Problem Set #2

Quiz, 5 questions

4/5 points (80%)



# **Congratulations! You passed!**

Next Item



1/1 point

1.

This question will give you further practice with the Master Method. Suppose the running time of an algorithm is governed by the recurrence  $T(n)=7*T(n/3)+n^2$ . What's the overall asymptotic running time (i.e., the value of T(n))?

- $\theta(n \log n)$
- $\theta(n^{2.81})$
- $\theta(n^2 \log n)$
- $\theta(n^2)$

### Correct

a=7, b=3, d=2. Since  $b^d > a$ , this is case 2 of the Master Method.



1/1 point

2

This question will give you further practice with the Master Method. Suppose the running time of an algorithm is governed by the recurrence  $T(n)=9*T(n/3)+n^2$ . What's the overall asymptotic running time (i.e., the value of T(n))?

- $\theta(n^2)$
- $\theta(n \log n)$
- $\theta(n^{3.17})$

### Correct

 $a = b^d = 9$ , so this is case 1 of the Master Method.

## Problem Set #2

4/5 points (80%)

Quiz, 5 questions



1/1 point

3

This question will give you further practice with the Master Method. Suppose the running time of an algorithm is governed by the recurrence T(n)=5\*T(n/3)+4n. What's the overall asymptotic running time (i.e., the value of T(n))?

- $\theta(n^{2.59})$
- hinspace hin
- $\theta(n^2)$
- $\theta(n\log(n))$
- $\theta(n^{\frac{\log 3}{\log 5}})$
- $igg( n^{\log_3(5)} ig)$

#### Correct

a = 5, b = 3, d = 1. Since  $a > b^d$ , this is case 3 of the Master Method.



0/1 point

4.

Consider the following pseudocode for calculating  $a^b$  (where a and b are positive integers)

```
FastPower(a,b) :
      if b = 1
3
        return a
4
      else
 5
        c := a*a
        ans := FastPower(c,[b/2])
 6
7
      if b is odd
8
        return a*ans
9
      else return ans
10
   end
```

Here [x] denotes the floor function, that is, the largest integer less than or equal to x.

Now assuming that you use a calculator that supports multiplication and division (i.e., you can do multiplications and divisions in constant time), what would be the overall asymptotic running time of the above algorithm (as a function of b)?

4/5 points (80%)

- $\Theta(b)$
- $\Theta(b\log(b))$

### This should not be selected

This would be even worse than the naive method of multiplying a by itself b times...



1/1 point

Choose the smallest correct upper bound on the solution to the following recurrence: T(1)=1 and  $T(n) \leq T(\lceil \sqrt{n} \rceil) + 1$  for n > 1. Here [x] denotes the "floor" function, which rounds down to the nearest integer. (Note that the Master Method does not apply.)

- $O(\sqrt{n})$
- $O(\log n)$
- O(1)
- $O(\log \log n)$

### Correct

Bingo! This answer may be easiest to see by writing n as  $2^{\log n}$  and then noting that every squareroot operation cuts the exponent in half.