

Lab5: Channel Coding

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<Problem 1>

- There are also .jpg file at the submitted folder.

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(d) Given that $g_1 = [1, 0, 1]$, $g_2 = [1, 1, 1]$

(a) Calculate the code rate R .

Ans: $R = \frac{1}{2}$

<sol> $R = \frac{k}{n}$: $b^{(k)} \rightarrow c^{(n)}$, for $g_1: g_1[0] \oplus g_1[2] = c^{(1)}$, for $g_2: g_2[0] \oplus g_2[1] \oplus g_2[2] = c^{(2)}$

$b^{(n)}$ to generate $c^{(n)}$ (with constraint length $K=3$). $\therefore k=1, n=2 \Rightarrow R = \frac{1}{2}$

(b) Draw the shift register structure.

<sol>

(c) Draw the state transition diagram.

<sol> Based on the graph in (b).

Input	State $b_{i-1}b_{i-2}$ (before)	output $c_i^{(1)}$	output $c_i^{(2)}$	State $b_{i-1}b_{i-2}$ (after)
0	00 (state A)	0	0	00 (state A)
0	01 (state B)	1	1	00 (state A)
0	10 (state C)	0	1	01 (state B)
0	11 (state D)	1	0	01 (state B)
1	00 (state A)	1	1	10 (state C)
1	01 (state B)	0	0	10 (state C)
1	10 (state C)	1	0	11 (state D)
1	11 (state D)	0	1	11 (state D)

⇒

input 0 $c_i^{(1)} c_i^{(2)}$ $d_H(d, 0)$
input 1 $c_i^{(1)} c_i^{(2)}$ $d_H(d, 1)$

(d) Draw the trellis diagram.

<sol>

(e) Illustrate the path on the trellis diagram, also list the decoded codewords $\hat{c} = [\hat{c}_i]$, and the decoded bits $\hat{b} = [\hat{b}_i]$ ($d = [0, 0, 1, 1, 0, 1, 0, 0, 1, 0, 0, 1]$)

<sol> Received y_i : 00, 11, 01, 00, 01, 01, 00
Decoded \hat{c} : 00, 11, 01, 00, 01, 11, 00
Decoded \hat{b} : 0, 1, 0, 1, 0, 0, 0

(f) Based on (e), calculate the $d_H(d, \hat{c})$.

<sol> Compare $d = 001101000100$, there's 1 bit different. $\therefore d_H(d, \hat{c}) = 1$
 $\hat{c} = 0011010001100$

<Problem 2>

- Construct a FSM table by declaring global variables and doing xor. (general case of $K == 3$)

```
global A 0 A 1 B 0 B 1 C 0 C 1 D 0 D 1;
A 0 = '';
A 1 = '';
B 0 = '';
B 1 = '';
C 0 = '';
C 1 = '';
D 0 = '';
D 1 = '';
a 0 = [0, 0, 0];
b 0 = [0, 0, 1];
c 0 = [0, 1, 0];
d 0 = [0, 1, 1];
a 1 = [1, 0, 0];
b 1 = [1, 0, 1];
c 1 = [1, 1, 0];
d 1 = [1, 1, 1];
```

(initialize)

```

for j = 1:size(impulse_response, 2)
    if (impulse_response(i, j) == 1)
        g = [g, j];
    end
end

```

(if g(i) == 1, do xor)

```

if (length(g) == 1)
    A_0 = [A_0, int2str(a_0(g(1)))];
    A_1 = [A_1, int2str(a_1(g(1)))];
    B_0 = [B_0, int2str(b_0(g(1)))];
    B_1 = [B_1, int2str(b_1(g(1)))];
    C_0 = [C_0, int2str(c_0(g(1)))];
    C_1 = [C_1, int2str(c_1(g(1)))];
    D_0 = [D_0, int2str(d_0(g(1)))];
    D_1 = [D_1, int2str(d_1(g(1)))];
end

```

(if only 1 bit needs to do xor, xor the bit with '0')

```

for k = 2:length(g)
    a = xor(a, a_0(g(k)));
    b = xor(b, b_0(g(k)));
    c = xor(c, c_0(g(k)));
    d = xor(d, d_0(g(k)));
    aa = xor(aa, a_1(g(k)));
    bb = xor(bb, b_1(g(k)));
    cc = xor(cc, c_1(g(k)));
    dd = xor(dd, d_1(g(k)));
end

