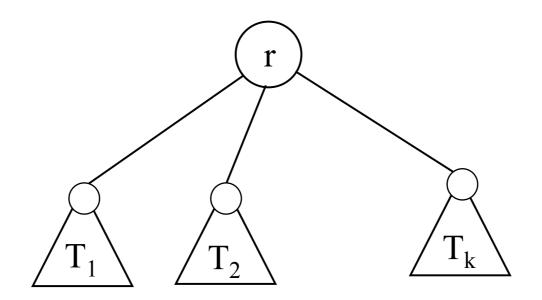
# Data Structure: Tree

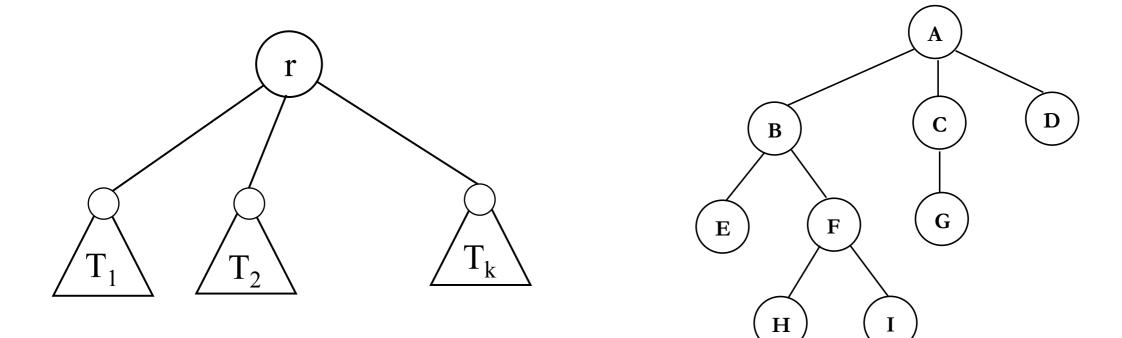
- a collection of nodes connected by edges without cycle
- by recursive definition:
  - an empty or
  - **a** root r and subtrees  $T_1, T_2, ..., T_k$  (disjoint sets) each of whose roots are connected to r by an edge



recursive definition of tree

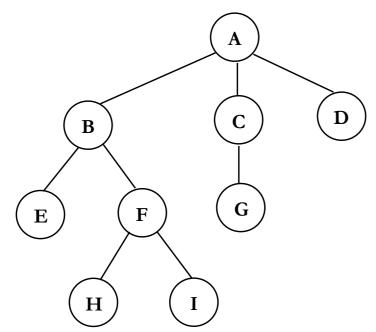
#### tree

- Each root of  $T_1, T_2, ..., T_k$  is a child of r, and r is the parent of each root.
- The roots of the subtrees are siblings of one another
- If there is an order among the  $T_i$ 's, the tree is an ordered tree.
- The degree of a node is the number of children it has.
- The degree of a tree is the maximum degree of the nodes.
- A leaf is a node of degree 0.



#### tree

- path between two nodes is a sequence of nodes  $n_1, n_2,... n_k$ , such that  $n_i$  is a parent of  $n_{i+1}$
- length of a path is the number of edges on the path (the path  $n_1$ ,  $n_2$ ,...  $n_k$ : length k-I)
- depth (level) of a node is the length of the (unique) path from the root to that node (root: level 0)
- height of a node is the length of the longest path from that node to a leaf (leaf: height 0)
- the height of a tree is the height of the root



#### representation of tree

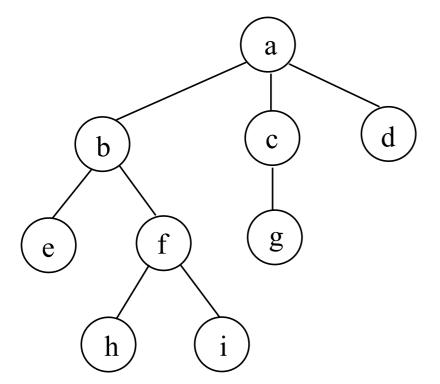
- $\blacksquare$  for any node x, there exists exactly one path from the root to x?
- tree can be empty with no node?
- how many edges are in a tree with n nodes?

