

Chapter 3

DATA TRANSMISSION

TRANSMISSION MODES

Data Link



- **Data Communications:** exchange of digitally encoded information between 2 DTEs
 - (i) **Information:** user data
 - (ii) **Error control:** detect and respond to corrupted data
 - **Corruption from physical layer** ---> bit erroneously recovered
 - **Error detection** ---> probability of accepting undetected errors must be low
 - (iii) **Flow control:** detect and respond to lost data
 - Buffer overflow
 - Intermediate storage problems
 - ...

Data Link: link between 2 directly connected DTEs (no inter-network)

Data Transmission: techniques to achieve reliable transfer of information across bit-serial links between 2 DTEs

Data Link

- Data Transmission Techniques
 - Encode: representing characters by bits
 - Decode: extracting characters from bit representations
 - Codeword: bit representation of characters
- EBCDIC: 8 bit code (used with IBM peripherals)
- ASCII: 7 bit code
- Types of characters
 - Printable characters ('q', 'w', ...)
 - Control characters
 - Format control (LF, CR, DEL, ESC, ...)
 - Information separators (FS, RS, file & record separator)
 - Transmission control character (SOH, STX, ETX, ACK, NAK, SYN)
- Internal representation of data: bits, bytes, words

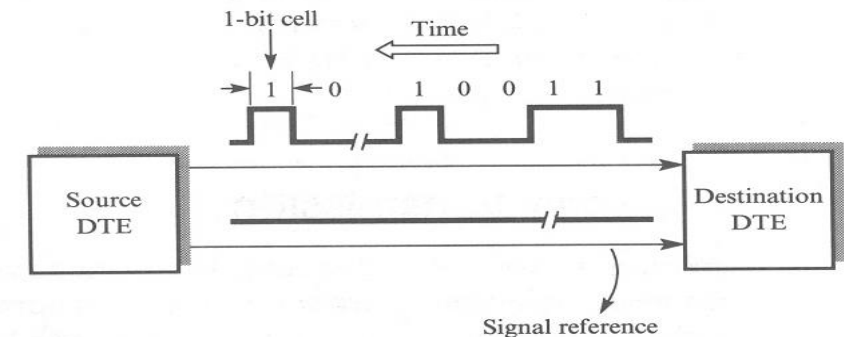
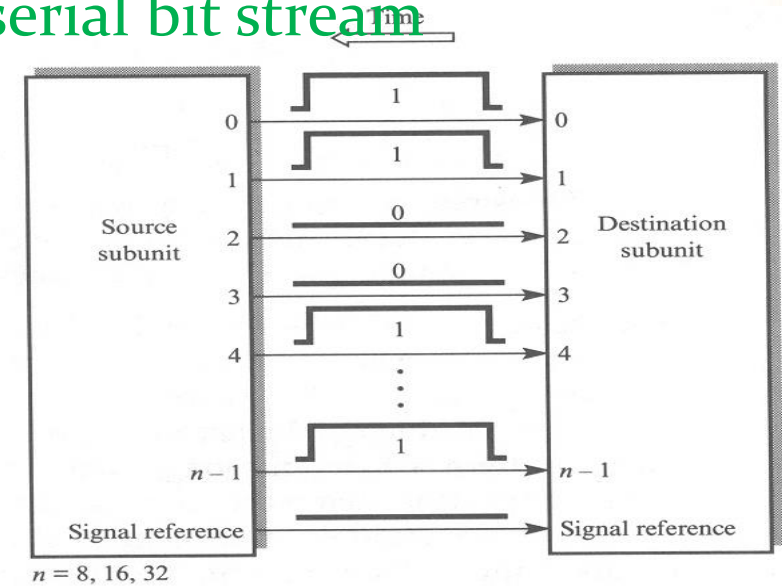
Data Transmission

- Bit serial transmission

- Inter-IC communications: usually **parallel**, short distance, low power
- Inter-device communications: remote devices use **serial bit stream**
 - Cabling, cost, signal power, noise & interference
- Bit representation:
 - Bipolar (serial)
 - Unipolar (parallel)

