# MPI Lecture 2: The 'nitty-gritty'





## Overview

- Point-to-Point Pic'n'Mix:
  - Synchronous message passing
  - Buffers & 'blocking' routines
  - Non-blocking patterns
- Support for Recurring Patterns of Communication:
  - Message exchange
  - Collective communication
- Reducing Message Traffic:
  - Using MPI derived types
  - Packed data messages



## Point-to-Point Pic'n'mix

MPI offers many, many options for point-to-point messages:

```
MPI_Send()
MPI_Ssend()
MPI_Bsend()
MPI_Isend()
MPI_Rsend()
MPI_Issend()
MPI_Issend()
MPI_Ibsend()
MPI_Irsend()
```

```
MPI_Recv()
MPI_IRecv()
```

and any sending routine can be paired with either receiving routine



# Many Concepts Bundled up

Blocking vs. Non-blocking

**Buffered or Unbuffered** 

Synchronous vs. Asynchronous

System or User-Buffer

Safe or Unsafe

Portable or not Portable



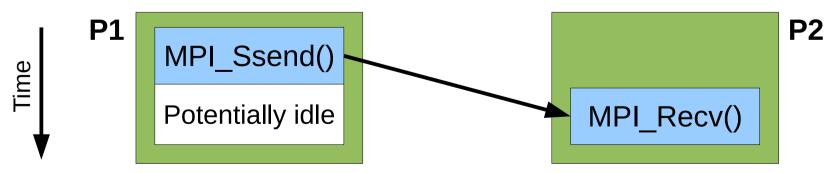
## OK, Let's Unravel it All

- Let's start with MPI\_Ssend(buffer, . . ).
- This stands for synchronous send.
- It is *blocking*, meaning that the call will not return until the function arguments are *safe for* re-use in the program.
- It will only return when a matching receive has been posted and data transmission has completed.



## MPI\_Ssend() Pros and Cons

- The call is **safe** and also **portable** (more of that in a moment).
- But we can anticipate that processes may be
   idle waiting for a matching receive to be posted
   to make the pair.
- Also we can anticipate that blocking, synchronous communication patterns will be most vulnerable to deadlock.





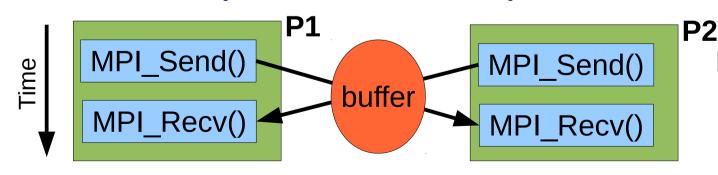
## 'Standard mode': MPI\_Send()

- Also blocking, but potentially asynchronous.
- Potentially? Why?
  - Some MPI implementations use system buffers, but some do not.
- If there is a system buffer, the call will return when the message has been copied from the user-space to the system buffer—i.e. the variable is safe for re-use.
- Without a system buffer, MPI\_Send() will behave just like MPI\_Ssend().



## MPI\_Send() Pros and Cons

- **Asynchronous** communication <u>may</u> allow our programs **run faster**. Why?
- We can use looser communication patterns.
- But copying large messages into a buffer will carry its own overheads too. No free lunch.
- Our programs may not be portable to other MPI implementations (look at 'deadlock.c'):

