



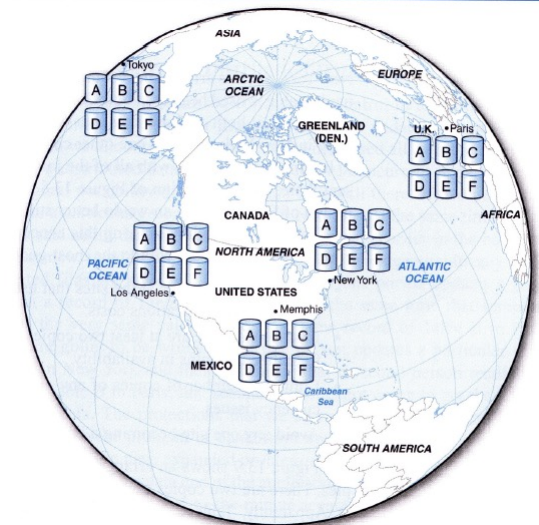
INFO90002 Database Systems & Information Modelling

Lecture 16: Distributed Databases

- What is a distributed database?
- Why are they used, and how they work
- Pros and cons of different approaches
 - material in this lecture is drawn from Hoffer et al. (2013) *Modern Database Management* 11th edition, chapter 12, available online at http://wps.prenhall.com/bp_hoffer_mdm_11/230/58943/15089539.cw/index.html
 - pictures on this page are from Gillenson (2005) *Fundamentals of Database Management Systems*



distributed database

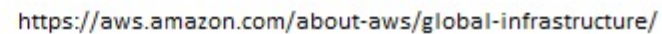


replicated database

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- **Distributed Database**
 - a single logical database physically spread across multiple computers in multiple locations that are connected by a data communications link
 - *appears to users as though it is one database*
- **Decentralized Database**
 - a collection of independent databases which are not networked together as one logical database
 - appears to users as though many databases
- We are concerned with *distributed* databases