## representation of tree

■ how can we implement a tree?

## representation of tree

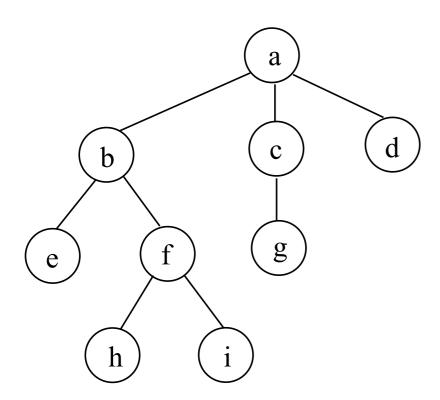
- how can we implement a tree?
  - linked list?
  - can we have pointers for the children nodes?
  - can we have fixed number of pointers to represent a tree?
    - for a tree of fixed number of degree?
    - else?



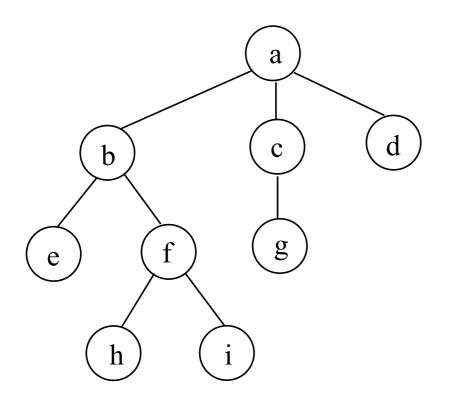
| data | link 1 | link 2 |  | link n |
|------|--------|--------|--|--------|
|------|--------|--------|--|--------|

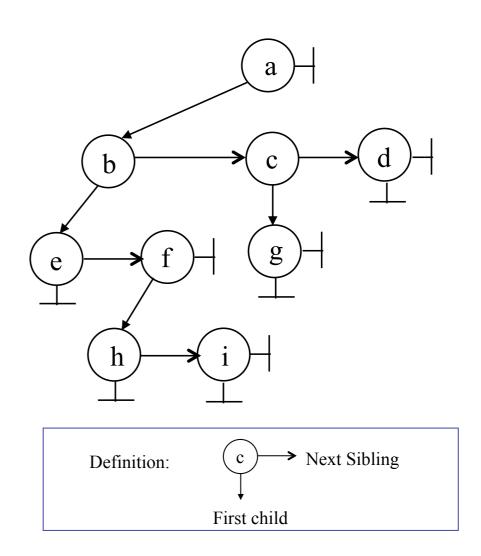
#### left child-right sibling representation

every node has at most one leftmost child and at most one closet right sibling

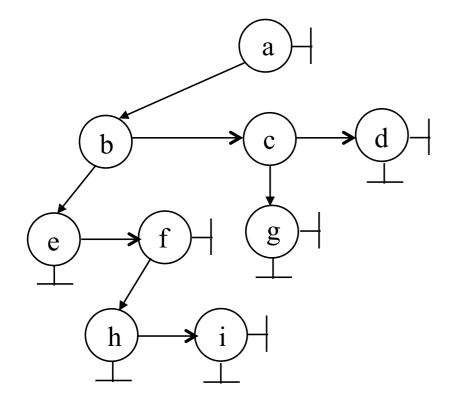


every node has at most one leftmost child and at most one closet right sibling





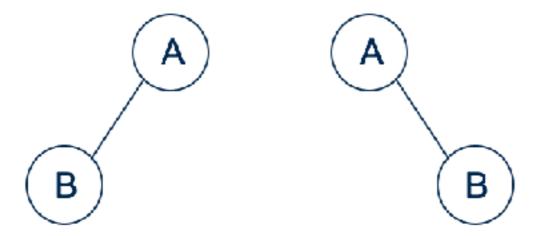
#### left child- right sibling representation



```
Definition: C Next Sibling
First child
```

```
struct TreeNode{
    ElementType Element;
    PtrToNode FirstChild;
    PtrToNode NextSibling;
    };
typedef struct TreeNode *PtrToNode;
```

