) { //if this is a root (start) of
      an SCC
    while (42) {
      int v = S.back(); S.pop_back();
      vis[v] = 0;
      SCC[v] = numSCC; //wich SCC this vertex belong
      if (u == v) break;
    numSCC++;
```

```
}
int main() {
    // read graph
    rep(i,0,V) {
        dfs_num[i] = -1;
        dfs_low[i] = vis[i] = 0;
        SCC[i] = -i;
    }
    dfsCounter = numSCC = 0;
    rep(i,0,V)
        if(dfs_num[i] == -1)
            tarjanSCC(i);
    return 0;
}
```

6.7 Kosaraju

```
#define MAXV 100000
vector<int> adj[2][MAXV]; // adj[0][] original, adj[1][]
   transposto
vector<int> S, dfs_num;
int N, numSCC, SCC[MAXV];
void Kosaraju(int u, int t, int comp) {
  Wdfs_num[u] = 1;
  if (t == 1) SCC[u] = comp;
  for (int j = 0; j < (int)adj[t][u].size(); j++) {</pre>
    int v = adj[t][u][j];
    if (dfs_num[v] == DFS_WHITE) Kosaraju(v, t, comp);
  S.push_back(u);
void doit(){ // chamar na main
  S.clear();
  dfs_num.assign(N, DFS_WHITE);
  for (int i = 0; i < N; i++)</pre>
   if (dfs_num[i] == DFS_WHITE) Kosaraju(i, 0, -1);
    numSCC = 0;
  dfs_num.assign(N, DFS_WHITE);
  for (int i = N - 1; i >= 0; i--)
    if (dfs_num[ S[i] ] == DFS_WHITE) {
      Kosaraju(S[i], 1, numSCC);
      numSCC++;
  printf("There are %d SCCs\n", numSCC);
```

```
// 2-sat - Codigo do problema X-Mart
// vertices indexado em 1
vector<int> adj[2*MAXV];
vector<int> radj[2*MAXV];
int seen[2 * MAXV], comp[2 * MAXV], order[2 * MAXV], ncomp,
   norder;
int N; // #variaveis
int n; // #vertices
#define NOT(x) ( (x<=N)? (x+N) : (x-N) )
#define quero 1
void add_edge(int a, int b, int opcao) {
  if (a > b) swap(a, b);
 if (b == 0) return;
  if (a == 0) {
    if (opcao == quero) adj[NOT(b)].pb(b);
                        adj[b].pb(NOT(b));
  }else{ // normal...
    if (opcao == quero) {
      adj[NOT(a)].pb(b);
      adj[NOT(b)].pb(a);
    }else{
      a = NOT(a); b = NOT(b);
      adj[NOT(a)].pb(b);
      adj[NOT(b)].pb(a);
    }
  }
void init(){
 rep(i, 0, n+1) {
   adj[i].clear();
    radj[i].clear();
 }
void dfs1(int u) {
  seen[u] = 1;
  rep(i,0,adj[u].size()) if (!seen[adj[u][i]]) dfs1(adj[u][i]
     );
  order[norder++] = u;
void dfs2(int u) {
  seen[u] = 1;
  rep(i,0,radj[u].size()) if (!seen[radj[u][i]]) dfs2( radj[u][
     il);
  comp[u] = ncomp;
void strongly connected components() {
  rep(v,1,n+1) rep(i,0,(int)adj[v].size()) radj[adj[v][i]].pb
     (v);
  norder = 0; memset(seen, 0, sizeof seen);
  rep(v,1,n+1) if (!seen[v]) dfs1(v);
  ncomp = 0; memset(seen, 0, sizeof seen);
```

```
for (int i = n-1, u = order[n-1]; i>=0; u = order[--i])
   if (!seen[u]) { dfs2(u); ncomp++; }
bool sat2(){
  strongly_connected_components();
 rep(i,1,n+1) if (comp[i] == comp[NOT(i)]) return false;
  return true;
int main(){
 int Clientes;
 while (cin >> Clientes >> N) {
   if (Clientes == 0 && N==0) break;
   n = 2 * N;
   init();
   int u, v;
   rep(i, 0, Clientes) {
     scanf("%d %d", &u, &v);
     add_edge(u, v, quero);
     scanf("%d %d", &u, &v);
     add_edge(u, v, !quero);
    sat2() ? printf("yes\n") : else printf("no\n");
 return 0;
```

6.9 LCA

```
/*Lowest Common Ancestor (LCA) entre dois vertices A, B de uma
arvore. LCA(A,B) = ancestral mais proximo de A e B.
O codigo abaixo tambem calcula a menor aresta do caminho entre
A \in B.
Para saber quantas arestas tem entre A e B basta fazer:
 level[A]+level[B]-2*level[lca(A,B)] Pode-se modificar para
 retorna a distancia entre A e B.
Como usar:
(1) ler a arvore em E[] e W[], chamar doit(raiz),
passando a raiz da arvore. Indexar em 0 os vertices
(2) A funcao retorna o LCA e a menor aresta entre A e B.
#define MAXV 101000
const int maxl = 20; // profundidade maxima 2^(maxl) > MAXV
int pai[MAXV][maxl+1]; // pai[v][i] = pai de v subindo 2^i
   arestas
int dist[MAXV][maxl+1]; // dist[v][i] = menor aresta de v
   subindo 2^i arestas
int level[MAXV]; // level[v] = #arestas de v ate a raiz
int N, M; // numero de vertices e arestas
vector<pair<int,int> > E[MAXV]; // {v,custo}
void dfs(int v, int p, int peso){
```

```
level[v] = level[p] + 1;
  pai[v][0] = p;
  dist[v][0] = peso; // aresta de v->pai[v]
  for (int i = 1; i <= maxl; i++) {</pre>
    pai[v][i] = pai[ pai[v][i-1] ][i-1]; // subindo 2^i arestas
    dist[v][i] = min(dist[v][i-1]), dist[pai[v][i-1]][i-1]);
         //(11)
    rep(i, 0, E[v].size()){
   int viz = E[v][i].first;
   int cost = E[v][i].second;
   if (viz == p) continue;
    dfs(viz, v, cost);
void doit(int root){
 level[root] = 0;
  for (int i = 0; i <= maxl; i++)</pre>
    pai[root][i] = root, dist[root][i] = INF;
  rep(i, 0, E[root].size()){
      int viz = E[root][i].first;
      int cost = E[root][i].second;
    dfs(viz, root, cost);
}
pair<int, int> lca(int a, int b) {
  int menor = INF; // valor da menor aresta do caminho a->b
  if (level[a] < level[b]) swap(a, b);</pre>
  for (int i = maxl; i >= 0; i--) {
   if (level[ pai[a][i] ] >= level[b]) {
     menor = min(menor, dist[a][i]);
      a = pai[a][i];
  if (a != b) {
   for (int i = maxl; i >= 0; i--) {
      if (pai[a][i] != pai[b][i]) {
        menor = min(menor, min(dist[a][i], dist[b][i]));
        a = pai[a][i]; b = pai[b][i];
    //ultimo salto
   menor = min(menor, min(dist[a][0], dist[b][0]));
    a = pai[a][0]; b = pai[b][0];
  return make_pair(a, menor);
```

```
// Encontra o casamento bipartido maximo entre os vertices
// grupo X e Y. x = [0, X-1], y = [0, Y-1]
// match[y] = x - contem quem esta casado com y
// Teorema de Konig - Num grafo bipartido, o matching eh igual
    ao minimum vertex cover
// Complexidade O(nm)
#define MAXV 1000
vector<int> adj[MAXV];
int match[MAXV], V, X, Y;
bool vis[MAXV];
int aug(int v) {
 if(vis[v]) return 0;
  vis[v] = true;
  rep(i, 0, adj[v].size()) {
    int r = adj[v][i];
    if (match[r] == -1 \mid \mid aug(match[r])){
      match[r] = v; // augmenting path
      return 1;
  return 0;
int matching(int X, int Y) {
 int V = X + Y;
  rep(i, 0, V) match[i] = -1;
  int mcbm = 0;
  rep(i, 0, X) {
    rep(j, 0, X) vis[j] = false;
    mcbm += aug(i);
  return mcbm;
```

6.11 Hopcroft Karp - Maximum Bipartite Matching

```
/*Encontra o casamento bipartido maximo de forma eficiente
1) Chamar a funcao init(L,R) L: #vertices da esquerda, R: #
   vertices da direita
2) Usar addEdge(Li,Ri) para adicionar a aresta Li -> Ri
3) maxMatching() retorna o casamento maximo
matching[]: matching[Rj] -> armazena Li
Complexidade do algoritmo:
O(sqrt(n) * m) */
#define MAXN1 3010
#define MAXN2 3010
#define MAXM 6020
int n1, n2, edges, last[MAXN1], prev[MAXM], head[MAXM];
int matching[MAXN2], dist[MAXN1], Q[MAXN1];
bool used[MAXN1], vis[MAXN1];
void init(int L, int R){
 n1 = L, n2 = R;
```

