

MASTER THESIS

Thesis submitted in partial fulfillment of the requirements for the degree of Master of Science in Engineering at the University of Applied Sciences Technikum Wien - Degree Program Mechatronics/Robotics

Virtualisierung eines Echtzeit-Betriebssystems zur Steuerung eines Roboters mit Schwerpunkt auf die Einhaltung der Echtzeit

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Wien, March 27, 2024

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Kurzfassung

Erstellung einer Echtzeit-Robotersteuerungsplattform unter Verwendung von Salamander OS, Xenomai, QEMU und PCV-521 in der Yocto-Umgebung. Die Plattform basiert auf Salamander OS und nutzt Xenomai für Echtzeit- Funktionen. Dazu muss im ersten Schritt die Virtualisierungsplattform evaluiert werden. (QEMU, Hyper-V, Virtual Box, etc.) Als weiterer Schritt folgt die Anbindung eines Roboters über eine VARAN-Bus Schnittstelle. Das gesamte System wird in der Yocto-Umgebung erstellt und konfiguriert. Das Hauptziel der Arbeit ist es, herauszufinden, wie die Integration von Echtzeit-Funktionen und effizienten Kommunikationssystemen in eine Robotersteuerungsplattform die Reaktionszeit und Zuverlässigkeit von Roboteranwendungen verbessern kann

Schlagworte: Schlagwort1, Schlagwort2, Schlagwort3, Schlagwort4

Abstract

Abstract

Keywords: Echtzeit, Virtualisierung, Xenomai, VARAN

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1 Introduction

In today's industrial production and automation, robot systems are well established and of crucial importance. Robots must react to their environment and perform time-critical tasks within strict time constraints. Delays or errors can have catastrophic consequences in some cases. Traditional operating systems, such as Windows or Linux, are often not suitable for these types of real-time requirements as they cannot guarantee deterministic execution times. Therefore, real-time operating systems are required that are specifically designed to react to events within fixed time limits and prioritise the execution of high-priority processes.

The core component of an RTOS that enables real-time capabilities is the kernel. The kernel is responsible for managing system resources, scheduling tasks, and ensuring deterministic behavior. It employs preemptive scheduling mechanisms to allow high-priority tasks to preempt lower-priority tasks, ensuring that time-critical tasks are not delayed. The kernel also implements priority-based scheduling algorithms, such as Rate Monotonic Scheduling (RMS) or Earliest Deadline First (EDF), to schedule tasks based on their priorities and timing constraints. Additionally, RTOS kernels are designed to minimize interrupt latency, which is crucial for real-time applications that require immediate response to external events.

In these RTOS systems, task scheduling is based on so-called priority-based preemptive scheduling. Each task in a software application is assigned a priority. A higher priority means that a faster response is required. Pre-emptive task scheduling ensures a very fast response. Preemptive means that the scheduler can stop a currently running task at any point if it recognises that another task needs to be executed immediately. The basic rule on which priority-based preemptive scheduling is based is that the task with the highest priority that is ready to run is always the task that must be executed. So if both a task with a lower priority and a task with a higher priority are ready to run, the scheduler ensures that the task with the higher priority runs first. The lower priority task is only executed once the higher priority task has been processed. Real-time systems are usually categorised as either soft or hard real-time systems. The difference lies exclusively in the consequences of a violation of the time limits.

Hard real-time is when the system stops operating if a deadline is missed, which can have catastrophic consequences. These consequences include endangering people, damaging machinery, impacting the health and integrity of the environment. Typical examples of this are some control systems in aeroplanes or cars, e.g. internal combustion engines.

Soft real-time exists when a system continues to function even if it cannot perform the tasks within a specified time. If the system has missed the deadline, this has no critical conse-

quences. The system continues to run, although it does so with undesirably lower output quality. Examples of this are multimedia systems where the occasional deviation from a playback rate of 25 frames/sec. leads to stuttering or similar.

1.1 Application Context

This master's thesis was written at SIGMATEK GmbH & Co KG [1].

