

All 4-variable functions can be perfectly quadratized with only 1 auxiliary variable

Nike Dattani¹ and Hou Tin Chau²

¹Harvard-Smithsonian Center for Astrophysics, USA

²Cambridge University, Department of Mathematics, UK

We prove that any function whose input is 4 binary variables and whose output is a real number, is perfectly equivalent to a *quadratic* function whose input is 5 binary variables and is minimized over the new variable. Our proof is constructive, so we provide quadratic functions that quadratize any 4-variable function, but there exists 5 different classes of 4-variable functions that each have their own 5-variable quadratization formula. Since we provide ‘perfect’ quadratizations, we can apply these formulas to any 4-variable subset of an n -variable function even if $n \gg 4$.

I. INTRODUCTION

Many problems can be solved by minimizing a real-valued degree- k function of binary variables with $k > 2$. Some examples include image de-blurring (where typically $k = 4$ but in general we can have $k = m^2$ with $m \geq 2$ being the length in pixels of the square-shaped mask) [1, 2], integer factoring (where typically $k = 4$) [3–10], and determining whether or not a number N is an m -color Ramsey number (where $k = \frac{mN(N-1)}{2}$) [11–13].

Solving such discrete optimization problems with $k > 2$ can be very difficult, and more algorithms have been developed for the $k = 2$ case (such as the algorithm known as “QPBO” and extensions of it [14]) than for the $k > 2$ case. Fortunately it is possible to turn any k -degree binary optimization problem into a 2-degree binary optimization problem, by a transformation called “quadratization” [15].

Quadratization methods exist which can turn an n -variable degree- k problem into an n -variable quadratic problem (i.e. the number of variables does not change) [8, 13, 16, 17], but not every function can be quadratized without adding some auxiliary variables (so the number of variables in the quadratic problem is usually much more than in the original degree- k problem). Coming up with better quadratizations (for example with fewer auxiliary variables) has been a very active area of research recently: The first quadratization method was published in 1975 [18], and some subsequent quadratization methods were published in 2004 [19], 2005 [20], and 2011 [1, 2, 21], but the rest of the methods were published in the last 5 years (from 2014–2018) [8, 13, 15–17, 22–32].

In the most recent of these papers [31, 32], a remarkable discovery was made, that degree- k monomials can be quadratized with only $\log_2 k - 1$ auxiliary variables. For many functions this can still be prohibitively costly though: If a 44-variable function has 1 million degree-5 terms and each term requires $\log_2 k - 1$ auxiliary variables for quadratization, the quadratic function will have more than 2 million variables (the search space increases from $2^{44} \approx 10^{13}$ to $2^{2,000,044} \approx 10^{602,073}$).

It was also shown in [31, 32] that sometimes the entire function of n variables can be quadratized with only $\log_2 k - 1$ auxiliary variables no matter how many terms and how many variables it contains (so a 44-variable degree-5 function with 1 million degree-5 terms would only need 2 auxiliary variables rather than 2 million!). However, it is only known how to do this very “compact” quadratization for a very specific class of functions called “at-least- k -of- n ” ($AkON$) functions, which includes the case of a single positive monomial term.

Learning from [31, 32] that it is possible to quadratize very large functions so compactly inspired us, and the fact that such “compact” quadratizations are only known for a very specific category of functions (the $AkON$ functions), motivated us to look for quadratizations that are “compact”, but also applicable to a much wider class of functions. The result of this study is the theorem described in the title of this paper, and explained in more detail in the section below. It allows up to 5 terms of a function (1 of them can be of degree-4 and the other 4 can be of degree-3) to be quadratized with only 1 auxiliary variable rather than 5 auxiliaries (which is what would be required if quadratizing each term individually with $\log_2 k - 1$ auxiliary variables for each term).

II. RESULTS

Theorem 1: All 4-variable functions of binary variables with real-valued coefficients can be quadratized perfectly with only 1-auxiliary variable.

