

CS 480

Introduction to Artificial Intelligence

February 8, 2022

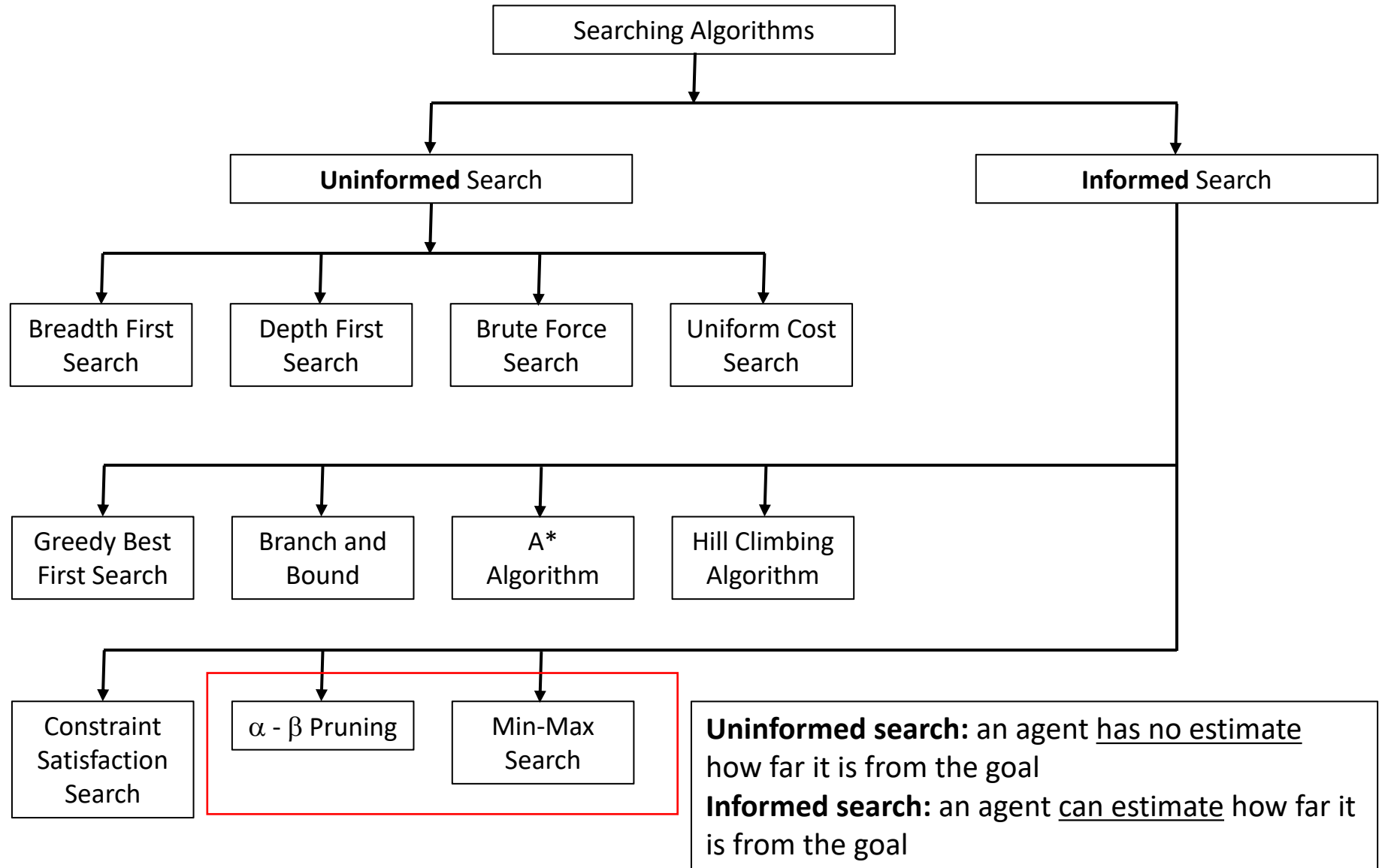
Announcements / Reminders

- **Written Assignment #01 is posted**
 - **due on Thursday (02/17/22) at 11:00 PM CST**
- **Contribute to the discussion on Blackboard, please**
- **Please follow the Week 04 To Do List instructions**
- **Midterm course review will be available next week**
 - **Please respond. Thank you!**

Plan for Today

- **Adversarial Search: MinMax / α - β Pruning**
- **Constraint Satisfaction Problems: Introduction**

Selected Searching Algorithms

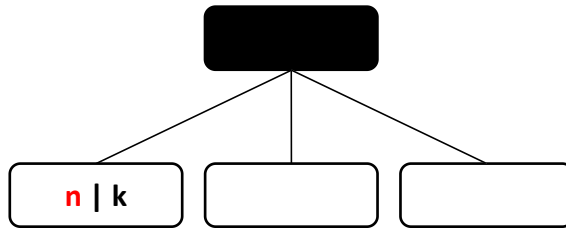


MinMax: Assigning MINMAX Values

CASE 1:

State **n** is Terminal Node

ISTERMINAL(**n**) = true
TOMOVE(**n**) = MAX or MIN



TERMINAL/LEAF

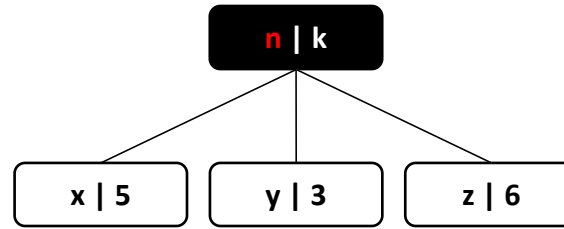
$$k = \text{MINMAX}(n) = \text{UTILITY}(n)$$

= utility value of this state for MAX Player

CASE 2:

State **n** is a Non-Terminal Node
and it is MIN Player's move

ISTERMINAL(**n**) = false
TOMOVE(**n**) = MIN



$$k = \text{MINMAX}(n) =$$

$$= \min_{a \in \text{ACTIONS}(n)} \text{MINMAX}(\text{RESULT}(n, a))$$

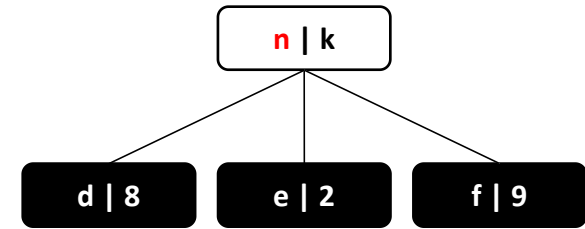
$$= \min(\text{MINMAX}(x), \text{MINMAX}(y), \text{MINMAX}(z))$$

$$= \min(5, 3, 6)$$

CASE 3:

State **n** is a Non-Terminal Node
and it is MAX Player's move

ISTERMINAL(**n**) = false
TOMOVE(**n**) = MAX



$$k = \text{MINMAX}(n) =$$

$$= \max_{a \in \text{ACTIONS}(n)} \text{MINMAX}(\text{RESULT}(n, a))$$

$$= \max(\text{MINMAX}(d), \text{MINMAX}(e), \text{MINMAX}(f))$$

$$= \max(8, 2, 9)$$

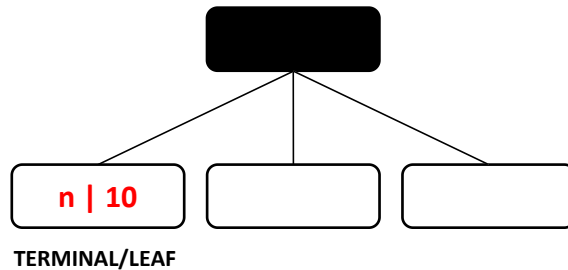
$$\text{MINMAX}(n) = \begin{cases} \text{UTILITY}(n, \text{MAX}), & \text{if } \text{ISTERMINAL}(n) \\ \max_{a \in \text{ACTIONS}(n)} \text{MINMAX}(\text{RESULT}(n, a)), & \text{if } \text{TOMOVE}(s) = \text{MAX} \\ \min_{a \in \text{ACTIONS}(n)} \text{MINMAX}(\text{RESULT}(n, a)), & \text{if } \text{TOMOVE}(s) = \text{MIN} \end{cases}$$

MinMax: Assigning MINMAX Values

CASE 1:

State **n** is Terminal Node

ISTERMINAL(**n**) = true
TOMOVE(**n**) = MAX or MIN



$$k = \text{MINMAX}(n) = \text{UTILITY}(n)$$

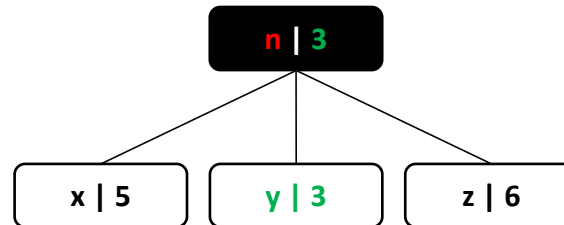
= utility value of this state for MAX Player

$$= 10$$

CASE 2:

State **n** is a Non-Terminal Node
and it is MIN Player's move

ISTERMINAL(**n**) = false
TOMOVE(**n**) = MIN



$$k = \text{MINMAX}(n) =$$

$$= \min_{a \in \text{ACTIONS}(n)} \text{MINMAX}(\text{RESULT}(n, a))$$

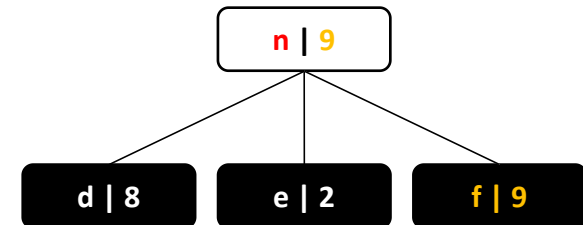
$$= \min(\text{MINMAX}(x), \text{MINMAX}(y), \text{MINMAX}(z))$$

$$= \min(5, 3, 6) = 3$$

CASE 3:

State **n** is a Non-Terminal Node
and it is MAX Player's move

ISTERMINAL(**n**) = false
TOMOVE(**n**) = MAX



$$k = \text{MINMAX}(n) =$$

$$= \max_{a \in \text{ACTIONS}(n)} \text{MINMAX}(\text{RESULT}(n, a))$$

$$= \max(\text{MINMAX}(d), \text{MINMAX}(e), \text{MINMAX}(f))$$

$$= \max(8, 2, 9) = 9$$

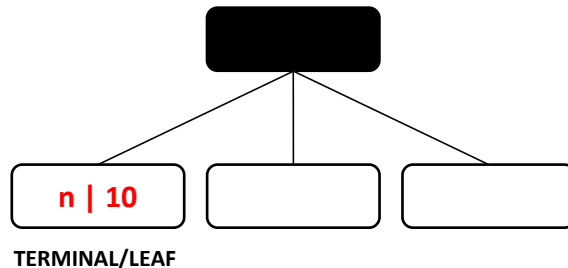
$$\text{MINMAX}(n) = \begin{cases} \text{UTILITY}(n, \text{MAX}), & \text{if } \text{ISTERMINAL}(n) \\ \max_{a \in \text{ACTIONS}(n)} \text{MINMAX}(\text{RESULT}(n, a)), & \text{if } \text{TOMOVE}(s) = \text{MAX} \\ \min_{a \in \text{ACTIONS}(n)} \text{MINMAX}(\text{RESULT}(n, a)), & \text{if } \text{TOMOVE}(s) = \text{MIN} \end{cases}$$

MinMax: Assigning MINMAX Values

CASE 1:

State **n** is Terminal Node

ISTERMINAL(**n**) = true
TOMOVE(**n**) = MAX or MIN



$$k = \text{MINMAX}(n) = \text{UTILITY}(n)$$

= utility value of this state for MAX Player

$$= 10$$

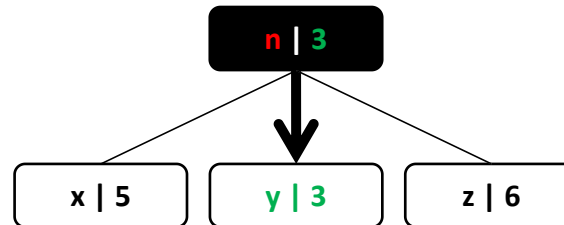
What does it mean?

Utility of node **n**, to MAX Player, is 10 (if the game gets here, this is what MAX Player will receive)

CASE 2:

State **n** is a Non-Terminal Node
and it is MIN Player's move

ISTERMINAL(**n**) = false
TOMOVE(**n**) = MIN



$$k = \text{MINMAX}(n) =$$

$$= \min_{a \in \text{ACTIONS}(n)} \text{MINMAX}(\text{RESULT}(n, a))$$

$$= \min(\text{MINMAX}(x), \text{MINMAX}(y), \text{MINMAX}(z))$$

$$= \min(5, 3, 6) = 3$$

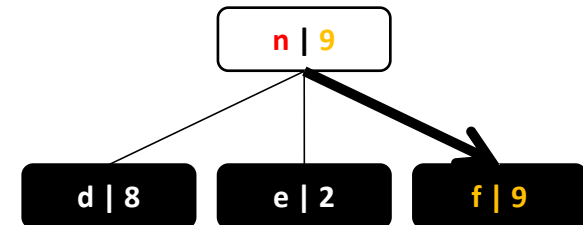
What does it mean?

At node **n**, MIN Player will choose a move from **n** to **y** to MINIMIZE MAX Player's utility

CASE 3:

State **n** is a Non-Terminal Node
and it is MAX Player's move

ISTERMINAL(**n**) = false
TOMOVE(**n**) = MAX



$$k = \text{MINMAX}(n) =$$

$$= \max_{a \in \text{ACTIONS}(n)} \text{MINMAX}(\text{RESULT}(n, a))$$

$$= \max(\text{MINMAX}(d), \text{MINMAX}(e), \text{MINMAX}(f))$$

$$= \max(8, 2, 9) = 9$$

What does it mean?

At node **n**, MAX Player will choose a move from **n** to **f** to MAXIMIZE MAX Player's utility

MinMax Algorithm: Pseudocode

function MINIMAX-SEARCH(*game*, *state*) **returns** *an action*

$\text{player} \leftarrow \text{game}.\text{TO-MOVE}(\text{state})$

$\text{value}, \text{move} \leftarrow \text{MAX-VALUE}(\text{game}, \text{state})$

return *move*

function MAX-VALUE(*game*, *state*) **returns** *a (utility, move) pair*

if *game*.IS-TERMINAL(*state*) **then return** *game*.UTILITY(*state*, *player*), *null*

$v \leftarrow -\infty$

for each *a* **in** *game*.ACTIONS(*state*) **do**

$v2, a2 \leftarrow \text{MIN-VALUE}(\text{game}, \text{game}.\text{RESULT}(\text{state}, a))$

if $v2 > v$ **then**

$v, \text{move} \leftarrow v2, a$

return *v*, *move*

function MIN-VALUE(*game*, *state*) **returns** *a (utility, move) pair*

if *game*.IS-TERMINAL(*state*) **then return** *game*.UTILITY(*state*, *player*), *null*

$v \leftarrow +\infty$

for each *a* **in** *game*.ACTIONS(*state*) **do**

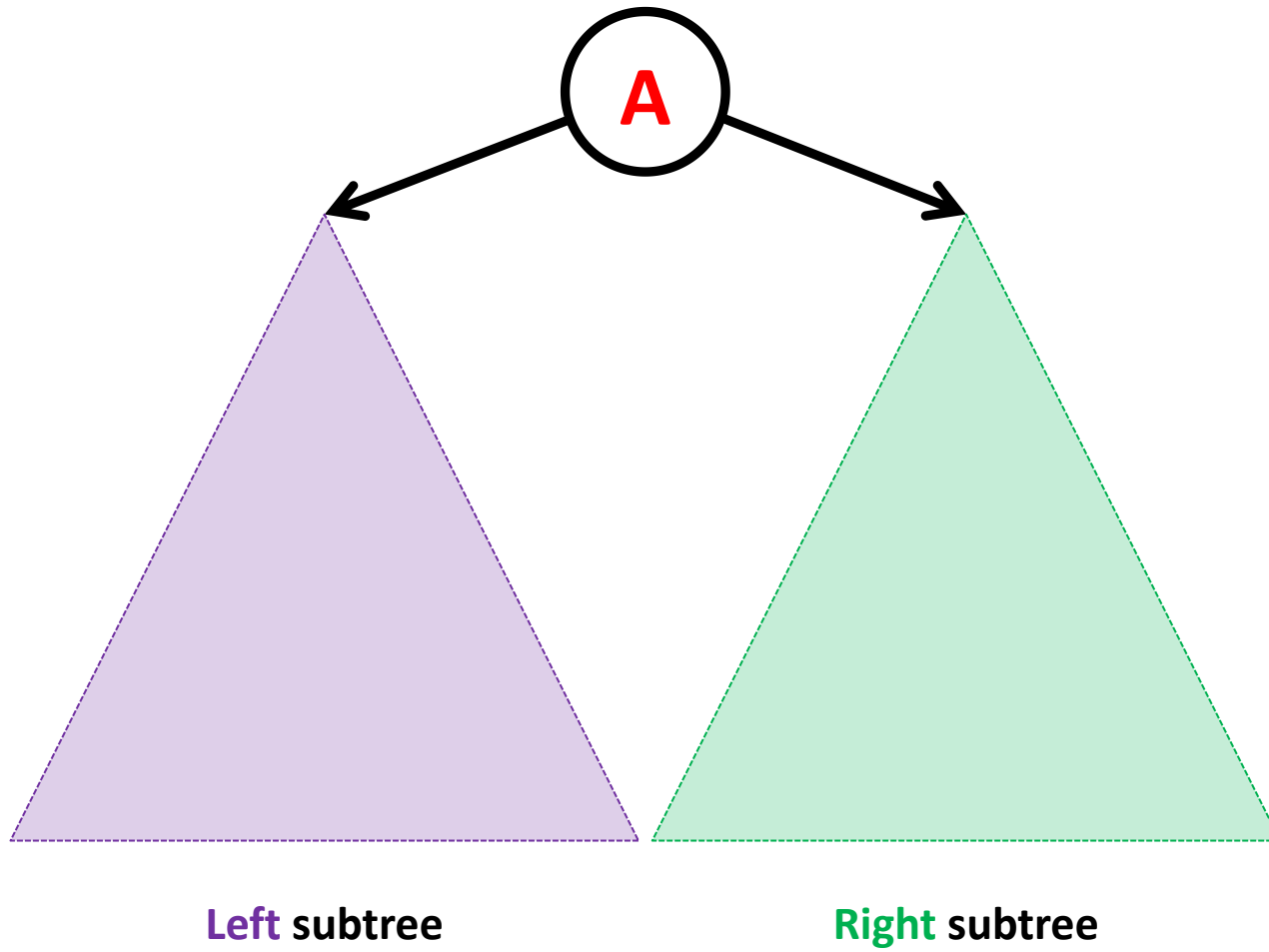
$v2, a2 \leftarrow \text{MAX-VALUE}(\text{game}, \text{game}.\text{RESULT}(\text{state}, a))$

if $v2 < v$ **then**

$v, \text{move} \leftarrow v2, a$

return *v*, *move*

Search Tree: Recursive Structure



MinMax Algorithm: Pseudocode

function MINIMAX-SEARCH(*game*, *state*) **returns** an action

$\text{player} \leftarrow \text{game}.\text{TO-MOVE}(\text{state})$

$\text{value}, \text{move} \leftarrow \text{MAX-VALUE}(\text{game}, \text{state})$

return *move*

function MAX-VALUE(*game*, *state*) **returns** a (*utility*, *move*) pair

if *game*.IS-TERMINAL(*state*) **then return** *game*.UTILITY(*state*, *player*), null

$v \leftarrow -\infty$

for each *a* **in** *game*.ACTIONS(*state*) **do**

$v2, a2 \leftarrow \text{MIN-VALUE}(\text{game}, \text{game}.\text{RESULT}(\text{state}, a))$

if $v2 > v$ **then**

$v, \text{move} \leftarrow v2, a$

return *v*, *move*

function MIN-VALUE(*game*, *state*) **returns** a (*utility*, *move*) pair

if *game*.IS-TERMINAL(*state*) **then return** *game*.UTILITY(*state*, *player*), null

$v \leftarrow +\infty$

for each *a* **in** *game*.ACTIONS(*state*) **do**

$v2, a2 \leftarrow \text{MAX-VALUE}(\text{game}, \text{game}.\text{RESULT}(\text{state}, a))$

if $v2 < v$ **then**

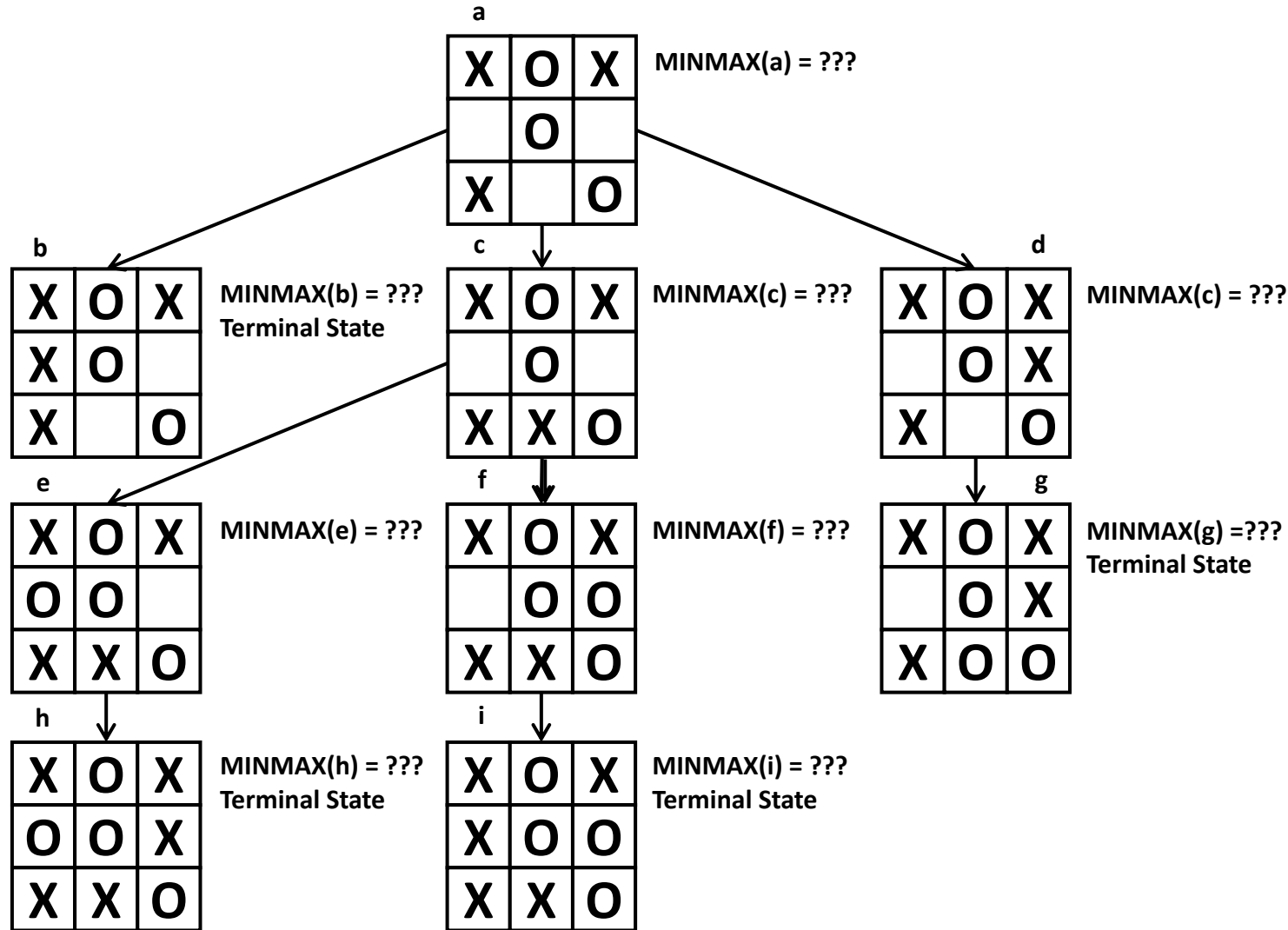
$v, \text{move} \leftarrow v2, a$

return *v*, *move*

RECURSION

A diagram illustrating the recursive nature of the MinMax algorithm. A central point has three arrows pointing outwards. One arrow points to the 'MAX-VALUE' function, another points to the 'MIN-VALUE' function, and a third points to the 'MINIMAX-SEARCH' function. This indicates that each of these functions can call itself or the other to solve subproblems.