```
edges = 0;
  fill(last, last + n1, -1);
void addEdge(int u, int v) {
  head[edges] = v;
  prev[edges] = last[u];
 last[u] = edges++;
void bfs(){
  fill(dist, dist + n1, -1);
  int sizeQ = 0;
  for (int u = 0; u < n1; ++u) {
    if (!used[u]) {
      Q[sizeQ++] = u;
      dist[u] = 0;
  for (int i = 0; i < sizeQ; i++) {</pre>
    int u1 = Q[i];
    for (int e = last[u1]; e >= 0; e = prev[e]) {
      int u2 = matching[head[e]];
      if (u^2 >= 0 \&\& dist[u^2] < 0) {
        dist[u2] = dist[u1] + 1;
        Q[sizeQ++] = u2;
    }
bool dfs(int u1) {
  vis[u1] = true;
  for (int e = last[u1]; e >= 0; e = prev[e]) {
    int v = head[e];
    int u2 = matching[v];
    if (u^2 < 0 \mid | !vis[u^2] \&\& dist[u^2] == dist[u^1] + 1 \&\& dfs(u^2)
        2)){
      matching[v] = u1;
      used[u1] = true;
      return true;
  return false;
int maxMatching() {
  fill(used, used + n1, false);
  fill (matching, matching + n^2, -1);
  for (int res = 0;;) {
    bfs();
    fill(vis, vis + n1, false);
    int f = 0;
    for (int u = 0; u < n1; ++u)
      if (!used[u] && dfs(u)) ++f;
    if (!f) return res;
    res += f;
```

6.12 Network Flow (lento)

```
#define MAXV 250
vector<int> edge[MAXV];
int cap[MAXV][MAXV];
bool vis[MAXV];
void init() {
  rep(i, 0, MAXV) edge[i].clear();
  memset(cap, 0, sizeof cap);
void add(int a, int b, int cap_ab, int cap_ba) {
  edge[a].pb(b), edge[b].pb(a);
  cap[a][b] += cap_ab, cap[b][a] += cap_ba;
int dfs(int src, int snk, int fl){
  if(vis[src]) return 0;
  if(snk == src) return fl;
  vis[src] = 1;
  rep(i, 0, edge[src].size()){
    int v = edge[src][i];
    int x = min(fl, cap[src][v]);
    if (x > 0) {
      x = dfs(v, snk, x);
      if( !x ) continue;
      cap[src][v] -= x;
      cap[v][src] += x;
      return x;
  return 0;
int flow(int src, int snk) {
  int ret = 0;
  while( 42 ) {
    memset(vis, 0, sizeof vis);
    int delta = dfs(src, snk, 1 << 30);
   if(!delta) break;
    ret += delta;
  return ret;
```

6.13 Network Flow - Dinic

```
/*
Grafo indexado em 1
```

```
Inicializar maxN, maxE.
    Chamar init() com #nos e source e sink
    montar o grafo chamando add(a,b,c1,c2), sendo c1 cap. de a
       ->b e c2 cap. de b->a
#define FOR(i,a,b) for(int i=a; i<=b; i++)</pre>
#define SET(c, v) memset(c, v, sizeof c)
const int maxN = 5000;
const int maxE = 70000;
const int inf = 1000000005;
int nnode, nedge, src, snk;
int Q[ maxN ], pro[ maxN ], fin[ maxN ], dist[ maxN ];
int flow[ maxE ], cap[ maxE ], to[ maxE ], next[ maxE ];
void init( int _nnode, int _src, int _snk ) {
  nnode = _nnode, nedge = 0, src = _src, snk = _snk;
 FOR(i, 1, nnode) fin[i] = -1;
void add( int a, int b, int c1, int c2 ) {
  to[nedge]=b, cap[nedge]=c1, flow[nedge]=0, next[nedge]=fin[a
     l, fin[a]=nedge++;
  to[nedge]=a, cap[nedge]=c2, flow[nedge]=0, next[nedge]=fin[b
     ], fin[b]=nedge++;
bool bfs() {
  SET ( dist, -1 );
 dist[src] = 0;
  int st = 0, en = 0;
  Q[en++] = src;
  while( st < en ) {</pre>
   int u = Q[ st++ ];
   for(int e = fin[u]; e >= 0; e = next[e] ) {
     int v = to[e];
     if( flow[e] < cap[e] && dist[v] == -1 ) {</pre>
       dist[v] = dist[u] + 1;
        Q[en++] = v;
      }
    }
  return dist[snk] != -1;
int dfs(int u, int fl) {
  if( u == snk ) return fl;
  for( int& e = pro[u]; e >= 0; e = next[e] ) {
    int v = to[e];
   if( flow[e] < cap[e] && dist[v] == dist[u]+1 ) {</pre>
      int x = dfs(v, min(cap[e] - flow[e], fl));
     if(x > 0) {
        flow[e] += x, flow[e^1] -= x;
        return x;
```

```
}
}
return 0;
}
LL dinic() {
   LL ret = 0;
   while( bfs() ) {
    FOR(i, 1, nnode) pro[i] = fin[i];
    while( 42 ) {
       int delta = dfs( src, inf );
       if( !delta ) break;
       ret += delta;
    }
}
return ret;
}
```

6.14 Min Cost Max Flow

```
// Criar o grafo chamando MCMF q(V), onde q eh o grafo e V a
    gtde de vertices (indexado em 0)
// chamar q.add(u,v,cap,cost) para add a aresta u->v, se for
   bidirecional, chamar denovo .: g.add(v,u,cap,cost)
struct MCMF {
  typedef int ctype;
  enum { MAXN = 550, INF = INT_MAX };
  struct Edge { int x, y; ctype cap, cost; };
  vector<Edge> E;
                      vector<int> adj[MAXN];
  int N, prev[MAXN]; ctype dist[MAXN], phi[MAXN];
 MCMF (int NN) : N(NN) {}
 void add(int x, int v, ctype cap, ctype cost) { // cost >= 0
   Edge e1=\{x,y,cap,cost\}, e2=\{y,x,0,-cost\};
   adj[e1.x].push_back(E.size()); E.push_back(e1);
   adj[e2.x].push_back(E.size()); E.push_back(e2);
  void mcmf(int s, int t, ctype &flowVal, ctype &flowCost) {
    flowVal = flowCost = 0; memset(phi, 0, sizeof(phi));
    while (true) {
     for (x = 0; x < N; x++) prev[x] = -1;
     for (x = 0; x < N; x++) dist[x] = INF;
      dist[s] = prev[s] = 0;
      set< pair<ctype, int> > Q;
      Q.insert(make_pair(dist[s], s));
      while (!Q.empty()) {
        x = Q.begin() -> second; Q.erase(Q.begin());
        // x is a vertex
```

```
// FOREACH(it, adj[x]) { // adj[] stores the index of a
             edge
        for (vector<int> :: iterator it = adj[x].begin(); it !=
            adj[x].end(); it++){
          const Edge &e = E[*it];
          if (e.cap <= 0) continue;</pre>
          ctype cc = e.cost + phi[x] - phi[e.y];
                                 // ***
          if (dist[x] + cc < dist[e.y]) {
            Q.erase(make_pair(dist[e.y], e.y));
            dist[e.y] = dist[x] + cc;
            prev[e.v] = *it;
            Q.insert(make_pair(dist[e.y], e.y));
      if (prev[t] == -1) break;
      ctvpe z = INF;
      for (x = t; x != s; x = E[prev[x]].x) z = min(z, E[prev[x]])
         11.cap);
      for (x = t; x != s; x = E[prev[x]].x)
       { E[prev[x]].cap -= z; E[prev[x]^1].cap += z; }
      flowVal += z:
      flowCost += z * (dist[t] - phi[s] + phi[t]);
      for (x = 0; x < N; x++) if (prev[x] != -1) phi[x] += dist
          [X]; // ***
} ;
```

6.15 Min Cost Max Flow (Stefano)

```
void init(int _V) {
  memset(last, -1, sizeof(last));
  V = _{V}; E = 0;
void add_edge(int u, int v, cap_type _cap, cost_type _cost) {
  to[E] = v, cap[E] = _cap;
  cost[E] = cost, next[E] = last[u];
  last[u] = E++;
  to[E] = u, cap[E] = 0;
  cost[E] = -\_cost, next[E] = last[v];
  last[v] = E++;
// only if there is initial negative cycle
void BellmanFord(int s, int t) {
  bool stop = false;
  for (int i = 0; i < V; ++i) dist[i] = INF;</pre>
  dist[s] = 0;
  for (int i = 1; i <= V && !stop; ++i) {</pre>
    stop = true;
    for (int j = 0; j<E; ++j) {</pre>
      int u = to[j^1], v = to[j];
      if(cap[j]>0 && dist[u]!=INF && dist[u]+cost[j]<dist[v]){</pre>
        stop = false;
        dist[v] = dist[u]+cost[j];
  for (int i = 0; i < V; ++i) if (dist[i]!=INF) pot[i] = dist[i];</pre>
void mcmf(int s, int t){
  flowVal = flowCost = 0;
  memset(pot, 0, sizeof(pot));
  BellmanFord(s,t);
  while(true) {
    memset (prev, -1, sizeof (prev));
    memset (visited, false, sizeof (visited));
    for(int i = 0;i<V;++i) dist[i] = INF;</pre>
    priority_queue< pair<cost_type, int> > Q;
    Q.push(make_pair(0,s));
    dist[s] = prev[s] = 0;
    while(!Q.empty()){
      int aux = Q.top().second;
      Q.pop();
```

```
if(visited[aux]) continue;
  visited[aux] = true;
  for(int e = last[aux];e!=-1;e = next[e]){
    if(cap[e] <= 0) continue;</pre>
    cost_type new_dist = dist[aux]+cost[e]+pot[aux]-pot[to[
    if (new_dist<dist[to[e]]) {</pre>
      dist[to[e]] = new_dist;
      prev[to[e]] = e;
      Q.push (make_pair(-new_dist, to[e]));
if (prev[t]==-1) break;
cap_type f = cap[prev[t]];
for (int i = t;i!=s;i = to[prev[i]^1]) f = min(f, cap[prev[i
for(int i = t;i!=s;i = to[prev[i]^1]){
  cap[prev[i]] -= f;
  cap[prev[i]^1] += f;
flowVal += f;
flowCost += f*(dist[t]-pot[s]+pot[t]);
for(int i = 0; i < V; ++i) if (prev[i]!=-1) pot[i] += dist[i];</pre>
```

6.16 Tree Isomorphism

