Cryptography Homework 1

2024 Spring Semester

21 CST H3Art

Exercise 2.1

Evaluate the following:

- (a) $7503 \mod 81$
- (b) $(-7503) \mod 81$
- (c) $81 \mod 7503$
- (d) $(-81) \mod 7503$

```
Solution:
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```
(a) Since 7503 = 92 \times 81 + 51, 0 < 51 < 81, 7503 \mod 81 = 51

(b) Since -7503 = -93 \times 81 + 30, 0 < 30 < 81, (-7503) \mod 81 = 30

(c) Since 81 = 0 \times 7503 + 81, 0 < 81 < 7503, 81 \mod 7503 = 81

(d) Since -81 = 1 \times 7503 + 7422, 0 < 7422 < 7503, (-81) \mod 7503 = 7422
```

Exercise 2.7

Determine the number of keys in an Affine Cipher over \mathbb{Z}_m for m=30,100 and 1225.

Solution:

To find the prime factors of a number, we used this code below.

```
num = int(input())
prime_factor = []

for i in range(2, num):
    while True:
        if num % i == 0:
            num /= i
            prime_factor.append(i)
        else:
            break

print(prime_factor)
```

(a) For m=30, to find the number of keys, we need to factor the prime factors of 30 to get $30=2\times3\times5$. Using the Euler function, we can calculate that there are $\Phi(30)=(2-1)(3-1)(5-1)=8$ numbers satisfying $\gcd(a,30)=1$. The number of keys in an Affine Cipher over \mathbb{Z}_{30} is $30\times8=240$.

- (b) For m=100, to find the number of keys, we need to factor the prime factors of 100 to get $100=2^2\times 5^2$. Using the Euler function, we can calculate that there are $\Phi(100)=(2^2-2)(5^2-5)=40$ numbers satisfying $\gcd(a,100)$. The number of keys in an Affine Cipher over \mathbb{Z}_{100} is $100\times 40=4000$.
- (c) For m=1225, to find the number of keys, we need to factor the prime factors of 1225 to get $1225=5^2\times 7^2$. Using the Euler function, we can calculate that there are $\Phi(1225)=(5^2-5)(7^2-7)=840$ numbers satisfying $\gcd(a,1225)$. The number of keys in an Affine Cipher over \mathbb{Z}_{1225} is $1225\times 840=1029000$.

List all the invertible elements in \mathbb{Z}_m for m=28,33, and 35.

Exercise 2.9

For $1 \le a \le 28$, determine $a^{-1} \bmod 29$ by trial and error.

Solution:

Since 29's factors are only 1 and 29, 29 is a prime. Therefore, for $1 \le a \le 28$, all $a^{-1}a \equiv 1 \mod 29$ have solution.

To determine $a^{-1} \mod 29$ by trial and error:

```
for i in range(1, 29):
    for j in range(1, 100):
        if (i * j) % 29 == 1:
            print('For {0}, the invert is {1}.'.format(i, j))
            break
```

Finally, all the results are as follows:

```
For 1 , the invert is 1.
For 2 , the invert is 15.
```

```
For 3 , the invert is 10.
For 4 , the invert is 22.
For 5 , the invert is 6.
For 6, the invert is 5.
For 7 , the invert is 25.
For 8 , the invert is 11.
For 9 , the invert is 13.
For 10 , the invert is 3.
For 11 , the invert is 8.
For 12 , the invert is 17.
For 13 , the invert is 9.
For 14 , the invert is 27.
For 15 , the invert is 2.
For 16 , the invert is 20.
For 17, the invert is 12.
For 18, the invert is 21.
For 19, the invert is 26.
For 20, the invert is 16.
For 21, the invert is 18.
For 22 , the invert is 4.
For 23 , the invert is 24.
For 24, the invert is 23.
For 25, the invert is 7.
For 26, the invert is 19.
For 27, the invert is 14.
For 28, the invert is 28.
```

Suppose that K=(5,21) is a key in an Affine Cipher over \mathbb{Z}_{29} .

- (a) Express the decryption function $d_K(y)$ in the form $d_K(y)=a'y+b'$, where $a',b'\in\mathbb{Z}_{29}$
- (b) Prove that $d_K(e_K(x)) = x$ for all $x \in \mathbb{Z}_{29}$.

