SM4 算法原理:

1. 常量定义:

SM4 FK, SM4 CK

S盒(8位输入,8位输出的非线性置换)

2. 加密流程:

明文分组(128位)

密钥扩展(生成32个轮密钥)

32 轮迭代变换

最终变换 (反序输出)

T-table 优化:

核心思路:预计算 S 盒与线性变换的组合结果,减少实时计算量

- 1. 预计算 T enc[i] = L(S(i) << 24) (加密用)
- 2. 预计算 T_key[i] = L'(S(i) << 24) (密钥扩展用)
- 3. 代码实现:

def _init_t_tables(self):

self.T_enc = [0] * 256 # 加密用 T 表

self.T_key = [0] * 256 # 密钥扩展用 T 表

for i in range (256):

 $s = SM4_SBOX[i]$

self.T enc[i] = self. 1 transform(s << 24)</pre>

self.T key[i] = self. 1 prime transform(s << 24)</pre>

```
def _round_function(self, x0: int, x1: int, x2: int, x3: int, rk: int) -> int:

tmp = x1 ^ x2 ^ x3 ^ rk

# 使用预计算 T 表加速

t = self. T_enc[tmp >> 24]

t ^= (self. T_enc[(tmp >> 16) & 0xFF] >> 8)

t ^= (self. T_enc[(tmp >> 8) & 0xFF] >> 16)

t ^= (self. T_enc[tmp & 0xFF] >> 24)

return x0 ^ t
```

AES-NI 优化:

AES-NI 是 Intel 处理器的加密加速指令集, 虽专为 AES 设计, 但可通过封装调用提升 SM4 性能。代码中通过检测硬件支持, 优先使用优化路径:

```
def _check_hardware_support(self):
    self.aesni_supported = False
    try:
```

from cryptography.hazmat.primitives.ciphers import

from cryptography.hazmat.backends import default_backend

```
backend = default_backend()

if hasattr(backend, 'has_aesni_support') and
backend.has_aesni_support():

self.aesni_supported = True

print("AES-NI 硬件加速支持已启用")

except ImportError:

pass # 降级为 T-table 实现
```

GFNI 优化:

GFNI 是新一代加密指令集,支持有限域运算和置换操作,可直接加速 S 盒和线性变换:

```
def _check_gfni_support(self):
    self.gfni_supported = False
    if os.name == 'posix':
        with open('/proc/cpuinfo', 'r') as f:
        cpuinfo = f.read()
        if 'gfni' in cpuinfo and 'vpbroadcastd' in
cpuinfo:
        self.gfni_supported = True
        # 检测到 GFNI 支持
```

GCM 优化:

1. 核心组件:

计数器模式 (CTR): 用于加密和解密

生成计数器块: CB = IV || 0x00000001

加密: $C_i = P_i \hat{E}(K, CB_i)$, $CB_{i+1} = CB_i + 1$

2. GHASH 函数:用于消息认证

基于 GF(2¹²⁸)乘法的哈希函数

输入: 关联数据、密文、长度信息

输出: 认证标签

实现代码:

def encrypt(self, key: bytes, iv: bytes, plaintext: bytes,
associated_data: bytes = b'') -> Tuple[bytes, bytes]:

生成哈希密钥 H = SM4(K, 0^128)

 $h = self. sm4. encrypt block (b' \x00' * 16, key)$

计数器模式加密

 $cb = iv + b' \times 00 \times 00 \times 00$

初始计数器块

ciphertext = []

for block in [plaintext[i:i+16] for i in range(0,
len(plaintext), 16)]:

ctr = self.sm4.encrypt_block(current_cb, key)

加密计数器

ciphertext_block = bytes([b ^ c for b, c in zip(block,

```
ctr[:len(block)])])
       ciphertext.append(ciphertext block)
       current cb = inc iv(current cb)
        # 计数器递增
   # 计算认证标签
   ghash_input = associated_data + b''.join(ciphertext)
   ghash_input += struct.pack(">QQ", len(associated_data)*8,
len(ciphertext)*8)
   s = self. ghash(h, ghash input)
         = bytes([b
                      c for b, c in
                                                   zip(s,
    tag
self. sm4. encrypt block (cb, key))])
   return b''. join(ciphertext), tag
   其中 GF (2~128) 乘法实现:
def _gf128_mul(a: int, b: int) -> int:
   p = 0x87
   # 不可约多项式: x^128 + x^7 + x^2 + x + 1
   result = 0
   for i in range (128):
       if b & 1:
           result ^= a
       a <<= 1
       if a & (1 << 128):
```

模不可约多项式

 $b \gg 1$

return result & ((1 << 128) - 1)

5. 性能测试:

标准测试向量

key = bytes.fromhex("0123456789abcdeffedcba9876543210")

plaintext =

 $\verb|bytes.from| hex("0123456789 abcdeffed cba9876543210")$

expected_ciphertext

bytes.fromhex("681edf34d206965e86b3e94f536e4246")

验证加密正确性

ciphertext = sm4_base.encrypt_block(plaintext, key)

assert ciphertext == expected_ciphertext, 基础实现加密错误