

# A STUDY ON EFFICIENT AUTHENTICATION METHOD FOR V2V COMMUNICATION

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A Study on Efficient Authentication Method for V2V Communication

# Content



## 1. VANET (VEHICULAR AD-HOC NETWORK)



DEVELOPMENT  
AND SECURITY  
PROBLEMS  
IN  
VANET

In the ever-changing Internet age, Internet are not limited to linking humans together. With the concept of the Internet of Things brought out, VANET, vehicular ad-hoc network, as a major branch of the Internet of Things is gradually emerging and has aroused the attention of scientific research institutions and automobile manufacturers.

Information security problems hidden in the VANET are concerned by people. The IEEE 1609.2 standard provides a reference for the vehicle manufactures in information security designation for VANET, but there are still some details in IEEE 1609.2 standard can be optimized.

## 2. CRYPTOLOGY BACKGROUND KNOWLEDGE



DIGITAL  
SIGNATURE

X.509 FORMAT  
DIGITAL  
CERTIFICATION

## 3. SYSTEM DESIGNATION

A STUDY ON EFFICIENT  
AUTHENTICATION  
METHOD FOR V2V  
COMMUNICATION



## *Using NS-3 to simulate communication system*

By utilizing NS-3 simulator to construct a ad-hoc network, importing OpenSSL to support cryptology algorithms, I perform simulations both on the improved system and the system applying original IEEE1609.2 standard.



SIMULATIONS

## 4. RESULTS AND ANALYSIS

# 1. VANET

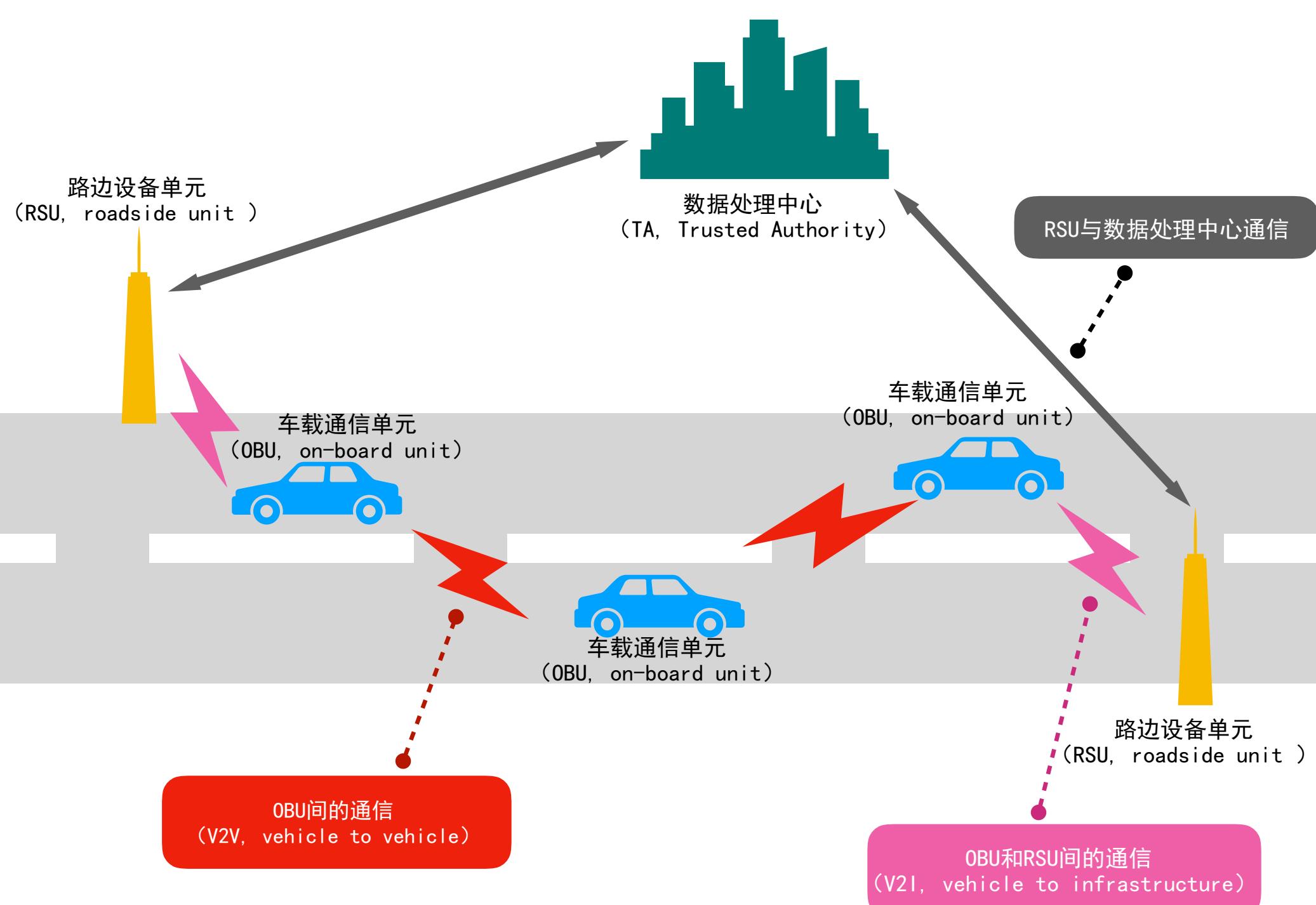
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A Study on Efficient Authentication Method for V2V Communication

# 1. VANET

## STRUCTURE OF VANET



VANET is consisted of 3 entities:

**OBU**, On-Board Unit

**RSU**, Roadside Unit

**TA**, Trusted Authority

Communications in VANET:

**V2V**, vehicle to vehicle communication

**V2I**, vehicle to infrastructure communication

communication between RSU and TA



# 1. VANET

## DEVELOPMENT AND SECURITY PROBLEMS IN VANET

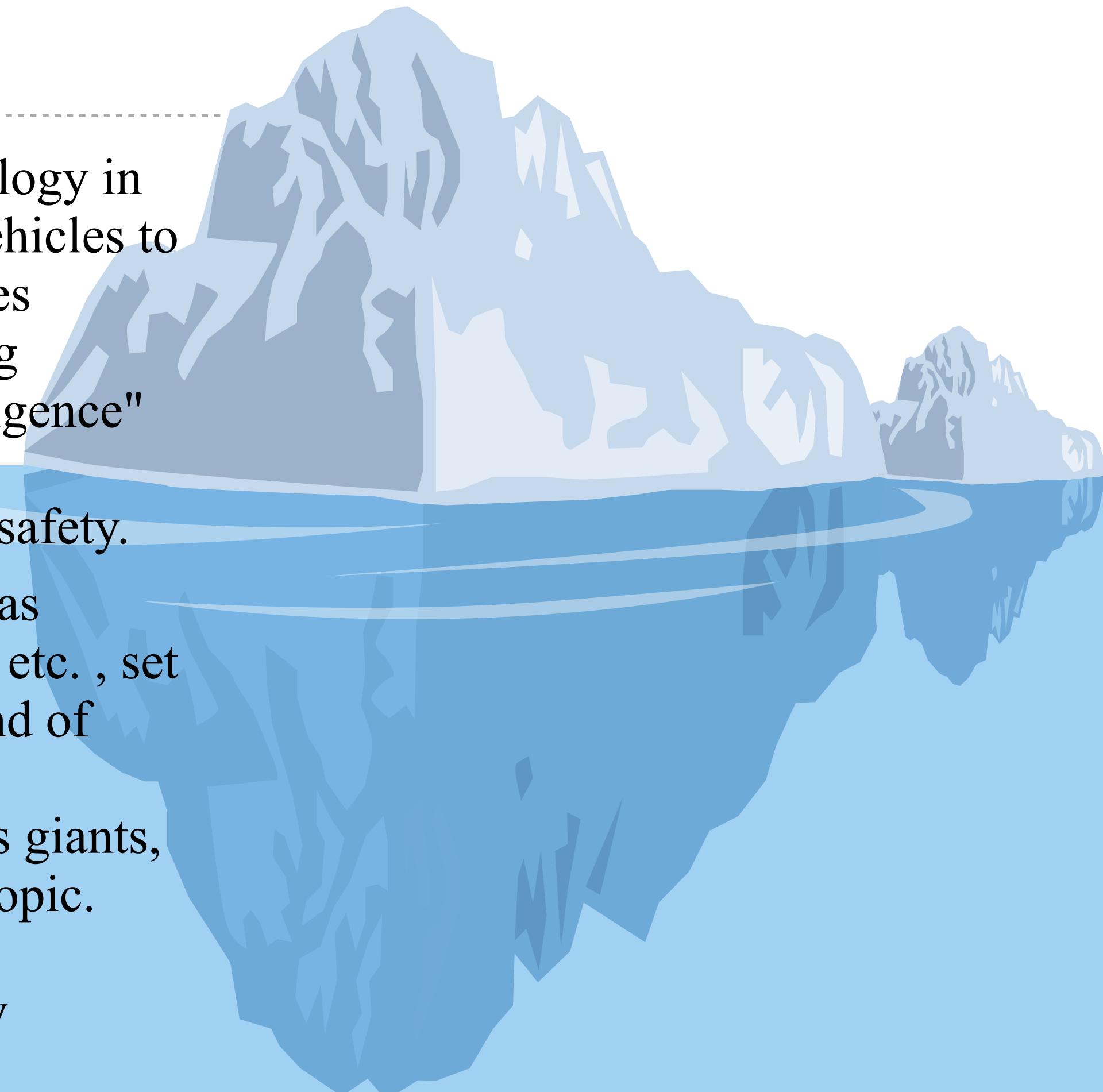
### DEVELOPMENT

VANET is the most promising technology in mobile Ad Hoc network. It enables vehicles to exchange information, and summarizes remote information by interconnecting vehicles, thereby enhances the "intelligence" of vehicles, and achieves the goal of improving road traffic efficiency and safety.

Numbers of famous companies, such as Google, Alibaba, Baidu, Tencent and etc., set up their VANET teams to join the trend of developing this frontier technique.

With the promotion of these industry's giants, VANET immediately becomes a hot topic.

Economic benefits of the business opportunities in VANET are gradually emerging.





