

# All 4-variable functions can be perfectly quadratized with only 1 auxiliary variable

Nike Dattani<sup>1</sup> and Hou Tin Chau<sup>2</sup>

<sup>1</sup>Harvard-Smithsonian Center for Astrophysics, USA

<sup>2</sup>Cambridge University, Department of Mathematics, UK

We prove that any function whose input is 4 binary variables and whose output is a real number, is perfectly equivalent to a function whose input is 5 binary variables and is minimized over the new variable. Our proof is constructive, so we provide quadratic functions that quadratize any 4-variable function, but there exists 8 different classes of 4-variable functions that each have their own 5-variable quadratization formula. Since we provide ‘perfect’ quadratizations, we can apply these formulas to any 4-variable subset of an  $n$ -variable function even if  $n \gg 4$ .

## I. INTRODUCTION

Many problems can be solved by minimizing a real-valued degree- $k$  function of binary variables with  $k > 2$ . Some examples include image de-blurring (where typically  $k = 4$  but in general we can have  $k = m^2$  with  $m \geq 2$  being the length in pixels of the square-shaped mask) [1, 2], integer factoring (where typically  $k = 4$ ) [3–10], and determining whether or not a number  $N$  is an  $m$ -color Ramsey number (where  $k = \frac{mN(N-1)}{2}$ ) [11–13].

Solving such discrete optimization problems with  $k > 2$  can be very difficult, and more algorithms have been developed for the  $k = 2$  case (such as the algorithm known as “QPBO” and extensions of it [14]) than for the  $k > 2$  case. Fortunately it is possible to turn any  $k$ -degree binary optimization problem into a 2-degree binary optimization problem, by a transformation called “quadratization” [15].

Quadratization methods exist which can turn an  $n$ -variable degree- $k$  problem into an  $n$ -variable quadratic problem (i.e. the number of variables does not change) [8, 13, 16, 17], but not every case can be quadratized without adding some auxiliary variables (so the number of variables in the quadratic problem is usually much more than in the original degree- $k$  problem). Coming up with better quadratizations (for example with fewer auxiliary variables) has been a very active area of research recently: The first quadratization method was published in 1975 [18], and some subsequent quadratization methods were published in 2004 [19], 2005 [20], and 2011 [1, 2, 21], but the rest of the methods were published in the last 5 years (from 2014–2018) [8, 13, 22–30].

In the most recent of these papers [29, 30], a remarkable discovery was made, that some entire functions (no matter how many terms and how many variables they contain) can be quadratized with only  $\log_2(k) - 1$  auxiliary variables. Unfortunately those functions are either extremely obscure functions that would rarely come up in a real-world problem, or they are just monomials (meaning that for a real-world problem containing many terms, each of these terms would be quadratized separately, with  $\log_2(k) - 1$  new auxiliary variables being added each time, and adding up to a potentially untractable number of total variables in the quadratic optimization problem).

This motivated us to look for quadratizations that are compact, but also applicable to many real-world problems. The result of this study is the theorem described in the title of this paper, and explained in more detail in the section below.

## II. RESULTS

**Theorem I: All 4-variable functions can be quadratized perfectly with only 1-auxiliary variable.**

By “perfect” quadratization we mean all  $2^4$  outcomes of the 4-variable function are exactly preserved when minimizing over the auxiliary variable in the 5-variable quadratic function. Therefore any 4-variable subset of an  $n$ -variable problem can be quadratized with only 1-auxiliary.

We prove the theorem by providing an explicit quadratization for various different cases, of the following function of binary variables  $b_i \in \{0, 1\}$ :

$$\alpha_{1234}b_1b_2b_3b_4 + \alpha_{123}b_1b_2b_3 + \alpha_{234}b_2b_3b_4 + \alpha_{134}b_1b_3b_4 + \alpha_{124}b_1b_2b_4. \quad (1)$$

**Case 1:**  $\alpha_i \leq 0, \forall i$

$$b_a \left( \alpha_{1234} \left( \sum_i b_i - 3 \right) + \sum_i \alpha_i \left( \sum_j b_j - 2 \right) \right) \quad (2)$$

**Case 2:**  $\alpha_{1234} \geq 0$

*Case 2.1:*  $-\alpha_{1234} \geq \alpha_{123}$

*Case 2.1.1:*  $-\alpha_{1234}/2 \leq \alpha_{234} \leq \alpha_{134} \leq \alpha_{124} \leq 0$

$$b_a (\alpha_{123} (-b_1 + b_2 + b_3) + \alpha_{234} (b_2 + b_3 - b_4) + \alpha_{134} (b_1 - b_3 + b_4 - 1) + \alpha_{124} (b_1 - b_2 + b_4 - 1) + \alpha_{1234} (b_1 + b_4 - 1)) \quad (3)$$

$$+ \alpha_{123} (-b_2 - b_3 + b_1b_2 + b_1b_3) + \alpha_{234} (-b_2 - b_3 + b_2b_4 + b_3b_4) \quad (4)$$

$$+ \alpha_{134} (1 - b_1 + b_3 - b_4 + b_1b_4) + \alpha_{124} (1 - b_1 + b_2 - b_4 + b_1b_4) + \alpha_{1234} (1 - b_1 - b_4 + b_1b_4) \quad (5)$$

