



University of
Southampton

Generating SPARK from Event-B

Providing Fundamental Safety and Security

Asieh Salehi, Thai Son Hoang, Dana Dghaym, Michael Butler and **Colin Snook**

Outline of talk

- Background
 - Event-B, CamilleX & Event-B notation extensions (including Records)
 - Spark
- Overview -
 - why generate SPARK from Event-B
 - From abstract record structures to SPARK
- SPARK code generation
 - Overview of translation rules (inc. records)
- SBB Electronic Voting Case study
 - using latest CamilleX and Records
 - SBB final refinement -> SPARK specifications

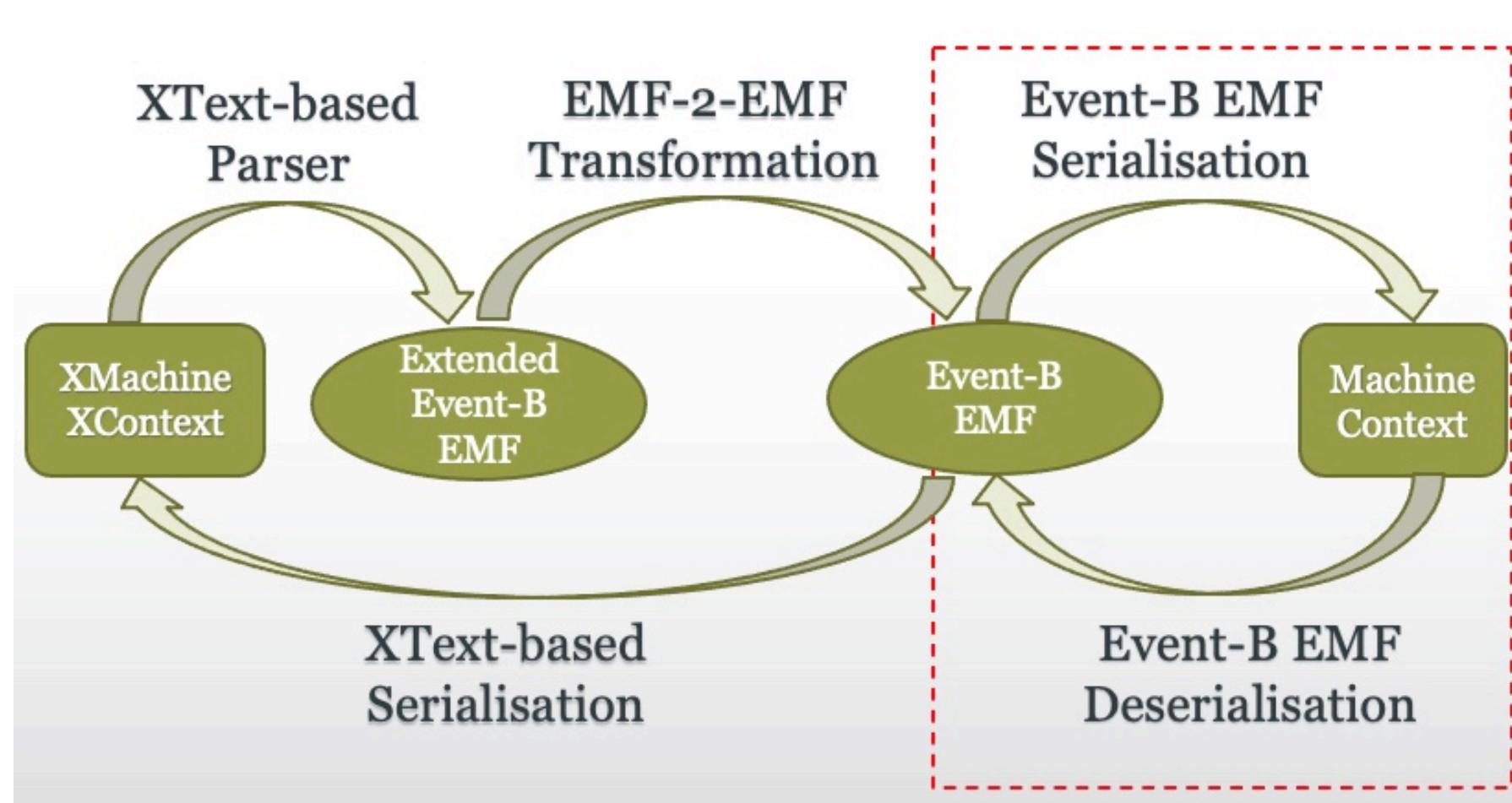
Background - Event-B

- Discrete transition systems
 - **Variables** representing states
 - **Guarded events** representing transitions
 - **Contexts**: Static part of the models (carrier sets, constants, etc.)
 - **Machines**: Dynamic part of the models (variables, events, etc.)
- First-order logic with set theory
- Refinement
 - Start with a simple abstract model
 - Add detail and design in small steps
- Verification by automatic theorem provers
- Validation by model checking
- Model checker also useful for liveness and debugging