- a finite set of nodes that is either
  - i) empty or
  - ii) a root node and two disjoint binary trees
- the tree on the left and the tree on the right are different



■ the maximum number of nodes on level i of a binary tree is 2<sup>i</sup>, i>=0

#### the proof by induction

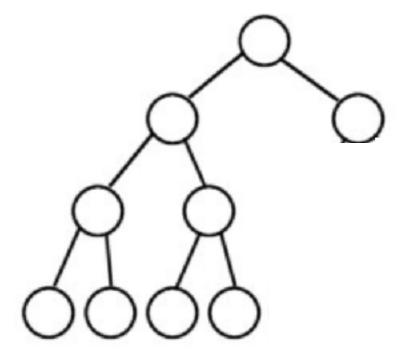
- base: for the root at level  $i=0, 2^0 = 1$
- induction hypothesis: assume that the maximum number of nodes on level i-1 > 0,  $2^{i-1}$
- induction step: on level i, 2\* (the maximum number of nodes on level i-1) =  $2*2^{i-1} = 2^{i}$
- the maximum number of nodes in a binary tree of depth k is  $2^{k+1}$ -1,  $k \ge 0$

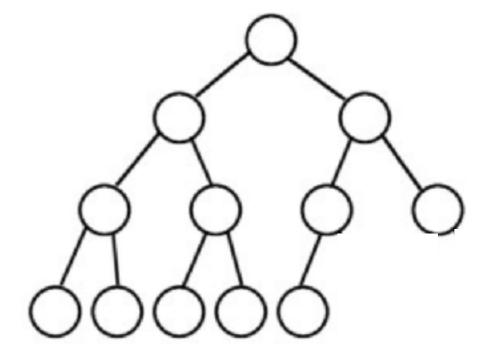
■ For any nonempty binary tree T, if  $n_0$  is the number of leaf nodes, and  $n_2$  is the number of nodes of degree 2, then  $n_0 = n_2 + 1$ 

 $n = n_0 + n_1 + n_2$ ,  $n_i$  is the number of nodes with i degree  $n_i$  is the number of nodes in the tree

 $n = B + I = n_1 + 2n_2 + I$ , B is the number of branches (edge)

- full binary tree is a binary tree in which every node has 0 or 2 children
- complete binary tree is a binary tree in which every level, except the last, is completely filled and the last level has all its nodes to the left side



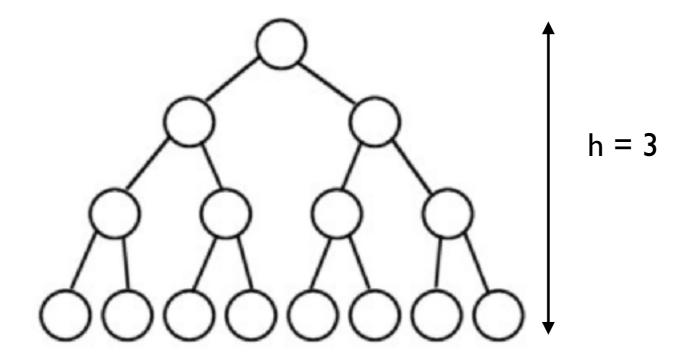


full binary tree

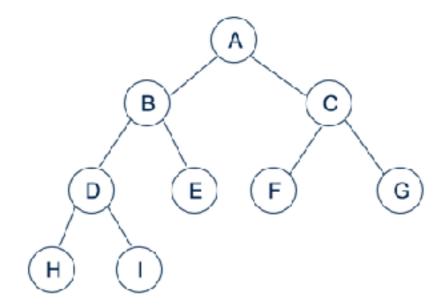
complete binary tree

- perfect binary tree of height h is a binary tree of height h having  $2^{h+1}$  1 nodes, (h >=0)
- the max number of nodes in the complete binary tree (height h) is 2 h+1 -1