*Case 2.1.2:*  $-\alpha_{1234}/2 \leq \alpha_{234} \leq \alpha_{134} \leq 0 \leq \alpha_{124}$

$$b_a (\alpha_{123} (b_1 - b_2 - b_3) + \alpha_{234} (-b_2 - b_3 + b_4) + \alpha_{134} (1 - b_1 + b_3 - b_4) + \alpha_{124} (1 - b_1 + b_2 - b_4) + \alpha_{1234} (1 - b_1 - b_4)) \quad (6)$$

$$+ \alpha_{123} (-b_1 + b_1b_2 + b_1b_3) + \alpha_{234} (-b_4 + b_2b_4 + b_3b_4) + (\alpha_{134} + \alpha_{124} + \alpha_{1234}) b_1b_4 \quad (7)$$

*Case 2.2:*  $-\frac{\alpha_{1234}}{2} \geq \alpha_{123}$

*Case 2.2.1:*  $\alpha_{1234} + \alpha_{123} + \alpha_{234} \leq 0$

*Case 2.2.1.1:*  $-\alpha_{1234} \leq \alpha_{123} \leq \alpha_{1234}/2 \leq \alpha_{234} \leq \alpha_{134} \leq \alpha_{124} \leq 0$

$$b_a (\alpha_{123} (-b_1 + b_2 + b_3) + \alpha_{234} (b_2 + b_3 - b_4) + \alpha_{134} (b_1 - b_3 + b_4 - 1) + \alpha_{124} (b_1 - b_2 + b_4 - 1) + \alpha_{1234} (b_1 + b_4 - 1)) \quad (8)$$

$$+ \alpha_{123} (-b_2 - b_3 + b_1b_2 + b_1b_3) + \alpha_{234} (-b_2 - b_3 + b_2b_4 + b_3b_4) \quad (9)$$

$$+ \alpha_{134} (1 - b_1 + b_3 - b_4 + b_1b_4) + \alpha_{124} (1 - b_1 + b_2 - b_4 + b_1b_4) + \alpha_{1234} (1 - b_1 - b_4 + b_1b_4) \quad (10)$$

*Case 2.2.1.2:*  $-\alpha_{1234} \leq \alpha_{123} \leq \alpha_{1234}/2 \leq \alpha_{234} \leq \alpha_{134} \leq 0 \leq \alpha_{124}$

*Case 2.2.1.3:*  $-\alpha_{1234} \leq \alpha_{123} \leq \alpha_{1234}/2 \leq \alpha_{234} \leq 0 \leq \alpha_{134} \leq \alpha_{124}$

*Case 2.2.2:*  $\alpha_{1234} + \alpha_{123} + \alpha_{234} \geq 0$

*Case 2.2.2.1:*  $-\alpha_{1234} \leq \alpha_{123} \leq \alpha_{1234}/2 \leq \alpha_{234} \leq \alpha_{134} \leq \alpha_{124} \leq 0$

*Case 2.2.2.2:*  $-\alpha_{1234} \leq \alpha_{123} \leq \alpha_{1234}/2 \leq \alpha_{234} \leq \alpha_{134} \leq 0 \leq \alpha_{124}$

*Case 2.2.2.3:*  $-\alpha_{1234} \leq \alpha_{123} \leq \alpha_{1234}/2 \leq \alpha_{234} \leq 0 \leq \alpha_{134} \leq \alpha_{124}$

*Case 2.3:*  $-\frac{\alpha_{1234}}{2} \leq \alpha_{123}$

*Case 2.3.1:*  $-\alpha_{1234}/2 \leq \alpha_{234} \leq \alpha_{134} \leq \alpha_{124} \leq 0$

Case 2.3.2:  $-\alpha_{1234}/2 \leq \alpha_{234} \leq \alpha_{134} \leq 0 \leq \alpha_{124}$

