

# Info of Exam

- Closed-book examination: 2025-01-08 08:00-09:50
- Q & A (Tencent Meeting): 2024-12-31, 08:30-10:00
- Problem types:
  - Fill-in-the-blank
  - True / False
  - Short questions
  - Long questions
- Grade system:
  - Final exam + Homework + Attendance

# Outline

- PPT courseware + Homework
- Notice: specific tools (e.g., Lex or Yacc) or sample languages (e.g., TINY or C-minus), will not be examined.

# Outline-1

- The phases of a compiler
- Write regular expressions
  - Pay attention to basic and extended operations
- Convert the regular expression into an NFA, and then convert the NFA into a DFA
  - Subset construction method
- Minimize DFA

# Outline-2

- Derivation and parse tree
- Write a Context-Free Grammar
- Chomsky hierarchy
- Construct a LL(1) parsing table
  - Table construction:  $M[N, T]$
  - $e \in \text{First}(\alpha)$ : fill  $A \rightarrow \alpha$  in  $M[A, e]$
  - $e \in \text{Follow}(A)$ : fill  $A \rightarrow \varepsilon$  in  $M[A, e]$
- Parse an input string corresponding to the LL(1) parsing table

# Outline-3

- Construct a LR(0)/SLR(1)/LR(1)/LALR(1) parsing table
  - Table construction: state (list) + input (subtable) + goto (subtable)
  - SLR(1):
    - Follow(A): reduce with  $A \rightarrow \alpha$
  - LR(1)
    - $b \in \text{First}(\gamma a)$ :  $\epsilon$ -transition from item  $[A \rightarrow \alpha \cdot B \gamma, a]$  to item  $[B \rightarrow \cdot \beta, b]$
- A grammar is LL or LR grammar?
  - No shift-reduce or reduce-reduce conflicts
- Parse an input string corresponding to the LR parsing table

# Exercises

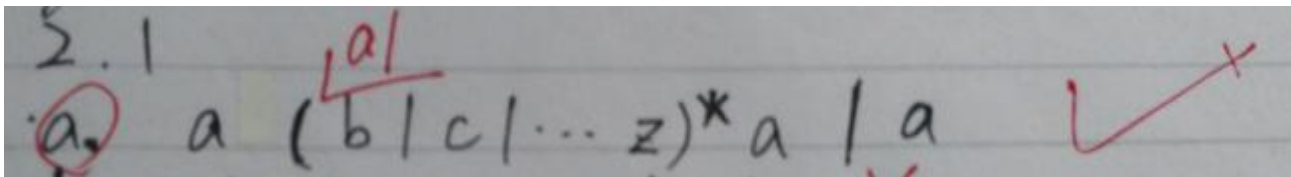
## 2.1

A. All strings of lowercase letters that begin and end in *a*.

Positive examples ( $x^+$ ): *a*, *aa*, *aba*,  
*abaabca*, ...

Negative examples ( $x^-$ ):  $\varepsilon$ , *ab*, *ba*, *abb*,...

Perfect answer:  $a[a-z]^*a \mid a$



Handwritten solution for problem 2.1. The text reads: 2.1 a  $a (b | c | \dots | z)^* a \mid a$ . There is a red checkmark at the end of the expression.

## 2.1


B. All strings of lowercase letters that either begin or end in *a* (or both).

$x^+$ : *a, aa, ab, ba, aba, aabcaca, ...*

$x^-$ :  $\varepsilon, b, bab, ...$

Perfect answer:  $a[a-z]^* \mid [a-z]^*a$

Wrong answer:

b.  $(a| \varepsilon)^* (a|b| \dots |z)^* (a| \varepsilon)^*$  

b.  $a(a|b| \dots |z)^* \cup (a|b| \dots |z)^*a$



## 2.1

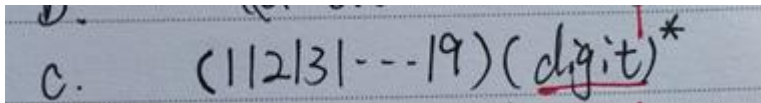
C. All strings of digits that contain no leading zeros.

$x^+$ : 0, 1, 2, 12, 123...

$x^-$ : 01, 002, 023...

Perfect answer:  $0|[1-9][0-9]^*$

Wrong answer:



D.  $(1|2|3|...|9)(\text{digit})^*$

C.  $(1|2|3|...|9)(\text{digit})^*$

## 2.1

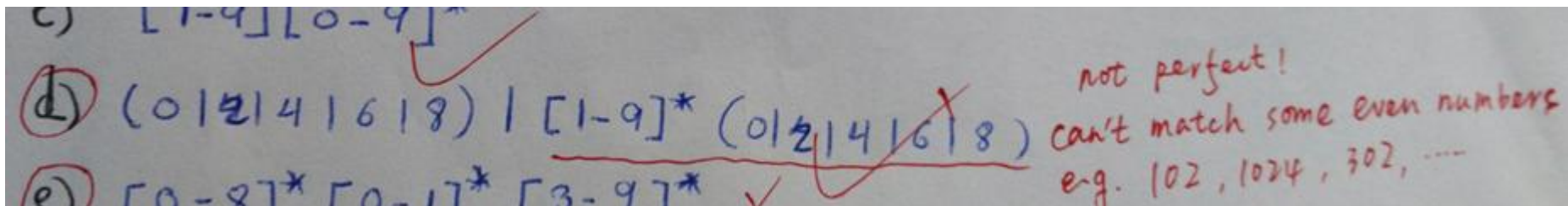
D. All strings of digits that represent even numbers.

$x^+$ : 2, 4, 12, 112...

$x^-$ : 1, 3, 13...

Perfect answer:  $[0|2|4|6|8][1-9][0-9]^*[0|2|4|6|8]$

Wrong answer:



# 2.1

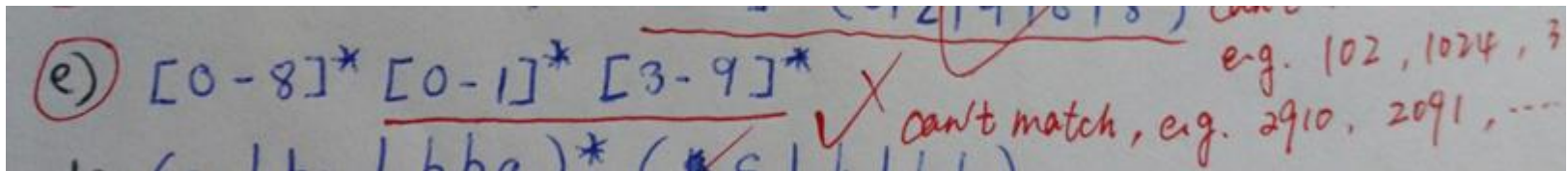
E. All strings of digits such that all the 2's occur before all the 9's.