$$2^{0} + 2^{1} + ... + 2^{h} = (2^{h+1} - 1)/(2 - 1) = 2^{h+1} - 1$$



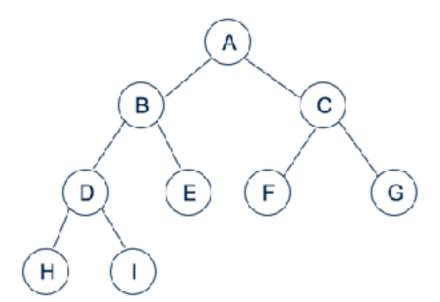
## binary tree: array representation



| [1] | Α |
|-----|---|
| [2] | В |
| [3] | С |
| [4] | D |
| [5] | Е |
| [6] | F |
| [7] | G |
| [8] | Н |
| [9] | I |

## binary tree: array representation

- if a complete binary tree with n nodes (i is the index) is represented sequentially,
  - leftChild(i) is at 2i for 2i <=n
  - rightChild(i) is at 2i + I for 2i + I <= n</pre>
  - parent(i) is at  $\lfloor i/2 \rfloor$  for i > 1



| [1] | Α   |
|-----|-----|
| [2] | В   |
| [3] | С   |
| [4] | D   |
| [5] | Е   |
| [6] | F   |
| [7] | G   |
| [8] | Н   |
| [9] | - 1 |

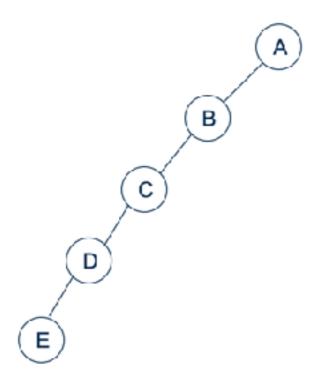
## binary tree: array representation

• if a complete binary tree with n nodes (i is the index) is represented sequentially,

leftChild(i) is at 2i for 2i <=n

rightChild(i) is at 2i + I for 2i + I <= n</pre>

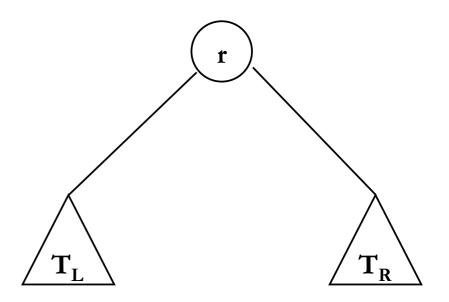
parent(i) is at Li/2\_ for i > I



| [1] | Α |  |
|-----|---|--|
| [2] | В |  |
| [3] | - |  |
| [4] | С |  |
| [5] | - |  |
| [6] | - |  |
| [7] | - |  |
| [8] | D |  |
| [9] | - |  |
|     |   |  |
| •   |   |  |
| -   |   |  |
| 16] | E |  |

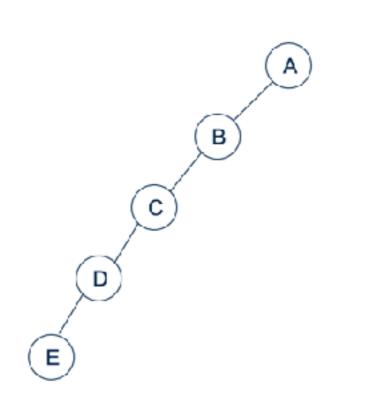
## binary tree: linked list representation

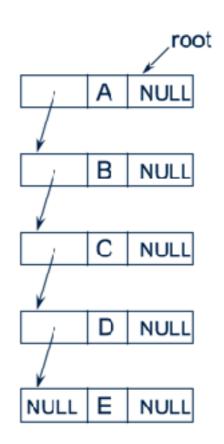
a tree in which each node has no more than 2 children (left subtree and right subtree)