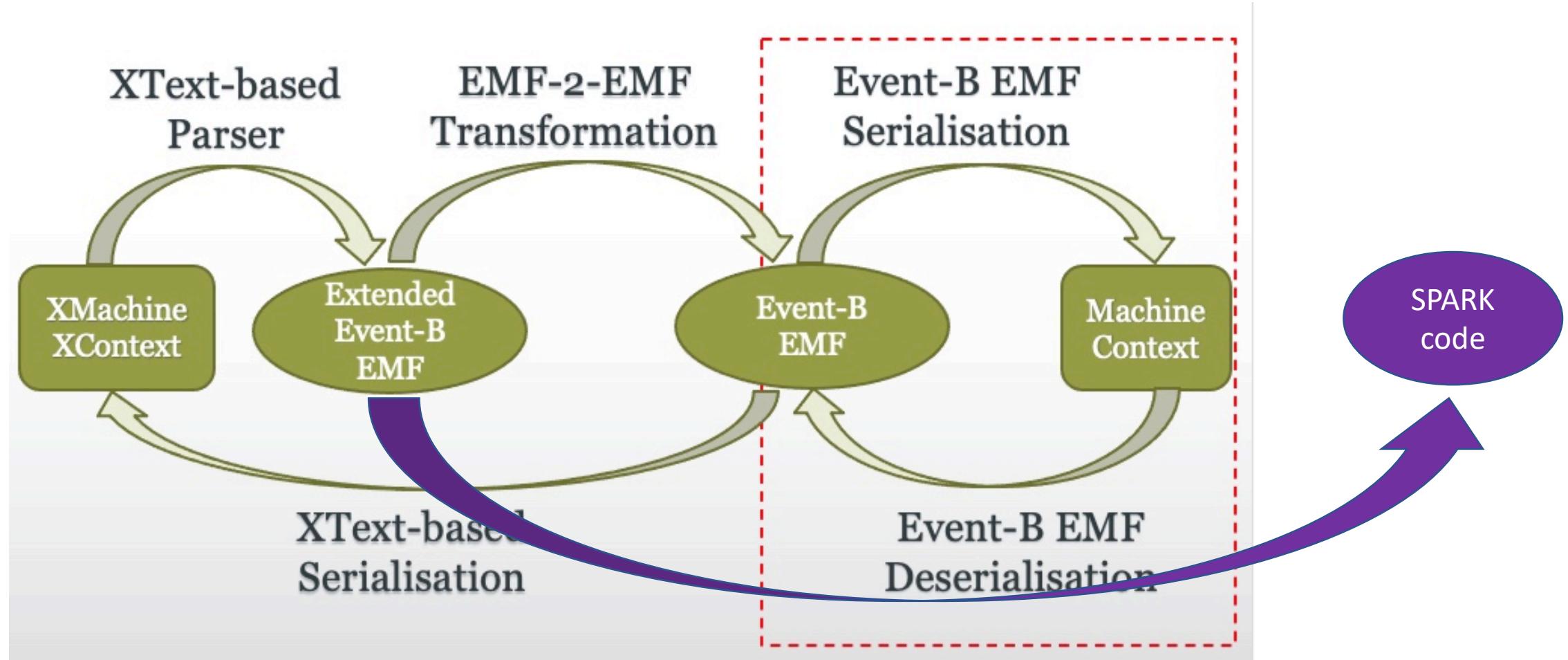
The Need for Textual Representation

- (True) Textual representation helps with teamworking
- Framework (e.g., XText) for developing IDE for DSLs.
- Design Principles:
 1. Reuse the existing Event-B tools of Rodin as much as possible.
 2. Support direct extension of the Event-B syntax to provide additional features.
 3. Provide compatibility with other kinds of 'higher-level' models that contribute to the overall model, e.g., UML-B diagrams.
- We make use of the Event-B EMF and EMF-2-EMF framework

