# V4VSockets: low-overhead intra-node communication in Xen.

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#### Intro

### Cloud computing

- application oriented
- fast, ease-of-use

#### Consolidation

- $\frac{vCPU}{physical cores} >> 1$
- multi/many-cores

The number of co-located VMs is drastically increasing



#### Introduction

### Applications

- stand-alone (flexibility)
- distributed (elasticity)

When it comes to down to numbers, the need for intra-node communication in small data center increases.

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#### Contribution

Design and Implement V4VSockets:

- efficient message exchange (one order of magnitude)
- isolation
- API compatible (Sockets)



#### Overview

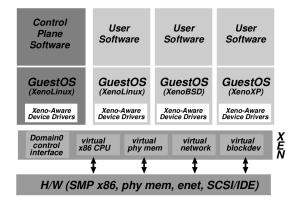
- Communication in Virtual Environments
  - Basic Concepts
- 2 I/O access in Virtualized Environments
  - Xen
  - Intra-node communication in Xen

- 3 V4VSockets
  - Architecture
  - Experimental Evaluation



#### Xen - Architecture

• hypervisor & privileged VM (driver domain) to access hardware





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# I/O Internals – Xen

#### Xen basics

- hypervisor driver domain runs as a Linux guest
- split driver model (frontend/backend)

#### Xen – Event channels

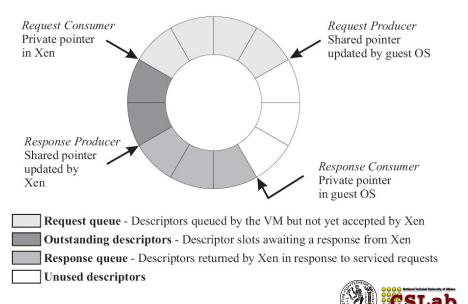
- notify Guest/Host about a pending transaction
- easy to setup bind to a specific "port"

#### Xen – Grant mechanism

- issue a page grant request
- the other end maps the grant (accept)
- this page is shared across the two domains



### Xen Ring buffers



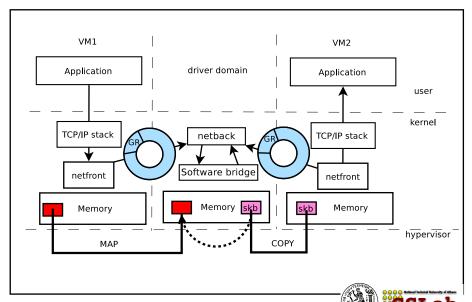
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### Intra-node communication in Xen

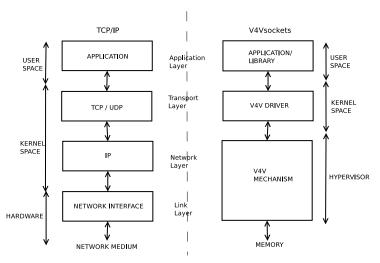


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### Application/Library layer

• forwards the relevant actions and arguments to the transport layer.

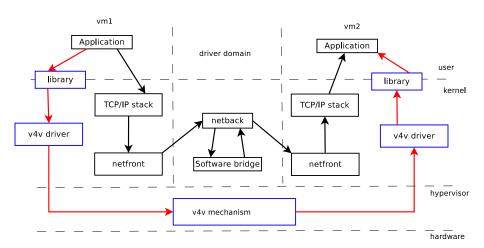
### Transport layer – V4V kernel driver

- handles the virtual connection semantics between peer VMs that need to communicate,
- is in charge of fragmenting and sending upper-layer packets by issuing hypercalls to the hypervisor (network layer), and
- provides a notification mechanism to the VM's user-space for receiving packets, as well as error control.

#### Network/Link layer - Hypervisor

- encapsulation of upper-layer messages to packets that will be transmitted to their destination, according to V4V semantics,
- packet delivery.

# V4Vsockets vs Netfront/Netback





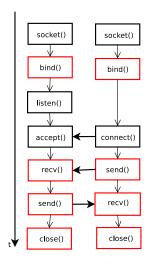
### V4VSockets

#### Communication mechanisms in V4VSockets

- system calls
- hypercalls
- event channels VIRQS

### V4V Rings

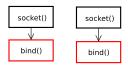
- circular buffers
- allocated in the VM address space by the VM kernel
- data exchange medium

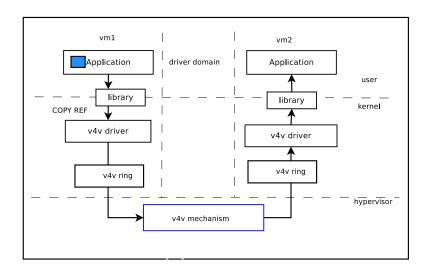


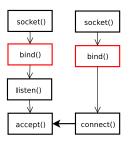


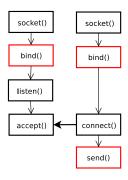
socket()

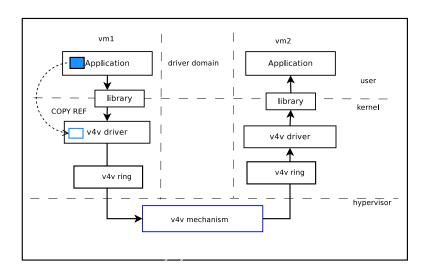
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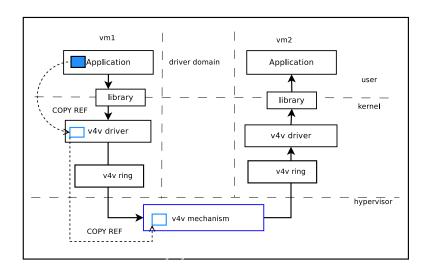