```
// Determina se duas arvores com indices distintos sao iguais
const int MAXV = 100010;
int N;
vector <int> A[MAXV], NA[MAXV], B[MAXV], NB[MAXV];
/* Armazenar em N o numero de vertices de cada arvore,
    salvar em A e em B as duas arvores indexadas em 1
    Antes de ler as arvores, limpar A[], NA[], B[], NB[]

*/
bool treeIsomorphism() {
    rep(i, 1, N+1) {
        rep(j, 0, A[i].size()) NA[i].pb( A[A[i][j]].size() );
        sort(NA[i].begin(), NA[i].end());

        rep(j, 0, B[i].size()) NB[i].pb( B[B[i][j]].size() );
        sort(NB[i].begin(), NB[i].end());
}
sort(NA + 1, NA + N + 1);
sort(NB + 1, NB + N + 1);
```

```
rep(i, 1, N+1) {
    if(NA[i] != NB[i])
      return false;
}
return true;
}
```

6.17 Stoer Wagner - Minimum Cut

```
/*
   Retorna o corte minimo do grafo
   (Conjunto de arestas que se removido, desconecta o grafo)
Input: n = \#vertices, q[i][j] = custo da aresta (i->j)
Output: Retorna o corte minimo
Complexidade: O(N^3)
// Maximum number of vertices in the graph
#define NN 101
// Maximum edge weight (MAXW * NN * NN must fit into an int)
#define MAXW 110
// Adjacency matrix and some internal arrays
int g[NN][NN], v[NN], w[NN], na[NN], n;
bool a[NN];
int stoer_wagner() {
 // init the remaining vertex set
  for( int i = 0; i < n; i++ ) v[i] = i;</pre>
  // run Stoer-Wagner
  int best = MAXW * n * n;
  while (n > 1)
    // initialize the set A and vertex weights
    a[v[0]] = true;
    for( int i = 1; i < n; i++ ){</pre>
      a[v[i]] = false;
      na[i - 1] = i;
      w[i] = q[v[0]][v[i]];
    // add the other vertices
    int prev = v[0];
    for( int i = 1; i < n; i++ ) {</pre>
     // find the most tightly connected non-A vertex
      int zj = -1;
      for( int j = 1; j < n; j++ )
       if( !a[v[j]] && ( zj < 0 || w[j] > w[zj] ) )
          zj = j;
      // add it to A
      a[v[zj]] = true;
      // last vertex?
      if ( i == n - 1 ) {
       // remember the cut weight
        best = min(best,w[zj]);
        // merge prev and v[zj]
```

6.18 Erdos Gallai

```
// Determina se existe um grafo tal que b[i] eh o grau do i-
    esimo vertice
// vertices indexado em 1. Apenas armazenar em b[1...N] e
    chamar EGL()
long long b[100005], n;
long long dmax, dmin, dsum, num_degs[100005];
bool basic_graphical_tests() { //Sort and perform some simple
    tests on the sequence
  int p = n;
  memset(num_degs, 0, (n+1)*sizeof(long long));
  dmax = dsum = n = 0;
  dmin = p;
  for (int d=1; d<=p; d++) {</pre>
    if(b[d] < 0 | | b[d] >= p)
      return false:
    else if(b[d] > 0){
      if(dmax < b[d]) dmax = b[d];
      if(dmin > b[d]) dmin = b[d];
      dsum = dsum + b[d];
      n++;
      num_degs[b[d]]++;
  if (dsum%2 || dsum > n*(n-1)) return false;
  return true;
bool EGL() {
  long long k, sum_deg, sum_nj, sum_jnj, run_size;
  if( !basic_graphical_tests()) return false;
  if (n == 0 \mid | 4*dmin*n >= (dmax+dmin+1)*(dmax+dmin+1)) return
      true;
  k = sum\_deg = sum\_nj = sum\_jnj = 0;
  for(int dk = dmax; dk >= dmin; dk--) {
```

```
if(dk < k+1) return true;

if(num_degs[dk] > 0) {
    run_size = num_degs[dk];
    if(dk < k+run_size)
        run_size = dk-k;
    sum_deg += run_size * dk;

for(int v=0; v<run_size; v++) {
        sum_nj += num_degs[k+v];
        sum_jnj += (k+v) * num_degs[k+v];
    }
    k += run_size;
    if(sum_deg > k*(n-1) - k*sum_nj + sum_jnj)
        return false;
}

return true;
}
```

6.19 Stable Marriage

```
/*Seja um conjunto de m homens e n mulheres, onde cada pessoa
  tem uma preferencia por outra de sexo oposto. O algoritmo
 produz o casamento estavel de cada homem com uma mulher.
Estavel:
- Cada homem se casara com uma mulher diferente (n >= m)
- Dois casais H1M1 e H2M2 nao serao instaveis.
Dois casais H1M1 e H2M2 sao instaveis se:
- H1 prefere M2 ao inves de M1, e
- M1 prefere H2 ao inves de H1.
Entrada
(1) m = \#homens, n = \#mulheres
(2) R[x][y] = i, i: eh a ordem de preferencia do homem y pela
Obs.: Quanto maior o valor de i menor eh a preferencia do homem
    y pela mulher x
(3) L[x][i] = y: A mulher y eh a i-esima preferencia do homem
Obs.: 0 \le i \le n-1, quanto menor o valor de i maior en a
   preferencia
do homem x pela mulher y
Saida
L2R[i]: a mulher do homem i (sempre entre 0 e n-1)
R2L[j]: o homem da mulher j (-1 se a mulher for solteira)
Complexidade O(m^2)
#define MAXM 1000
#define MAXW 1000
int L[MAXM][MAXW];
```

```
int R[MAXW][MAXM];
int L2R[MAXM], R2L[MAXW];
int m, n;
int p[MAXM];
void stableMarriage () {
  static int p[MAXM];
 memset (R2L, -1, sizeof (R2L));
 memset(p, 0, sizeof(p));
  for (int i = 0; i < m; ++i) {
    int man = i;
    while (man >= 0) {
      int wom;
      while (42) {
        wom = L[man][p[man]++];
        if (R2L[wom] < 0 \mid \mid R[wom][man] > R[wom][R2L[wom]])
      int hubby = R2L[wom];
      R2L[L2R[man] = wom] = man;
      man = hubby;
```

6.20 Hungarian Max Bipartite Matching with Cost

```
/*Encontra o casamento bipartido maximo/minimo com peso nas
   arestas
  Como usar:
  Criar o grafo:
  Hungarian G(L, R, ehMaximo)
  L = #vertices a esquerda
  R = #vertices a direita
  ehMaximo = variavel booleana que indica se eh casamento
     maximo ou minimo
  Adicionar arestas:
  G.add\_edge(x, y, peso)
  x = vertice da esquerda no intervalo [0, L-1]
 y = vertice da direita no intervalo [0, R-1]
  peso = custo da aresta
obs: tomar cuidado com multiplas arestas.
Resultado:
match_value = soma dos pesos dos casamentos
pairs = quantidade de pares (x-y) casados
xy[x] = vertice \ y \ casado \ com \ x
yx[y] = vertice x casado com y
Complexidade do algoritmo: O(n^3)
Problemas resolvidos: SCITIES (SPOJ)
```

```
*/
struct Hungarian{
  enum{ MAXN = 150, INF = 0x3f3f3f3f3f};
 int cost[MAXN][MAXN];
 int xy[MAXN], yx[MAXN];
 bool S[MAXN], T[MAXN];
 int lx[MAXN], ly[MAXN], slack[MAXN], slackx[MAXN], prev[MAXN
 int match_value, pairs;
 bool ehMaximo;
  int n;
  Hungarian(int L, int R, bool _ehMaximo=true) {
   n = max(L, R);
   ehMaximo = _ehMaximo;
   if (ehMaximo) memset(cost, 0, sizeof cost);
              memset(cost, INF, sizeof cost);
 void add_edge(int x, int y, int peso){
   if (!ehMaximo) peso *= (-1);
    cost[x][y] = peso;
  int solve(){
   match_value = 0;
   pairs = 0;
   memset (xy, -1, sizeof(xy));
   memset (yx, -1, sizeof(yx));
   init_labels();
    augment();
    for (int x = 0; x < n; ++x)
     match_value += cost[x][xy[x]];
    return match_value;
 void init_labels () {
   memset(lx, 0, sizeof(lx));
   memset(ly, 0, sizeof(ly));
   for (int x = 0; x < n; ++x)
     for (int y = 0; y < n; ++y)
        lx[x] = max(lx[x], cost[x][y]);
 void augment () {
   if (pairs == n) return;
   int x, y, root;
   int q[MAXN], wr = 0, rd = 0;
   memset(S, false, sizeof(S));
   memset(T, false, sizeof(T));
   memset(prev, -1, sizeof(prev));
    for (x = 0; x < n; ++x)
```

```
if (xy[x] == -1) {
    q[wr++] = root = x;
    prev[x] = -2;
    S[x] = true;
    break:
for (y = 0; y < n; ++y) {
  slack[y] = lx[root] + ly[y] - cost[root][y];
  slackx[y] = root;
while (true) {
  while (rd < wr) {</pre>
    x = q[rd++];
    for (y = 0; y < n; ++y)
      if (cost[x][y] == lx[x] + ly[y] && !T[y]){
        if (yx[y] == -1)
          break;
        T[y] = true;
        q[wr++] = yx[y];
        add(yx[y], x);
    if (y < n)
      break;
  if (y < n) break;</pre>
  update_labels();
  wr = rd = 0;
  for (y = 0; y < n; ++y)
    if (!T[y] && slack[y] == 0) {
      if (yx[y] == -1) {
        x = slackx[y];
        break;
      else{
        T[y] = true;
        if (!S[yx[y]]){
          q[wr++] = yx[y];
          add(yx[y], slackx[y]);
  if (y < n) break;</pre>
if (y < n) {
  ++pairs;
  for (int cx = x, cy = y, ty; cx != -2; cx = prev[cx], cy
     = ty) {
    ty = xy[cx];
    yx[cy] = cx;
    xy[cx] = cy;
  augment();
```