1.2 State of the art

1.3 Problem and task definition

In robotics, accurate and timely control is crucial to ensure precise movements and reliable interaction with the environment. However, existing robot control systems often have limitations in terms of real-time capabilities and communication efficiency, which can have a detrimental effect on response time and reliability. The challenge is to develop a powerful real-time robot control platform that overcomes these limitations and improves the performance of robotic applications.

1.4 Objective

The main objective of this work is to create a real-time robot control platform that integrates Salamander OS, Xenomai, QEMU and PCV-521 in the Yocto environment.

This integration is expected to lead to a significant improvement in the response time and reliability of robot applications.

2 Methodology

In diesem Abschnitt werden sämtliche theoretischen Konzepte und Randbedingungen sowie praktische Methoden, die zur Erreichung der Zielsetzung dieser Masterarbeit beigetragen haben, detailliert beschrieben.

Zu Beginn der Arbeit wurde eine ausführliche Analyse der einzelnen Virtualisierungsmöglichkeiten von Salamander 4 durchgeführt. Im Besonderen wurde hier die Virtualisierungsperformance von Ubuntu 22.04, Windows 10 und WSL unter QEMU verglichen.

First, a comprehensive literature review is conducted to understand the current trends and challenges in real-time robot control. Based on the literature study, a suitable virtualisation platform is selected.

After the selection of the virtualisation platform, the robot control platform is implemented. This step includes the installation and configuration of Salamander OS, Xenomai, QEMU and PCV-521 in the Yocto environment. Once the platform has been implemented, the robot is connected via a VARAN bus interface. Finally, the platform is evaluated to determine how the integration of real-time functions and efficient communication systems improves the response time and reliability of robot applications.

- Evaluation der Virtualisierungsplattform: Ich werde verschiedene Virtualisierungsplattformen wie QEMU, Hyper-V, Virtual Box usw. evaluieren. Dies ist ein wichtiger Schritt, um die beste Plattform für meine Anforderungen zu finden.

- Erstellung und Konfiguration des Systems in der Yocto-Umgebung: Ich werde das Yocto-Framework verwenden, um mein Embedded Linux System zu erstellen und zu konfigurieren. Yocto bietet viele Tools und Funktionen, die mir bei der Erstellung und Konfiguration meines Systems helfen können.

- Anbindung eines Roboters über eine VARAN-Bus Schnittstelle: Ich plane, einen Roboter in mein System zu integrieren. Ich werde eine VARAN-Bus Schnittstelle verwenden, um eine schnelle und zuverlässige Kommunikation zwischen dem Roboter und dem Steuerungssystem zu gewährleisten.

Trace-cmd was used for tracing the Linux kernel. It can record various kernel events such as interrupts, scheduler decisions, file system activity, function calls in real time. Trace-cmd helped in getting detailed insights into system behaviour and identify reasons for latency [2].

The data that was recorded by trace-cmd was then fed into Kernelshark, which is a graphical front-end tool [3]. It visualizes the recorded kernel trace data in a readable way on an interactive timeline, which facilitated the process of identifying patterns and correlations between events.

By further filtering the displayed events according to specific criteria such as processes, event types or time ranges, the latency issues were analyzed.

Real-time operating system capabilities were provided by Xenomai, which is real-time development framework that extends the Linux kernel. It enables low-latency and deterministic execution of time-critical tasks. Xenomai introduces a dual-kernel approach with a real-time kernel coexisting alongside Linux. A key utility within the Xenomai suite is the latency tool, which benchmarks the timer latency - the time it takes for the kernel to respond to timer interrupts or task activations. The tool creates real-time tasks or interrupt handlers and measures the latency between expected and actual execution times [4].

3 Salamander Operating System

Salamander 4 is the proprietary operating system of SIGMATEK. It is based on Linux version 5.15.94 and integrates Xenomai 3.2, a real-time development environment [4]. Salamander 4 is a 64-bit system, which refers to the x86_64 architecture. The real-time behaviour is achieved through the use of Symmetric Multi-Processing (SMP) and Preemptive Scheduling (PREEMPT). In addition, it utilises IRQPIPE to process interrupts in a way that meets the real-time requirements of the system. The output of the command `uname -a` can be observed in code 1.

```
1 root@sigmatek-core2:~# uname -a
2 Linux sigmatek-core2 5.15.94 #1 SMP PREEMPT IRQPIPE Tue Feb 14 18:18:05 UTC
   2023 x86_64 GNU/Linux
```

Code 1: System information

Xenomai consists of 3 parts. These can be found in the Table 1.

