Shell Lab

前言

shell lab实验是对于CSAPP第八章的内容的代码补充,需要在开始之前阅读掌握异常控制流(ECF)这一部分内容.中文版第三版524页开始

实验内容

本次实验推荐大家在本机WSL中完成

开始之前

在正式开始实验之前我们先从一个宏观的角度来审视一下这次实验

```
(base) kamilu@LZX:~/shlab$ ls
Makefile mysplit.c trace02.txt trace06.txt trace10.txt trace14.txt tshref
README mystop.c trace03.txt trace07.txt trace11.txt trace15.txt tshref.out
myint.c sdriver.pl trace04.txt trace08.txt trace12.txt trace16.txt
myspin.c trace01.txt trace05.txt trace09.txt trace13.txt tsh.c
```

其中出现了几个my开头的c文件,mysplit.c mystop.c myspin.c myint.c 这些文件的代码相当简单,正如其文件名所示.这些函数并没有参与到本次shell的编译之中,你可以在Makefile中看到它们会被单独编译得到对应的可执行文件,这些文件会被用于测试

接着你可以看到trace开头的16个文件,这些文件是本次实验使用的测试文件,用于验证你是否正确的实现了shell. 从最简单的trace01.txt到复杂的trace16.txt,在逐步完成每一个阶段的任务之后,你会得到一个更加强大,更加全面的shell.

然后你可以发现一个可执行文件 tshref,这是一个shell的参考文件,如果你不确定shell的运行结果或者输出结果,你可以通过这个已经编写好的shell来辅助你判断正确的结果

最后是tsh.c,这是你本次实验需要的主文件,你只需要编写修改这个文件中的代码来完成一个shell

接下来编译整个实验代码

```
make
```

运行之后你得到了tsh,这是一个最初始的shell,运行此文件你会进入一个命令行,尝试输入一些文字,没有什么反应.使用ctrl + d 退出

除此之外还得到了提到的 my开头的几个可执行文件,我们暂时先不去管他们,当然你也可以尝试运行一下,不过没什么输出效果就是了.

接下来我们浏览一下本次实验的核心 tsh.c 这个文件

开头引入了几个C库文件,接着可以看到定义的宏.最大行数,参数上限,任务上限,任务的ID上限等等.

```
/* Misc manifest constants */
#define MAXLINE    1024    /* max line size */
#define MAXARGS    128    /* max args on a command line */
#define MAXJOBS    16    /* max jobs at any point in time */
#define MAXJID    1<<16    /* max job ID */</pre>
```

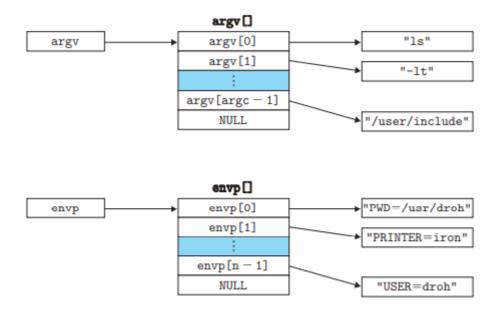
接着是对任务JOB state的一些宏,我们放到后面去说明

接着是一些全局变量

其中 extern char **environ 这个变量是环境变量,大概率会在 execve 函数调用的时候被作为第三个参数传入;jobs则是任务列表

```
#include <unistd.h>
int execve(const char *filename, const char *argv, const char *envp[]);
// 如果成功不返回,如果错误返回-1
```

execve函数加载并运行可执行目标文件filename,并且带参数列表argv和环境变量列表envp,与fork返回两次不同.execve调用一次并且不返回,只有出现错误,比如找到不到filename才会返回到调用程序



其中参数列表的数据结构如上图所示

argv变量指向一个null结尾的指针数组,每个指针指向一个参数字符串,通常来说argv[0]是可执行目标文件的名字 envp变量指向一个null结尾的指针数据,每个指针指向一个字符串,每个串都是类似"name=value"的键值对

接下来是我们本次实验需要完成的一些函数,这里已经体现列了出来

```
/* Here are the functions that you will implement */
void eval(char *cmdline);
int builtin_cmd(char **argv);
void do_bgfg(char **argv);
void waitfg(pid_t pid);

void sigchld_handler(int sig);
void sigtstp_handler(int sig);
void sigint_handler(int sig);
```

然后是一些本次实验中作者提供的一些辅助函数,我们需要在完成这几个函数的同时利用到作者提供的这些辅助 函数来执行或者或许信息,具体的函数内容我们放到对应的出现位置再去讲解

接下来就是main函数了,先判断一下有没有携带-h-v-p参数,如果使用了那么执行相关函数.接下来使用Signal将四个中断信号注册到对应的方法.initjob将所有的任务初始化,进入while(1)死循环,输出prompt即tsh>,刷新标准输出流然后进入eval,接着再次刷新,进入循环.

现在我们的eval函数暂时未完成,所以整个程序运行的流程大致如此.