By “perfect” quadratization we mean all 2^4 output values of the 4-variable function are exactly preserved when minimizing over the auxiliary variable in the 5-variable quadratic function. Therefore any 4-variable subset of an n -variable problem can be quadratized with only 1-auxiliary variable.

We prove the theorem by providing an explicit quadratization for various different cases, of the following function of binary variables $b_i \in \{0, 1\}$ with real-valued coefficients α :

$$\alpha_{1234}b_1b_2b_3b_4 + \alpha_{123}b_1b_2b_3 + \alpha_{124}b_1b_2b_4 + \alpha_{134}b_1b_3b_4 + \alpha_{234}b_2b_3b_4. \quad (1)$$

Since $\alpha_{123}, \alpha_{124}, \alpha_{134}$ and α_{234} are completely symmetric (they can be switched with each other and have their subscripts relabeled without any effect on the function), we can order them however we desire, so for convenience we choose for the rest of this paper: $\alpha_{123} \leq \alpha_{124} \leq \alpha_{134} \leq \alpha_{234}$.

We will now provide 5 different quadratization formulas for Eq. 1 (Lemmas 1-5), which each are only valid for their own specific conditions on the α coefficients; but we will then prove with Lemmas 6-8 that these 5 cases for the coefficients, cover every possible case. The explicit quadratizations for Lemmas 1-5 are given below, but their proofs take up a lot of space so they are given in the Appendix.

Lemma 1: If $\alpha_{1234} \geq 0$, and either $\alpha_{ijk} \geq -\frac{\alpha_{1234}}{2}$ for all ijk , or $-\alpha_{1234} \leq \alpha_{123} \leq -\frac{\alpha_{1234}}{2} \leq 0 \leq \alpha_{234} \leq \alpha_{134} \leq \alpha_{124}$, then Eq. 1 is perfectly quadratized by:

$$\left(3\alpha_{1234} + \sum_{ijk} \alpha_{ijk}\right) b_a + \alpha_{1234} \sum_{ij} b_i b_j + \sum_{ij} \sum_{k \notin ij} \alpha_{ijk} b_i b_j - \sum_i \left(2\alpha_{1234} + \sum_{jk, i \neq jk} \alpha_{ijk}\right) b_i b_a. \quad (2)$$

Lemma 2: If $\alpha_{1234} \leq 0$ and $\alpha_{ijk} \leq 0$, then Eq. 1 is perfectly quadratized by:

$$\left(\alpha_{1234} \left(\sum_i b_i - 3\right) + \sum_{ijk} \alpha_{ijk} \left(\sum_{l \in ijk} b_l - 2\right)\right) b_a. \quad (3)$$

Lemma 3: If $\alpha_{1234} \geq 0$, $\alpha_{123} \leq -\alpha_{1234}$, and $-\frac{\alpha_{1234}}{2} \leq \alpha_{124} \leq \alpha_{134} \leq 0 \leq \alpha_{234}$, then Eq. 1 is perfectly quadratized by:

$$\alpha_{1234} - \sum_i (\alpha_{12i} + \alpha_{1234}) b_i + \sum_i \alpha_{i34} b_a + \sum_{ijk} \alpha_{ijk} b_i b_j + \alpha_{1234} b_3 b_4 - \sum_{\substack{i=p,q \\ p,q=1,2 \text{ or } 3,4}} \left(\sum_{\substack{j=r,s \\ r,s=3,4 \text{ or } 1,2}} \alpha_{pqj} - \alpha_{irs}\right) b_i b_a. \quad (4)$$