$x^+$ : 2, 9, 29, 229, 12526996...

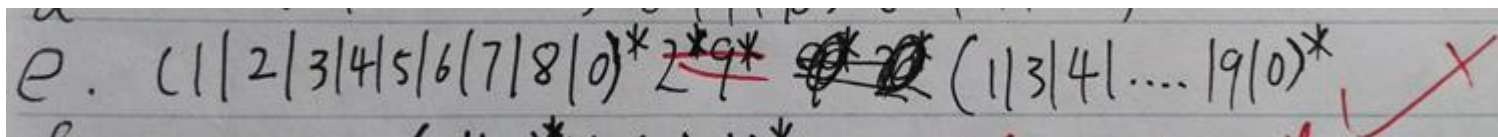
$x^-$ : 92, 992, 1913222...

Perfect answer:  $[0-8]^*[0|1|3|4|5|6|7|8|9]^*$

Wrong answer:



(e)  $[0-8]^* [0-1]^* [3-9]^*$   
can't match, e.g. 2910, 2091, ...  
e.g. 102, 1024, 3



e.  $(1|2|3|4|5|6|7|8|0)^* 2^* 9^* (1|3|4|...|9|0)^*$

# 2.1

F. All strings of  $a$ 's and  $b$ 's that contain no three consecutive  $b$ 's.

$x^+$ :  $\varepsilon, b, bb, abb, bba, abbaaaba, bab, bbabb...$

$x^-$ :  $bbb, bbbb, abbb, ababaabbba...$

Perfect answer:  $(\varepsilon|b|bb)(a|ab|abb)^*$  or  $(a|ba|bba)^*(\varepsilon|b|bb)$

f.  $(a|ba)^*bb(a|ab)^*$  X can't produce, e.g. bbabb

g.  $a^*((b|bb)a)^*$  X

e.  $((abb|bab|bba)a)^* \cup ((aab|baa)a)^*$   
can't produce e.g. b, bb, ... X  
 ~~$(aa)^*a((a|b|a|e) \cup (b|e)(bb|a)^*(b|e))^*$~~  X

## 2.1

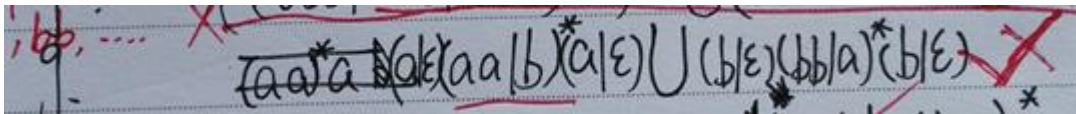
G. All strings of  $a$ 's and  $b$ 's that contain an odd number of  $a$ 's or an odd number of  $b$ 's (or both).

$x^+$ :  $a, b, aaa, bbb, ab, aaabb, bababaaba...$

$x^-$ :  $\varepsilon, aa, bb, abab, aaaabb...$

Perfect answer:  $b^*a(b^*ab^*a)^*b^*|a^*b(a^*ba^*b)^*a^*$

Wrong answer:



# 2.1

H. All strings of  $a$ 's and  $b$ 's that contain an even number of  $a$ 's and an even number of  $b$ 's

$x^+$ :  $\varepsilon, aa, bb, abab, aaaabb, bababaabaa...$

$x^-$ :  $a, b, aaa, bbb, ab, aaabb, bababaa...$

Perfect answer:  $((aa|bb)|(ab|ba)(aa|bb)^*(ab|ba))^*$

Wrong answer:

~~$(b^*|a^*b^*a^*)^*|b^*|a^*(a^*b^*a^*b^*)^*a^*$~~  even a's and even b X

~~$(b|ab^*a)^*|(a|ba^*b)^*$~~

~~$(aabb|abab|bbba|baba|baab|abba)^*$~~

## 2.1

- I. All strings of  $a$ 's and  $b$ 's that contain exactly as many  $a$ 's as  $b$ 's.

$x^+$ :  $\varepsilon, abab, bababaabab...$

$x^-$ :  $a, b, aaa, bbb, aaaabb...$

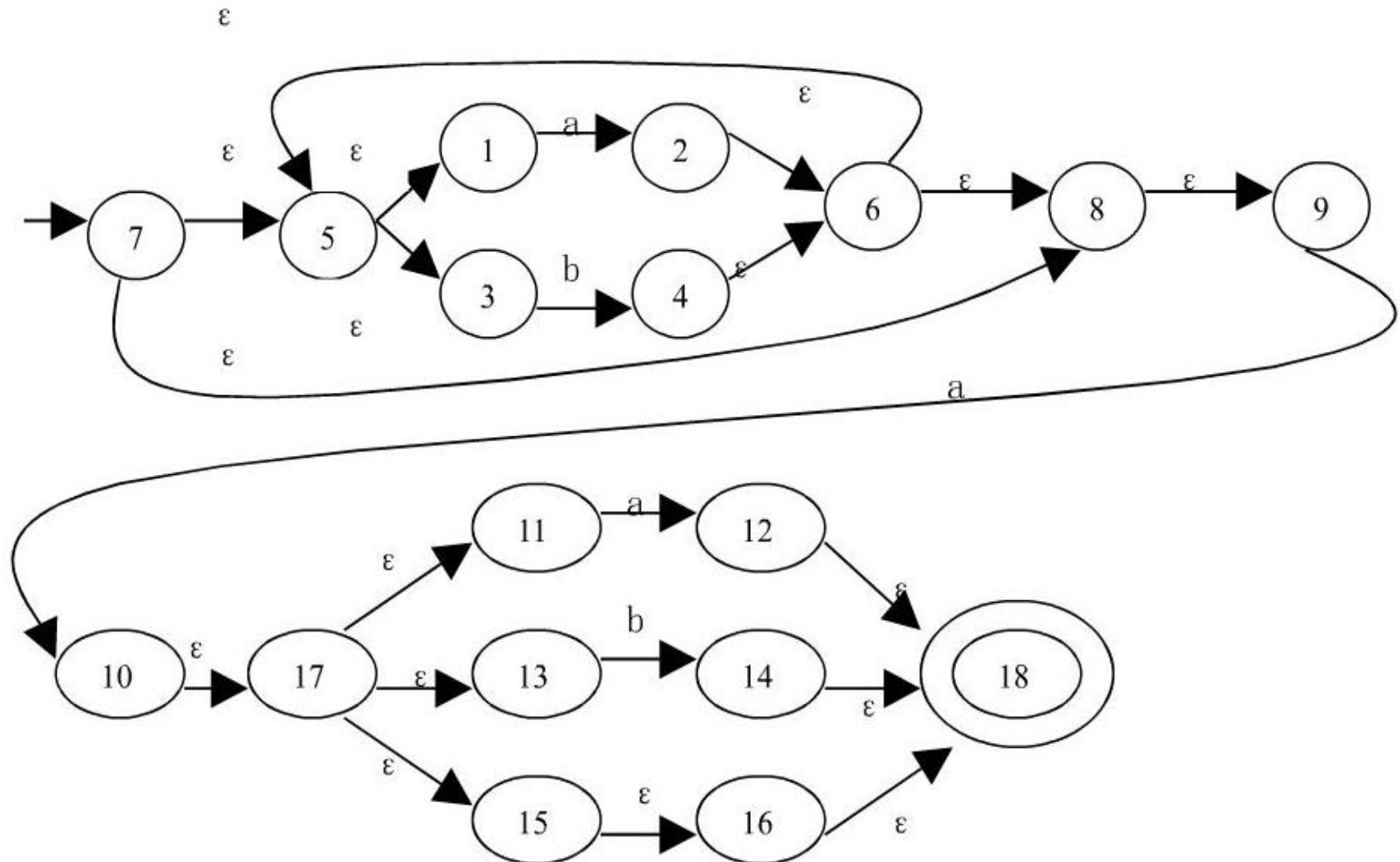
None. Because a regular expression cannot count.