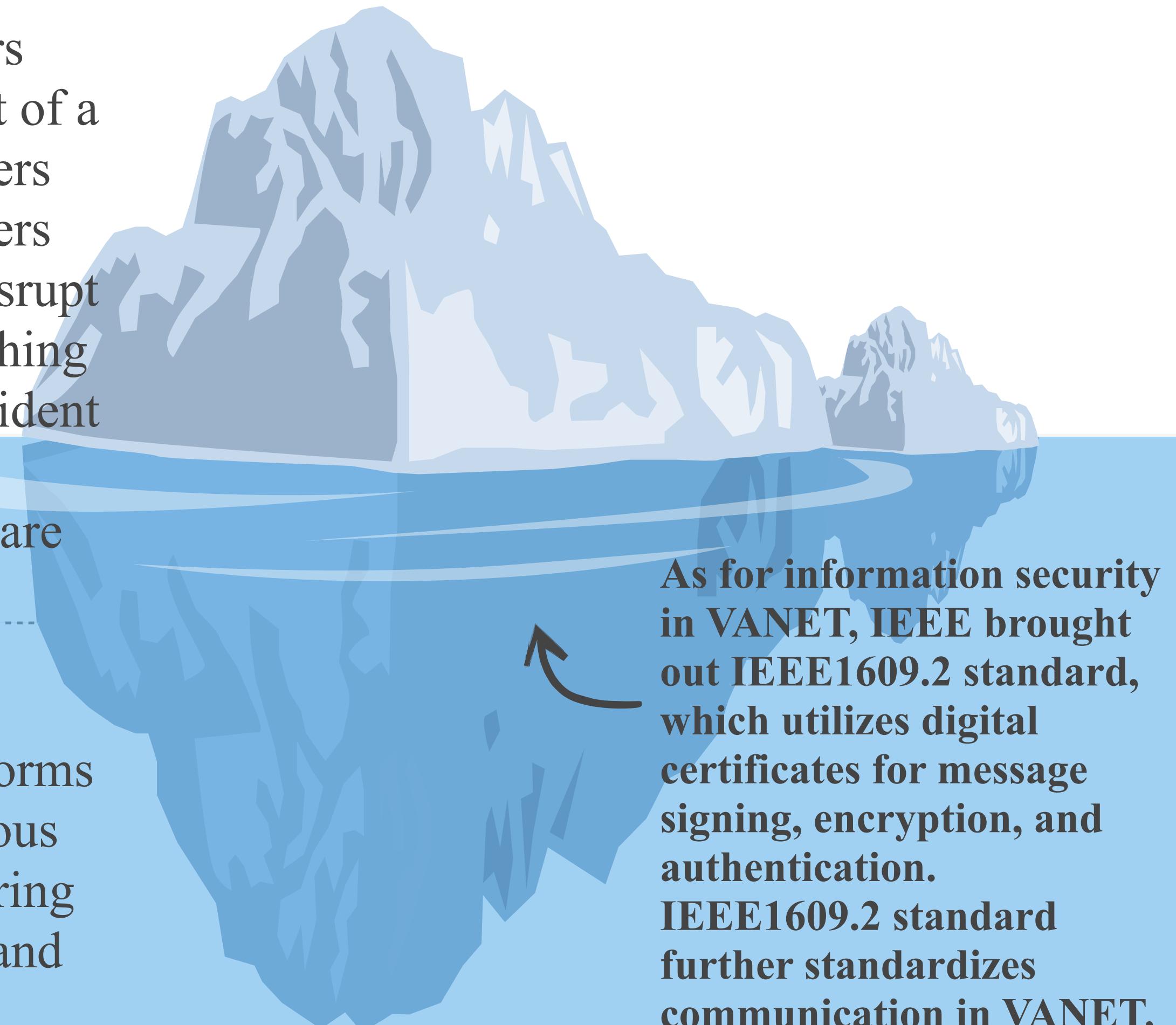
# 1. VANET

## DEVELOPMENT AND SECURITY PROBLEMS IN VANET

During the communication, attackers may modify the information content of a data package to mislead data receivers about the source of message; attackers may also create fake messages to disrupt the traffic. In the process of establishing VANET, the need for reproduce accident scene and the importance of tracing correct time sequence for accidents are also increasing.

## SECURITY ISSUES

Since the V2V communication performs in wireless network, false or erroneous traffic injection and message tampering may pose a threat to the credibility and security of the VANET.