MinMax Algorithm: Tic Tac Toe



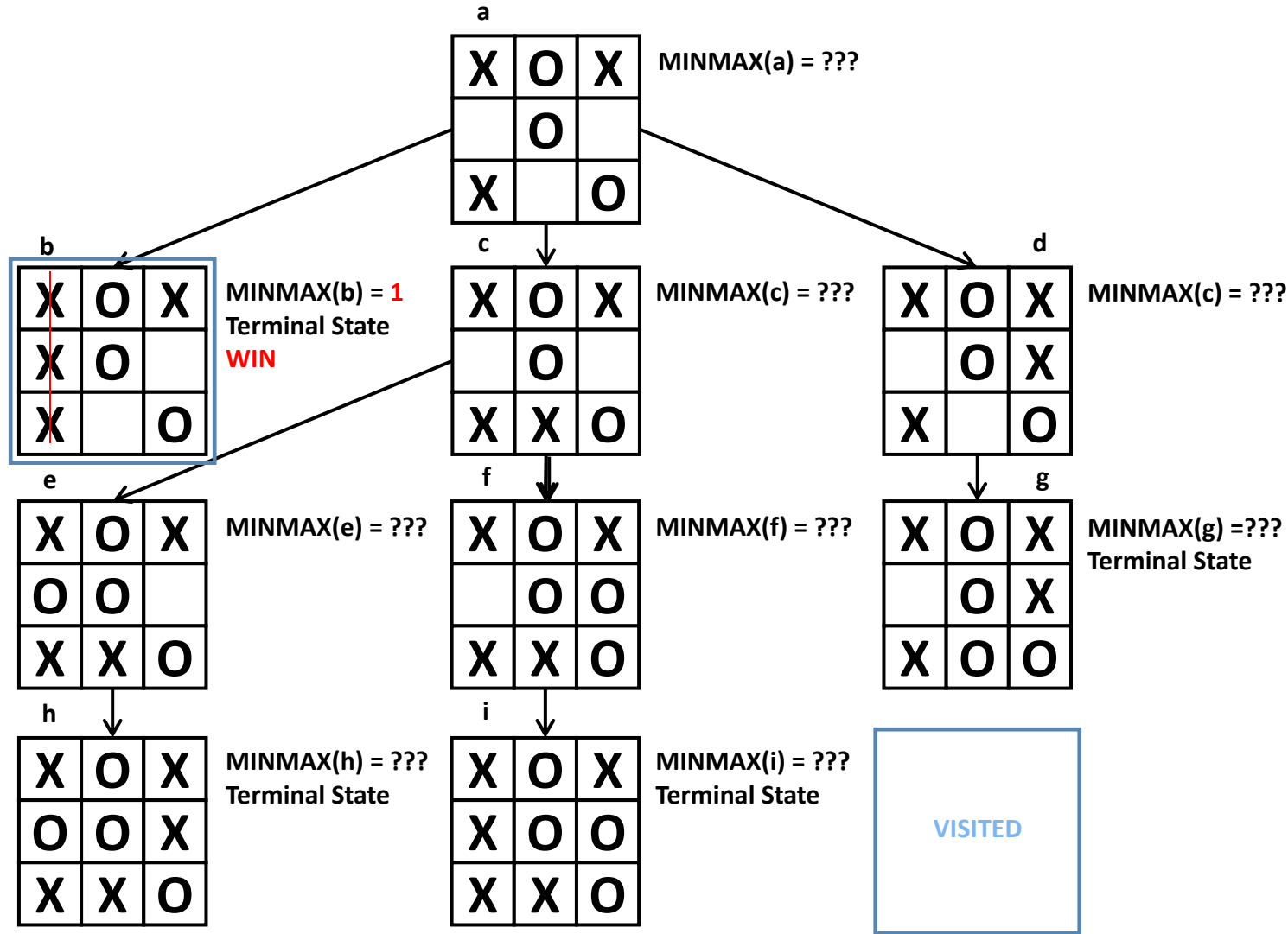
X's move
(MAX Player)

O's move
(MIN Player)

X's move
(MAX Player)

O's move
(MIN Player)

MinMax Algorithm: Tic Tac Toe



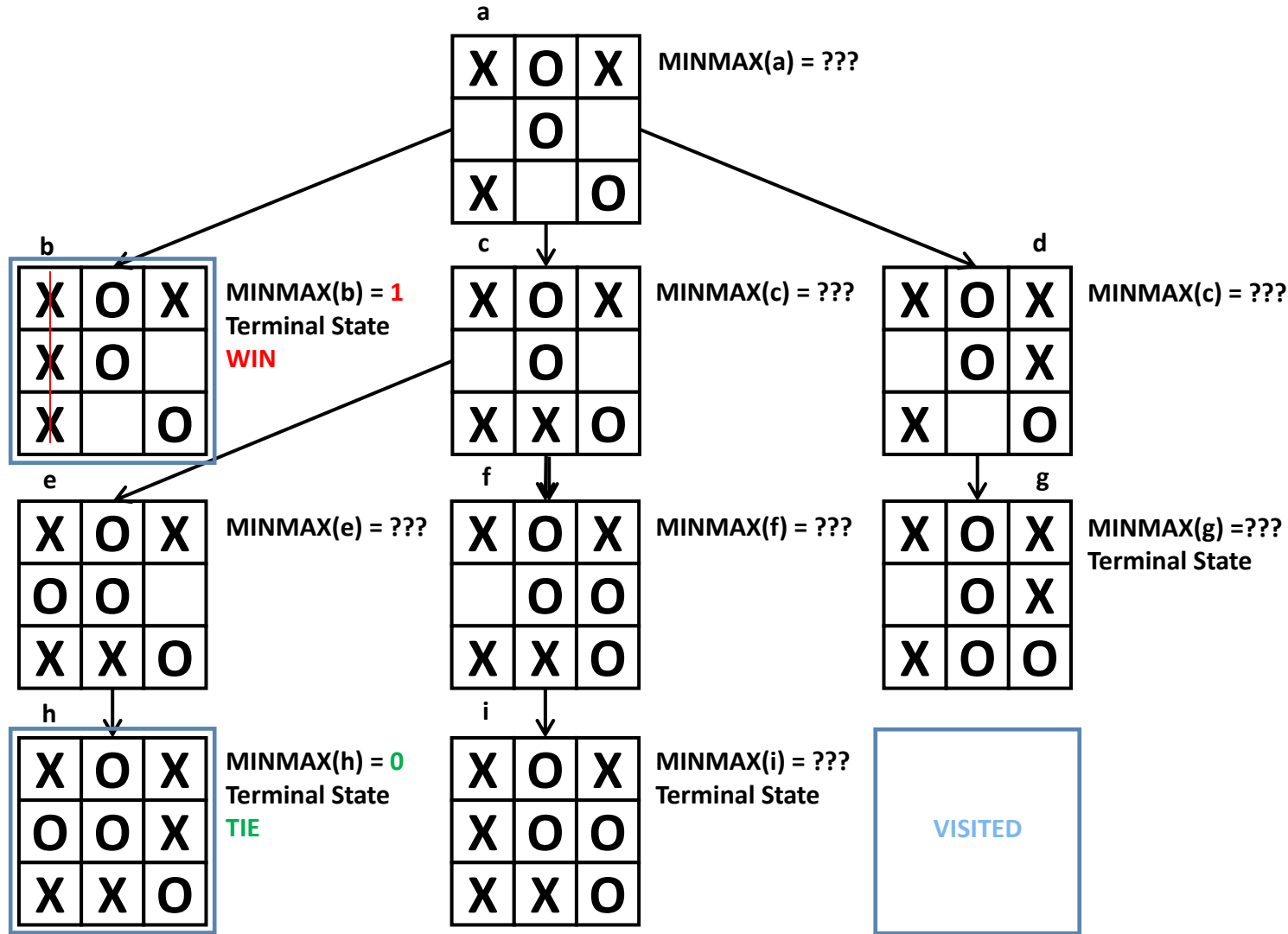
X's move
(MAX Player)

O's move
(MIN Player)

X's move
(MAX Player)

O's move
(MIN Player)

MinMax Algorithm: Tic Tac Toe



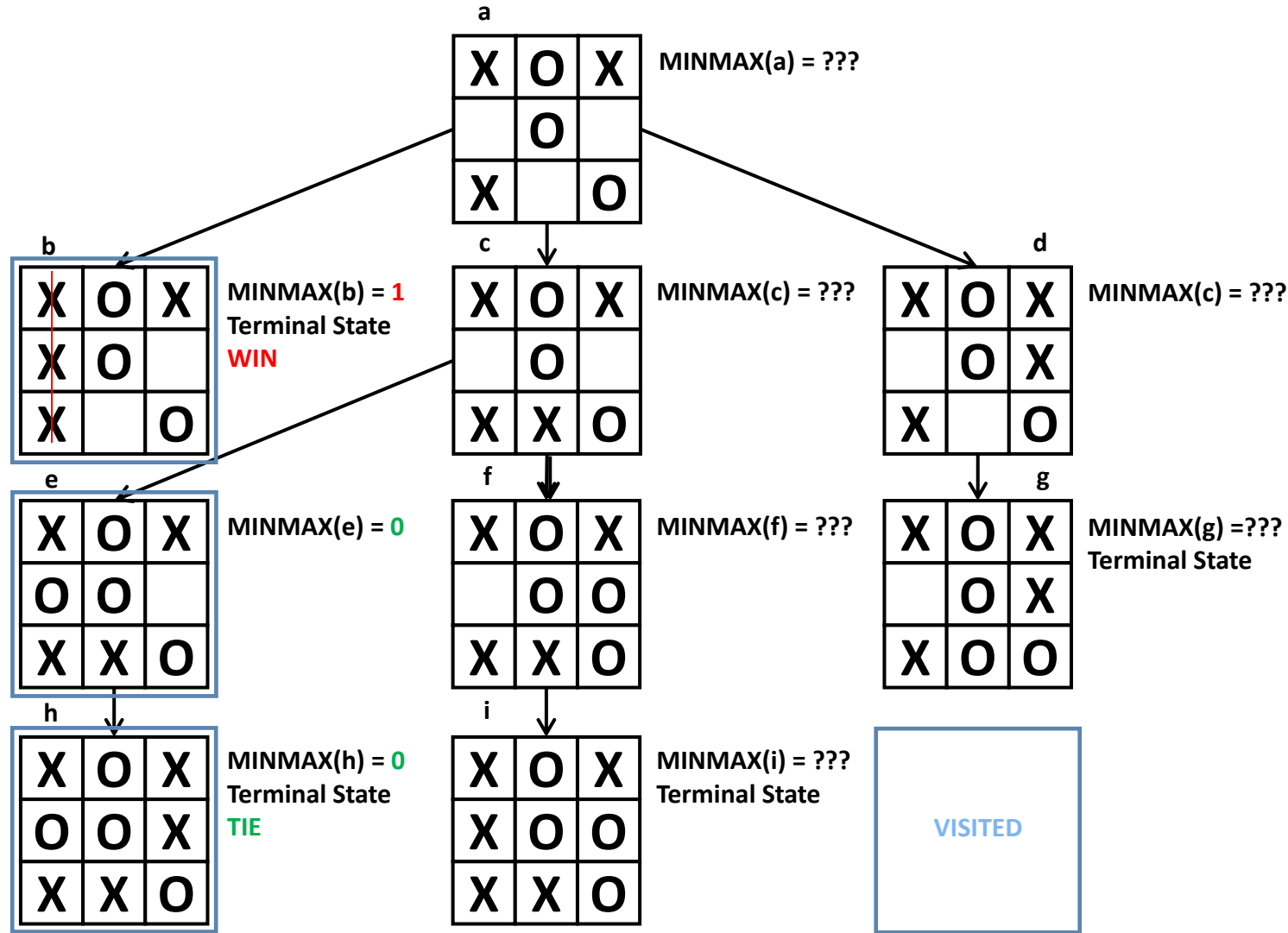
X's move
(MAX Player)

O's move
(MIN Player)

X's move
(MAX Player)

O's move
(MIN Player)

MinMax Algorithm: Tic Tac Toe



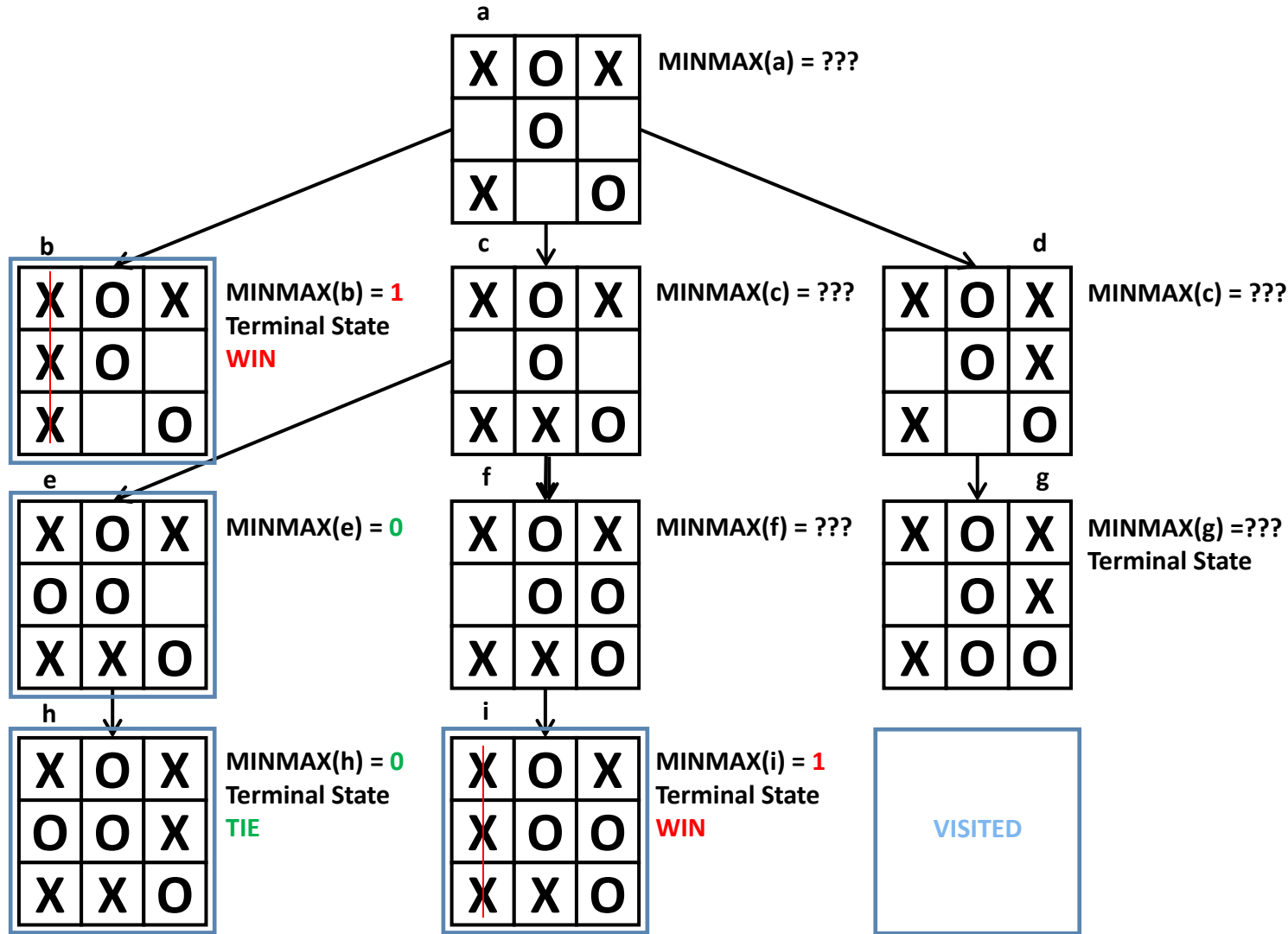
X's move
(MAX Player)

O's move
(MIN Player)

X's move
(MAX Player)

O's move
(MIN Player)

MinMax Algorithm: Tic Tac Toe



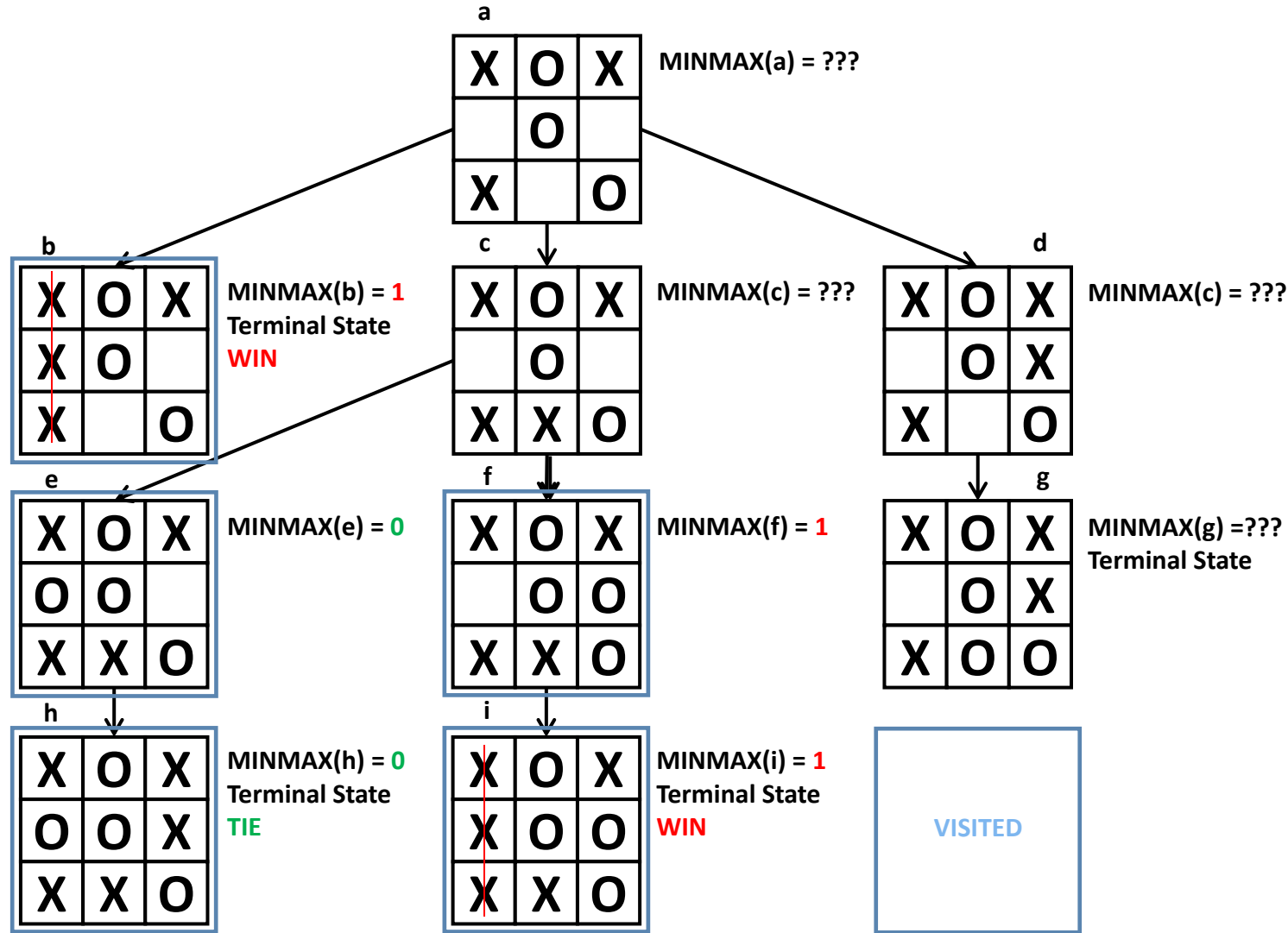
X's move
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X's move
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O's move
(MIN Player)