2.12 a. Use Thompson's construction to convert the regular expression  $(a|b)^*a(a|b|\epsilon)$  into an NFA.

b. Convert the NFA of part (a) into a DFA using the subset construction.

[Solution]

a. An NFA of the regular expression  $(a|b)^*a(a|b|\epsilon)$





b. The subsets constructed as follows:

$$\underline{\{7\}} = \{7, 5, 1, 3, 8, 9\}$$

$$\underline{\{7\}}_a = \{2, 10\}$$

$$\underline{\{7\}}_b = \{4\}$$

$$\underline{\{2, 10\}} = \{2, 6, 5, 1, 3, 8, 9, 10, 17, 11, 13, 15, 16, 18\}$$

$$\underline{\{2, 10\}}_a = \{2, 10, 12\}$$

$$\underline{\{2, 10\}}_b = \{4, 14\}$$

$$\underline{\{2, 10, 12\}} = \{2, 6, 5, 1, 3, 8, 9, 12, 18, 10, 17, 11, 13, 15, 16\}$$

$$\underline{\{2, 10, 12\}}_a = \{2, 10, 12\}$$

$$\underline{\{2, 10, 12\}}_b = \{4, 14\}$$

$$\underline{\{4, 14\}} = \{4, 6, 5, 1, 3, 8, 9, 14, 18\}$$

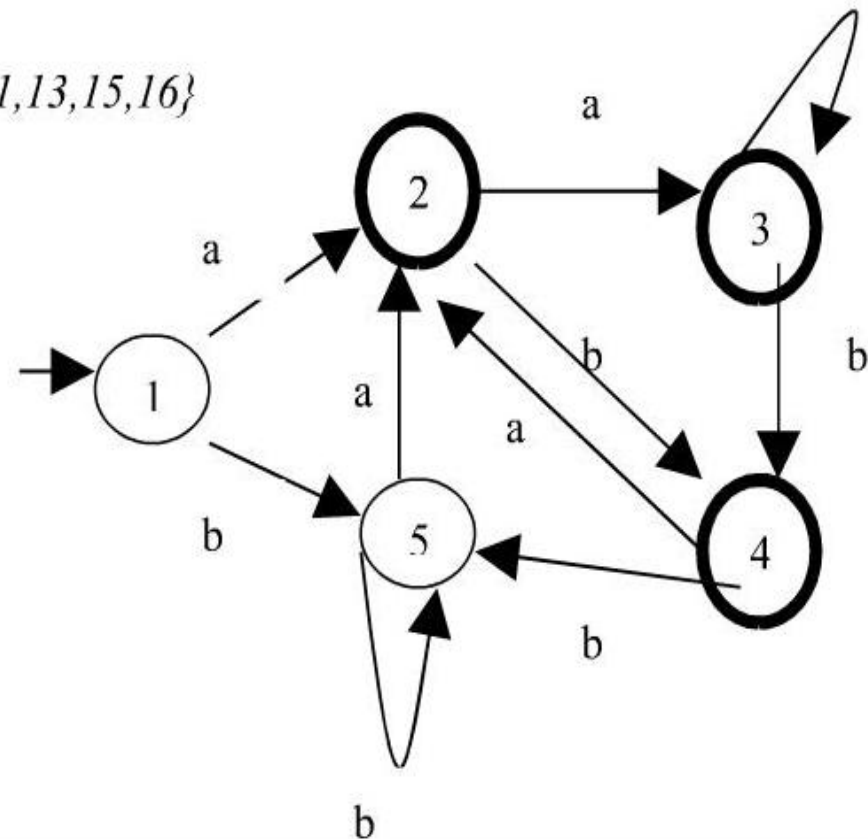
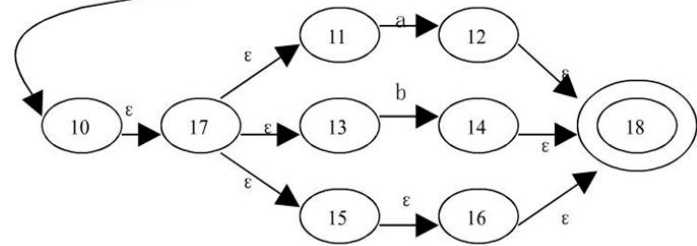
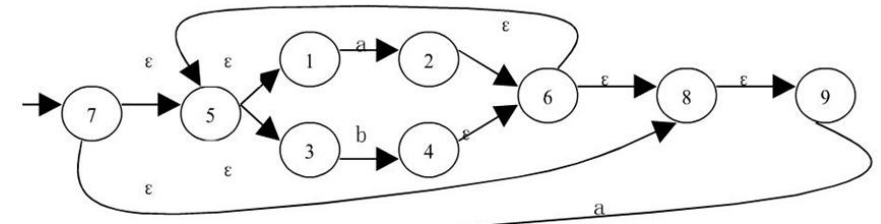
$$\underline{\{4, 14\}}_a = \{2, 10\}$$

