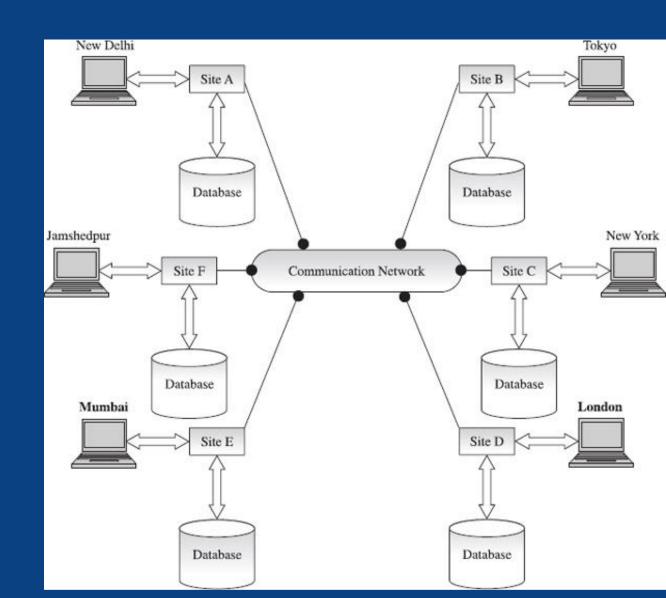


# Distributed Databases

**Database Systems & Information Modelling INFO90002** 

Week 10 – Distributed Databases Dr Tanya Linden David Eccles





#### This lecture discusses...

What is a distributed database?

Why are they used, and how they work

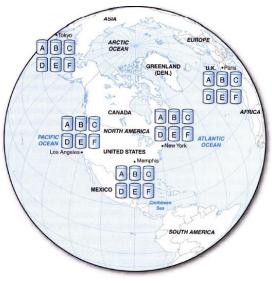
Pros and cons of different approaches

Material in this lecture is drawn from Hoffer et al. (2013) *Modern Database Management* 11<sup>th</sup> edition, chapter 12, available online at <a href="http://wps.prenhall.com/bp">http://wps.prenhall.com/bp</a> hoffer mdm 11/230/58943/15089539.cw/index.html

Images on this page are from Gillenson (2005) Fundamentals of Database Management Systems



distributed database



replicated database



#### Distributed Database

- a single logical database physically spread across multiple computers in multiple locations that are connected by a data communications link
- appears to users as though it is one database

#### **Decentralized Database**

- a collection of independent databases which are not networked together as one logical database
- appears to users as many databases

We are concerned with *distributed* databases



# Example – Amazon AWS

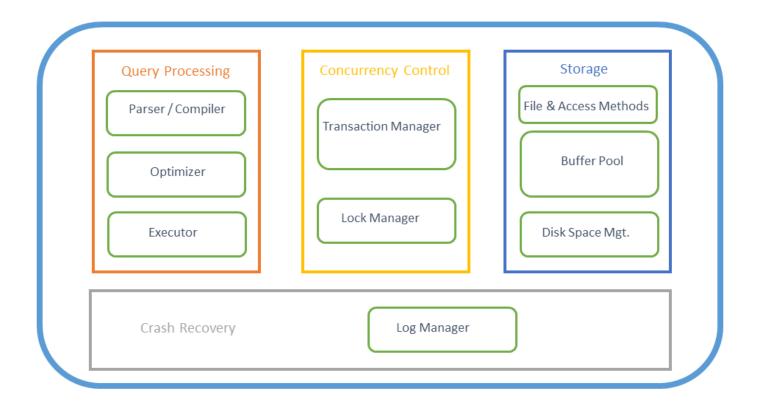
#### Global Infrastructure





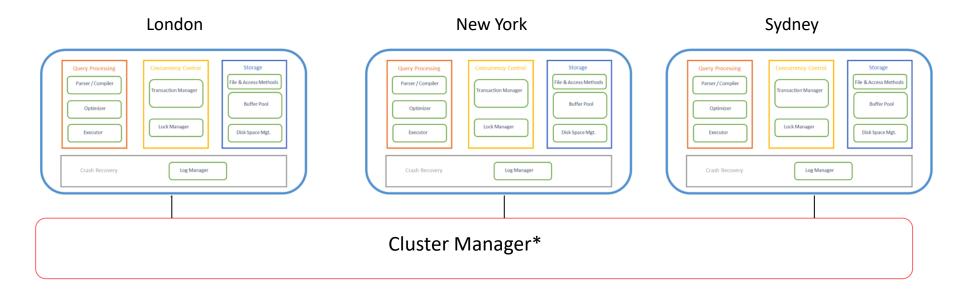
# **Database Memory Structure (1 Server)**