Table 1: Xenomai architecture

| Teil | Beschreibung |
|----------------|--|
| i-pipe | Kernelerweiterung für das Domain-Konzept |
| Xenomai Kernel | Benutzt die i-pipe, und hängt sich als root-Domain ein |
| Xenomai User | Programme (LRT) verwenden diese Bibliothek, um Xenomai Funktionen verwenden zu können. |

Xenomai beruht auf einem Domain-Konzept, das bedeutet, dass alle IRQ an die erste Domain gesendet werden. (root – Domain / Xenomai) Nur wenn diese nichts mehr zu tun hat, dann darf die 2. Domain arbeiten. Das bedeutet, erst wenn alle Xenomai Task in einem Wartezustand sind, arbeiten die Linux-Tasks.

In der Regel unterbricht der Prozessor beim IRQ-Handling seine aktuellen Aktivitäten, um einen Interrupt zu bearbeiten, während die IRQ-Behandlung von Xenomai einen Interrupt-Pipeline-Mechanismus verwendet, der das gleichzeitige Abrufen und Vorbereiten eines anderen Interrupts ermöglicht, während ein Interrupt bearbeitet wird, was die Leistung verbessert und die Latenzzeit verringert.

Was Xenomai4 von seinem Vorgänger Xenomai3 unterscheidet, ist die vollständige Neugestaltung der Ausführungsphase mit hoher Priorität. Dies geschah aus Gründen der Portabilität und

Wartungsfreundlichkeit: I-pipe - die zweite Iteration der ursprünglichen Adeos-Interrupt-Pipeline
- wurde vollständig durch Dovetail ersetzt.

Table 2: Domain specific functions

| Xenomai spezifische Funktionen | Linux spezifische Funktionen |
|---------------------------------------|-------------------------------------|
| Tasks | Dateizugriffe |
| Mutexes, Semaphoren, Events | Netzwerk |

Ein Aufruf dieser Funktionen erfordert die entsprechende Domain. Wenn der Task in der falschen Domain läuft, dann wird ein Domain-Wechsel forciert. Ein Domainwechsel von Xenomai nach Linux geht relativ einfach. Aber der Wechsel von Linux nach Xenomai braucht Unterstützung, und dafür ist die Hilfe des Gatekeepers notwendig. Das bedeutet, der Gatekeeper hilft einem Task von Linux nach Xenomai zu wechseln.

3.1 Task priorities

Es gibt grundsätzlich 4 Gruppen

Table 3: Overview of the priority groups and their relationships

| Prioritätsgruppe | Bereich |
|------------------------------|----------------|
| Xenomai Priorität | 0 bis 99 |
| Linux RT Priorität | 1 bis 99 |
| Linux (Nice Level) Priorität | -20 bis 19 |
| RTK Priorität | 0 bis 14 |

3.2 Memory Management

Es gibt verschiedene Speicherbereiche

Linux/System/Programm Speicher Der Speicher, den Linux und Programme belegt haben. Dieser Speicher ist intern in viele Teile aufgeteilt. (DMA, ...)

LRT-Heap Speicher Speicher den der LRT verwendet, oder welcher über ein CIL Funktionen angefordert wird.

App Heap, App Code, ...