MinMax Algorithm: Tic Tac Toe



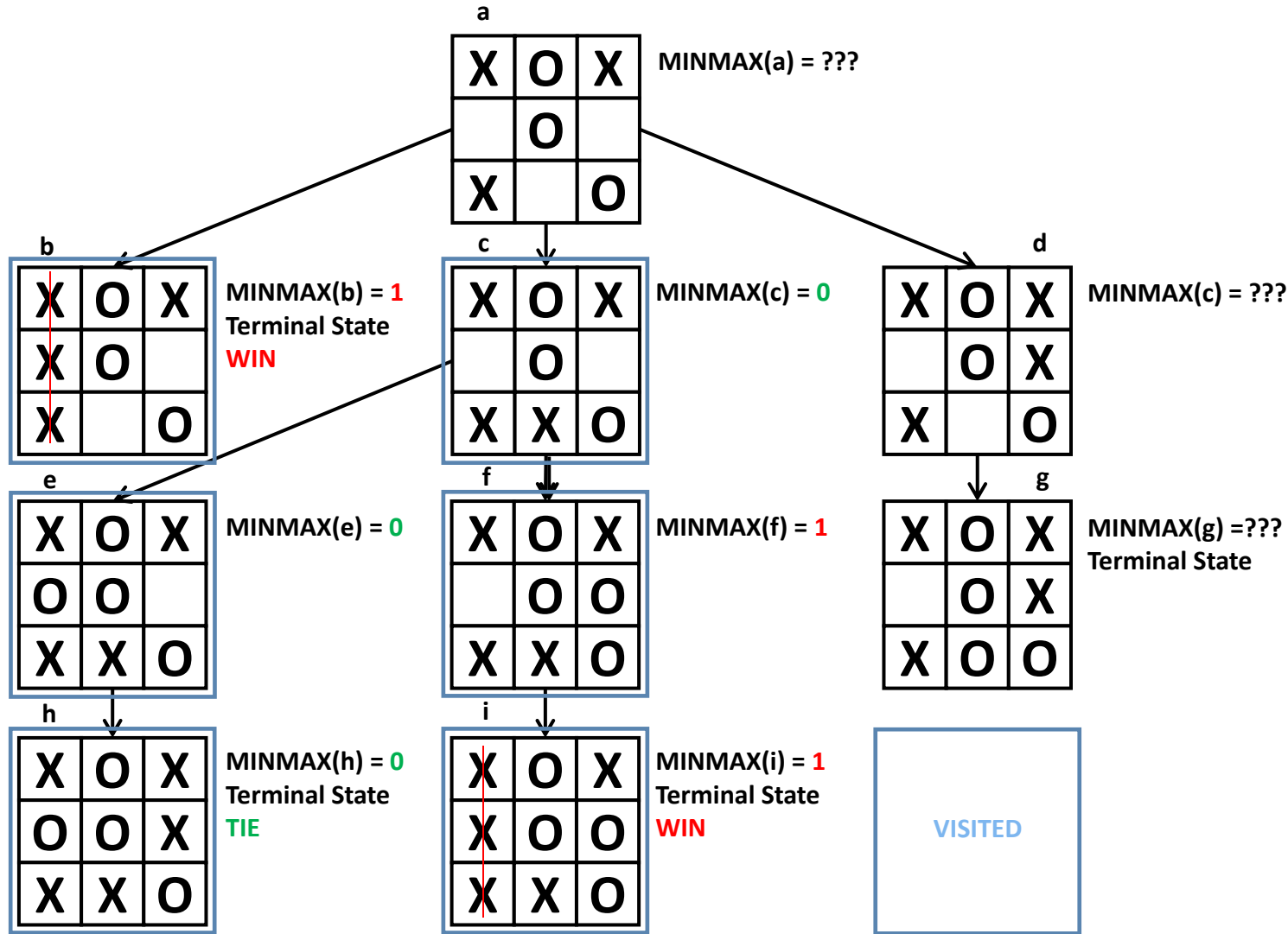
X's move
(MAX Player)

O's move
(MIN Player)

X's move
(MAX Player)

O's move
(MIN Player)

MinMax Algorithm: Tic Tac Toe



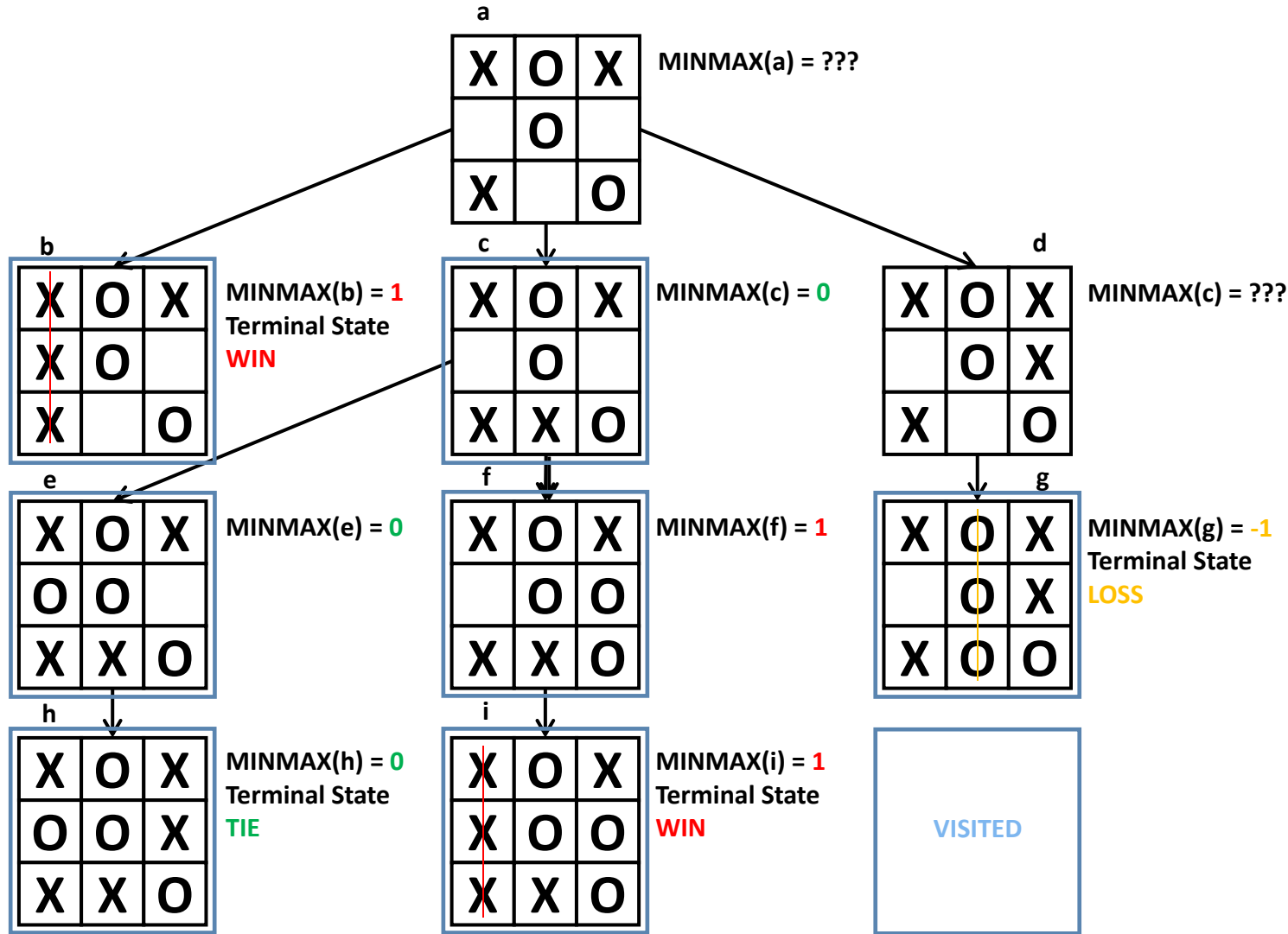
X's move
(MAX Player)

O's move
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X's move
(MAX Player)

O's move
(MIN Player)

MinMax Algorithm: Tic Tac Toe



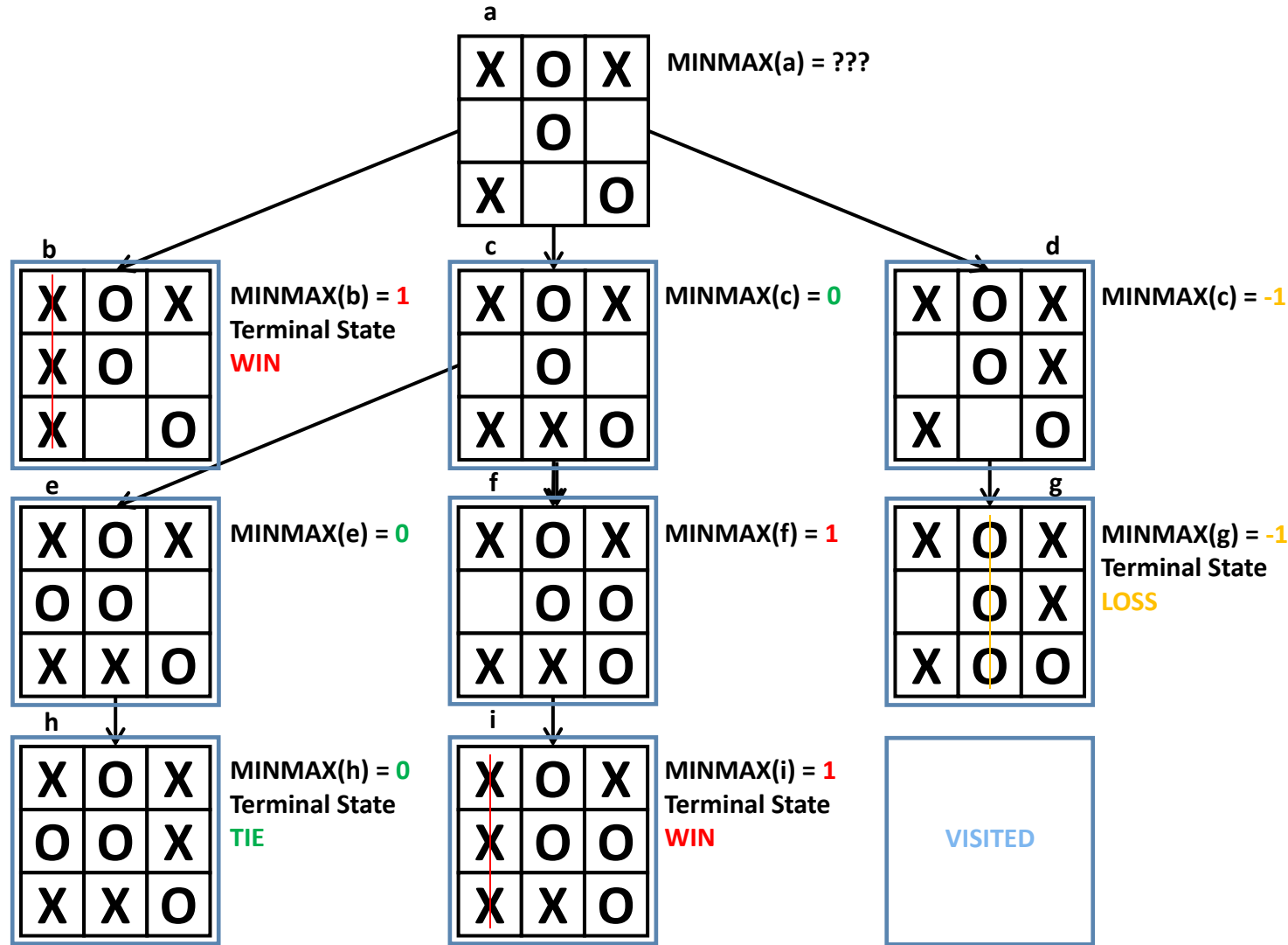
X's move
(MAX Player)

O's move
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X's move
(MAX Player)

O's move
(MIN Player)

MinMax Algorithm: Tic Tac Toe



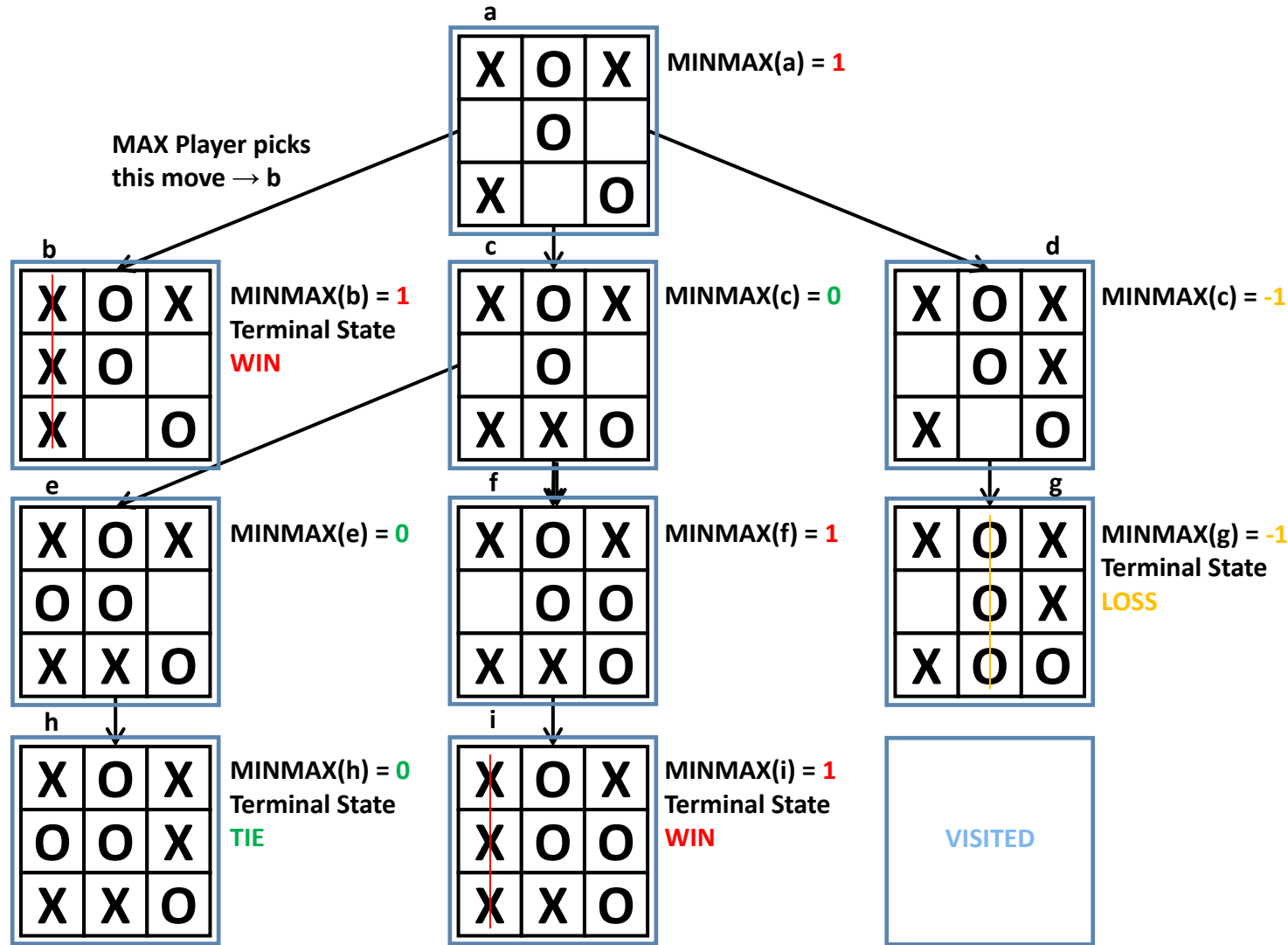
X's move
(MAX Player)

O's move
(MIN Player)

X's move
(MAX Player)

O's move
(MIN Player)

MinMax Algorithm: Tic Tac Toe



X's move
(MAX Player)

O's move
(MIN Player)

X's move
(MAX Player)

O's move
(MIN Player)

MinMax with α - β : Pseudocode

function ALPHA-BETA-SEARCH(*game*, *state*) **returns** an action

$\text{player} \leftarrow \text{game.TO-MOVE}(\text{state})$

$\text{value}, \text{move} \leftarrow \text{MAX-VALUE}(\text{game}, \text{state}, -\infty, +\infty)$

return *move*

function MAX-VALUE(*game*, *state*, α , β) **returns** a (*utility*, *move*) pair

if *game.IS-TERMINAL*(*state*) **then return** *game.UTILITY*(*state*, *player*), null

$v \leftarrow -\infty$

for each *a* **in** *game.ACTIONS*(*state*) **do**

$v2, a2 \leftarrow \text{MIN-VALUE}(\text{game}, \text{game.RESULT}(\text{state}, a), \alpha, \beta)$

if $v2 > v$ **then**

$v, \text{move} \leftarrow v2, a$

$\alpha \leftarrow \text{MAX}(\alpha, v)$

if $v \geq \beta$ **then return** *v*, *move*

return *v*, *move*

function MIN-VALUE(*game*, *state*, α , β) **returns** a (*utility*, *move*) pair

if *game.IS-TERMINAL*(*state*) **then return** *game.UTILITY*(*state*, *player*), null

$v \leftarrow +\infty$

for each *a* **in** *game.ACTIONS*(*state*) **do**

$v2, a2 \leftarrow \text{MAX-VALUE}(\text{game}, \text{game.RESULT}(\text{state}, a), \alpha, \beta)$

if $v2 < v$ **then**

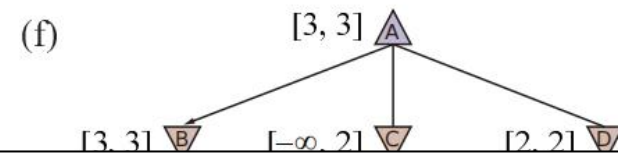
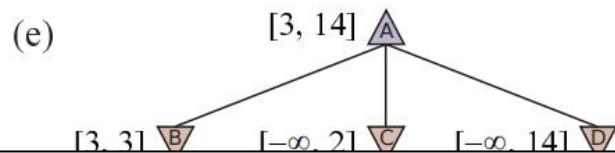
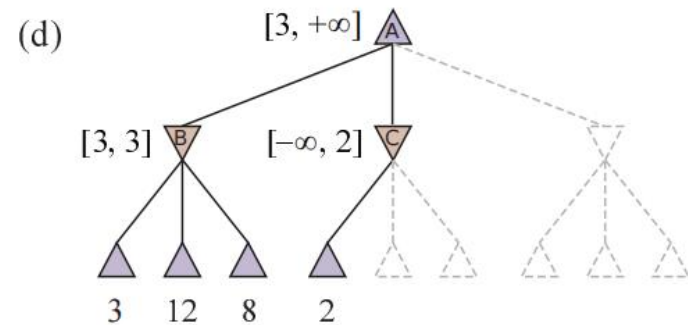
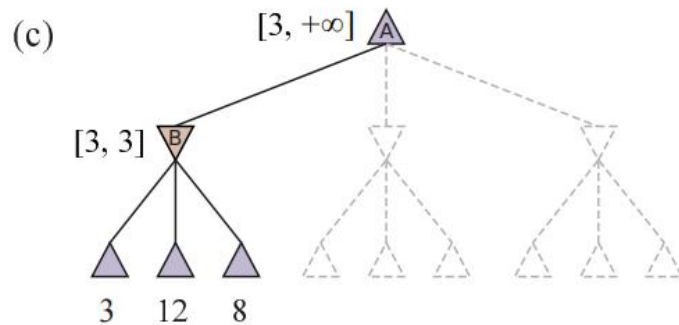
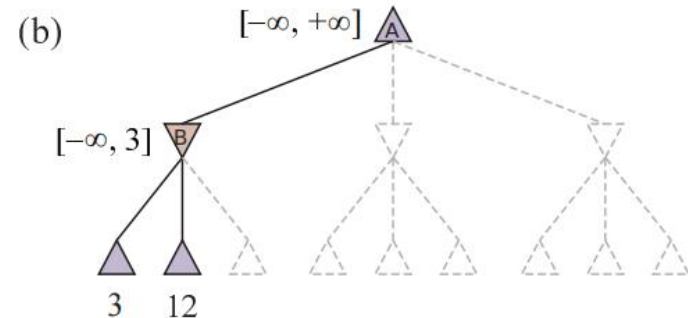
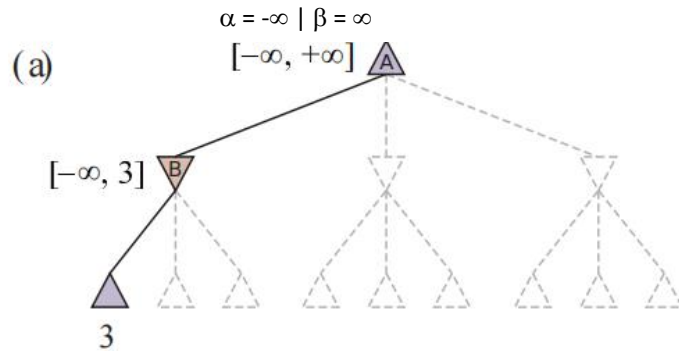
$v, \text{move} \leftarrow v2, a$

$\beta \leftarrow \text{MIN}(\beta, v)$

if $v \leq \alpha$ **then return** *v*, *move*

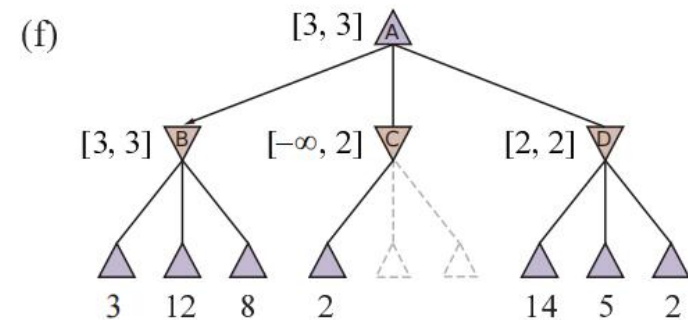
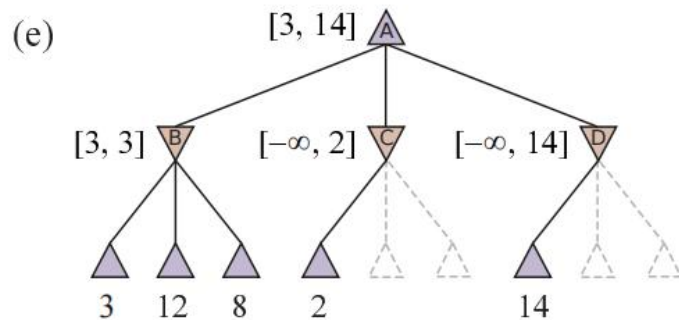
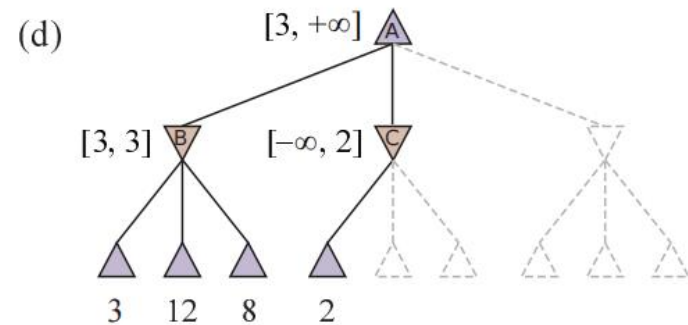
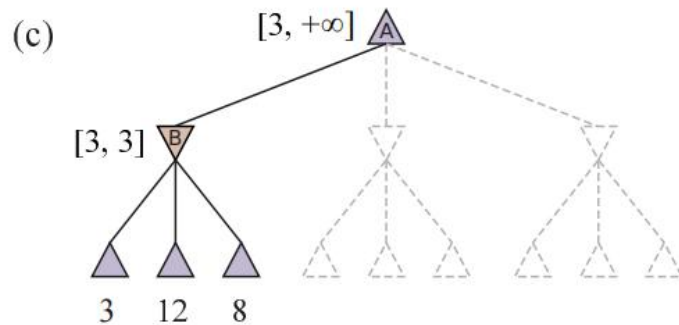
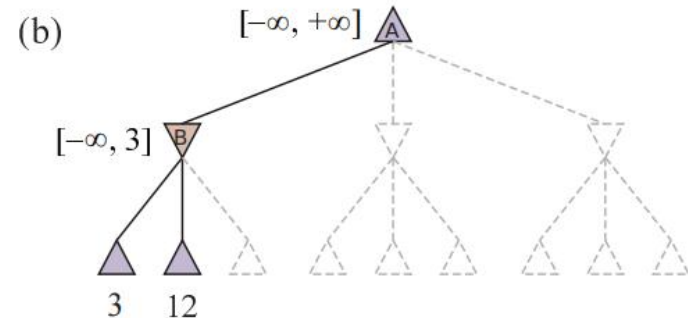
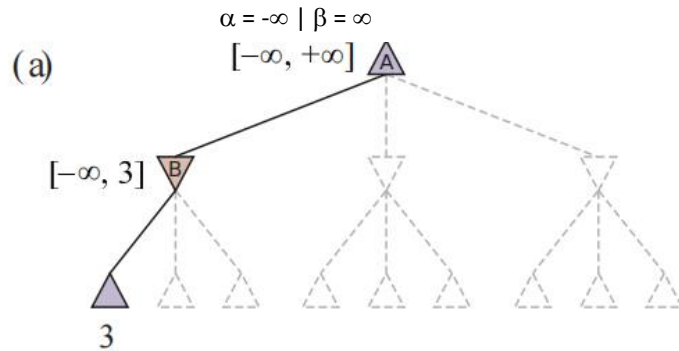
return *v*, *move*

Example MinMax with α - β Pruning



α : the value of the best (highest-value) choice we have found so far at any choice point along the path for MAX player ("at least")
 β : the value of the best (lowest-value) choice we have found so far at any choice point along the path for MIN player ("at most")

Example MinMax with α - β Pruning



MinMax with α - β : Pseudocode

function ALPHA-BETA-SEARCH(*game*, *state*) **returns** an action

player \leftarrow *game*.TO-MOVE(*state*)

value, *move* \leftarrow MAX-VALUE(*game*, *state*, $-\infty$, $+\infty$)

return *move*

function MAX-VALUE(*game*, *state*, α , β) **returns** a (*utility*, *move*) pair

if *game*.IS-TERMINAL(*state*) **then return** *game*.UTILITY(*state*, *player*), null

v $\leftarrow -\infty$

for each *a* **in** *game*.ACTIONS(*state*) **do**

v2, *a2* \leftarrow MIN-VALUE(*game*, *game*.RESULT(*state*, *a*), α , β)

if *v2* > *v* **then**

v, *move* \leftarrow *v2*, *a*

$\alpha \leftarrow$ MAX(α , *v*)

if *v* \geq β **then return** *v*, *move*

return *v*, *move*

function MIN-VALUE(*game*, *state*, α , β) **returns** a (*utility*, *move*) pair

if *game*.IS-TERMINAL(*state*) **then return** *game*.UTILITY(*state*, *player*), null

v $\leftarrow +\infty$

for each *a* **in** *game*.ACTIONS(*state*) **do**

v2, *a2* \leftarrow MAX-VALUE(*game*, *game*.RESULT(*state*, *a*), α , β)

if *v2* < *v* **then**

v, *move* \leftarrow *v2*, *a*

$\beta \leftarrow$ MIN(β , *v*)

if *v* \leq α **then return** *v*, *move*

return *v*, *move*

RECURSION

A diagram consisting of three arrows originating from a single point on the right, labeled 'RECURSION'. One arrow points upwards and to the left towards the MAX-VALUE function call in the ALPHA-BETA-SEARCH function. A second arrow points horizontally to the left towards the MIN-VALUE function call within the MAX-VALUE function. A third arrow points downwards and to the left towards the MAX-VALUE function call within the MIN-VALUE function.