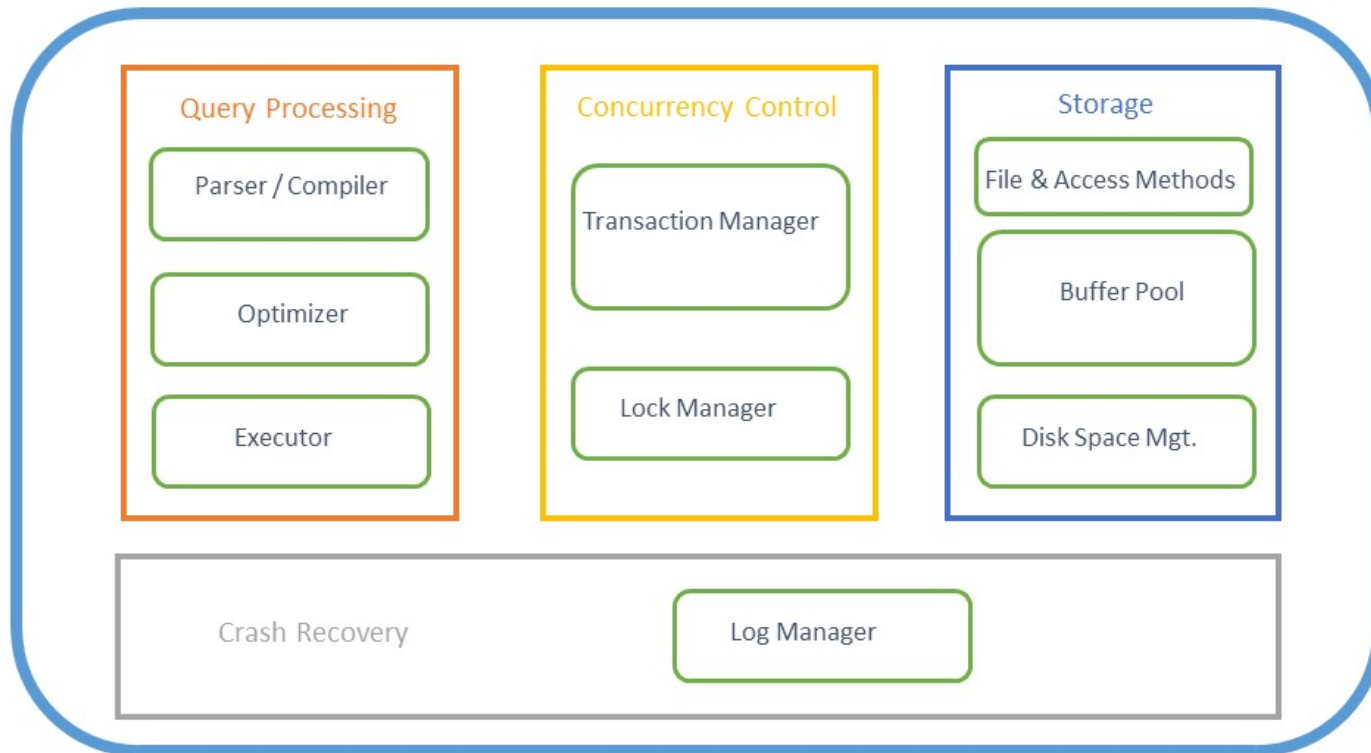
Global Infrastructure



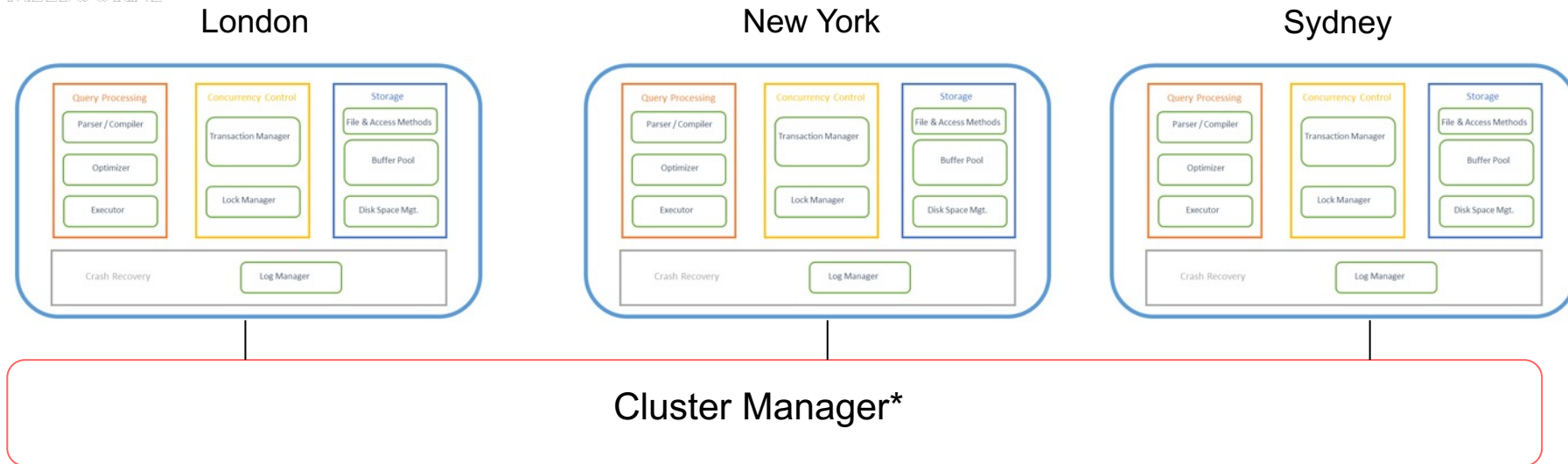
<https://aws.amazon.com/about-aws/global-infrastructure/>

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“zones” available within that data centre providing redundancy and availability



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- Each Physical Server has one of these memory structures
- Often accessing their own and shared physical storage between all physical servers
- Cluster Manager coordinates communication between physical servers

* May be called something else in different vendor databases

- Good fit for geographically distributed organizations / users
 - Utilize the internet
- Data located near site with greatest demand
 - ESPN Weekend Sports Scores