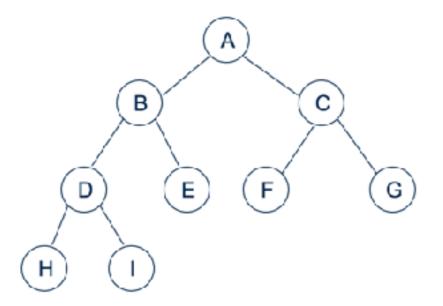


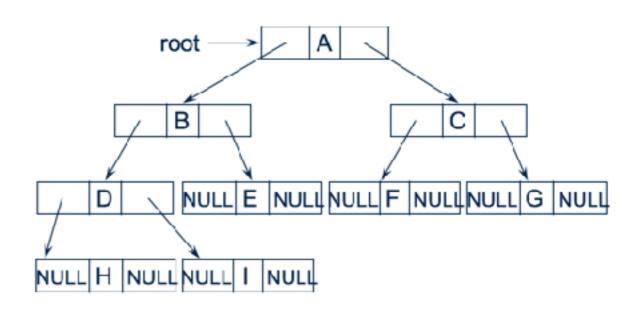
| Left | Element | Right |
|------|---------|-------|
|------|---------|-------|

## binary tree: linked list representation



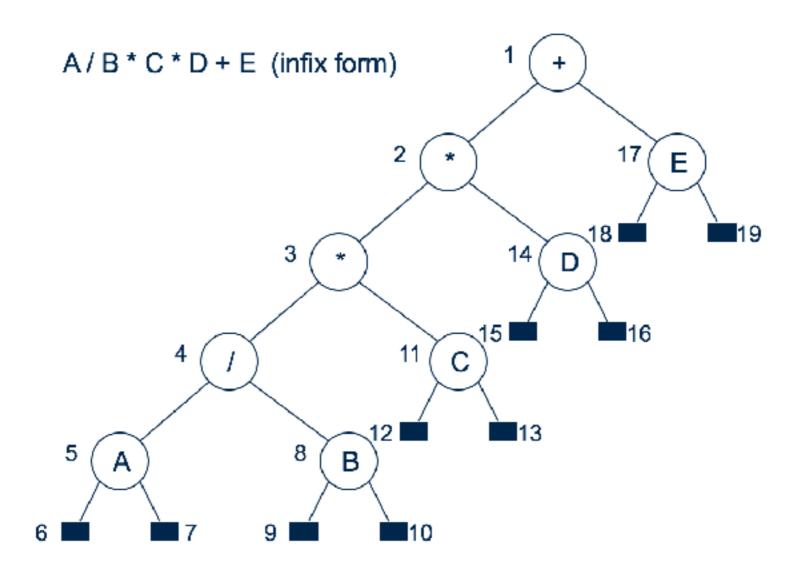






## application of binary tree

Expression Tree: intermediate representation for expressions used by the compiler



inorder traversal

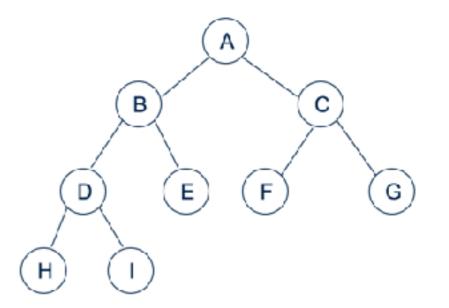
```
void inorder(Tree ptr) {
   if(ptr) {
       inorder(ptr->left_child);
       printf("%d", ptr->data);
       inorder(ptr->right_child);
```

| call of | : | value in     | action | call of | value in     | action |
|---------|---|--------------|--------|---------|--------------|--------|
| inorde  | r | root         |        | inorder | root         |        |
|         | 1 | +            |        | 11      | С            |        |
| 2       | 2 | *            |        | 12      | NULL         |        |
| 3       | 3 | *            |        | 11      | C            | printf |
| 4       | 4 | /            |        | 13      | NULL         | _      |
|         | 5 | A            |        | 2       | *            | printf |
| (       | 6 | NULL         |        | 14      | D            | -      |
|         | 5 | $\mathbf{A}$ | printf | 15      | NULL         |        |
|         | 7 | NULL         |        | 14      | D            | printf |
| -       | 4 | /            | printf | 16      | NULL         | _      |
|         | 8 | В            |        | 1       | +            | printf |
| 9       | 9 | NULL         |        | 17      | Е            | _      |
|         | 8 | В            | printf | 18      | NULL         |        |
| 10      | 0 | NULL         |        | 17      | $\mathbf{E}$ | printf |
| 3       | 3 | *            | printf | 19      | NULL         |        |