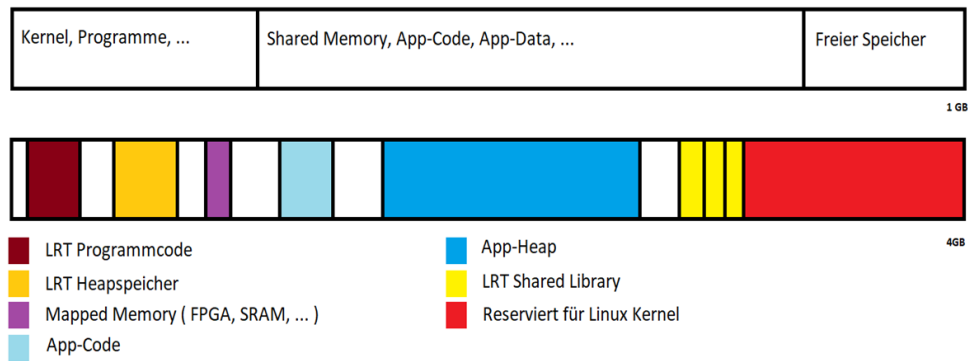


Figure 1: Memory Management

Eine LASAL CPU besteht aus den folgenden Software-Modulen:

- Operating system
- Loader
- Hardware-Klassen

Die Schnittstelle zwischen den einzelnen Modulen wird in Abbildung 2 durch einen Pfeil gekennzeichnet.

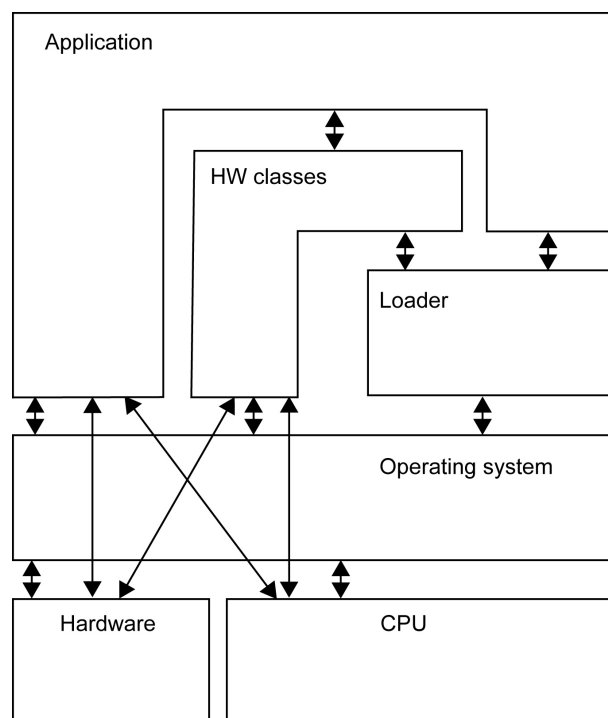


Figure 2: LASAL CPU

4 Comparison of Real-Time Latency

In the initial phase, a comparative latency analysis was conducted between the hardware version and the virtualized version of Salamander 4. For this purpose, the latency tool of the Xenomai test suite was used. The latency was measured under two conditions, idle and CPU-stressed. The goal was to optimize the latency of the virtualisation of Salamander 4 OS to closely match that of the bare metal version.

Table 4: Kernel and Patches

| Kernel | Patches |
|--------|---------|
| . | . |
| . | . |

After analyzing the initial latency of both versions, Trace-cmd and Kernelshark were used to further inspect the reasons that caused this divergence.

4.1 Salamander 4 Bare Metal

Salamander 4 Bare Metal refers to the proprietary hardware of SIGMATEK used to employ the custom operating system, including

4.2 Salamander 4 Virtualisation

In addition to providing Salamander 4 on its own hardware, SIGMATEK has also developed a virtualised version of this operating system. It was developed using Yocto, an open source project that allows customised Linux distributions to be created for embedded systems [5]. The virtualisation runs in a QEMU environment, which is an open source tool for hardware virtualisation [6]. With the help of the script depicted in code 3, Salamander 4 is started together with the necessary hardware components in the QEMU environment. This makes it possible to run Salamander 4 on a variety of host systems, regardless of the specific hardware of the host. Upon generating the necessary files, Yocto generates a QEMU folder with the following components shown in code 2.

```

1  sigma_ibo@localhost:~/Desktop/salamander-image$ ls -l
2  bzImage
3  drive-c
4  ovmf.code.qcow2
5  qemu_def_bkp.sh
6  qemu_def.sh
7  salamander-image-sigmatek-core2.ext4
8  stek-drive-c-image-sigmatek-core2.tar.gz
9  vmlinux

