

Introduction to Theoretical Computer Science, Fall 2024

Quiz 2 Solutions

Q1. The following CFG generates A .

$$S \rightarrow AB$$

$$B \rightarrow B0|B1|e$$

$$A \rightarrow 0A0|1A1|\#B$$

Q2. We give a reduction from ALL_{PDA} to EQ_{PDA} . Let P be a PDA. Let P^* be a PDA with $L(P^*) = \Sigma^*$. P^* can be easily constructed. Then we have that " P " $\in \text{ALL}_{\text{PDA}}$ if and only if " P " " P^* " $\in \text{EQ}_{\text{PDA}}$. This completes the reduction.

Q3. Suppose that A is recursively enumerable. Let M_A be the TM that semidecides A . We shall use M_A to construct a TM M_H that halts on H_{TM} . Let x be an arbitrary string. Note that either $x \in H_{\text{TM}}$ or $x \in \overline{H_{\text{TM}}}$. Or equivalently, either $0x \in A$ or $1x \in A$. Therefore, M_A halts on exactly one of $0x$ and $1x$. We can run M_A on $0x$ and $1x$ in parallel, and see on which input it halts. If M_A halts on $0x$, then $x \in H_{\text{TM}}$. If M_A halts on $1x$, then $x \notin H_{\text{TM}}$.