

CALL-BY-NAME, CALL-BY-VALUE, CALL-BY-NEED, AND THE LINEAR LAMBDA CALCULUS

SHORT TALK

Quentin Schroeder

MPRI - Université Paris-Cité

IDEA

MOTIVATION

Problem

We want to understand how to evaluate expressions in the lambda calculus.

Call by Name

1

MOTIVATION

Problem

We want to understand how to evaluate expressions in the lambda calculus.

- · Call by Name
- · Call by Value

MOTIVATION

Problem

We want to understand how to evaluate expressions in the lambda calculus.

- · Call by Name
- · Call by Value
- · Call by Need

CALL BY VALUE (OCAML)

```
let rec dumb_works n = if false then dumb_works n else 42
let () = print_int (dumb_works 0)
```

CALL BY VALUE (OCAML)

```
let rec dumb_works n = if false then dumb_works n else 42
let () = print_int (dumb_works 0)

let branch b l r = if b then l else r
let rec dumb_breaks n = branch false (dumb_breaks n) 42

let () = print_int (dumb_breaks 0)
```

CALL BY NEED (HASKELL)

```
dumb n = if (n == 0 && n == 1) then dumb n else 42
forceEvalBranch b l r = if b then l else r

dumb2 n = forceEvalBranch (n == 0 && n == 1) (dumb2 n) 42

main = do
    print $ dumb 0
    print $ dumb 1
    print $ dumb2 0
    print $ dumb2 1
```

CALL BY NEED (HASKELL)

Example

```
dumb n = if (n == 0 && n == 1) then dumb n else 42
forceEvalBranch b l r = if b then l else r

dumb2 n = forceEvalBranch (n == 0 && n == 1) (dumb2 n) 42
main = do
    print $ dumb 0
    print $ dumb 1
    print $ dumb 2 0
    print $ dumb2 1
```

```
project x y = x
loop x = loop x
main = print $ project 2 (loop 3)
```

CALL BY NAME (???)

Is there a call by name programming language?

Yes ... kind of (e.g. Algol 60)

Example

Consider the following:

```
let rec h n = if n = o then 1 else n * h (n - 1) * h (n - 1)
let printalot m =
   let () = print_endline (string_of_int m) in
   let () = print_endline (string_of_int m) in
   let () = print_endline (string_of_int m) in
   let () = print_endline (string_of_int m)
```

Preliminaries

PRELIMINARIES

Lambda Calculus

- · model of computation
- basis for functional programming languages

PRELIMINARIES

Linear Logic

- a resource sensitive logic
- · can be used to priority in evaluation of proof terms

PRELIMINARIES

Linear Logic

- · a resource sensitive logic
- can be used to priority in evaluation of proof terms

A massive leap of faith

- build a linear lambda calculus [Maraist et al.(1995)Maraist, Odersky, Turner, and Wadler,]
- show that it can be used to model execution strategies

Why do we care?

- learn about the underlying structure of evaluation strategies
- learn how to improve our current programming languages (e.g. Rust)
- learn about them as models for complexity theory [Forster et al.(2019)Forster, Kunze, and Roth,]

Call by Value

- · evaluate the argument before the function
- the argument is evaluated at most once

Call by Value

- evaluate the argument before the function
- the argument is evaluated at most once

Call by Value via "suspending computation"

- · suspend computation of all the innermost functions calls
- · evaluate from last to first

```
let f x y = x + y
let () = print_int (f (1 + 2) (3 + 4))
```

Example

```
let f x y = x + y
let () = print_int (f (1 + 2) (3 + 4))
```

• suspend f(1 + 2)(3 + 4)

```
let f x y = x + y
let () = print_int (f (1 + 2) (3 + 4))
```

- suspend f (1 + 2) (3 + 4)
- suspend 1 + 2

```
let f x y = x + y
let () = print_int (f (1 + 2) (3 + 4))
```

- suspend f (1 + 2) (3 + 4)
- suspend 1 + 2
- suspend 3 + 4

```
let f x y = x + y
let () = print_int (f (1 + 2) (3 + 4))
```

- suspend f (1 + 2) (3 + 4)
- suspend 1 + 2
- suspend 3 + 4
- evaluate 3 + 4

```
let f x y = x + y
let () = print_int (f (1 + 2) (3 + 4))
```

- suspend f (1 + 2) (3 + 4)
- suspend 1 + 2
- suspend 3 + 4
- evaluate 3 + 4
- evaluate 1 + 2

```
let f x y = x + y
let () = print_int (f (1 + 2) (3 + 4))
```

- suspend f (1 + 2) (3 + 4)
- suspend 1 + 2
- suspend 3 + 4
- evaluate 3 + 4
- evaluate 1 + 2
- evaluate f 3 7



CONCLUSION

In short

- · Call by Name
- · Call by Value
- Call by Need

For the next time

- Model Call by Name and Call by Need in this model
- Show interpretations have nice properties (soundness, completeness, etc.)



REFERENCES I



Yannick Forster, Fabian Kunze, and Marc Roth.

The weak call-by-value lambda-calculus is reasonable for both time and space.

Proc. ACM Program. Lang., 4(POPL), dec 2019.

doi: 10.1145/3371095.

URL https://doi.org/10.1145/3371095.



John Maraist, Martin Odersky, David N. Turner, and Philip Wadler.

Call-by-name, Call-by-value, Call-by-need, and the Linear Lambda Calculus.

Electron. Notes Theor. Comput. Sci., 1:370–392, January 1995. ISSN 1571-0661.

doi: 10.1016/S1571-0661(04)00022-2.