Lemma 4: If $\alpha_{1234} \geq 0$, $\alpha_{123} \leq -\frac{\alpha_{1234}}{2} \leq \alpha_{124} \leq \alpha_{134} \leq \alpha_{234} \leq 0$, and $\alpha_{123} + \alpha_{124} \leq -\alpha_{1234}$, then Eq. 1 is perfectly quadratized by:

$$\begin{aligned} \alpha_{1234} - \sum_i (\alpha_{12i} + \alpha_{1234}) b_i + \sum_i \alpha_{i34} b_a + \sum_{ijk} \alpha_{ijk} b_i b_j + \alpha_{1234} b_3 b_4 - \sum_{\substack{i=p,q \\ p,q=1,2 \text{ or } 3,4}} \left(\sum_{\substack{j=r,s \\ r,s=3,4 \text{ or } 1,2}} \alpha_{pqj} - \alpha_{irs}\right) b_i b_a \\ - \sum_i \alpha_{12i} (b_1 + b_2) - \sum_i \alpha_{i34} (b_3 + b_4 - 1 - b_i) - \alpha_{1234} (b_3 + b_4 - 1) b_a. \end{aligned} \quad (5)$$

Lemma 5: If $\alpha_{1234} \geq 0$, $-\alpha_{1234} \leq \alpha_{123} \leq -\frac{\alpha_{1234}}{2} \leq \alpha_{124} \leq \alpha_{134} \leq \alpha_{234} \leq 0$, and $\alpha_{123} + \alpha_{124} \geq -\alpha_{1234}$, then Eq. 1 is perfectly quadratized by:

$$\sum_{ijk} \alpha_{ijk} \sum_{lm \subset ijk} b_l b_m + \alpha_{1234} \sum_{ij} b_i b_j + b_a \left(\sum_{ijk} \alpha_{ijk} \left(1 - \sum_{l \in ijk} b_l\right) + \alpha_{1234} \left(3 - 2 \sum_i b_i\right)\right). \quad (7)$$

In the Appendix, Lemmas 1-5 are each proven for their own specific conditions on the coefficients α . However with ‘bit-flipping’ (a strategy described in [1] and on Pg. 27 of the current version of [15]) we can extend their applicability to more general conditions for which a laborious proof was not performed explicitly. Lemma 6 will describe the effect of flipping one bit in Eq. 1, and Lemma 7 will describe the effect of flipping two. Since the function is completely symmetric with respect to the four variables b_1, b_2, b_3 and b_4 , these Lemmas depend only on the number of bits flipped and not at all on which bits are flipped.

Lemma 6: If one bit is flipped ($b_1 \rightarrow \bar{b}_1 \equiv 1 - b_1$) everywhere in Eq. 1, then the function remains exactly the same except with $\bar{\alpha}_{1234} \equiv -\alpha_{1234}$, $\bar{\alpha}_{123} \equiv -\alpha_{123}$, $\bar{\alpha}_{124} \equiv -\alpha_{124}$, $\bar{\alpha}_{134} \equiv -\alpha_{134}$, $\bar{\alpha}_{234} \equiv \alpha_{234} + \alpha_{1234}$.

Proof: We start with Eq. 1 but with every occurrence of b_1 replaced by its flipped version:

$$\alpha_{1234}\bar{b}_1b_2b_3b_4 + \alpha_{123}\bar{b}_1b_2b_3 + \alpha_{124}\bar{b}_1b_2b_4 + \alpha_{134}\bar{b}_1b_3b_4 + \alpha_{234}b_2b_3b_4. \quad (8)$$

Expanding \bar{b}_1 as $1 - b_1$, and completely expanding the expressions for each term of Eq. 8, we get:

$$\alpha_{1234}(b_2b_3b_4 - b_1b_2b_3b_4) + \alpha_{123}(b_2b_3 - b_1b_2b_3) + \alpha_{124}(b_2b_4 - b_1b_2b_4) + \alpha_{134}(b_3b_4 - b_1b_3b_4) + \alpha_{234}b_2b_3b_4 \quad (9)$$

We can now regroup everything in Eq. 9 such that it is back in the form of Eq. 1, except with new coefficients:

$$-\alpha_{1234}b_1b_2b_3b_4 - \alpha_{123}b_1b_2b_3 - \alpha_{124}b_1b_2b_4 - \alpha_{134}b_1b_3b_4 + (\alpha_{1234} + \alpha_{234})b_2b_3b_4 \quad (10)$$

$$= \bar{\alpha}_{1234}b_1b_2b_3b_4 + \bar{\alpha}_{123}b_1b_2b_3 + \bar{\alpha}_{124}b_1b_2b_4 + \bar{\alpha}_{134}b_1b_3b_4 + \bar{\alpha}_{234}b_2b_3b_4. \quad (11)$$

