Subgraph Matching: on Compression and Computation

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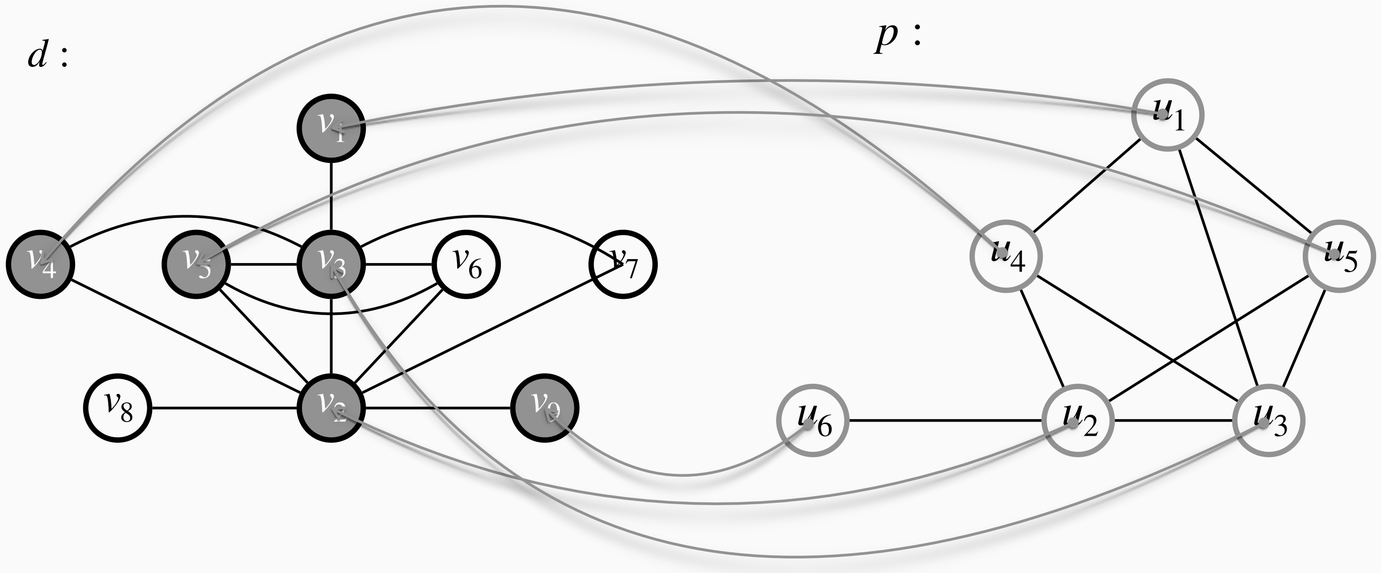
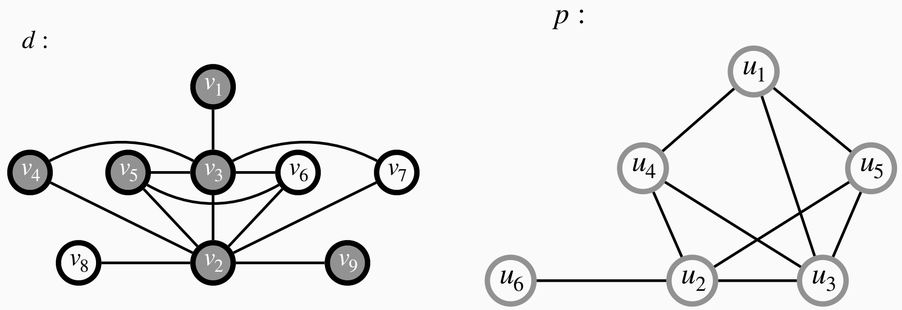
Subgraph Matching. Given a target graph *d* and a pattern graph *p*, report the set *I* of all the subgraphs in the target graph *d* that are isomorphic to the pattern graph *p*.



$$f : u_1 \mapsto v_1, u_2 \mapsto v_2, \cdots, u_5 \mapsto v_5, u_6 \mapsto v_9$$

Motivation

Output Crisis



dataset	size(d) in MB	P's Match Result
ego-Gplus(GP)	390	270.15TB
web-BerkStan(WB)	211	14.45TB
as-Skitter(AS)	355	8.46TB
soc-LiveJournal(LJ)	1373	9.14TB
uk-2002(UK)	9539	246.59TB

Current Method

- No output : repetitive computation
- Output the resulting set : expensive I/O cost

Motivation

Current Method

- No output : repetitive computation
- Output the resulting set : expensive I/O cost

Output solution: Is there an ideal **compression** on the instance set *I* ?