$$\underline{\{4, 14\}}_b = \{4\}$$

$$\underline{\{4\}} = \{4, 6, 5, 1, 3, 8, 9\}$$

$$\underline{\{4\}}_a = \{2, 10\}$$

$$\underline{\{4\}}_b = \{4\}$$



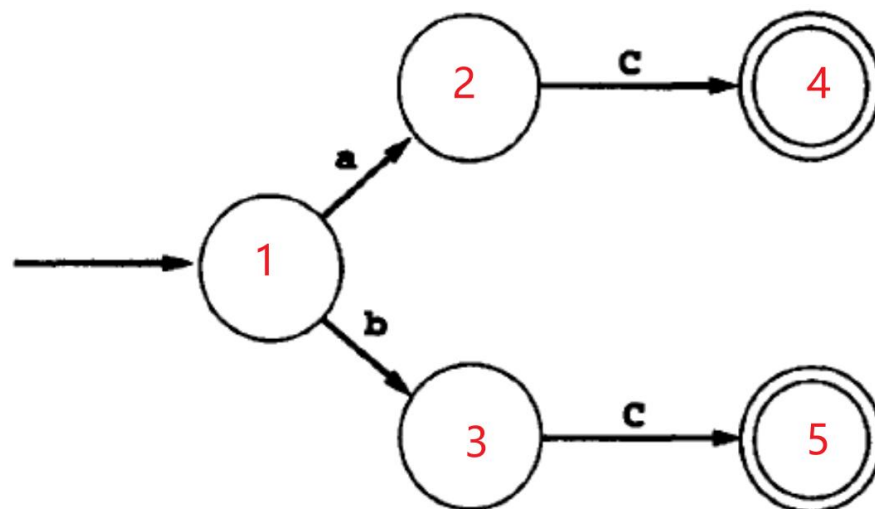
(b) From (a), we can get that:

| $S$       |  | $S_a$     | $S_b$ |
|-----------|--|-----------|-------|
| 7         | 7, 5, 1, 3, 8, 9   | 2, 10     | 4     |
| 2, 10     | 2, 6, 5, 1, 3, 8, 9, 10, 17, 11,<br>13, 15, 16, 18                     | 2, 10, 12 | 4, 14 |
| 2, 10, 12 | 2, 6, 5, 1, 3, 8, 9, <del>10</del> , 10, 17, 11,<br>13, 15, 16, 18, 12 | 2, 10, 12 | 4, 14 |
| 4, 14     | 4, 6, 5, 1, 3, 8, 9, 14, 18  | 2, 10     | 4     |
| 4         | 4, 6, 5, 1, 3, 8, 9  | 2, 10     | 4     |

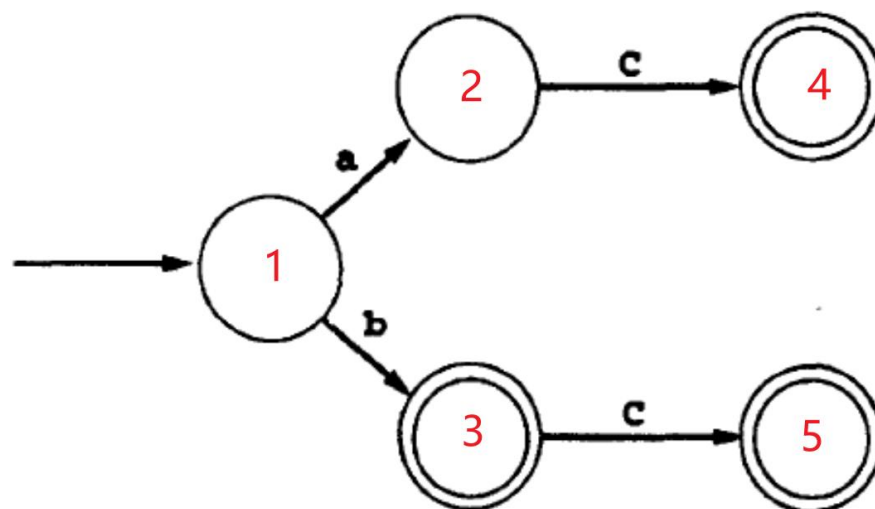
Therefore the answer is:

**2.16** Apply the state minimization algorithm of Section 2.4.4 to the following DFAs:

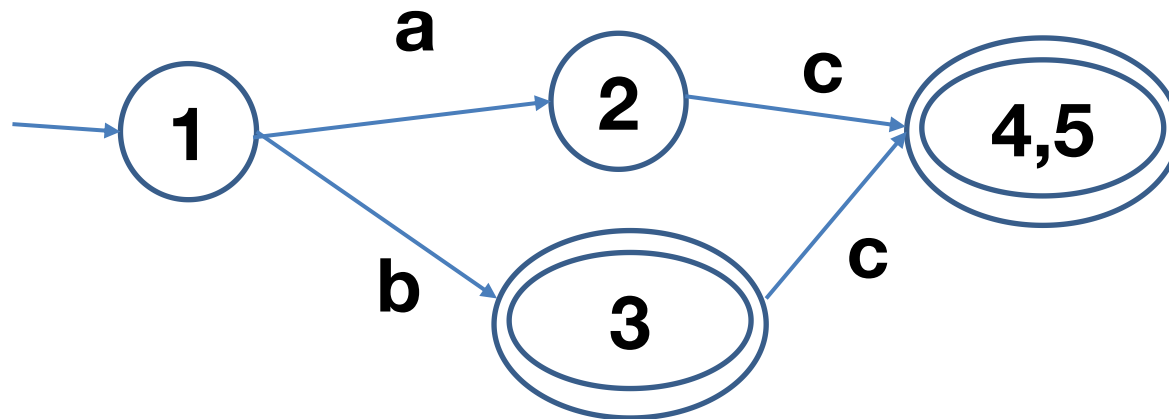
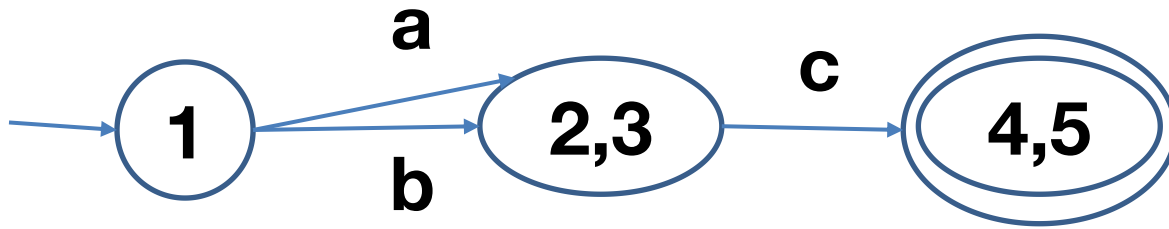
**a.**