---

- [1] H. Ishikawa, *IEEE Transactions on Pattern Analysis and Machine Intelligence* **33**, 1234 (2011).
- [2] A. Fix, A. Gruber, E. Boros, and R. Zabih, in *2011 International Conference on Computer Vision* (IEEE, 2011) pp. 1020–1027.
- [3] N. S. Dattani and N. Bryans, (2014), [arXiv:1411.6758](https://arxiv.org/abs/1411.6758).
- [4] C. J. C. Burges, *Microsoft Research MSR-TR-200* (2002).
- [5] X. Peng, Z. Liao, N. Xu, G. Qin, X. Zhou, D. Suter, and J. Du, *Physical Review Letters* **101**, 220405 (2008).
- [6] G. Schaller and R. Schützhold, *Quantum Information & Computation* **10**, 109 (2010).
- [7] N. Xu, J. Zhu, D. Lu, X. Zhou, X. Peng, and J. Du, *Physical Review Letters* **108**, 130501 (2012).
- [8] R. Tanburn, E. Okada, and N. Dattani, *Reducing multi-qubit interactions in adiabatic quantum computation without adding auxiliary qubits. Part 1: The "deduc-reduc" method and its application to quantum factorization of numbers* (2015) [arXiv:1508.04816](https://arxiv.org/abs/1508.04816).
- [9] O. Lunt, R. Tanburn, E. Okada, and N. S. Dattani, *Physical Review A* (in preparation) (2015).
- [10] Z. Li, N. S. Dattani, X. Chen, X. Liu, H. Wang, R. Tanburn, H. Chen, X. Peng, and J. Du, <http://arxiv.org/abs/1706.08061> (2017), [arXiv:1706.08061](https://arxiv.org/abs/1706.08061).
- [11] F. Gaitan and L. Clark, (2012), [arXiv:arXiv:1103.1345v3](https://arxiv.org/abs/1103.1345v3).
- [12] Z. Bian, F. Chudak, W. G. Macready, L. Clark, and F. Gaitan, *Physical Review Letters* **111**, 130505 (2013).
- [13] E. Okada, R. Tanburn, and N. S. Dattani, *Reducing multi-qubit interactions in adiabatic quantum computation without adding auxiliary qubits. Part 2: The "split-reduc" method and its application to quantum determination of Ramsey numbers* (2015) [arXiv:1508.07190](https://arxiv.org/abs/1508.07190).
- [14] C. Rother, V. Kolmogorov, V. Lempitsky, and M. Szummer, in *2007 IEEE Conference on Computer Vision and Pattern Recognition* (IEEE, 2007) pp. 1–8.
- [15] N. Dattani, *Quadratization in discrete optimization and quantum mechanics* (2019) [arXiv:1901.04405](https://arxiv.org/abs/1901.04405).
- [16] H. Ishikawa, in *2014 IEEE Conference on Computer Vision and Pattern Recognition* (IEEE, 2014) pp. 1362–1369.
- [17] R. Dridi and H. Alghassi, *Scientific Reports* **7**, 43048 (2017).
- [18] I. G. Rosenberg, *Cahiers du Centre d'Études de Recherche Operationnelle* **17**, 71 (1975).
- [19] V. Kolmogorov and R. Zabih, *IEEE Transactions on Pattern Analysis and Machine Intelligence* **26**, 147 (2004).
- [20] D. Freedman and P. Drineas, in *2005 IEEE Computer Society Conference on Computer Vision and Pattern Recognition (CVPR'05)*, Vol. 2 (IEEE, 2005) pp. 939–946.
- [21] A. C. Gallagher, D. Batra, and D. Parikh, in *CVPR 2011* (IEEE, 2011) pp. 1857–1864.
- [22] M. Anthony, E. Boros, Y. Crama, and A. Gruber, (2014), [arXiv:1404.6535](https://arxiv.org/abs/1404.6535).
- [23] M. Anthony, E. Boros, Y. Crama, and A. Gruber, (2015).
- [24] M. Anthony, E. Boros, Y. Crama, and A. Gruber, *Discrete Applied Mathematics* **203**, 1 (2016).
- [25] M. Leib, P. Zoller, and W. Lechner, *Quantum Science and Technology* **1**, 15008 (2016).
- [26] A. Rocchetto, S. C. Benjamin, and Y. Li, *Science Advances* **2** (2016), 10.1126/sciadv.1601246.
- [27] M. Anthony, E. Boros, Y. Crama, and A. Gruber, *Mathematical Programming* **162**, 115 (2017).
- [28] N. Chancellor, S. Zohren, and P. A. Warburton, *npj Quantum Information* **3**, 21 (2017).
- [29] E. Boros, Y. Crama, and E. Rodríguez-Heck, *Quadratizations of symmetric pseudo-Boolean functions: sub-linear bounds on the number of auxiliary variables*, Tech. Rep. (2018).
- [30] E. Boros, Y. Crama, and E. Rodríguez Heck, *Compact quadratizations for pseudo-Boolean functions* (2018).