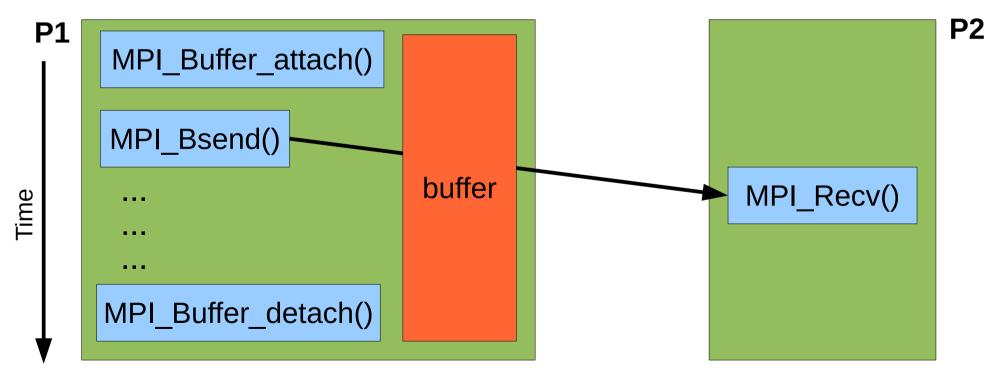


Deadlocks without a system buffer, or if system buffer overfills



## User Space Buffers: MPI\_Bsend()

- A buffered (and so asynchronous), blocking send.
- The buffer is explicitly allocated and managed (attached & detached) by the programmer.





## MPI\_Bsend() Pros and Cons

- Asynchronous and portable, as user guarantees that buffer is available and of sufficient size.
- But memory allocation is expensive and attaching & detaching will carry overheads.
   You'd better be sure that it's worth it!
- Adds extra complexity to your code.
- Different MPI implementations vary in rules around attach and detach calls, so potential portability problems.



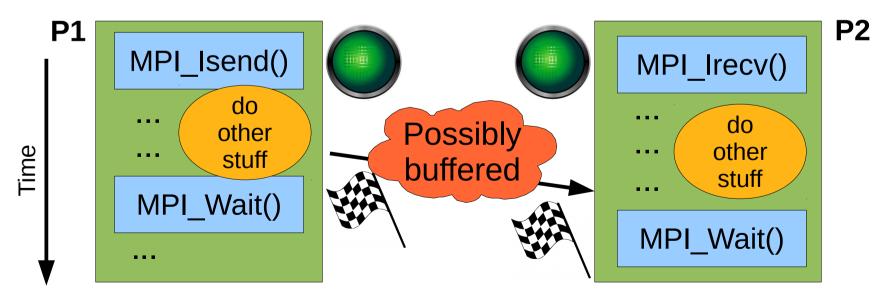
# Non-blocking MPI\_Isend()

- ..and counterpart MPI\_Irecv().
- Is non-blocking and so returns (almost) immediately.
- Communication is only requested at that point.
- Must test or wait for completion later in your code (via extra request argument).
- Function arguments will not be safe for re-use until the communication is complete.
- 'Standard mode' behaviour applies:
  - i.e. there are also explicitly buffered and synchronous (non-buffered) versions.



## MPI\_Isend() Pros and Cons

- Best chance yet to overlap computation and communication. (This manoeuvre is potentially key for good scaling..)
- But code is more complex and we run the risk of introducing some hard to find bugs. So you'd better be sure that it's worth it..





#### There's more!

- To explore at your own leisure..;)
- Notably:
  - 'Ready' send: MPI\_Rsend() (recv already posted)
  - Persistent communications: MPI\_Send\_init() & MPI\_Start().
- Example programs for point-to-point are given in examples 2-4 in the MPI practical.