**b.**



2.16



# 3.20

- 3.20 a.** Write a regular expression that generates the same language as the following grammar:

$$\begin{aligned} A &\rightarrow aA \mid B \mid \varepsilon \\ B &\rightarrow bB \mid A \end{aligned}$$

- b.** Write a grammar that generates the same language as the following regular expression:

$$(a \mid c \mid ba \mid bc)^* (b \mid \varepsilon)$$

1.  $(a|b)^*$

2.  $A \rightarrow Bb \mid B$

$$B \rightarrow aB \mid cB \mid baB \mid bcB \mid \varepsilon$$

Or  $A \rightarrow aA \mid cA \mid baA \mid bcA \mid B$

$$B \rightarrow b \mid \varepsilon$$

**4.9** Consider the following grammar (similar, but not identical to the grammar of Exercise 4.8):

$$\begin{aligned}lexp &\rightarrow atom \mid list \\atom &\rightarrow \textbf{number} \mid \textbf{identifier} \\list &\rightarrow ( lexp\text{-}seq ) \\lexp\text{-}seq &\rightarrow lexp , lexp\text{-}seq \mid lexp\end{aligned}$$


Start  
symbol

- Left factor this grammar.
- Construct First and Follow sets for the nonterminals of the resulting grammar.
- Show that the resulting grammar is LL(1).
- Construct the LL(1) parsing table for the resulting grammar.

# 4.9

- $\text{lexp} \rightarrow \text{atom} | \text{list}$   
 $\text{atom} \rightarrow \text{number} | \text{identifier}$   
 $\text{list} \rightarrow (\text{lexp-seq})$   
 $\text{lexp-seq} \rightarrow \text{lexp lexp-seq}'$   
 $\text{lexp-seq}' \rightarrow , \text{lexp-seq} | \varepsilon$

$\$ \rightarrow \text{Follow}(\text{lexp-seq})$

$\text{Follow}(\text{lexp}) \rightarrow \text{follow}(\text{atom})$  or  $\text{follow}(\text{list})$

$) \rightarrow \text{follow}(\text{lexp-seq})$

$\text{First}(\text{lexp-seq}') \text{ except for } \varepsilon$

$\rightarrow \text{follow}(\text{lexp})$

$\text{Follow}(\text{lexp-seq}) \rightarrow \text{follow}(\text{lexp})$

$\text{Follow}(\text{lexp-seq}) = \text{Follow}(\text{lexp-seq}')$

| Nonterminal | First set             | Follow set |
|-------------|-----------------------|------------|
| lexp        | number, identifier, ( | \$, ), ,   |
| atom        | number, identifier    | \$, ), ,   |
| list        | (                     | \$, ), ,   |
| lexp-seq    | number, identifier, ( | \$, )      |
| lexp-seq'   | ,, $\varepsilon$      | \$, )      |



| Nonterminal | First set             | Follow set |
|-------------|-----------------------|------------|
| lexp        | number, identifier, ( | \$, ), ,   |
| atom        | number, identifier    | \$, ), ,   |
| list        | (                     | \$, ), ,   |
| lexp-seq    | number, identifier, ( | \$, )      |
| lexp-seq'   | ,, $\epsilon$         | \$, )      |

| M[N,T]    | number                                   | identifier                               | (  | )                                | ,                                  | \$                               |
|-----------|--|--|--|----------------------------------|------------------------------------|----------------------------------|
| lexp      | lexp $\rightarrow$<br>atom               | lexp $\rightarrow$<br>atom               | lexp $\rightarrow$<br>list               |                                  |                                    |                                  |
| atom      | atom $\rightarrow$<br>number             | atom $\rightarrow$<br>identifier         |  |                                  |                                    |                                  |
| list      |  |  | list $\rightarrow$<br>(lexp-seq)         |                                  |                                    |                                  |
| lexp-seq  | lexp-seq $\rightarrow$ lexp<br>lexp-seq' | lexp-seq $\rightarrow$ lexp<br>lexp-seq' | lexp-seq $\rightarrow$ lexp<br>lexp-seq' |                                  |                                    |                                  |
| lexp-seq' |  |  |  | lexp-seq' $\rightarrow \epsilon$ | lexp-seq' $\rightarrow$ , lexp-seq | lexp-seq' $\rightarrow \epsilon$ |



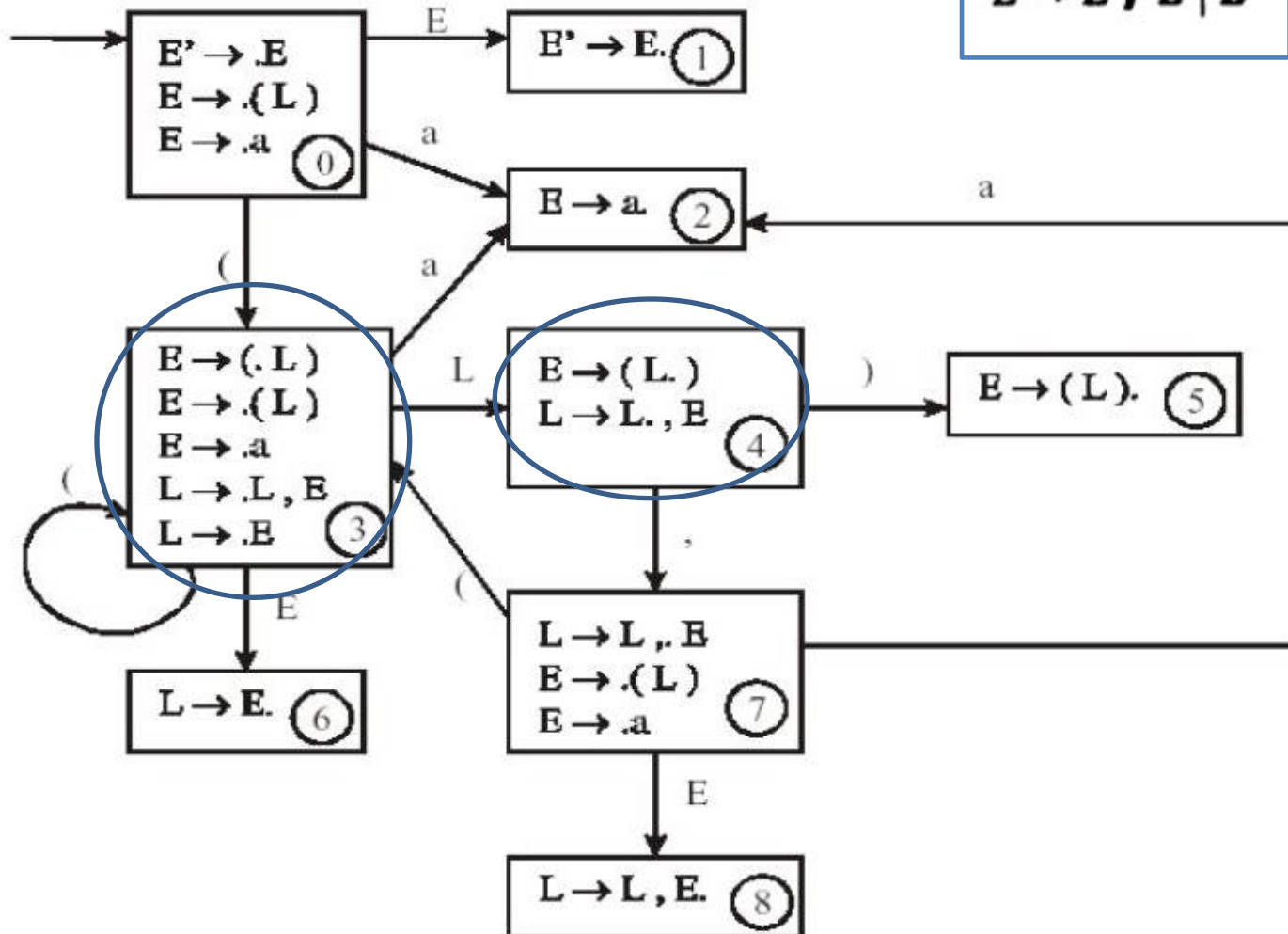
**5.1** Consider the following grammar:

$$\begin{aligned} E &\rightarrow ( L ) \mid a \\ L &\rightarrow L , E \mid E \end{aligned}$$

- a.** Construct the DFA of LR(0) items for this grammar.
- b.** Construct the SLR(1) parsing table.
- c.** Show the parsing stack and the actions of an SLR(1) parser for the input string  $((a), a, (a, a))$ .
- d.** Is this grammar an LR(0) grammar? If not, describe the LR(0) conflict. If so, construct the LR(0) parsing table, and describe how a parse might differ from an SLR(1) parse.

# 5.1

## a. DFA of LR(0) items



# 5.1

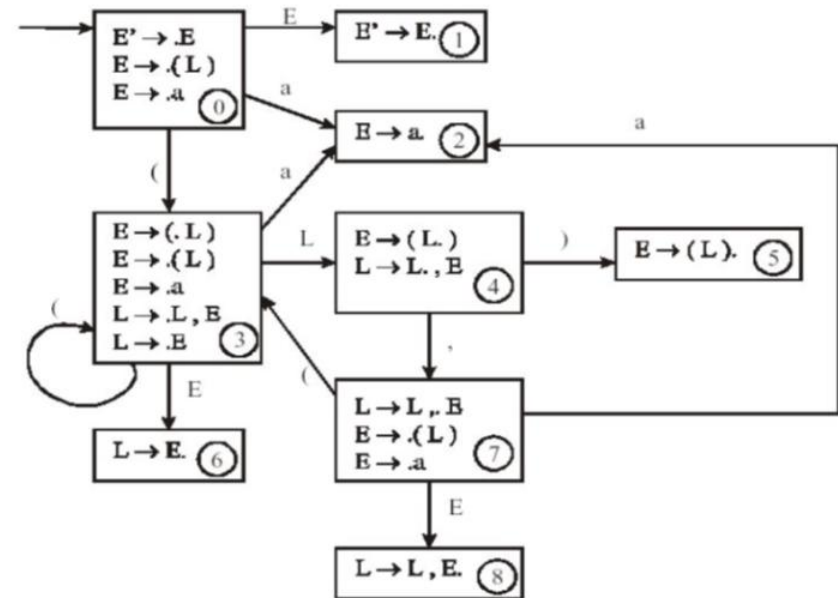
b. SLR(1) parsing table

$\text{Follow}(E') = \{\$ \}$

$\text{Follow}(E) = \{ \}, \text{ , } \$ \}$

$\text{Follow}(L) = \{ \}, \text{ , } \}$

$E \rightarrow (L) \mid a$   
 $L \rightarrow L, E \mid E$



| State | Input |             |    |             |            | Goto |   |
|-------|-------|-------------|----|-------------|------------|------|---|
|       | (     | )           | a  | ,           | \$         | E    | L |
| 0     | s3    |             | s2 |             |            | 1    |   |
| 1     |       |             |    |             | accept     |      |   |
| 2     |       | r(E → a)    |    | r(E → a)    | r(E → a)   |      |   |
| 3     | s3    |             | s2 |             |            | 6    | 4 |
| 4     |       | s5          |    | s7          |            |      |   |
| 5     |       | r(E → (L))  |    | r(E → (L))  | r(E → (L)) |      |   |
| 6     |       | r(L → E)    |    | r(L → E)    |            |      |   |
| 7     | s3    |             | s2 |             |            | 8    |   |
| 8     |       | r(L → L, E) |    | r(L → L, E) |            |      |   |

# c. parse of ((a),a,(a,a))

|    | Parsing stack     | Input           | Action     |
|----|-------------------|-----------------|------------|
| 1  | \$0               | ((a),a,(a,a))\$ | s3         |
| 2  | \$0(3             | (a),a,(a,a))\$  | s3         |
| 3  | \$0(3(3           | a),a,(a,a))\$   | s2         |
| 4  | \$0(3(3a2         | )a,(a,a))\$     | r(E → a)   |
| 5  | \$0(3(3E6         | )a,(a,a))\$     | r(L → E)   |
| 6  | \$0(3(3L4         | )a,(a,a))\$     | s5         |
| 7  | \$0(3(3L4)5       | ,a,(a,a))\$     | r(E → (L)) |
| 8  | \$0(3E6           | ,a,(a,a))\$     | r(L → E)   |
| 9  | \$0(3L4           | ,a,(a,a))\$     | s7         |
| 10 | \$0(3L4,7         | a,(a,a))\$      | s2         |
| 11 | \$0(3L4,7a2       | ,(a,a))\$       | r(E → a)   |
| 12 | \$0(3L4,7E8       | ,(a,a))\$       | r(L → L,E) |
| 13 | \$0(3L4           | ,(a,a))\$       | s7         |
| 14 | \$0(3L4,7         | (a,a))\$        | s3         |
| 15 | \$0(3L4,7(3       | a,a))\$         | s2         |
| 16 | \$0(3L4,7(3a2     | ,a))\$          | r(E → a)   |
| 17 | \$0(3L4,7(3E6     | ,a))\$          | r(L → E)   |
| 18 | \$0(3L4,7(3L4     | ,a))\$          | s7         |
| 19 | \$0(3L4,7(3L4,7   | a))\$           | s2         |
| 20 | \$0(3L4,7(3L4,7a2 | ))\$            | r(E → a)   |
| 21 | \$0(3L4,7(3L4,7E8 | ))\$            | r(L → L,E) |
| 22 | \$0(33L4,7(3L4    | ))\$            | s5         |
| 23 | \$0(33L4,7(3L4)5  | )\$             | r(E → (L)) |
| 24 | \$0(33L4,7E8      | )\$             | r(L → L,E) |
| 25 | \$0(3L4           | )\$             | s5         |
| 26 | \$0(3L4)5         | \$              | r(E → (L)) |
| 27 | \$0E1             | \$              | accept     |

