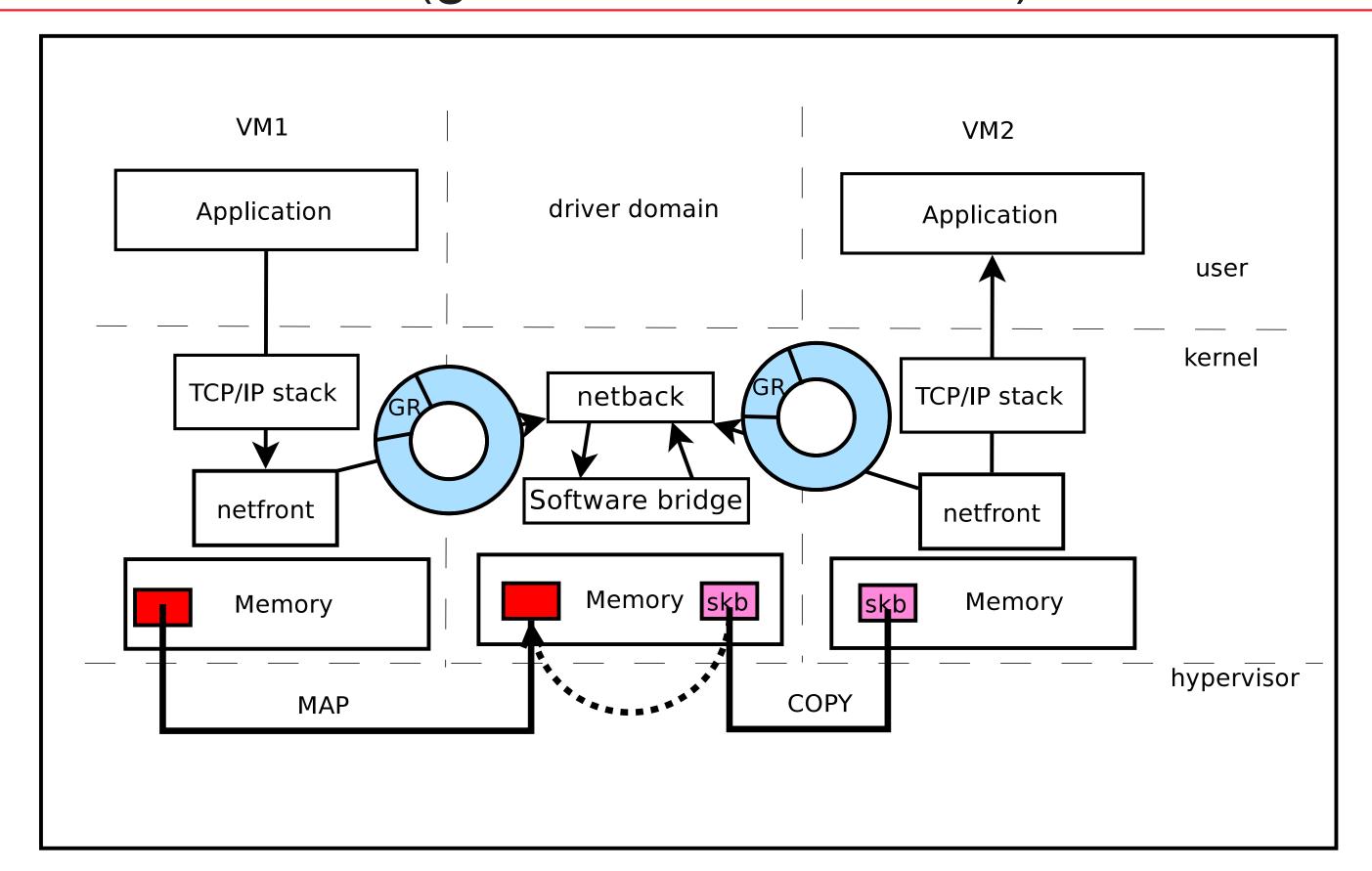


V4VSOCKETS

LOW OVERHEAD INTRA-NODE COMMUNICATION IN XEN



I/O Path in Xen (generic environment)



Intra-node communication suffers from severe overheads:

