Algorithm Analysis

Name: Maiqi Hou, Course: CSI 3344

1.

Description

Creating a priority queue

Data Structure

A array and queue

Algorithm

For i = 0 to size (Set a size arbitrarily according to the ASCII code)

If array[ASCII Code] not equal zero

Push into the priority queue

Analysis

Input N	Array's size

Basic Operation	`void priorityQ(int array[]) {`
	`for (int i = 0; i < 128; i++) { // n = 128`
	`if (array[i] != 0) { // n`
	`Node* n = new Node(i, array[i]); // n`
	`pq.push(n); // n`
	`}`
	`}`
	`}`

Summation or Recurrence Relation	T(n) = O(n)

Worst Case Analysis

T(n) = O(n)

Best Case Analysis

T(n) = O(n)

2.

Description

Concerte a Huffman Tree

Data Structure

Priority queue

Algorithm

Priority queue q

While q's size remians 1

Declare Node left = q's top, then pop.

Delcare Node right = q's top, then pop.

Declare new Node = (left node' weight plus right node's weight)

New node's right and left equal to Node left and Node right

Push new node into q

Analysis

Input N	Priority queue's size

Basic Operation	`while (q.size() != 1) {`
	`Node* left = q.top();`
	`q.pop();`
	`Node* right = q.top();`
	`q.pop();`
	`Node* node = new Node(128, left->weight + right->weight);`
	`node->left = left;`
	`node->right = right;`
	`q.push(node);`
	`}`

Summation or Recurrence Relation	T(n) = O(log n)

Worst Case Analysis

T(n) = O(log n)

Best Case Analysis

T(n) = O(log n)

3.

Description

Record the code of the character

Data Structure

Map and array

Algorithm

If root's left exists

array [index]= 0

Recurrsive root's left, index +1

If root's right exists

Array = 1

Recurrsive root's right,index +1

If root's right and left not exists

Map [root's character] = array

Analysis

Input N	Number of node = n

Basic Operation	`void huffcode(Node* t, map<int, string>& m, int index, char a[])`
	`{`
	`if (t->left) {`
	`a[index] = '0';`
	`huffcode(t->left, m, index+1,a);`
	`}`
	`if (t->right) {`
	`a[index] = '1';`
	`huffcode(t->right, m, index+1,a);`

	}
	if (!t->left && !t->right) {
	string chara;
	for (int i = 0; i < index; i++) {
	chara += a[i];
	}
	m[t->ch] = chara;
	}
	}

Summation or Recurrence Relation	T(n) = O(n^2)

Worst Case Analysis

T(n) = O(n^2)

Best Case Analysis

T(n) = O(n)

4.

Description

Pass the map to main, storing each character's code

Data Structure

map

Algorithm

None

Analysis

Input N	None

Basic Operation	void Encode(map<int, string>& table) {
	table = m;
	}

Summation or Recurrence Relation	T(n) = 1;

Worst Case Analysis

T(n) = 1

Best Case Analysis

T(n) = 1

5.

Description

Read compressed file, read the character's counts

Data Structure

array

Algorithm

For i = 0 to n = 128(Seting size according to ASCII code)

Analysis

Input N	N = Seting size according to ASCII code

Basic Operation	`void Store(int array[], string s, ofstream& a)`
	`for (int i = 0; i < 128; i++) {`
	`if (array[i] != 0) {`
	`a << i << " " << array[i] << endl;`

Summation or Recurrence Relation	T(n) = O(n)

Worst Case Analysis

T(n) = O(n)

Best Case Analysis

T(n) = O(n)

6.

Description

Recreate huffman tree

Data Structure

priority queue

Algorithm

While file not terminate

File read int type character and number of character

Push into a priority queue

Analysis

Input N	Number of file's lines

Basic Operation	`void ReCreateTree(string filen) {`
	`ifstream f;`
	`f.open(filen);`
	`int index = 0, n = 0, ti = 0;`
	`while (!f.eof()) {`
	`f >> index >> n;`
	`Node* node = new Node(index, n);`
	`pq.push(node);`
	`}`
	`ti = index;`
	`}`
	`}`

Summation or Recurrence Relation	T(n) = O(log n)

Worst Case Analysis

T(n) = O(log n)

Best Case Analysis

T(n) = O(log n)

7.

Description

Codes decode each character

Data Structure

None

Algorithm

For i = 0 to string s' size

If s[i] = '0'

Node point to node's left

Else

Node point to node's right

If node's right and left not exists

Return character

Analysis

Input N	String's size = n

Basic Operation	void Decode(string s, string& c) {
	for (int i = 0; i < s.size(); i++) {
	if (s[i] == '0') {
	temp = temp->left;
	}
	else if (s[i] == '1'){
	temp = temp->right;
	}
	if (!temp->left && !temp->right) {
	c += char(temp->ch);
	temp = pq.top();
	}
	}

	}
Summation or Recurrence Relation	T(n) = n

Worst Case Analysis

T(n) = O(n)

Best Case Analysis

T(n) = O(n)