

Programming Shared-Memory Parallel Systems - OpenMP

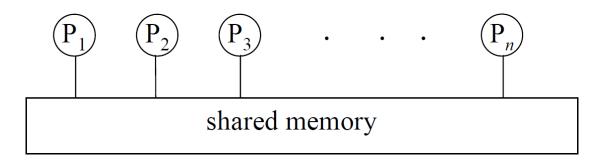
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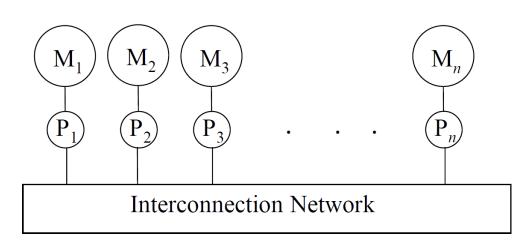
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Parallel and Distributed Computing

 The difference includes whether the processes communicate using shared or a distributed memory





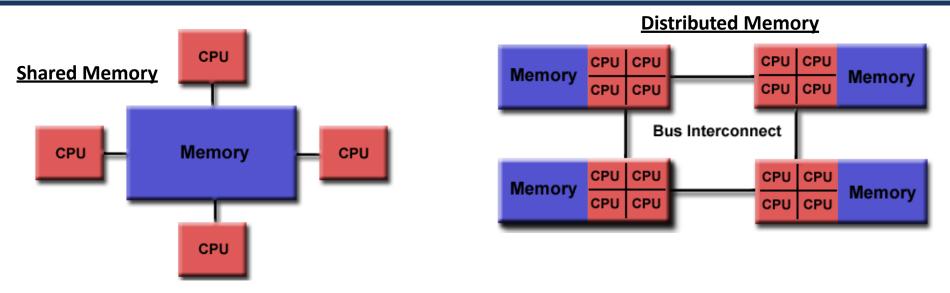


OpenMP

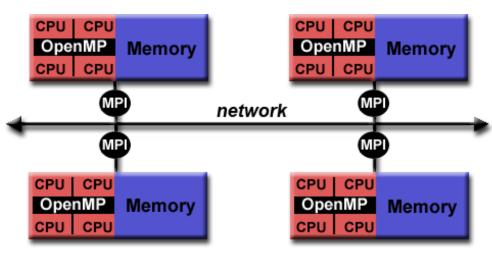
- Programming of shared memory systems
- An API for Fortran and C/C++
 - Directives
 - Runtime routines
 - Environment variables



Memory Models



Hybrid Memory





Goals

- Standardization
 - Provide a standard among a <u>variety of shared</u>
 <u>memory architectures</u> (platforms)
- High-level interfaces to thread programming
- Multi-vendor support
- Multi-OS support (Unix, Windows, Mac, etc.)
- The MP in OpenMP is for Multi-Processing
- Don't confuse OpenMP with Open MPI! :)



Release History

Date	Version
Oct 1997	Fortran 1.0
Oct 1998	C/C++ 1.0
Nov 1999	Fortran 1.1
Nov 2000	Fortran 2.0
Mar 2002	C/C++ 2.0
May 2005	OpenMP 2.5
May 2008	OpenMP 3.0
Jul 2011	OpenMP 3.1
Jul 2013	OpenMP 4.0
Nov 2015	OpenMP 4.5
Nov 2018	OpenMP 5.0

Programming Shared Memory Systems

- Explicit Parallelism
 - For example, pthreads

- Programmer Directed
 - For example, OpenMP

Hello World - pthreads based version

```
#include <pthread.h>
#include <stdio.h>
void* thrfunc(void* arg) {
        printf("hello from thread %d\n", *(int*)arg);
int main(void) {
        pthread t thread[4];
        pthread attr t attr;
        int arg[4] = \{0,1,2,3\};
        int i;
        // setup joinable threads
        pthread attr init(&attr);
        pthread_attr_setdetachstate(&attr,PTHREAD_CREATE_JOINABLE);
        // create N threads
        for(i=0; i<4; i++)
                 pthread create(&thread[i], &attr, thrfunc, (void*)&arg[i]);
        // wait for the N threads to finish
        for(i=0; i<4; i++)
                 pthread_join(thread[i], NULL);
```