Lossless:

The compression should be lossless

d-optimal:

The decompression should be optimal, that is, the resulting set I of a subgraph matching restored in $O(\frac{|I|}{B})$ I/Os.

Motivation

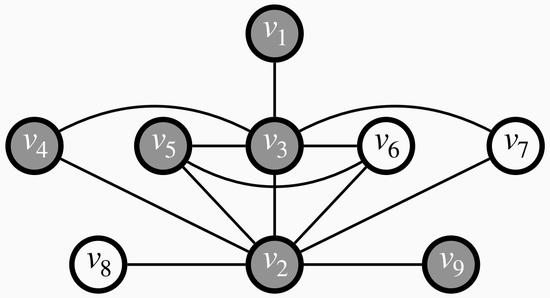
Possible Bijection between p and d given a image for u1, u2, u3(Uncompressed)

d\p	u_1	u_2	u_3	u_4	u_5	u_6
v_1						
v_2						
v_3						
v_4						
v_5						
v_6						
v_7						
v_8						
v_9						

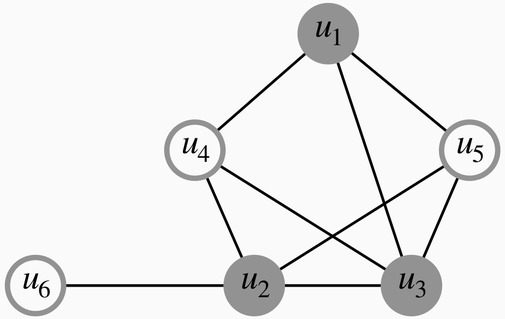
$$|u_4| \times |u_5| \times |u_6|$$

54 Possible Combination ! = **54*6*4 bytes**

d :



p :



Possible Bijection between p and d given bijection for u1, u2, u3(Compressed)

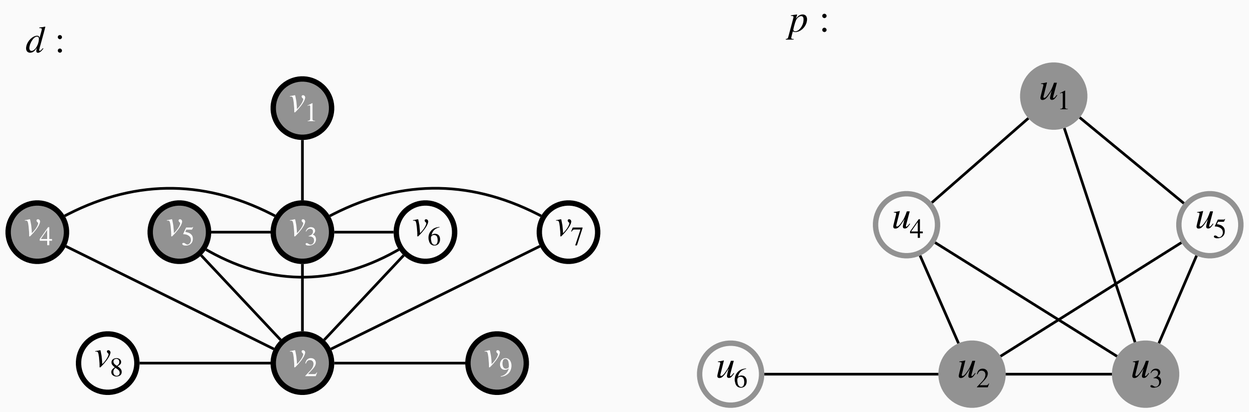
V(p)	u_1	u_2	u_3	u_4	u_5	u_6
Images	v_1	v_2	v_3	v_4,v_5,v_6	v_5,v_6,v_7	v_4,v_5, …,v_9