Lemma 7: If two bits are flipped ($b_1 \rightarrow \bar{b}_1 \equiv 1 - b_1$, $b_2 \rightarrow \bar{b}_2 \equiv 1 - b_2$) everywhere in Eq. 1, the function remains exactly the same except with $\bar{\alpha}_{134} \equiv -(\alpha_{134} + \alpha_{1234})$, $\bar{\alpha}_{234} \equiv -(\alpha_{234} + \alpha_{1234})$, and some extra quadratic terms: $f_{\text{quadratic},2}(b_1, b_2, b_3, b_4) \equiv \alpha_{1234}b_3b_4 + \alpha_{123}(b_3 - b_1b_3 - b_2b_3) + \alpha_{124}(b_4 - b_1b_4 - b_2b_4) + (\alpha_{134} + \alpha_{234})b_3b_4$.

Proof: We start with Eq. 1 but with every occurrence of b_1 and b_2 replaced by their flipped versions:

$$\alpha_{1234}\bar{b}_1\bar{b}_2b_3b_4 + \alpha_{123}\bar{b}_1\bar{b}_2b_3 + \alpha_{124}\bar{b}_1\bar{b}_2b_4 + \alpha_{134}\bar{b}_1b_3b_4 + \alpha_{234}\bar{b}_2b_3b_4. \quad (12)$$

Expanding \bar{b}_1 and \bar{b}_2 as $1 - b_1$ and $1 - b_2$ respectively, and completely expanding the expressions for each term of Eq. 12, we get:

$$\alpha_{1234}(b_3b_4 - b_2b_3b_4 - b_1b_3b_4 + b_1b_2b_3b_4) + \alpha_{123}(b_3 - b_1b_3 - b_2b_3 + b_1b_2b_3) + \alpha_{124}(b_4 - b_1b_4 - b_2b_4 + b_1b_2b_4) + \alpha_{134}(b_3b_4 - b_1b_3b_4) + \alpha_{234}(b_3b_4 - b_2b_3b_4). \quad (13)$$

We can now regroup everything in Eq. 13 such that it is back in the form of Eq. 1, except with new coefficients for two of the terms, and some extra quadratics terms:

$$\alpha_{1234}b_1b_2b_3b_4 + \alpha_{123}b_1b_2b_3 + \alpha_{124}b_1b_2b_4 - (\alpha_{134} + \alpha_{1234})b_1b_3b_4 - (\alpha_{234} + \alpha_{1234})b_2b_3b_4 \quad (14)$$

$$+ \alpha_{1234}b_3b_4 + \alpha_{123}(b_3 - b_1b_3 - b_2b_3) + \alpha_{124}(b_4 - b_1b_4 - b_2b_4) + (\alpha_{134} + \alpha_{234})b_3b_4. \quad (15)$$

Lemma 6 allows us to assume from now on that $\alpha_{1234} \geq 0$, because every case with $\alpha_{1234} < 0$ can be turned into a case with $\alpha_{1234} > 0$ by flipping only one bit. With $\alpha_{1234} \geq 0$, we can categorize all cubic coefficients α_{ijk} according to whether they are $\leq -\alpha_{1234}$, or $\leq -\frac{\alpha_{1234}}{2}$, or whether they are simply just ≤ 0 or ≥ 0 . There's then 35 different cases for how the four cubic coefficients α_{ijk} can fit into the four different non-overlapping intervals that can be made on the number line with $-\alpha_{1234}$, $-\frac{\alpha_{1234}}{2}$ and 0 as partition points. For some of these cases, Lemma 1, 3, 4 or 5 can be applied immediately. Due to the conditions used to prove Lemma 2, it cannot be applied directly to any of the 35 cases, but thanks to Lemmas 6 and 7, Lemma 2 can be applied to two of the 35 cases after bit-flipping appropriately. Lemma 1, 3, 4 or 5 can be applied for the rest of the 35 cases if 2, 3, or 4 bits are flipped (meaning either one application of Lemma 7, one application of Lemma 6 combined with one application of Lemma 7, or two applications of Lemma 7, is done). This means that one of Lemmas 1-5 can be applied for all of the 35 possible cases, as long as Lemmas 6 and/or 7 are applied appropriately. Table I summarizes which bits have to be flipped using Lemma 6 and/or Lemma 7, and which of Lemmas 1-5 can be applied, for each of the 35 possible cases.