```

(else do xor with others)

```

A_0 = [A_0, int2str(a)];
A_1 = [A_1, int2str(aa)];
B_0 = [B_0, int2str(b)];
B_1 = [B_1, int2str(bb)];
C_0 = [C_0, int2str(c)];
C_1 = [C_1, int2str(cc)];
D_0 = [D_0, int2str(d)];
D_1 = [D_1, int2str(dd)];

```

(finally we can get the output bits)

- (a) **Encode binary data**

- **result** of encoding {1, 0, 1, 1, 0}, which is the same as lecture ppt.

```
>> lab5_q2
Encoded data:
111001100110010
```

(mine)

Input bits b_i	1	0	1	1	0
Coded bits c_i	111	001	100	110	010

(lecture ppt)

- **main function:** input a binary sequence, and call the function “FSM_Table” to output the encoded bits sequence.

```
% conv_enc
function enc_data = conv_enc(bin, imp)
    % coded bits
    enc_data = [];
    % initial state
    K = size(imp, 2);
    state = '';
    for i = 1:K-1
        state = [state, '0'];
    end
    % For every input bits
    for i = 1:length(bin)
        [state, enc_data] = FSM_Table(state, bin(i), enc_data);
    end
end
```

- **subfunction:** FSM_Table(finite state machine table): Based on the constructed table to output a list containing “state $b(i-1)b(i-2)$ ” and “encoded bits”, dividing the scenario into 4 parts -- state A/B/C/D.

Input	State $b_i b_{i-2}$ (Before)	Output $c_i^{(1)}$	Output $c_i^{(2)}$	Output $c_i^{(3)}$	State $b_i b_{i-2}$ (After)
0	00 (State A)	0	0	0	00 (State A)
0	01 (State B)	0	1	1	00 (State A)
0	10 (State C)	0	0	1	01 (State B)
0	11 (State D)	0	1	0	01 (State B)
1	00 (State A)	1	1	1	10 (State C)
1	01 (State B)	1	0	0	10 (State C)
1	10 (State C)	1	1	0	11 (State D)
1	11 (State D)	1	0	1	11 (State D)

(state table)

```

% State A: 00
if (strcmp(state, '00') == 1)
    if (bits == 0)
        s = '00';
        e = [enc_string, A_0];
        return;
    end
    if (bits == 1)
        s = '10';
        e = [enc_string, A_1];
        return;
    end
end
end

```

(State A)

```

% State B: 01
if (strcmp(state, '01') == 1)
    if (bits == 0)
        s = '00';
        e = [enc_string, B_0];
        return;
    end
    if (bits == 1)
        s = '10';
        e = [enc_string, B_1];
        return;
    end
end
end

```

(State B)

```

% State C: 10
if (strcmp(state, '10') == 1)
    if (bits == 0)
        s = '01';
        e = [enc_string, C_0];
        return;
    end
    if (bits == 1)
        s = '11';
        e = [enc_string, C_1];
        return;
    end
end
end

```

(State C)

```

% State D: 11
if (strcmp(state, '11') == 1)
    if (bits == 0)
        s = '01';
        e = [enc_string, D_0];
        return;
    end
    if (bits == 1)
        s = '11';
        e = [enc_string, D_1];
        return;
    end
end
end

```

(State D)

- (b) Decode the received $y(i)$

- **result** of decoding “d”

$$d = [0 \ 1 \ 0 \ 0 \ 0 \ 0 \ 1 \ 0 \ 1 \ 1 \ 1 \ 1 \ 0 \ 0 \ 1 \ 0 \ 1 \ 1 \ 0 \ 0 \ 0].$$

Decoded data:

0 0 1 1 1 0 1

- **result** of 1-(e) via Matlab & handwriting ({0, 1, 0, 1, 0, 0, 0})

```
>> lab5_q2b_verify_1e
```

Decoded data of 1-(e):

0 1 0 1 0 0 0

(Matlab)

(e) Illustrate the path on the trellis diagram, also list the decoded codewords $\hat{c} = [\hat{c}_i]$, and the decoded bits $\hat{b} = [\hat{b}_i]$ ($d = [0, 0, 1, 1, 0, 1, 0, 0, 1, 0, 1, 0]$)