Only (3+(3+3+6))*4 bytes

Methods

Vertex-Cover Based Compression(VCBC)

- Identify a vertex-cover Vc of p
- Group the matched subgraphs by their images on the vertex cover nodes
- Compress : list the image sets of each node in V(p) for each group



Theorem(Lossless)

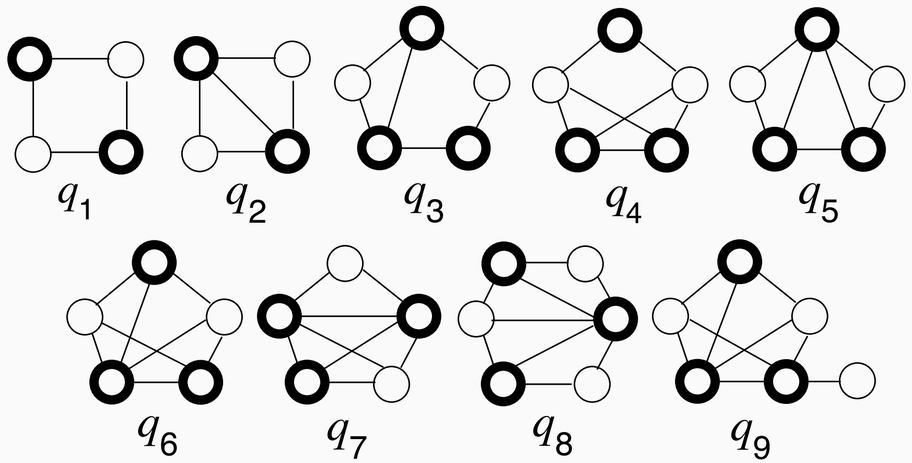
The vertex-cover based compression is lossless, that is, for a given V_c , $\bar{I} = I$.

Theorem(d-optimal)

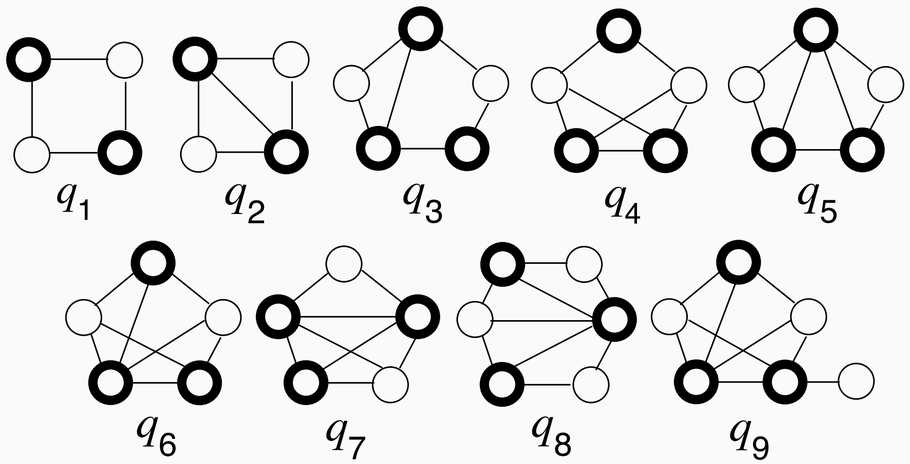
The vertex-cover based compression is decompression optimal, that is, the resulting set I of a subgraph matching can be restored in $O(\frac{|I|}{B})$ I/Os.

dataset	|V(d)| * 10^6	|E(d)| * 10^6	avg-deg	Degen-Eracy	size(d) in MB
ego-Gplus(GP)	0.1	12.2	244	1504	390
web-BerkStan(WB)	0.7	6.6	19	402	211
as-Skitter(AS)	1.7	11.1	13	222	355
soc-LiveJournal(LJ)	4.8	42.9	18	746	1373
uk-2002(UK)	18.5	298.1	32	1886	9539