- ⇒ inefficient data paths
- ⇒ driver domain handles packet forwarding
- ⇒ unnecessary TCP/IP stack crossing and fragmentation

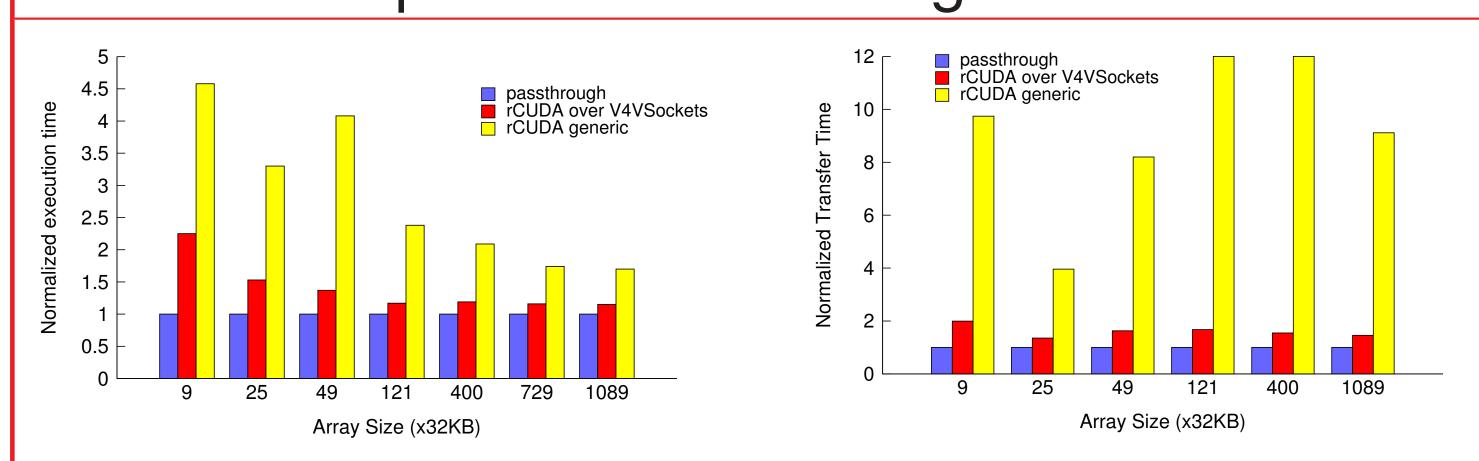
Key features

V4VSockets is an efficient, socket-compliant, high performance intra node communication framework.

V4VSockets features:

- ⇒ optimized data path (data are copied from / to the VM kernel memory without the need to share pages between VMs)
- ⇒ low-overhead framework (no intermediary VM driver domain no scheduling implications are involved)
- ⇒ secure (no security implications data cross the hypervisor and either get dropped or pushed forward using V4V semantics)
- ⇒ ultra low latency and high bandwidth
- ⇒ Open-source, available @ https://github.com/HPSI/v4v