Solution:

```
(a) Since e_K(x) = (5x+21) \mod 29 = 5x+21 and d_K(y) = a^{-1}(y-b) \mod 29, where a=5, b=21. According to aa^{-1} \equiv 1 \mod 29, the value of a^{-1} is 6. Thus, d_K(y) = 6(y-21) \mod 29 = 6y-10.
```

(b) Proof:

$$d_K(e_K(x)) = a^{-1}(ax + b - b) \mod 29$$

= $6 \times 5x \mod 29$
= $30x \mod 29$
= x

Exercise 2.15(a)

Determine the inverses of the following matrices over \mathbb{Z}_{26} :

(a)
$$\begin{pmatrix} 2 & 5 \\ 9 & 5 \end{pmatrix}$$

Solution:

Suppose $\mathbf{A} = \begin{pmatrix} 2 & 5 \\ 9 & 5 \end{pmatrix}$, to find \mathbf{A}^{-1} over \mathbb{Z}_{26} , we have an equation $\mathbf{A}^{-1} = (\det \mathbf{A})^{-1} \mathbf{A}^*$, where \mathbf{A}^* is adjugate matrix of \mathbf{A} .

$$\mathbf{A}^* = \begin{pmatrix} (-1)^{1+1} \times a_{2,2} & (-1)^{1+2} \times a_{1,2} \\ (-1)^{2+1} \times a_{2,1} & (-1)^{2+2} \times a_{1,1} \end{pmatrix} \mod 26$$

$$= \begin{pmatrix} (-1)^{1+1} \times 5 & (-1)^{1+2} \times 5 \\ (-1)^{2+1} \times 9 & (-1)^{2+2} \times 2 \end{pmatrix} \mod 26$$

$$= \begin{pmatrix} 5 & -5 \\ -9 & 2 \end{pmatrix} \mod 26$$

$$= \begin{pmatrix} 5 & 21 \\ 17 & 2 \end{pmatrix}$$

and for $\det \mathbf{A}$ over \mathbb{Z}_{26} :

$$det(\mathbf{A}) = det \begin{pmatrix} 2 & 5 \\ 9 & 5 \end{pmatrix} \mod 26$$
$$= (2 \times 5 - 5 \times 9) \mod 26$$
$$= -35 \mod 26$$
$$= 17$$

Since $17^{-1} \mod 26 = 23$

$$\mathbf{A}^{-1} = 23\mathbf{A}^*$$

$$= 23 \begin{pmatrix} 5 & 21 \\ 17 & 2 \end{pmatrix} \mod 26$$

$$= \begin{pmatrix} 11 & 15 \\ 1 & 20 \end{pmatrix}$$

Therefore, the inverse of the matrix $\begin{pmatrix} 2 & 5 \\ 9 & 5 \end{pmatrix}$ over \mathbb{Z}_{26} is $\begin{pmatrix} 11 & 15 \\ 1 & 20 \end{pmatrix}$.

Exercise 2.16

(a) Suppose that π is the following permutation of $\{1,...,8\}$:

x	1	2	3	4	5	6	7	8
$\pi(x)$	4	1	6	2	7	3	8	5

Compute the permutation π^{-1} .

(b) Decrypt the following ciphertext, for a *Permutation Cipher* with m=8, which was encrypted using the key π :

TGEEMNELNNTDROEOAAHDOETCSHAEIRLM.

Solution:

(a) The permutation π^{-1} is as follows:

x	1	2	3	4	5	6	7	8
$\pi^{-1}(x)$	2	4	6	1	8	3	5	7

(b) We need to partition the ciphertext into 8 letters per group:

TGEEMNEL NNTDROEO AAHDOETC SHAEIRLM.

Next we only need to use $\pi^{-1}(x)$ to get:

GENTLEME NDONOTRE ADEACHOT HERSMAIL.

Therefore, the plaintext is:

GENTLEMEN DO NOT READ EACH OTHERS MAIL.

Exercise 2.18

Consider the following linear recurrence over \mathbb{Z}_2 of degree four:

$$z_{i+4} = (z_i + z_{i+1} + z_{i+2} + z_{i+3}) \bmod 2$$

 $i \geq 0$. For each of the 16 possible initialization vectors $(z_0, z_1, z_2, z_3) \in (\mathbb{Z}_2)^4$, determine the period of the resulting keystream.