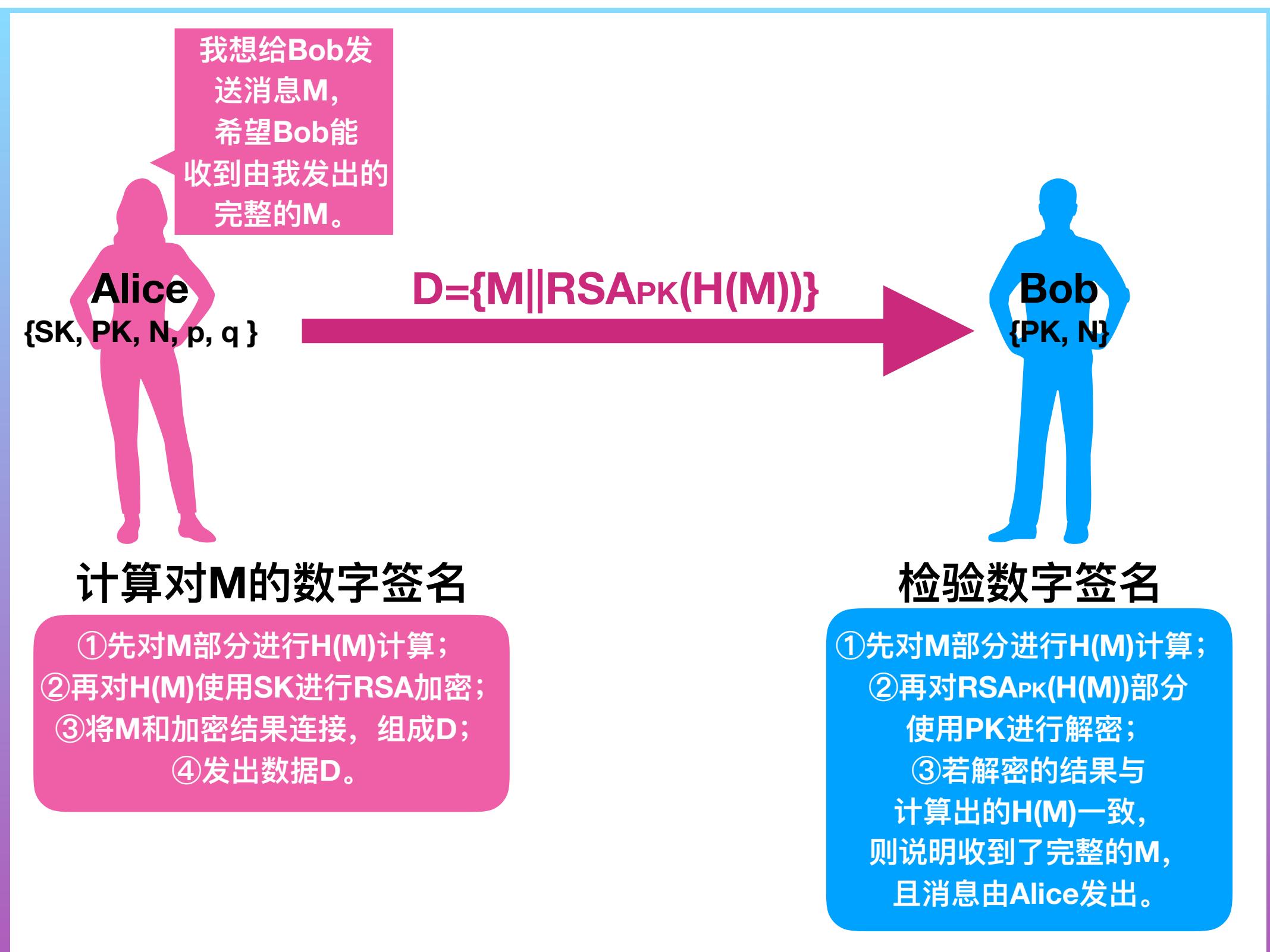


## 2. CRYPTOLOGY BACKGROUND KNOWLEDGE

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### DIGITAL SIGNATURE



### RSA

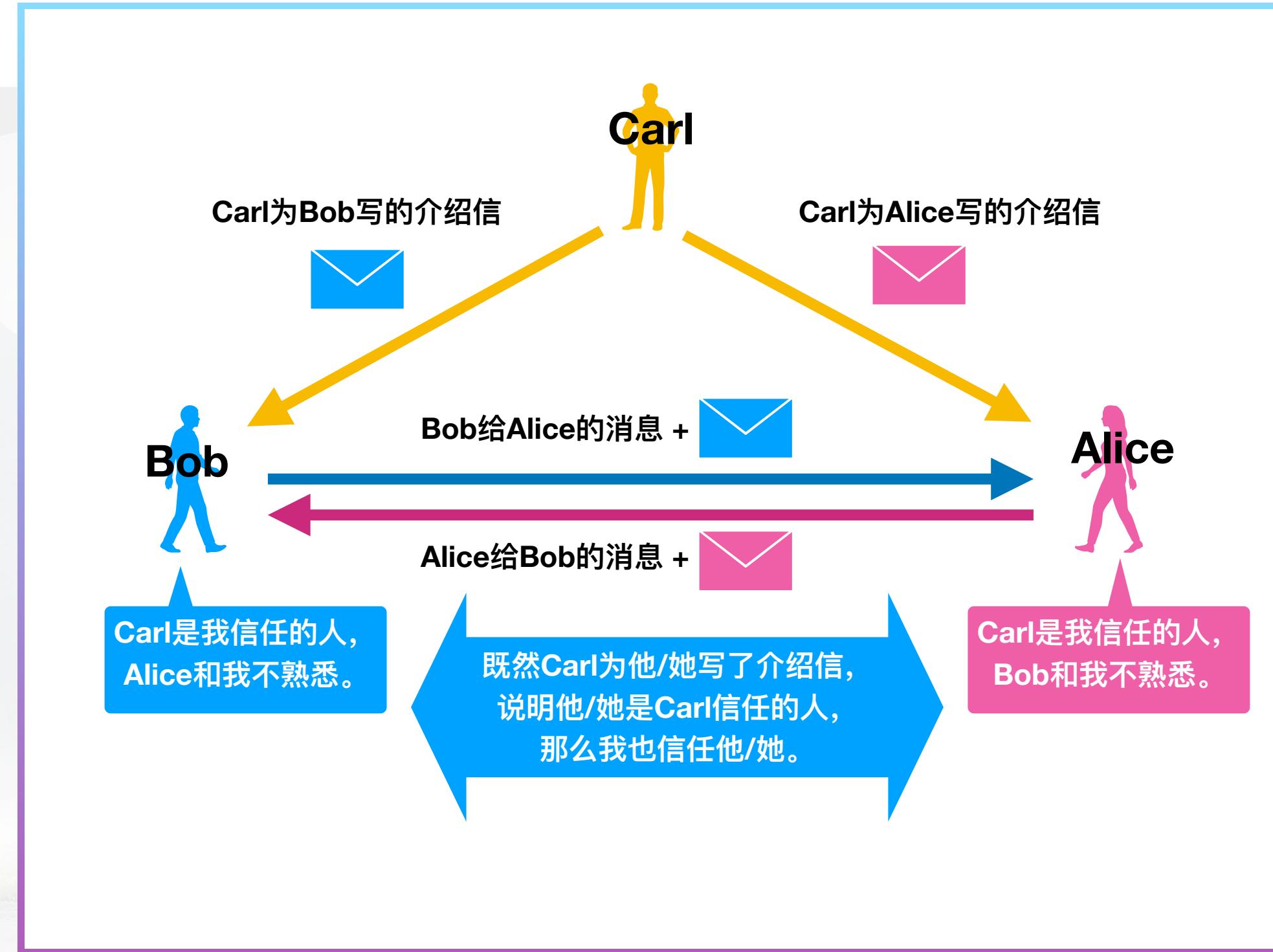
1. Randomly choose 2 distinct prime number  $p$  and  $q$ ;
2. Compute  $N$  and  $\varphi(N)$ .  $N = p \times q$ ,  $\varphi(N) = (p - 1) \times (q - 1)$ ,  $\varphi(N)$  is the Euler totient function of  $N$ ;
3. Generate a random number  $e$  as the private key, where  $1 < e < \varphi(N)$  and  $\gcd(e, \varphi(N)) = 1$ ;
4. Compute the public key  $d$ , where  $d \cdot e \equiv 1 \pmod{\varphi(N)}$ ;
5. Use  $M$  to present the original data, and use  $C$  to present the encrypted data. Relationship between  $M$  and  $C$  is:

$$C \equiv M^e \pmod{N}, \quad M \equiv C^d \pmod{N}.$$



## 2. CRYPTOLOGY BACKGROUND KNOWLEDGE

### X.509 FORMAT DIGITAL CERTIFICATION



消息密文  $\text{Enc}_{\text{pubb}}(M)$

签名  $\text{Sig}_{\text{priva}}(\text{H}(\text{Enc}_{\text{privb}}(M)))$

数字证书

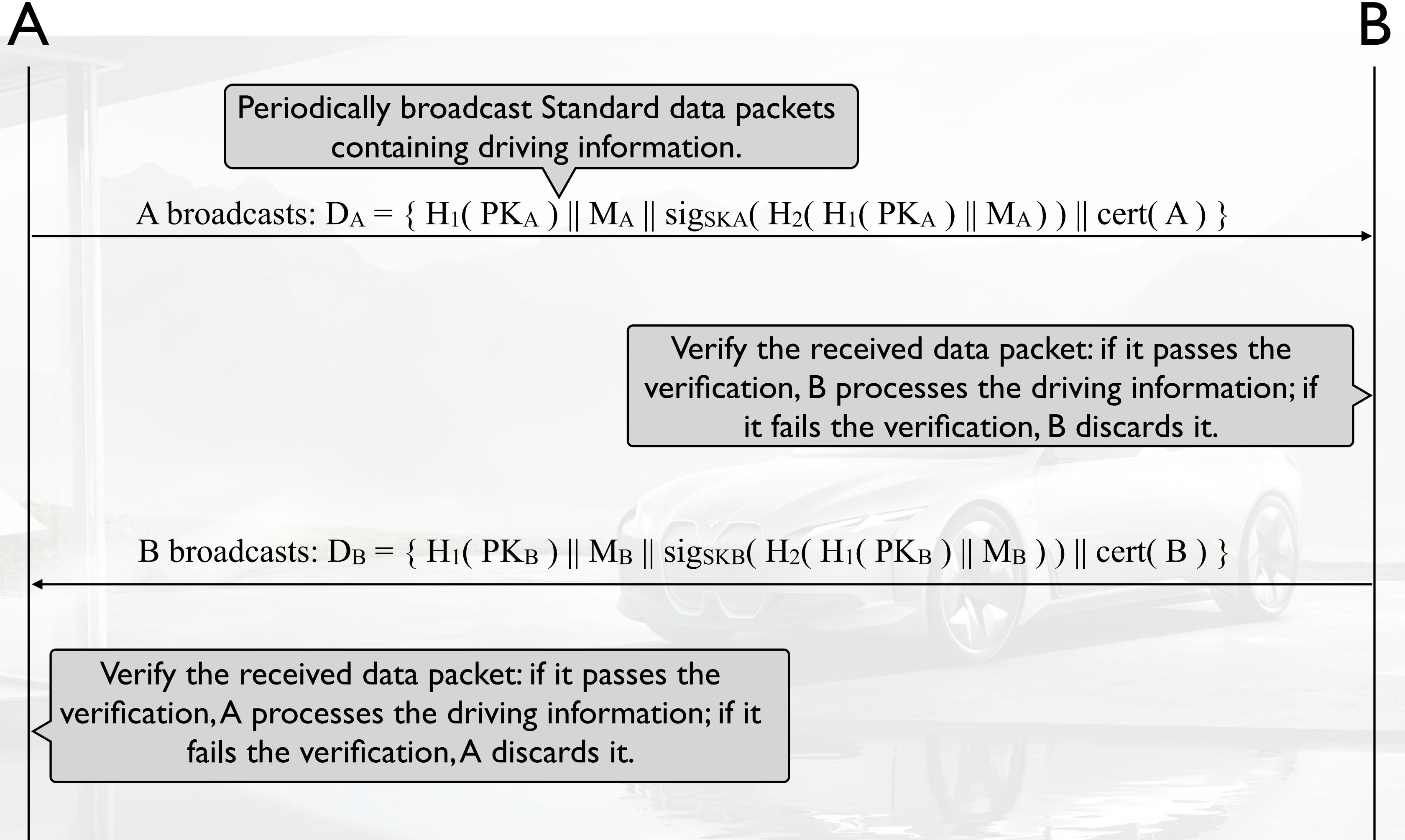
CA is trusted by entities on Internet, and the entity carries with digital certification signed by trusted CA is regarded as trusted by data receiver.

### 3. SYSTEM DESIGNATION

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### IEEE 1609.2 STANDARD



### 3. SYSTEM DESIGNATION

#### A STUDY ON EFFICIENT AUTHENTICATION METHOD FOR V2V COMMUNICATION

A

Periodically broadcast standard data packet(S packet) containing driving information.

B

A broadcasts S packet:  $D_{A1} = \{ H_1(PK_A) \parallel M_{A1} \parallel \text{sig}_{SKA}(H_2(H_1(PK_A) \parallel M_{A1})) \}$

B search its cache to find whether contains the public key of A. If B has, then verifies  $D_{A1}$ ; else B broadcasts the require/response data packet(R packet).

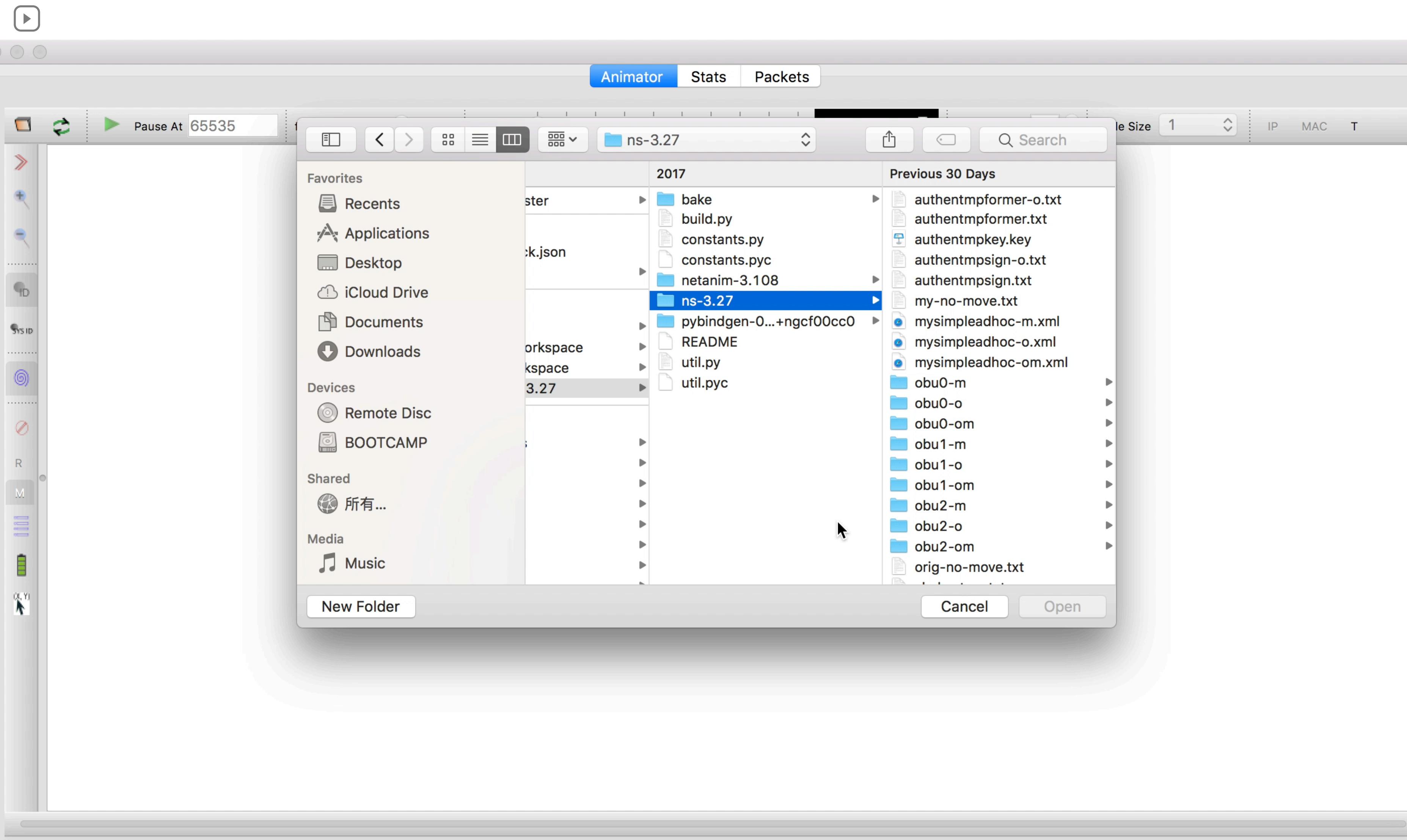
B broadcasts the R packet:

$D_{B1} = \{ H_1(PK_B) \parallel M_{B1} \parallel \text{sig}_{SKB}(H_2(H_1(PK_B) \parallel M_{B1})) \parallel \text{cert}(B) \}$

A verifies  $D_{B1}$ ; if it passes verification, then if A doesn't have B's public key, A saves it and broadcasts R packet. If A already has B's public key, then A directly processes B's driving data. If it doesn't pass the verification, then A discards it.

A broadcasts R data packet:  $D_{A2} = \{ H_1(PK_A) \parallel M_{A2} \parallel \text{sig}_{SKA}(H_2(H_1(PK_A) \parallel M_{A2})) \parallel \text{cert}(A) \}$

If A doesn't have B's public key information, then broadcasts R packet to response.



eclipse-c-workspace - /Users/apple/ns-allinone-3.27/ns-3.27/scratch/my-move-count.cc - Eclipse

个人收藏 隔空 桌面 iCloud 文稿 最近设备 远程 BOOT 可咩. 共享的 所有 标记 黄色 橙色 礼物 绿色 蓝色 所有

Project my-move-count.cc orig-move-count.cc mysimpleadhoc.cc mysimpleadhoc-orig.cc

```
219     lniw.close();
220     std::ifstream in1("obu2-m/cacert2.pem", std::ios::binary | std::ios::in);
221     in1.seekg(0, std::ios::beg);
222     in1.read ((char*)cer, 2312);
223     in1.close();
224     uint32_t cercount=0;
225     for(cercount=0;cercount<2312;cercount++){
226         buffer[392+2+cercount]=cer[cercount];
227     }
228 }
229 else {
230     socket->Close();
231     return;
232 }
233 uint8_t packet1[MTU]={0};
234 uint32_t i=0;
235 for(i=0;i<MTU;i++){
236     packet1[i]=buffer[i];
237 }
238 for(i=0;i<QLEN-MTU+2;i++){
239     packet2[i+2]=buffer[MTU+i];
240 }

241 Ptr<Packet> packets1=Create<Packet>((uint8_t const*)packet1,MTU);
242 Ptr<Packet> packets2=Create<Packet>((uint8_t const*)packet2,QLEN-MTU+4);

243 socket->Send (packets1);
244 PSend++;
245 NS_LOG_UNCOND (receivetime<<" Node "<<id<<" broadcasts one Q packet1!"<<" size = "<<packets1->GetSize());
246 NS_LOG_UNCOND (receivetime<<" PSend = "<<PSend<<" PRecv = "<<PRecv);
247
248 socket->Send (packets2);
```

Writable Smart Insert 200 : 20

## 4. RESULTS AND ANALYSIS

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## 4. RESULTS AND ANALYSIS

### TIME CONSUMING

S packets

| (发送方,接收方)   | 传输距离 (m) | 发送时间 ( $10^{-9}$ s) | 接收时间 ( $10^{-9}$ s) | 间隔时间 ( $10^{-9}$ s) | 平均时长 ( $10^{-9}$ s) |
|-------------|----------|---------------------|---------------------|---------------------|---------------------|
| (OBU0,OBU1) | 7.71     | 1000000000.0        | 1003848023.0        | 3848023.0           | 3848023.0           |
|             |          | 5000000000.0        | 5003848023.0        | 3848023.0           |                     |
|             |          | 9000000000.0        | 9003848023.0        | 3848023.0           |                     |
|             |          | 2000000000.0        | 2003848023.0        | 3848023.0           |                     |
|             |          | 6000000000.0        | 6003848023.0        | 3848023.0           |                     |
|             |          | 1000000000.0        | 11003848023.0       | 3848023.0           |                     |
| (OBU1,OBU0) | 5.0      | 2000000000.0        | 2003848016.0        | 3848016.0           | 3848016.0           |
|             |          | 6000000000.0        | 6003848016.0        | 3848016.0           |                     |
|             |          | 1000000000.0        | 10003848016.0       | 3848016.0           |                     |
|             |          | 3000000000.0        | 3003848016.0        | 3848016.0           |                     |
|             |          | 7000000000.0        | 7003848016.0        | 3848016.0           |                     |
|             |          | 11000000000.0       | 11003848016.0       | 3848016.0           |                     |
| (OBU2,OBU0) | 11.18    | 3000000000.0        | 3003848037.0        | 3848037.0           | 3848037.0           |
|             |          | 7000000000.0        | 7003848037.0        | 3848037.0           |                     |
|             |          | 11000000000.0       | 11003848037.0       | 3848037.0           |                     |
| (OBU0,OBU2) | 11.18    | 1000000000.0        | 1003848037.0        | 3848037.0           | 3848037.0           |
|             |          | 5000000000.0        | 5003848037.0        | 3848037.0           |                     |
|             |          | 9000000000.0        | 9003848037.0        | 3848037.0           |                     |