#### Remember this?





# **Distributed Memory Structures**



Each Physical Server has one of these memory structures

Often accessing their own and shared physical storage between all physical servers

Cluster Manager coordinates communication between physical servers

\* May be called something else in different vendor databases



# **Distributed DBMS Advantages**

Good fit for geographically distributed organisations / users

Utilize the internet

Data located near site with greatest demand

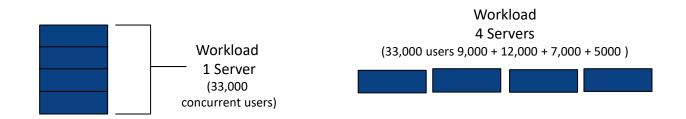
• ESPN Weekend Sports Scores



Faster data access (to local data)

Faster data processing

workload is shared between each physical server





# Distributed DBMS Advantages (cont.)

#### Allows modular growth

Add new servers as load increases

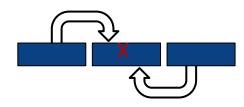
#### Increased *reliability* and *availability*

Lower danger of a single-point of failure (SPOF)



#### Supports database recovery

- Data is replicated across multiple sites
- Recovery Logs replicated





# **Disadvantages**

#### Complexity of management and control

- database or/and application must stitch together data across sites
  - What is the current version of the record (row and column) and where is it?
  - Who is waiting to update that information and where are they?
  - How does the logic display this to the web and application server?

#### Data integrity

- additional exposure to improper updating
  - If two users in two locations update the record at the exact same time who decides which statement should "win"?
  - Solution: Transaction Manager or Master-slave design

#### Security

- many server sites -> higher chance of breach
  - Multiple access sites require protection from both cyber and physical attacks (including protection of network and storage infrastructure)



#### Lack of standards

different Relational DDBMS vendors use different protocols

Increased training and maintenance costs

- more complex IT infrastructure
  - Increased Disk storage (\$)
  - Fast intra and inter network infrastructure (\$\$\$)
  - Clustering software (\$\$\$\$)
  - Network Speed (\$\$\$\$)

Increased storage requirements

Replication model



# **Objectives and Trade-offs**

#### Location transparency

a user does not need to know where particular data are stored

#### Local autonomy

a node can continue to function for local users if connectivity to the network is lost

#### Trade-offs

- Availability vs Consistency
- Synchronous vs Asynchronous updates



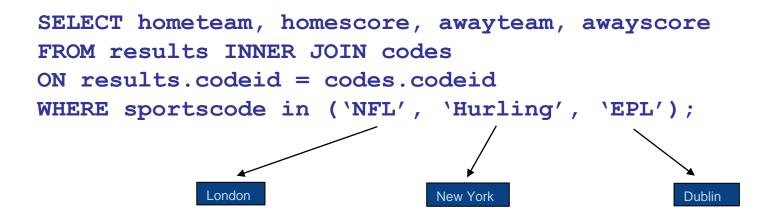
## **Location Transparency**

A user (or program) accessing data do not need to know the location of the data in the network of DBMS's

Requests to retrieve or update data from any site are automatically forwarded by the system to the site or sites related to the processing request

All data in the network appears to users as a single logical database stored at one site

A single query can join data from tables in multiple sites





## **Local Autonomy**

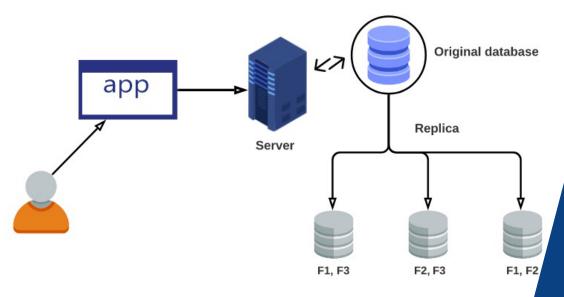
Users can administer their local database