Solution:

- Start with (0,0,0,0), the keystream is $0,0,0,0,0,0,0,0,0,0,0,0,\cdots$, the period of the resulting keystream is 1.
- Start with (1,0,0,0), the keystream is $1,0,0,0,1,1,0,0,1,1,0,\cdots$, the period of the resulting keystream is 5.
- Start with (0,1,0,0), the keystream is $0,1,0,0,1,0,1,0,1,0,1,\cdots$, the period of the resulting keystream is 5
- Start with (0,0,1,0), the keystream is $0,0,1,0,1,0,0,1,0,0,\cdots$, the period of the resulting keystream is 5.
- Start with (0,0,0,1), the keystream is $0,0,0,1,1,0,0,0,1,1,0,0,\cdots$, the period of the resulting keystream is 5
- Start with (1, 1, 0, 0), the keystream is $1, 1, 0, 0, 0, 1, 1, 0, 0, 0, 1, 1, \cdots$, the period of the resulting keystream is 5.
- Start with (1,0,1,0), the keystream is $1,0,1,0,0,1,0,1,0,0,1,0,\cdots$, the period of the resulting keystream is 5.
- Start with (1,0,0,1), the keystream is $1,0,0,1,0,1,0,1,0,1,0,\cdots$, the period of the resulting keystream is 5.
- Start with (0, 1, 1, 0), the keystream is $0, 1, 1, 0, 0, 0, 1, 1, 0, 0, 0, 1, \cdots$, the period of the resulting keystream is 5.
- Start with (0,1,0,1), the keystream is $0,1,0,1,0,0,1,0,0,1,\cdots$, the period of the resulting keystream is 5.
- Start with (0,0,1,1), the keystream is $0,0,1,1,0,0,0,1,1,0,0,0,\cdots$, the period of the resulting keystream is 5.
- Start with (1,1,1,0), the keystream is $1,1,1,0,1,1,1,0,1,1,1,\cdots$, the period of the resulting keystream is 5
- Start with (1, 1, 0, 1), the keystream is $1, 1, 0, 1, 1, 1, 1, 0, 1, 1, 1, 1, \dots$, the period of the resulting keystream is 5.

- Start with (1,0,1,1), the keystream is $1,0,1,1,1,1,0,1,1,1,0,\cdots$, the period of the resulting keystream is 5.
- Start with (0,1,1,1), the keystream is $0,1,1,1,1,0,1,1,1,0,1,\cdots$, the period of the resulting keystream is 5
- Start with (1, 1, 1, 1), the keystream is $1, 1, 1, 1, 0, 1, 1, 1, 0, 1, 1, \cdots$, the period of the resulting keystream is 5.

Suppose we are told that the plaintext

breathtaking

yields the ciphertext

RUPOTENTOIFV

where the *Hill Cipher* is used (but m is not specified). Determine the encryption matrix.

Solution:

Firstly, we suppose m=1, then $e_{\mathbf{K}}(\mathbf{b})=e_{\mathbf{K}}(1)=(\mathbf{R})=(17)$. The encrption equation is as follows:

$$(1)$$
K = (17)

To test whether $\mathbf{K}=(17)$, we also have $e_{\mathbf{K}}(\mathbf{r})=e_{\mathbf{K}}(17)=(\mathrm{U})=(20)$, since

$$(17)(17) \mod 26 = (3) \neq (20)$$

The ${\bf K}$ we calculated above is incorrect, $m \neq 1$.

Secondly, suppose m=2, the first two processes of Hill Cipher encryption is as follows:

$$e_{\mathbf{K}}(\mathbf{b}, \mathbf{r}) = e_{\mathbf{K}}(1, 17) = (\mathbf{R}, \mathbf{U}) = (17, 20)$$

$$e_{\mathbf{K}}(\mathbf{e}, \mathbf{a}) = e_{\mathbf{K}}(4, 0) = (\mathbf{P}, \mathbf{O}) = (15, 14)$$

Therefore, we can form a matrix equation:

$$\begin{pmatrix} 1 & 17 \\ 4 & 0 \end{pmatrix} \mathbf{K} = \begin{pmatrix} 17 & 20 \\ 15 & 14 \end{pmatrix}$$

The next thing to do is to find the inverse of \mathbf{K} , namely \mathbf{K}^{-1} .

However, when we calculate $\det \begin{pmatrix} 1 & 17 \\ 4 & 0 \end{pmatrix} \mod 26 = (1 \times 0 - 17 \times 4) \mod 26 = 10$, we cannot find the result of $10^{-1} \mod 26$, it means $\det \begin{pmatrix} 1 & 17 \\ 4 & 0 \end{pmatrix}$ doesn't have inverse over \mathbb{Z}_{26} , m can't be 2.