- Communications modes

- Simplex: 1-way (pager)
- Half duplex: 2 way alternate (push to talk)
- Full duplex: 2 way simultaneous



Transmission Modes

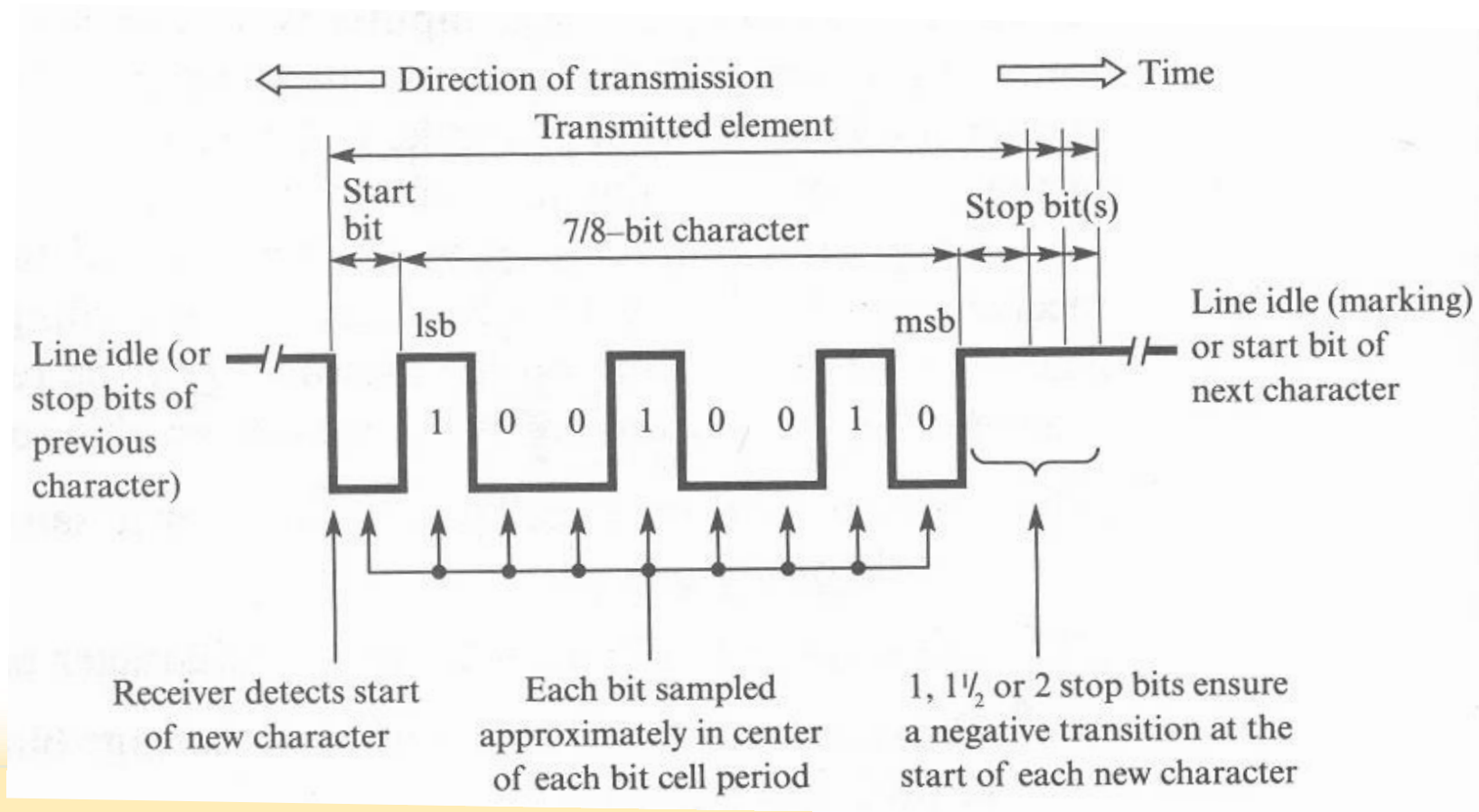
- Data often transmitted between 2 DTEs in 8 bit units
- Receiving DTE receives signal levels which vary according to bit pattern
- Type of **synchronization** depends on **clocking**
 - **Asynchronous transmission**: local independent clocks
 - **Synchronous transmission**: global clock or synchronized local clocks
- **Synchronization** tasks for receiver to interpret signal
 1. **Bit synchronization**
 - Start of each bit cell period (clock) (to sample at mid-point)
 2. **Byte synchronization**
 - Start & end of each word (byte)
 3. **Frame synchronization**
 - Start & end of each block (frame)