The CamilleX Framework



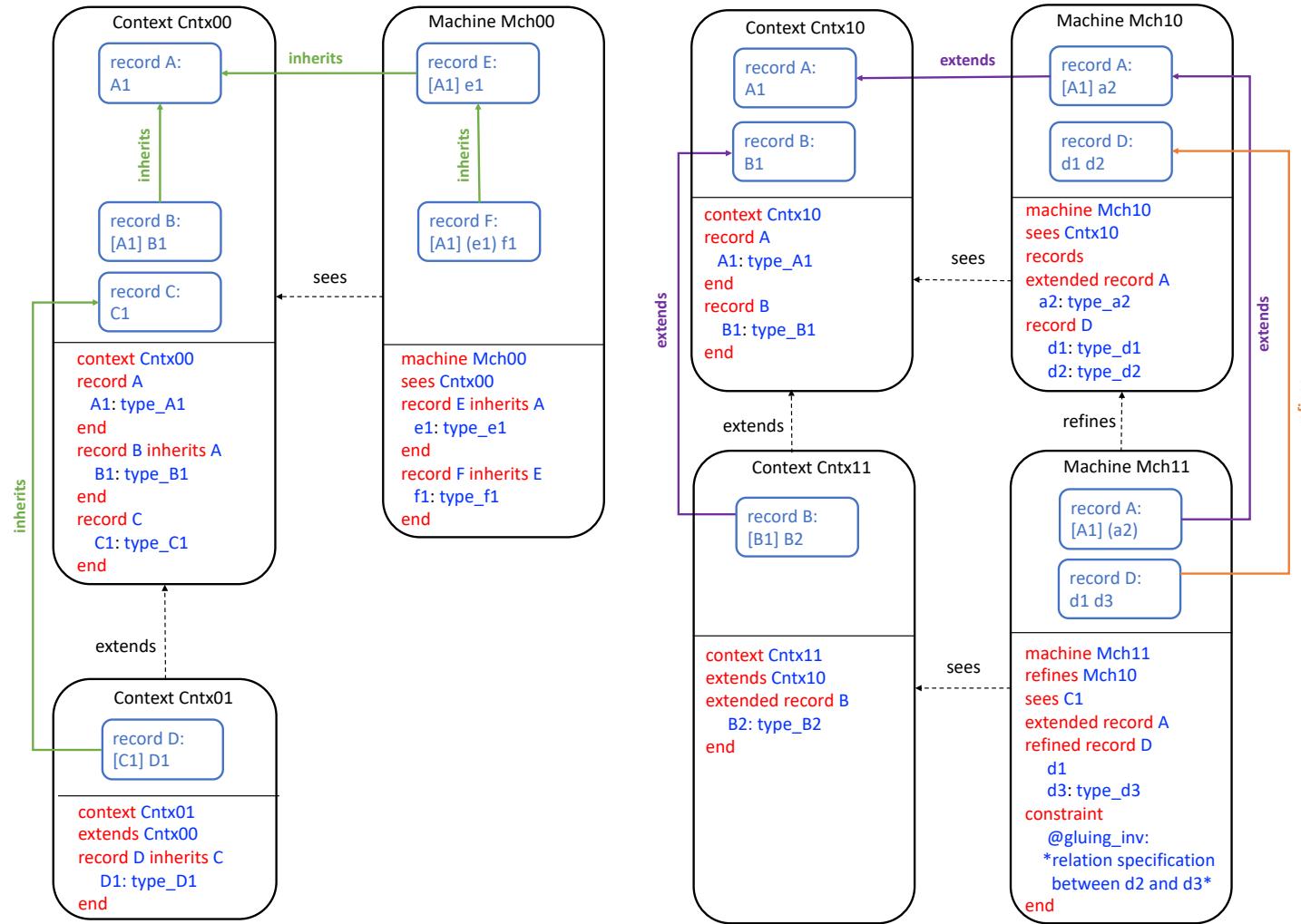
The CamilleX Framework



Records – an important extension to Event-B

- Inherits
 - Subtyping records
 - Implicit fields
- Refines
 - Replacing fields
 - Adding fields
- Extends
 - (Only) Adding fields
- Constraints
 - Properties of record instances

Ongoing research :
decomposition of records to prepare for translation to SPARK records



Background - SPARK

- Subset of Ada programming language
- Assertions, pre-conditions, post-conditions
- Targeted at highly reliable software
- Formal verification to prove the absence of runtime errors:
 - arithmetic overflow, buffer overflow and division-by-zero.
- Applied over many years
 - e.g. aircraft systems, control systems and rail systems.

Motivation

- Abstraction to isolate important properties
- Refinement to add detail and design
- Resulting in Validated and Verified models
- ...But how can we ensure the code complies with the models
- Answer: generate Spark outline code
 - With pre/post-conditions that match events of the model
 - Assertions for run time checking

Steps : From abstract concept to Spark Implementation

- Abstract model of **concept**
- Refinements that introduce more **detailed requirements**
- Refinements that introduce **design decisions**
- Decomposition ... ***Controller*** + Environment
- Refine ***Controller*** to prepare for code generation
- Generate **SPARK** from ***Controller***

Overview of Event-B to SPARK Translation Rules

- Component Translation

- (All) Context → specification package
 - Context **extends** → context packages **use** (and all extended context packages)
- (Last) Refined Machine → specification and body packages
 - Machine **sees** context → **use** contexts packages (and all extended context packages)

machine mch_name **sees** ctx_name

```
with ctx_name; use ctx_name;  
package mach_name  
with SPARK_Mode=> On  
end mach_name;
```

```
package body mch_name with SPARK_Mode is ...  
end mach_name;
```

Overview of Event-B to SPARK Translation Rules

- Constant Translation
 - Non function constants → constant, type depends on the axiom definitions

```
constants const_name
axioms
const_name ∈ const_type
const_name = const_value
```

```
const_name : constant const_type := const_val;
```

- Function type constant → function with return type depending on the range of the function and the function parameters are the domain of the Event-B function.

```
constants cnst_name
axioms cnst_name ∈ dom → ran
```

```
function cnst_name (p_dom : in dom) return ran;
```

Overview of Event-B to SPARK Translation Rules

- Variable Translation
 - Variable → Global variable,
 - initialised according to the INITIALISATION event actions
- Record Translation
 - CamilleX Record → SPARK record
 - with all Event-B record fields (direct and implicit)

```
record rec_name
  field_name : field_type
```

```
type rec_name is
  field_name : field_type ;
end record;
```