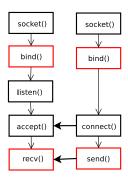


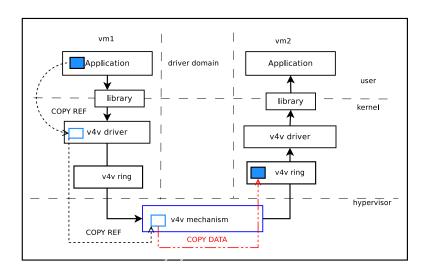


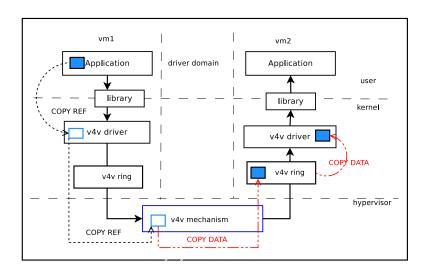


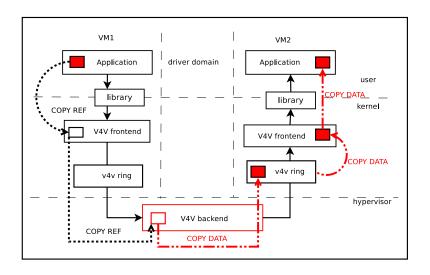




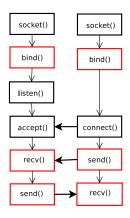


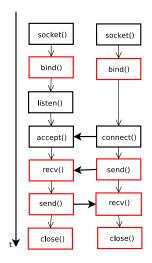












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### Performance Results

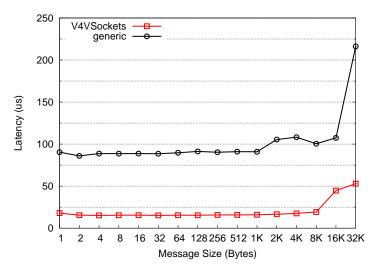
#### Testbed

- $\bullet$  2x {Intel Xeon @2.4Ghz}, Intel 5520, 48GB memory, Generic 10GbE
- Xen 4.5-unstable, Debian GNU/Linux (Linux kernel 3.14.2)
- generic micro-benchmark: pingpong

#### Cases:

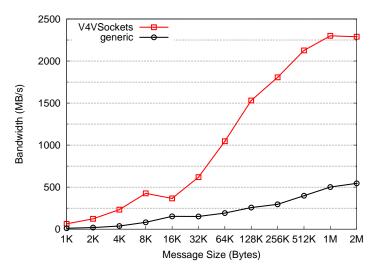
- TCP sockets
- V4V Stream

### Experimental evaluation – latency



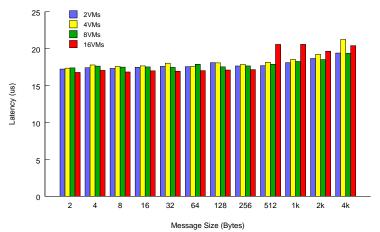


### Experimental evaluation – Bandwidth



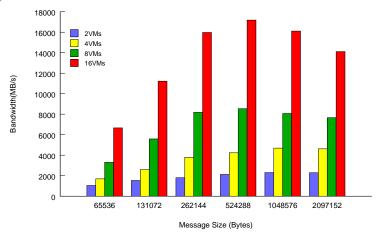


# Experimental evaluation – Scaling factor (up to 16 VMs)

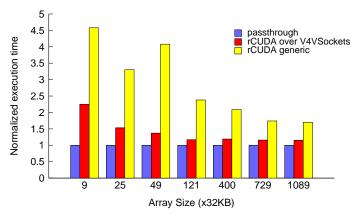




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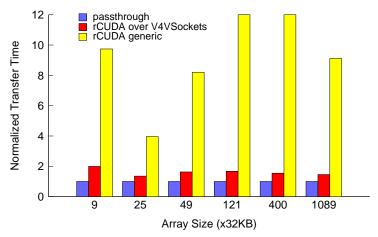


### Experimental evaluation – GPU stencil





### Experimental evaluation – GPU stencil



#### Summary

#### Intra-node communication in VM environments

- split driver model generic
- V4VSockets: framework for low-overhead intra-node communication
  - is not based on a driver domain
  - does not use shared memory between guests (map/grant mechanism)
  - uses memory copies, hypercalls and event channels
- better throughput (efficient data path, bypass the complex TCP/IP stack)
- hypercall overheads (small, negligible if correctly finetuned)
- scalability (no privileged guest involved in communication)
- isolation (no shared memory)
- extensible



#### Future endaevors

- cpu utilization overheads
- NUMA and multihierarchical memory architectures
- map instead of copy (study the systems behavior of providing a shared memory space between VMs)
- Optimize away copies on the data path.

Available online as open-source

https://github.com/HPSI/V4VSockets

Thanks!

Questions?

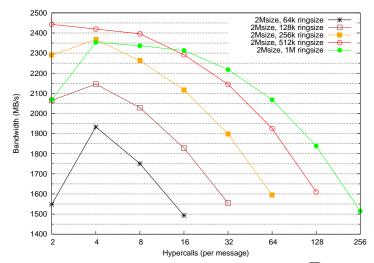




# Backup

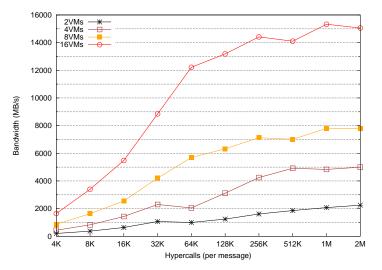


# Experimental evaluation – 2M message size vs Hypercalls





## Experimental evaluation – Datagram scalability





## Experimental evaluation – Hypercalls per message

