

Counting number of distinct real roots of Pham system (Assignment 1)

Mountain Chan / B1-2

October 28, 2019

1 Problem

Definition 1 (Pham system) A system $f \in \mathbb{C}[x_1, \dots, x_n]^n$ is called a Pham system if it has the following form:

$$\begin{aligned} f_1 &= x_1^{d_1} + \sum_{e_1 + \dots + e_n < d_1} a_{1,e} x_1^{e_1} \dots x_n^{e_n} \\ &\vdots \\ f_n &= x_1^{d_n} + \sum_{e_1 + \dots + e_n < d_n} a_{n,e} x_1^{e_1} \dots x_n^{e_n} \end{aligned}$$

Remark 2 The pham system has exactly $m = d_1 \dots d_n$ many complex roots (counting multiplicities). Thus we will study the number of real roots of a real pham system.

Problem 3 Devise an algorithm with the following specification:

In: $f \in \mathbb{R}[x_1, \dots, x_n]^n$, Pham system

Out: N , the number of distinct real roots of f

2 Theory

Definition 4 (Discriminant matrix) Let $f \in \mathbb{C}[x_1, \dots, x_n]^n$ be a Pham system of degrees d_1, \dots, d_n . Let $m = d_1 \dots d_n$. Let $z_1, \dots, z_m \in \mathbb{C}^n$ be the complex roots of f . The discriminant matrix Q of f is defined by

$$Q = VV^t$$

where

$$V = \begin{bmatrix} \omega_1(z_1) & \dots & \omega_1(z_m) \\ \vdots & & \vdots \\ \omega_m(z_1) & \dots & \omega_m(z_m) \end{bmatrix}$$

where

$$\omega = \{x_1^{e_1} \dots x_n^{e_n} : e_1 < d_1, \dots, e_n < d_n\}$$

Theorem 5 (Multivariate Hermite) # of distinct real roots of $f = \sigma(Q)$

Proof.

$$\begin{aligned}
Q_{ij} &= \sum_{k=1}^m \omega_i(z_k) \omega_j(z_k) \\
Q_{ij} &= \sum_{p=1}^s \omega(\alpha_p) \mu_p \omega_j(\alpha_p) + \sum_{q=1}^t v_q (\omega_i(\beta_q) \omega_j(\beta_q)) + \sum_{q=1}^t v_q \overline{\omega_i(\beta_q)} \overline{\omega_j(\beta_q)} \\
Q_{ij} &= \sum_{p=1}^s \omega_i(\alpha_p) \mu_p \omega_j(\alpha_p) + \sum_{q=1}^t v_q (\omega_i(\beta_q) \omega_j(\beta_q) \overline{\omega_i(\beta_q)} \overline{\omega_j(\beta_q)}) \\
Q_{ij} &= \sum_{p=1}^s \omega_i(\alpha_p) \mu_p \omega_j(\alpha_p) + \sum_{q=1}^t v_q (\omega_i(\beta_q) \omega_j(\beta_q) \overline{\omega_i(\beta_q)} \overline{\omega_j(\beta_q)}) \\
Q_{ij} &= \sum_{p=1}^s \omega_i(\alpha_p) \mu_p \omega_j(\alpha_p) + \sum_{q=1}^t 2v_q ((\operatorname{Re}(\omega_i(\beta_q)) \operatorname{Re}(\omega_j(\beta_q)) - \operatorname{Im}(\omega_i(\beta_q)) \operatorname{Im}(\omega_j(\beta_q))) \\
Q_{ij} &= \sum_{p=1}^s \omega_i(\alpha_p) \mu_p \omega_j(\alpha_p) + \sum_{q=1}^t \operatorname{Re}(\omega_i(\beta_q)) (2v_q \operatorname{Re}(\omega_j(\beta_q)) + \sum_{q=1}^t \operatorname{Im}(\omega_i(\beta_q)) (-2v_q) \operatorname{Im}(\omega_j(\beta_q))
\end{aligned} \tag{1}$$

$$D = \begin{bmatrix} \mu_1 & & & & & & & & & \\ & \ddots & & & & & & & & \\ & & \mu_s & & & & & & & \\ & & & 2v_1 & & & & & & \\ & & & & \ddots & & & & & \\ & & & & & 2v_t & & & & \\ & & & & & & -2v_1 & & & \\ & & & & & & & \ddots & & \\ & & & & & & & & -2v_t & \end{bmatrix}$$