... and the OpenMP version

```
#include <omp.h>
#include <stdio.h>
int main(int argc, char* argv[])
   #pragma omp parallel {
   printf("Hello World... from thread = %d\n", omp_get_thread_num());
Compilation: $ gcc -fopenmp hello.c -o hello
           $ export OMP_NUM_THREADS=4
Execution:
           $ ./hello
```

Demo: hello.c



Compiling

```
Intel (icc, ifort, icpc)

-openmp

PGI (pgcc, pgf90, ...)

-mp

GNU (gcc, gfortran, g++)

-fopenmp
```

OpenMP - User Interface Model

- Shared Memory with thread based parallelism
- Not a new language
- Compiler directives, library calls, and environment variables extend the base language
 - f77, f90, f95, C, C++
- Not automatic parallelization
 - User explicitly specifies parallelism
 - NOTE: <u>Compiler does not ignore user directives even if</u> wrong



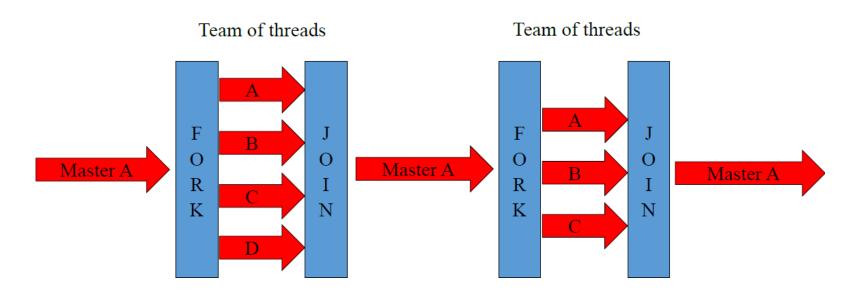
OpenMP - Syntax

- Parallelism is highlighted using compiler directives or pragmas
- For C and C++, the pragmas take the form: #pragma omp construct [clause [clause]...]
- Any compiler (even if it does not have OpenMP support)
 can compile the program (with no parallelism though)



Fork*/Join Execution Model

- An OpenMP program starts as a single thread (master thread).
- Additional threads (*Team*) are created when the <u>master hits a</u> <u>parallel region</u>.
- When <u>all threads finished the parallel region</u>, the new threads are given back to the runtime or operating system



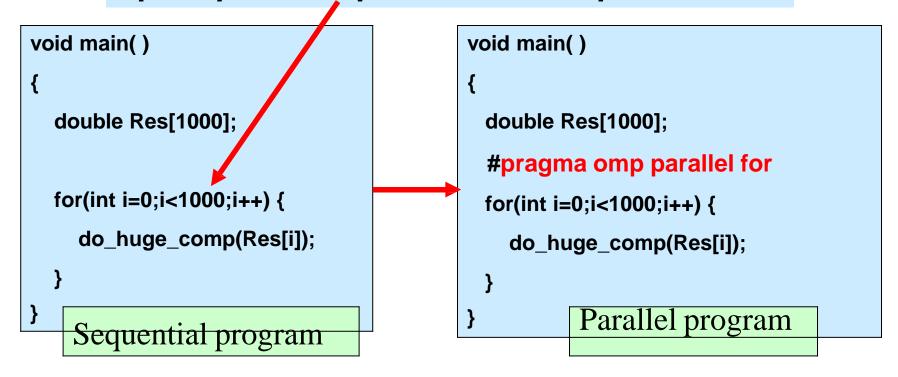
^{*}Not to be confused with fork() system call



Using OpenMP

- OpenMP is usually used to parallelize loops:
 - Find most time consuming loops
 - Split them among threads

Split-up this loop between multiple threads





OpenMP Directives



OpenMP - Directives

- OpenMP compiler directives are used for various purposes:
 - Spawning a parallel region
 - Dividing blocks of code among threads
 - Distributing loop iterations between threads