Table I. All possible cases of 4-variable functions with $\alpha_{1234} \geq 0$.

$\alpha_{ijk} \leq -\alpha_{1234}$	$-\alpha_{1234} \leq \alpha_{ijk} \leq -\frac{\alpha_{1234}}{2}$	$-\frac{\alpha_{1234}}{2} \leq \alpha_{ijk} \leq 0$	$0 \leq \alpha_{ijk}$	Bits flipped	Quadratization
			$\alpha_{123}, \alpha_{124}, \alpha_{134}, \alpha_{234}$	-	Lemma 1
		α_{123}	$\alpha_{124}, \alpha_{134}, \alpha_{234}$	-	Lemma 1
α_{123}	α_{123}		$\alpha_{124}, \alpha_{134}, \alpha_{234}$	-	Lemma 1
		$\alpha_{123}, \alpha_{124}$	$\alpha_{124}, \alpha_{134}, \alpha_{234}$	b_4	Lemma 2
	α_{123}	$\alpha_{123}, \alpha_{124}$	$\alpha_{134}, \alpha_{234}$	-	Lemma 1
α_{123}	α_{123}	α_{124}	$\alpha_{134}, \alpha_{234}$	b_2, b_4	Lemma 3
		α_{124}	$\alpha_{134}, \alpha_{234}$	b_3, b_4	Lemma 1
	$\alpha_{123}, \alpha_{124}$		$\alpha_{134}, \alpha_{234}$	b_3, b_4	Lemma 1
α_{123}	$\alpha_{123}, \alpha_{124}$		$\alpha_{134}, \alpha_{234}$	b_3, b_4	Lemma 1
$\alpha_{123}, \alpha_{124}$	α_{124}		$\alpha_{134}, \alpha_{234}$	b_3, b_4	Lemma 1
		$\alpha_{123}, \alpha_{124}, \alpha_{134}$	α_{234}	-	Lemma 1
	α_{123}	$\alpha_{124}, \alpha_{134}$	α_{234}	b_1, b_4	Lemma 4
α_{123}		$\alpha_{124}, \alpha_{134}$	α_{234}	-	Lemma 3
	$\alpha_{123}, \alpha_{124}$	α_{134}	α_{234}	b_3, b_4	Lemma 1
α_{123}	α_{124}	α_{134}	α_{234}	b_3, b_4	Lemma 1
$\alpha_{123}, \alpha_{124}$		α_{134}	α_{234}	b_3, b_4	Lemma 1
	$\alpha_{123}, \alpha_{124}, \alpha_{134}$		α_{234}	b_1, b_2, b_3, b_4	Lemma 4
α_{123}	$\alpha_{124}, \alpha_{134}$		α_{234}	b_2, b_3	Lemma 3
$\alpha_{123}, \alpha_{124}$	α_{134}		α_{234}	b_3, b_4	Lemma 1
$\alpha_{123}, \alpha_{124}, \alpha_{134}$			α_{234}	b_2, b_3, b_4	Lemma 2
	α_{123}	$\alpha_{123}, \alpha_{124}, \alpha_{134}, \alpha_{234}$		-	Lemma 1
α_{123}		$\alpha_{124}, \alpha_{134}, \alpha_{234}$		-	Lemma 4, 5
		$\alpha_{124}, \alpha_{134}, \alpha_{234}$		-	Lemma 4
	$\alpha_{123}, \alpha_{124}$	$\alpha_{134}, \alpha_{234}$		b_3, b_4	Lemma 1
α_{123}	α_{124}	$\alpha_{134}, \alpha_{234}$		b_3, b_4	Lemma 1
$\alpha_{123}, \alpha_{124}$		$\alpha_{134}, \alpha_{234}$		b_3, b_4	Lemma 1
	$\alpha_{123}, \alpha_{124}, \alpha_{134}$	α_{234}		b_3, b_4	Lemma 4, 5
α_{123}	$\alpha_{124}, \alpha_{134}$	α_{234}		b_2, b_3	Lemma 4
$\alpha_{123}, \alpha_{124}$	α_{134}	α_{234}		b_2, b_3	Lemma 3
$\alpha_{123}, \alpha_{124}, \alpha_{134}$		α_{234}		b_1, b_2, b_3, b_4	Lemma 1
	$\alpha_{123}, \alpha_{124}, \alpha_{134}, \alpha_{234}$			b_1, b_2, b_3, b_4	Lemma 1
α_{123}	$\alpha_{124}, \alpha_{134}, \alpha_{234}$			b_1, b_2, b_3, b_4	Lemma 1
$\alpha_{123}, \alpha_{124}$	$\alpha_{134}, \alpha_{234}$			b_1, b_2, b_3, b_4	Lemma 1
$\alpha_{123}, \alpha_{124}, \alpha_{134}$	α_{234}			b_1, b_2, b_3, b_4	Lemma 1
$\alpha_{123}, \alpha_{124}, \alpha_{134}, \alpha_{234}$				b_1, b_2, b_3, b_4	Lemma 1

