

## CS3402 Tutorial 8:

1. Which of the following schedules is (conflict) serializable? For each serializable schedule, determine the equivalent serial schedules.

- (a)  $r_1(X); r_3(X); w_1(X); r_2(X); w_3(X);$   
 (b)  $r_1(X); r_3(X); w_3(X); w_1(X); r_2(X);$   
 (c)  $r_3(X); r_2(X); w_3(X); r_1(X); w_1(X);$

2. Consider the following concurrent schedule. Draw the serialization graph for the schedule. Is it conflict serializable?

<b>Ta</b>	<b>Tb</b>	<b>Tc</b>
	Read(x)	
Write(y)		
		Read(y)
	Write(y)	
Write(x)		
	Commit	
		Write(z)
Commit		
		Commit

3. Consider the following concurrent schedule. Draw the serialization graph for the schedule. Is it conflict serializable?

<b>Ta</b>	<b>Tb</b>	<b>Tc</b>	<b>Td</b>
	Read(z)		
			Read(z)
			Write(z)
Write(y)			
		Read(z)	
	Write(y)		
		Read(y)	
Write(x)			
	Commit		
Commit			
		Commit	
			Commit

4. Consider two transactions as specified below:

T1: Read1(x);  $x := x - 1$ ;  $y := x - 1$ ; Write1(y)

T2: Read2(y);  $y := y + 1$ ;  $x := y + 1$ ; Write2(x)

If the initial values of x and y in the database are 5 and 10 respectively, give the final values of x and y for each schedule below.

(1) Read2(y);  $y := y + 1$ ;  $x := y + 1$ ; Write2(x); Read1(x);  $x := x - 1$ ;  $y := x - 1$ ; Write1(y).

(2) Read1(x);  $x := x - 1$ ;  $y := x - 1$ ; Read2(y);  $y := y + 1$ ;  $x := y + 1$ ; Write1(y); Write2(x).

(3) Read1(x);  $x := x - 1$ ;  $y := x - 1$ ; Write1(y); Read2(y);  $y := y + 1$ ;  $x := y + 1$ ; Write2(x).

(4) For the above three schedules, which schedules are correct, and which are wrong? Briefly explain your answer.