- control local data
- administer security
- log transactions
- recover when local failures occur
- provide full access to local data

Being able to operate locally when connections to other databases fail



#### Partial replication in DBMS



[FI, F2, F3 are the different fragments of the main database]

# Distributed Options

Data Replication
Horizontal Partitioning
Vertical Partitioning



# **Distribution options**

#### Data replication

Data copied across sites

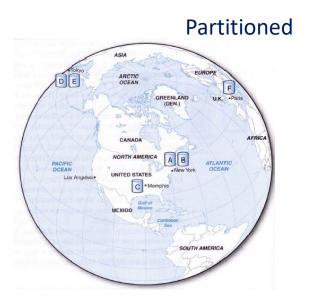
#### Horizontal partitioning

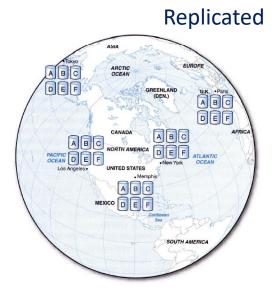
Table rows distributed across sites

#### Vertical partitioning

Table columns distributed across sites

#### Combinations of the above







# **Replication - advantages**

High reliability due to redundant copies of data

Fast access to data at the location where it is most accessed

May avoid complicated distributed integrity routines

replicated data is refreshed at scheduled intervals

Decoupled nodes don't affect data availability

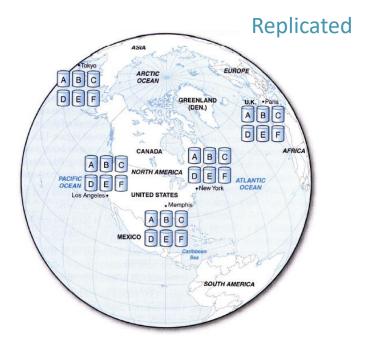
transactions proceed even if some nodes are down

Reduced network traffic at prime time

if updates can be delayed

This is currently popular as a way of achieving high availability for global systems.

Most SQL and NoSQL databases offer replication





# **Replication - disadvantages**

#### Need more storage space

Each server stores a copy of the row

#### Data Integrity:

retrieve incorrect data if updates have not arrived



Centralised Database
One database in one server
(1 copy of data)



Distributed Database
One database in 4 physical servers
(4 copies of data)



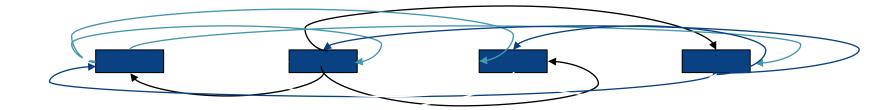
Data Size



# **Replication - disadvantages**

#### Takes time for update operations

- High tolerance for out-of-date data may be required
- Updates may cause performance problems for busy nodes



#### Network communication capabilities

- Updates can place heavy demand on telecommunications/networks
- Cost high speed networks are expensive (\$\$\$\$)



# Synchronous updates

Data is continuously kept up to date

Users anywhere can access data and get the same answer.

If any copy of a data item is updated anywhere on the network, the same update is immediately applied to all other copies or the update is aborted.

Ensures data integrity and minimises the complexity of knowing where the most recent copy of data is located.

Can result in slow response time and high network usage

- The DDBMS spends time checking that an update is accurately and completely propagated across the network.
- The committed updated record must be identical in all servers



# **Asynchronous updates**

Some delay in propagating data updates to remote databases

- Some degree of at least temporary inconsistency is tolerated
- May be ok if it is temporary and well managed

Acceptable response time

Updates happen locally and data replicas are synchronized in batches and predetermined intervals

May be more complex to plan and design

Need to ensure the right level of data integrity and consistency

Suits some information systems more than others

Compare commerce/finance systems with social media



# Horizontal partitioning

#### Different rows of a table at different sites

#### Advantages

- data stored close to where it is used
  - efficiency
- local access optimization
  - better performance
- only relevant data is stored locally
  - security
- unions across partitions
  - ease of query

#### Disadvantages

- accessing data across partitions
  - inconsistent access speed
- no data replication
  - backup vulnerability (SPOF)