# 5.1

d. Is LR(0) grammar?

Yes. No conflicts.

| State | Action | Rule                | Input |   |   |   | Goto |   |
|-------|--------|---------------------|-------|---|---|---|------|---|
|       |        |                     | (     | a | ) | , | E    | L |
| 0     | shift  |                     | 3     | 2 |   |   | 1    |   |
| 1     | reduce | $E' \rightarrow E$  |       |   |   |   |      |   |
| 2     | reduce | $E \rightarrow a$   |       |   |   |   |      |   |
| 3     | shift  |                     | 3     | 2 |   |   | 6    | 4 |
| 4     | shift  |                     |       |   | 5 | 7 |      |   |
| 5     | reduce | $E \rightarrow (L)$ |       |   |   |   |      |   |
| 6     | reduce | $L \rightarrow E$   |       |   |   |   |      |   |
| 7     | shift  |                     | 3     | 2 |   |   | 8    |   |
| 8     | reduce | $L \rightarrow L,E$ |       |   |   |   |      |   |

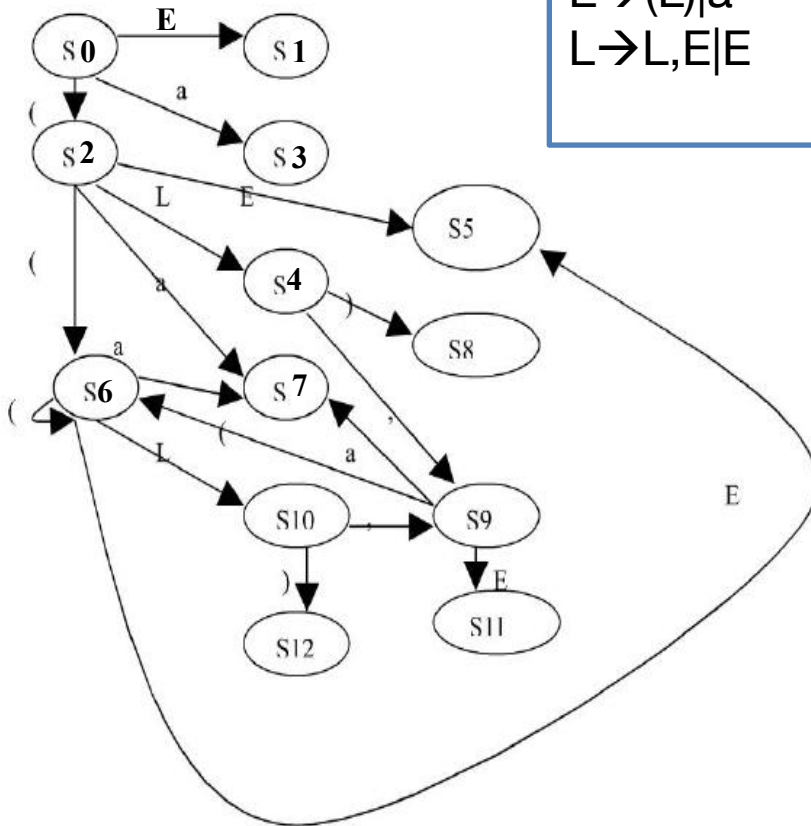
**5.2** Consider the grammar of the previous exercise.

- a.** Construct the DFA of LR(1) items for this grammar.
- b.** Construct the general LR(1) parsing table.
- c.** Construct the DFA of LALR(1) items for this grammar.
- d.** Construct the LALR(1) parsing table.

**e.** Describe any differences that might occur between the actions of a general LR(1) parser and an LALR(1) parser.

# 5.2

## a. DFA of LR(1) items



State 0:  $[E' \rightarrow \cdot E, \$]$

$[E \rightarrow \cdot (L), \$]$

$[E \rightarrow \cdot a, \$]$

State 2:  $[E \rightarrow ( \cdot L), \$]$

$[L \rightarrow \cdot L, E, )]$

$[L \rightarrow \cdot E, )]$

$[L \rightarrow \cdot L, E, , ]$

$[L \rightarrow \cdot E, , ]$

$[E \rightarrow \cdot (L), )]$

$[E \rightarrow \cdot a, )]$

$[E \rightarrow \cdot (L), , ]$

$[E \rightarrow \cdot a, , ]$

State 6:  $[E \rightarrow (L \cdot), )]$

$[E \rightarrow (L \cdot), , ]$

$[L \rightarrow \cdot L, E, )]$

$[L \rightarrow \cdot E, )]$

$[L \rightarrow \cdot L, E, , ]$

$[L \rightarrow \cdot E, , ]$

$[E \rightarrow \cdot (L), )]$

$[E \rightarrow \cdot a, )]$

$[E \rightarrow \cdot (L), , ]$

$[E \rightarrow \cdot a, , ]$

State 10:  $[E \rightarrow (L \cdot), )]$

$[E \rightarrow (L \cdot), , ]$

$[L \rightarrow L \cdot, E, )]$

$[L \rightarrow L \cdot, E, , ]$

State 1:  $[E' \rightarrow E \cdot, \$]$

State 3:  $[E \rightarrow a \cdot, \$]$

State 4:  $[E \rightarrow (L \cdot), \$]$

$[L \rightarrow L \cdot, E, )]$

$[L \rightarrow L \cdot, E, , ]$

State 5:  $[L \rightarrow E \cdot, )]$

$[L \rightarrow E \cdot, , ]$

State 7:  $[E \rightarrow a \cdot, )]$

$[E \rightarrow a \cdot, , ]$

State 8:  $[E \rightarrow (L \cdot), \$]$

State 9:  $[L \rightarrow L \cdot, E, )]$

$[E \rightarrow \cdot (L), )]$

$[E \rightarrow \cdot a, )]$

$[L \rightarrow L \cdot, E, , ]$

$[E \rightarrow \cdot (L), , ]$

$[E \rightarrow \cdot a, , ]$

State 11:  $[L \rightarrow L \cdot, E, )]$

$[L \rightarrow L \cdot, E, , ]$

State 12:  $[E \rightarrow (L \cdot), , ]$

$[E \rightarrow (L \cdot), , ]$

# 5.2

## b. LR(1) parsing table

b.