R packets

| (发送方,接收方)   | 传输距离  | 发送时间 ( $10^{-9}$ s) | 接收时间 ( $10^{-9}$ s) | 间隔时间 ( $10^{-9}$ s) | 平均时长 ( $10^{-9}$ s) |
|-------------|-------|---------------------|---------------------|---------------------|---------------------|
| (OBU0,OBU1) | 7.71  | 1000000000.0        | 1023510023.0        | 23510023.0          | 31140055.6          |
|             |       | 3023450023.0        | 3059804134.0        | 36354111.0          |                     |
|             |       | 8211050357.0        | 8234480380.0        | 23430023.0          |                     |
|             |       | 2000000000.0        | 2023450023.0        | 23450023.0          |                     |
|             |       | 2023510023.0        | 2070420108.0        | 46910085.0          |                     |
|             |       | 6190318272.0        | 6224134341.0        | 33186069.0          |                     |
| (OBU1,OBU2) | 5.0   | 2000000000.0        | 2023450016.0        | 23450016.0          | 31674665.0          |
|             |       | 2023510023.0        | 2070420101.0        | 46910078.0          |                     |
|             |       | 3059874085.0        | 3093930133.0        | 34056048.0          |                     |
|             |       | 2023510037.0        | 2059874085.0        | 36363748.0          |                     |
|             |       | 3000000000.0        | 3023410016.0        | 23410016.0          |                     |
|             |       | 4093930133.0        | 4119788217.0        | 25858084.0          |                     |
| (OBU2,OBU0) | 11.18 | 3000000000.0        | 3023410037.0        | 23410037.0          | 25570730.4          |
|             |       | 4093930133.0        | 4119788196.0        | 25858063.0          |                     |
|             |       | 6142938256.0        | 6166168293.0        | 23230037.0          |                     |
|             |       | 1000000000.0        | 1023510037.0        | 23510037.0          |                     |
|             |       | 3023450023.0        | 3059804148.0        | 36354125.0          |                     |
|             |       | 7224134341.0        | 7245196424.0        | 21062083.0          |                     |

## 4. RESULTS AND ANALYSIS

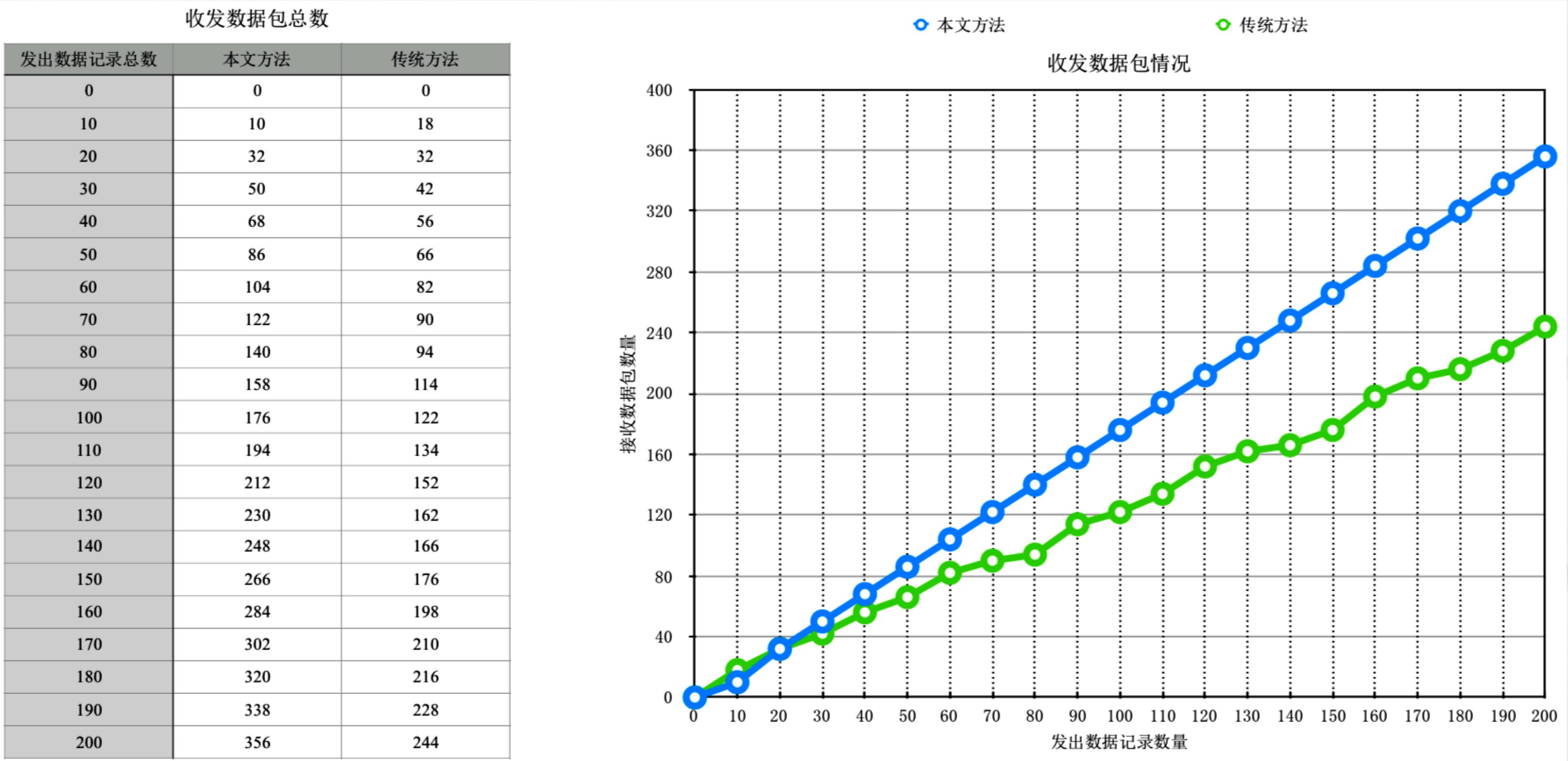
### TIME CONSUMING

#### Comprehensive Analysis

| 数据收发平均时长                                  | OBUs 与 OBU1 之间 | OBUs 与 OBU2 之间 | OBUs 与 OBU3 之间 |
|---|----------------|----------------|----------------|
| 标准数据包 ( $10^{-9}$ s)                      | 3848023.0      | 3848037.0      | 3848016.0      |
| 传统数据包 ( $10^{-9}$ s)                      | 31140055.6     | 25570730.4     | 31674665.0     |
| $\frac{\text{传统数据收发时长}}{\text{标准数据收发时长}}$ | 8.09           | 6.65           | 8.23           |