— ...

```
#pragma omp parallel private(var)
```

Supported Clauses for the Parallel Construct

Valid Clauses:

```
if (logical expression)
num_threads (integer)
private (list of variables)
firstprivate (list of variables)
shared (list of variables)
default (none|shared|private *fortran only*)
copyin (list of variables)
reduction (operator: list)
```



OpenMP Constructs

- OpenMP constructs can be divided into 5 categories:
 - 1. Parallel Regions
 - 2. Work-sharing
 - 3. Data Environment
 - 4. Synchronization
 - 5. Runtime functions/environment variables



OpenMP: Parallel Regions

- You create threads in OpenMP with "omp parallel" pragma
- For example: a 4-thread based Parallel region:

```
int A[10];
omp_set_num_threads(4);
#pragma omp parallel
{
  int ID =omp_get_thread_num();
  fun1(ID,A);
}
```

- Implicit barrier at the end of parallel block
- Each thread calls fun1(ID,A) for ID = 0 to 3
- Each thread executes the <u>same code</u> within the block



The parallel directive

- A parallel region is a block of code that will be executed by multiple threads
- When (in <u>serial program</u>) a <u>PARALLEL directive</u> is <u>found</u>, a <u>team of threads</u> is <u>created</u> and <u>main-thread</u> (serial execution thread) <u>becomes</u> the <u>master of the team</u>
- Master thread has id or number 0 (within that team)
- The code is duplicated and all threads will execute that code
- There is an implicit barrier at the end of a parallel region
- Master thread continues execution after this point



The parallel directive

- Some common clauses include:
 - if (expression)
 - private (list)
 - shared (list)
 - num_threads (integer-expression)



How Many Threads?

- The number of threads in a parallel region is determined by the following factors, in order of precedence:
 - 1. Evaluation of the if clause
 - 2. Setting of the NUM_THREADS clause
 - Use of the omp_set_num_threads() library function
 - 4. Setting of the **OMP_NUM_THREADS** environment variable
 - Implementation default: Usually the number of CPUs on a node
- Threads are numbered from 0 (master thread) to N-1



IF clause

```
#pragma omp parallel if (scalar_expression)
```

- Execute in parallel <u>if expression is true</u>
 - Otherwise serial execution

NUM_THREADS clause

```
#pragma omp parallel if(np>1) num_threads(np)
{
    ...
}
```

- Execute in parallel <u>if expression is true</u>
- Executes using np number of threads

omp_set_num_threads() function

```
#define TOTAL_THREADS 8
int main()
{
    omp_set_num_threads(TOTAL_THREADS);
    #pragma omp parallel
    {
        . . .
    }
    . . .
```

Execute in parallel using 8 threads

```
$ export OMP_NUM_THREADS=4
$ echo $OMP_NUM_THREADS
```

Sets and displays the value of the environment variable OMP_NUM_THREADS



Execution Status in Parallel Region

```
int omp_in_parallel()
```

- Returns non-zero: if execution is in parallel region
- Returns zero: if execution in non-parallel region

Demo: PRegion.c



- Shared data are accessible by all threads
- A reference a[5] to a shared array accesses the same address in all threads
- Private data are accessible only by a thread
 - Each thread has its own copy
- The default is shared



```
int main(int argc, char* argv[])
    int threadData = 10;
    // Beginning of parallel region
    #pragma omp parallel private(threadData)
        threadData = 200;
    // Ending of parallel region
    printf("Value: %d\n", threadData);
```

Demo: SPData.c

#pragma omp parallel shared(list)

- Default behavior
- List will be shared
- Each thread access the same memory location
- Initial value (<u>for the first thread</u>) will be same as before the region
- Final value will be updated by the <u>last thread</u> leaving the region
- Problems: <u>Data Race</u>



```
#pragma omp parallel private (list)
```

- Data <u>local to thread</u>
- You <u>should not rely</u> on any <u>initial</u> and <u>terminal value</u>
 (after execution of the parallel region)
- Separate "Stack Memory" for each thread's private data
- No storage associated with original object (even with same name for data-items)

Use firstprivate and/or lastprivate clause to override



```
#pragma omp parallel firstprivate (list)
```