```
void add (int x, int prevx) {
    S[x] = true;
    prev[x] = prevx;
    for (int y = 0; y < n; ++y)
     if (lx[x] + ly[y] - cost[x][y] < slack[y]){
        slack[y] = lx[x] + ly[y] - cost[x][y];
        slackx[y] = x;
  }
  void update_labels () {
    int x, y, delta = INF;
    for (y = 0; y < n; ++y)
      if (!T[v])
        delta = min(delta, slack[y]);
    for (x = 0; x < n; ++x)
      if (S[x])
        lx[x] -= delta;
    for (y = 0; y < n; ++y)
      if (T[v])
        ly[y] += delta;
    for (y = 0; y < n; ++y)
      if (!T[y])
        slack[y] -= delta;
  int casouComX(int x){
    return xy[x];
  int casouComY(int y) {
    return yx[y];
};
// O codigo abaixo resolve o problema scities (Spoj)
int main(){
  int casos; cin >> casos;
  while (casos--) {
   int L, R;
    cin >> L >> R;
    Hungarian G(L, R, true);
    int x, y, w, aux[L][R];
    memset(aux, 0, sizeof aux);
    while (scanf("%d %d %d", &x, &y, &w) != EOF) {
      if (x==0 && y==0 && w==0) break;
      aux[x-1][y-1] += w;
    for (int x = 0; x < L; x++) {
      for (int y = 0; y < R; y++) {
        if (aux[x][y] != 0) {
          G.add\_edge(x,y, aux[x][y]);
```

```
}
}
printf("%d\n", G.solve());
}
return 0;
}
```

7 Estruturas de Dados

7.1 BIT

```
11 bit[MAXN]; int N;
void update(int i, 11 v) {
   for (; i <= N; i += i&-i)
      bit[i] += v;
}
11 query(int i) {
    11 sum=0;
   for(; i; i -= i&-i)
      sum += BIT[i];
   return sum;
}</pre>
```

7.2 BIT 2D

```
1l bit[MAXL][MAXC]; int L, C;
void update(int x, int y, int v) {
    for (; x <= L; x += x&-x)
        for (int yy = y; yy <= C; yy += yy&-yy)
        bit[x][yy] += v;
}
ll query(int x, int y) {
    ll sum = 0;
    for (; x; x -= x&-x)
        for (int yy = y; yy; yy -= yy&-yy)
        sum += bit[x][yy];
    return sum;
}</pre>
```

7.3 RMQ

```
#define MAXN 100000
int A[MAXN], T[4*MAXN];
int N; // #number of elements in A[]

void build(int no, int a, int b) {
   if (a == b) {
      T[no] = a;
   }else{
```

```
int m = (a + b)/2;
        build(2*no, a, m);
        build(2*no+1, m+1, b);
        int p1 = T[2*no], p2 = T[2*no+1];
        T[no] = (A[p1] \le A[p2]) ? p1 : p2;
// Faz A[i] = v
void update(int no, int a, int b, int i, int v) {
  if (a > i || b < i) return;
    if (a == i && b == i) {
   A[i] = v;
    T[no] = i;
    return:
  int m = (a + b)/2;
  update(2*no, a, m, i, v);
  update(2*no+1, m+1, b, i, v);
  int p1 = T[2*no], p2 = T[2*no+1];
  T[no] = (A[p1] \le A[p2]) ? p1 : p2;
// Retorna o idx do menor valor do intervalo [i, j]
int guery(int no, int a, int b, int i, int j){
 if (a > j || b < i) return -1;
   if (a >= i && b <= j) return T[no];</pre>
    int m = (a + b)/2;
  int p1 = query(2*no, a, m, i, j);
  int p^2 = query(2*no+1, m+1, b, i, j);
   if (p1 == -1) return p2;
   if (p2 == -1) return p1;
    return (A[p1] \le A[p2]) ? p1 : p2;
```

7.4 RMQ-Aplicações

```
//3. Determinar qual o no mais a esquerda na arvore tem valor
    maior que K, e adicionar K a ele.
void update(int no, int b, int e, int K) {
    if (b == e) { A[b] += K; T[no] = b; return; }
    int m = (b+e)/2;
    if (A[T[no*2]] >= K) update(no*2, b, m, K);
    else if(A[T[no*2 + 1] >= K) update(no*2+1, m+1, e, K);
    if (A[T[2*no]] >= A[T[2*no + 1]]) T[no] = T[2*no];
    else T[no] = T[2*no + 1];
}
```

7.5 Seg Tree com Lazy

```
11 v[MAXN], tree[4*MAXN], lazy[4*MAXN];
void build_tree(int no, int a, int b) {
  if (a == b) {
    tree[no] = v[a];
    return;
  int m = (a + b)/2;
  build_tree(2*no, a, m);
  build tree(2*no+1, m+1,b);
  tree[no] = tree[no\star2] + tree[no\star2+1];
void propagate(int no, int a, int b) {
  if(lazy[no]){
    tree[no] += (b - a + 1) * lazy[no];
    if(a != b) {
      lazy[2*no] += lazy[no];
      lazy[2*no+1] += lazy[no];
    lazy[no] = 0;
void update(int no, int a, int b, int i, int j, ll value) {
  propagate(no, a, b);
  if(a > j || b < i) return;
  if(a >= i && b <= j){
   lazy[no] += value;
    propagate(no, a, b);
    return;
  int m = (a + b)/2;
  update(2*no, a, m, i, j, value);
  update (2*no+1, m+1, b, i, j, value);
  tree[no] = tree[no\star2] + tree[no\star2+1];
ll query(int no, int a, int b, int i, int j){
  propagate (no, a, b);
  if(a > j || b < i) return 0;
```

```
if(a >= i && b <= j) return tree[no];
int m = (a + b)/2;
ll q1 = query(2*no, a, m, i, j);
ll q2 = query(2*no+1, m+1, b, i, j);
return q1+q2;</pre>
```

7.6 Union-Find

```
#define MAXV 100000
int V, pai[MAXV], rank[MAXV], size[MAXV];
void init(){
  rep(i, 0, V)
    pai[i] = i, rank[i] = 0, size[i] = 1;
void find(int v) {
  if(v != pai[v]) pai[v] = find(pai[v]);
  return pai[v];
void join(int u, int v){
  u = find(u); v = find(v);
  if(u == v) return;
  if(rank[u] < rank[v]) swap(u, v);</pre>
  pai[v] = u; // add v na arvore de u
  size[u] += size[v];
  if (rank[u] == rank[v]) rank[u]++;
bool same_set(int u, int v){
  return find(u) == find(v);
```

7.7 Treap

```
#include <bits/stdc++.h>
using namespace std;

#define rep(i, a, b) for(int i=a; i<b; i++)

typedef struct node{
  int prior, size;
  int val; //value stored in the array
  int sum; //whatever info you want to maintain in segtree for
        each node
  int lazy; //whatever lazy update you want to do
  int rev;
  struct node *l,*r;
}node;

typedef node* pnode;
int sz(pnode t) {
  return t?t->size:0;
}
```

```
void upd_sz(pnode t) {
                                                                           operation(t);
  if (t) t->size=sz(t->1) +1+sz(t->r);
                                                                         pnode init(int val){
void lazy(pnode t) {
                                                                           pnode ret = new node;
  if(!t || t->lazv==-1)return;
                                                                           ret->prior=rand(); ret->size=1;
  t->val=t->lazy;//operation of lazy
                                                                           ret->val=val;
  t \rightarrow sum = t \rightarrow lazy * sz(t);
                                                                           ret->sum=val;ret->lazy=-1; ret->rev=0;
  if (t->1) t->1->lazy=t->lazy; //propagate lazy
                                                                           ret->l=NULL, ret->r=NULL;
  if (t->r) t->r->lazy=t->lazy;
                                                                           return ret;
  t \rightarrow lazy = -1;
                                                                         int range_query(pnode t,int l,int r){//[l,r]
void reset(pnode t){
                                                                           pnode L, mid, R;
  if(t)t->sum = t->val;//no need to reset lazy coz when we call
                                                                           split(t, L, mid, l-1);
       this lazy would itself be propagated
                                                                            split(mid,t,R,r-1);//note: r-1!!
                                                                           int ans = t->sum;
void combine(pnode& t,pnode l,pnode r){//combining two ranges
                                                                           merge(mid, L, t);
                                                                           merge(t, mid, R);
    of segtree
  if(!l || !r)return void(t = 1?l:r);
                                                                           return ans:
  t \rightarrow sum = 1 \rightarrow sum + r \rightarrow sum;
                                                                         void range_update(pnode t, int l, int r, int val) {//[l, r]
void operation(pnode t){//operation of segtree
                                                                           pnode L.mid.R:
  if(!t)return;
                                                                           split(t, L, mid, l-1);
  reset(t); //reset the value of current node assuming it now
                                                                           split(mid,t,R,r-1);//note: r-1!!
      represents a single element of the array
                                                                           t->lazy=val; //lazy_update
  lazy(t->1); lazy(t->r); //imp: propagate lazy before combining t
                                                                           merge(mid, L, t);
      ->1, t->r;
                                                                           merge(t, mid, R);
  combine (t, t->1, t);
  combine (t, t, t->r);
                                                                         void reverse(pnode t, int 1, int r){
                                                                           pnode L, mid, R;
void push(pnode t) {
                                                                           split(t, L, mid, l-1);
  if(!t || !t->rev) return;
                                                                           split (mid, mid, R, r-l);
                                                                           mid->rev^=true;
  t->rev=false;
  swap(t->1, t->r);
                                                                           merge(t, L, mid);
  if(t->1) t->1->rev ^= true;
                                                                           merge(t, t, R);
  if(t->r) t->r->rev ^= true;
                                                                         void output(pnode t) {
void split(pnode t,pnode &1,pnode &r,int pos,int add=0) {
                                                                           if (!t) return;
  if(!t)return void(l=r=NULL);
                                                                           push (t);
  push(t); lazy(t);
                                                                           lazy(t);
  int curr pos = add + sz(t->1);
                                                                           output (t->1);
  if(curr_pos<=pos) //element at pos goes to left subtree(1)</pre>
                                                                           printf ("%d ", t->val);
    split(t->r,t->r,r,pos,curr_pos+1),l=t;
                                                                           output (t->r);
  else
    split(t->1,1,t->1,pos,add),r=t;
                                                                         int valor(int val){
  upd_sz(t);
                                                                           return val&1?0:1;
  operation(t);
void merge(pnode &t,pnode 1,pnode r) { //1->leftarray,r->
    rightarray, t->resulting array
                                                                         int main(){
  push(1);push(r);lazy(1);lazy(r);
                                                                           int P, Q;
  if(!l || !r) t = l?l:r;
                                                                           while(scanf("%d %d", &P, &Q) !=EOF){
  else if(l->prior>r->prior)merge(l->r,l->r,r),t=1;
                                                                             pnode tree=NULL, T1=NULL, T2=NULL, T3=NULL; int val;
        merge(r->1,1,r->1), t=r;
                                                                             rep(i, 0, P){
                                                                                scanf("%d", &val);
  upd_sz(t);
```