<<sol> Received y_i : 00, 11, 01, 00, 01, 01, 00

Decoded \hat{c}_i : 00, 11, 01, 00, 01, 11, 00

Decoded \hat{b}_i : 0, 1, 0, 1, 0, 0, 0

(handwriting)

- **main function:** 4 state -- A/B/C/D, call the function “ham_dis”, which outputs “[next_state, encoded bits, hamming_distance]”, if next input ‘0’ or ‘1’ all has the same hamming distance, it will set the next_state = 2 and call another function “next_ham” to compare next route and determine the current way.

```

% State A
dist = 0;
if (strcmp(state, '00') == 1)
    [next_state, state, dist] = ham_dis(bits, A_0, A_1, '00');
    if (next_state == 2)
        [next_state, state, dist] = next_ham(bits_next, A_0, A_1, C_0, C_1, '00');
        dec_data = [dec_data, next_state];
        continue;
    else
        dec_data = [dec_data, next_state];
        continue;
    end
end
end

```

(State A)

```

% State B
if (strcmp(state, '01') == 1)
    [next_state, state, dist] = ham_dis(bits, B_0, B_1, '01');
    if (next_state == 2)
        [next_state, state, dist] = next_ham(bits_next, A_0, A_1, C_0, C_1, '01');
        dec_data = [dec_data, next_state];
        continue;
    else
        dec_data = [dec_data, next_state];
        continue;
    end
end
end

```

(State B)

```

% State C
if (strcmp(state, '10') == 1)
    [next_state, state, dist] = ham_dis(bits, C_0, C_1, '10');
    if (next_state == 2)
        [next_state, state, dist] = next_ham(bits_next, B_0, B_1, D_0, D_1, '10');
        dec_data = [dec_data, next_state];
        continue;
    else
        dec_data = [dec_data, next_state];
        continue;
    end
end
end

```

(State C)

```

% State D
if (strcmp(state, '11') == 1)
    [next_state, state, dist] = ham_dis(bits, D_0, D_1, '11');
    if (next_state == 2)
        [next_state, state, dist] = next_ham(bits_next, B_0, B_1, D_0, D_1, '11');
        dec_data = [dec_data, next_state];
        continue;
    else
        dec_data = [dec_data, next_state];
        continue;
    end
end
end

```

(State D)

- **subfunction 1:** ham_dis, if there is a bit different from the correct decoded bit, accumulating it by 1. Divide it into 2 scenario -- input '0' and '1'


```

for i = 1:length(bits)
    if (strcmp(bits(i), str_1(i)) == 0)
        diff_0 = diff_0 + 1;
    end
    if (strcmp(bits(i), str_2(i)) == 0)
        diff_1 = diff_1 + 1;
    end
end

```

```

% decode to 0
if (diff_0 < diff_1)
    next_state = 0;
    dist = diff_0;
    % State A
    if (strcmp(state, '00') == 1)
        output = '00';
    end
    % State B
    if (strcmp(state, '01') == 1)
        output = '00';
    end
    % State C
    if (strcmp(state, '10') == 1)
        output = '01';
    end
    % State D
    if (strcmp(state, '11') == 1)
        output = '01';
    end
end

```

```

% decode to 1
if (diff_0 > diff_1)
    next_state = 1;
    dist = diff_1;
    % State A
    if (strcmp(state, '00') == 1)
        output = '10';
    end
    % State B
    if (strcmp(state, '01') == 1)
        output = '10';
    end
    % State C
    if (strcmp(state, '10') == 1)
        output = '11';
    end
    % State D
    if (strcmp(state, '11') == 1)
        output = '11';
    end
end

```

- **subfunction 2:** next_ham, according to the next two possible nodes two other routes respectively, record their hamming distance and recursively call.

```

[n_1, s_1, d_1] = ham_dis(bit, str_1_1, str_1_2, sta);
[n_2, s_2, d_2] = ham_dis(bit, str_2_1, str_2_2, sta);
if (d_1 < d_2)
    n = n_1;
    s = s_1;
    d = d_1;
elseif (d_1 > d_2)
    n = n_2;
    s = s_2;
    d = d_2;
end

```

- **(c) Simulation**

- Generate 100003 (>100000) a random (0, 1) with the equiprobability.

```

r = mod(reshape(randperm(1*100002), 1, 100002), 2);