ID	Team	City	Code	Region	League
1	Arsenal	London	Football	Europe	EPL
2	Jets	NYC	Grid Iron	Americas	NFL
3	Carlton FC	Melbourne	Aussie Rules	APAC	AFL
4	Racing92	Paris	Rugby	Europe	Top14
5	Yankees	NYC	Baseball	Americas	MLB
6	Swifts	Sydney	Netball	APAC	ANZ

Team table -21-



# **Example horizontal partitioning**

ID	Team	City	Code	Region	League
1	Arsenal	London	Football	Europe	EPL
2	Jets	NYC	Grid Iron	Americas	NFL
3	Carlton FC	Melbourne	Aussie Rules	APAC	AFL
4	Racing92	Paris	Rugby	Europe	Top14
5	Yankees	NYC	Baseball	Americas	MLB
6	Swifts	Sydney	Netball	APAC	ANZ

#### London

#### Team

1,Arsenal,London, Football, Europe, EPL

4,Racing92, Paris, Rugby, Europe, Top14

#### Melbourne

#### Team

3, CarltonFC, Melbourne, Aussie Rules, APAC, AFL

6,Swifts, Sydney, Netball, APAC, ANZ

#### New York

#### Team

2, Jets, NYC, Grid Iron, Americas, NFL

5, Yankees, NYC, Baseball, Americas, MLB



# Vertical partitioning

Different columns of a table at different sites

Advantages and disadvantages are the same as for horizontal partitioning

- except
  - combining data across partitions is more difficult because it requires joins (instead of unions)

ID	Firstname	Lastname	Team	League	Photo	Biography
110	Luc	Ducalon	4	Top14		Ipso locum
120	Vasil	Kakokan	4	Top14		Ipso locum est
130	Donacca	Ryan	4	Top14	<null></null>	
210	Edwin	Maka	4	Top14		

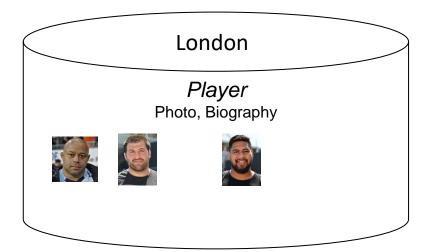


# **Example vertical partitioning**

ID	Firstname	Lastname	Team	League	Photo	Biography
110	Luc	Ducalon	4	Top14	E C	Ipso locum
120	Vasil	Kakokan	4	Top14		Ipso locum est
130	Donacca	Ryan	4	Top14	<null></null>	
210	Edwin	Maka	4	Top14		

Vertical Partitioning based on column requirements

# Player ID, First, Lastname, Team, League 110, Luc, Ducalon, 4, Top14 120, Vasil, Kakokan, 4, Top14 130, Donacca, Ryan, 4, Top14 210 Edwin Maka, 4, Top14





# Comparing 5 configurations

Centralised database, distributed access

• DB is at one location, and accessed from everywhere

Replication with periodic (asynchronous) snapshot update

Many locations, each data copy updated periodically

Replication with near real-time synchronization of updates

Many locations, each data copy updated in near real time

Partitioned, integrated, one logical database

Data partitioned across many sites, within a logical database, and a single DBMS

Partitioned, independent, non-integrated segments

- Data partitioned across many sites.
- Independent, non-integrated segments
- Multiple DBMS, multiple computers



# **Comparing Configurations**

	Reliability	Expandability	Communication Overhead	Management	Data Consistency
Centralised	POOR Depends on central server.	POOR Single Server is limited by memory & storage maximums.	VERY HIGH Traffic heads to one centralised location.	EXCELLENT One very large site is easier to manage.	EXCELLENT All users always see the same data.
Replicated with Snapshots	GOOD Redundancy and tolerated delays in data synch.	VERY GOOD Cheap to scale up with new servers.	LOW to MEDIUM Intermittent bursts of network traffic (but not constant flooding of network).	VERY GOOD Each copy is alike.	MEDIUM Update delays are tolerable with snapshot catch ups for data consistency.
Synchronised Replication	EXCELLENT Redundancy and minimal delays.	VERY GOOD  Low cost and only linear growth in synchronisation.	MEDIUM  Constant messages to maintain synchronisation.	MEDIUM  Data collusions need to be resolved and need good design and management.	VERY GOOD Close to precise consistency.
Integrated Partitions	GOOD Effective use of partitioning and redundancy.	VERY GOOD  New nodes only get the data they need and no need to change DB design.	LOW to MEDIUM  Most queries are local, but global queries to create temporary comms load.	DIFFICULT Distributed table updates require tight precise coordination.	VERY POOR Requires considerable effort and inconsistencies are not tolerated.
Decentralised Independent Partitions	GOOD Depends on local DB availability.	GOOD  New sites are independent of all other sites.	LOW Little or no traffic needs to be communicated across the network.	VERY GOOD  Easy – as each site is independent of the other sites and minimal need to share data.	LOW  No guarantee of consistency  – therefore high chance of consistency.

