

~~Recall~~

## :- ARRAY :-

It is the collection of similar data elements stored in contiguous memory.

Formula to calculate memory :-

$$\text{address}(a[i]) = b + i \times w$$

$w$  = word size

$b$  = base address

$i$  = index no.

Ex - int  $a[5]$ ;

5000	5002	5004	5006	5008
$a[0]$	$a[1]$	$a[2]$	$a[3]$	$a[4]$

$$\begin{aligned}\text{address of } a[3] &= b + i \times w \\ &= 5000 + 3 \times 2 \\ &= 5006\end{aligned}$$

operations on array :-

- 1) Traversing 2) Insertion 3) Deletion 4) Searching 5) Sorting

① Traversing :- This operation is used to visit all the elements in an array.

'C' code

```
void traverse(int a[], int n)
{
    int i;
    for (i = 0; i < n; i++)
        process(a);
}
```

Algorithm

- ```
void traverse(a[]),
1) [initialize counter]
   set  $i = 0$ 
2) Repeat step 3 & 4 till  $i < n$ 
3) Apply process to  $a$ 
4) set  $i = i + 1$ 
   end of for loop
5) Exit.
```

② Insertion :- The operation is used to insert an element into an array provide it is not full.

'C' code

```
void insert(int a[], int n, int pos, int item)
{
    int i;
    for (i = n - 1; i >= pos; i--)
        a[i + 1] = a[i];
    a[pos] = item;
    n = n + 1;
}
```



### Algorithm:

void insert(int a[], int n, int pos, int item)

- 1) [Initialize Counter] set  $i = 0$
- 2) Repeat step 3 & 4 till  $i \geq pos$
- 3) set  $a[i+1] = a[i]$
- 4) set  $i = i + 1$   
end of for loop
- 5) set  $a[pos] = item$
- 6) set  $n = n + 1$

3) Deletion: The operation is used to delete an element from an array.

C code:  
void delete(int a[], int n, int pos) {  
    int i;  
    for( $i = pos$ ;  $i < n-1$ ;  $i++$ )  
         $a[i] = a[i+1]$ ;  
     $n = n - 1$ ;  
}

### Algorithm

void delete(int a[], int n, int pos)

- 1) [Initialize counter] set  $i = pos$
- 2) Repeat step 3 & 4 till  $i < n-1$
- 3) set  $a[i] = a[i+1]$
- 4) set  $i = i + 1$
- 5) set  $n = n - 1$
- 6) Exit

### 4) Searching:

void search(int a[], int n, int item)

{  
    int i, c = 0;  
    for( $i = 0$ ;  $i < n$ ;  $i++$ )  
    {  
        if( $a[i] == item$ )  
             $c++$ ;  
    }

if( $c > 0$ )

printf("Element found")

else  
    printf("Element Not Found");

}

### Algorithm:

void search(int a[], int n, int item)

- 1) [Initialize counter]  $i = 0, c = 0$
- 2) Repeat step 3 to 5 till  $i < n$
- 3) if  $a[i] == item$
- 4) set  $c = c + 1$
- 5) set  $i = i + 1$
- 6) if  $c > 0$
- 7) write element found
- 8) else  
    write Not found
- 9) Exit.



### 5) Sorting :

```
void Sort (int a[], int n)
{
    int i, j, temp;
    for (i = 0; i < n-1; i++)
    {
        for (j = i+1; j < n; j++)
        {
            if (a[i] > a[j])
            {
                temp = a[i];
                a[i] = a[j];
                a[j] = temp;
            }
        }
    }
}
```

### Algorithm

- ```
void Sort (a[], n)
1) initialize i = 0
2) Repeat step 3 to 10 till i < n-1
3) Repeat set j = i+1
4) Repeat step 5 to 9 till j < n
5) if a[i] > a[j]
6) set temp = a[i]
7) set a[i] = temp
8) set a[j] = temp
9) set j = j+1
   end of for loop
10) set i = i+1
    end of for loop
11) Exit.
```

### Memory Representation of a Matrix :-

There are 2 convention of storing any matrix in memory.