Target Graph



Pattern Graph

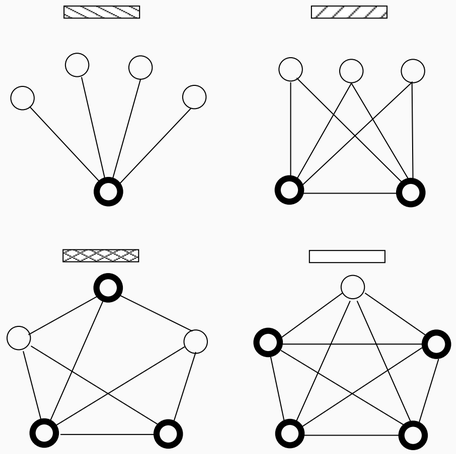


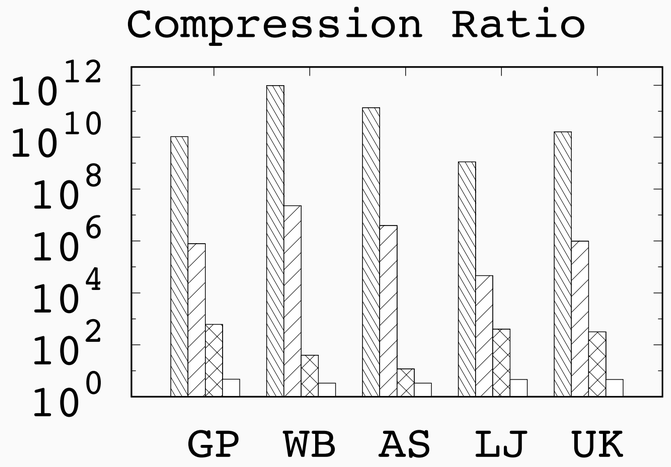
Pattern Graph

Compression Ratio

d\p	Q1	Q2	Q3	Q4	Q5	Q6	Q7	Q8	Q9
GP	333	1435	1263	409	1016	601	862	636433	23871
WB	17	2031	93	27	107	39	127	176833	93842
AS	23	790	80	9	76	12	66	39979	12724
LJ	19	342	581	201	362	400	440	147317	45336
UK	40	787	350	156	348	315	483	238077	130367

Compression Ratio - Freedom





Computation Solution: will the computation of subgraph matching be optimized based on the compression technique?