GPU stencil performance through rCUDA



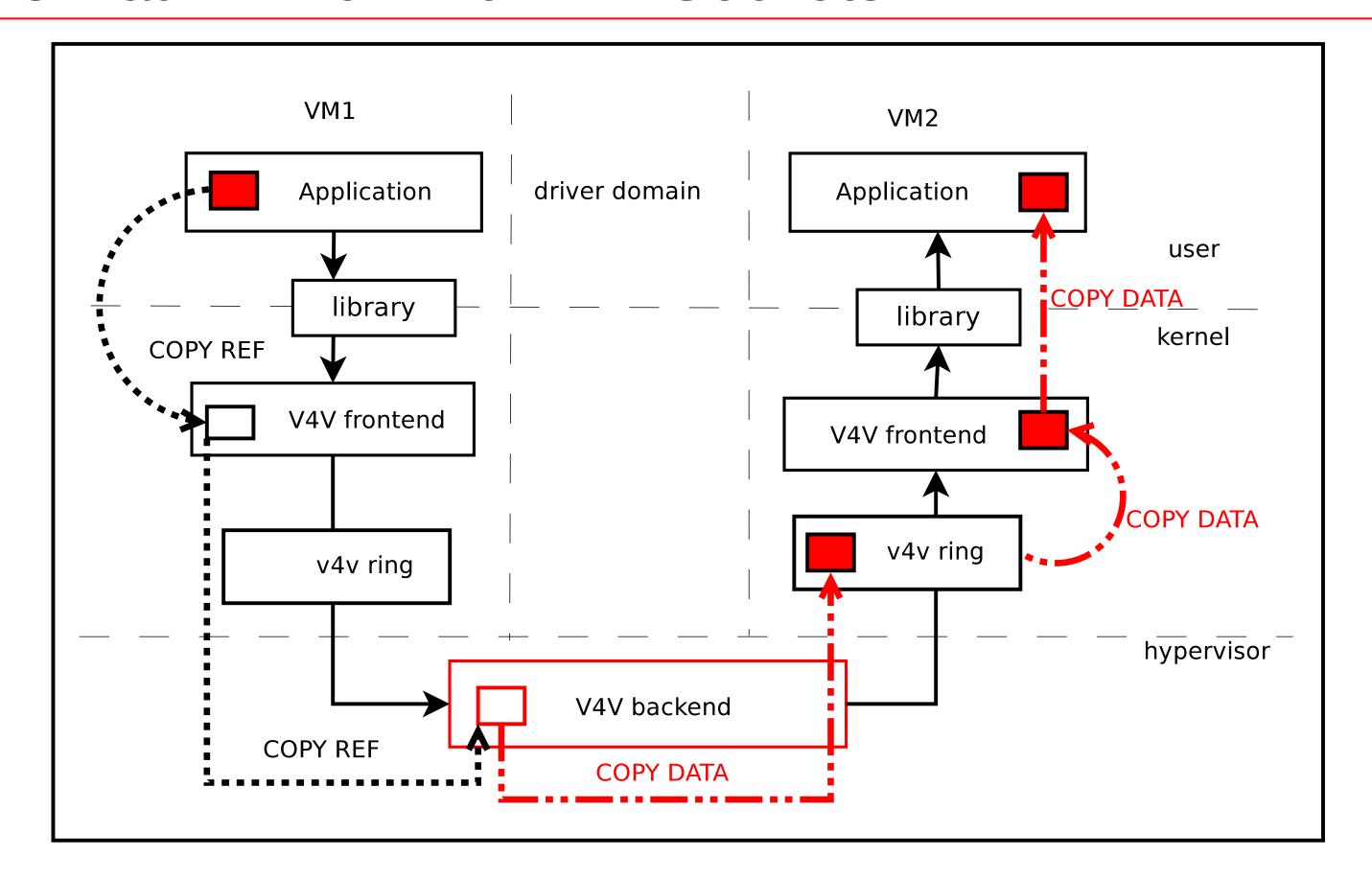
We run the matrix-matrix product benchmark: (a) natively via GPU passthrough, (b) via rCUDA over TCP/IP sockets and (c) via rCUDA over V4VSockets. Steps include: (i) 2x input matrix copies, (ii) GPU product compute, (iii) 1x output matrix copy back. V4VSockets:

- ⇒ adds a minimum overhead of 15% (compared to native execution)
- \Rightarrow boosts transfer throughput by 7x (at best) compared to TCP/IP

Work in Progress

- ⇒ Strengthen our implementation towards a more user-friendly approach.
- ⇒ Thoroughly examine the CPU utilization overheads imposed by V4VSockets.
- ⇒ Polish the peer discovery framework to adaptively use V4VSockets over generic sockets.
- ⇒ Perform an elaborate performance evaluation of the GPU sharing framework we have developed.