Overview of Event-B to SPARK Translation Rules

- Event Translation
 - Event → Procedure *(except Initialisation)*
 - Event Guard → Pre-condition
 - Event Action → Post-condition & Procedure body
 - Event Parameter → Procedure Parameter
 - (where output/input/in_out is deduced from guards and actions)
- Note that we have already proved invariants in the Event-B.. No need to translate invariants to SPARK
 - *Or is there! Some industrial partners have suggested that it may still be useful.. E.g. to catch problems caused by interrupts.*

Refinement of SBB example

From abstract concept to spark

- Ballot
 - Paper (voter,vote)
 - Paper (voter,vote,time)
 - Paper (~~voter,vote~~,time,encrypted)
 - Paper (~~voter,vote~~,time,encrypted, mac)
 - Decompose -> smart ballot box + voters/attackers
 - Refine SBB data towards arrays etc.
 - Generate SPARK for SBB

Refinement of SBB example

From abstract concept to spark

- Ballot
- Paper (**voter,vote**)
 - Replace ballot with its physical representation :
Paper - fields for voter and vote
- Paper (voter,vote,time)
 - This introduces the possibility of invalid papers.. Copying faking etc.
- Paper (~~voter,vote~~,time,encrypted)
- Paper (~~voter,vote~~,time,encrypted, mac)
- Decompose -> smart ballot box + voters/attackers
- Refine SBB data towards arrays etc.
- Generate SPARK for SBB

Refinement of SBB example

From abstract concept to spark

- Ballot
- Paper (voter,vote)
- Paper (voter,vote,**time**)
 - Introduce new field : time
 - Voting papers can expire,
 - Reduces opportunity for validity threats
- Paper (~~voter,vote~~,time,encrypted)
- Paper (~~voter,vote~~,time,encrypted, mac)
- Decompose -> smart ballot box + voters/attackers
- Refine SBB data towards arrays etc.
- Generate SPARK for SBB

Refinement of SBB example

From abstract concept to spark

- Ballot Refine voter,vote with encrypted.
- Paper (voter,vote) Provide confidentiality
- Paper (voter,vote,time)
- Paper (~~voter,vote~~,time,**encrypted**)
- Paper (~~voter,vote~~,time,encrypted, mac)
- Decompose -> smart ballot box + voters/attackers
- Refine SBB data towards arrays etc.
- Generate SPARK for SBB

Refinement of SBB example

From abstract concept to spark

- Ballot
 - Paper (voter,vote)
 - Paper (voter,vote,time)
 - Paper (~~voter,vote~~,time,encrypted)
 - Paper (~~voter,vote~~,time,encrypted, **mac**)
 - Decompose -> smart ballot box + voters/attackers
 - Refine SBB data towards arrays etc.
 - Generate SPARK for SBB
- Introduce mac (algorithm for hashing)
Enables checking validity of vote..
E.g. if an attacker tries to alter the vote

Refinement of SBB example

From abstract concept to spark

- Ballot
 - Paper (voter,vote)
 - Paper (voter,vote,time)
 - Paper (~~voter,vote~~,time,encrypted)
 - Paper (~~voter,vote~~,time,encrypted, mac)
 - Decompose -> smart ballot box + voters/attackers
 - Refine SBB data towards arrays etc.
 - Generate SPARK for SBB
- Event-B model is a closed system
 - Some parts of model are the controller
 - Others the environment being controlled

*Future work - How to decompose records sets e.g.
only cast_papers are in the SBB system*

Refinement of SBB example

From abstract concept to spark

- Ballot
- Paper (voter,vote)
- Paper (voter,vote,time)
- Paper (~~voter,vote~~,time,encrypted)
- Paper (~~voter,vote~~,time,encrypted, mac)
- Decompose -> smart ballot box + voters/attackers
- Refine SBB data towards arrays etc.
- Generate SPARK for SBB
 - *Data Refinement from abstract SET into Array*
 - *Array can be modelled as a **Total function** from 0..n to set*
 - *Event-B records can have optional fields SPARK we can only use total functions -*
 - *Define a **null value** for optional field so that all records are total*

Refinement of SBB example

From abstract concept to spark

- Ballot
- Paper (voter,vote)
- Paper (voter,vote,time)
- Paper (~~voter,vote~~,time,encrypted)
- Paper (~~voter,vote~~,time,encrypted, mac)
- Decompose -> smart ballot box + voters/attackers
- Refine SBB data towards arrays etc.
- Generate SPARK for SBB

Example: Application to Smart Ballot Box Model

```

event cast_paper
refines cast_paper
any
paper
where
@grd1: paper  $\in$  BARCODE
@grd2: cast_count  $\in$  0 .. max_votes -1
.....
then
@act1: cast_arr(cast_count) := paper
@act2: cast_count := cast_count + 1
end
  