```

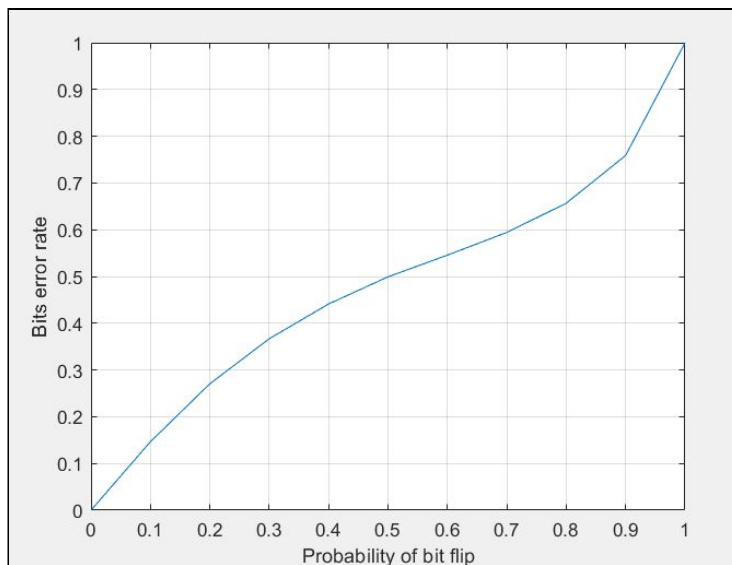
- Encode the list with my FIR and convert the string bits to binary bits.

```
e = conv_enc(r, impulse_response);
e_double = [];
for i = 1:length(e)
    temp = str2double(e(i));
    e_double = [e_double, temp];
end
```

- Use the build-in function “bsc()” and “biterr()” to flip the encoded bits with a probability $p=[0, 1]$ (for every 0.1). Calculate the bit error rate

```
% Bit Error
p = 0;
bit_error = [];
for i = 1:11
    p = (i-1)*0.1;
    BSC = bsc(e_double, p); % Binary symmetric channel
    dec_r = conv_dec(BSC, impulse_response);
    [num, err] = biterr(r, dec_r);
    bit_error = [bit_error, err];
end
```

- Plot the result.



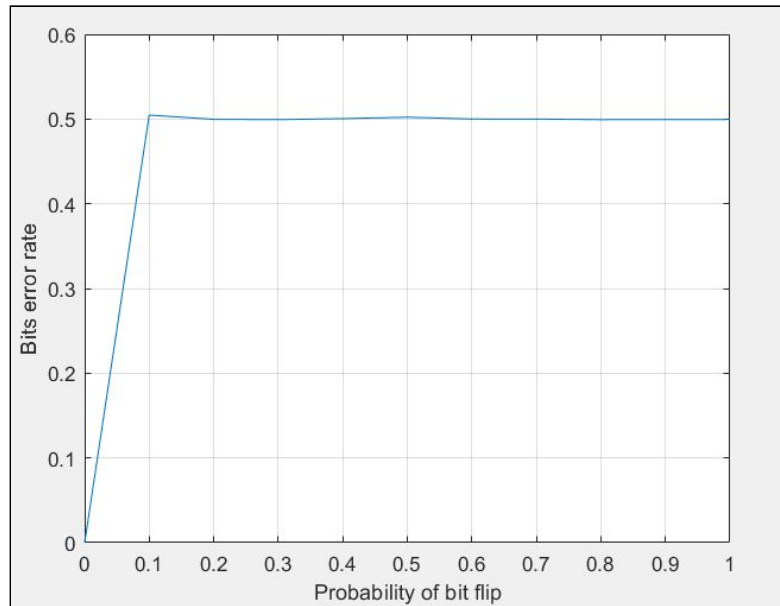
<Problem 3>

- Simulate as 2-(c) with the given generators (random 100002 (>100000) numbers). We can find that for 2-(c), the bit error rate will increase until 1 (i.e.

all the bits will be flipped), while in 3, the bit error rates are oscillated around $p=0.5$.

$$\mathbf{g}_1 = \begin{bmatrix} 1 & 1 & 0 \end{bmatrix},$$

$$\mathbf{g}_2 = \begin{bmatrix} 1 & 0 & 1 \end{bmatrix}.$$



(result)