MinMax with α - β : Pseudocode

```
function ALPHA-BETA-SEARCH(game, state) returns an action  
  player  $\leftarrow$  game.TO-MOVE(state)  
  value, move  $\leftarrow$  MAX-VALUE(game, state,  $-\infty$ ,  $+\infty$ )  
  return move
```

```
function MAX-VALUE(game, state,  $\alpha$ ,  $\beta$ ) returns a (utility, move) pair  
  if game.IS-TERMINAL(state) then return game.UTILITY(state, player), null  
  v  $\leftarrow$   $-\infty$   
  for each a in game.ACTIONS(state) do  
    v2, a2  $\leftarrow$  MIN-VALUE(game, game.RESULT(state, a),  $\alpha$ ,  $\beta$ )  
    if v2 > v then  
      v, move  $\leftarrow$  v2, a  
       $\alpha \leftarrow$  MAX( $\alpha$ , v)  
    if v  $\geq$   $\beta$  then return v, move  
  return v, move
```

MAX Player's move

```
function MIN-VALUE(game, state,  $\alpha$ ,  $\beta$ ) returns a (utility, move) pair  
  if game.IS-TERMINAL(state) then return game.UTILITY(state, player), null  
  v  $\leftarrow$   $+\infty$   
  for each a in game.ACTIONS(state) do  
    v2, a2  $\leftarrow$  MAX-VALUE(game, game.RESULT(state, a),  $\alpha$ ,  $\beta$ )  
    if v2 < v then  
      v, move  $\leftarrow$  v2, a  
       $\beta \leftarrow$  MIN( $\beta$ , v)  
    if v  $\leq$   $\alpha$  then return v, move  
  return v, move
```

MIN Player's move

MinMax with α - β : Pseudocode

```
function ALPHA-BETA-SEARCH(game, state) returns an action  
  player  $\leftarrow$  game.TO-MOVE(state)  
  value, move  $\leftarrow$  MAX-VALUE(game, state,  $-\infty$ ,  $+\infty$ )  
  return move
```

```
function MAX-VALUE(game, state,  $\alpha$ ,  $\beta$ ) returns a (utility, move) pair  
  if game.IS-TERMINAL(state) then return game.UTILITY(state, player), null  
  v  $\leftarrow$   $-\infty$   
  for each a in game.ACTIONS(state) do  
    v2, a2  $\leftarrow$  MIN-VALUE(game, game.RESULT(state, a),  $\alpha$ ,  $\beta$ )  
    if v2 > v then  
      v, move  $\leftarrow$  v2, a  
       $\alpha \leftarrow$  MAX( $\alpha$ , v)  
    if v  $\geq$   $\beta$  then return v, move  
  return v, move
```

Go through all legal
actions/moves
(subtrees) **recursively**

MAX Player's move

```
function MIN-VALUE(game, state,  $\alpha$ ,  $\beta$ ) returns a (utility, move) pair  
  if game.IS-TERMINAL(state) then return game.UTILITY(state, player), null  
  v  $\leftarrow$   $+\infty$   
  for each a in game.ACTIONS(state) do  
    v2, a2  $\leftarrow$  MAX-VALUE(game, game.RESULT(state, a),  $\alpha$ ,  $\beta$ )  
    if v2 < v then  
      v, move  $\leftarrow$  v2, a  
       $\beta \leftarrow$  MIN( $\beta$ , v)  
    if v  $\leq$   $\alpha$  then return v, move  
  return v, move
```

Go through all legal
actions/moves
(subtrees) **recursively**

MIN Player's move

MinMax with α - β : Pseudocode

```
function ALPHA-BETA-SEARCH(game, state) returns an action  
  player  $\leftarrow$  game.TO-MOVE(state)  
  value, move  $\leftarrow$  MAX-VALUE(game, state,  $-\infty$ ,  $+\infty$ )  
  return move
```

```
function MAX-VALUE(game, state,  $\alpha$ ,  $\beta$ ) returns a (utility, move) pair  
  if game.IS-TERMINAL(state) then return game.UTILITY(state, player), null  
  v  $\leftarrow$   $-\infty$   
  for each a in game.ACTIONS(state) do  
    v2, a2  $\leftarrow$  MIN-VALUE(game, game.RESULT(state, a),  $\alpha$ ,  $\beta$ )  
    if v2 > v then                                If higher MINMAX(subtree) value found  
      v, move  $\leftarrow$  v2, a                               store a as the best move  
       $\alpha \leftarrow$  MAX( $\alpha$ , v)                               update bound  $\alpha$  (within this recursive call only!)  
    if v  $\geq$   $\beta$  then return v, move  
  return v, move
```

Go through all legal
actions/moves
(subtrees) **recursively**

MAX Player's move

```
function MIN-VALUE(game, state,  $\alpha$ ,  $\beta$ ) returns a (utility, move) pair  
  if game.IS-TERMINAL(state) then return game.UTILITY(state, player), null  
  v  $\leftarrow$   $+\infty$   
  for each a in game.ACTIONS(state) do  
    v2, a2  $\leftarrow$  MAX-VALUE(game, game.RESULT(state, a),  $\alpha$ ,  $\beta$ )  
    if v2 < v then                                If lower MINMAX(subtree) value found:  
      v, move  $\leftarrow$  v2, a                               store a as the best move  
       $\beta \leftarrow$  MIN( $\beta$ , v)                               update bound  $\beta$  (within this recursive call only!)  
    if v  $\leq$   $\alpha$  then return v, move  
  return v, move
```

Go through all legal
actions/moves
(subtrees) **recursively**

MIN Player's move

MinMax with α - β : Pseudocode

```
function ALPHA-BETA-SEARCH(game, state) returns an action  
  player  $\leftarrow$  game.TO-MOVE(state)  
  value, move  $\leftarrow$  MAX-VALUE(game, state,  $-\infty$ ,  $+\infty$ )  
  return move
```

```
function MAX-VALUE(game, state,  $\alpha$ ,  $\beta$ ) returns a (utility, move) pair  
  if game.IS-TERMINAL(state) then return game.UTILITY(state, player), null  
  v  $\leftarrow$   $-\infty$   
  for each a in game.ACTIONS(state) do  
    v2, a2  $\leftarrow$  MIN-VALUE(game, game.RESULT(state, a),  $\alpha$ ,  $\beta$ )  
    if v2 > v then                                If higher MINMAX(subtree) value found  
      v, move  $\leftarrow$  v2, a                               store a as the best move  
       $\alpha \leftarrow$  MAX( $\alpha$ , v)                               update bound  $\alpha$  (within this recursive call only!)  
    if v  $\geq$   $\beta$  then return v, move  
  return v, move
```

Go through all legal
actions/moves
(subtrees) **recursively**

MAX Player **does NOT**
change bound β here!

MAX Player's move

```
function MIN-VALUE(game, state,  $\alpha$ ,  $\beta$ ) returns a (utility, move) pair  
  if game.IS-TERMINAL(state) then return game.UTILITY(state, player), null  
  v  $\leftarrow$   $+\infty$   
  for each a in game.ACTIONS(state) do  
    v2, a2  $\leftarrow$  MAX-VALUE(game, game.RESULT(state, a),  $\alpha$ ,  $\beta$ )  
    if v2 < v then                                If lower MINMAX(subtree) value found:  
      v, move  $\leftarrow$  v2, a                               store a as the best move  
       $\beta \leftarrow$  MIN( $\beta$ , v)                               update bound  $\beta$  (within this recursive call only!)  
    if v  $\leq$   $\alpha$  then return v, move  
  return v, move
```

Go through all legal
actions/moves
(subtrees) **recursively**

MIN Player **does NOT**
change bound α here!

MIN Player's move

MinMax with α - β : Example

n | MINMAX(n)

MAX player state / move / turn

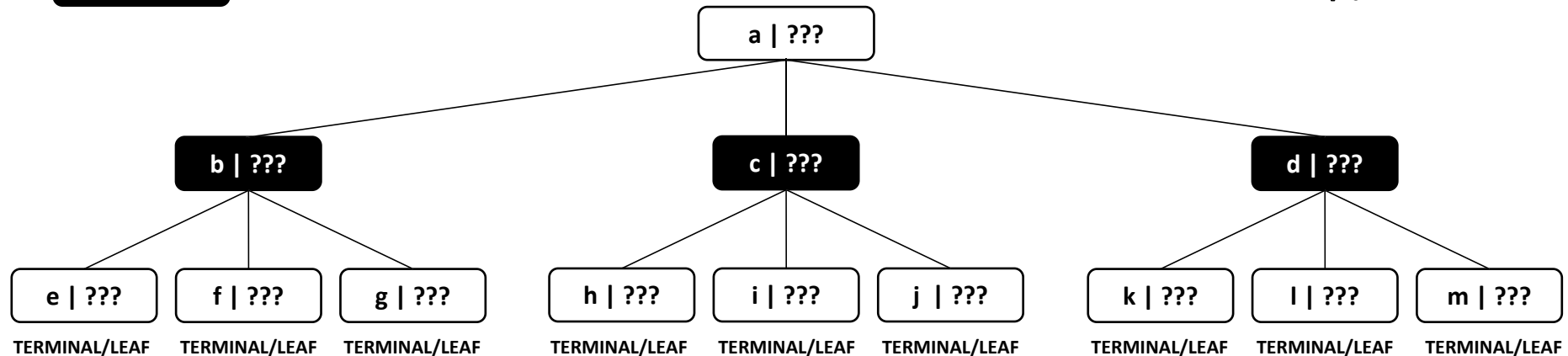
n | MINMAX(n)

MIN player state / move / turn

MINMAX(a) can be

at least: $\alpha_a = -\infty$

at most: $\beta_a = \infty$



MAX Player wants to maximize the utility of the game by choosing the right move from state a to one of three successor states: b, c, d. It will choose the state with maximum MAXMIN value. Currently:

- MINMAX(b) = UNKNOWN | MINMAX(c) = UNKNOWN | MINMAX(d) = UNKNOWN
- MAX Player's decision: not enough information yet.

MINMAX(a) = $\max(\text{MINMAX}(b), \text{MINMAX}(c), \text{MINMAX}(d)) = \max(???, ???, ???) \rightarrow$ can't be established

MinMax with α - β : Example

n | MINMAX(n)

MAX player state / move / turn

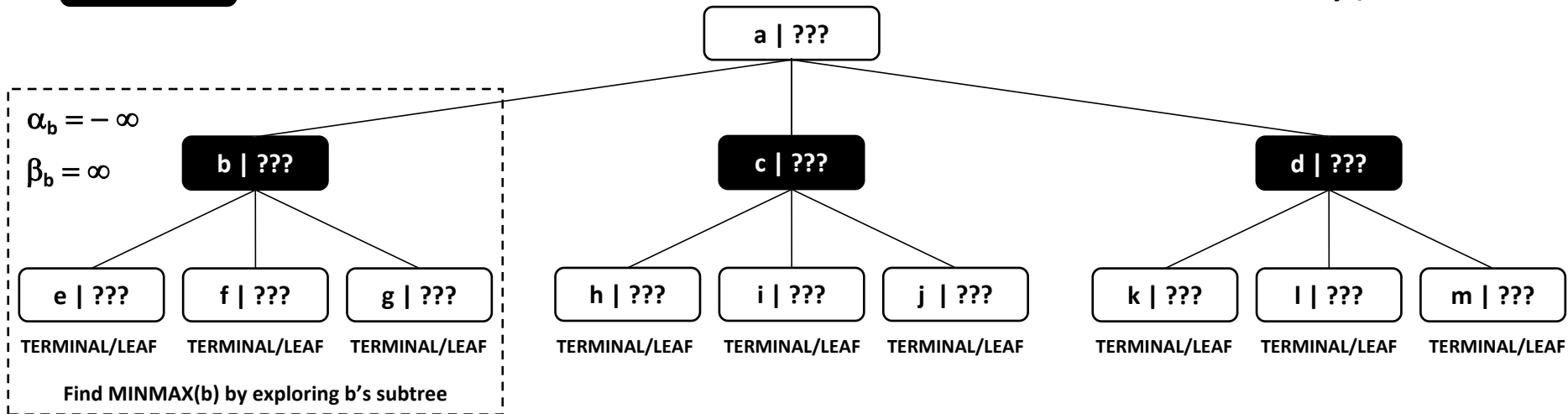
n | MINMAX(n)

MIN player state / move / turn

MINMAX(a) can be

at least: $\alpha_a = -\infty$

at most: $\beta_a = \infty$



MAX Player wants to maximize the utility of the game by choosing the right move from state a to one of three successor states: b, c, d. It will choose the state with maximum MAXMIN value. Currently:

- $\text{MINMAX}(b) = \text{UNKNOWN}$ | $\text{MINMAX}(c) = \text{UNKNOWN}$ | $\text{MINMAX}(d) = \text{UNKNOWN}$
 - MAX Player's decision: not enough information yet.
- $\text{MINMAX}(a) = \max(\text{MINMAX}(b), \text{MINMAX}(c), \text{MINMAX}(d)) = \max(???, ???, ???) \rightarrow \text{can't be established}$

MIN Player needs to explore b's subtree:

- MIN Player (at node b) has not seen any successor MINMAX values yet $\rightarrow \min \text{ MINMAX seen: } v = \infty$
- $v > \alpha_a$ ($\infty > -\infty$) \rightarrow we can keep exploring b's subtree

MinMax with α - β : Example

n | MINMAX(n)

MAX player state / move / turn

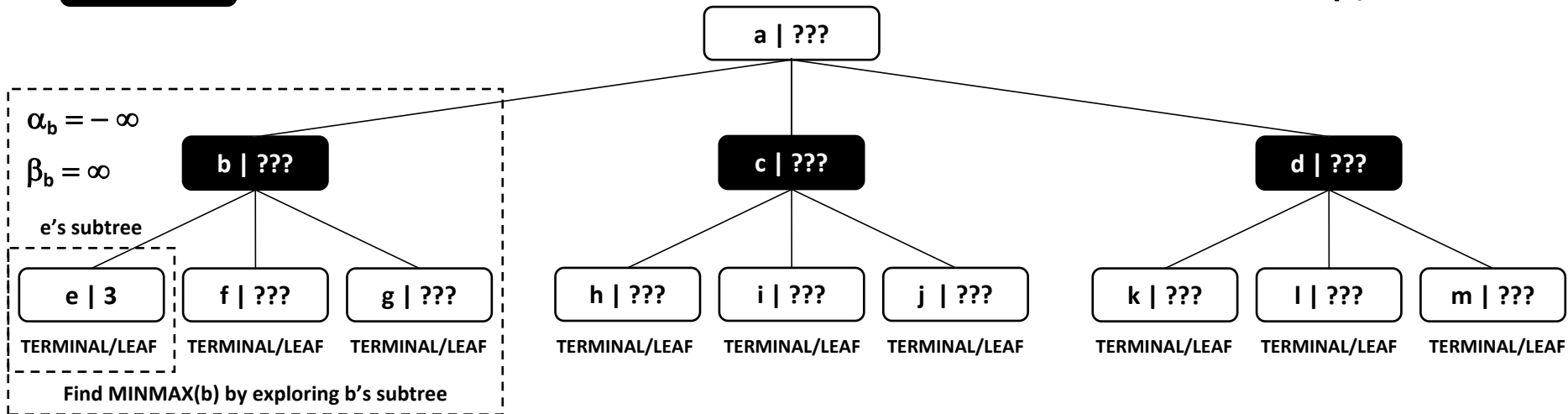
n | MINMAX(n)

MIN player state / move / turn

MINMAX(a) can be

at least: $\alpha_a = -\infty$

at most: $\beta_a = \infty$



MAX Player wants to maximize the utility of the game by choosing the right move from state a to one of three successor states: b, c, d. It will choose the state with maximum MAXMIN value. Currently:

- MINMAX(b) = UNKNOWN | MINMAX(c) = UNKNOWN | MINMAX(d) = UNKNOWN
 - MAX Player's decision: not enough information yet.
- MINMAX(a) = $\max(\text{MINMAX}(b), \text{MINMAX}(c), \text{MINMAX}(d)) = \max(???, ???, ???) \rightarrow$ can't be established

MIN Player needs to explore b's subtree:

- We need to analyze e's subtree
- Node e is a terminal node (Case 1) $\rightarrow \text{MINMAX}(e) = \text{UTILITY}(e) = 3$ | $v_2 = \text{MINMAX}(e) = 3$

MinMax with α - β : Example

n | MINMAX(n)

MAX player state / move / turn

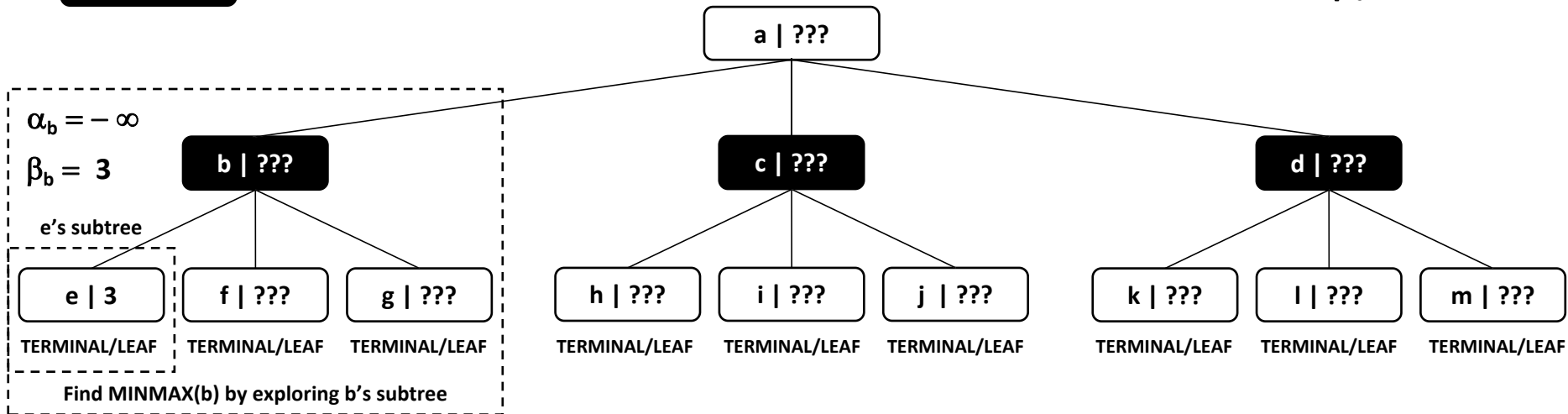
n | MINMAX(n)

MIN player state / move / turn

MINMAX(a) can be

at least: $\alpha_a = -\infty$

at most: $\beta_a = \infty$



MAX Player wants to maximize the utility of the game by choosing the right move from state a to one of three successor states: b, c, d. It will choose the state with maximum MAXMIN value. Currently:

- MINMAX(b) = UNKNOWN | MINMAX(c) = UNKNOWN | MINMAX(d) = UNKNOWN
 - MAX Player's decision: not enough information yet.
- MINMAX(a) = $\max(\text{MINMAX}(b), \text{MINMAX}(c), \text{MINMAX}(d)) = \max(???, ???, ???) \rightarrow$ can't be established

MIN Player needs to explore b's subtree:

- $v_2 < v$ ($3 < \infty$) $\rightarrow v = v_2 = 3 \rightarrow \beta_b = \min(\beta_b, v) = \min(\infty, 3) = 3$
- $v > \alpha_a$ ($3 > -\infty$) \rightarrow we can keep exploring b's subtree

MinMax with α - β : Example

n | MINMAX(n)