AFL - Melbourne



EPL - London

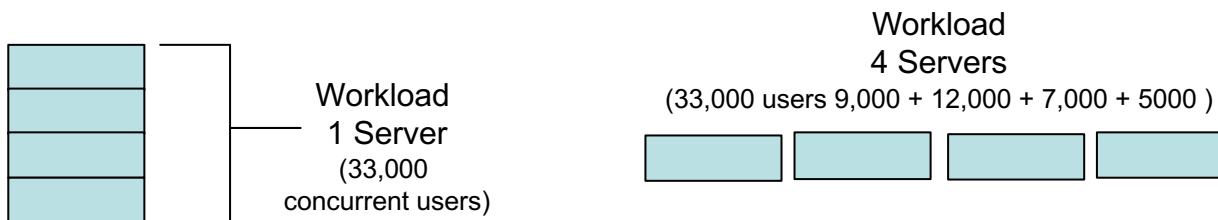


NFL - New York



Hurling - Dublin

- Faster data access (to local data)
- Faster data processing
 - workload is shared between each physical server



- Allows modular growth

- add new servers as load increases



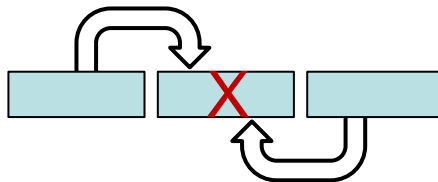
- Increased reliability and availability

- less danger of a single-point of failure (SPOF)



- Supports database recovery

- data is replicated across multiple sites



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- Complexity of management and control
 - database or/and application must stitch together data across sites
 - Who and where is the current version of the record (row & column)?
 - Who is waiting to update that information and where are they?
 - How does the logic display this to the web & application server?
- Data integrity
 - additional exposure to improper updating
 - If two users in two locations update the record at the exact same time who decides which statement should “win”?
 - Solution: Transaction Manager or Master-slave design
- Security
 - many server sites -> higher chance of breach
 - Multiple access sites require protection including network and storage infrastructure from both cyber & physical attacks

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- Lack of standards
 - different Relational DDBMS vendors use different protocols
- Increased training & maintenance costs
 - more complex IT infrastructure
 - Increased Disk storage (\$)
 - Fast intra and inter network infrastructure (\$\$\$)
 - Clustering software (\$\$\$\$)
 - Network Speed (\$\$\$\$\$)
- Increased storage requirements
 - Replication model

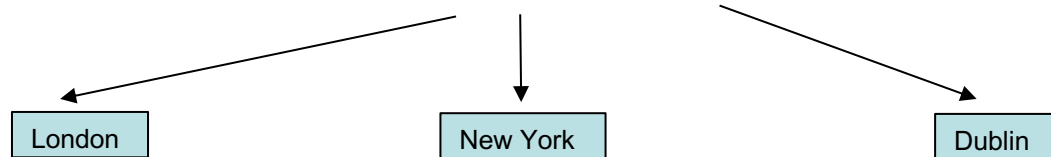
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- Location transparency
 - a user does not need to know where particular data are stored
- Local autonomy
 - a node can continue to function for local users if connectivity to the network is lost
- Trade-offs
 - Availability vs Consistency
 - Synchronous vs Asynchronous updates

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- A user (or program) accessing data do not need to know the location of the data in the network of DBMS's
- Requests to retrieve or update data from any site are automatically forwarded by the system to the site or sites related to the processing request
- All data in the network appears as a single logical database stored at one site to the users
- A single query can join data from tables in multiple sites

```
SELECT hometeam, homescore, awayteam, awayscore  
FROM results INNER JOIN codes  
ON results.codeid = codes.codeid  
WHERE sportscore in ('NFL', 'Hurling', 'EPL');
```



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- Users can administer their local database
 - control local data (e.g Hurling results)
 - administer security
 - log transactions
 - recover when local failures occur
 - provide full access to local data
- Being able to operate locally when connections to other databases fail



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Distributed Options

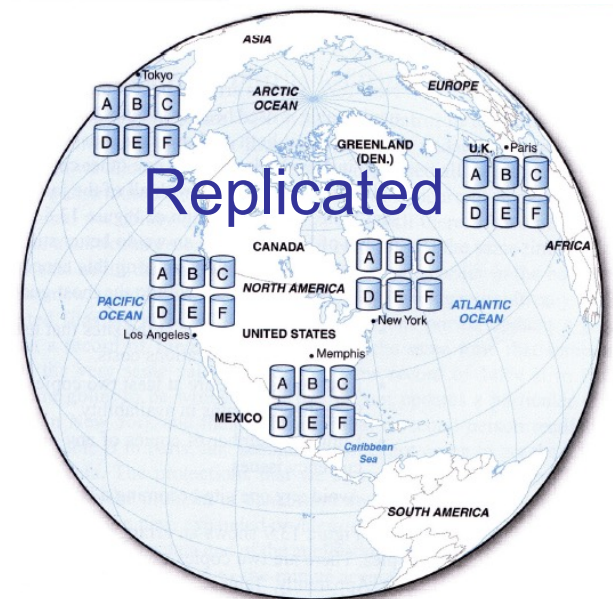
Data Replication
Horizontal Partitioning
Vertical Partitioning