Nextly, suppose m=3, the first three processes of Hill Cipher encryption is as follows:

$$e_{\mathbf{K}}(\mathbf{b}, \mathbf{r}, \mathbf{e}) = e_{\mathbf{K}}(1, 17, 4) = (\mathbf{R}, \mathbf{U}, \mathbf{P}) = (17, 20, 15)$$

$$e_{\mathbf{K}}(\mathbf{a}, \mathbf{t}, \mathbf{h}) = e_{\mathbf{K}}(0, 19, 7) = (0, T, \mathbf{E}) = (14, 19, 4)$$

$$e_{\mathbf{K}}(t, a, k) = e_{\mathbf{K}}(19, 0, 10) = (N, T, O) = (13, 19, 14)$$

We can form another matrix equation:

$$\begin{pmatrix} 1 & 17 & 4 \\ 0 & 19 & 7 \\ 19 & 0 & 10 \end{pmatrix} \mathbf{K} = \begin{pmatrix} 17 & 20 & 15 \\ 14 & 19 & 4 \\ 13 & 19 & 14 \end{pmatrix}$$

Next, assume
$$\mathbf{A} = \begin{pmatrix} 1 & 17 & 4 \\ 0 & 19 & 7 \\ 19 & 0 & 10 \end{pmatrix}$$
, we calculate $\det \mathbf{A} = \det \begin{pmatrix} 1 & 17 & 4 \\ 0 & 19 & 7 \\ 19 & 0 & 10 \end{pmatrix} \mod 26 = 19$, since

 $19^{-1} \bmod 26 = 11$, finally we need to find the adjugate matrix ${\bf A}^*$ to get ${\bf A}^{-1}$:

$$\mathbf{A}_{11}^* = (-1)^{1+1} (19 \times 10 - 7 \times 0) \bmod 26 = 8$$

$$\mathbf{A}_{12}^* = (-1)^{2+1} (17 \times 10 - 4 \times 0) \bmod 26 = 12$$

$$\mathbf{A}_{13}^* = (-1)^{3+1} (17 \times 7 - 4 \times 19) \bmod 26 = 17$$

$$\mathbf{A}_{21}^* = (-1)^{1+2} (0 \times 10 - 7 \times 19) \bmod 26 = 3$$

$$\mathbf{A}_{22}^* = (-1)^{2+2} (1 \times 10 - 4 \times 19) \bmod 26 = 12$$

$$\mathbf{A}_{23}^* = (-1)^{3+2} (1 \times 7 - 4 \times 0) \bmod 26 = 19$$

$$\mathbf{A}_{31}^* = (-1)^{1+3} (0 \times 0 - 19 \times 19) \bmod 26 = 3$$

$$\mathbf{A}_{32}^* = (-1)^{2+3} (1 \times 0 - 17 \times 19) \bmod 26 = 11$$

$$\mathbf{A}_{33}^* = (-1)^{3+3} (1 \times 19 - 17 \times 0) \bmod 26 = 19$$

Therefore,
$$\mathbf{A}^* = \begin{pmatrix} 8 & 12 & 17 \\ 3 & 12 & 19 \\ 3 & 11 & 19 \end{pmatrix}$$

$$\mathbf{A}^{-1} = 11\mathbf{A}^*$$

$$= 11 \begin{pmatrix} 8 & 12 & 17 \\ 3 & 12 & 19 \\ 3 & 11 & 19 \end{pmatrix} \mod 26$$

$$= \begin{pmatrix} 10 & 2 & 5 \\ 7 & 2 & 1 \\ 7 & 17 & 1 \end{pmatrix}$$

After finding \mathbf{A}^{-1} , the key matrix \mathbf{K} can be found:

$$\mathbf{K} = \mathbf{A}^{-1} \begin{pmatrix} 17 & 20 & 15 \\ 14 & 19 & 4 \\ 13 & 19 & 14 \end{pmatrix}$$

$$= \begin{pmatrix} 10 & 2 & 5 \\ 7 & 2 & 1 \\ 7 & 17 & 1 \end{pmatrix} \begin{pmatrix} 17 & 20 & 15 \\ 14 & 19 & 4 \\ 13 & 19 & 14 \end{pmatrix} \mod 26$$

$$= \begin{pmatrix} 3 & 21 & 20 \\ 4 & 15 & 23 \\ 6 & 14 & 5 \end{pmatrix}$$

Let's test whether the ${\bf K}$ can encrypt the final three letters "ing" to "INV", namely satisfy $e_K(i,n,g)=e_K(8,13,6)=(I,N,V)=(8,5,21)$.

$$\begin{pmatrix} 8 & 13 & 6 \end{pmatrix} \begin{pmatrix} 3 & 21 & 20 \\ 4 & 15 & 23 \\ 6 & 14 & 5 \end{pmatrix} \mod 26 = \begin{pmatrix} 8 & 5 & 21 \end{pmatrix}$$

So we can conclude that
$$m=3$$
 and the key matrix $\mathbf{K}=\begin{pmatrix} 3 & 21 & 20 \\ 4 & 15 & 23 \\ 6 & 14 & 5 \end{pmatrix}$.