- Variables in list are private
- Initialized with the value the variable had before entering the construct

```
#pragma omp parallel for lastprivate (list)
```

- Used in "for" loops
- Variables in list are private
- The thread that executes the final iteration of the loop updates the value (of the variables in the list)



- Alter the default behavior
- To implement customized access behavior

Shared and Private Data – Example (1/4)

```
#include <omp.h>
#include <stdio.h>
int main() {
 int i;
 const int N = 1000;
 int a = 50;
 int b = 0;
#pragma omp parallel for default(shared)
 for (i=0; i<N; i++) {
   b = a + i;
 printf("a=%d b=%d (expected a=50 b=1049)\n", a, b);
}
```

Demo: SPDE1.c

Shared and Private Data – Example (2/4)

```
#include <omp.h>
#include <stdio.h>
int main() {
  int i;
 const int N = 1000;
 int a = 50;
 int b = 0;
#pragma omp parallel for default(none) private(i) private(a) private(b)
  for (i=0; i<N; i++) {
   b = a + i:
 printf("a=%d b=%d (expected a=50 b=1049)\n", a, b);
```

Demo: SPDE1.c

Shared and Private Data – Example (3/4)

```
#include <omp.h>
#include <stdio.h>
int main() {
  int i;
  const int N = 1000;
  int a = 50;
  int b = 0;
#pragma omp parallel for default(none) private(i) private(a) lastprivate(b)
  for (i=0; i<N; i++) {</pre>
   b = a + i;
 printf("a=%d b=%d (expected a=50 b=1049)\n", a, b);
```

Demo: SPDE1.c

Shared and Private Data – Example (4/4)

```
#include <omp.h>
#include <stdio.h>
int main() {
 int i;
 const int N = 1000;
 int a = 50;
 int b = 0;
#pragma omp parallel for default(none) private(i) firstprivate(a) lastprivate(b)
  for (i=0; i<N; i++) {
   b = a + i:
 printf("a=%d b=%d (expected a=50 b=1049)\n", a, b);
```

Demo: SPDE1.c



Getting ID of Current Thread

```
int main(int argc, char* argv[])
  int iam, nthreads;
 #pragma omp parallel private(iam,nthreads) num threads(2)
  {
      iam = omp get thread num();
      nthreads = omp get num threads();
      printf("ThradID %d, out of %d threads\n", iam, nthreads);
      if (iam == 0)
          printf("Here is the Master Thread.\n");
      else
          printf("Here is another thread.\n");
```

Demo: CTID.c



Work-Sharing Constructs

- If all the threads are doing the same thing, what is the advantage then?
- Within each "Team" threads are assigned IDs, with master thread assigned ID 0
 - omp_get_thread_num() //to get thread number

Can we use this to distribute tasks amongst the "team" members?

- Work-sharing constructs distribute the specified work to all threads within the current team
- Work-sharing constructs do not launch new threads

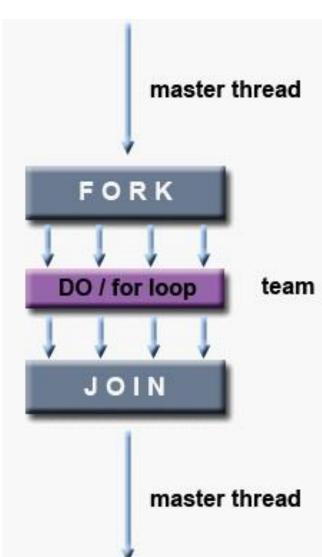


For Work-Sharing Construct

for - shares iterations of a loop across the team

#pragma omp for [clause ...] newline

There is an implicit synchronization after #pragma omp for





For Work-Sharing Construct

 SCHEDULE clause describes how iterations of the loop are divided among the threads in the team



Chunks of specified size assigned round-robin



Chunks of specified size are assigned when thread finishes previous chunk (work-Stealing mechanism)



Do/For Work-Sharing Construct

```
int main(int argc, char* argv[])
      int i, a[10];
      #pragma omp parallel num threads(2)
             #pragma omp for schedule(static, 2)
             for (i=0; i<10;i++)
                    a[i] = omp get thread num();
      for (i=0; i<10;i++)
             printf("%d",a[i]);
```