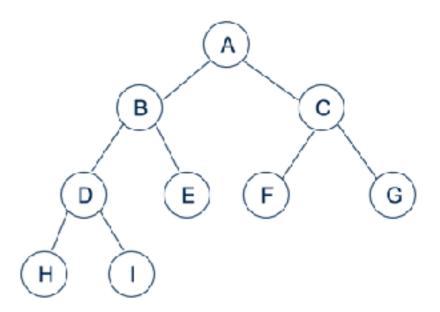
```
void preorder(Tree ptr) {
   if(ptr) {
       printf("%d", ptr->data);
       preorder(ptr->left_child);
       preorder(ptr->right_child);
 void postorder(Tree ptr) {
    if(ptr) {
        postorder(ptr->left_child);
        postorder(ptr->right_child);
        printf("%d", ptr->data);
```

iterative in-order traversal using stack

```
void iterInorder (Tree node) {
     int top = -1
     Tree stack[MAX_SIZE];
     for (;;) {
          for (; node; node = node -> leftChild)
               push(node);
          node = pop();  // pop parent
          if (!node) break;
          printf("%d", node -> data);
          node = node -> rightChild;
```



level-order traversal

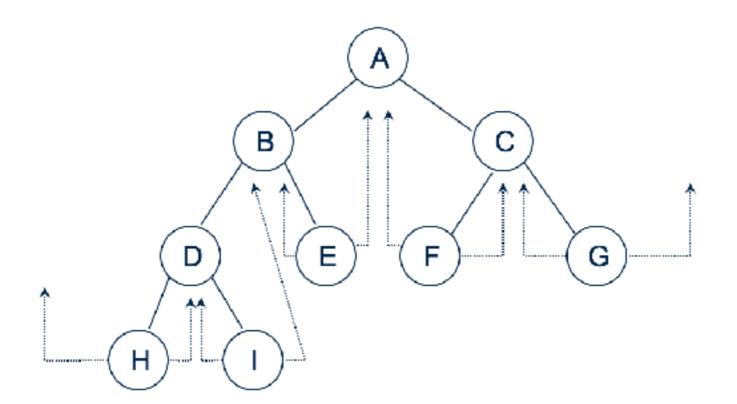


level-order traversal

```
void levelOrder (Tree ptr) {
     int front = rear = 0;
     Tree queue[MAX];
     if (! node)
                 return;
     addq(ptr);
     for (;;) {
           ptr = deleteq();
           if (ptr) {
                printf("%d", ptr->data);
                if (ptr -> leftChild)
                      addq(ptr -> leftChild);
                if (ptr -> rightChild)
                      addq(ptr -> rightChild);
           else break;
```

## threaded binary trees

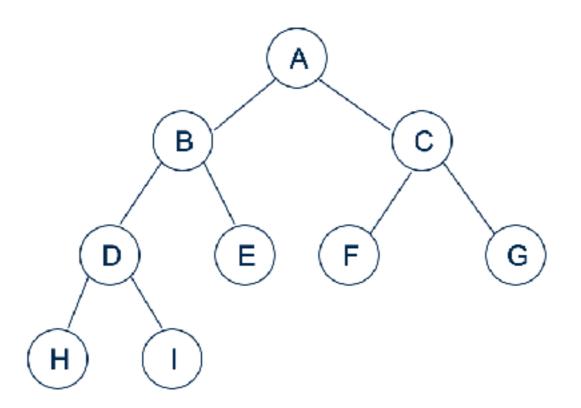
- there are n+1 null links out of 2n total links
- replace the null links by pointers, called threads to other nodes in the tree
  - if ptr -> leftChild is null, replace the null with a pointer to the node that would be visited before ptr in an in-order traversal
  - if ptr -> rightChild is null, replace the null with a pointer to the node that would be visited after ptr in an in-order traversal