We describe another stream cipher, which incorporates one of the ideas from the *Enigma* machime used by Germany in World War II. Suppose that π is a fixed permutation of \mathbb{Z}_{26} . The key is an element $K \in \mathbb{Z}_{26}$. For all integers $i \geq 1$, the keystream element $z_i \in \mathbb{Z}_{26}$ is defined according to the rule $z_i = (K+i-1) \mod 26$. Encryption and decryption are performed using the permutations π and π^{-1} , respectively, as follows:

$$e_z(x) = \pi(x) + z \mod 26$$

and

$$d_z(y)=\pi^{-1}(y-z \bmod 26)$$

where $z \in \mathbb{Z}_{26}$.

Suppose that π is the following permutation of \mathbb{Z}_{26} :

\boldsymbol{x}	0	1	2	3	4	5	6	7	8	9	10	11	12
$\pi(x)$	23	13	24	0	7	15	14	6	25	16	22	1	19

x	13	14	15	16	17	18	19	20	21	22	23	24	25
$\pi(x)$	18	5	11	17	2	21	12	20	4	10	9	3	8

The following ciphertext has been encrypted using this stream cipher; use exhaustive key search to decrypt it:

$\begin{array}{l} WRTCNRLDSAFARWKXFTXCZRNHNYPDTZ\\ UUKMPLUSOXNEUDOKLXRMCBKGRCCURR \end{array}$

Solution:

Based on the decryption principles of stream cipher and permutation ciphers, I constructed the following Python code, which can help me decrypt using the exhaustive key search:

```
import math
PI_inv = {
    0: 3, 1: 11, 2: 17, 3: 24, 4: 21, 5: 14, 6: 7, 7: 4, 8: 25, 9: 23, 10: 22,
   11: 15, 12: 19, 13: 1, 14: 6, 15: 5, 16: 9, 17: 16, 18: 13, 19: 12, 20: 20,
    21: 18, 22: 10, 23: 0, 24: 2, 25: 8
}
ciphertext = 'WRTCNRLDSAFARWKXFTXCZRNHNYPDTZUUKMPLUSOXNEUDOKLXRMCBKGRCCURR'
def decrypt(ciphertext: str, K: int) -> str:
    plaintext = []
    for index, ch in enumerate(ciphertext):
        plain_ch = (ord(ch) - 65) - (K + index) % 26
        if plain_ch < 0:</pre>
            plaintext.append(chr(PI_inv.get(26 + plain_ch) + 65))
        else:
            plaintext.append(chr(PI_inv.get(plain_ch) + 65))
    return ''.join(plaintext)
```

```
def main() -> None:
    for K in range(0, 26):
        print('Current K = {0}, the decrypted text is {1}.'.format(K, decrypt(ciphertext, K)))

if __name__ == '__main__':
    main()
```