Avoid generating the instance set

Directly Generate Compressed Instance Set

Directly Generate Compressed Instance Set

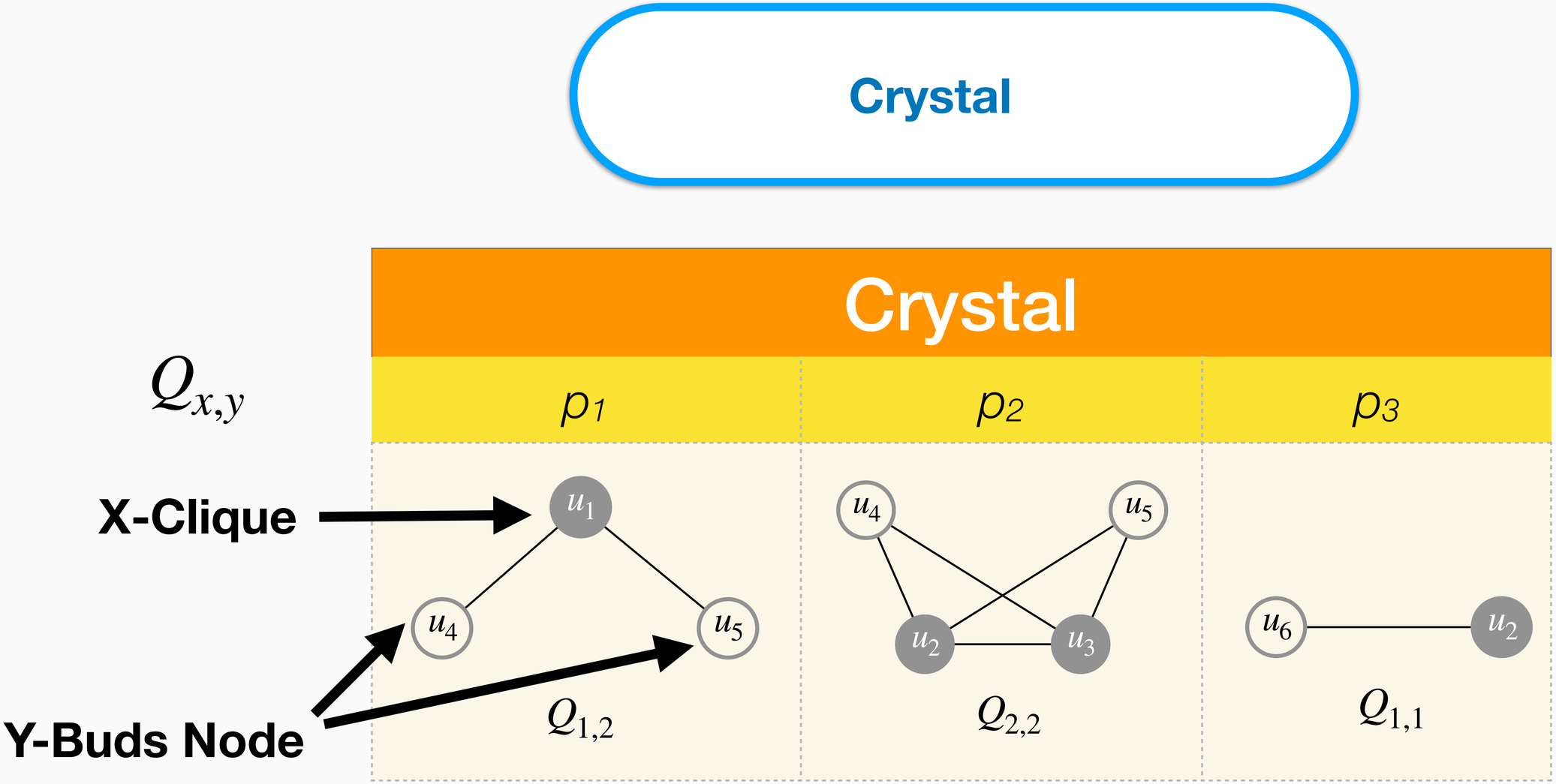
Crystal