#### (1) Row-major-order :-

In row major order elements of a matrix are stored on a row by row basis, that is all the elements of first row are filled, then in second row and so on.

$a[3][4] \rightarrow$

$a[0][0] \rightarrow a[0][1] \rightarrow a[0][2] \rightarrow a[0][3]$

$a[1][0] \rightarrow a[1][1] \rightarrow a[1][2] \rightarrow a[1][3]$

$a[2][0] \rightarrow a[2][1] \rightarrow a[2][2] \rightarrow a[2][3]$

$a[0][0]$	$a[0][1]$	$a[0][2]$	$a[0][3]$	$a[1][0]$	$a[1][1]$	$a[1][2]$	$a[1][3]$	$a[2][0]$	$a[2][1]$	$a[2][2]$	$a[2][3]$
-----------	-----------	-----------	-----------	-----------	-----------	-----------	-----------	-----------	-----------	-----------	-----------

if array of  $a[m][n]$  is there & to



Calculate the address of  $a[i][j]$  formula is

$$\text{add}(a[i][j]) = b + ((i - \text{LBR}) \times n + (j - \text{LBC})) \times w$$

$b$  = base address

$i$  = row no

$j$  = col no

$w$  = word size

$n$  = total no of columns

$\text{LBR}$  = Lower bound row  
(starting index of row)

$\text{LBC}$  = Lower bound col  
(starting index of column)

Q) Calculate the address of  $a[7][8]$  element of matrix  $a[20][25]$ ?