Demo: ForConst.c



Do/For Work-Sharing Construct

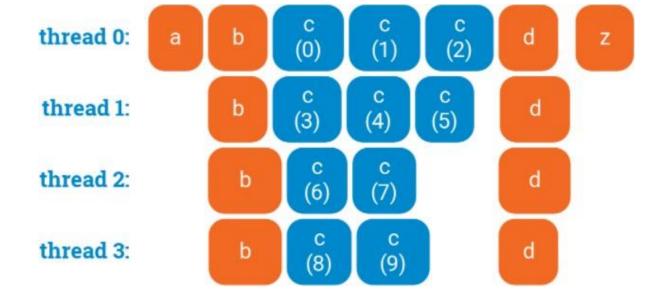
```
int main(int argc, char* argv[])
    int sum, counter, inputList[6] = \{11,45,3,5,12,-3\};
    #pragma omp parallel num threads(2)
       #pragma omp for schedule(static, 3)
      for (counter=0; counter<6; counter++) {</pre>
            printf("%d adding %d to the
            sum\n",omp get thread num(), inputList[counter]);
            sum+=inputList[counter];
       } //end of for
    } //end of parallel section
    printf("The summed up Value: %d", sum);
```

Demo: ForConst2.c



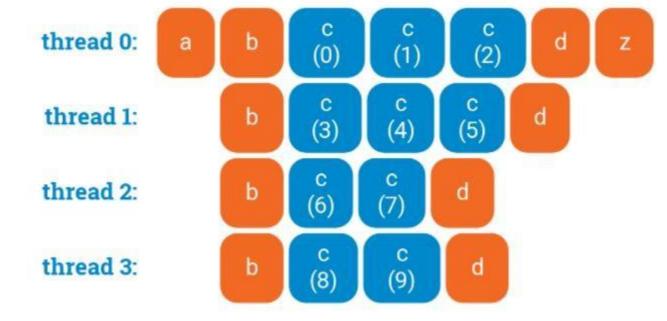
For Work-Sharing –Synchronized

```
a();
#pragma omp parallel
{
    b();
    #pragma omp for
    for (int i = 0; i < 10; ++i) {
        c(i);
    }
    d();
}
z();</pre>
```



For Work-Sharing – Non Synchronized

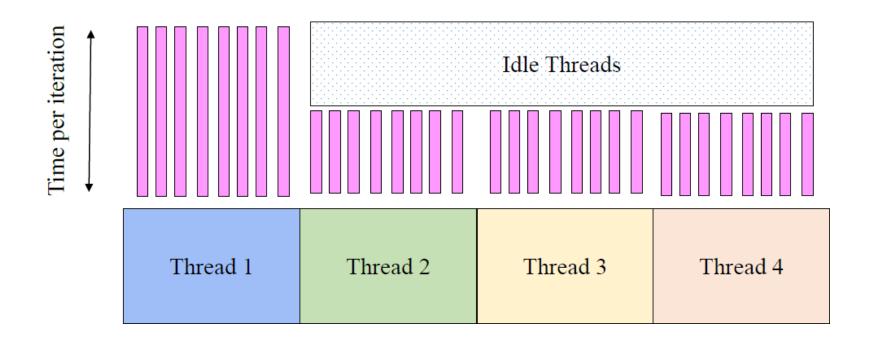
```
a();
#pragma omp parallel
    b();
    #pragma omp for nowait
    for (int i = 0; i < 10; ++i) {
        c(i);
    d();
z();
```





Problems with Static Scheduling

- What happens if loop iterations do not take the same amount of time?
 - Load imbalance





Dynamic Scheduling

- Fixed size chunks assigned on the fly
 - Work-stealing mechanism
- Disadvantage: more overhead as compared to Static



Demo: LoopSched.c



Threads share Global variables!

```
#include <pthread.h>
#include <iostream>
#include <unistd.h>
using namespace std;
#define NUM THREADS 10
int sharedData = 0;
void* incrementData(void* arg) {
     sharedData++;
     pthread exit(NULL);
int main()
    pthread t threadID;
    for (int counter=0; counter<NUM THREADS;counter++) {</pre>
         pthread_create(&threadID, NULL, incrementData, NULL);
    cout << "ThreadCount:" << sharedData <<endl;</pre>
    pthread exit(NULL);
```

The output for the pthread version?