```

Code 2: Contents of QEMU folder for Salamander 4

```

1  #!/bin/sh
2
3  if [ ! -d drive-c/ ]; then
4      echo "Filling_drive-c/"
5      mkdir drive-c/
6      tar -C drive-c/ -xf stek-drive-c-image-sigmatek-core2.tar.gz
7  fi
8
9  exec qemu-system-x86_64 -M pc,accel=kvm -kernel ./bzImage \
10 -m 2048 -drive
    file=salamander-image-sigmatek-core2.ext4,format=raw,media=disk \
11 -append "console=ttyS0_console=ttyl_root=/dev/sda_rw_panic=1_
    sigmatek_lrt.QEMU=1_ip=dhcp_rootfstype=ext4_schedstats=enable" \
12 -net nic,model=e1000,netdev=e1000 -netdev bridge,id=e1000,br=nm-bridge \
13 -fsdev local,security_model=none,id=fsdev0,path=drive-c -device
    virtio-9p-pci,id=fs0,fsdev=fsdev0,mount_tag=/mnt/drive-C \
14 -device vhost-vsock-pci,guest-cid=3,id=vsock0 \
15 -drive if=pflash,format=qcow2,file=ovmf.code.qcow2 \
16 -no-reboot -nographic

```

Code 3: QEMU script for starting Salamander 4 virtualisation

Here is a description of the used components:

- **bzImage**: Compressed Linux kernel image, loaded by QEMU at system start.
- **ovmf.code.qcow2**: Firmware file for QEMU, enables UEFI boot process.
- **qemu_def.sh**: Shell script, starts QEMU with correct parameters to boot Salamander 4 OS.
- **stek-drive-c-image-sigmatek-core2.tar.gz**: Archive containing files for C drive, unpacked and copied to drive-c/ directory by qemu_def.sh script.
- **drive-c**: Directory, serves as C drive for QEMU system, created and filled by qemu_def.sh script..

- **salamander-image-sigmathek-core2.ext4**: Root file system for Salamander 4 OS, used as hard drive for QEMU system.
- **vmlinux**: Uncompressed Linux kernel image, typically used for debugging, contains debugging symbols not present in bzImage.