$$\text{add}(a[7][8]) = b + ((i - 0) \times n + (j - 0)) \times w$$

Let starting address is 1000, & word size is 2 byte

$$= 1000 + (7 \times 25 + 8) \times 2$$

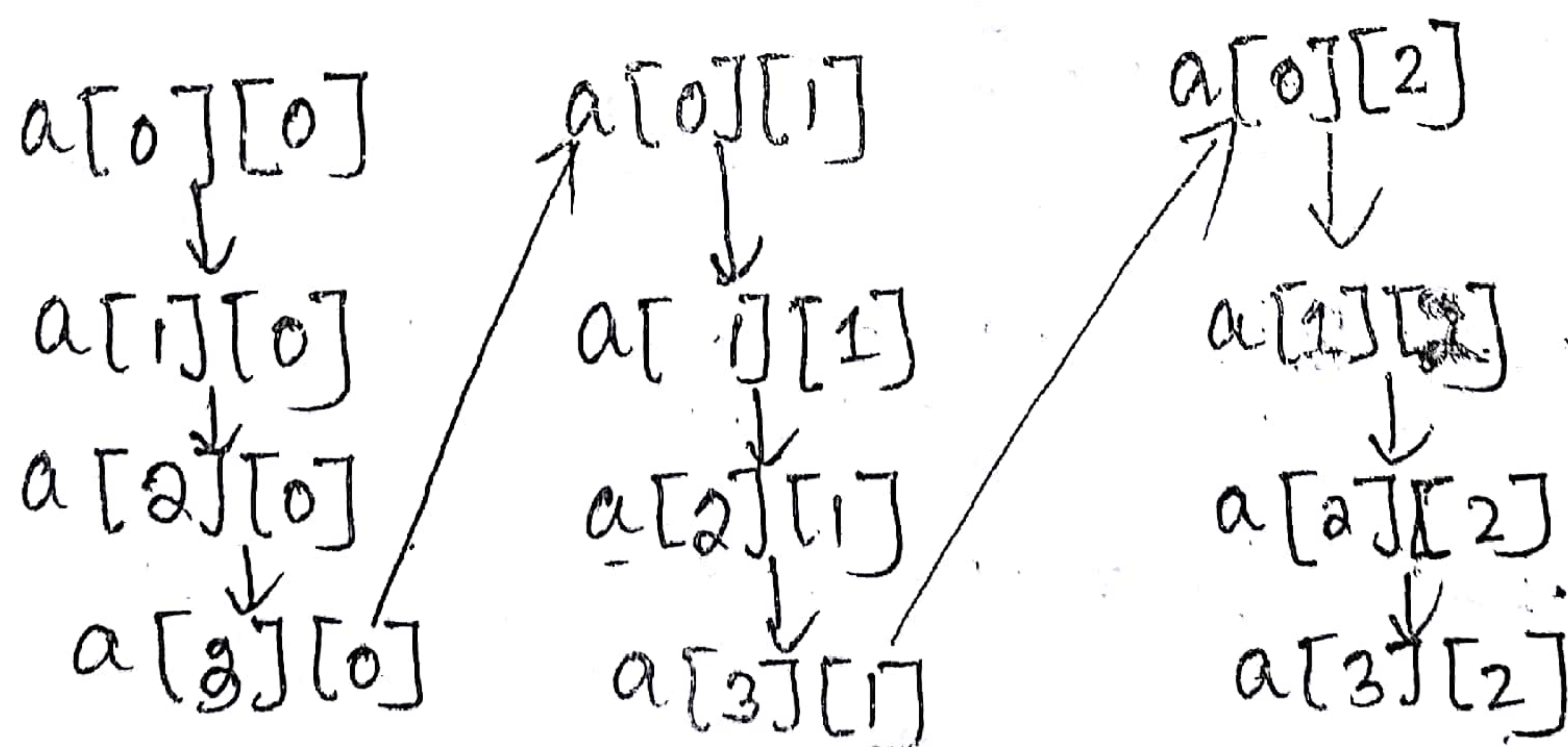
$$= 1000 + (175 + 8) \times 2$$

$$= 1366$$

(ii) column-major order:-

In this elements are stored column by column that is all the elements in first column are stored in their order of rows, then the 2nd column & so on

Ex:  $a[4][3]$



$a[0][0]$	$a[1][0]$	$a[2][0]$	$a[3][0]$	$a[0][1]$	$a[1][1]$	$a[2][1]$	$a[3][1]$	$a[0][2]$	$a[1][2]$	$a[2][2]$	$a[3][2]$
-----------	-----------	-----------	-----------	-----------	-----------	-----------	-----------	-----------	-----------	-----------	-----------

If array  $a[m][n]$  is present & address of  $a[i][j]$



element of array  $a[m][n]$  is

$$\text{add}(a[i][j]) = b + ((j - \text{LBC}) \times m + (i - \text{LBR})) \times w$$

here  $m$  is total no. of row.

Q) Suppose an array  $\text{VAL}[15][10]$  is stored in the memory with each element requiring 4 bytes of storage in C programming. If the base address of array  $\text{VAL}$  is 1500, determine the location of  $\text{VAL}[12][9]$  when the array  $\text{VAL}$  is stored (i) row-major order  
(ii) column-major order

Ans:

$m = 15$	$w = 4$	$\text{LBR} = 0$
$n = 10$	$i = 12$	$\text{LBC} = 0$
$b = 1500$	$j = 9$	

(i)  $\text{add}(a[i][j]) = b + ((i - \text{LBR}) \times n + j) \times w$   
 $= 1500 + (12 \times 10 + 9) \times 4$   
 $= 2016$

(ii)  $\text{add}(a[i][j]) = b + ((j - \text{LBC}) \times m + (i - \text{LBR})) \times w$   
 $= 1500 + (9 \times 15 + 12) \times 4$   
 $= 2088$

Sparse Matrix :-

- > A matrix  $A$  is said to be sparse if many of its elements are zero.
- > A matrix that is not sparse is called dense.
- > It is not possible to define an exact boundary between dense and sparse matrix.