>./globalData

ThreadCount:10

>./globalData

ThreadCount:8



ThreadCount: A better implementation

```
#include <pthread.h>
#include <iostream>
#include <unistd.h>
using namespace std;
#define NUM_THREADS 100
int sharedData = 0;
void* incrementData(void* arg)
    sharedData++;
    pthread exit(NULL); }
int main()
   pthread t threadID[NUM THREADS];
   for (int counter=0; counter<NUM THREADS; counter++) {</pre>
      pthread create(&threadID[counter], NULL, incrementData, NULL);
   //waiting for all threads
   int statusReturned;
   for (int counter=0; counter<NUM_THREADS;counter++) {</pre>
      pthread join(threadID[counter], NULL);
   cout << "ThreadCount:" << sharedData <<endl;</pre>
   pthread_exit(NULL);
```



Is the problem solved?

- Unfortunately, not yet :(
- The output from running it with 1000 threads is as below:

```
>./6join
ThreadCount:990
>./6join
ThreadCount:978
>./6join
ThreadCount:1000
>
```

- Reasons?
- What can be done?

ThreadCount: OpenMP Implementation

```
int main(int argc, char* argv[])
  int threadCount=0;
  #pragma omp parallel num threads(100)
    int myLocalCount = threadCount;
    sleep(1);
    myLocalCount++;
    threadCount = myLocalCount;
  printf("Total Number of Threads: %d\n", threadCount);
```

Demo: TCount1.c

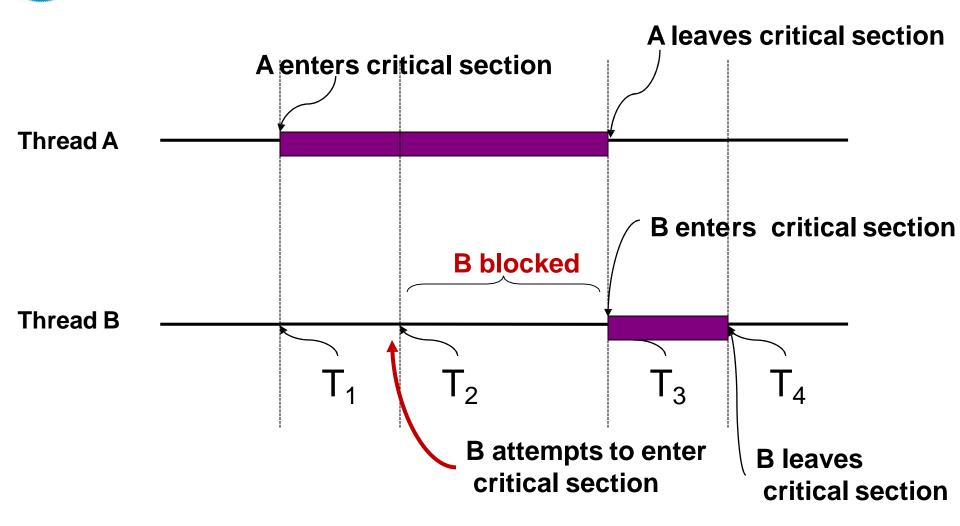


Critical-Section (CS) Problem

- n processes all competing to use some shared data
- ➤ Each process has a code segment, called critical section, in which the shared data is accessed
- Problem (ensures that):
 - -Two process are not allowed to execute in their critical section at the same time
 - Access to the critical section must be an atomic action