Core-Crystal Decomposition

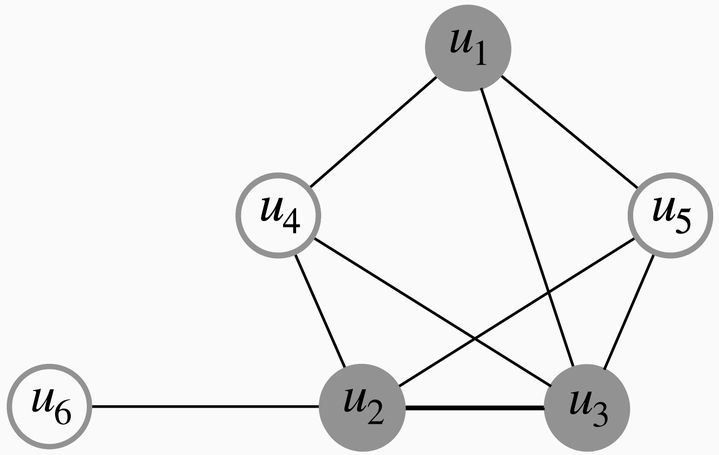
One-off Assembly

Methods

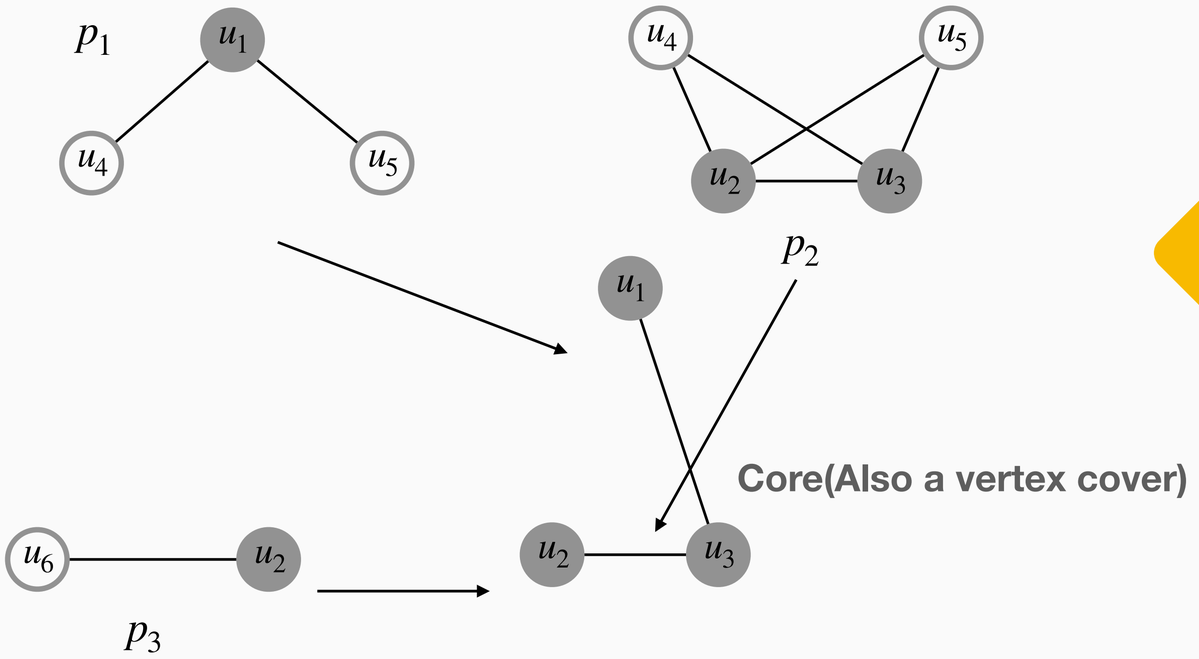
Crystal Based Computation (CBF)