除此之外我们注意到Makefile中提供了很方便的测试指令,比如想要测试trace11是否通过,可以使用

```
make test11
```

检测是否正确,可以使用

```
make rtest11
```

这里由于两次pid不同所以输出会有一点点差异,我们可以对比两次结果来判断是否正确实现了shell

trace01

```
#
# trace01.txt - Properly terminate on EOF.
#
CLOSE
WAIT
```

第一关的输入是CLOSE和WAIT,这两个命令并没有什么作用,所以什么也不需要改动

trace02

```
#
# trace02.txt - Process builtin quit command.
#
quit
WAIT
```

第二关需要我们实现内置命令"quit",当输入quit之后退出当前shell.我们完善一下eval函数解析命令行

```
void eval(char *cmdline)
{
    char *argv[MAXARGS];
   char buf[MAXLINE];
   int bg;
    strcpy(buf,cmdline);
    bg=parseline(buf,argv);
    if (argv[0] == NULL) return;
    builtin_cmd(argv);//读取到quit就退出了
   return;
}
int builtin_cmd(char **argv)
{
    if(!strcmp(argv[0],"quit")) exit(0);
   return 0; /* not a builtin command */
}
```

这里的parseline是一个解析命令行的函数,可以将传入的字符串通过空格拆分,并将拆分后的结果保存在argv中

它的返回值用于判断是否是一个后台进程,如果是返回1否则返回0

然后判断一下是否是空输入,是则直接返回

接下来实现builtin cmd对于内置命令的解析,判断第一个参数是否是"quit",如果是则直接退出

trace03

```
#
# trace03.txt - Run a foreground job.
#
/bin/echo tsh> quit
quit
```

第三关加入了一个新的测试命令,其中/bin/echo是Linux中一个文件,一个可执行程序.它的作用是输出后面的字符串.你可以直接在命令行中使用这个指令

```
echo hello world
```

这里直接使用echo是因为默认会把/bin目录加入到环境变量之中,它与使用/bin/echo是完全等价的

所以如果想要在我们的shell中执行/bin/echo,我们需要在解析完成之后创建一个子进程,然后交由操作系统去执行

```
void eval(char *cmdline)
    char *argv[MAXARGS];
    char buf[MAXLINE];
    int bg;
    pid t pid;
    strcpy(buf,cmdline);
    bg=parseline(buf,argv);
    if (argv[0] == NULL) return;
    if (!builtin cmd(argv)) {
        if ((pid = fork() == ∅)) {
            if (execve(argv[0],argv,environ) < 0) {</pre>
                 printf("%s: Command not found\n",argv[0]);
                 exit(0);
            }
        }
    }
    return;
}
```

先使用builtin_cmd判断是否是内置的指令,如果不是则fork一个子进程,使用execve函数去执行当前的参数.execve函数的三个参数分别是argv[0]即"/bin/echo"这个程序的名字,argv其余参数包括需要输出的字符串,environ环境变量信息,这个值直接使用整个shell中的环境变量信息

trace04

```
# trace04.txt - Run a background job.
/bin/echo -e tsh> ./myspin 1 \046
./myspin 1 &
```

第四关首先是-e,这是echo的一个参数,用于激活转义字符,这里我们不需要考虑太多

然后是运行了一个程序myspin,参数是1,并且在后台执行.

这里的myspin函数实现很简单,就是休眠n秒,n为传入的参数.这里是休眠1s

这里我先放出修改后的代码,然后我们一点一点看

```
void eval(char *cmdline) {
  char *argv[MAXARGS];
  char buf[MAXLINE];
 int bg;
 pid t pid;
  strcpy(buf, cmdline);
  bg = parseline(buf, argv);
 if (argv[0] == NULL)
   return;
  sigset_t mask_all, mask_one, prev;
  sigfillset(&mask all);
  sigemptyset(&mask_one);
  sigaddset(&mask_one, SIGCHLD);
  if (!builtin cmd(argv)) {
    sigprocmask(SIG_BLOCK, &mask_one, &prev);
    if ((pid = fork()) == 0) {
     sigprocmask(SIG SETMASK, &prev, NULL);
     setpgid(∅, ∅);
     if (execve(argv[0], argv, environ) < 0) {
       printf("%s: Command not found\n", argv[0]);
       exit(0);
     }
    }
    if (!bg) {
      // 对于前台进程,添加job后解除阻塞,并通过waitfg等待子进程结束后回收
      sigprocmask(SIG_BLOCK, &mask_all, NULL);
      addjob(jobs, pid, FG, cmdline);
     sigprocmask(SIG_SETMASK, &prev, NULL);
     waitfg(pid);
    } else {
      // 后台进程不需要等待子进程,进程结束之后收到SIGCHLD信号回收即可
      sigprocmask(SIG_BLOCK, &mask_all, NULL);
      addjob(jobs, pid, BG, cmdline);
      sigprocmask(SIG_SETMASK, &prev, NULL);
```

```
printf("[%d] (%d) %s", pid2jid(pid), pid, cmdline);
    }
  }
  return;
}
void waitfg(pid_t pid) {
  sigset_t mask;
  sigemptyset(&mask);
  while (pid == fgpid(jobs))
    sigsuspend(&mask);
  return;
void sigchld_handler(int sig) {
 pid_t pid;
  int status;
  sigset_t mask_all, prev;
  sigfillset(&mask all);
  while ((pid = waitpid(-1, &status, WNOHANG | WUNTRACED)) > 0) {
    if (WIFEXITED(status)) // 正常退出 delete
      sigprocmask(SIG_BLOCK, &mask_all, &prev);
      deletejob(jobs, pid);
      sigprocmask(SIG_SETMASK, &prev, NULL);
    }
  }
  return;
}
```