- Data replication
 - Data copied across sites
- Horizontal partitioning
 - Table rows distributed across sites
- Vertical partitioning
 - Table columns distributed across sites
- Combinations of the above

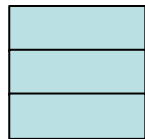


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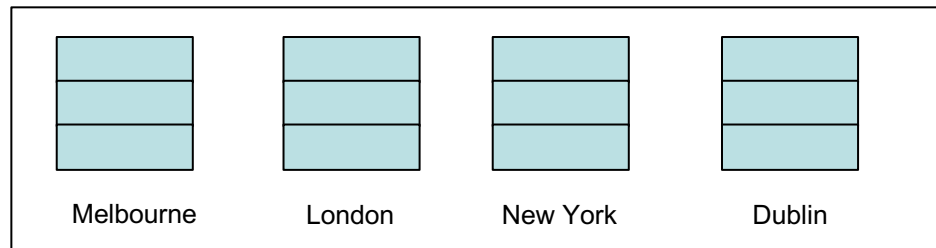
- High reliability due to redundant copies of data
- Fast access to data at the location where it is most accessed
- May avoid complicated distributed integrity routines
 - replicated data is refreshed at scheduled intervals
- Decoupled nodes don't affect data availability
 - transactions proceed even if some nodes are down
- Reduced network traffic at prime time
 - if updates can be delayed
- This is currently popular as a way of achieving high availability for global systems.
 - Most SQL & NoSQL databases offer replication



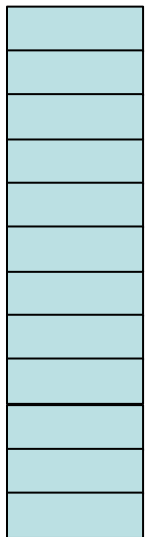
- Need more storage space
 - Each server stores a copy of the row
- Data Integrity:
 - high tolerance for out-of-date data may be required
 - updates may cause performance problems for busy nodes
 - retrieve incorrect data if updates have not arrived



Centralised Database
One database in one server
(1 copy of data)



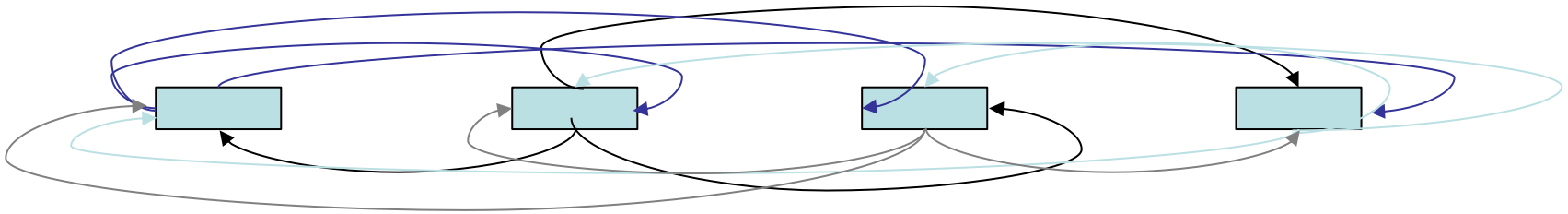
Distributed Database
One database in 4 physical servers
(4 copies of data)



Data Size

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- Takes time for update operations
 - high tolerance for out-of-date data may be required
 - updates may cause performance problems for busy nodes



- Network communication capabilities
 - updates can place heavy demand on telecommunications/networks
 - Cost high speed networks are expensive (\$\$\$\$\$)

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- Data is continuously kept up to date
 - users anywhere can access data and get the same answer.
- If any copy of a data item is updated anywhere on the network, the same update is immediately applied to all other copies or it is aborted.
- Ensures data integrity and minimizes the complexity of knowing where the most recent copy of data is located.
- Can result in slow response time and high network usage
 - the DDBMS spends time checking that an update is accurately and completely propagated across the network.
 - The committed updated record must be identical in all servers