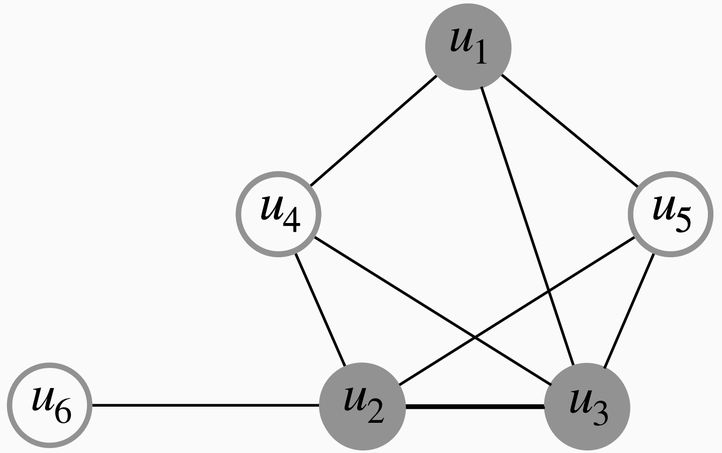


Core-Crystal Decomposition



Core-Crystal Decomposition

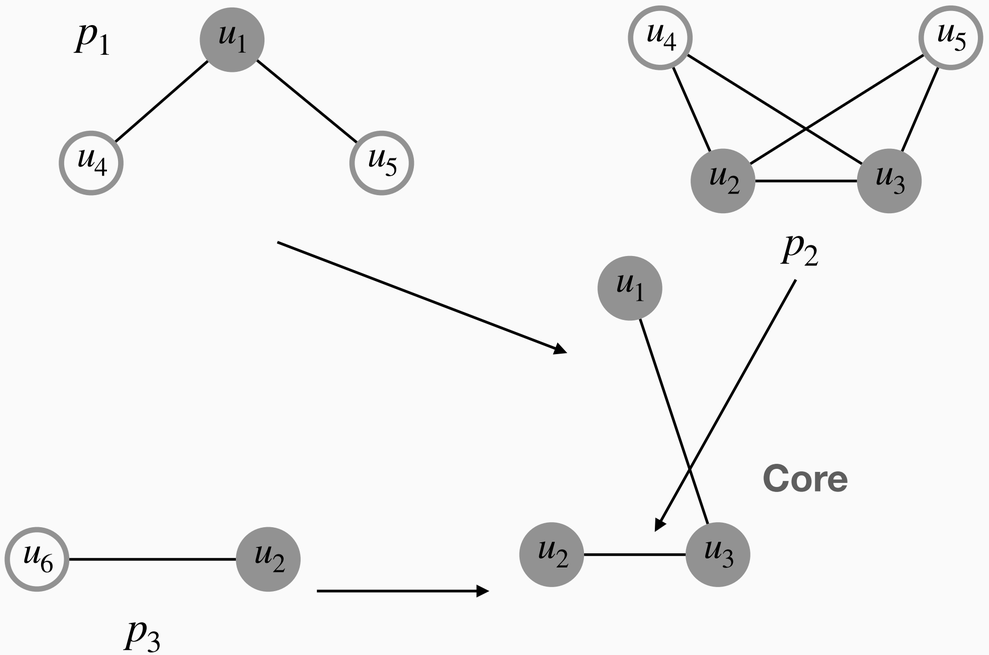




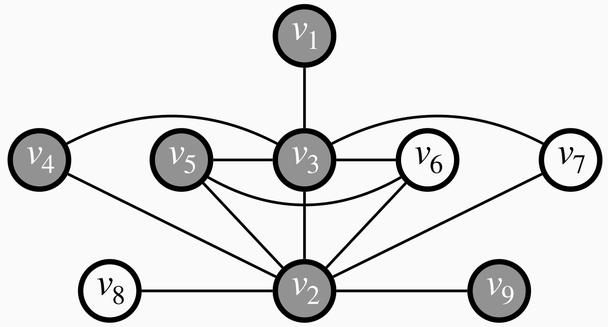
Methods

Crystal Based Computation (CBF)

One-off Assembly



d :



	u_1	u_2	u_3	u_4	u_5	u_6
Core	v_1	v_2	v_3			
p_1	v_1			$\{v_3,v_4,v_5,v_6,v_7\}$		
p_2		v_2	v_3	$\{v_4,v_5,v_6,v_7\}$		
p_3		v_2				$\{v_3,..,v_9\}$
P						

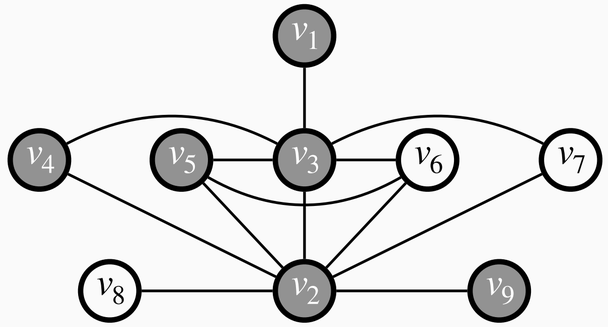
Mapping for core {v1,v2,v3} and crystal

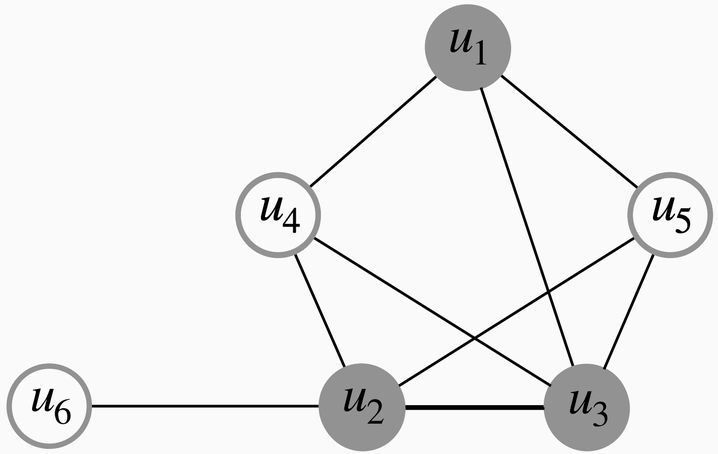
Methods

Crystal Based Computation (CBF)

One-off Assembly

d :