如果你对这部分代码感到疑惑,请重新阅读第八章

首先是加入了信号和阻塞,在fork之前阻塞SIGCHLD信号,以防子进程和父进程竞争.判断该进程是前台进程还是后台进程之后使用addjob添加对应的job类型,然后解除阻塞,再去响应子进程结束信号SIGCHLD,确保addjob操作deletejob操作顺序对应

如果是前台进程比如echo,那么使用waitfg等待前台进程结束,waitfg的具体实现就是调用fgpid找到当前的前台进程,然后判断处于前台的进程pid是否是该pid,即前台进程仍然没有执行结束,那么挂起当前进程sigsuspend等待结束

如果后台进程比如myspin,那么使用addjob加入任务队列之后就不需要关心了,使用printf输出一句话即可(与make rtest04的结果相同)

这里使用了开头定义的宏BG FG

对于子进程结束的信号SIGCHLD的处理则完成对应的sigchld_handler,这里的waitpid使用的第三个参数位置是WNOHANG | WUNTRACED代表立即返回,然后判断一下返回的状态,如果是正常退出那么deletejob删除对应任务即可

除此之外还需要注意在创建子进程之后使用了setpgid(0,0),这里的作用是将子进程的gpid组设置与父进程pid相同

运行make test04后结果如下,其中pid可能不同,导致pid2jid的结果不同

这里简单解释一下,2同英文to,所以日常编程之中经常使用2来代指to表示转换,这里的pid2jid类似pid_to_jid

```
./sdriver.pl -t trace04.txt -s ./tsh -a "-p"
#
# trace04.txt - Run a background job.
#
tsh> ./myspin 1 &
[1] (980) ./myspin 1 &
```

trace05

```
#
# trace05.txt - Process jobs builtin command.
#
/bin/echo -e tsh> ./myspin 2 \046
./myspin 2 &

/bin/echo -e tsh> ./myspin 3 \046
./myspin 3 &

/bin/echo tsh> jobs
jobs
```

第五关是分别运行了前台echo,后台myspin,前台echo,后台myspin,前台echo,前台jobs

这里需要实现一个内置命令jobs,功能是显示目前任务列表中的所有任务以及所有属性,这里可以利用listjob函数

```
int builtin_cmd(char **argv) {
   if (!strcmp(argv[0], "quit")) {
      exit(0);
   }
   if (!strcmp(argv[0], "jobs")) {
      listjobs(jobs);
      return 1;
   }
   return 0; /* not a builtin command */
}
```

trace06

```
#
# trace06.txt - Forward SIGINT to foreground job.
#
/bin/echo -e tsh> ./myspin 4
```

```
./myspin 4

SLEEP 2
INT
```

第六关加入了INT,表示接收到了中断信号SIGINT(即CTRL_C),那么结束前台进程,此时需要我们完成sigint_handler

```
void sigchld_handler(int sig) {
 pid_t pid;
 int status;
 sigset_t mask_all, prev;
 sigfillset(&mask_all);
 while ((pid = waitpid(-1, &status, WNOHANG | WUNTRACED)) > ∅) {
    if (WIFEXITED(status)) // 正常退出 delete
   {
      sigprocmask(SIG_BLOCK, &mask_all, &prev);
      deletejob(jobs, pid);
      sigprocmask(SIG_SETMASK, &prev, NULL);
    } else if (WIFSIGNALED(status)) // 信号退出 delete
      sigprocmask(SIG_BLOCK, &mask_all, &prev);
      printf("Job [%d] (%d) terminated by signal %d\n", pid2jid(pid), pid,
WTERMSIG(status));
      deletejob(jobs, pid);
      sigprocmask(SIG_SETMASK, &prev, NULL);
   }
  }
 return;
}
void sigint_handler(int sig) {
   sigset_t mask_all, prev;
    sigfillset(&mask_all);
    sigprocmask(SIG_SETMASK, &mask_all, &prev);
    pid_t pid = fgpid(jobs);
    sigprocmask(SIG_SETMASK, &prev, NULL);
    if (pid > 0) {
        kill(-pid,sig);
    }
    return;
}
```

这里修改了两个地方,首先是完成了sigint_handler,如果接收到了INT信号,则通过fgpid找到当前的前台进程,使用kill杀死这个进程.这里kill的第一个参数是负数,代表杀死这个进程组的所有进程,对子进程及其所有后代都发送终止信号

另外值得一提的是对于INT终止类型我们将printf的输出内容写在sigchld_handler中而不是sigint_handler,实际上如果你写在sigint_handler中对于这道题来说也是可以的,但是意义完全不同

• 如果写在sigint_handler,那么代表每次处理INT中断的时候输出这句话,与子进程无关

如果写在sigchld_handler,那么代表是子进程结束了,并且捕捉到是来自WIFSIGNALED的信号退出方式,对应INT,所以输出这句话

后者的写法才是正确的,前者的写法会导致来自外界的中断信号无法别识别,你会在trace16中看到差距 此外应该注意对于WIFSIGNALED的处理中应该先printf后delete,其实顺序无所谓,只是为了和tshref的输出相同 trace07