MAX player state / move / turn

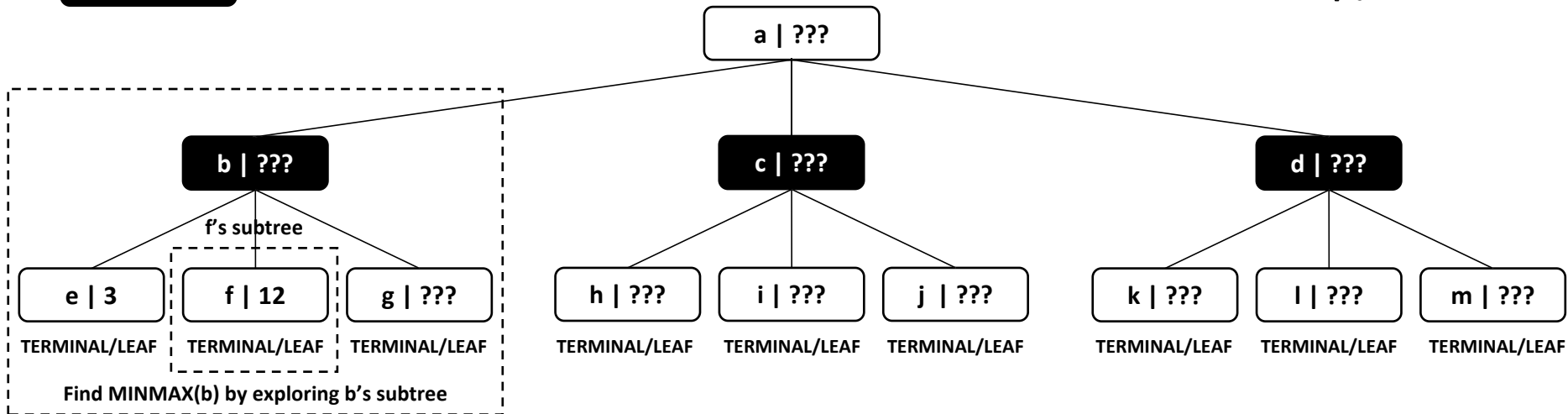
n | MINMAX(n)

MIN player state / move / turn

MINMAX(a) can be

at least: $\alpha_a = -\infty$

at most: $\beta_a = \infty$



MAX Player wants to maximize the utility of the game by choosing the right move from state **a** to one of three successor states: **b**, **c**, **d**. It will choose the state with maximum MAXMIN value. Currently:

- $\text{MINMAX}(b) = \text{UNKNOWN}$ | $\text{MINMAX}(c) = \text{UNKNOWN}$ | $\text{MINMAX}(d) = \text{UNKNOWN}$
 - **MAX Player's decision:** not enough information yet.
- $\text{MINMAX}(a) = \max(\text{MINMAX}(b), \text{MINMAX}(c), \text{MINMAX}(d)) = \max(???, ???, ???) \rightarrow \text{can't be established}$

MIN Player needs to explore **b's** subtree:

- We need to analyze **f's** subtree
- Node **f** is a terminal node (Case 1) $\rightarrow \text{MINMAX}(f) = \text{UTILITY}(f) = 12$ | $v_2 = \text{MINMAX}(f) = 12$

MinMax with α - β : Example

n | MINMAX(n)

MAX player state / move / turn

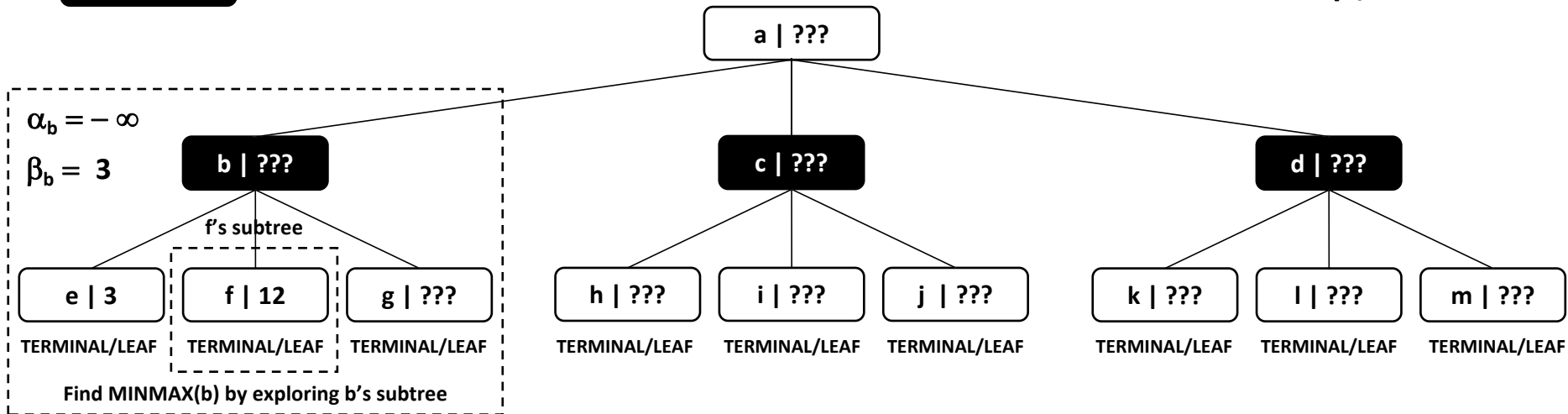
n | MINMAX(n)

MIN player state / move / turn

MINMAX(a) can be

at least: $\alpha_a = -\infty$

at most: $\beta_a = \infty$



MAX Player wants to maximize the utility of the game by choosing the right move from state a to one of three successor states: b, c, d. It will choose the state with maximum MAXMIN value. Currently:

- MINMAX(b) = UNKNOWN | MINMAX(c) = UNKNOWN | MINMAX(d) = UNKNOWN
 - MAX Player's decision: not enough information yet.
- MINMAX(a) = $\max(\text{MINMAX}(b), \text{MINMAX}(c), \text{MINMAX}(d)) = \max(???, ???, ???) \rightarrow$ can't be established

MIN Player needs to explore b's subtree:

- $v_2 > v$ ($12 > 3$) \rightarrow MINMAX(f) is not "better" than MINMAX(e) \rightarrow no changes
- $v > \alpha_a$ ($3 > -\infty$) \rightarrow we can keep exploring b's subtree

MinMax with α - β : Example

n | MINMAX(n)

MAX player state / move / turn

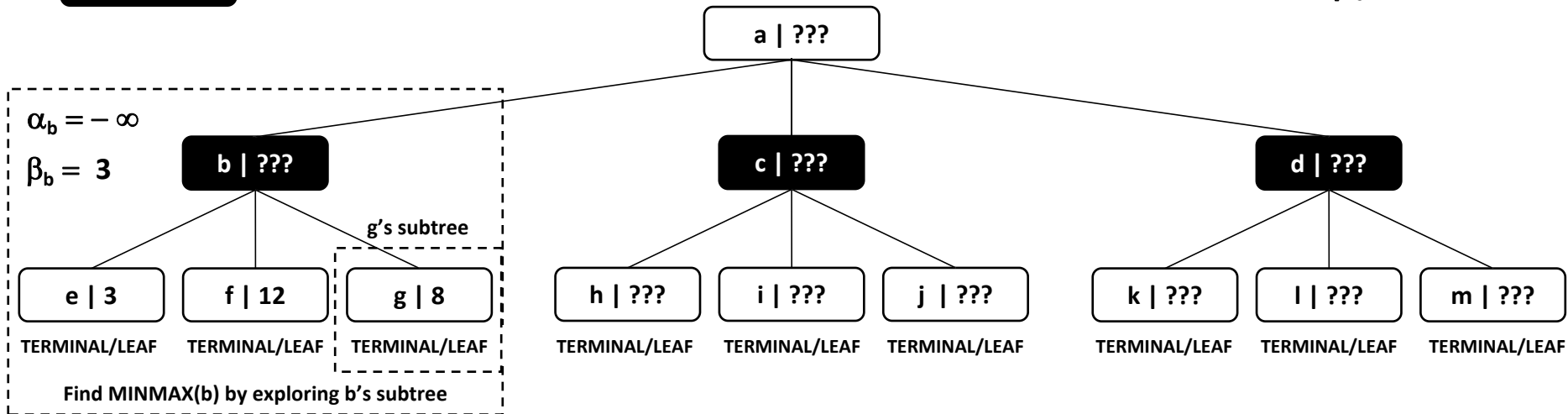
n | MINMAX(n)

MIN player state / move / turn

MINMAX(a) can be

at least: $\alpha_a = -\infty$

at most: $\beta_a = \infty$



MAX Player wants to maximize the utility of the game by choosing the right move from state a to one of three successor states: b, c, d. It will choose the state with maximum MAXMIN value. Currently:

- MINMAX(b) = UNKNOWN | MINMAX(c) = UNKNOWN | MINMAX(d) = UNKNOWN
 - MAX Player's decision: not enough information yet.
- MINMAX(a) = $\max(\text{MINMAX}(b), \text{MINMAX}(c), \text{MINMAX}(d)) = \max(???, ???, ???) \rightarrow$ can't be established

MIN Player needs to explore b's subtree:

- We need to analyze g's subtree
- Node g is a terminal node (Case 1) $\rightarrow \text{MINMAX}(g) = \text{UTILITY}(g) = 8 \mid v_2 = \text{MINMAX}(g) = 8$

MinMax with α - β : Example

n | MINMAX(n)

MAX player state / move / turn

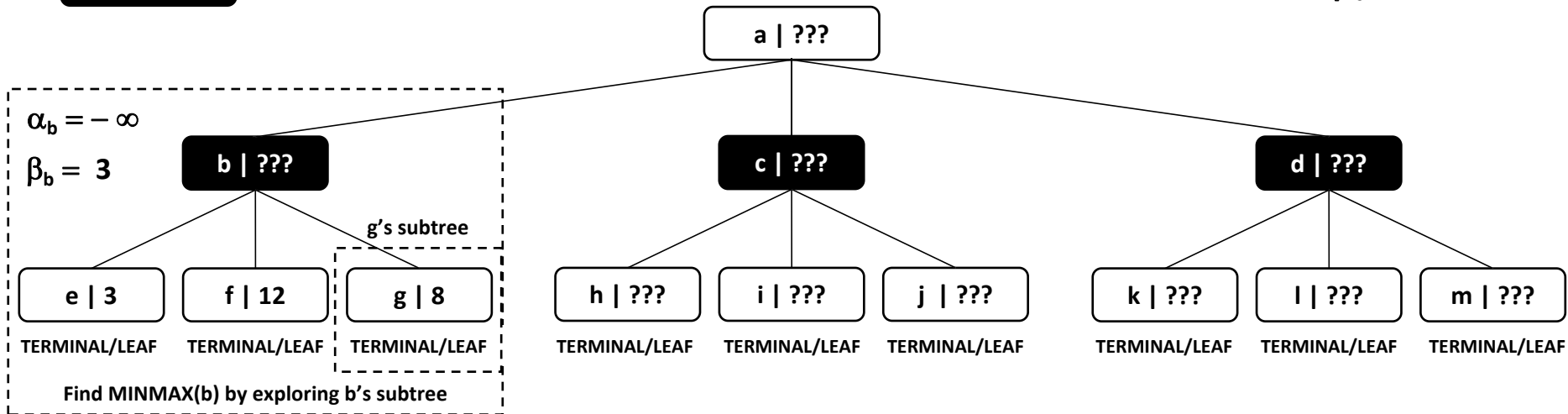
n | MINMAX(n)

MIN player state / move / turn

MINMAX(a) can be

at least: $\alpha_a = -\infty$

at most: $\beta_a = \infty$



MAX Player wants to maximize the utility of the game by choosing the right move from state a to one of three successor states: b, c, d. It will choose the state with maximum MAXMIN value. Currently:

- MINMAX(b) = UNKNOWN | MINMAX(c) = UNKNOWN | MINMAX(d) = UNKNOWN
 - MAX Player's decision: not enough information yet.
- MINMAX(a) = max(MINMAX(b), MINMAX(c), MINMAX(d)) = max(???, ???, ???) → can't be established

MIN Player needs to explore b's subtree:

- $v_2 > v$ ($8 > 3$) → MINMAX(g) is not "better" than MINMAX(e) → no changes
- $v > \alpha_a$ ($3 > -\infty$) → we could keep exploring b's subtree, but all b's subtrees are explored now

MinMax with α - β : Example

n | MINMAX(n)

MAX player state / move / turn

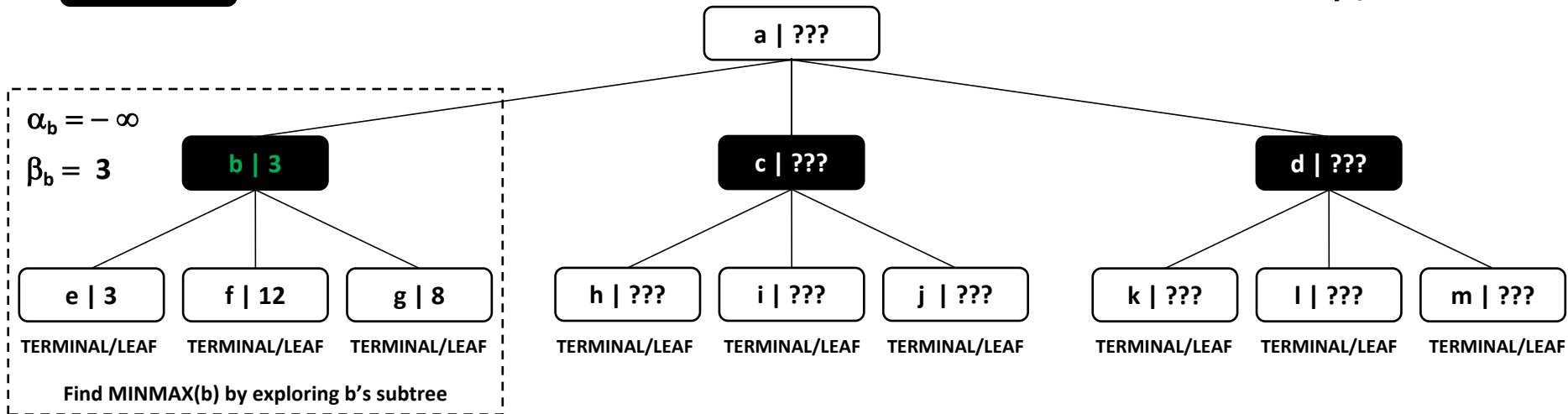
n | MINMAX(n)

MIN player state / move / turn

MINMAX(a) can be

at least: $\alpha_a = -\infty$

at most: $\beta_a = \infty$



MAX Player wants to maximize the utility of the game by choosing the right move from state a to one of three successor states: b, c, d. It will choose the state with maximum MAXMIN value. Currently:

- MINMAX(b) = UNKNOWN | MINMAX(c) = UNKNOWN | MINMAX(d) = UNKNOWN
 - MAX Player's decision: not enough information yet.
- MINMAX(a) = $\max(\text{MINMAX}(b), \text{MINMAX}(c), \text{MINMAX}(d)) = \max(???, ???, ???) \rightarrow$ can't be established

MIN Player explored entire b's subtree:

- MINMAX(b) = $\min(\text{MINMAX}(e), \text{MINMAX}(f), \text{MINMAX}(g)) = 3$ (Case 2)
- $v > \alpha_a$ ($3 > -\infty$) \rightarrow we could keep exploring b's subtree, but all b's subtrees are explored now

MinMax with α - β : Example

n | MINMAX(n)

MAX player state / move / turn

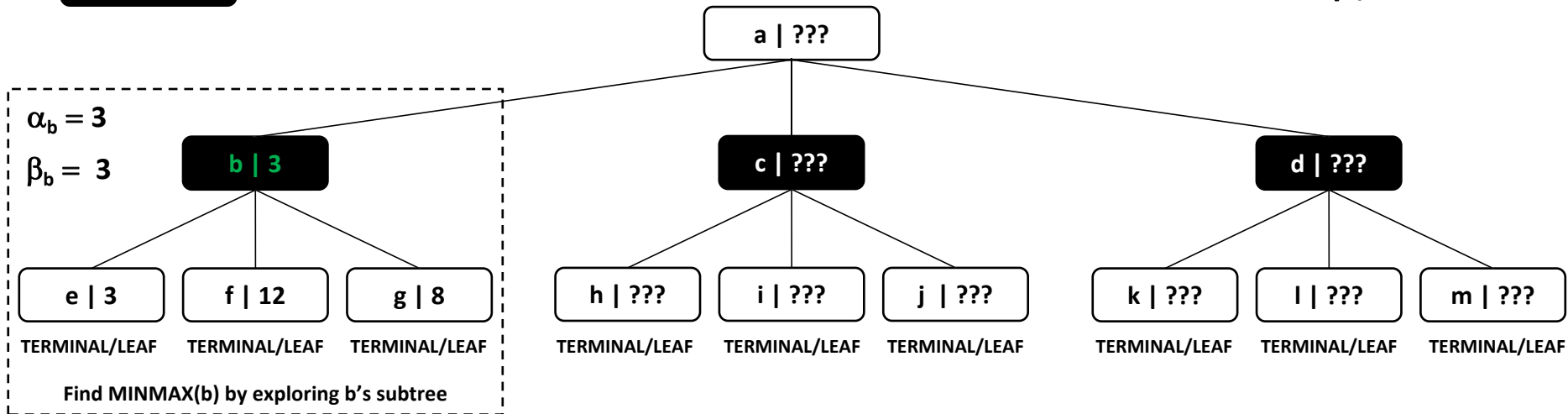
n | MINMAX(n)

MIN player state / move / turn

MINMAX(a) can be

at least: $\alpha_a = -\infty$

at most: $\beta_a = \infty$



MAX Player wants to maximize the utility of the game by choosing the right move from state a to one of three successor states: b, c, d. It will choose the state with maximum MAXMIN value. Currently:

- MINMAX(b) = UNKNOWN | MINMAX(c) = UNKNOWN | MINMAX(d) = UNKNOWN
 - MAX Player's decision: not enough information yet.
- MINMAX(a) = $\max(\text{MINMAX}(b), \text{MINMAX}(c), \text{MINMAX}(d)) = \max(???, ???, ???) \rightarrow$ can't be established

MIN Player explored entire b's subtree:

- We know the exact value of MINMAX(b) $\rightarrow \alpha_b = 3$

MinMax with α - β : Example

n | MINMAX(n)

MAX player state / move / turn

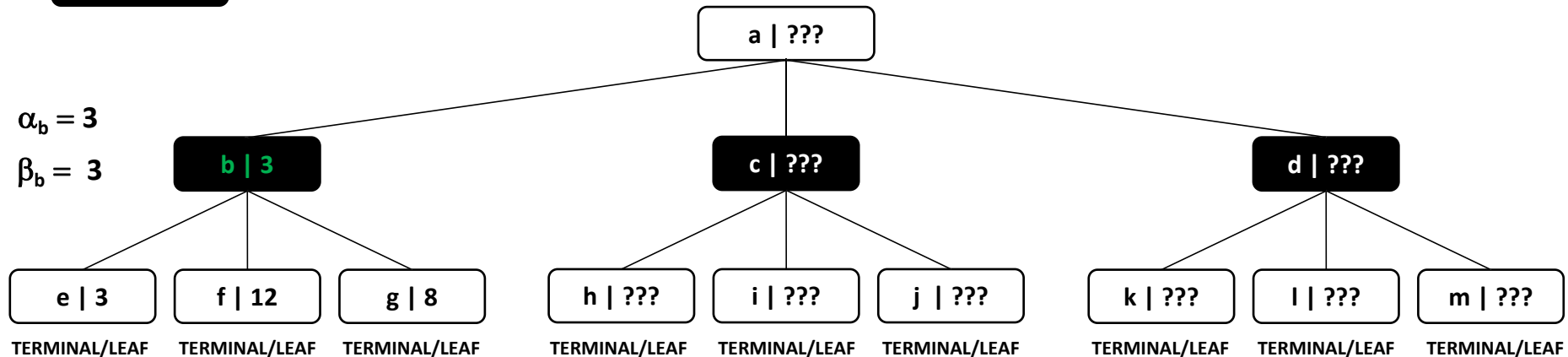
n | MINMAX(n)

MIN player state / move / turn

MINMAX(a) can be

at least: $\alpha_a = 3$

at most: $\beta_a = \infty$



MAX Player wants to maximize the utility of the game by choosing the right move from state a to one of three successor states: b, c, d. It will choose the state with maximum MAXMIN value. Currently:

- **MINMAX(b) = 3** | MINMAX(c) = UNKNOWN | MINMAX(d) = UNKNOWN
- MAX Player's decision: not enough information yet.

MINMAX(a) = $\max(\text{MINMAX}(b), \text{MINMAX}(c), \text{MINMAX}(d)) = \max(3, ???, ???) \rightarrow$ can't be established, **but**
MAX Player now knows that it will be AT LEAST 3 (3 OR HIGHER) $\rightarrow \alpha_a = 3$

MinMax with α - β : Example

n | MINMAX(n)

MAX player state / move / turn

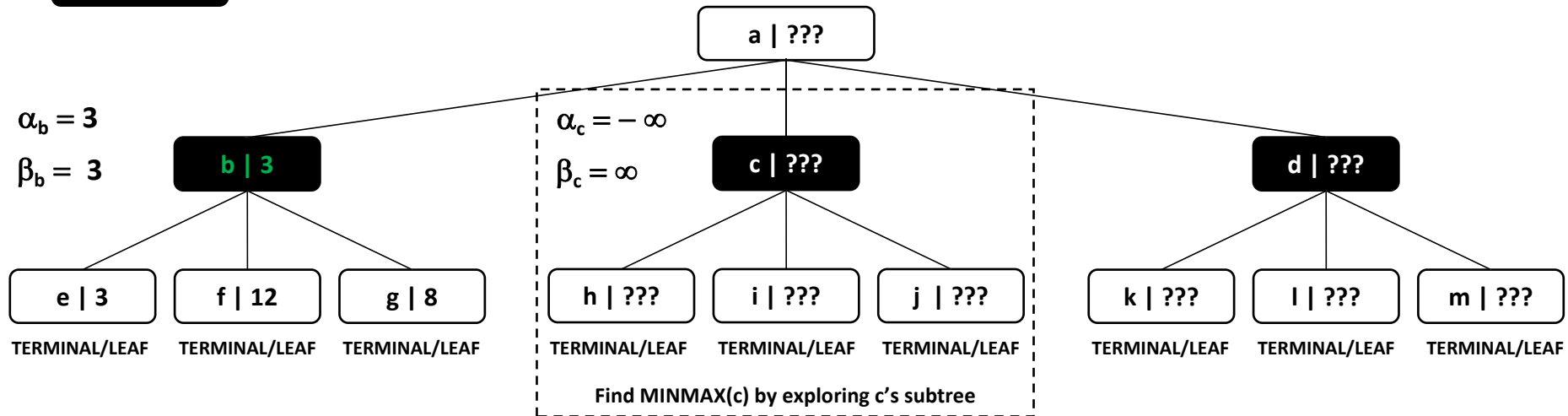
n | MINMAX(n)

MIN player state / move / turn

MINMAX(a) can be

at least: $\alpha_a = 3$

at most: $\beta_a = \infty$



MAX Player wants to maximize the utility of the game by choosing the right move from state a to one of three successor states: b, c, d. It will choose the state with maximum MAXMIN value. Currently:

- **MINMAX(b) = 3** | MINMAX(c) = UNKNOWN | MINMAX(d) = UNKNOWN
- MAX Player's decision: not enough information yet.