```
split(tree, T1, T2, i);
   merge(T1, T1, init(valor(val)));
   merge(tree, T1, T2);
}
while(Q--){ }
}
```

7.8 Seg Tree 2D

```
struct node{
  int qt;
  int f1, f2, f3, f4;
node new_node(){
  node ret;
  ret.qt=ret.f1=ret.f2=ret.f3=ret.f4=0;
  return ret;
vector <node> tree;
int cnt = 0;
bool inRange(int x1, int x2, int y1, int y2, int a1, int a2,
    int b1, int b2) {
  if (x2 < x1 \mid | y2 < y1) return false;
  if (x2 < a1 \mid | x1 > a2) return false;
  if(y2 < b1 || y1 > b2) return false;
  return true;
}
void update(int no, int x1, int x2, int y1, int y2, int a1, int
     a2, int b1, int b2, int val) {
  if (no == cnt) tree[cnt++] = new node();
  if (x1>=a1 && x2<=a2 && y1>=b1 && y2<=b2) {
    tree[no].qt = val;
    return;
  int f1=0, f2=0, f3=0, f4=0;
  if (inRange (x1, (x1+x2)/2, y1, (y1+y2)/2, a1, a2, b1, b2)) {
    if(!tree[no].fl) tree[no].fl = cnt;
    update(tree[no].f1, x1, (x1+x2)/2, y1, (y1+y2)/2, a1, a2, b
        1, b2, val);
  if (inRange(x1, (x1+x2)/2, (y1+y2)/2+1, y2, a1, a2, b1, b2)) {
    if(!tree[no].f2) tree[no].f2 = cnt;
    update(tree[no].f2, x1, (x1+x2)/2, (y1+y2)/2+1, y2, a1, a2,
         b_1, b_2, val);
  if (inRange((x1+x2)/2+1, x2, y1, (y1+y2)/2, a1, a2, b1, b2)){
```

```
if(!tree[no].f3) tree[no].f3 = cnt;
    update(tree[no].f3, (x1+x2)/2+1, x2, y1, (y1+y2)/2, a1, a2,
         b_1, b_2, val);
  if (inRange((x1+x2)/2+1, x2, (y1+y2)/2+1, y2, a1, a2, b1, b2))
    if(!tree[no].f4) tree[no].f4 = cnt;
    update(tree[no].f4, (x1+x2)/2+1, x2, (y1+y2)/2+1, y2, a1, a
        2, b1, b2, val);
  if(tree[no].fl) fl=tree[tree[no].fl].qt;
  if(tree[no].f2) f2=tree[tree[no].f2].qt;
  if(tree[no].f3) f3=tree[tree[no].f3].qt;
  if (tree[no].f4) f4=tree[tree[no].f4].qt;
 tree[no].qt=f_1+f_2+f_3+f_4;
int query(int no, int x1, int x2, int y1, int y2, int a1, int a
   2, int b1, int b2) {
  if(!inRange(x1, x2, y1, y2, a1, a2, b1, b2) || no>=cnt ||
      tree[no].qt==0) return 0;
  if(x1>=a1 && x2<=a2 && y1>=b1 && y2<=b2) return tree[no].qt;</pre>
  int f1=0, f2=0, f3=0, f4=0;
  if (tree[no].f1) f1 = query(tree[no].f1, x1, (x1+x2)/2, y1, (y))
      1+y^2)/2, a1, a2, b1, b2);
  if (tree[no].f2) f2 = query(tree[no].f2, x1, (x1+x2)/2, (y1+y2)
      )/2+1, y2, a1, a2, b1, b2);
  if(tree[no].f3) f3 = query(tree[no].f3, (x1+x2)/2+1, x2, y1,
      (y_1+y_2)/2, a1, a2, b1, b2);
  if (tree[no].f4) f4 = query(tree[no].f4, (x1+x2)/2+1, x2, (y1+x2)/2+1)
      y^2)/2+1, y^2, a1, a2, b1, b2);
  return f1+f2+f3+f4;
void erase(){
  tree.clear();
  vector <node> xua;
  swap(tree, xua);
  tree.resize(1000010);
  cnt = 0;
```

7.9 Polyce

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
```

```
#include <ext/rope>
using namespace std;
using namespace __gnu_cxx;
using namespace __qnu_pbds;
typedef tree <int, null_type, less<int>, rb_tree_tag,
    tree_order_statistics_node_update> ordered_set;
int main(){
  ordered set T; ordered set :: iterator it;
  T.insert(); T.erase(); int k = T.find_by_order(); int kk = T.
      order_of_key();
  rope <int> v; rope <int> :: iterator it;
  v.push_back();
  rope \langle int \rangle cur = v.substr(l, r - l + 1);
  for( it = cur.mutable_begin(); it != cur.mutable_end(); it++
     ) cout << *it << " ";
  v.erase(1, r - 1 + 1);
  v.insert(v.mutable_begin(), cur);
```

7.10 KD2

```
struct point{
  int x, y, z;
  point (int x=0, int y=0, int z=0): x(x), y(y), z(z) {}
  point operator-(point q) {
    return point(x-q.x, y-q.y, z-q.z);
  int operator*(point q){
    return x*q.x + y*q.y + z*q.z;
};
typedef vector<point> polygon;
priority_queue <double> vans;
int NN, CC, KK, DD;
struct KDTreeNode{
  point p;
  int level;
  KDTreeNode *below, *above;
  KDTreeNode (const point& q, int lev1) {
    p = q;
    level = levl;
    below = above = 0;
  ~KDTreeNode(){
    delete below, above;
```

```
int diff (const point& pt) {
    switch (level)
      case 0:
        return pt.x - p.x;
      case 1:
        return pt.y - p.y;
      case 2:
        return pt.z - p.z;
    return 0;
  11 distSq (point& q) {
    return (p-q) * (p-q);
  int rangeCount (point& pt, ll K) {
    int count = (distSq(pt) <= K*K) ? 1 : 0;</pre>
    if(count) vans.push(-sqrt(distSq(pt)));
    int d = diff(pt);
    if (~d <= K && above != 0)
      count += above->rangeCount(pt, K);
    if (d <= K && below != 0)
      count += below->rangeCount(pt, K);
    return count;
};
class KDTree {
  public:
    polygon P;
    KDTreeNode *root;
    int dimention;
    KDTree() {}
    KDTree (polygon &poly, int D) {
      P = poly;
      dimention = D;
      root = 0;
      build();
    ~KDTree(){
      delete root;
    //count the number of pairs that has a distance less than K
    11 countPairs(11 K) {
      11 count = 0;
      rep(i, 0, P.size())
        count += root->rangeCount(P[i], K) - 1;
      return count;
  protected:
    void build() {
      //random_shuffle(all(P));
```

```
rep(i, 0, P.size()){
        root = insert(root, P[i], -1);
    KDTreeNode *insert(KDTreeNode* t, const point& pt, int
       parentLevel) {
      if (t == 0) {
        t = new KDTreeNode (pt, (parentLevel+1) % dimention);
      else{
        int d = t->diff(pt);
        if (d <= 0) t->below = insert (t->below, pt, t->level);
        else t->above = insert (t->above, pt, t->level);
      return t;
    }
} ;
int main(){
  point e; e.z=0;
 polygon p;
  set <ii> st;
  while (scanf ("%d %d %d %d", &NN, &CC, &KK, &DD)!=EOF) {
    p.clear(); KK=min(NN, KK); st.clear();
    rep(i, 0, NN) {
      scanf("%d %d", &e.x, &e.y);
      st.insert(mp(e.x, e.y));
      p.pb(e);
    KDTree tree(p, 2);
    int ans=0;
    rep(i, 0, CC) {
      scanf("%d %d", &e.x, &e.y);
      if(st.count(mp(e.x, e.y))) continue;
      11 at=0;
      rep(i, 0, 30) {
        at=ll(1)<<i; while(!vans.empty()) vans.pop();
        int aux=tree.root->rangeCount(e, at);
        if(aux>=KK) break;
      double sum=0.0;
      rep(i, 0, KK) {
        sum+=-vans.top();
        vans.pop();
      if(sum>=DD) ans++;
    printf("%d\n", ans);
```

```
return 0;
```