第七关没有任何新的命令,检测你之前是否是正确完成了所有的函数,有没有投机取巧如果你之前都正确完成了,那么不需要改动直接通过

trace08

```
#
# trace08.txt - Forward SIGTSTP only to foreground job.
#
/bin/echo -e tsh> ./myspin 4 \046
./myspin 4 &
/bin/echo -e tsh> ./myspin 5
./myspin 5

SLEEP 2
TSTP
/bin/echo tsh> jobs
jobs
```

第八关加入了TSTP(CTRL_Z),当接收到了TSTP中断信号(即CTRL_Z),将前台进程挂起,然后输出被挂起的任务,和INT类似,完成两个函数

```
void sigchld_handler(int sig) {
  pid_t pid;
  int status;
  sigset_t mask_all, prev;
  sigfillset(&mask_all);
  while ((pid = waitpid(-1, &status, WNOHANG | WUNTRACED)) > 0) {
    if (WIFEXITED(status)) // 正常退出 delete
    {
       sigprocmask(SIG_BLOCK, &mask_all, &prev);
       deletejob(jobs, pid);
       sigprocmask(SIG_SETMASK, &prev, NULL);
    } else if (WIFSIGNALED(status)) // 信号退出 delete
    {
       sigprocmask(SIG_BLOCK, &mask_all, &prev);
       printf("Job [%d] (%d) terminated by signal %d\n", pid2jid(pid), pid,
       WTERMSIG(status));
       deletejob(jobs, pid);
```

```
sigprocmask(SIG_SETMASK, &prev, NULL);
    } else if (WIFSTOPPED(status)) {
      sigprocmask(SIG_BLOCK, &mask_all, &prev);
      struct job_t *job=getjobpid(jobs,pid);
      job->state=ST;
      printf("Job [%d] (%d) stopped by signal
%d\n",pid2jid(pid),pid,WSTOPSIG(status));
      sigprocmask(SIG_SETMASK, &prev, NULL);
    }
  }
  return;
}
void sigtstp_handler(int sig) {
    sigset_t mask_all, prev;
    sigfillset(&mask_all);
    sigprocmask(SIG_SETMASK, &mask_all, &prev);
    pid_t pid = fgpid(jobs);
    sigprocmask(SIG_SETMASK, &prev, NULL);
    if (pid > 0) {
        kill(-pid,sig);
    return;
}
```

其中sigtstp_handler与sigint_handler实现完全相同,在sigchld_handler中找到对于TSTP(WIFSTOPPED)的信号,通过getjobpid找到对应任务并将其state修改为ST即可

trace09

```
#
# trace09.txt - Process bg builtin command
#
/bin/echo -e tsh> ./myspin 4 \046
./myspin 4 &

/bin/echo -e tsh> ./myspin 5
./myspin 5

SLEEP 2
TSTP
/bin/echo tsh> jobs
jobs
/bin/echo tsh> bg %2
bg %2
/bin/echo tsh> jobs
jobs
```

其中加入了bg命令,我们可以在开头的注释中找到对应命令所表示的含义

```
/*
 * Jobs states: FG (foreground), BG (background), ST (stopped)
 * Job state transitions and enabling actions:
 * FG -> ST : ctrl-z
 * ST -> FG : fg command
 * ST -> BG : bg command
 * BG -> FG : fg command
 * At most 1 job can be in the FG state.
 */
```

其中bg command代表将一个ST进程变为BG进程,即将挂起的进程调入后台执行,我们需要完成builtin_cmd和do_bgfg函数

```
int builtin_cmd(char **argv) {
 if (!strcmp(argv[0], "quit")) {
    exit(0);
  }
  if (!strcmp(argv[0],"jobs")) {
   listjobs(jobs);
    return 1;
  }
  if (!strcmp(argv[0], "bg")) {
    do_bgfg(argv);
   return 1;
  }
  return 0; /* not a builtin command */
void do_bgfg(char **argv) {
    int jid;
    struct job t *job;
    if (!strcmp(argv[0], "bg")) {
        jid = atoi(&argv[1][1]);
        job=getjobjid(jobs,jid);
        job->state=BG;
        kill(-(job->pid),SIGCONT);
        printf("[%d] (%d) %s", jid, job->pid, job->cmdline);
    }
    return;
}
```

这里使用了atoi的C库函数,将字符串转为int类型,因为bg的命令的输入格式是 % + jid, 所以使用argv[1][1]获取该jid的字符串的起始地址

通过getjobjid找到对应的job后修改state为BG,同时发送SIGCONT信号使其执行.单纯修改state毫无意义,必须发送信号使其真正由挂起状态转入运行状态

trace10

第十关类似,只不过是fg命令

```
/*
 * Jobs states: FG (foreground), BG (background), ST (stopped)
 * Job state transitions and enabling actions:
 * FG -> ST : ctrl-z
 * ST -> FG : fg command
 * ST -> BG : bg command
 * BG -> FG : fg command
 * At most 1 job can be in the FG state.
 */
```

