

## **QUESTIONS ON SERIALIZABILITY IN TRANSACTION**

### **Testing Conflict and View Serializability**

To test whether a schedule is conflict-serializable, you can use the precedence graph (or conflict graph). The process involves the following steps:

- Step 1.

Identify Conflicting Operations: Conflicts occur when two transactions access the same data item, and at least one of them is a write operation. Common conflicts are:

- Write-Read (WR): A transaction writes a data item, and another transaction reads it.
- Read-Write (RW): A transaction reads a data item, and another transaction writes to it.
- Write-Write (WW): Two transactions write to the same data item.

- Step 2.

Construct the Precedence Graph:

- Create a node for each transaction in the schedule.
- Draw a directed edge from transaction  $T_i$  to transaction  $T_j$  if  $T_i$  performs a conflicting operation before  $T_j$ .

- Step 3.

Check for Cycles:

- If the precedence graph has no cycles, the schedule is conflict-serializable.
- If there are cycles, the schedule is not conflict-serializable.

#### **Q.1**

You are given a schedule as below. Check whether schedule is conflict serializable or not  
schedule S:

S:  $r_1(x)$   $r_1(y)$   $w_2(x)$   $w_1(x)$   $r_2(y)$

#### **Q.2** Consider another schedule S1:

S1:  $r_1(x)$   $r_3(y)$   $w_1(x)$   $w_2(y)$   $r_3(x)$   $w_2(x)$

#### **Q.3** Check whether the given schedule S is conflict serializable or not-

S :  $R_1(A)$  ,  $R_2(A)$  ,  $R_1(B)$  ,  $R_2(B)$  ,  $R_3(B)$  ,  $W_1(A)$  ,  $W_2(B)$

#### **Q4.** Check whether the given schedule S is conflict serializable and recoverable or not-

T1	T2	T3	T4
W(X) Commit	R(X)  W(Y) R(Z) Commit	W(X) Commit	R(X) R(Y) Commit

**Q5.** Check whether the given schedule S is conflict serializable or not.

T1	T2	T3	T4
W(B)	R(A)  W(A)  W(B)	R(A)  R(B)	R(A)

**Q 6.** Given schedule S, check whether the given schedule S is conflict serializable or not.

T1	T2
R(A)   W(A) R(B) W(B)	R(A) W(A) R(B)   W(B)

**Q.7** Check given schedule S1 & S2 is view Serializable or not  
S1: R1(A),W1(A),R2(A),W2(A),R1(B),W1(B),R2(B),W2(B)  
S2:R1(A),W1(A),R1(B),W1(B),R2(A),W2(A),R2(B),W2(B)

**Q.8** Check whether schedule is view equivalent or not?

T1	T2	T3
Read(A)	Write(A)	Write(A)
Write(A)		

**Schedule S**

T1	T2	T3
Read(A)	Write(A)	Write(A)
Write(A)		

**Schedule S1**

**Q.9** Check given schedule S1 & S3 is view Serializable or not

S1: R1(A), W1(A), R2(A), W2(A), R1(B), W1(B), R2(B), W2(B)

S3: R2(A), W2(A), R2(B), W2(B), R1(A), W1(A), R1(B), W1(B),

**Q.10** Check given schedule S is Conflict Serializable or not

S2: r1(x), r2(x), r2(y), w2(y), r1(y), w1(x)

**Q.11** Check given schedule S is Conflict Serializable or not

S1: r1(x), r1(y), r2(x), r2(y), w2(y), w1(x)