$\text{MINMAX}(a) = \max(3, \text{MINMAX}(c), \text{MINMAX}(d)) = \max(3, ???, ???) \rightarrow \text{can't be established}$

MIN Player needs to explore c's subtree:

- MIN Player (at node c) has not seen any successor MINMAX values yet $\rightarrow \min \text{ MINMAX seen: } v = \infty$
- $v > \alpha_a$ ($\infty > 3$) \rightarrow we can keep exploring c's subtree

MinMax with α - β : Example

n | MINMAX(n)

MAX player state / move / turn

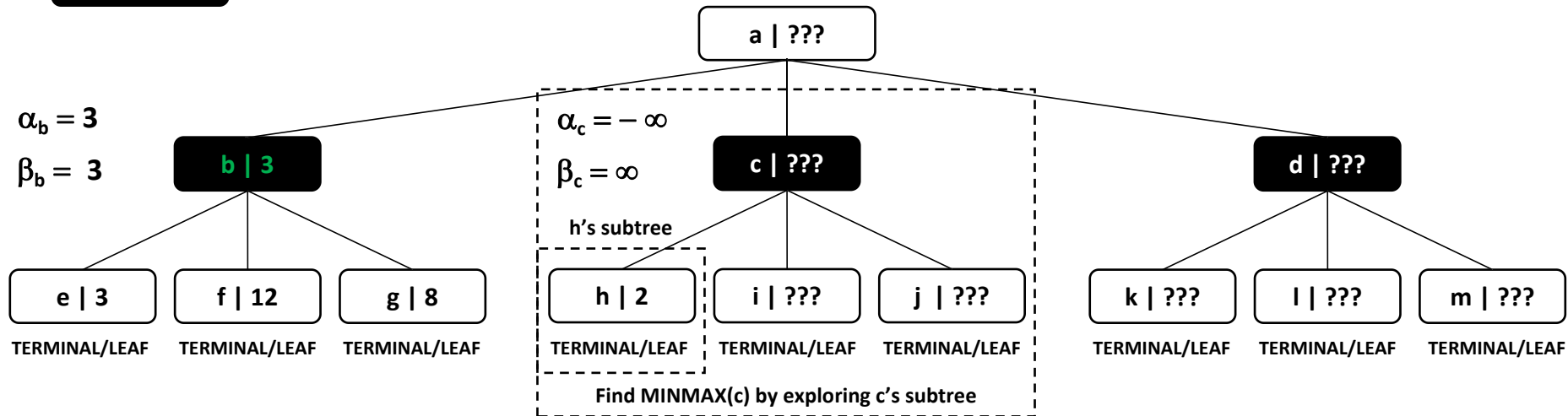
n | MINMAX(n)

MIN player state / move / turn

MINMAX(a) can be

at least: $\alpha_a = 3$

at most: $\beta_a = \infty$



MAX Player wants to maximize the utility of the game by choosing the right move from state a to one of three successor states: b, c, d. It will choose the state with maximum MAXMIN value. Currently:

- **MINMAX(b) = 3** | MINMAX(c) = UNKNOWN | MINMAX(d) = UNKNOWN
- MAX Player's decision: not enough information yet.
 $\text{MINMAX}(a) = \max(3, \text{MINMAX}(c), \text{MINMAX}(d)) = \max(3, ???, ???) \rightarrow \text{can't be established}$

MIN Player needs to explore c's subtree:

- We need to analyze h's subtree
- Node h is a terminal node (Case 1) $\rightarrow \text{MINMAX}(h) = \text{UTILITY}(h) = 2$ | $v_2 = \text{MINMAX}(h) = 2$

MinMax with α - β : Example

n | MINMAX(n)

MAX player state / move / turn

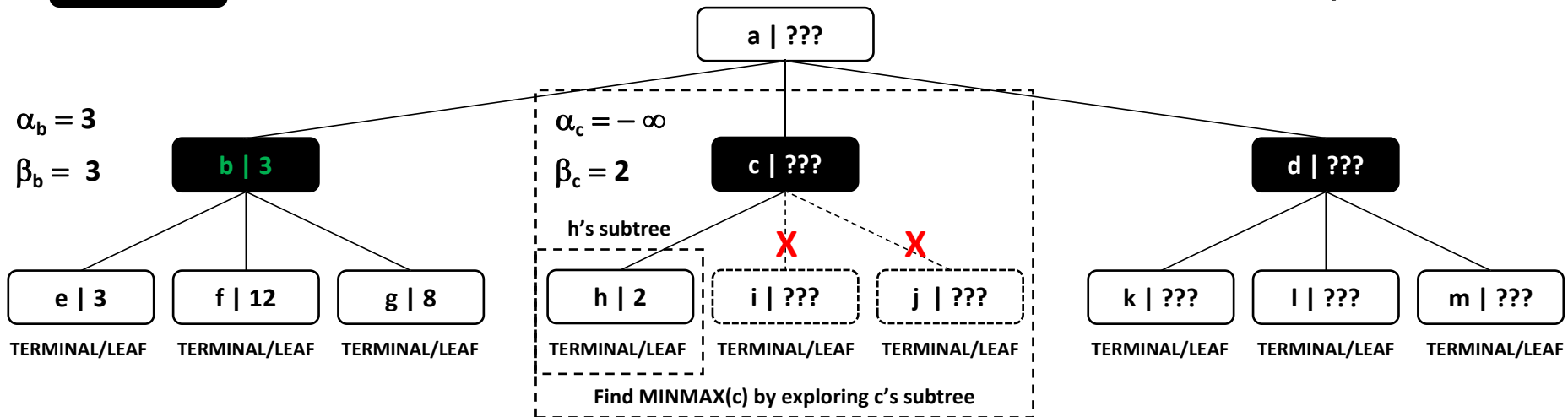
n | MINMAX(n)

MIN player state / move / turn

MINMAX(a) can be

at least: $\alpha_a = 3$

at most: $\beta_a = \infty$



MAX Player wants to maximize the utility of the game by choosing the right move from state a to one of three successor states: b, c, d. It will choose the state with maximum MAXMIN value. Currently:

- MINMAX(b) = 3 | MINMAX(c) = UNKNOWN | MINMAX(d) = UNKNOWN

- MAX Player's decision: not enough information yet.

$\text{MINMAX}(a) = \max(3, \text{MINMAX}(c), \text{MINMAX}(d)) = \max(3, ???, ???) \rightarrow \text{can't be established}$

MIN Player needs to explore c's subtree:

- $v_2 < v$ ($2 < \infty$) $\rightarrow v = v_2 = 2 \rightarrow \beta_c = \min(\beta_c, v) = \min(\infty, 2) = 2$

- $v < \alpha_a$ ($2 < 3$) \rightarrow we cannot keep exploring c's subtree \rightarrow prune remaining branches

MinMax with α - β : Example

n | MINMAX(n)

MAX player state / move / turn

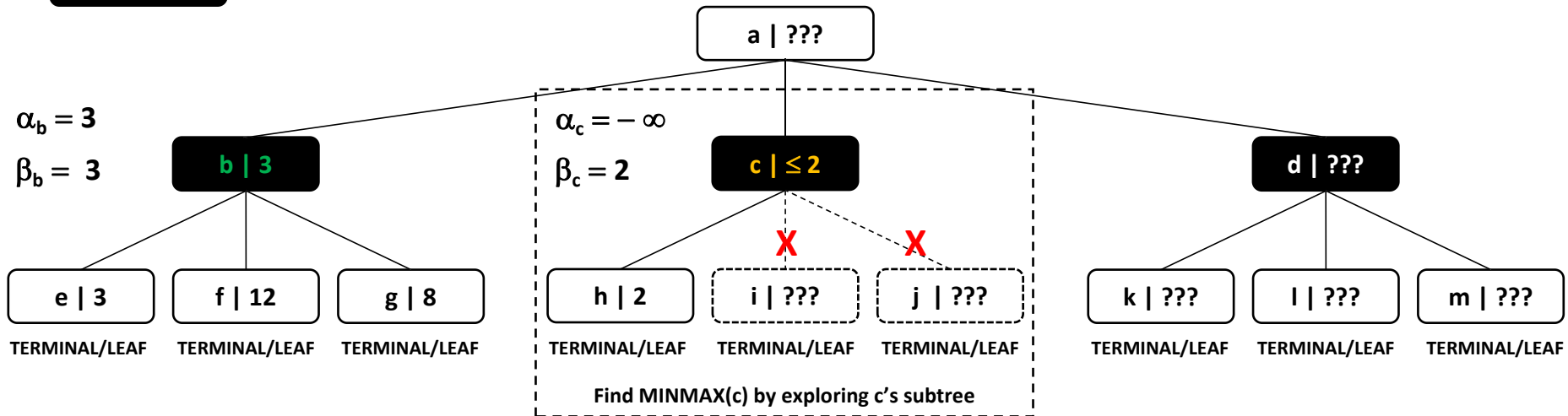
n | MINMAX(n)

MIN player state / move / turn

MINMAX(a) can be

at least: $\alpha_a = 3$

at most: $\beta_a = \infty$



MAX Player wants to maximize the utility of the game by choosing the right move from state a to one of three successor states: b, c, d. It will choose the state with maximum MAXMIN value. Currently:

- **MINMAX(b) = 3** | MINMAX(c) = UNKNOWN | MINMAX(d) = UNKNOWN
- MAX Player's decision: not enough information yet.
 $\text{MINMAX}(a) = \max(3, \text{MINMAX}(c), \text{MINMAX}(d)) = \max(3, ???, ???) \rightarrow \text{can't be established}$

MIN Player explored c's subtree as far as it was necessary:

- We know that **MINMAX(c) ≤ 2**

MinMax with α - β : Example

n | MINMAX(n)

MAX player state / move / turn

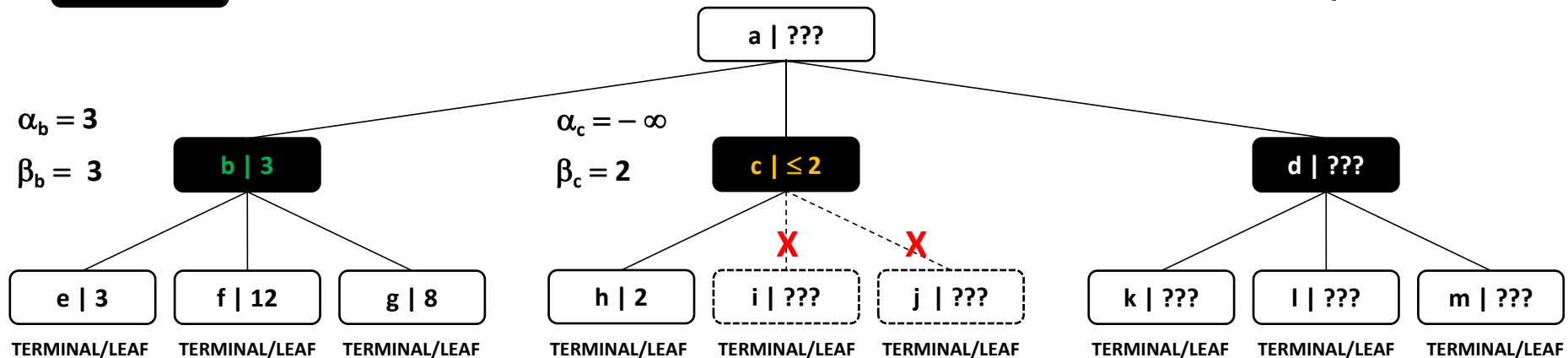
n | MINMAX(n)

MIN player state / move / turn

MINMAX(a) can be

at least: $\alpha_a = 3$

at most: $\beta_a = \infty$



MAX Player wants to maximize the utility of the game by choosing the right move from state a to one of three successor states: b, c, d. It will choose the state with maximum MAXMIN value. Currently:

- **MINMAX(b) = 3** | **MINMAX(c) = ≤ 2** | MINMAX(d) = UNKNOWN
- MAX Player's decision: not enough information yet.

$\text{MINMAX}(a) = \max(3, \text{MINMAX}(c), \text{MINMAX}(d)) = \max(3, \leq 2, ???) \rightarrow \text{can't be established}$

MinMax with α - β : Example

n | MINMAX(n)

MAX player state / move / turn

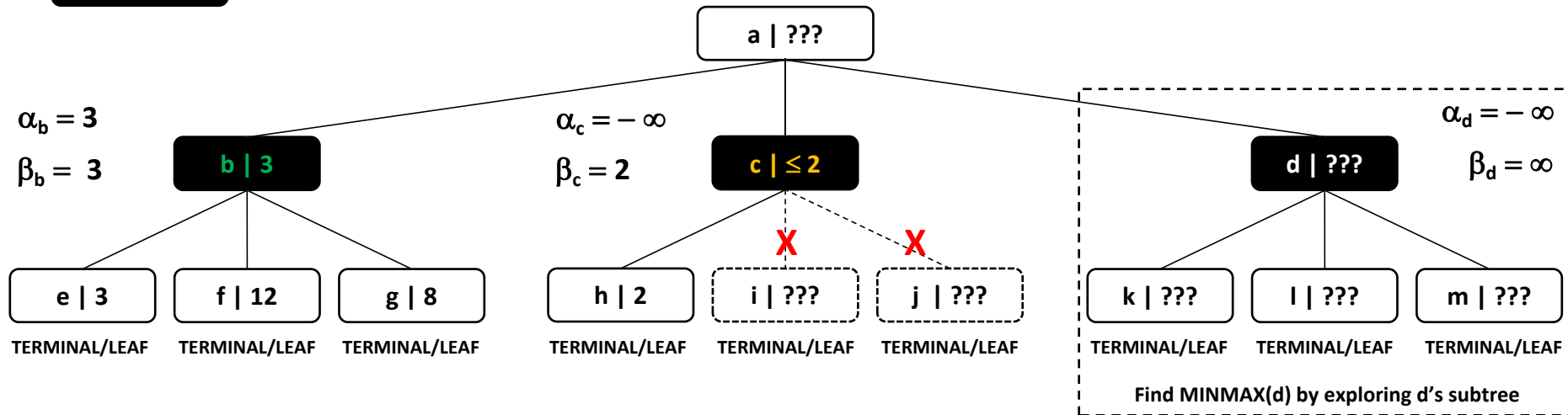
n | MINMAX(n)

MIN player state / move / turn

MINMAX(a) can be

at least: $\alpha_a = 3$

at most: $\beta_a = \infty$



MAX Player wants to maximize the utility of the game by choosing the right move from state a to one of three successor states: b, c, d. It will choose the state with maximum MAXMIN value. Currently:

- **MINMAX(b) = 3** | **MINMAX(c) = ≤ 2** | MINMAX(d) = UNKNOWN
- MAX Player's decision: not enough information yet.

$\text{MINMAX}(a) = \max(3, \text{MINMAX}(c), \text{MINMAX}(d)) = \max(3, \leq 2, ???) \rightarrow \text{can't be established}$

MIN Player needs to explore d's subtree:

- MIN Player (at node d) has not seen any successor MINMAX values yet $\rightarrow \min \text{MINMAX seen: } v = \infty$
- $v > \alpha_a$ ($\infty > 3$) \rightarrow we can keep exploring d's subtree

MinMax with α - β : Example

n | MINMAX(n)

MAX player state / move / turn

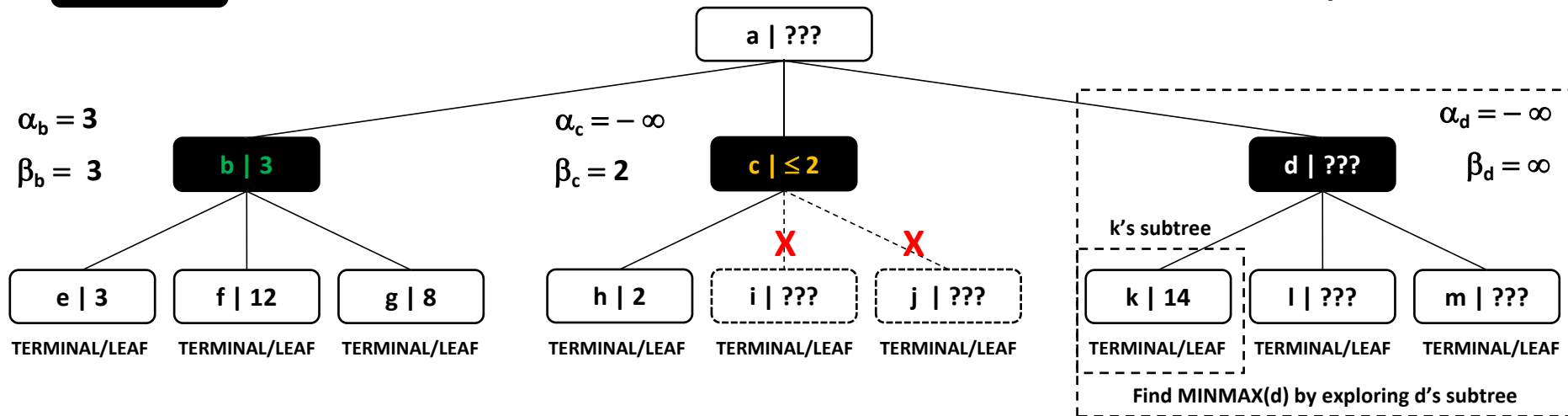
n | MINMAX(n)

MIN player state / move / turn

MINMAX(a) can be

at least: $\alpha_a = 3$

at most: $\beta_a = \infty$



MAX Player wants to maximize the utility of the game by choosing the right move from state a to one of three successor states: b, c, d. It will choose the state with maximum MAXMIN value. Currently:

- **MINMAX(b) = 3** | **MINMAX(c) = ≤ 2** | MINMAX(d) = UNKNOWN
- MAX Player's decision: not enough information yet.

$\text{MINMAX}(a) = \max(3, \text{MINMAX}(c), \text{MINMAX}(d)) = \max(3, \leq 2, ???) \rightarrow \text{can't be established}$

MIN Player needs to explore d's subtree:

- We need to analyze k's subtree
- Node k is a terminal node (Case 1) $\rightarrow \text{MINMAX}(k) = \text{UTILITY}(k) = 14$ | v2 = MINMAX(k) = 14

MinMax with α - β : Example

n | MINMAX(n)

MAX player state / move / turn

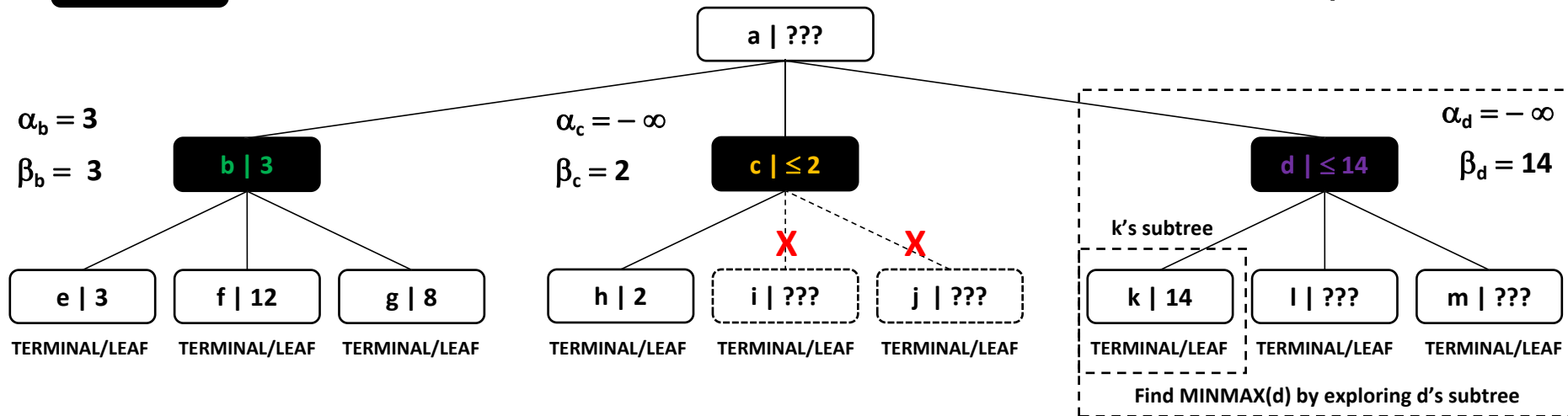
n | MINMAX(n)

MIN player state / move / turn

MINMAX(a) can be

at least: $\alpha_a = 3$

at most: $\beta_a = \infty$



MAX Player wants to maximize the utility of the game by choosing the right move from state a to one of three successor states: b, c, d. It will choose the state with maximum MAXMIN value. Currently:

- **MINMAX(b) = 3** | **MINMAX(c) = ≤ 2** | MINMAX(d) = UNKNOWN
- MAX Player's decision: not enough information yet.

$\text{MINMAX}(a) = \max(3, \text{MINMAX}(c), \text{MINMAX}(d)) = \max(3, \leq 2, ???) \rightarrow \text{can't be established}$

MIN Player needs to explore d's subtree:

- $v_2 < v$ ($14 < \infty$) $\rightarrow v = v_2 = 14 \rightarrow \beta_d = \min(\beta_d, v) = \min(\infty, 14) = 14$
- $v > \alpha_a$ ($14 > 3$) \rightarrow we can keep exploring d's subtree \rightarrow we also know that **MINMAX(d) ≤ 14**

MinMax with α - β : Example

n | MINMAX(n)

MAX player state / move / turn

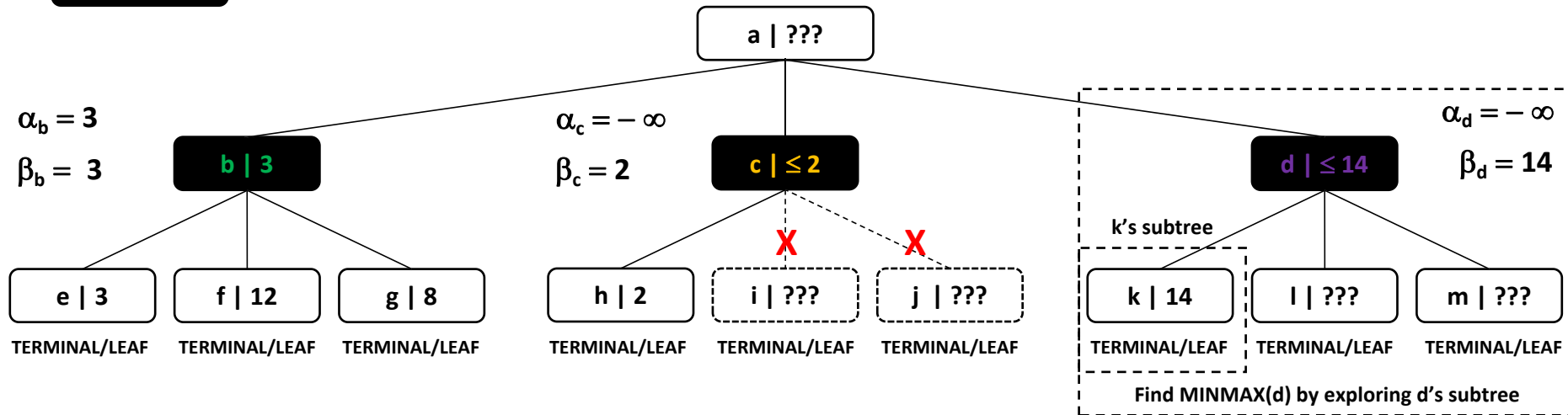
n | MINMAX(n)

MIN player state / move / turn

MINMAX(a) can be

at least: $\alpha_a = 3$

at most: $\beta_a = 14$



MAX Player wants to maximize the utility of the game by choosing the right move from state a to one of three successor states: b, c, d. It will choose the state with maximum MAXMIN value. Currently:

- **MINMAX(b) = 3** | **MINMAX(c) = ≤ 2** | **MINMAX(d) = ≤ 14**

- MAX Player's decision: not enough information yet.