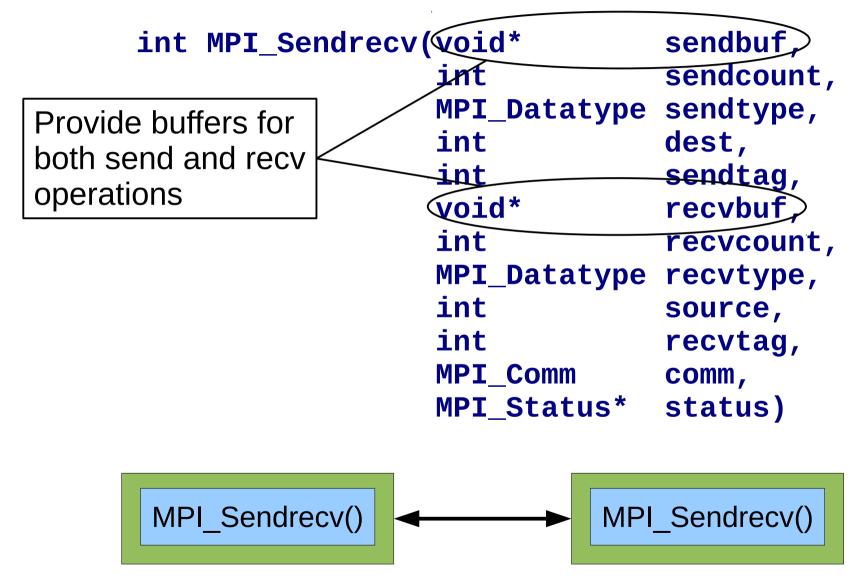


## Support for Recurring Patterns

- We've seen many options for point-to-point communications.
- In principle, point-to-point patterns are all that we need to write any program.
- However, we see some compound patterns of communication cropping up again and again.
- MPI is obliging enough to provide routines for these which save us effort and make our programs faster.



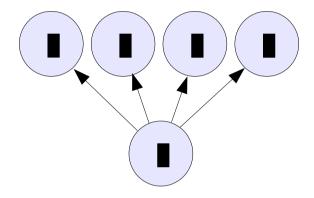
# For example: Message Exchange



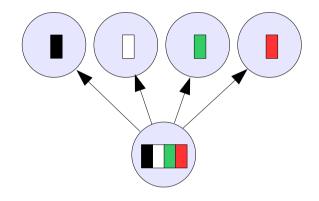


## **Collective Communications**

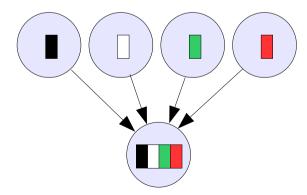
#### broadcast



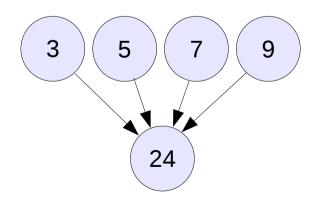
#### scatter



#### gather

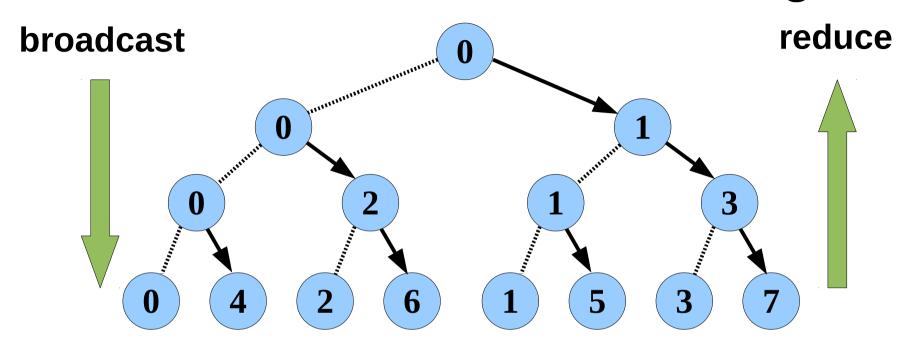


#### reduction





## Tree-Based Load Balancing



- Time taken: 3 units rather than 7, for an 8 process cohort.
- Big benefits of  $O(log_2N)$  vs. O(N) as the cohort size increases,  $log_2(1024) = 10$ . A saving of a factor of 100!
- With system buffering, collective comms are asynchronous, but are a point of synchronisation if no system buffer is available.
- Usage will be shown in example 5 of the practical.
- Benefits are amplified for higher-latency/lower-bandwidth networks



## Collective Comms: There's more...

- MPI offers several more options for collective communication.
- These include:
  - MPI\_Gatherv()
  - MPI\_Alltoall()
  - MPI\_Allreduce()
  - MPI\_Allgather()
  - MPI\_Reduce\_scatter()
- Should they come in handy for you..



## Message Traffic & Latency

- Given that message latencies and communication overheads always exist, it is sensible to try to reduce message traffic.
- One way to do this is to combine information of different datatypes into a single message.
- MPI gives us two options for this:
  - Using MPI derived types.
  - Packed data messages.
- However, beware of keeping a process waiting for data as it will lead to idle time. Potential benefits depend on context..