Asynchronous Transmission

- Applications
 - Most useful for data with irregular arrival times (human actions)
 - Transmission line with long idle states (marking)
- Examples:
 - Terminal & keyboard I/O
 - Block character transfers between 2 computers
- Characterized by:
 - No direct clock information between receiver and transmitter
 - Low bit rates, long idle times
 - Relatively coarse and inefficient method
 - Coarse bit synchronization

Asynchronous Transmission

- Each **byte** treated **independently** for clock synchronization purposes
- Receiver synchronizes at the **start** of each byte
- Uses **start bit** & **stop bit** for each byte

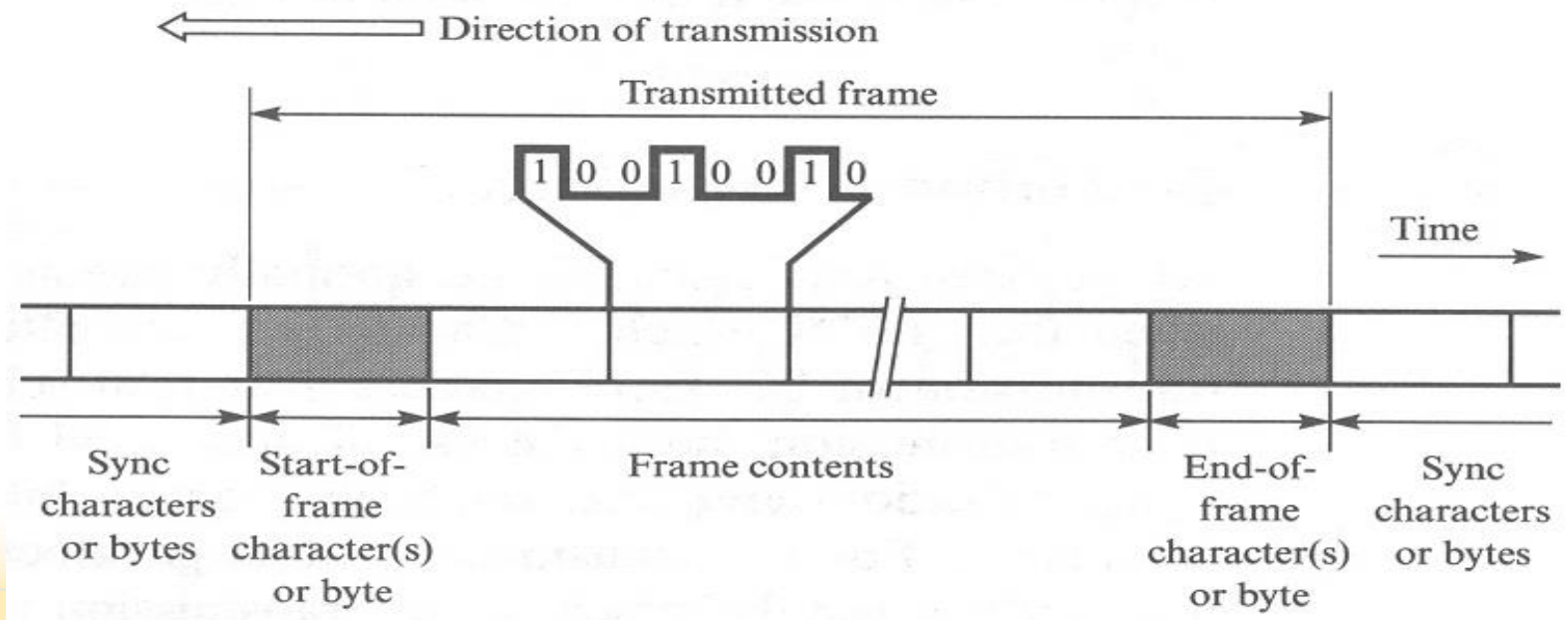


Asynchronous Transmission

- Bit synchronization
 - Stop & start bits have opposite polarity
 - At least 1 transition between each byte
 - Start bit: 1st 1 -> 0 transition after idle period
 - 1, 1.5 or 2 stop bits used to ensure end of byte
- Byte synchronization
 - Bytes are encapsulated in start bit & stop bit
 - Receiver has to re-synchronize at each byte
- Frame synchronization
 - Blocks are encapsulated by transmission control characters, SOF & EOF (start & end of frame)
 - Ensures receiver can distinguish new frame or end of frame after idle period

Synchronous Transmission

- Higher bit rates
- Unit of transmission = frame
- Continuous transmissions (frame transmitted as a contiguous bit stream)
- Receiver synchronizes each bit for duration of frame
- Contains overhead control bits and payload data



Synchronous Transmission

- Receiver synchronization
 - Bit synchronization
 - Bits are **explicitly** encoded in signal
 - Byte synchronization
 - **Reserved** bytes precedes each frame
 - Ensure receiver interprets bit stream on **correct byte** boundary
 - Frame synchronization
 - Frame is encapsulated between **reserved bytes**
 - Unique **SOF** & **EOF** identifiers to guarantee frame doesn't contain embedded **SOF** and **EOF**
 - **Inter-frame synchronization** (with inter-frame idle time)
 1. Each frame is preceded by **synch bytes**, so receiver can **regain** synchronization
 2. Synchronous **idle** bytes continuously transmitted to allow receiver to **retain** bit & byte synchronization

Ex. of Asynchronous vs. Synchronous Transmission

- How much **overhead** with **100** byte message?