$\text{MINMAX}(a) = \max(3, \text{MINMAX}(c), \text{MINMAX}(d)) = \max(3, \leq 2, \leq 14) \rightarrow \text{can't be established}$

MIN Player needs to explore d's subtree:

- we know that **MINMAX(d) ≤ 14** \rightarrow this tells us that $\text{MINMAX}(a)$ cannot be $> 14 \rightarrow \beta_a = 14$

MinMax with α - β : Example

n | MINMAX(n)

MAX player state / move / turn

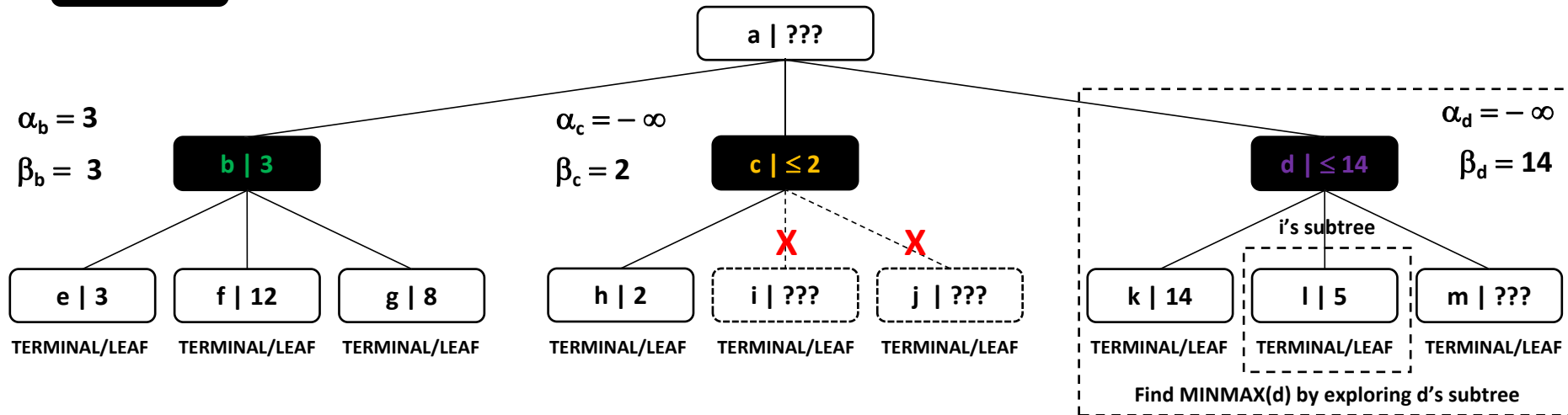
n | MINMAX(n)

MIN player state / move / turn

MINMAX(a) can be

at least: $\alpha_a = 3$

at most: $\beta_a = 14$



MAX Player wants to maximize the utility of the game by choosing the right move from state a to one of three successor states: b, c, d. It will choose the state with maximum MAXMIN value. Currently:

▪ $\text{MINMAX}(b) = 3$ | $\text{MINMAX}(c) = \leq 2$ | $\text{MINMAX}(d) = \leq 14$

▪ MAX Player's decision: not enough information yet.

$\text{MINMAX}(a) = \max(3, \text{MINMAX}(c), \text{MINMAX}(d)) = \max(3, \leq 2, \leq 14) \rightarrow \text{can't be established}$

MIN Player needs to explore d's subtree:

▪ We need to analyze l's subtree

▪ Node l is a terminal node (Case 1) $\rightarrow \text{MINMAX}(l) = \text{UTILITY}(l) = 5$ | $v_2 = \text{MINMAX}(l) = 5$

MinMax with α - β : Example

n | MINMAX(n)

MAX player state / move / turn

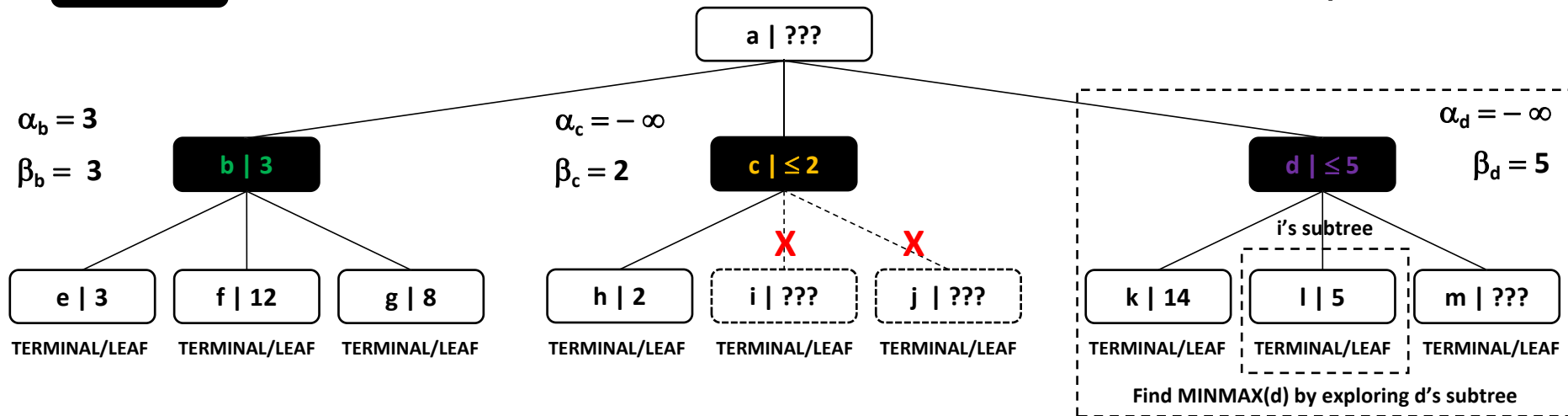
n | MINMAX(n)

MIN player state / move / turn

MINMAX(a) can be

at least: $\alpha_a = 3$

at most: $\beta_a = 14$



MAX Player wants to maximize the utility of the game by choosing the right move from state a to one of three successor states: b, c, d. It will choose the state with maximum MAXMIN value. Currently:

▪ $\text{MINMAX}(b) = 3$ | $\text{MINMAX}(c) = \leq 2$ | $\text{MINMAX}(d) = \leq 14$

▪ MAX Player's decision: not enough information yet.

$\text{MINMAX}(a) = \max(3, \text{MINMAX}(c), \text{MINMAX}(d)) = \max(3, \leq 2, \leq 14) \rightarrow \text{can't be established}$

MIN Player needs to explore d's subtree:

▪ $v_2 < v$ ($5 < 14$) $\rightarrow v = v_2 = 5 \rightarrow \beta_d = \min(\beta_d, v) = \min(\infty, 5) = 5$

▪ $v > \alpha_a$ ($5 > 3$) \rightarrow we can keep exploring d's subtree \rightarrow we also know that $\text{MINMAX}(d) \leq 5$

MinMax with α - β : Example

n | MINMAX(n)

MAX player state / move / turn

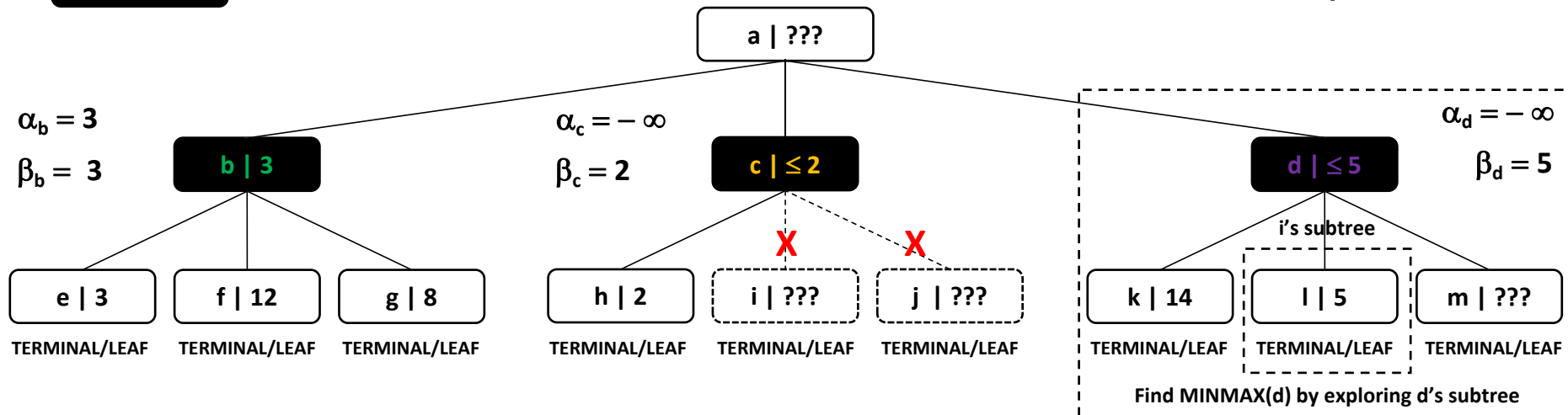
n | MINMAX(n)

MIN player state / move / turn

MINMAX(a) can be

at least: $\alpha_a = 3$

at most: $\beta_a = 5$



MAX Player wants to maximize the utility of the game by choosing the right move from state a to one of three successor states: b, c, d. It will choose the state with maximum MAXMIN value. Currently:

- $\text{MINMAX}(b) = 3$ | $\text{MINMAX}(c) = \leq 2$ | $\text{MINMAX}(d) = \leq 5$
- MAX Player's decision: not enough information yet.

$$\text{MINMAX}(a) = \max(3, \text{MINMAX}(c), \text{MINMAX}(d)) = \max(3, \leq 2, \leq 5) \rightarrow \text{can't be established}$$

MIN Player needs to explore d's subtree:

- we know that $\text{MINMAX}(d) \leq 5 \rightarrow$ this tells us that $\text{MINMAX}(a)$ cannot be $> 5 \rightarrow \beta_a = 5$

MinMax with α - β : Example

n | MINMAX(n)

MAX player state / move / turn

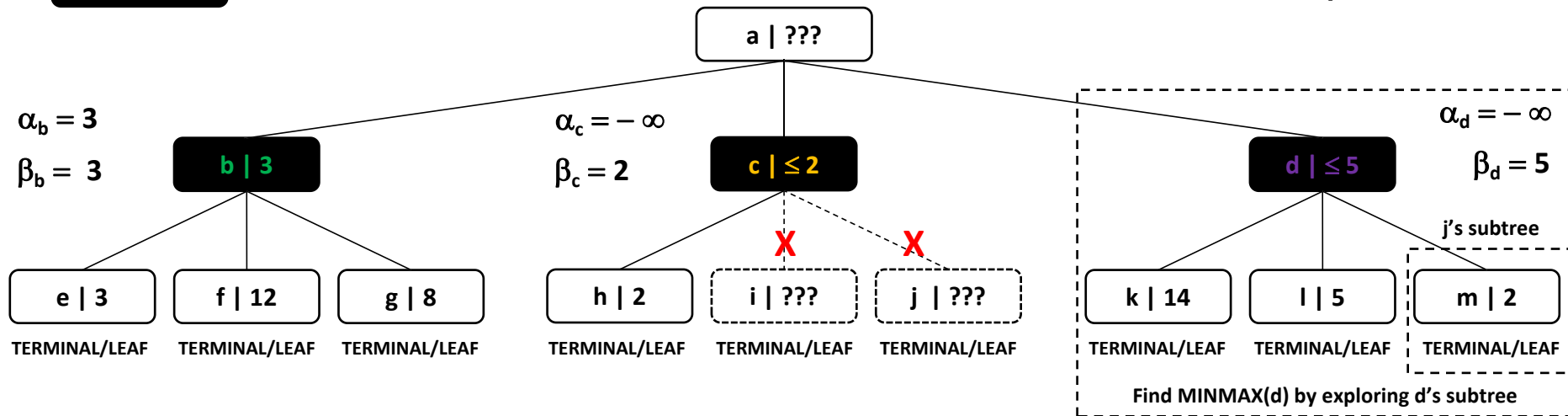
n | MINMAX(n)

MIN player state / move / turn

MINMAX(a) can be

at least: $\alpha_a = 3$

at most: $\beta_a = 5$



MAX Player wants to maximize the utility of the game by choosing the right move from state a to one of three successor states: b, c, d. It will choose the state with maximum MAXMIN value. Currently:

- $\text{MINMAX}(b) = 3$ | $\text{MINMAX}(c) = \leq 2$ | $\text{MINMAX}(d) = \leq 5$
- MAX Player's decision: not enough information yet.

$$\text{MINMAX}(a) = \max(3, \text{MINMAX}(c), \text{MINMAX}(d)) = \max(3, \leq 2, \leq 5) \rightarrow \text{can't be established}$$

MIN Player needs to explore d's subtree:

- We need to analyze m's subtree
- Node m is a terminal node (Case 1) $\rightarrow \text{MINMAX}(m) = \text{UTILITY}(m) = 2$ | $v_2 = \text{MINMAX}(m) = 2$

MinMax with α - β : Example

n | MINMAX(n)

MAX player state / move / turn

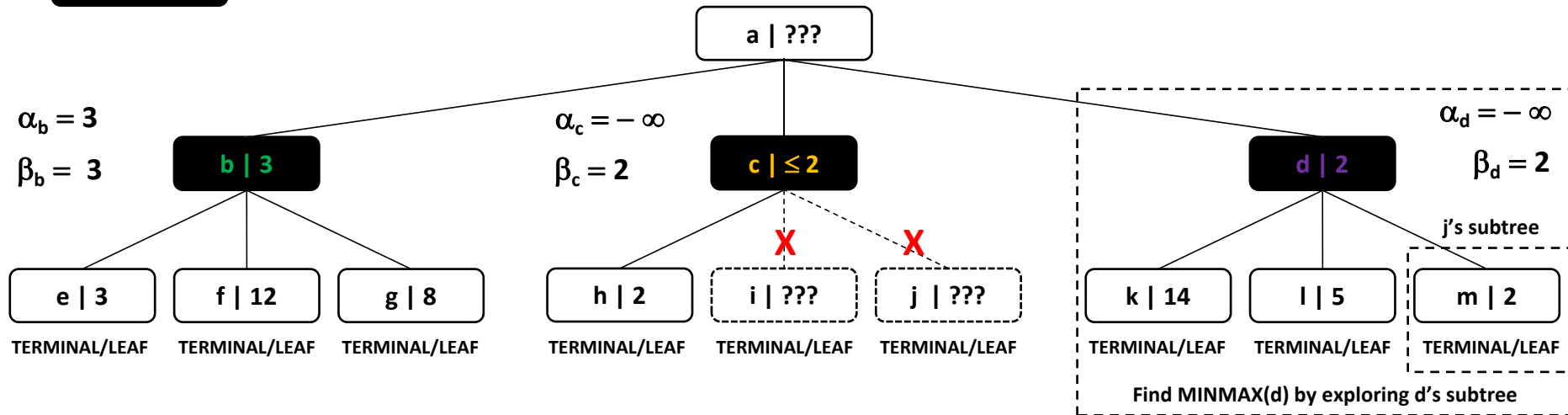
n | MINMAX(n)

MIN player state / move / turn

MINMAX(a) can be

at least: $\alpha_a = 3$

at most: $\beta_a = 5$



MAX Player wants to maximize the utility of the game by choosing the right move from state a to one of three successor states: b, c, d. It will choose the state with maximum MAXMIN value. Currently:

- $\text{MINMAX}(b) = 3$ | $\text{MINMAX}(c) = \leq 2$ | $\text{MINMAX}(d) = \leq 5$
- MAX Player's decision: not enough information yet.

$$\text{MINMAX}(a) = \max(3, \text{MINMAX}(c), \text{MINMAX}(d)) = \max(3, \leq 2, \leq 5) \rightarrow \text{can't be established}$$

MIN Player needs to explore d's subtree:

- $v_2 < v$ ($2 < 5$) $\rightarrow v = v_2 = 2 \rightarrow \beta_d = \min(\beta_d, v) = \min(\infty, 2) = 2$
- $v < \alpha_a$ ($2 < 3$) \rightarrow we cannot keep exploring d's subtree \rightarrow we also know that $\text{MINMAX}(d) = 2$

MinMax with α - β : Example

n | MINMAX(n)

MAX player state / move / turn

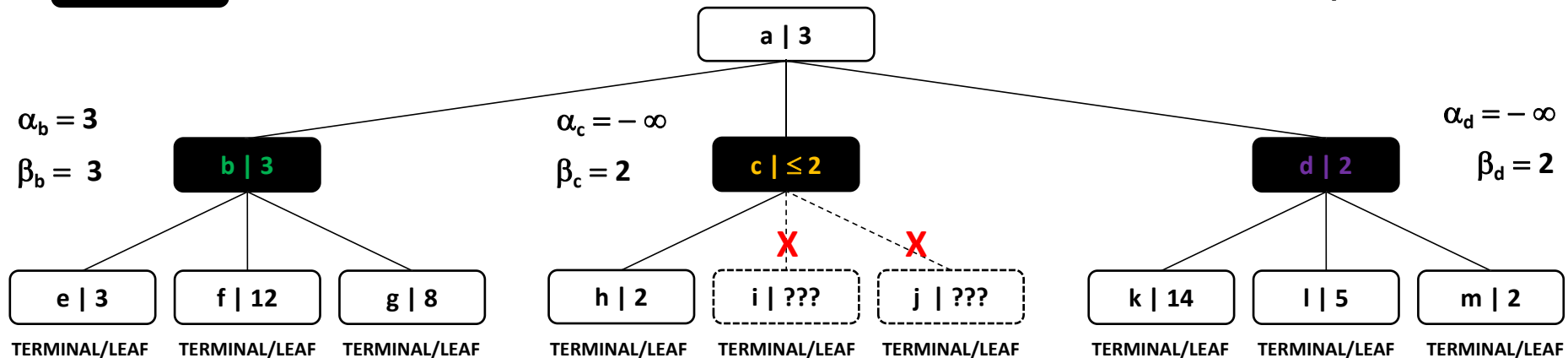
n | MINMAX(n)

MIN player state / move / turn

MINMAX(a) can be

at least: $\alpha_a = 3$

at most: $\beta_a = 3$



Find MINMAX(d) by exploring d's subtree

MAX Player wants to maximize the utility of the game by choosing the right move from state a to one of three successor states: b, c, d. It will choose the state with maximum MAXMIN value. Currently:

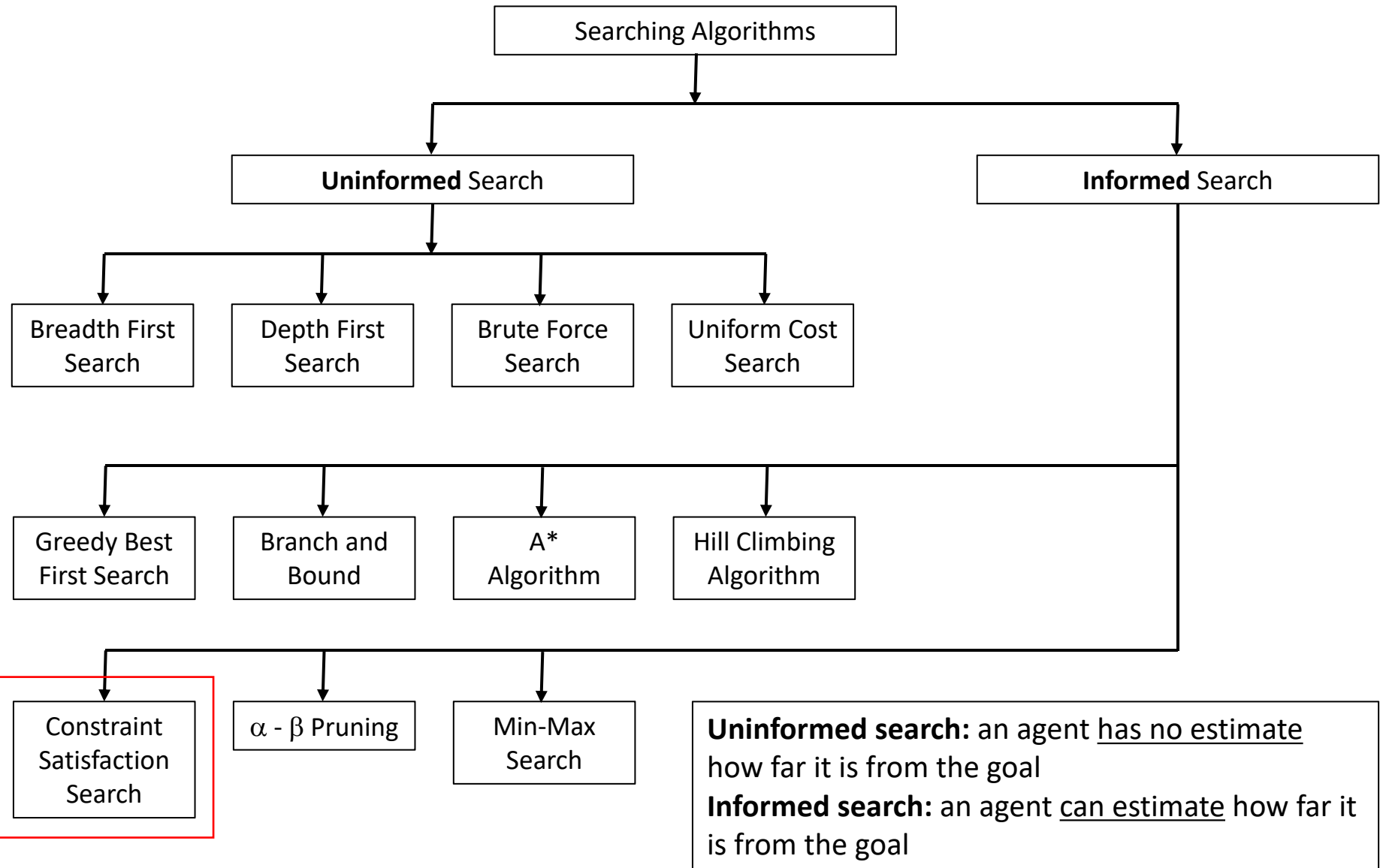
▪ **MINMAX(b) = 3** | **MINMAX(c) = ≤ 2** | **MINMAX(d) = 2**