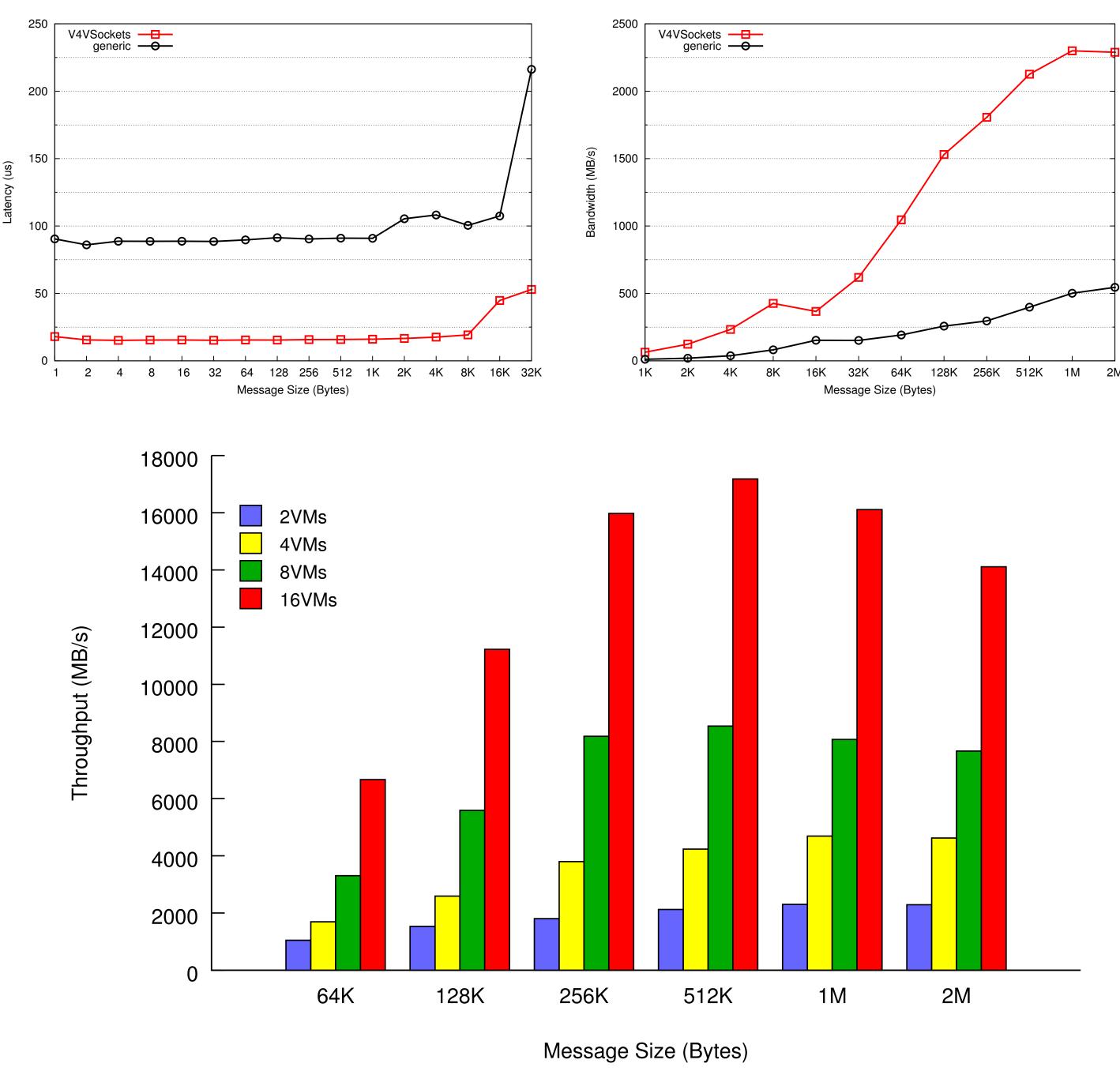
I/O Path in Xen with V4VSockets



V4VSockets is built as a full-stack protocol framework that supports p2p communication between VMs.

- ⇒ Application layer: the socket interface.
- ⇒ *Transport layer*: VM kernel driver.
- ⇒ *Network/Link layer*: the hypervisor, providing encapsulation of upper-layer messages to V4V messages, and packet delivery.

Ping-pong benchmark on 2x Xeon X5650, 48 GB



V4VSockets:

- ⇒ improves latency for small messages by 81%
- ⇒ achieves 2299 MB/s for large messages (1 MB) vs. 501 MB/s
- ⇒ scales efficiently with the number of VMs, both in terms of latency and bandwidth
- \Rightarrow achieves aggregate throughput \approx 17 GB/s for 512 KB messages when 16 VMs exchange data in pairs

Contact info and Acknowledgments

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