 - **Asynchronous** transmission with **1** start bit, **2** stop bits
 - 100 data bytes + SOF byte + EOF byte
 - **Efficiency** = $\frac{100 \times 8}{(100+2) \times (8+1+2)} = \frac{800}{1122} \cong 71.30 \%$
 - **Overhead** = $1122 - 800 = 322$ bits
 - **Synchronous** transmission
 - 100 data bytes + 2 synch characters + SOF byte + EOF byte
 - **Efficiency** = $\frac{100 \times 8}{(100+2+2) \times 8} = \frac{800}{832} \cong 96.15 \%$
 - **Overhead** = $832 - 800 = 32$ bits

Error Control

- Detect transmission **errors** (bit corruptions)
- Manage/Respond to **detected** errors
- **Bit corruption**: Signal level representing a bit is altered ---> incorrectly interpreted by receiver
- **Error detection**: depends on transmission techniques
 - **Asynchronous transmission**: detect errors in **bytes**
 - **Parity bit** attached to each byte
 - Each **byte** treated as **separate** entity
 - **Synchronous transmission**: detect errors over entire **frame**
 - Frame is basic **unit** of transmission
 - **Probability** of bit errors increases with frame **size**
 - **Check sequence** attached to frames ---> much more effective than **parity** bits
 - **Transmitter**: compute **check bits**, append to frame data, then send frame
 - **Receiver**: recompute **check bits**, compare with transmitted check bits

Flow Control

- How to manage **lost** data
- **Received data**
 - Must be temporarily stored (**input buffer**) at destination
 - On a network, may be intermediate storage on **switches**, **routers**, ...
- **Overflow**
 - Amount of received data **exceeds** available storage space ---> **data is lost**
 - If DTEs operate at **different** rates, faster device can **overflow** slower device
- **Lost data**: data that is sent but **never** received by intended recipient
 - Can be from **overflow**
 - In a network, **corruption** of **control** information can result in data being lost

Data Link Protocols

- **Protocol:**
 - Set of rules adhered to by all communicating entities to ensure successful exchange of information across serial data link
- **Components of data link protocol**
 - **Flow** control
 - **Error** control
 - Data **format**: bits/element, signal encoding
 - **Type** & **order** of message transmissions for reliable information transfer
- Example: connection set-up messages that ensure both DTEs are ready