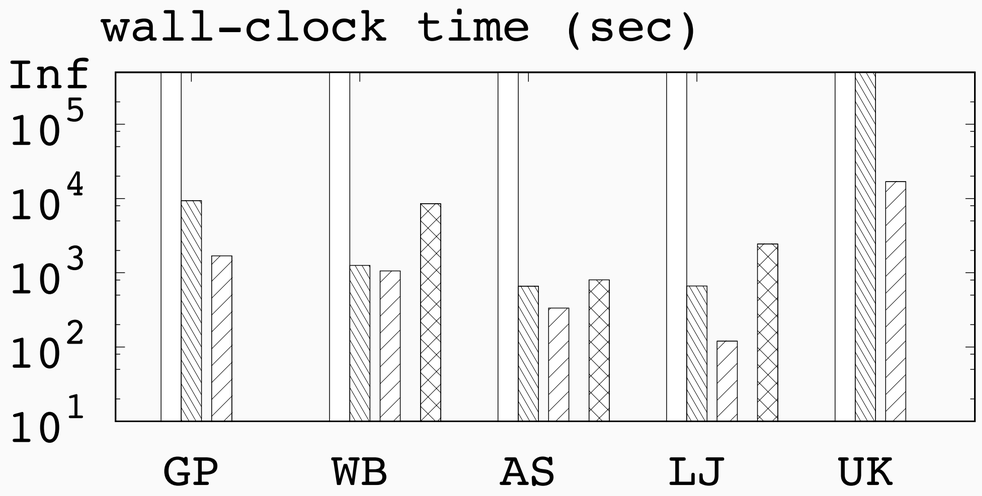


	u_1	u_2	u_3	u_4	u_5	u_6
Core	v_1	v_2	v_3			
p_1	v_1			$\{v_3,v_4,v_5,v_6,v_7\}$		
p_2		v_2	v_3	$\{v_4,v_5,v_6,v_7\}$		
p_3		v_2				$\{v_3,..,v_9\}$
P	v_1	v_2	v_3	$\{v_4,v_5,v_6,v_7\}$		$\{v_3,..,v_9\}$

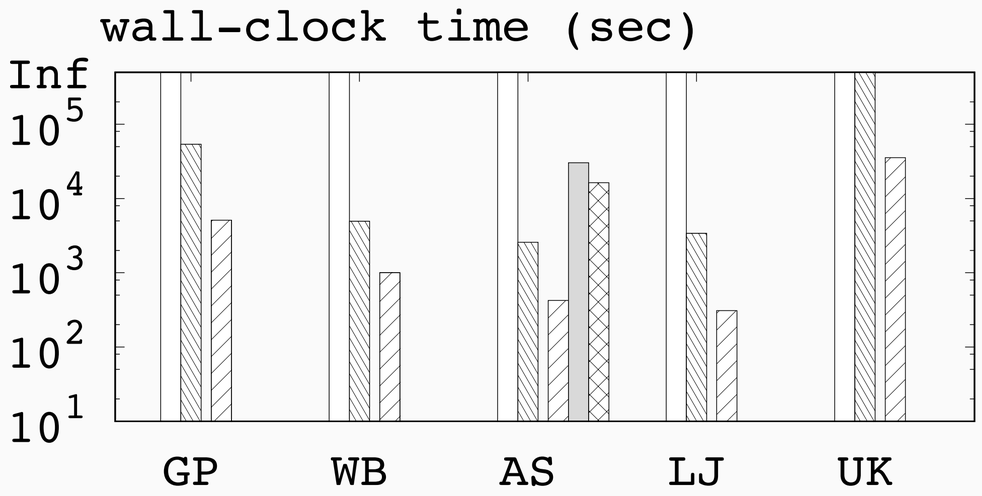
Mapping for core {v1,v2,v3} and crystal

- **Crystal(Distributed) :** our methods in distributed settings
- **Crystal-1(single machine) :** our methods using one thread
- **DualSim(single-machine) :** the state-of-the-art DFS-style solution
- **TwigTwin(Distributed) :** the state-of-the-art DFS-style solution
- **SEED(Distributed) :** the state-of-the-art DFS-style solution


DualSim
Crystal-1
Crystal
TwigTwin
SEED


wall-clock time (sec)
Inf
10^5
10^4
10^3
10^2
10^1
GP
WB
AS
LJ
UK

Q8


wall-clock time (sec)
Inf
10^5
10^4
10^3
10^2
10^1
GP
WB
AS
LJ
UK

Q9

Summary

- **Storage Solution For Subgraph Matching:**
  Vertex-Cover Based Compression(VCBC)

- **Direct Compressed Code Computation:**
  Crystal-Based Framework(CBF)

Thank You !

Questions ?