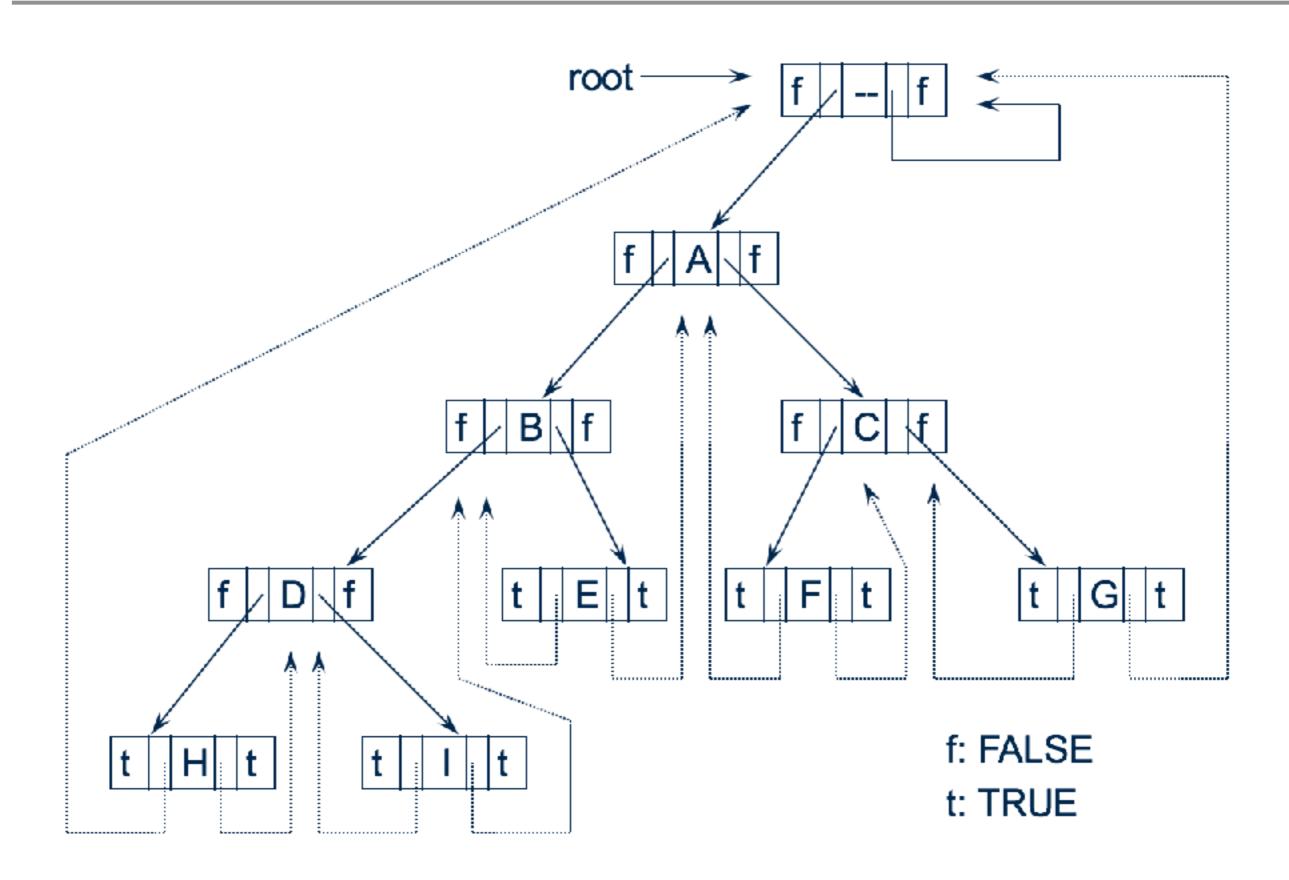
#### threaded binary trees

- How to distinguish actual pointers and threads?
  - →add two additional fields to the node structure
    - if ptr->left\_thread = true, ptr->left\_child contains thread
    - if ptr->left\_thread = false, ptr->left\_child contains a pointer to the left child

```
typedef struct threaded_tree *threaded_ptr;
typedef struct threaded_tree {
    short int left_thread;
    threaded_ptr left_child;
    char data;
    threaded_ptr right_child;
    short int right_thread;
};
```