可以看到fg命令有两种情况,分别是将一个挂起进程调入前台,和将一个后台进程调入前台

```
int builtin_cmd(char **argv) {
 if (!strcmp(argv[0], "quit")) {
    exit(0);
 }
 if (!strcmp(argv[0],"jobs")) {
   listjobs(jobs);
   return 1;
 }
 if (!strcmp(argv[0], "bg") || !strcmp(argv[0], "fg")) {
   do_bgfg(argv);
   return 1;
 }
 return 0; /* not a builtin command */
}
void do_bgfg(char **argv) {
    int jid;
    struct job_t *job;
    if (!strcmp(argv[0], "bg")) {
       jid = atoi(&argv[1][1]);
        job=getjobjid(jobs,jid);
        job->state=BG;
        kill(-(job->pid),SIGCONT);
        printf("[%d] (%d) %s", jid, job->pid, job->cmdline);
    } else if (!strcmp(argv[0], "fg")) {
        jid=atoi(&argv[1][1]);
        job=getjobjid(jobs, jid);
        if (job->state == ST) {
            //如果是挂起程序就重启并且转到前台,等待结束
            job->state=FG;
            kill(-(job->pid),SIGCONT);
            waitfg(job->pid);
        } else if (job->state == BG) {
```

```
//如果是后台程序就转到前台并等待结束
    job->state=FG;
    waitfg(job->pid);
    }
}
return;
}
```

前台程序和后台程序主要差异就是shell会不会等待你结束,前台shell主动等待waitfg,然后回收掉,后台就是什么时候结束什么时候返回SIGCHILD,然后回收掉

trace11

不需要修改

trace12

不需要修改

trace13

不需要修改

trace14

本关主要是测试所有的命令,判断是否正确。这一关没有技术难关,主要就是对照输出结果完善你的shell对于 错误的处理

注意大小写! 注意标点! 注意空格!

```
void eval(char **cmdline) {
    if (execve(argv[₀], argv, environ) < ₀) {</pre>
        printf("%s: Command not found\n", argv[0]);
        exit(0);
      }
}
void do_bgfg(char **argv) {
    int jid;
    struct job_t *job;
    if(argv[1]==NULL){
        printf("%s command requires PID or %%jobid argument\n",argv[0]);
        return;
    if(argv[1][0]=='%'){
        jid=atoi(&argv[1][1]);
        job=getjobjid(jobs,jid);
        if(job==NULL){
```

```
printf("%%%d: No such job\n",jid);
           return ;
       }
   } else if(isdigit(argv[1][0])){
       jid = atoi(argv[1]);
       job=getjobjid(jobs,jid);
       if(job==NULL){
           printf("(%d): No such process\n",jid);
           return ;
       }
   } else {
       printf("%s: argument must be a PID or %%jobid\n",argv[0]);
       return ;
   if (!strcmp(argv[0], "bg")) {
       job->state=BG;
       kill(-(job->pid),SIGCONT);
       printf("[%d] (%d) %s", jid, job->pid, job->cmdline);
   } else if (!strcmp(argv[0], "fg")) {
       if (job->state == ST) {
           //如果是挂起程序就重启并且转到前台,等待结束
           job->state=FG;
           kill(-(job->pid),SIGCONT);
           waitfg(job->pid);
       } else if (job->state == BG) {
           //如果是后台程序就转到前台并等待结束
           job->state=FG;
           waitfg(job->pid);
       }
   return;
}
```