```

```

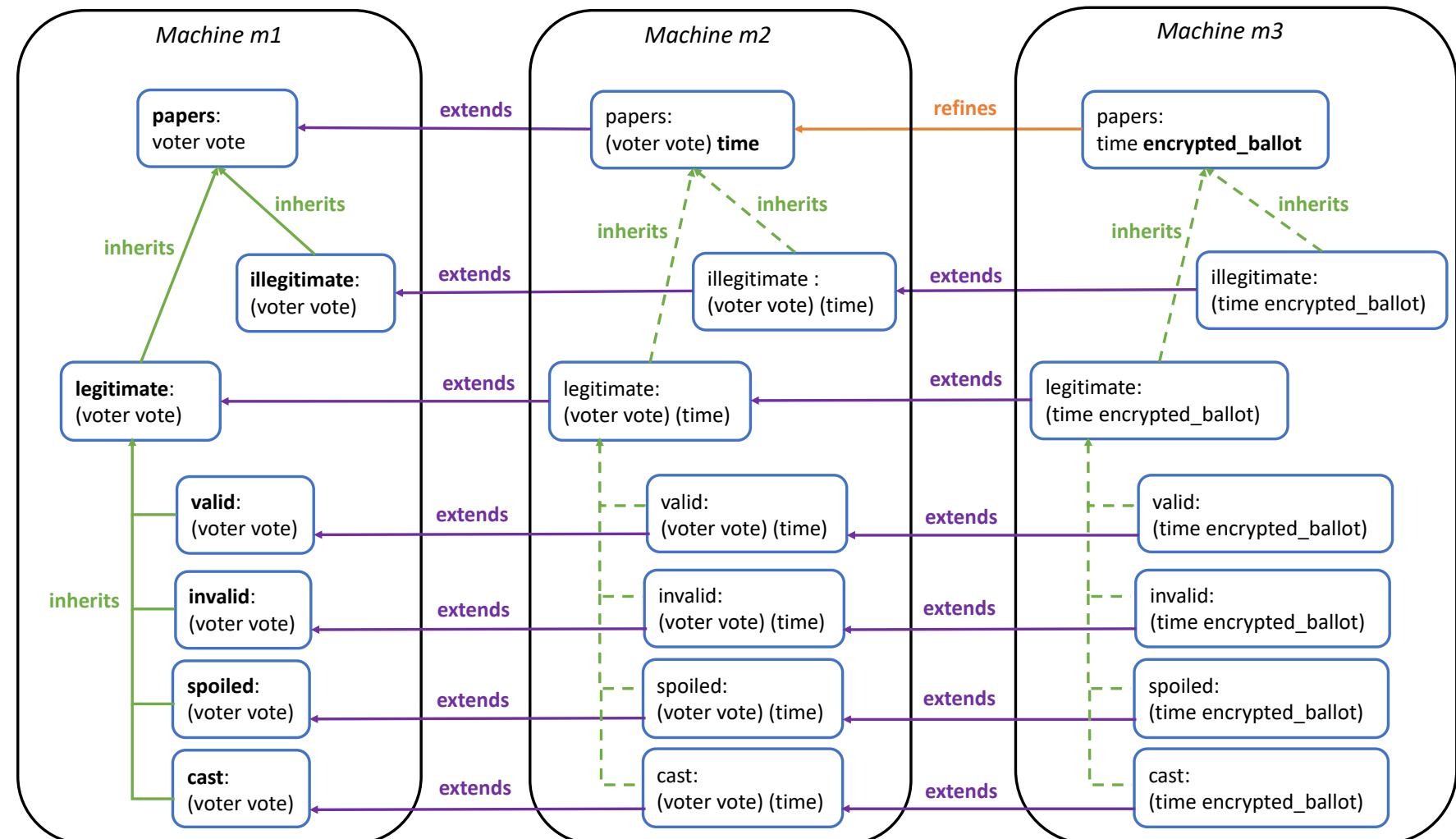
procedure cast(paper : in barcode) with
Global => (Proof_In => (spoiled_arr, curr_time, spoil_count),
In_Out => (cast_arr, cast_count)),
Pre => cast_count in 0 .. Max_Votes-1,
and then not already_cast(paper)
...
Post => already_cast(paper)
and then cast_count = cast_count' old + 1);
  