## 4. RESULTS AND ANALYSIS

### EFFICIENCY

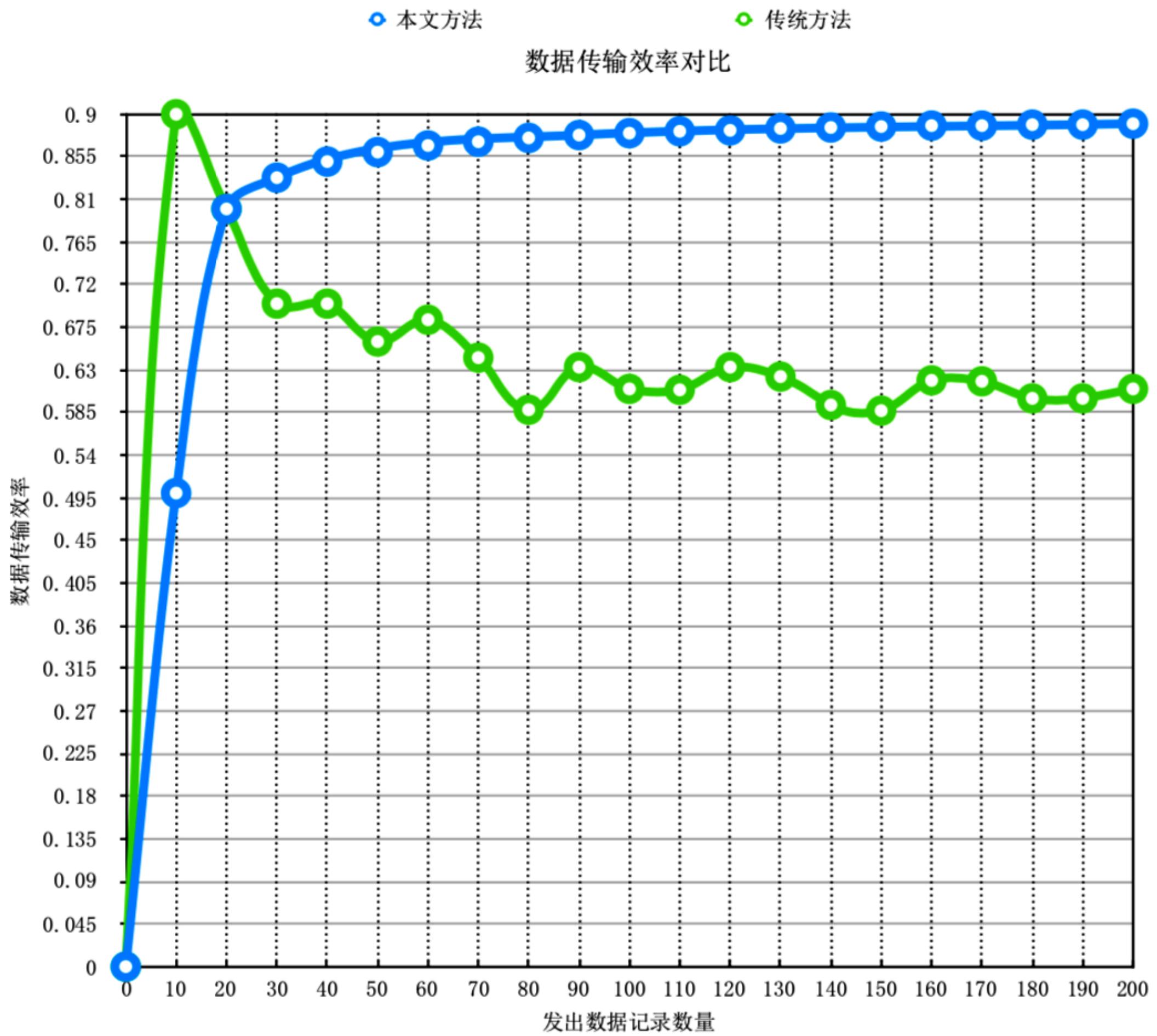




## 4. RESULTS AND ANALYSIS

### EFFICIENCY

| 发出数据记录总数 | 本文方法  | 传统方法  |
|----------|-------|-------|
| 0        | 0     | 0     |
| 10       | 0.5   | 0.9   |
| 20       | 0.8   | 0.8   |
| 30       | 0.833 | 0.7   |
| 40       | 0.85  | 0.7   |
| 50       | 0.86  | 0.66  |
| 60       | 0.867 | 0.683 |
| 70       | 0.871 | 0.643 |
| 80       | 0.875 | 0.588 |
| 90       | 0.878 | 0.633 |
| 100      | 0.88  | 0.61  |
| 110      | 0.882 | 0.609 |
| 120      | 0.883 | 0.633 |
| 130      | 0.885 | 0.623 |
| 140      | 0.886 | 0.593 |
| 150      | 0.887 | 0.587 |
| 160      | 0.888 | 0.619 |
| 170      | 0.888 | 0.618 |
| 180      | 0.889 | 0.6   |
| 190      | 0.889 | 0.6   |
| 200      | 0.89  | 0.61  |



## 4. RESULTS AND ANALYSIS

### CONCLUSION



**save more than  
80% of the  
transmission  
time**



**increase the  
efficiency of data  
transmission by  
45%**

It can be seen that the method proposed in this paper improves the communication speed and stability of communication while ensuring the data security of communication. The new method is more conducive to the transmission of important data in V2V. By using the method proposed in this paper, traffic data loss in the VANET will be alleviated, and V2V communication in the reliable VANET will bring greater convenience to people. It will also provide tremendous assistance to the traffic management department and the government in traffic management and accident recovery.

# Q&A

## A STUDY ON EFFICIENT AUTHENTICATION METHOD FOR V2V COMMUNICATION

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