trace15

还是大小写和标点空格

trace16

这里使用了myint,对于INT handler的处理需要注意,这一点我们在前文trace06已经提及过了,这里不再赘述

下附完整代码

```
/*
 * tsh - A tiny shell program with job control
 *
 * lzx 2019300003075
 */
#include <ctype.h>
#include <errno.h>
```

```
#include <signal.h>
#include <stdio.h>
#include <stdlib.h>
#include <string.h>
#include <sys/types.h>
#include <sys/wait.h>
#include <unistd.h>
/* Misc manifest constants */
#define MAXLINE 1024 /* max line size */
#define MAXARGS 128  /* max args on a command line */
#define MAXJOBS 16  /* max jobs at any point in time */
#define MAXJID 1 << 16 /* max job ID */</pre>
/* Job states */
#define UNDEF 0 /* undefined */
#define FG 1 /* running in foreground */
#define BG 2 /* running in background */
#define ST 3 /* stopped */
* Jobs states: FG (foreground), BG (background), ST (stopped)
* Job state transitions and enabling actions:
   FG -> ST : ctrl-z
     ST -> FG : fg command
     ST -> BG : bg command
 * BG -> FG : fg command
 * At most 1 job can be in the FG state.
*/
/* Global variables */
extern char **environ; /* defined in libc */
char prompt[] = "tsh> "; /* command line prompt (DO NOT CHANGE) */
/* The job struct */
struct job_t {
 pid_t pid;
                     /* job PID */
                     /* job ID [1, 2, ...] */
 int jid;
                     /* UNDEF, BG, FG, or ST */
 int state;
 char cmdline[MAXLINE]; /* command line */
};
struct job t jobs[MAXJOBS]; /* The job list */
/* End global variables */
/* Function prototypes */
/* Here are the functions that you will implement */
void eval(char *cmdline);
int builtin cmd(char **argv);
void do_bgfg(char **argv);
void waitfg(pid_t pid);
```

```
void sigchld_handler(int sig);
void sigtstp_handler(int sig);
void sigint_handler(int sig);
/* Here are helper routines that we've provided for you */
int parseline(const char *cmdline, char **argv);
void sigquit_handler(int sig);
void clearjob(struct job_t *job);
void initjobs(struct job_t *jobs);
int maxjid(struct job_t *jobs);
int addjob(struct job_t *jobs, pid_t pid, int state, char *cmdline);
int deletejob(struct job_t *jobs, pid_t pid);
pid_t fgpid(struct job_t *jobs);
struct job_t *getjobpid(struct job_t *jobs, pid_t pid);
struct job_t *getjobjid(struct job_t *jobs, int jid);
int pid2jid(pid_t pid);
void listjobs(struct job_t *jobs);
void usage(void);
void unix_error(char *msg);
void app_error(char *msg);
typedef void handler_t(int);
handler_t *Signal(int signum, handler_t *handler);
 * main - The shell's main routine
int main(int argc, char **argv) {
 char c;
  char cmdline[MAXLINE];
 int emit prompt = 1; /* emit prompt (default) */
  /* Redirect stderr to stdout (so that driver will get all output
  * on the pipe connected to stdout) */
 dup2(1, 2);
  /* Parse the command line */
  while ((c = getopt(argc, argv, "hvp")) != EOF) {
    switch (c) {
    case 'h': /* print help message */
      usage();
    case 'v': /* emit additional diagnostic info */
      verbose = 1;
      break;
    case 'p':
                   /* don't print a prompt */
      emit_prompt = 0; /* handy for automatic testing */
      break;
    default:
      usage();
    }
  }
```

```
/* Install the signal handlers */
  /* These are the ones you will need to implement */
 Signal(SIGINT, sigint_handler); /* ctrl-c */
 Signal(SIGTSTP, sigtstp_handler); /* ctrl-z */
 Signal(SIGCHLD, sigchld_handler); /* Terminated or stopped child */
 /* This one provides a clean way to kill the shell */
 Signal(SIGQUIT, sigquit_handler);
  /* Initialize the job list */
 initjobs(jobs);
  /* Execute the shell's read/eval loop */
 while (1) {
   /* Read command line */
   if (emit prompt) {
     printf("%s", prompt);
     fflush(stdout);
    if ((fgets(cmdline, MAXLINE, stdin) == NULL) && ferror(stdin))
      app_error("fgets error");
    if (feof(stdin)) { /* End of file (ctrl-d) */
     fflush(stdout);
     exit(0);
    }
   /* Evaluate the command line */
    eval(cmdline);
   fflush(stdout);
   fflush(stdout);
 exit(0); /* control never reaches here */
}
 * eval - Evaluate the command line that the user has just typed in
* If the user has requested a built-in command (quit, jobs, bg or fg)
 * then execute it immediately. Otherwise, fork a child process and
 * run the job in the context of the child. If the job is running in
 * the foreground, wait for it to terminate and then return. Note:
 * each child process must have a unique process group ID so that our
 * background children don't receive SIGINT (SIGTSTP) from the kernel
 * when we type ctrl-c (ctrl-z) at the keyboard.
 */
void eval(char *cmdline) {
 char *argv[MAXARGS];
 char buf[MAXLINE];
 int bg;
 pid_t pid;
  strcpy(buf, cmdline);
```

```
bg = parseline(buf, argv);
 if (argv[0] == NULL)
   return;
 sigset t mask all, mask one, prev;
 sigfillset(&mask all);
 sigemptyset(&mask_one);
 sigaddset(&mask one, SIGCHLD);
 if (!builtin_cmd(argv)) {
   sigprocmask(SIG_BLOCK, &mask_one, &prev);
   if ((pid = fork()) == 0) {
     sigprocmask(SIG_SETMASK, &prev, NULL);
     setpgid(∅, ∅);
     if (execve(argv[₀], argv, environ) < ₀) {