8 Strings

8.1 KMP

```
/*obs: A funcao strstr (char* text, char* pattern) da
   biblioteca <cstring> implementa KMP (C-ANSI). A funcao
   retorna a primeira ocorrencia do padrao no texto, KMP
   retorna todas.
nres -> O numero de ocorrencias do padrao no texto
res[] -> posicoes das nres ocorrencias do padrao no texto
Complexidade do algoritmo: O(n+m) */
#define MAXN 100001
int pi[MAXN], res[MAXN], nres;
void kmp(string text, string pattern) {
  nres = 0;
  pi[0] = -1;
  rep(i, 1, pattern.size()){
        pi[i] = pi[i-1];
        while (pi[i] >= 0 && pattern[pi[i]+1] != pattern[i])
            pi[i] = pi[pi[i]];
        if (pattern[pi[i]+1] == pattern[i])
            ++pi[i];
  int k = -1; //k+1 eh o tamanho do match atual
    rep(i, 0, text.size()){
        while (k \ge 0 \&\& pattern[k+1] != text[i])
            k = pi[k];
        if (pattern[k+1] == text[i])
            ++k;
        if (k+1 == pattern.size()) {
            res[nres++] = i-k;
            k = pi[k];
```

8.2 Aho Corasick

```
const int cc = 26;
const int MAX = 100;

int cnt;
int sig[MAX][cc];
int term[MAX];
int T[MAX];
int v[MAX];
int v[MAX];
```

```
return c - '0';
void add( string s, int id ){
  int x = 0;
 rep(i, 0, s.size()){
   int c = C(s[i]);
   if(!sig[x][c]){
     term[cnt] = 0;
      sig[x][c] = cnt++;
   x = sig[x][c];
  term[x] = 1;
  v[id] = x;
void aho() {
  queue < int > q;
 rep(i, 0, cc){
   int x = siq[0][i];
   if( !x ) continue;
    q.push(x);
   T[x] = 0;
  while( !q.empty() ){
   int u = q.front();
   q.pop();
    rep(i, 0, cc){
      int x = sig[u][i];
      if( !x ) continue;
      int v = T[u];
      while( v && !sig[v][i] ) v = T[v];
      v = sig[v][i];
     T[x] = v;
      term[x] += term[v];
      q.push(x);
// Conta a quantidade de palavras de exatamente l caracteres
   que se pode formar com um determinado alfabeto, dado que
   algumas palavras so "proibidas"
int mod = 1e9+7;
ll pd[100][MAX];
11 solve( int pos, int no ) {
 if( pos == 0 ) return 1;
  if( pd[pos][no] != -1 ) return pd[pos][no];
 ll ans = 0;
  rep(i, 0, cc){
   int v = no;
```

```
while( v && !sig[v][i] ) v = T[v];
    v = siq[v][i];
    if( term[v] ) continue;
    ans = (ans + solve(pos-1, v)) mod;
  return pd[pos][no] = ans;
void Qttd_de_Palavras() {
  while(1){
    memset( sig, 0, sizeof sig );
    memset (pd, -1, sizeof pd);
    cnt = 1;
    int l = readInt();
    if( !l ) break;
    int n = readInt();
    string pattern;
   rep(i, 0, n){
     cin >> pattern;
     add(pattern, i);
    aho();
    ll ans = 0;
    rep(i, 1, 1+1)
     ans = (ans + solve(i, 0)) mod;
    printf("%d\n", ans);
// Verifica quais padres ocorreram em um texto
int alc[MAX];
void busca( string s ) {
 int x = 0;
 rep(i, 0, s.size()){
   int c = C(s[i]);
   while ( x \& \& !siq[x][c] ) x = T[x];
   x = sig[x][c];
    alc[x] = 1;
void Ql_Ocorreu() {
  string pattern, text;
  while( getline(cin, text )){
    if( text == "*" ) break;
    memset( siq, 0, sizeof siq );
    memset( alc, 0, sizeof alc );
    cnt = 1;
    int n;
    cin >> n;
    rep(i, 0, n){
```

```
cin >> pattern;
      add( pattern, i );
    aho();
    busca (text);
    for( int i = cnt-1; i >= 0; i-- ){
      if(alc[i]) alc[T[i]] = 1;
    rep(i, 0, n){
      int u = v[i];
      if( alc[u] ) printf("Ocorreu\n");
      else printf("Nao ocorreu\n");
  }
// Total de ocorrencias de cada padro em uma string, mesmo com
    sufixos iquais
11 busca2( string s ) {
 11 x = 0, cont = 0;
  rep(i, 0, s.size()){
    int c = C(s[i]);
   while (x \& \& !sig[x][c]) x = T[x];
   x = siq[x][c];
    cont += term[x];
  return cont;
void Onts vezes Ocorreu() {
  string text, pattern;
  while( cin >> text ) {
   if( text == "*" ) break;
   memset( sig, 0, sizeof sig );
   cnt = 1;
   int n = readInt();
   rep(i, 0, n){
     cin >> pattern;
      add(pattern, i);
    aho();
    rep(i, 1, 10) debug(T[i])
    cout << busca2( text ) << endl;</pre>
  }
//Encontra a primeira ocorrencia de cada padro em uma string
void busca3( string s ) {
  int x = 0;
  rep(i, 0, s.size()){
   int c = C(s[i]);
   while (x \& \& !sig[x][c]) x = T[x];
   x = sig[x][c];
    if( !alc[x] ) alc[x] = i+1;
```

```
void Onde_Ocorreu() {
  string pattern, text;
 int tam[1000];
 while( cin >> text ) {
   if( text == "*" ) break;
   memset( sig, 0, sizeof sig );
   memset( alc, 0, sizeof alc );
   cnt = 1;
   int n;cin >> n;
   rep(i, 0, n){
     cin >> pattern;
     tam[i] = pattern.size();
     add(pattern, i);
    aho();
   busca3 (text);
    for( int i = cnt-1; i >= 0; i-- ){
      alc[ T[i] ] = min( alc[i], alc[ T[i] ] );
   rep(i, 0, n){
     int u = v[i];
     if( alc[u] != INF ) {
        int k = alc[u] - tam[i] + 1;
        printf("De %d a %d\n", k, alc[u]);
      else
        printf("Nao ocorreu\n");
```

8.3 Suffix Array

```
#define MAX 100010
#define MAX_N 100010
char T[MAX_N];
ll n;
int RA[MAX_N], tempRA[MAX_N];
int SA[MAX_N], tempSA[MAX_N];
int c[MAX_N];
int Phi[MAX_N], PLCP[MAX_N], LCP[MAX_N];

void countingSort(int k) {
  int i, sum, maxi = max((ll)300, n);
  memset(c, 0, sizeof c);
  for (i = 0; i < n; i++)
    c[i + k < n ? RA[i + k] : 0]++;
  for (i = sum = 0; i < maxi; i++) {
    int t = c[i]; c[i] = sum; sum += t;</pre>
```

```
for (i = 0; i < n; i++)</pre>
    tempSA[c[SA[i]+k < n ? RA[SA[i]+k] : 0]++] = SA[i];
  for (i = 0; i < n; i++)
    SA[i] = tempSA[i];
void constructSA() {
  int i, k, r;
  for (i = 0; i < n; i++) RA[i] = T[i];
  for (i = 0; i < n; i++) SA[i] = i;
  for (k = 1; k < n; k <<= 1) {
    countingSort(k);
    countingSort(0);
    tempRA[SA[0]] = r = 0;
    for (i = 1; i < n; i++)
      tempRA[SA[i]] = (RA[SA[i]] == RA[SA[i-1]] \&\& RA[SA[i]+k]
          == RA[SA[i-1]+k]) ? r : ++r;
    for (i = 0; i < n; i++)
      RA[i] = tempRA[i];
    if (RA[SA[n-1]] == n-1) break;
}
void computeLCP() {
  int i, L;
  Phi[SA[0]] = -1;
  for (i = 1; i < n; i++)
    Phi[SA[i]] = SA[i-1];
  for (i = L = 0; i < n; i++) {
    if (Phi[i] == -1) { PLCP[i] = 0; continue; }
    while (T[i + L] == T[Phi[i] + L]) L++;
    PLCP[i] = L;
    L = \max(L-1, 0);
  for (i = 0; i < n; i++) {
    LCP[i] = PLCP[SA[i]];
int main(){
  // concatenar $ no final
```