Critical Section



Mutual Exclusion

At any given time, only one Thread is in the critical section



...back to threads counting

```
int sharedData = 0:
pthread mutex t mutexIncrement;
void* incrementData(void* arg)
    pthread mutex lock(&mutexIncrement);
    sharedData++;
    pthread mutex unlock(&mutexIncrement);
    pthread exit(NULL);
int main()
  pthread mutex init(&mutexIncrement, NULL);
  pthread t threadID[NUM THREADS];
   for (int counter=0; counter<NUM THREADS; counter++) {</pre>
      pthread create(&threadID[counter], NULL, incrementData,
      NULL);
   //waiting for all threads
   int statusReturned;
   for (int counter=0; counter<NUM THREADS; counter++)</pre>
      { pthread join(threadID[counter], NULL);
   cout << "ThreadCount:" << sharedData</pre>
   <<endl; pthread exit(NULL);</pre>
```

OpenMP - Synchronization Constructs

- The CRITICAL directive specifies a region of code that must be executed by only one thread at a time
- If a thread is currently executing inside a CRITICAL region and another thread attempts to execute it, it will block until the first thread exits that CRITICAL region.

pragma omp critical [name]

. . .



... back to threadCount

```
int main(int argc, char* argv[])
   int threadCount;
   #pragma omp parallel num_threads(5)
      #pragma omp critical
          int myLocalCount = threadCount;
          sleep(1);
          myLocalCount++;
          threadCount = myLocalCount;
   printf("Total Number of Threads: %d\n", threadCount);
```

Demo: TCount2.c

OpenMP - Synchronization Constructs

 The MASTER directive specifies a region that is to be executed only by the master thread of the team

All other threads on the team skip this section of code

#pragma omp master

. . .

Demo: MasterOnly.c

OpenMP - Synchronization Constructs

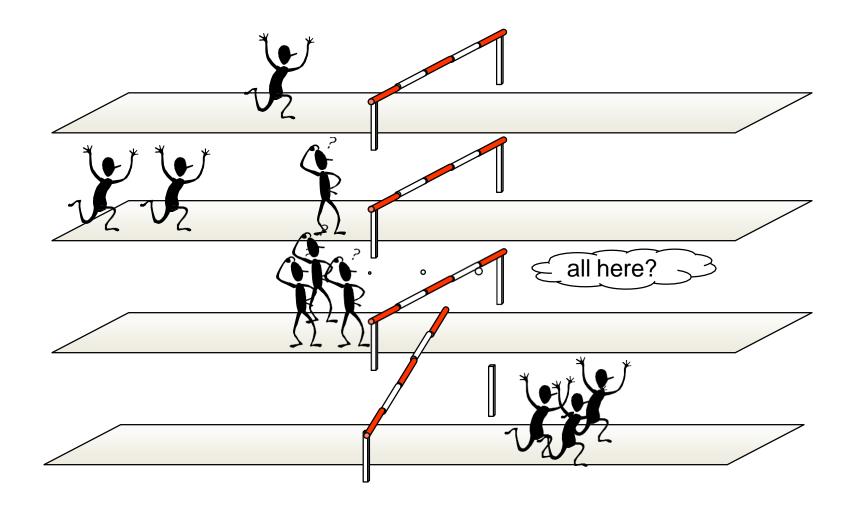
- When a BARRIER directive is reached, a thread will wait at that point until all other threads have reached that barrier
- All threads then resume executing in parallel the code that follows the barrier.

#pragma omp barrier

...



Barrier Synchronization



Demo: Barrier.c

Reduction (Data-sharing Attribute Clause)

- The REDUCTION clause performs a reduction operation on the variables that appear in the list
- A <u>private copy</u> for <u>each list variable</u> is <u>created</u> and <u>initialized</u> for <u>each thread</u>
- At the <u>end of the reduction</u>, the <u>reduction variable</u> (all private copies) is examined and the <u>shared variable</u>'s final result is written.

```
#pragma omp operator: list
```

. . .

operator can be +, -, *, &&, | |, max, min ...

Reduction (Data-sharing Attribute Clause)

```
int main(int argc, char* argv[])
  srand(time(NULL));
   int winner;
  #pragma omp parallel reduction(max:winner) num threads(10)
        winner = (rand() \% 1000) + omp get thread num();
        printf("Thread: %d has Chosen: %d\n",
omp get thread num(),winner);
  printf("Winner: %d\n", winner);
```

Demo: Reduction.c



Any Questions?