       printf("%s: Command not found\n", argv[0]);
       exit(0);
     }
   }
   if (!bg) {
     // 对于前台进程·添加job后解除阻塞·并通过waitfg等待子进程结束后回收
     sigprocmask(SIG_BLOCK, &mask_all, NULL);
     addjob(jobs, pid, FG, cmdline);
     sigprocmask(SIG_SETMASK, &prev, NULL);
     waitfg(pid);
   } else {
     // 后台进程不需要等待子进程,进程结束之后收到SIGCHLD信号回收即可
     sigprocmask(SIG_BLOCK, &mask_all, NULL);
     addjob(jobs, pid, BG, cmdline);
     sigprocmask(SIG_SETMASK, &prev, NULL);
     printf("[%d] (%d) %s", pid2jid(pid), pid, cmdline);
   }
 }
 return;
}
 * parseline - Parse the command line and build the argv array.
* Characters enclosed in single quotes are treated as a single
* argument. Return true if the user has requested a BG job, false if
 * the user has requested a FG job.
int parseline(const char *cmdline, char **argv) {
 static char array[MAXLINE]; /* holds local copy of command line */
 char *delim;
                           /* points to first space delimiter */
 int argc;
                           /* number of args */
                           /* background job? */
 int bg;
 strcpy(buf, cmdline);
 buf[strlen(buf) - 1] = ' '; /* replace trailing '\n' with space */
 while (*buf && (*buf == ' ')) /* ignore leading spaces */
   buf++;
```

```
/* Build the argv list */
 argc = 0;
 if (*buf == '\'') {
   buf++;
   delim = strchr(buf, '\'');
 } else {
   delim = strchr(buf, ' ');
 while (delim) {
   argv[argc++] = buf;
   *delim = '\0';
   buf = delim + 1;
   while (*buf && (*buf == ' ')) /* ignore spaces */
     buf++;
   if (*buf == '\'') {
     buf++;
     delim = strchr(buf, '\'');
   } else {
     delim = strchr(buf, ' ');
   }
 argv[argc] = NULL;
 if (argc == 0) /* ignore blank line */
   return 1;
 /* should the job run in the background? */
 if ((bg = (*argv[argc - 1] == '&')) != 0) {
   argv[--argc] = NULL;
 }
 return bg;
}
* builtin cmd - If the user has typed a built-in command then execute
* it immediately.
*/
int builtin cmd(char **argv) {
 if (!strcmp(argv[0], "quit")) {
   exit(0);
 }
 if (!strcmp(argv[0],"jobs")) {
   listjobs(jobs);
   return 1;
 }
 if (!strcmp(argv[0], "bg") || !strcmp(argv[0], "fg")) {
   do_bgfg(argv);
   return 1;
 }
 return 0; /* not a builtin command */
```

```
* do_bgfg - Execute the builtin bg and fg commands
*/
void do_bgfg(char **argv) {
    int jid;
    struct job_t *job;
    if(argv[1]==NULL){
        printf("%s command requires PID or %%jobid argument\n",argv[0]);
        return;
    }
    if(argv[1][0]=='%'){
        jid=atoi(&argv[1][1]);
        job=getjobjid(jobs,jid);
        if(job==NULL){
            printf("%%%d: No such job\n",jid);
            return ;
        }
    } else if(isdigit(argv[1][0])){
        jid = atoi(argv[1]);
        job=getjobjid(jobs,jid);
       if(job==NULL){
            printf("(%d): No such process\n",jid);
            return;
        }
    } else {
        printf("%s: argument must be a PID or %%jobid\n",argv[0]);
        return;
    if (!strcmp(argv[0], "bg")) {
        job->state=BG;
        kill(-(job->pid),SIGCONT);
        printf("[%d] (%d) %s", jid, job->pid, job->cmdline);
    } else if (!strcmp(argv[0], "fg")) {
        if (job->state == ST) {
           //如果是挂起程序就重启并且转到前台,等待结束
            job->state=FG;
           kill(-(job->pid),SIGCONT);
           waitfg(job->pid);
        } else if (job->state == BG) {
            //如果是后台程序就转到前台并等待结束
            job->state=FG;
           waitfg(job->pid);
        }
    }
    return;
}
 * waitfg - Block until process pid is no longer the foreground process
void waitfg(pid t pid) {
```

```
sigset_t mask;
  sigemptyset(&mask);
 while (pid == fgpid(jobs))
    sigsuspend(&mask);
 return;
}
/***********
 * Signal handlers
 *************/
 * sigchld_handler - The kernel sends a SIGCHLD to the shell whenever
       a child job terminates (becomes a zombie), or stops because it
      received a SIGSTOP or SIGTSTP signal. The handler reaps all
       available zombie children, but doesn't wait for any other
      currently running children to terminate.
void sigchld_handler(int sig) {
 pid_t pid;
 int status;
 sigset_t mask_all, prev;
 sigfillset(&mask_all);
 while ((pid = waitpid(-1, &status, WNOHANG | WUNTRACED)) > 0) {
   if (WIFEXITED(status)) // 正常退出 delete
      sigprocmask(SIG_BLOCK, &mask_all, &prev);
      deletejob(jobs, pid);
      sigprocmask(SIG_SETMASK, &prev, NULL);
    } else if (WIFSIGNALED(status)) // 信号退出 delete
    {
      sigprocmask(SIG_BLOCK, &mask_all, &prev);
      printf("Job [%d] (%d) terminated by signal %d\n", pid2jid(pid), pid,
WTERMSIG(status));
      deletejob(jobs, pid);
      sigprocmask(SIG_SETMASK, &prev, NULL);
    } else if (WIFSTOPPED(status)) {
      sigprocmask(SIG BLOCK, &mask all, &prev);
      struct job_t *job=getjobpid(jobs,pid);
      job->state=ST;
      printf("Job [%d] (%d) stopped by signal
%d\n",pid2jid(pid),pid,WSTOPSIG(status));
      sigprocmask(SIG_SETMASK, &prev, NULL);
   }
 }
 return;
}
 * sigint_handler - The kernel sends a SIGINT to the shell whenver the
    user types ctrl-c at the keyboard. Catch it and send it along