Figure 3

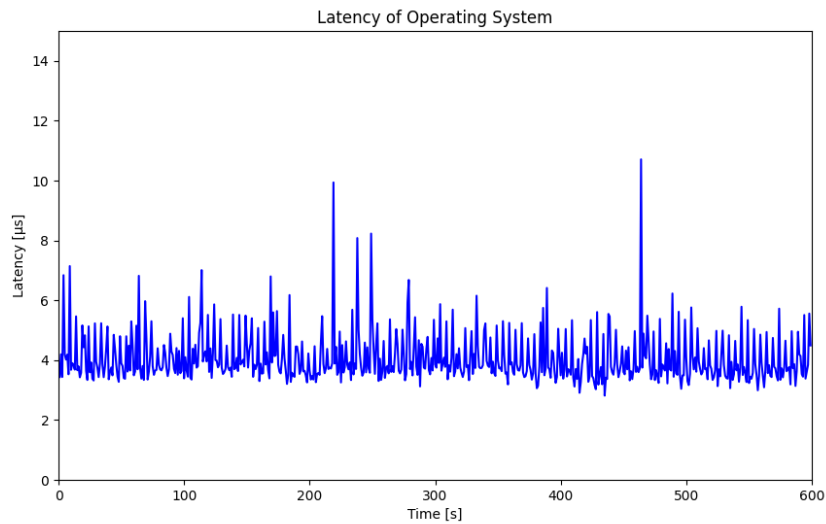


Figure 3: Latency hardware

Figure 4

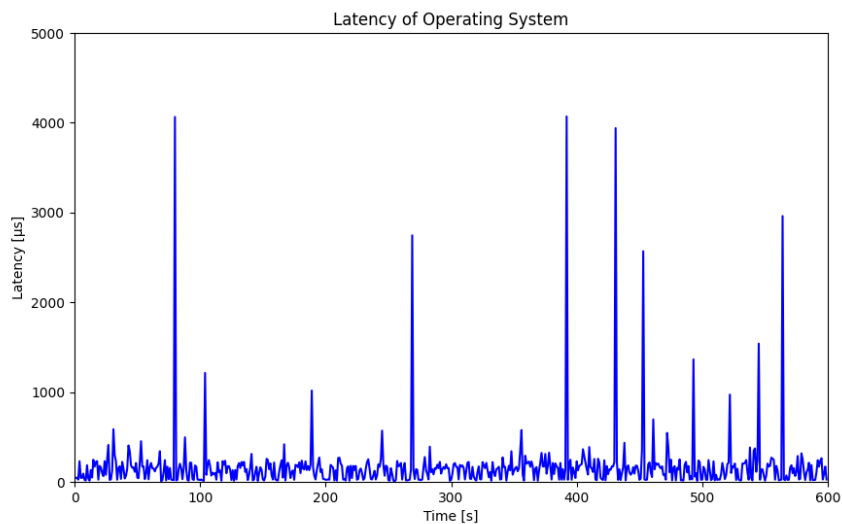


Figure 4: Latency no taskset

Figure 5

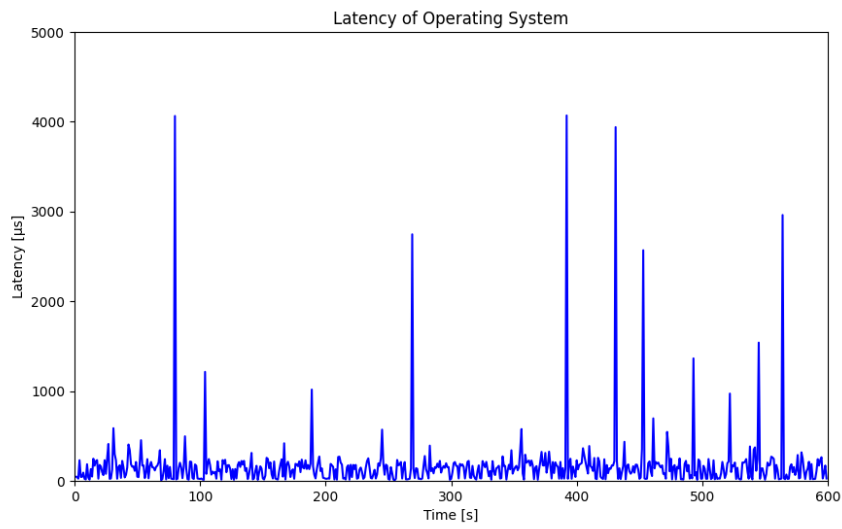


Figure 5: Latency taskset

Figure 6

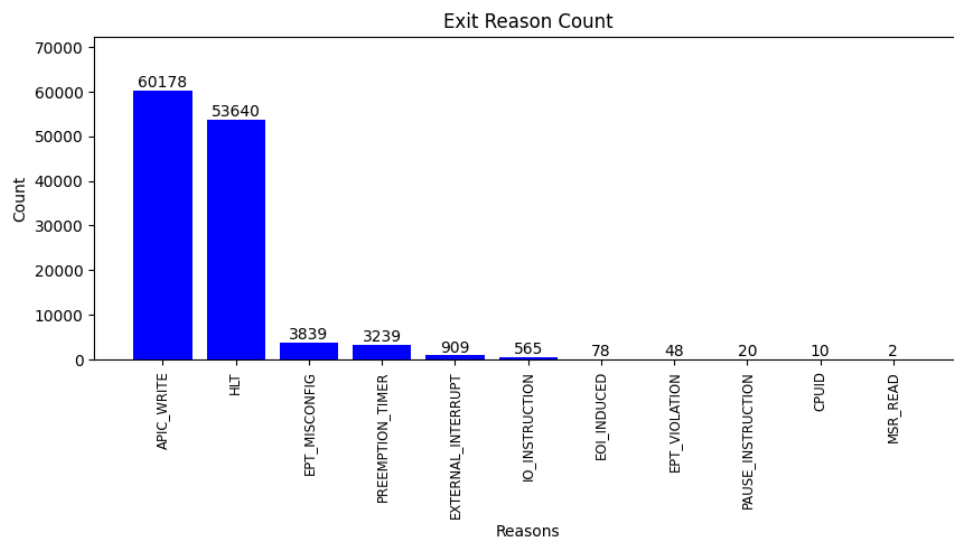


Figure 6: kvm exits

Figure 7

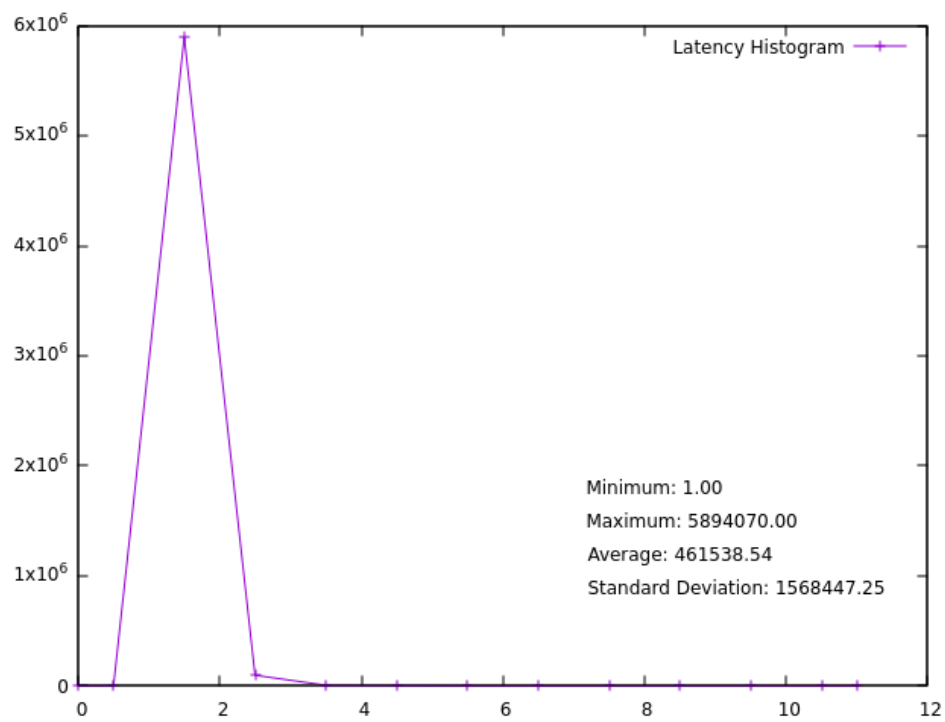


Figure 7: gnuplot latency hardware

Figure 8

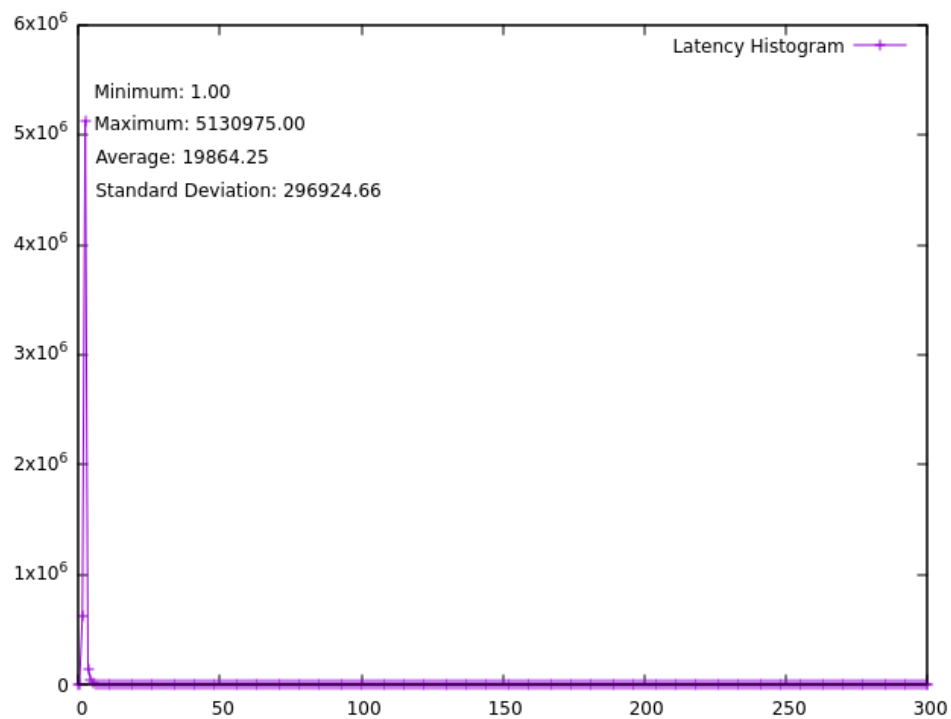


Figure 8: gnuplot latency no taskset

Figure 9

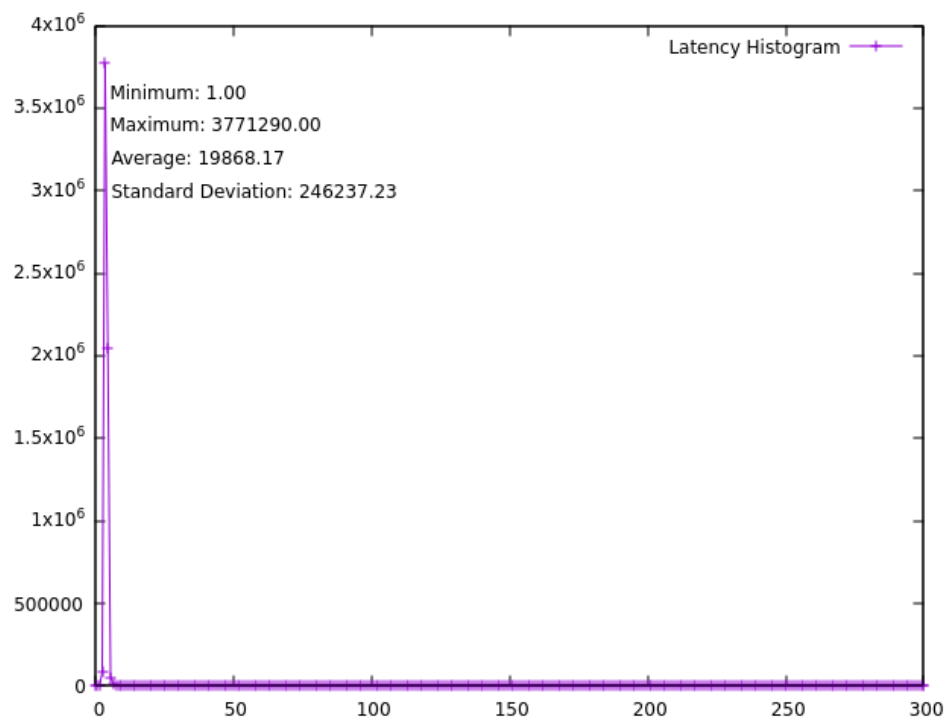


Figure 9: gnuplot latency with taskset

5 Results

VARAN-Bus PCV-521

Windows

Ubuntu

WSL

In Xenomai 3, you need:

The Linux kernel with the Dovetail interface, which allows the kernel to handle real-time tasks with low latency. The Cobalt real-time core, which is responsible for scheduling and executing real-time tasks. In Xenomai 4, you need: The Linux kernel with the EVL interface, which, like Dovetail, enables the kernel to handle real-time tasks with low latency. The libe vl library, which provides a user interface to the EVL core, allowing applications to use the real-time capabilities provided by EVL. So, while both versions require a Linux kernel with a real-time interface (Dovetail for Xenomai 3, EVL for Xenomai 4), Xenomai 3 uses a real-time core (Cobalt), while Xenomai 4 uses a library (libe vl) to provide applications with access to its real-time capabilities.

6 Discussion

7 Summary and Outlook

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List of Abbreviations

CPU Central Processing Unit

QEMU Quick Emulator

A Anhang A

B Anhang B