# threaded binary trees

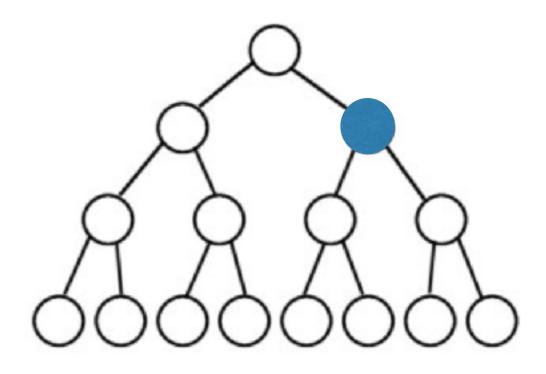


- find the in-order successor of ptr without using stack
  - if ptr -> right\_thread = TRUE, ptr -> right\_child
  - otherwise follow a path of left\_child links from the right\_child of ptr until we reach a node with left thread = TRUE

```
threaded_ptr insucc(threaded_ptr tree) {
    threaded_ptr temp;
                                                                 root
    temp = tree->right_child;
    if (!tree->right_thread)
        while (!temp->left_thread)
             temp = temp->left_child;
    return temp;
                                                           f B f
                                                                E
                                                                          t
                                                                             F
                                                                                             G t
                                                              t
                                                                                     f: FALSE
                                             H
                                                                                     t: TRUE
```

- find the in-order successor of ptr without using stack
  - if ptr -> right\_thread = TRUE, ptr -> right\_child
  - otherwise follow a path of left\_child links from the right\_child of ptr until we reach a node with left thread = TRUE

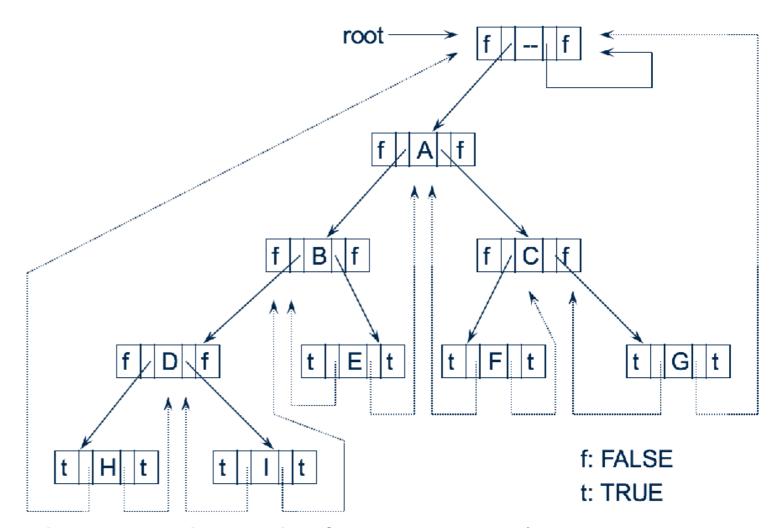
```
threaded_ptr insucc(threaded_ptr tree) {
    threaded_ptr temp;
    temp = tree->right_child;
    if (!tree->right_thread)
        while (!temp->left_thread)
        temp = temp->left_child;
    return temp;
}
```



Which node will be returned if blue node is passed into the function insucc?

- find the in-order successor of ptr without using stack
  - if ptr -> right\_thread = TRUE, ptr -> right\_child
  - otherwise follow a path of left\_child links from the right\_child of ptr until we reach a node with left thread = TRUE

```
threaded_ptr insucc(threaded_ptr tree) {
    threaded_ptr temp;
    temp = tree->right_child;
    if (!tree->right_thread)
        while (!temp->left_thread)
        temp = temp->left_child;
    return temp;
}
```



Which node will be returned if root node is passed into the function insucc?

```
void tinorder(threaded_ptr tree) {
     threaded_ptr temp = tree;
     for (;;) {
           temp = insucc(temp);
           if (temp == tree) break;
           printf("%3c", temp->data);
                                                     root
                                                  В
                                                     E
                                                             t
                                                                              G t
                                                                       f: FALSE
                                              I | t
                                           t
                                  H
                                                                       t: TRUE
```