Ex -

$$\begin{bmatrix} 7 & 0 & 0 & 1 & 0 & 2 \\ 0 & 1 & 9 & 0 & 0 & 0 \\ 0 & 0 & 0 & 7 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 \\ 8 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 3 & 0 & 0 & 0 \end{bmatrix}$$



# Array Representation of sparse Matrix (Triplet form),

- > In array rep<sup>n</sup> an array triplet of type  $\langle \text{row}, \text{col}, \text{element} \rangle$
- > It is used to store the elements which are nonZero where first field of the triplet is used ~~is used~~ to record row position, second to column position, and third one to record the nonZero element of sparse matrix

	Row	col	NonZero element
A[0]	6	6	8
[1]	0	0	7
[2]	0	3	1
[3]	0	5	2
[4]	1	1	1
[5]	1	2	9
[6]	2	3	7
[7]	4	0	8
[8]	5	2	3

Array matrix of sparse matrix would be

A[9][3]   
 1st col for row   
 2nd col for col   
 3rd col for nonZero values.

8 1   
 non zero value total row, col, nonZero values   
 is used for storing

Program:

WAP for sparse matrix-

```
#include <stdio.h>
#include <conio.h>
void main()
{
    int x[10][10], y[10][3];
    int i, j, col, row, k=1, l=0, m=0;
    clrscr();
    printf("Enter row & col size");
    scanf("%d %d", &row, &col);
    for(i=0; i<row; i++)
    {
        for(j=0; j<col; j++)
        {
            scanf("%d", &x[i][j]);
        }
    }
}
```



It is now the matrix  $u$  of size  $n \times n$ ;

```
for(i=0; i<row; i++)
```

```
{ for(j=0; j<col; j++)
```

```
{ printf("x.d", u[i][j]);
```

```
if (u[i][j] != 0)
```

```
col++;
```

```
else
```

```
row++;
```

```
}
```

```
} printf("n");
```

```
if (n < 3)
```

```
{ y[0][0] = row;
```

```
y[0][1] = col;
```

```
y[0][2] = n;
```

```
for(i=0; i<row; i++)
```

```
{ for(j=0; j<col; j++)
```

```
{ if (u[i][j] != 0)
```

```
{ y[k][0] = i;
```

```
y[k][1] = j;
```

```
y[k][2] = u[i][j];
```

```
k++;
```

```
}
```

```
}
```

```
}
```

```
printf("In the sparse matrix to n");
```

```
for(i=0; i<k; i++)
```

```
{ for(j=0; j<3; j++)
```

```
{ printf("x.d", y[i][j]);
```

```
} printf("n");
```

```
}
```



```

    }
else
{
    pf("Not a sparse matrix");
}

```

```

    getch();
}

```

WAP to delete all the occurrence of an element  $x_0$  in an array.

```

void main()
{
    int i, a[10], n, el, j;
    clrscr();
    pf("Enter the size\n");
    sf("%d", &n);
    pf("Enter elements\n");
    for(i=0; i<n; i++)
        sf("%d", &a[i]);
    pf("Enter the element to delete\n");
    sf("%d", &el);
    while(i<n)
    {
        if(a[i]==el)
        {
            for(j=i; j<n-1; j++)
                a[j]=a[j+1];
            i=0;
            n=n-1;
        }
        else
            i++;
    }
    for(i=0; i<n; i++)
        pf("%d", a[i]);
    getch();
}

```

OR void main()

```

{
    int a[10], p[10], i, n, el, c=0, k;
    clrscr();
    pf("Enter size\n");
    sf("%d", &n);
    pf("Enter element\n");
    for(i=0; i<n; i++)
        sf("%d", &a[i]);
    pf("Enter the element to be deleted\n");
    sf("%d", &el);
    for(i=0; i<n; i++)
    {
        if(a[i]==el)
        {
            p[k]=i;
            c++; k++;
        }
    }
    for(i=0; i<c; i++)
    {
        p[i]=p[i]-i;
        for(j=p[i]; j<n-1; j++)
            a[j]=a[j+1];
        n=n-1;
    }
    for(i=0; i<n; i++)
        pf("%d", a[i]);
    getch();
}

```