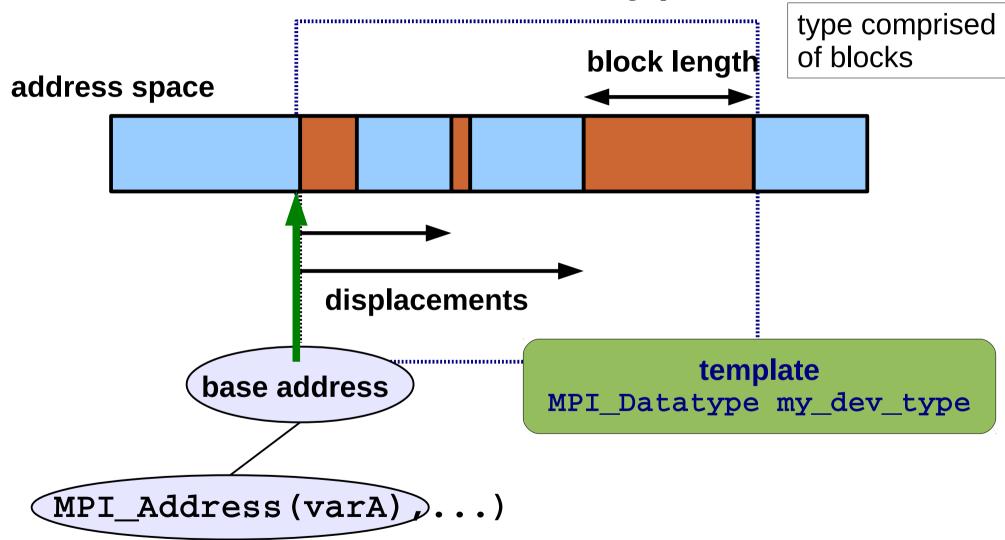


## MPI Derived Types

- C derived types are not guaranteed to be stored as contiguous elements in memory.
- Thus we cannot send them as messages.
- The MPI solution is to construct a sequence of pairs: {(t<sub>0</sub>, d<sub>0</sub>), (t<sub>1</sub>, d<sub>1</sub>), ..., (t<sub>n-1</sub>, d<sub>n-1</sub>)}, where:
  - t<sub>x</sub> is the datatype,
  - $\mathbf{d}_{\mathbf{x}}$  is the displacement in bytes from a base address



# MPI Derived Types



See example 6



## Example 6

```
block_lengths[0] = block_lengths[2] = 1;
block lengths[1] = STRLEN;
typelist[0] = MPI_FLOAT;
typelist[1] = MPI CHAR;
typelist[2] = MPI INT:
displacements[0] = 0;
MPI Address(&station freq, &base address);
MPI_Address(&station_name, &address);
displacements[1] = address - base address;
MPI Address(&station preset num, &address);
displacements[2] = address - base address;
MPI Type struct(NTYPES, block lengths, displacements,
         typelist, &my dev type);
MPI_Type_commit(&my_dev_type);
```



## MPI Derived Types

- Block lengths allow for elements to be in arrays.
- Note: to ensure portability, addresses are stored in the type MPI\_Aint and we call MPI\_Address() to determine the address of an element. (i.e., don't use C's '&'.)
- We pass the base address as the message and the derived type as the datatype to the sending routine and the MPI library constructs the sequence of bytes for us.



## Packed Data Messages

- Another option is to fill a buffer (array) with copies of variables using calls MPI\_Pack().
  - Mixed types and multiples allowed.
- The constructed buffer is contiguous in memory and so can be transmitted as a regular message.
- The buffer can be unpacked at the receiving end with analogous calls to MPI\_unpack().
- Since this approach involves (many) copies, using a derived type may be more efficient.



## MPI2: Recap

- Many point-to-point options. Aspects to consider include blocking, buffering and (a)synchronous communication.
- Buffering raises portability issues.
- Collective communication functions can be very convenient and also highly efficient.
- MPI offers approaches to send compound messages, which can save on network traffic and ameliorate cumulative latency costs.