▪ MAX Player's decision: **choose move b, because:**

$$\text{MINMAX}(a) = \max(3, \text{MINMAX}(c), \text{MINMAX}(d)) = \max(3, \leq 2, 2) = 3$$

▪ Since we know MINMAX(a), we can update β_a for completeness $\rightarrow \beta_a = 3$

Selected Searching Algorithms



Constraint Satisfaction Problem

A Constraint Satisfaction Problem (CSP) consists of three components:

- **a set of variables $X = \{X_1, \dots, X_n\}$**
- **a set of domains $D = \{D_1, \dots, D_n\}$**
- **a set of constraints C that specify allowable combinations of values**
- **A domain D_i is a set of allowable values $\{v_1, \dots, v_k\}$ for variable X_i**
- **A constraint C_j is a $\langle \text{scope}, \text{relation} \rangle$ pair, for example $\langle (X_1, X_2), X_1 > X_2 \rangle$**

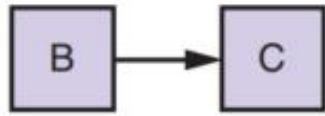
Constraint Satisfaction Problem

The goal is to **find an assignment** (variable = value):

$$\{X_1 = v_1, \dots, X_n = v_n\}$$

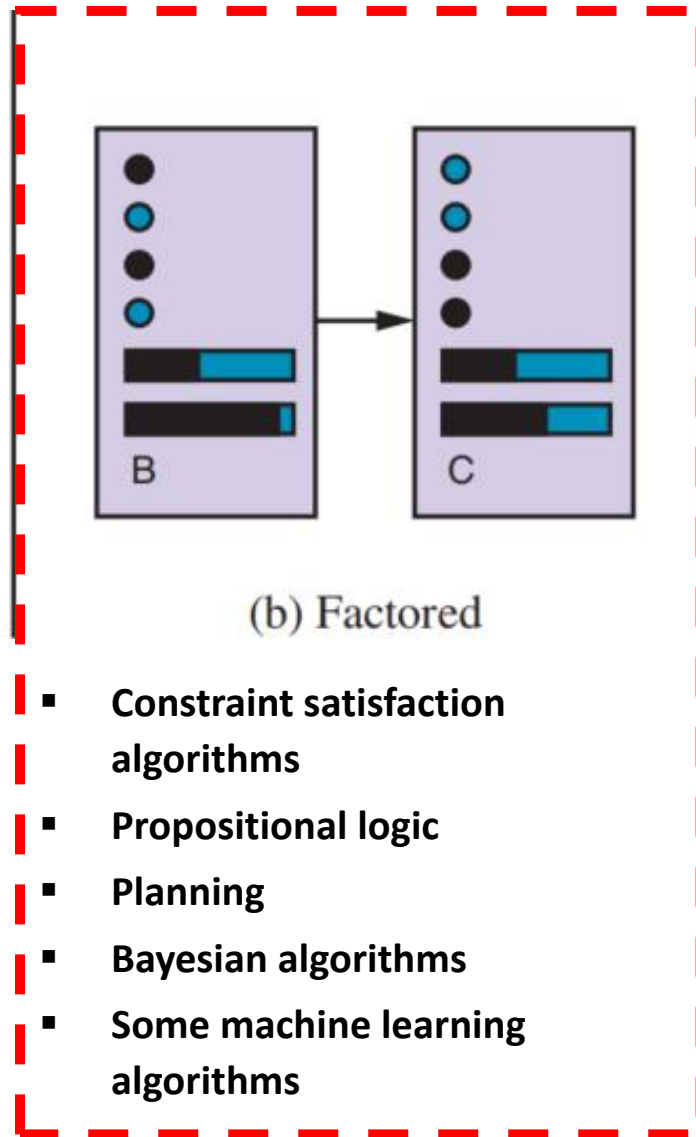
- If NO constraints violated: **consistent** assignment
- If ALL variables have a value: **complete** assignment
- If SOME variables have NO value: **partial** assignment
- SOLUTION: **consistent** and **complete** assignment
- PARTIAL SOLUTION: **consistent** and **partial** assignment

State Representations



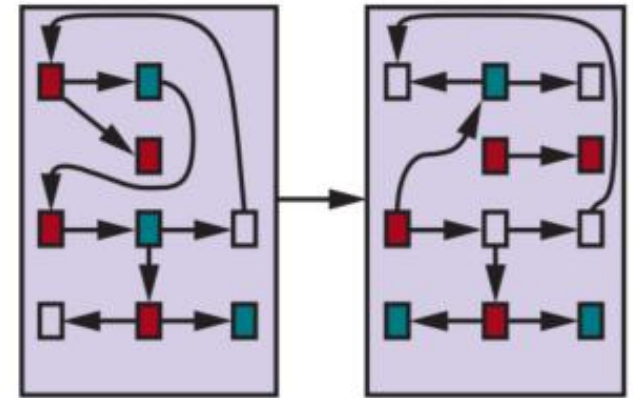
(a) Atomic

- Searching
- Hidden Markov models
- Markov decision process
- Finite state machines



(b) Factored

- Constraint satisfaction algorithms
- Propositional logic
- Planning
- Bayesian algorithms
- Some machine learning algorithms



(c) Structured

- Relational database algorithms
- First-order logic
- First-order probability models
- Natural language understanding (some)

CSP Example: Map Coloring

Problem:



Variables:

$X = \{WA, NT, Q, NSW, V, SA, T\}$

Variable Domains:

$D_{WA} = \{RED, GREEN, BLUE\}$

$D_{NT} = \{RED, GREEN, BLUE\}$

$D_Q = \{RED, GREEN, BLUE\}$

$D_{NSW} = \{RED, GREEN, BLUE\}$

$D_V = \{RED, GREEN, BLUE\}$

$D_{SA} = \{RED, GREEN, BLUE\}$

$D_T = \{RED, GREEN, BLUE\}$

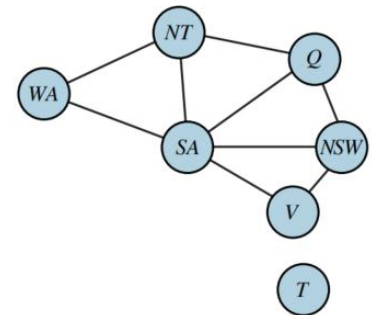
Color this map in a way that no two neighbors have same color

Constraints (Rules):

- Neighboring regions have to have **DISTINCT** colors:

$CONSTRAINTS = C = \{SA \neq WA, SA \neq NT, SA \neq Q, SA \neq NSW, SA \neq V, WA \neq NT, NT \neq Q, Q \neq NSW, NSW \neq V\}$

Constraint Graph:



CSP Example: Sudoku (3x3 for now)

Problem:

$x_{1,1}$	$x_{1,2}$	$x_{1,3}$
$x_{2,1}$	$x_{2,2}$	$x_{2,3}$
$x_{3,1}$	$x_{3,2}$	$x_{3,3}$

Variables:

$X = \{x_{1,1}, x_{1,2}, x_{1,3}, x_{2,1}, x_{2,2}, x_{2,3}, x_{3,1}, x_{3,2}, x_{3,3}\}$

Variable Domains:

$D_{x_{1,1}} = \{1, 2, 3, 4, 5, 6, 7, 8, 9\}$

$D_{x_{1,2}} = \{1, 2, 3, 4, 5, 6, 7, 8, 9\}$

$D_{x_{1,3}} = \{1, 2, 3, 4, 5, 6, 7, 8, 9\}$

$D_{x_{2,1}} = \{1, 2, 3, 4, 5, 6, 7, 8, 9\}$

$D_{x_{2,2}} = \{1, 2, 3, 4, 5, 6, 7, 8, 9\}$

$D_{x_{2,3}} = \{1, 2, 3, 4, 5, 6, 7, 8, 9\}$

$D_{x_{3,1}} = \{1, 2, 3, 4, 5, 6, 7, 8, 9\}$

$D_{x_{3,2}} = \{1, 2, 3, 4, 5, 6, 7, 8, 9\}$

$D_{x_{3,3}} = \{1, 2, 3, 4, 5, 6, 7, 8, 9\}$

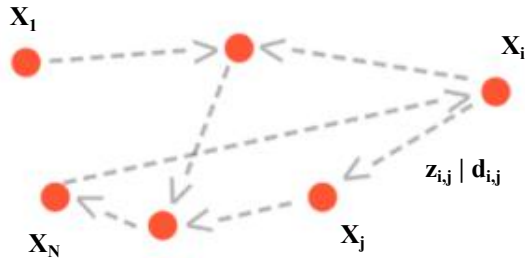
Constraints (Rules):

- Each value $\{1, 2, 3, 4, 5, 6, 7, 8, 9\}$ can appear EXACTLY once:

$CONSTRAINTS = C = \{x_{1,1} \neq x_{1,2}, x_{1,1} \neq x_{1,3}, x_{1,1} \neq x_{2,1}, x_{1,1} \neq x_{2,2}, x_{1,1} \neq x_{2,3}, x_{1,2} \neq x_{1,3}, x_{1,2} \neq x_{2,1}, x_{1,2} \neq x_{2,2}, x_{1,2} \neq x_{2,3}, x_{1,2} \neq x_{3,1}, x_{1,2} \neq x_{3,2}, x_{1,3} \neq x_{2,1}, x_{1,3} \neq x_{2,2}, x_{1,3} \neq x_{2,3}, x_{1,3} \neq x_{3,1}, x_{1,3} \neq x_{3,2}, x_{1,3} \neq x_{3,3}, x_{2,1} \neq x_{2,2}, x_{2,1} \neq x_{2,3}, x_{2,1} \neq x_{3,1}, x_{2,1} \neq x_{3,2}, x_{2,1} \neq x_{3,3}, x_{2,2} \neq x_{2,3}, x_{2,2} \neq x_{3,1}, x_{2,2} \neq x_{3,2}, x_{2,2} \neq x_{3,3}, x_{2,3} \neq x_{3,1}, x_{2,3} \neq x_{3,2}, x_{2,3} \neq x_{3,3}, x_{3,1} \neq x_{3,2}, x_{3,1} \neq x_{3,3}, x_{3,2} \neq x_{3,3}\}$

CSP Example: Traveling Salesman

Problem:



Variables:

$$Z = \{z_{1,2}, z_{1,3}, \dots, z_{N-1,N}\}$$

$$D = \{d_{1,2}, d_{1,3}, \dots, d_{N-1,N}\}$$

Variable Domains:

$$D_{z_{i,j}} = \{\text{traveled}, \text{notTraveled}\}$$

or better:

$$D_{z_{i,j}} = \{1, 0\}$$

$$D_{d_{i,j}} = \mathbb{R}_+$$

There are:

- N cities (vertices)
- $N(N-1)$ links (edges)
- Each link has some positive cost d
- Total path (tour) cost is $COST$

Constraints (Rules):

- Exit each city EXACTLY once:

$$\sum_{j=1}^N z_{i,j} = 1$$

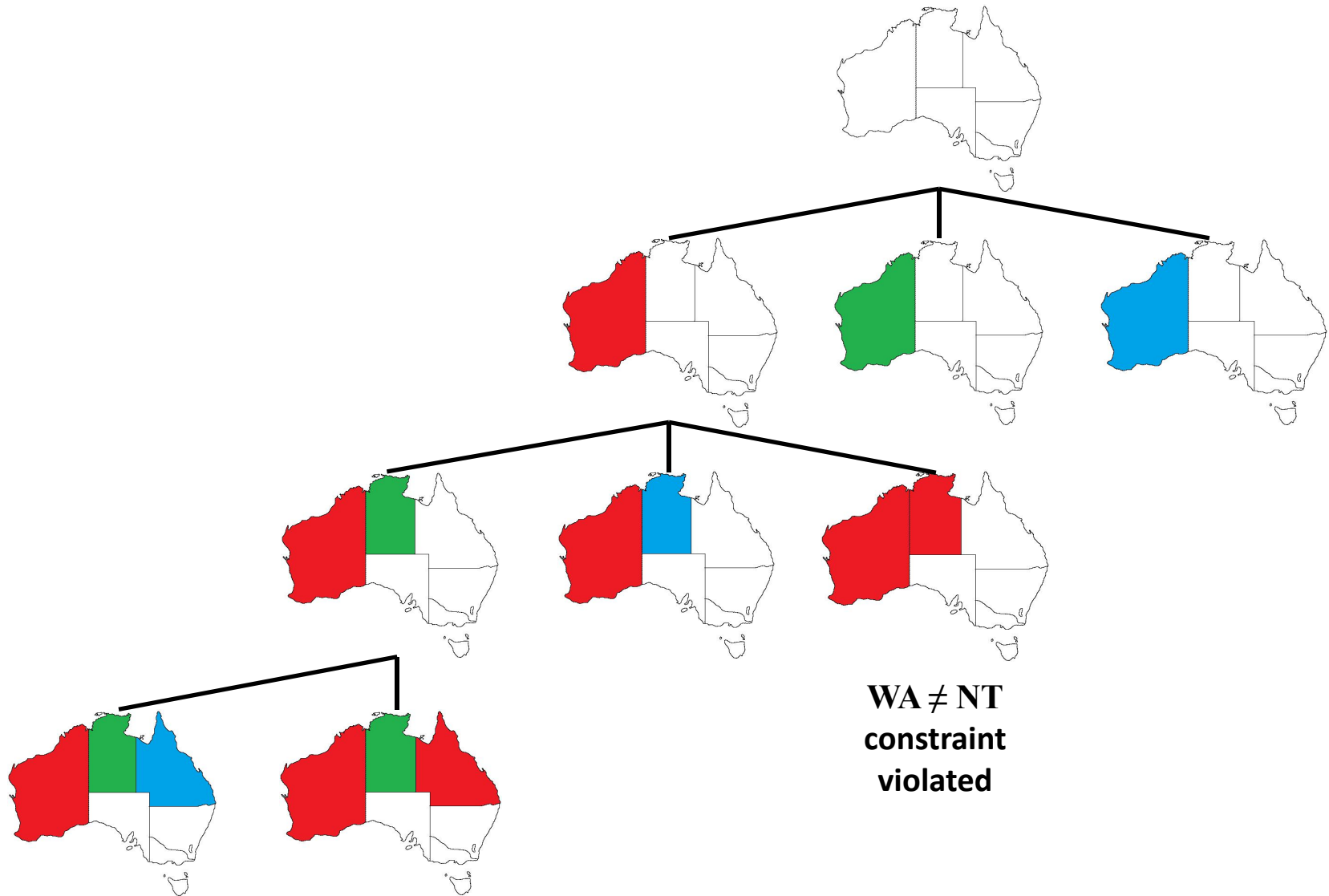
- Enter each city EXACTLY once:

$$\sum_{i=1}^N z_{i,j} = 1$$

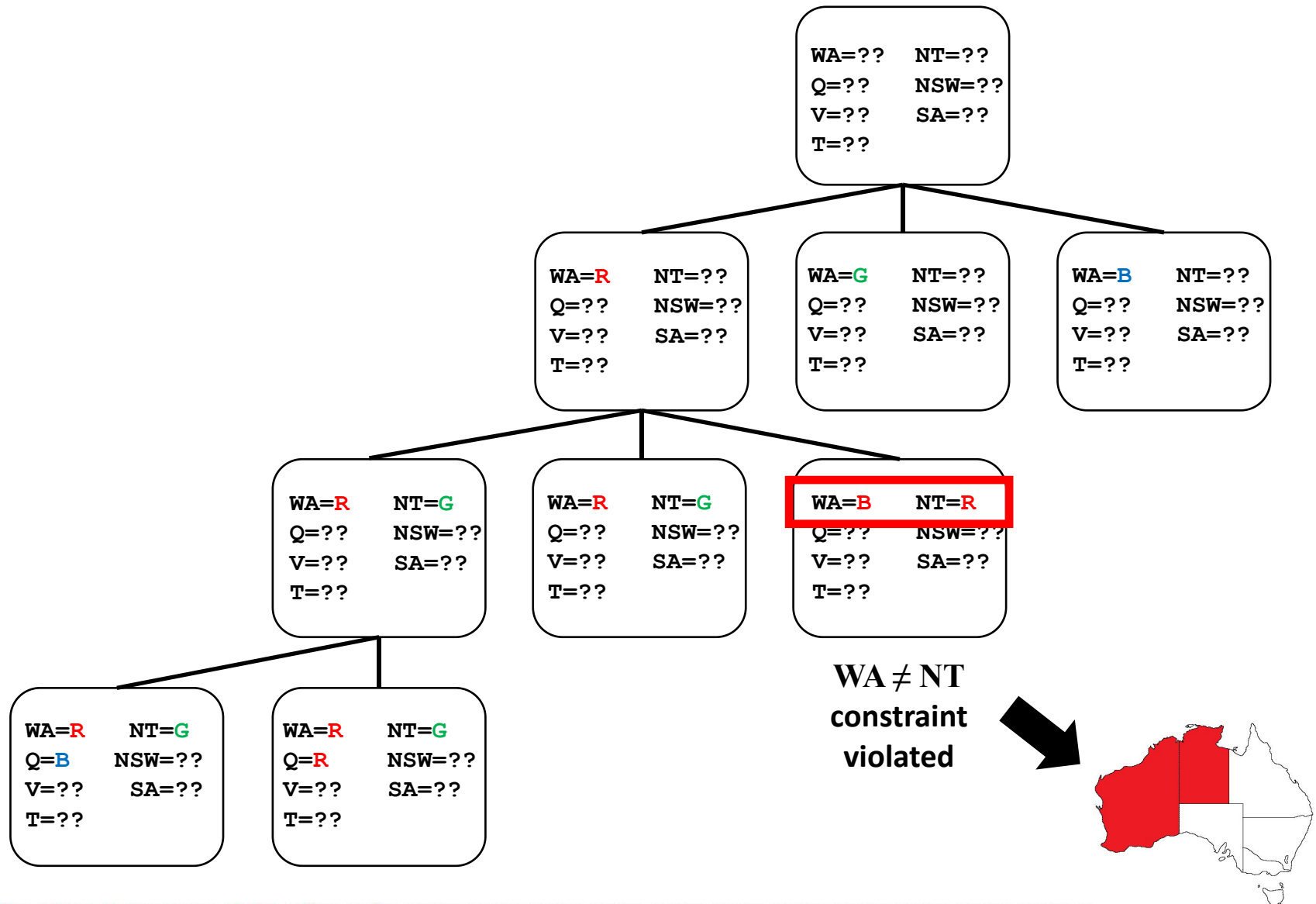
- Cost of tour is at most C :

$$\sum_{i=1}^N \sum_{j=1}^N z_{i,j} d_{i,j} \leq COST$$

CSP as a Tree Search Problem



CSP as a Tree Search Problem



CSP: Variable Types

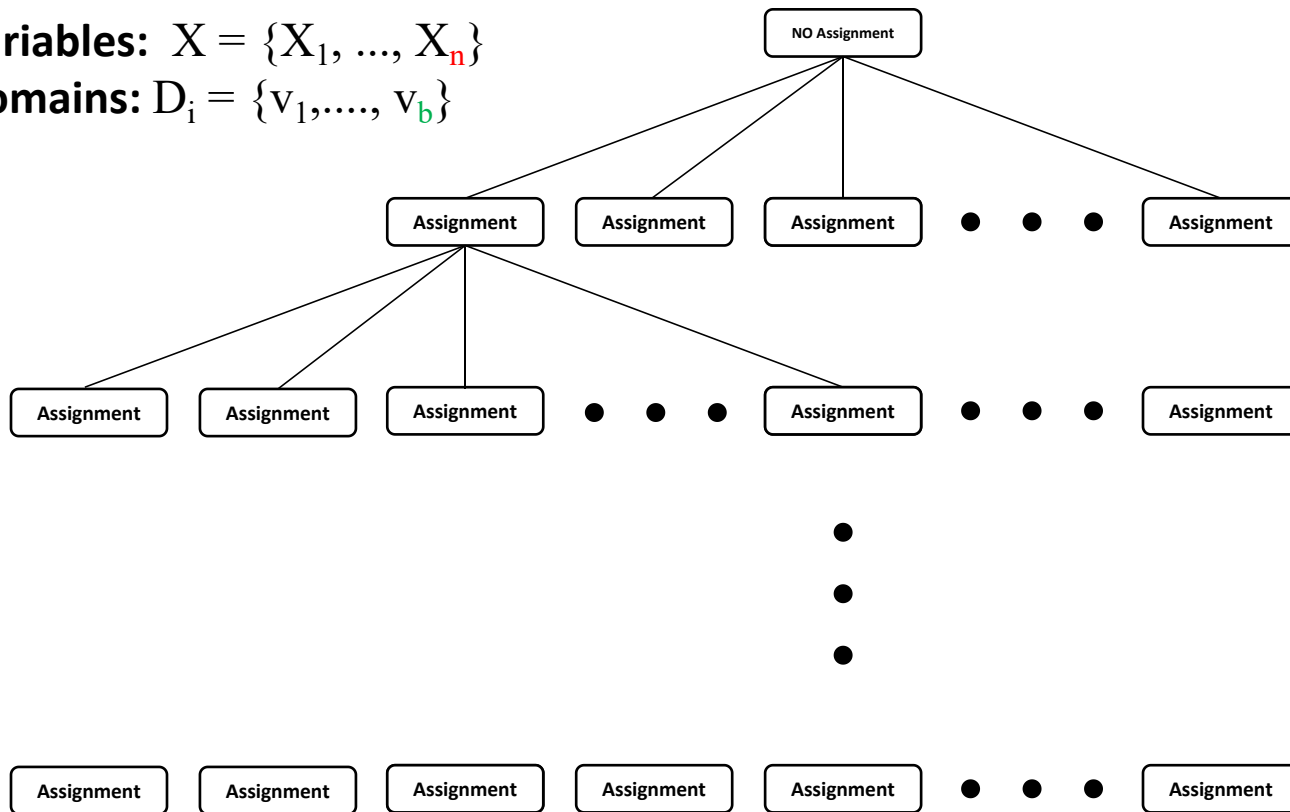
- Domains can be:
 - finite, for example: $\{1, 2, 3, 5, 8, 20\}$ (simpler)
 - infinite, for example: a set of all integers
- Variables can be:
 - discrete, for example: $X = \{X_1, \dots, X_n\}$ (simpler)
 - continuous, for example: R_+
- Constraints can be:
 - unary (involve single variable), for example: $X_1 = 5$
 - binary (involve two variables), for example: $X_1 = X_2$
 - higher order (involve > 2 variables), for example: $X_1 = X_2 * X_3$
- Soft constraints (preferences: green over blue) possible

CSP Search Tree: Idea

CSP Problem:

Variables: $X = \{X_1, \dots, X_n\}$

Domains: $D_i = \{v_1, \dots, v_b\}$



**0 variable
assigned**

**1 variables
assigned**

**2 variables
assigned**

•
•
•

ALL (n**) variables
assigned**

Tree leaves are COMPLETE assignments

The sequence of variable assignments does NOT matter*

***(when you disregard performance)**