Legend: the darker the colour the better the feature



#### Functions of a distributed DBMS

Locate data with a distributed catalog (meta data)

Determine location from which to retrieve data and process query components

DBMS translation between nodes with different local DBMSs (using middleware)

Data consistency (via multiphase commit protocols)

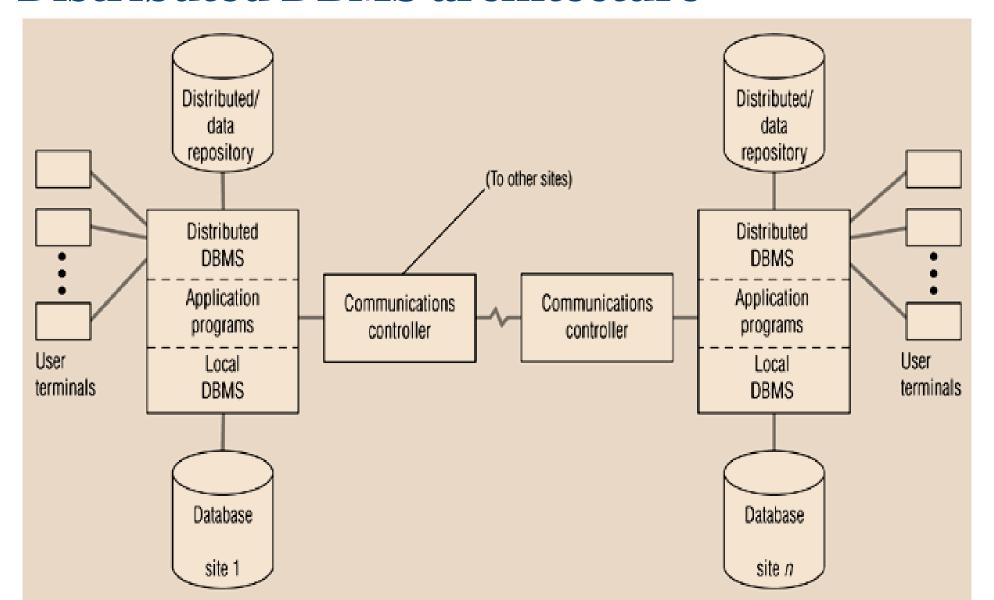
Global primary key control

Scalability

Security, concurrency, query optimisation, failure recovery



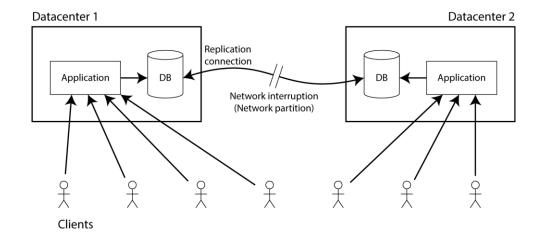
### Distributed DBMS architecture





# **Network partitions**

- Imagine you have a synchronously-updating, replicated database
- Now imagine that the link between 2 nodes is interrupted



- What are your choices?
  - shut down the system (to avoid inconsistency)
  - keep it available to users (and accept inconsistency)



# What's examinable

#### **Distributed Database**

**Advantages and Disadvantages** 

**Replicated Databases** 

**Advantages and Disadvantages** 

**Synchronous vs Asynchronous** 

Difference between

**Advantages and Disadvantages** 

**Partitioning Options** 

Vertical, Physical, Vertical and Physical

The five configurations

**Advantages and Disadvantages** 

\* All material is examinable—these are the suggested key skills you would need to demonstrate in an exam scenario



# Thank you