     to the foreground job.
 */
void sigint handler(int sig) {
```

```
sigset_t mask_all, prev;
    sigfillset(&mask_all);
    sigprocmask(SIG_SETMASK, &mask_all, &prev);
    pid_t pid = fgpid(jobs);
    sigprocmask(SIG_SETMASK, &prev, NULL);
    if (pid > 0) {
        kill(-pid,sig);
    }
    return;
}
 * sigtstp_handler - The kernel sends a SIGTSTP to the shell whenever
      the user types ctrl-z at the keyboard. Catch it and suspend the
       foreground job by sending it a SIGTSTP.
 */
void sigtstp_handler(int sig) {
    sigset_t mask_all, prev;
    sigfillset(&mask_all);
    sigprocmask(SIG_SETMASK, &mask_all, &prev);
    pid_t pid = fgpid(jobs);
    sigprocmask(SIG_SETMASK, &prev, NULL);
    if (pid > 0) {
        kill(-pid,sig);
    }
    return;
}
/***********
 * End signal handlers
 *******************
/****************
 * Helper routines that manipulate the job list
/* clearjob - Clear the entries in a job struct */
void clearjob(struct job_t *job) {
 job \rightarrow pid = 0;
 job \rightarrow jid = 0;
 job->state = UNDEF;
  job \rightarrow cmdline[0] = ' (0';
}
/* initjobs - Initialize the job list */
void initjobs(struct job_t *jobs) {
 int i;
 for (i = 0; i < MAXJOBS; i++)
    clearjob(&jobs[i]);
}
/* maxjid - Returns largest allocated job ID */
int maxjid(struct job t *jobs) {
```

```
int i, max = 0;
 for (i = 0; i < MAXJOBS; i++)
    if (jobs[i].jid > max)
      max = jobs[i].jid;
 return max;
}
/* addjob - Add a job to the job list */
int addjob(struct job_t *jobs, pid_t pid, int state, char *cmdline) {
 int i;
 if (pid < 1)
   return 0;
  for (i = 0; i < MAXJOBS; i++) {
    if (jobs[i].pid == 0) {
      jobs[i].pid = pid;
      jobs[i].state = state;
      jobs[i].jid = nextjid++;
      if (nextjid > MAXJOBS)
        nextjid = 1;
      strcpy(jobs[i].cmdline, cmdline);
      if (verbose) {
       printf("Added job [%d] %d %s\n", jobs[i].jid, jobs[i].pid,
               jobs[i].cmdline);
      }
     return 1;
    }
  }
  printf("Tried to create too many jobs\n");
  return 0;
}
/* deletejob - Delete a job whose PID=pid from the job list */
int deletejob(struct job_t *jobs, pid_t pid) {
 int i;
 if (pid < 1)
    return 0;
 for (i = 0; i < MAXJOBS; i++) {
   if (jobs[i].pid == pid) {
      clearjob(&jobs[i]);
      nextjid = maxjid(jobs) + 1;
      return 1;
    }
  }
  return 0;
/* fgpid - Return PID of current foreground job, 0 if no such job */
pid_t fgpid(struct job_t *jobs) {
 int i;
```

```
for (i = 0; i < MAXJOBS; i++)
    if (jobs[i].state == FG)
      return jobs[i].pid;
  return 0;
}
/* getjobpid - Find a job (by PID) on the job list */
struct job_t *getjobpid(struct job_t *jobs, pid_t pid) {
 int i;
 if (pid < 1)
   return NULL;
 for (i = 0; i < MAXJOBS; i++)
   if (jobs[i].pid == pid)
      return &jobs[i];
 return NULL;
}
/* getjobjid - Find a job (by JID) on the job list */
struct job_t *getjobjid(struct job_t *jobs, int jid) {
 int i;
 if (jid < 1)
   return NULL;
 for (i = 0; i < MAXJOBS; i++)
   if (jobs[i].jid == jid)
      return &jobs[i];
 return NULL;
}
/* pid2jid - Map process ID to job ID */
int pid2jid(pid_t pid) {
 int i;
 if (pid < 1)
   return 0;
 for (i = 0; i < MAXJOBS; i++)
   if (jobs[i].pid == pid) {
      return jobs[i].jid;
    }
 return 0;
}
/* listjobs - Print the job list */
void listjobs(struct job_t *jobs) {
 int i;
  for (i = 0; i < MAXJOBS; i++) {
    if (jobs[i].pid != 0) {
      printf("[%d] (%d) ", jobs[i].jid, jobs[i].pid);
      switch (jobs[i].state) {
      case BG:
        printf("Running ");
```

```
break;
      case FG:
        printf("Foreground ");
       break;
      case ST:
        printf("Stopped ");
       break;
      default:
        printf("listjobs: Internal error: job[%d].state=%d ", i, jobs[i].state);
     printf("%s", jobs[i].cmdline);
   }
 }
}
 * end job list helper routines
 ****************************
/***************
 * Other helper routines
 ******************
 * usage - print a help message
*/
void usage(void) {
 printf("Usage: shell [-hvp]\n");
 printf(" -h print this message\n");
 printf(" -v print additional diagnostic information\n");
 printf(" -p do not emit a command prompt\n");
 exit(1);
}
* unix_error - unix-style error routine
void unix_error(char *msg) {
 fprintf(stdout, "%s: %s\n", msg, strerror(errno));
 exit(1);
}
 * app_error - application-style error routine
void app_error(char *msg) {
 fprintf(stdout, "%s\n", msg);
 exit(1);
}
* Signal - wrapper for the sigaction function
handler_t *Signal(int signum, handler_t *handler) {
  struct sigaction action, old action;
```

```
action.sa_handler = handler;
sigemptyset(&action.sa_mask); /* block sigs of type being handled */
action.sa_flags = SA_RESTART; /* restart syscalls if possible */

if (sigaction(signum, &action, &old_action) < 0)
    unix_error("Signal error");
    return (old_action.sa_handler);
}

/*
    * sigquit_handler - The driver program can gracefully terminate the
    * child shell by sending it a SIGQUIT signal.
    */
void sigquit_handler(int sig) {
    printf("Terminating after receipt of SIGQUIT signal\n");
    exit(1);
}</pre>
```