Table II. Reduced version of Table I

$\alpha_{ijk} \leq -\alpha_{1234}$	$-\alpha_{1234} \leq \alpha_{ijk} \leq -\frac{\alpha_{1234}}{2}$	$-\frac{\alpha_{1234}}{2} \leq \alpha_{ijk} \leq 0$	$0 \leq \alpha_{ijk}$	Bits flipped	Quadratization
		$\alpha_{123}, \alpha_{124}, \alpha_{134}, \alpha_{234}$		-	Lemma 1
	α_{123}		$\alpha_{124}, \alpha_{134}, \alpha_{234}$	-	Lemma 1
α_{123}			$\alpha_{124}, \alpha_{134}, \alpha_{234}$	b_4	Lemma 2
	α_{123}	α_{124}	$\alpha_{134}, \alpha_{234}$	b_2, b_3	Lemma 3
α_{123}		α_{124}	$\alpha_{134}, \alpha_{234}$	b_3, b_4, b_3, b_4	Lemma 1
$\alpha_{123}, \alpha_{124}$		$\alpha_{134}, \alpha_{234}$		b_3, b_4, b_3, b_4	Lemma 1
	α_{123}	$\alpha_{124}, \alpha_{134}$	α_{234}	b_1, b_4	Lemma 4
α_{123}		$\alpha_{124}, \alpha_{134}$	α_{234}	-	Lemma 3
	$\alpha_{123}, \alpha_{124}, \alpha_{134}$		α_{234}	b_1, b_2, b_3, b_4	Lemma 4
α_{123}	$\alpha_{124}, \alpha_{134}$		α_{234}	b_2, b_3	Lemma 3
$\alpha_{123}, \alpha_{124}$	α_{134}		α_{234}	b_3, b_4	Lemma 1
$\alpha_{123}, \alpha_{124}, \alpha_{134}$			α_{234}	b_2, b_3, b_4	Lemma 2
	α_{123}	$\alpha_{124}, \alpha_{134}, \alpha_{234}$		-	Lemma 4 & 5
α_{123}		$\alpha_{124}, \alpha_{134}, \alpha_{234}$		-	Lemma 4
	$\alpha_{123}, \alpha_{124}, \alpha_{134}$	α_{234}		b_3, b_4	Lemma 4 & 5
α_{123}	$\alpha_{124}, \alpha_{134}$	α_{234}		b_2, b_3	Lemma 4
$\alpha_{123}, \alpha_{124}$	α_{134}	α_{234}		b_2, b_3	Lemma 3
$\alpha_{123}, \alpha_{124}, \alpha_{134}$		α_{234}		b_1, b_2, b_3, b_4	Lemma 1
$\alpha_{123}, \alpha_{124}, \alpha_{134}, \alpha_{234}$				b_1, b_2, b_3, b_4	Lemma 1