```

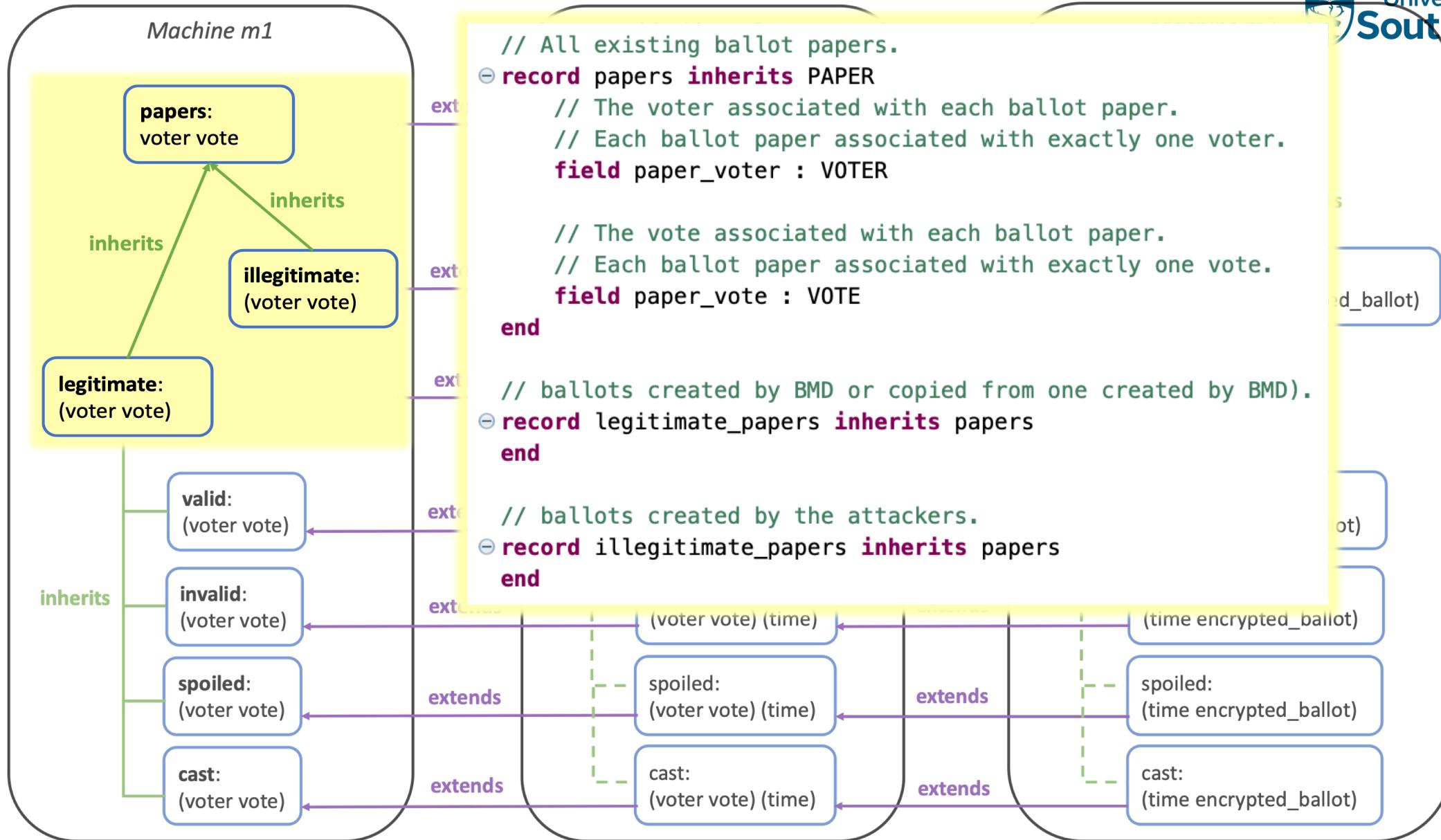
```

procedure cast(paper : in barcode) is
begin
  cast_arr(cast_count) := paper;
  cast_count := cast_count + 1;
end cast;
  
```

QUESTIONS?