r1:  $E \rightarrow (L)$

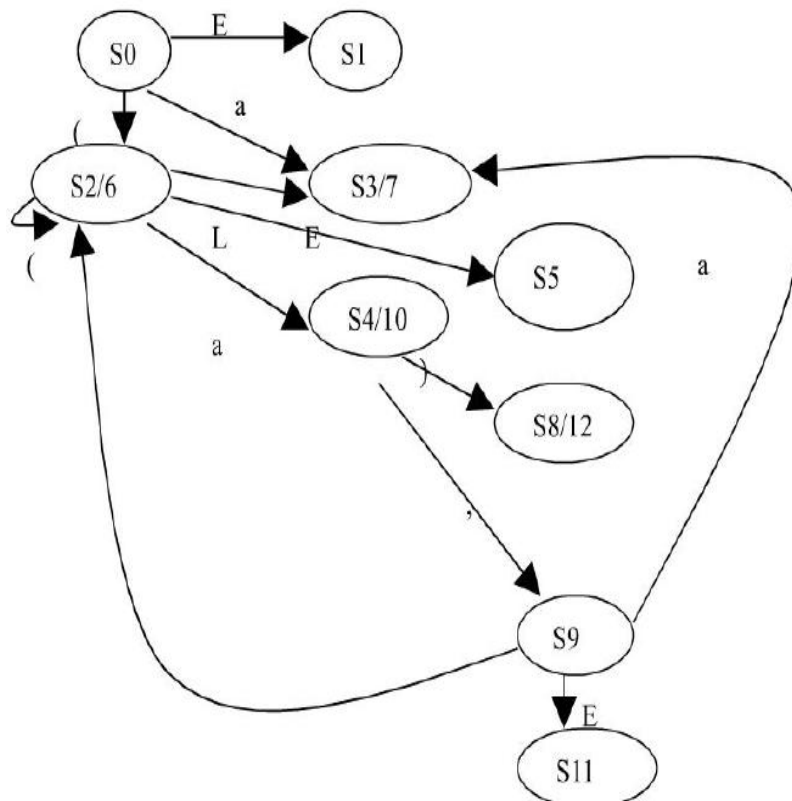
r2:  $E \rightarrow a$    r3:  $L \rightarrow L, E$    r4:  $L \rightarrow E$

| State | Input |    |     |    |        | Goto |    |
|-------|-------|----|-----|----|--------|------|----|
|       | (     | a  | )   | ,  | \$     | L    | E  |
| 0     | S2    | S3 |     |    |        |      | 1  |
| 1     |       |    |     |    | Accept |      |    |
| 2     | S6    | S7 |     |    |        | 4    | 5  |
| 3     |       |    |     |    | r2     |      |    |
| 4     |       |    | S8  | S9 |        |      |    |
| 5     |       |    | r4  | r4 |        |      |    |
| 6     | S6    | S7 |     |    |        | 10   | 5  |
| 7     |       |    | r2  | r2 |        |      |    |
| 8     |       |    |     |    | r1     |      |    |
| 9     | S6    | S7 |     |    |        |      | 11 |
| 10    |       |    | S12 | S9 |        |      |    |
| 11    |       |    | r3  | r3 |        |      |    |
| 12    |       |    | r1  | r1 |        |      |    |



# 5.2

## c. DFA of LALR(1) items



State 0:  $[E' \rightarrow .E, \$]$

$[E \rightarrow .(L), \$]$

$[E \rightarrow .a, \$]$

State 2/6:  $[E \rightarrow (L), \$ / )/, ]]$

$[L \rightarrow .L, E, )]$

$[L \rightarrow .E, )]$

$[L \rightarrow .L, E, , ]]$

$[L \rightarrow .E, , ]]$

$[E \rightarrow .(L), , ]]$

$[E \rightarrow .a, , ]]$

$[E \rightarrow .(L), , ]]$

$[E \rightarrow .a, , ]]$

State 8/12:  $[E \rightarrow (L), \$ / )/, ]]$

State 9:  $[L \rightarrow L, E, )]$

$[E \rightarrow .(L), )]$

$[E \rightarrow .a, )]$

$[L \rightarrow L, E, , ]]$

$[E \rightarrow .(L), , ]]$

$[E \rightarrow .a, , ]]$

State 1:  $[E' \rightarrow E., \$]$

State 3/7:  $[E \rightarrow a., \$ / )/, ]]$

State 4/10:  $[E \rightarrow (L.), \$ / )/, ]]$

$[L \rightarrow L., E, )]$

$[L \rightarrow L., E, , ]]$

State 5:  $[L \rightarrow E., )]$

$[L \rightarrow E., , ]]$

State 11:  $[L \rightarrow L, E., )]$

$[L \rightarrow L, E., , ]]$

# 5.2

## d. LALR(1) parsing table

d.  $r1: E \rightarrow (L)$   $r2: E \rightarrow a$   $r3: L \rightarrow L, E$   $r4: L \rightarrow E$

| State | Input |      |       |    |        | Goto |    |
|-------|-------|------|-------|----|--------|------|----|
|       | (     | a    | )     | ,  | \$     | L    | E  |
| 0     | S2/6  | S3/7 |       |    |        |      | 1  |
| 1     |       |      |       |    | Accept |      |    |
| 2/6   | S2/6  | S3/7 |       |    |        | 4/10 | 5  |
| 3/7   |       |      | r2    | r2 | r2     |      |    |
| 4/10  |       |      | S8/12 | S9 |        |      |    |
| 5     |       |      | r4    | r4 |        |      |    |
| 8/12  |       |      | r1    | r1 | r1     |      |    |
| 9     | S2/6  | S3/7 |       |    |        |      | 11 |
| 10    |       |      | S8/12 | S9 |        |      |    |
| 11    |       |      | r3    | r3 |        |      |    |