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- Some delay in propagating data updates to remote databases
 - some degree of at least temporary inconsistency is tolerated
 - may be ok it is temporary and well managed
- Acceptable response time
 - updates happen locally and data replicas are synchronized in batches and predetermined intervals
- May be more complex to plan and design
 - need to ensure the right level of data integrity and consistency
- Suits some information systems more than others
 - compare commerce/finance systems with social media

- Different rows of a table at different sites
- Advantages
 - data stored close to where it is used
 - efficiency
 - local access optimization
 - better performance
 - only relevant data is stored locally
 - security
 - unions across partitions
 - ease of query
- Disadvantages
 - accessing data across partitions
 - inconsistent access speed
 - no data replication
 - backup vulnerability (SPOF)

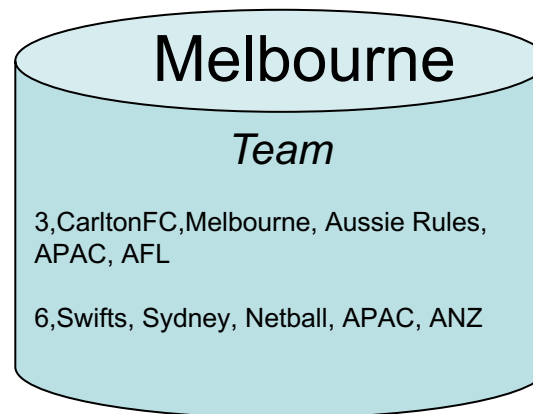
Team table

ID	Team	City	Code	Region	League
1	Arsenal	London	Football	Europe	EPL
2	Jets	NYC	Grid Iron	Americas	NFL
3	Carlton FC	Melbourne	Aussie Rules	APAC	AFL
4	Racing92	Paris	Rugby	Europe	Top14
5	Yankees	NYC	Baseball	Americas	MLB
6	Swifts	Sydney	Netball	APAC	ANZ

Example horizontal partitioning

ID	Team	City	Code	Region	League
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Horizontal Partitioning based on Region






- Different columns of a table at different sites
- Advantages and disadvantages are the same as for horizontal partitioning
 - except
 - combining data across partitions is more difficult because it requires joins (instead of unions)

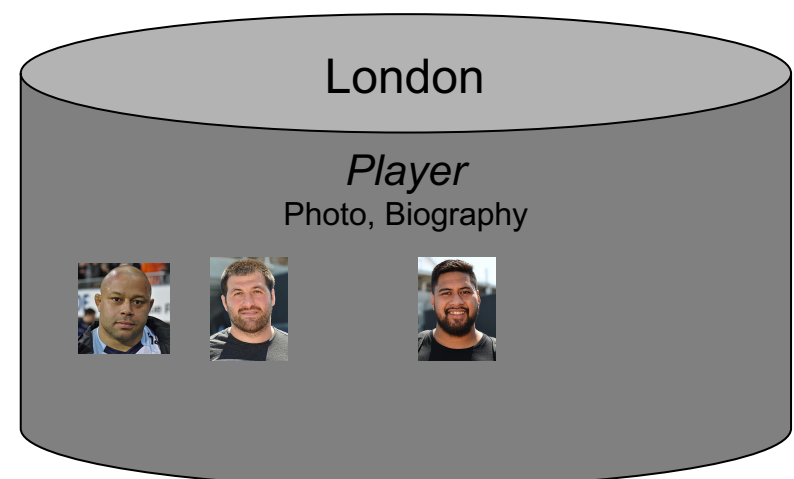
Player table

ID	Firstname	Lastname	Team	League	Photo	Biography
110	Luc	Ducalon	4	Top14		Ipso locum
120	Vasil	Kakokan	4	Top14		Ipso locum est
130	Donacca	Ryan	4	Top14	<null>	
210	Edwin	Maka	4	Top14		

Example vertical partitioning

ID	Firstname	Lastname	Team	League	Photo	Biography
110	Luc	Ducalon	4	Top14		Ipso locum
120	Vasil	Kakokan	4	Top14		Ipso locum est
130	Donacca	Ryan	4	Top14	<null>	
210	Edwin	Maka	4	Top14		