8.4 Suffix Array (Gugu)

```
#include <bits/stdc++.h>
//LIFE IS NOT A PROBLEM TO BE SOLVED
using namespace std;
#define rep(i,a,b) for(int i = int(a); i < int(b) ; i++)
typedef long long int ll;</pre>
```

```
const int MAX = 100010;
int gap, tam, sa[MAX], pos[MAX], lcp[MAX], tmp[MAX];
bool sufixCmp(int i, int j) {
  if(pos[i]!=pos[j]) return pos[i]<pos[j];</pre>
  i+=qap, j+=qap;
  return (i<tam && j<tam) ? pos[i]<pos[j] : i>j;
void buildSA(char s[]) {
  tam=strlen(s);
  for(int i=0; i<tam; i++)</pre>
    sa[i]=i, pos[i]=s[i], tmp[i]=0;
  for(gap=1;;gap*=2){
    sort(sa, sa+tam, sufixCmp);
    tmp[0]=0;
    for (int i=0; i<tam-1; i++)</pre>
      tmp[i+1] = tmp[i] + sufixCmp(sa[i], sa[i+1]);
    for(int i=0; i<tam; i++)
      pos[sa[i]]=tmp[i];
    if(tmp[tam-1] == tam-1) break;
11 buildLCP(char s[]) {
  11 \text{ sum}=0;
  for (int i=0, k=0; i<tam; i++) {</pre>
    if (pos[i] == tam-1) continue;
    for (int j=sa[pos[i]+1]; s[i+k]==s[j+k];) k++;
    lcp[pos[i]+1] = k;
    sum+=k;
    if(k>0) k--;
  return sum;
void PrintAll(char s[]){
  printf("SA\ttam\tLCP\tSuffix\n");
  rep(i, 0, tam)
    printf("%2d\t%2d\t%2d\t%s\n", sa[i], tam-sa[i], lcp[i], s +
11 num_subs(ll m) {
  return (11) tam* (tam+1) /2 - m;
11 num_subsrn() {
  11 \text{ ret}=0;
  rep(i, 1, tam)
    if(lcp[i]>lcp[i-1])
      ret+=lcp[i]-lcp[i-1];
  return ret;
void printans(char s[], int n) {
  int maior=0, id=-1;
  rep(i, 0, tam)
```

```
if(lcp[i]>n && lcp[i]>maior)
     maior=lcp[i], id=i;
  if(id==-1) printf("*");
 else
    rep(i, sa[id], sa[id]+maior)
     printf("%c", s[i]);
 printf("\n");
char s[MAX];
int main(){
  while(1){
    scanf("%s", s);
   if(s[0]=='*') break;
   buildSA(s);
   11 m=buildLCP(s);
   PrintAll(s); //printa sa, lcp, suffixs
   //printf("%lld\n", num_subs(m)); //numero de substrings nao
         repetidas
    //printf("%lld\n", num_subsrn()); //numero de substrings
       que se repete
    //printans(s, 2); //maior substring de tamanho maior ou
       igual a n que se repete
```

8.5 Rolling Hash

```
//Permite encontrar um hash de uma substring de S. precompute O
    (n), my_hash O(1)
#define NN 1000006
const 11 mod = 1e9+7; //modulo do hash
const 11 x = 33; //num. primo > que o maior caracter de S.
11 H[NN], X[NN];
11 V(char c) { return c-'A'; }
ll my_hash(int i, int j){
  ll ret = H[j];
  if(!i) return ret;
  return ((ret-(H[i-1]*X[j-i+1])%mod)+mod)%mod;
void precompute(string s){
 X[0] = 1;
  rep(i, 1, NN)
   X[i] = (X[i-1]*x) % mod;
  H[0] = V(s[0]);
  rep(i, 1, s.size())
    H[i] = ((H[i-1]*x)%mod + V(s[i]))%mod;
```

8.6 Longest Common Prefix with Hash

```
// Longest Commom Prefix between S[i..] and S[j..]
int lop(int i, int j, int tam) {
  int lo = 0, hi = tam, ans;
  while(lo <= hi) {
    int mid = (lo+hi)/2;
    if(my_hash(i, i+mid-1) == my_hash(j, j+mid-1)) {
      ans = mid;
      lo = mid+1;
    }
    else hi = mid-1;
  }
  return ans;
}</pre>
```

8.7 Minimum Lexicographic Rotation

```
/*Retorna a menor string lexicografica de s
Necessario my_hash() e lcp()*/
string min_lex_rot(string s) {
   int t = s.size();
   precompute(s); // hashing
   s += s;
   int idx = 0;
   for(int i = 1; i < t; i++) {
        // tam do prefix comum
        int len = lcp(i, idx, t);
        if (s[i+len] < s[idx+len])
            idx = i;
   }
   return s.substr(idx, t);
}</pre>
```

8.8 Longest Palindrome (Manacher algorithm)

```
string preProcess(string s) {
   int n = s.length();
   if (n == 0) return "^$";

   string ret = "^";
   for (int i = 0; i < n; i++)
     ret += "#" + s.substr(i, 1);
   ret += "#$";
   return ret;
}

string longestPalindrome(string s) {
   L=C=s.size();
   string T = preProcess(s);
   int n = T.length();
   int *P = new int[n];
   int C = 0, R = 0;</pre>
```

```
for (int i = 1; i < n-1; i++) {</pre>
 int i mirror = 2*C-i;
 P[i] = (R > i) ? min(R-i, P[i_mirror]) : 0;
 while (T[i + 1 + P[i]] == T[i - 1 - P[i]])
   P[i]++;
  if (i + P[i] > R) {
   C = i;
   R = i + P[i];
}
int maxLen = 0;
int centerIndex = 0;
for (int i = 1; i < n-1; i++) {
 if(!P[i]) continue;
 if (P[i] > maxLen) {
   maxLen = P[i];
    centerIndex = i;
delete[] P;
return s.substr((centerIndex - 1 - maxLen)/2, maxLen);
```

8.9 Autômato de Sufixos

```
struct state{
  int len, link;
 int next[26];
};
const int MAXN = 200020;
state st[2*MAXN]; //vetor que armazena os estados
int sz; //contador do numero de estados
int last; //numero do estado que corresponde ao texto todo
void sa_init(){
  sz = 1;
 last = 0:
  st[0].len = 0;
  st[0].link = -1;
  rep(i, 0, 26) st[0].next[i]=0;
  //limpa o mapeamento de transicoes
void sa_extend (int c, ll &ans) {
  int cur = sz++; // novo estado a ser criado
  st[cur].len = st[last].len + 1;
  rep(i, 0, 26) st[cur].next[i]=0;
  int p; // variavel que itera sobre os estados terminais
  for(p=last; p!=-1 && !st[p].next[c]; p=st[p].link){
    st[p].next[c] = cur;
```

```
if(p == -1) \{ / / \text{ nao ocorreu transicao } c \text{ nos estados terminais} \}
    st[cur].link = 0;
    ans+=st[cur].len;
  }else{ // ocorreu transicao c no estado p
    int q = st[p].next[c];
    if(st[p].len + 1 == st[q].len) {
      st[cur].link = q;
    }else{
      int clone = sz++; // criacao do vertice clone de q
      st[clone].len = st[p].len + 1;
      rep(i, 0, 26) st[clone].next[i]=st[q].next[i];
      st[clone].link = st[q].link;
      for(; p!=-1 && st[p].next[c]==q; p=st[p].link){//
          atualização das transicoes c
        st[p].next[c] = clone;
      st[q].link = st[cur].link = clone;
    ans+=st[cur].len-st[st[cur].link].len;
  // atualizacao do estado que corresponde ao texto
  last = cur;
bool busca_automato(int m, string p) {
  int i, pos=0;
  for( i=0 ; i<m ; i++ ) {</pre>
   if( st[pos].next[p[i]]==0 ){
      return false;
    }else{
      pos=st[pos].next[p[i]];
  return true;
int maior_tamanho_em_comum( string s, string t ){
  11 nothing=0;
  // Constroi o automato com o primeiro texto
  sa_init();
  for (int i=0; i < (int) s.size(); i++) sa_extend (s[i]-'a',</pre>
      nothing);
  int estado = 0, tamanho = 0, maior = 0;
  // Passando pelos caracteres do segundo texto
  for (int i=0; i<(int)t.size(); ++i){</pre>
    while( estado && ! st[estado].next[t[i]-'a'] ){
      estado = st[estado].link;
      tamanho = st[estado].len;
   if (st[estado].next[t[i]-'a']) {
      estado = st[estado].next[t[i]-'a'];
      tamanho++;
```

```
if(tamanho > maior) {
    maior = tamanho;
}

return maior;
}

int main() {
    char s[MAXN];
    char p[MAXN];
    while(gets(s)) {
        sa_init();
        int tam=strlen(s); ll ans=0;
        rep(i, 0, tam) {
            sa_extend(s[i]-'a', ans);
        }
        gets(p);
        printf("%d\n", maior_tamanho_em_comum(s, p));
}

return 0;
}
```

9 PD

9.1 Soma acumulada 2D

```
/*Retorna o somatorio dos elementos de uma submatriz em O(1)
Submatriz definida por canto superior esquerdo [a,b] e canto
    inferior
direito (x,y) .: a \le x \&\& b \le y
#define MAXN 1000
int N, M; // linhas colunas
long long V[MAXN+2][MAXN+2]; // matriz da entrada
long long S[MAXN+2][MAXN+2]; // matriz com as somas acumuladas
// precomputa as somas em O(N*M)
void precal() {
  rep(x, 0, N)
    rep(y, 0, M)
      S[x][y] = V[x][y];
      if (x>0) S[x][y] += S[x-1][y];
      if (y>0) S[x][y] += S[x][y-1];
      if (x>0 \&\& y>0) S[x][y] -= S[x-1][y-1];
    }
// retorna a soma da submatriz
long long sum(int a, int b, int x, int y) {
  long long soma = S[x][y];
  if (a>0) soma -= S[a-1][y];
  if (b>0) soma -= S[x][b-1];
```

```
if (a>0 && b>0) soma += S[a-1][b-1];
return soma;
```