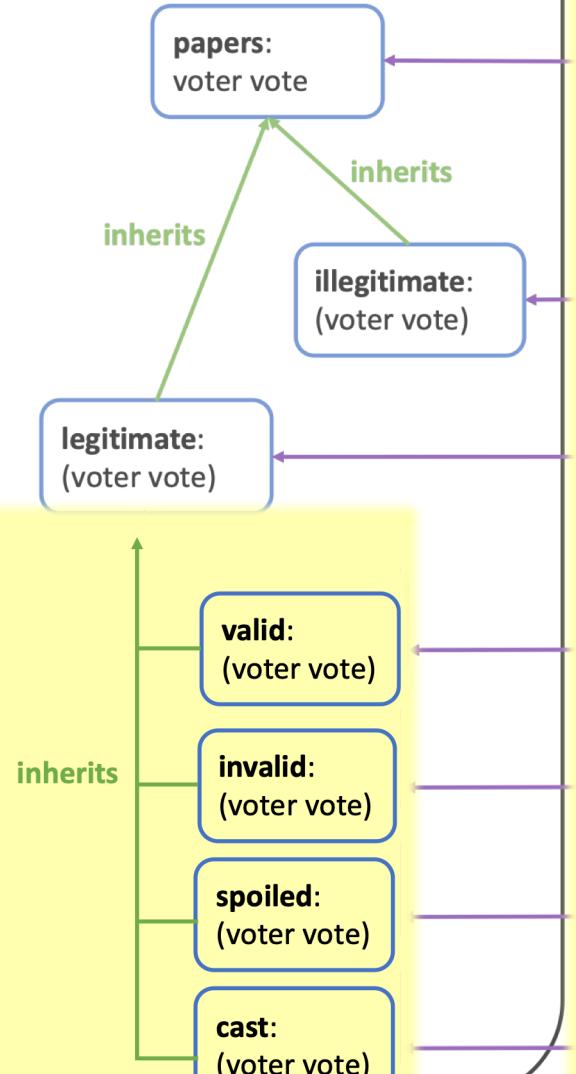
Case Study – SBB Electronic Voting



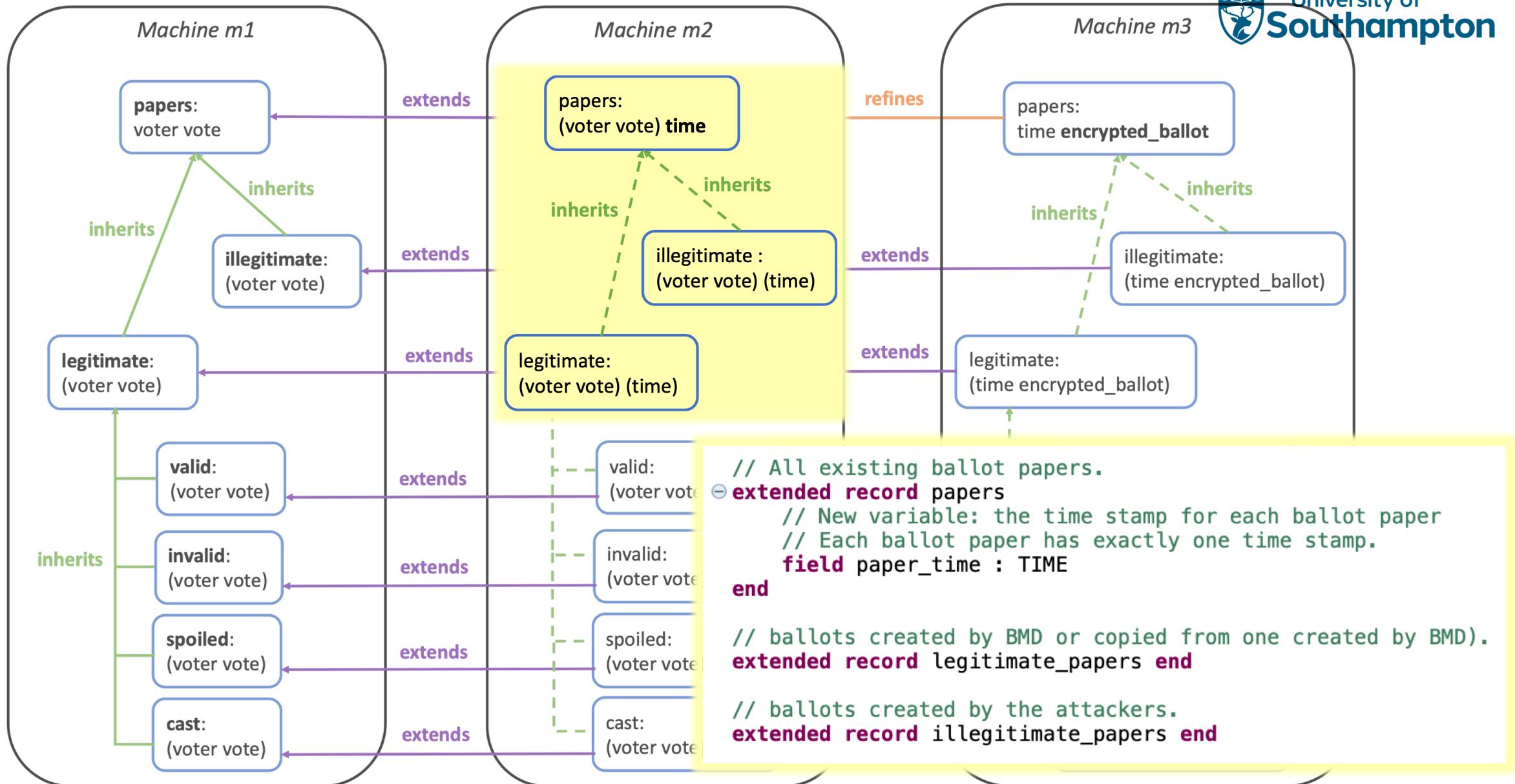


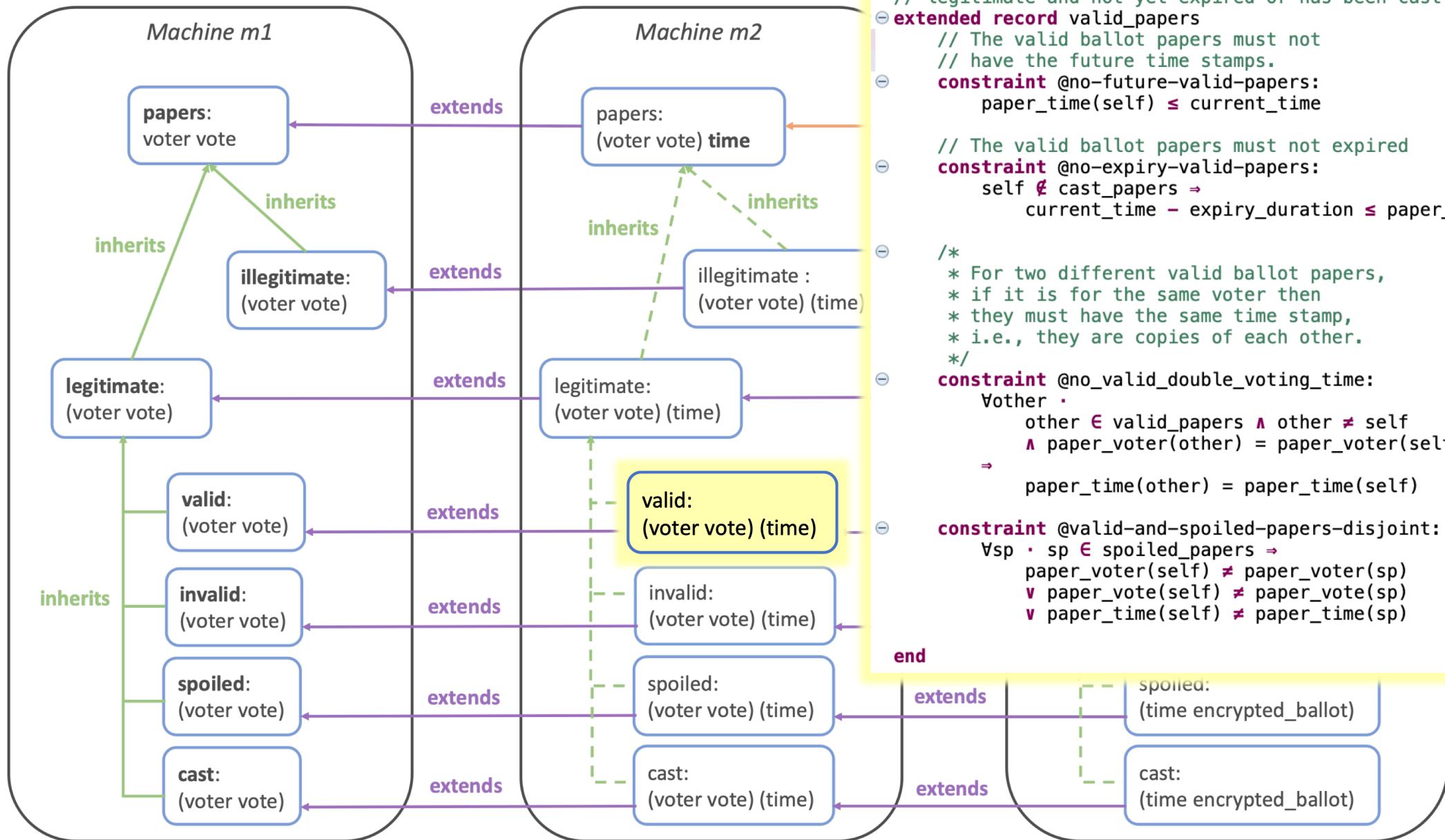


Machine m1



```
// The set of spoiled ballot paper.  
record spoiled_papers inherits legitimate_papers  
end  
  
// legitimate and not yet expired or has been cast before expired or not spoiled  
record valid_papers inherits legitimate_papers  
/*  
 * For two different valid ballot papers, if it is for the same voter then  
 * they must have the same vote, i.e., they are copies of each other.  
 */  
constraint @no_valid_double_voting_vote:  
    @other .  
        other ∈ valid_papers ∧ other ≠ self  
        ∧ paper_voter(other) = paper_voter(self)  
    =>  
        paper_vote(other) = paper_vote(self)  
end  
  
// legitimate but becomes invalid since expired or a copy has been cast or  
// spoiled.  
record invalid_papers inherits legitimate_papers  
end  
  
// legitimate and already cast  
record cast_papers inherits legitimate_papers  
/*  
 * If a voter already cast a ballot paper, they cannot have any valid  
 * ballot paper.  
 */  
constraint @no_cast_double_voting_vote:  
    paper_voter(self) ≠ paper_voter[valid_papers]  
end
```





```

// legitimate and not yet expired or has been cast before expired.
@extended_record valid_papers
  // The valid ballot papers must not
  // have the future time stamps.
  constraint @no-future-valid-papers:
    paper_time(self) ≤ current_time

  // The valid ballot papers must not expired
  constraint @no-expiry-valid-papers:
    self ∉ cast_papers ⇒
      current_time - expiry_duration ≤ paper_time(self)

/*
 * For two different valid ballot papers,
 * if it is for the same voter then
 * they must have the same time stamp,
 * i.e., they are copies of each other.
 */
constraint @no_valid_double_voting_time:
  ∀other .
  other ∈ valid_papers ∧ other ≠ self
  ∧ paper_voter(other) = paper_voter(self)
  ⇒
  paper_time(other) = paper_time(self)

constraint @valid-and-spoiled-papers-disjoint:
  ∀sp · sp ∈ spoiled_papers ⇒
  paper_voter(self) ≠ paper_voter(sp)
  ∨ paper_vote(self) ≠ paper_vote(sp)
  ∨ paper_time(self) ≠ paper_time(sp)
end

```