Vertical Partitioning based on column requirements



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- **Centralized database, distributed access**
 - DB is at one location, and accessed from everywhere
- **Replication with periodic (asynchronous) snapshot update**
 - many locations, each data copy updated periodically
- **Replication with near real-time synchronization of updates**
 - many locations, each data copy updated in near real time
- **Partitioned, integrated, one logical database**
 - data partitioned across at many sites, within a logical database, and a single DBMS
- **Partitioned, independent, nonintegrated segments**
 - data partitioned across many sites.
 - independent, non-integrated segments
 - multiple DBMS, multiple computers

- Replication with periodic (asynchronous) snapshot update
 - many locations, each data copy updated periodically
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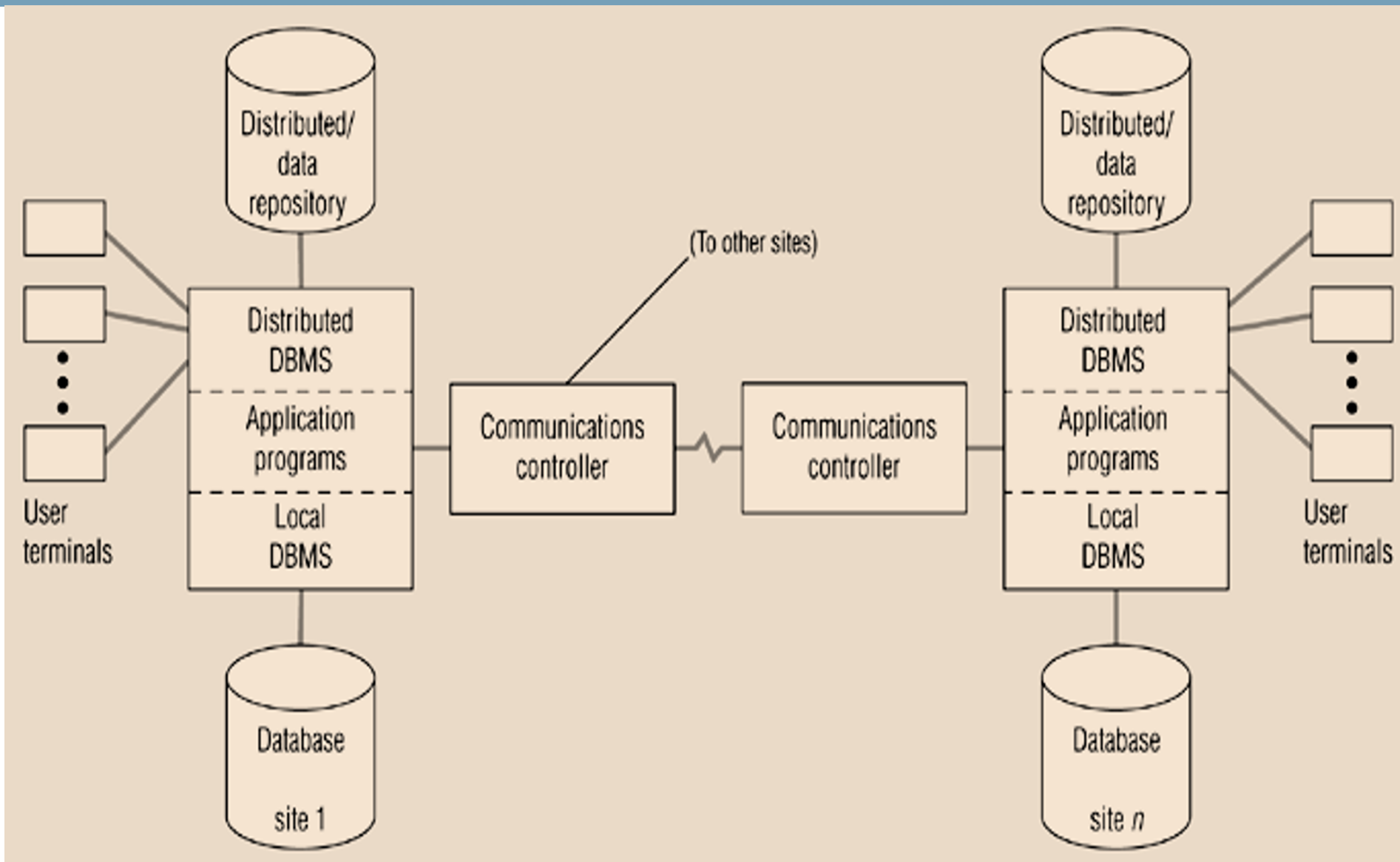
Comparing Configurations