III. DISCUSSION

A. Example

To quadratize the 15-term, 5-variable, degree-4 function with all terms being at least cubic:

$$5b_1b_2b_3b_4 + 4b_1b_2b_3b_5 + 3b_1b_2b_4b_5 + 2b_1b_3b_4b_5 + b_2b_3b_4b_5 - 3b_1b_2b_3 - b_1b_2b_4 - 5b_1b_2b_5 - b_1b_3b_4 - b_1b_3b_5 - b_1b_4b_5 - b_2b_3b_4 - b_2b_3b_5 - 4b_2b_4b_5 - 3b_3b_4b_5, \quad (16)$$

normally we would need at least 15 auxiliary variables (at least one for each term). Instead we quadratize the entire function with only 5 auxiliary variables. To do this we will apply Theorem 1 for 5 sub-functions (displayed below in five different colors) that contain only 4 variables:

$$5b_1b_2b_3b_4 + 4b_1b_2b_3b_5 + 3b_1b_2b_4b_5 + 2b_1b_3b_4b_5 + b_2b_3b_4b_5 - 3b_1b_2b_3 - b_1b_2b_4 - 5b_1b_2b_5 - b_1b_3b_4 - b_1b_3b_5 - b_1b_4b_5 - b_2b_3b_4 - b_2b_3b_5 - 4b_2b_4b_5 - 3b_3b_4b_5. \quad (17)$$

We now quadratize these 5 sub-functions with only 1 auxiliary variable for each sub-function:

$$5b_1b_2b_3b_4 - 3b_1b_2b_3 - b_1b_2b_4 - b_1b_3b_4 - b_2b_3b_4 \rightarrow b_1b_2 + b_1b_3 + 3b_1b_4 + b_2b_3 + 3b_2b_4 + 3b_3b_4 - b_{a_1}(5b_1 + 5b_2 + 5b_3 + 7b_4 - 9) \quad (18)$$

$$4b_1b_2b_3b_5 - 5b_1b_2b_5 - b_1b_3b_5 - b_2b_3b_5 \rightarrow -3b_1 + 6b_2 - 3b_3 + 5b_5 - 5b_1b_2 + 3b_1b_3 - 5b_1b_5 - b_2b_3 - b_3b_5 \quad (19)$$

$$-b_{a_2}(-8b_1 + 6b_2 - 4b_3 + 5b_5 + 3) + 3 \quad (20)$$

$$3b_1b_2b_4b_5 - b_1b_4b_5 - 4b_2b_4b_5 \rightarrow b_1 + 4b_2 + 3b_1b_2 - b_1b_4 - b_1b_5 - 4b_2b_4 - 4b_2b_5 + b_{a_3}(-4b_1 - 7b_2 + 5b_4 + 5b_5 + 3) \quad (21)$$

$$2b_1b_3b_4b_5 - 3b_3b_4b_5 \rightarrow b_{a_4}(2b_1 - 3b_3 - 3b_4 - 3b_5 + 6) \quad (22)$$

$$b_2b_3b_4b_5 \rightarrow b_2b_3 + b_2b_4 + b_2b_5 + b_3b_4 + b_3b_5 + b_4b_5 + b_{a_5}(3 - 2b_2 - 2b_3 - 2b_4 - 2b_5) \quad (23)$$

The final quadratic function contains only 10 variables (the 5 original ones and the 5 new auxiliary variables), whereas applying the previous state-of-the-art quadratization techniques would in the best case result in a quadratic function with 20 variables (the 5 original ones and the 15 new auxiliary variables).

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