Running the above code, the output is:

```
Current K = 0, the decrypted text is KJQIXTOKWQSFOXKZFROXOKQWFIFRQKJFVOERWERWIFVTKQQRSFVRWOFICFPW.
Current K = 1, the decrypted text is SFJCZPVSXJUGVZSEGLVZVSJXGCGLJSFGYVHLXHLXCGYPSJJLUGYLXVGCAGWX.
Current K = 2, the decrypted text is UGFAEWYUZFMBYEUHBDYEYUFZBABDFUGBRYODZODZABRWUFFDMBRDZYBAKBXZ.
Current K = 3, the decrypted text is MBGKHXRMEGNTRHMOTIRHRMGETKTIGMBTLRVIEVIEKTLXMGGINTLIERTKSTZE.
{\tt Current~K~=~4,~the~decrypted~text~is~NTBSOZLNHBQPLONVPCLOLNBHPSPCBNTPDLYCHYCHSPDZNBBCQPDCHLPSUPEH.}
Current K = 5, the decrypted text is QPTUVEDQOTJWDVQYWADVDQTOWUWATQPWIDRAORAOUWIEQTTAJWIAODWUMWHO.
Current K = 6, the decrypted text is JWPMYHIJVPFXIYJRXKIYIJPVXMXKPJWXCILKVLKVMXCHJPPKFXCKVIXMNXOV.
Current K = 7, the decrypted text is FXWNROCFYWGZCRFLZSCRCFWYZNZSWFXZACDSYDSYNZAOFWWSGZASYCZNQZVY.
Current K = 8, the decrypted text is GZXQLVAGRXBEALGDEUALAGXREQEUXGZEKAIURIURQEKVGXXUBEKURAEQJEYR.
Current K = 9, the decrypted text is BEZJDYKBLZTHKDBIHMKDKBZLHJHMZBEHSKCMLCMLJHSYBZZMTHSMLKHJFHRL.
Current K = 10, the decrypted text is THEFIRSTDEPOSITCONSISTEDOFONETHOUSANDANDFOURTEENPOUNDSOFGOLD.
Current K = 11, the decrypted text is POHGCLUPIHWVUCPAVQUCUPHIVGVQHPOVMUKQIKQIGVMLPHHQWVMQIUVGBVDI.
Current K = 12, the decrypted text is WVOBADMWCOXYMAWKYJMAMWOCYBYJOWVYNMSJCSJCBYNDWOOJXYNJCMYBTYIC.
Current K = 13, the decrypted text is XYVTKINXAVZRNKXSRFNKNXVARTRFVXYRQNUFAUFATRQIXVVFZRQFANRTPRCA.
Current K = 14, the decrypted text is ZRYPSCQZKYELQSZULGQSQZYKLPLGYZRLJQMGKMGKPLJCZYYGELJGKQLPWLAK.
{\tt Current~K~=~15,~the~decrypted~text~is~ELRWUAJESRHDJUEMDBJUJERSDWDBRELDFJNBSNBSWDFAERRBHDFBSJDWXDKS.}
\texttt{Current} \ \ \texttt{K} = \texttt{16}, \ \ \texttt{the} \ \ \texttt{decrypted} \ \ \ \texttt{text} \ \ \textbf{is} \ \ \texttt{HDLXMKFHULOIFMHNITFMFHLUIXITLHDIGFQTUQTUXIGKHLLTOIGTUFIXZISU}.
Current K = 17, the decrypted text is OIDZNSGOMDVCGNOQCPGNGODMCZCPDOICBGJPMJPMZCBSODDPVCBPMGCZECUM.
Current K = 18, the decrypted text is VCIEQUBVNIYABQVJAWBQBVINAEAWIVCATBFWNFWNEATUVIIWYATWNBAEHAMN.
Current K = 19, the decrypted text is YACHJMTYQCRKTJYFKXTJTYCQKHKXCYAKPTGXQGXQHKPMYCCXRKPXQTKHOKNQ.
Current K = 20, the decrypted text is RKAOFNPRJALSPFRGSZPFPRAJSOSZARKSWPBZJBZJOSWNRAAZLSWZJPSOVSQJ.
Current K = 21, the decrypted text is LSKVGQWLFKDUWGLBUEWGWLKFUVUEKLSUXWTEFTEFVUXQLKKEDUXEFWUVYUJF.
Current K = 22, the decrypted text is DUSYBJXDGSIMXBDTMHXBXDSGMYMHSDUMZXPHGPHGYMZJDSSHIMZHGXMYRMFG.
Current K = 23, the decrypted text is IMURTFZIBUCNZTIPNOZTZIUBNRNOUIMNEZWOBWOBRNEFIUUOCNEOBZNRLNGB.
{\tt Current} \ {\tt K = 24, the decrypted text is CNMLPGECTMAQEPCWQVEPECMTQLQVMCNQHEXVTXVTLQHGCMMVAQHVTEQLDQBT.}
Current K = 25, the decrypted text is AQNDWBHAPNKJHWAXJYHWHANPJDJYNAQJOHZYPZYPDJOBANNYKJOYPHJDIJTP.
```

Through observation, we can get that the plaintext should be this row:

Current K = 10, the decrypted text is THEFIRSTDEPOSITCONSISTEDOFONETHOUSANDANDFOURTEENPOUNDSOFGOLD.

Therefore, the plaintext is as follows:

THE FIRST DEPOSIT CONSISTED OF ONE THOUSAND AND FOURTEEN POUNDS OF GOLD