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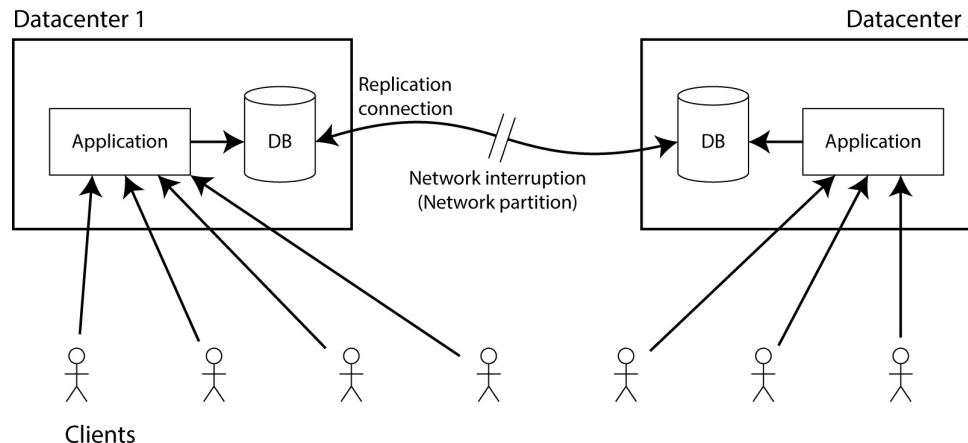
	Reliability	Expandability	Communication Overhead	Management	Data Consistency
Centralised	POOR Depends on central server.	POOR Single Server is limited by memory & storage maximums.	VERY HIGH Traffic heads to one centralised location.	EXCELLENT One very large site is easier to manage.	EXCELLENT All users always see the same data.
Replicated with Snapshots	GOOD Redundancy and tolerated delays in data synch.	VERY GOOD Cheap to scale up with new servers.	LOW to MEDIUM Intermittent bursts of network traffic (but not constant flooding of network).	VERY GOOD Each copy is alike.	MEDIUM Update delays are tolerable with snapshot catch ups for data consistency.
Synchronised Replication	EXCELLENT Redundancy and minimal delays.	VERY GOOD Low cost and only linear growth in synchronisation.	MEDIUM Constant messages to maintain synchronisation.	MEDIUM Data collisions need to be resolved and need good design and management.	VERY GOOD Close to precise consistency.
Integrated Partitions	GOOD Effective use of partitioning and redundancy.	VERY GOOD New nodes only get the data they need and no need to change DB design.	LOW to MEDIUM Most queries are local, but global queries to create temporary comms load.	DIFFICULT Distributed table updates require tight precise coordination.	VERY POOR Requires considerable effort and inconsistencies are not tolerated.
Decentralised Independent Partitions	GOOD Depends on local DB availability.	GOOD New sites are independent of all other sites.	LOW Little or no traffic needs to be communicated across the network.	VERY GOOD Easy – as each site is independent of the other sites and minimal need to share data.	LOW No guarantee of consistency – therefore high chance of consistency.

- Locate data with a distributed catalog (meta data)
- Determine location from which to retrieve data and process query components
- DBMS translation between nodes with different local DBMSs (using middleware)
- Data consistency (via multiphase commit protocols)
- Global primary key control
- Scalability
- Security, concurrency, query optimization, failure recovery

Distributed DBMS architecture



- Imagine you have a synchronously-updating, replicated database –
- Now imagine that the link between 2 nodes is interrupted



- What are your choices?
 - shut down the system (to avoid inconsistency)
 - keep it available to users (and accept inconsistency)
- More about this in the NoSQL lecture ...

(diagram from <https://martin.kleppmann.com/2015/05/11/please-stop-calling-databases-cp-or-ap.html>)



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- **Distributed Database**
 - Advantages & Disadvantages
- **Replicated Databases**
 - Advantages & Disadvantages
- **Synchronous v Asynchronous**
 - Difference between
 - Advantages & Disadvantages
- **Partitioning Options**
 - Vertical, Physical, Vertical and Physical
- **The five configurations**
 - Advantages & Disadvantages



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L16 Distributed Databases