9.2 Knuth Optimization

```
int N, B, C, yep, save[MAXN][MAXN], sav[MAXN];
11 n[MAXN], mc[MAXN][MAXN], se[MAXN], sd[MAXN], pd[MAXN][MAXN];
11 solve(int i, int k) {
  if(i==N) return 0;
  if(k==1) return pd[i][k] = mc[i][N-1];
  if (pd[i][k]!=-1) return pd[i][k];
  ll ret=LINF; int ini=i, fim=N-k+1, best = -1;
  if(i && save[i-1][k]) ini=save[i-1][k];
  if (save[i][k-1]) fim=save[i][k-1]+1;
  rep(l, ini, fim) {
    11 \text{ aux} = \text{solve}(1+1, k-1) + mc[i][1];
    if(ret>aux) {
      best=1;
      ret=aux;
  save[i][k]=best;
  return pd[i][k]=ret;
int main()
  rep(i, 0, N) scanf("%lld", &n[i]);
  se[0]=n[0];
  rep(i, 1, N) se[i] = se[i-1] + n[i];
  sd[N-1]=n[N-1];
  for (int i=N-2; i>=0; i--) sd[i]=sd[i+1]+n[i];
  rep(i, 1, N) pd[0][i]=pd[0][i-1]+se[i-1];
  for (int i=N-2; i>=0; i--) pd[N-1][i]=pd[N-1][i+1]+sd[i+1];
  rep(i, 1, N)
    rep(j, i+1, N)
      pd[i][j]=pd[i-1][j]-n[i-1]*(j-i+1);
  for(int i=N-2; i>=0; i--) {
    for(int j=i-1; j>=0; j--)
      pd[i][j]=pd[i+1][j]-n[i+1]*(i-j+1);
  rep(i, 0, N) {
    if(pd[i][i+1] < pd[i+1][i]) mc[i][i+1]=pd[i][i+1], save[i][</pre>
```

```
i+1]=i+1;
  else
                  mc[i][i+1]=pd[i+1][i], save[i][i+1]=i;
  rep(j, i+2, N){
    int ini=save[i][j-1];
    mc[i][j]=pd[i][ini]+pd[j][ini], save[i][j]=ini;
    rep(k, ini+1, j+1) {
      ll a=pd[i][k]+pd[j][k];
      if (mc[i][j] <= a) break;
      mc[i][j]=a;
      save[i][j]=k;
  rep(j, 0, N+1){
    pd[i][j] = -1,
    save[i][i]=0;
}
rep(j, 0, N+1) pd[N][j]=-1, save[N][j]=0;
solve();
return 0:
```

9.3 Convex Hull Trick

```
bool bad(int 11, int 12, int 13) {
  return (B[13]-B[11]) * (M[11]-M[12]) < (B[12]-B[11]) * (M[11]-M[13
      1);
void add(long long m, long long b) {
  M.push_back(m);
  B.push back(b);
  while (M.size() >= 3 \& \&bad(M.size() - 3, M.size() - 2, M.size() - 1)) {
    M.erase(M.end()-2);
    B.erase (B.end()-2);
  }
long long query(long long x) {
  if (pointer>=M.size())
    pointer=M.size()-1;
  while (pointer<M.size()-1&&
    M[pointer+1] *x+B[pointer+1] <M[pointer] *x+B[pointer])</pre>
    pointer++;
  return M[pointer] *x+B[pointer];
struct hux{
  int a, b, id;
};
bool my_sort(hux a, hux b) {
```

```
return a.b!=b.b ? a.b > b.b : a.a > b.a;
const ll LINF = 1LL << 52;</pre>
const double EPS = 1e-9;
const int MAXV = 100010;
double intersept(hux a, hux b) {
  return double(b.b-a.b)/(a.a-b.a);
vector < pair< double, double > > convex_hux( const vector <hux</pre>
   > &v ) {
  int p=0, n=v.size(), bestai=v[0].a;
  double cross=0.0;
  pair<double, int> aux;
  priority_queue < pair<double, int> > pq;
  vector < pair< double, double > > ret(n+1, mp(-1, -1));
  pq.push(mp(cross, p));
  ret[v[p].id].F=cross,
  ret[v[p].id].S=LINF;
  rep(i, 1, n) {
    aux=pq.top();
    cross=aux.F,
    p=aux.S;
    if(v[i].a <= bestai) continue; bestai=v[i].a;</pre>
    double new_cross = intersept(v[i], v[p]);
    while( new_cross <= cross + EPS ) {</pre>
      pq.pop();
      ret[v[p].id]=mp(-1.0, -1.0);
      aux=pq.top();
      cross=aux.F,
      p=aux.S;
      new_cross = intersept(v[i], v[p]);
    pq.push(mp(new_cross, i));
    ret[v[p].id].S=new cross;
    ret[v[i].id].F=new cross;
    ret[v[i].id].S=LINF;
  }
  //rep(i, 0, n) cout << ret[i].F << " " << ret[i].S << "\n";
```

```
return ret;
```

9.4 Longest Increasing Subsequence

```
//Maior subsequencia crescente
#define MAX N 100
int vet[MAX_N], P[MAX_N], N;
void reconstruct_print(int end){
  int x = end;
  stack<int> s:
  while (P[x] >= 0) {
    s.push(vet[x]);
   x = P[x];
  printf("%d", vet[x]);
  while(!s.empty()){
   printf(", %d", s.top());
    s.pop();
int lis(){
 int L[MAX_N], L_id[MAX_N];
  int li = 0, lf = 0; //lis ini, lis end
  rep(i, 0, N) {
   int pos = lower_bound(L, L+li, vet[i]) - L;
   L[pos] = vet[i];
   L_id[pos] = i;
   P[i] = pos ? L_id[pos - 1] : -1;
   if (pos + 1 > li) {
     li = pos + 1;
      lf = i;
  reconstruct_print(lf);
  return li;
```

9.5 Kadane 1D

```
// Encontra maior soma contigua positiva num vetor em O(N)
// [s,f] contem o intervalo de maior soma
int KadanelD(int vet[], int N, int &s, int &f) {
  int ret = -INF, sum, saux;
  sum = s = f = saux = 0;
  rep(i, 0, N) {
    sum += vet[i];
    if(sum > ret) { ret = sum; s = saux; f = i; }
    if(sum < 0) { sum = 0; saux = i+1; }
}
return ret;</pre>
```

9.6 Kadane 2D

```
//maior soma de uma sub-matriz a partir de valores positivos
//[x1,y1] = upper - left, [x2,y2] = bottom - right
int L, C, pd[MAX_L], mat[MAX_L][MAX_C];
int x1, y1, x2, y2;
int Kadane2D(){
 int ret = 0, aux;
  rep(left, 0, C){
    rep(i, 0, L) pd[i] = 0;
    rep(right, left, C){
      rep(i, 0, L) pd[i] += mat[i][right];
      int sum = aux = 0;
      rep(i, 0, L) { // Kadane1D
        sum += pd[i];
        if(sum>ret)
          ret=sum, x1=aux, y1=left, x2=i, y2=right;
        if(sum < 0) sum = 0, aux = i+1;
    }
  return ret;
```

9.7 Merge Sort com num de Inversoes

```
void merge(int* arr, int size1, int size2, ll &inversions) {
  int temp[size1+size2+2];
  int ptr1=0, ptr2=0;
  while (ptr1+ptr2 < size1+size2) {</pre>
    if (ptr1 < size1 && arr[ptr1] <= arr[size1+ptr2] || ptr1 <</pre>
        size1 \&\& ptr2 >= size2)
      temp[ptr1+ptr2] = arr[ptr1++];
    if (ptr2 < size2 && arr[size1+ptr2] < arr[ptr1] || ptr2 <</pre>
        size2 \&\& ptr1 >= size1) {
      temp[ptr1+ptr2] = arr[size1+ptr2++];
      inversions += size1-ptr1;
  for (int i=0; i < size1+size2; i++)</pre>
    arr[i] = temp[i];
void mergeSort(int* arr, int size, ll &inversions) {
  if (size == 1) return;
  int size1 = size/2, size2 = size-size1;
```

```
mergeSort(arr, size1, inversions);
mergeSort(arr+size1, size2, inversions);
merge(arr, size1, size2, inversions);
}
```

9.8 Quick Sort

```
//No main, chamar quicksort(array, 0, tam-1);
int partition(int s[], int l, int h) {
  int i, p, firsthigh;
 p = h;
  firsthigh = 1;
  for (i=1; i < h; i++)</pre>
   if (s[i] < s[p])
      swap(s[i], s[firsthigh]);
      firsthigh++;
  swap(s[i], s[firsthigh]);
  return firsthigh;
void quicksort(int s[], int l, int h){
  int p;
  if ((h-1) > 0) {
   p = partition(s, l, h);
    quicksort (s, l, p-1);
    quicksort (s, p+1, h);
```

10 Miscelânia

10.1 Calendário

```
// converts Gregorian date to integer (Julian day number)
int dateToInt (int m, int d, int y) {
 return
 1461 * (y + 4800 + (m - 14) / 12) / 4 +
 367 * (m - 2 - (m - 14) / 12 * 12) / 12 -
 3 * ((y + 4900 + (m - 14) / 12) / 100) / 4 +
 d - 32075;
// converts integer (Julian day number) to Gregorian date:
   month/day/year
void intToDate (int jd, int &m, int &d, int &y) {
 int x, n, i, j;
 x = jd + 68569;
 n = 4 * x / 146097;
 x = (146097 * n + 3) / 4;
 i = (4000 * (x + 1)) / 1461001;
 x = 1461 * i / 4 - 31;
 j = 80 * x / 2447;
 d = x - 2447 * j / 80;
 x = i / 11;
 m = j + 2 - 12 * x;
 v = 100 * (n - 49) + i + x;
// converts integer (Julian day number) to day of week
string intToDay (int jd) {
 return dayOfWeek[jd % 7];
int main () {
  int jd = dateToInt (3, 24, 2004);
  int m, d, y;
 intToDate (jd, m, d, y);
  string day = intToDay (jd);
 // expected output:
      2453089
       3/24/2004
        Wed
  cout << jd << endl
  << m << "/" << d << "/" << y << endl
  << day << endl;
```