$$T = \begin{bmatrix} \omega_0(\alpha_1) & \cdots & \omega_0(\alpha_s) & \operatorname{Re}(\omega_0(\beta_1)) & \cdots & \operatorname{Re}(\omega_0(\beta_t)) & \operatorname{Im}(\omega_0(\beta_1)) & \cdots & \operatorname{Im}(\omega_0(\beta_t)) \\ \vdots & \ddots & \vdots & \vdots & \ddots & \vdots & \vdots & \ddots & \vdots \\ \omega_m(\alpha_1) & \cdots & \omega_m(\alpha_s) & \operatorname{Re}(\omega_m(\beta_1)) & \cdots & \operatorname{Re}(\omega_m(\beta_t)) & \operatorname{Im}(\omega_m(\beta_1)) & \cdots & \operatorname{Im}(\omega_m(\beta_t)) \end{bmatrix}$$

$\sigma(D) = (s + t) - t = s = \text{number of distinct real roots}$

By Sly. Law of Inertia, $\sigma(Q) = \sigma(D)$, hence number of distinct real roots of $f = \sigma(D) = s$ ■

Remark 6 *Hermite's theorem is "useless" because*

1. *Computation of Q requires computing the roots of f .*
2. *Computation of $\sigma(Q)$ requires computing the roots of the characteristic polynomial of Q .*

Hence, to make Hermite's theorem useful, we need to find ways to

1. *Compute Q **without** computing the roots of f .*
2. *Compute $\sigma(Q)$ **without** computing the roots of the characteristic polynomial of Q .*

In the following, we will tackle the challenges one by one.

Definition 7 (Mutipilcation matrix) Let $g \in \mathbb{C}[x_1, \dots, x_n]$. Then the multiplication matrix M_g for g is defined by

$$g \begin{bmatrix} \omega_1 \\ \vdots \\ \omega_m \end{bmatrix} \equiv_f M_g \begin{bmatrix} \omega_1 \\ \vdots \\ \omega_m \end{bmatrix}$$

Theorem 8 We have $Q_{ij} = \text{tr} M_{\omega_i \omega_j}$.

Proof.

1. Let $f(z) = 0$

$$g(z) \begin{bmatrix} \omega_1(z) \\ \vdots \\ \omega_m(z) \end{bmatrix} = M_g \begin{bmatrix} \omega_1(z) \\ \vdots \\ \omega_m(z) \end{bmatrix}$$

$$Q_{ij} = \sum_{k=1}^m \omega_i(z_k) \omega_j(z_k)$$

$$2. \quad Q_{ij} = \sum_{k=1}^m (\omega_i \omega_j)(z_k)$$

$$Q_{ij} = \sum_{k=1}^m g(z_k)$$

$$Q_{ij} = \text{tr} M_{\omega_i \omega_j}$$

■

Theorem 9 We have

$$Q_{ij} = \sum_{\mu\nu} M_{i\mu\nu} M_{j\nu\mu}$$

where

$$\omega_i \omega_j \equiv_f \sum_k M_{ijk} \omega_k$$

Proof.

$$\begin{aligned} Q_{ij} &= \sum_{\mu} (M_{\omega_i \omega_j})_{\mu\mu} \\ &= \sum_{\mu} (M_{\omega_i} M_{\omega_j})_{\mu\mu} \end{aligned} \tag{2}$$

■

3 Algorithm

Algorithm 10 (DiscriminantMatrix)

In: $f \in \mathbb{R}[x_1, \dots, x_n]^n$, Pham system

Out: Q , the discriminant matrix of f

1. ω = Monomial basis of function f
2. p = dimension of matrix

3. $Q_{ij} = \text{trace of matrix } M$
4. $Q_{ij} = \text{tr} M_{\omega_i \omega_j}$
5. Construct Q with dimension $m \times m$ from Q_{ij}
6. Return Q

Remark 11 Recall that Descartes' theorem is exact when all the roots are real.

Algorithm 12 (Signature)

In: $M \in \mathbb{R}^{m \times m}$, symmetric

Out: $S = \sigma(M)$

1. $C = |\lambda I - M|$, the characteristic polynomial of Q
2. $m_+ =$ the sign variation count of the coefficients of C
3. $m_0 =$ the least exponent in C
4. $m_- = m - m_+ - m_0$
5. $S = m_+ - m_-$
6. Return S

Algorithm 13 (NumberDistinctRealRoots)

In: $f \in \mathbb{R}[x_1, \dots, x_n]^n$, Pham system

Out: N , the number of distinct real roots of f

1. $Q = \text{DiscriminantMatrix}(f)$
2. $N = \text{Signature}(Q)$
3. Return N