

## CS361 Questions: Week 4

### Lecture 53

1. Why is it important for a digital signature to be non reusable?
  - a. So the signature cannot be detached and reused for a different message.
2. Why is it the hash of the message typically signed, rather than the message itself?
  - a. Public key encryption is expensive – hashes are cheaper. The hash is signed for authentication reasons.
3. What assurance does R gain from the interchange on slide 4?
  - a. R knows that the message was sent from S.

### Lecture 54

1. What is the importance of certificate authorities?
  - a. A certificate is a letter of introduction – vouching for the accuracy of the binding.
2. In the example on slide 5, why does X sign the hash of the first message with its private key?
  - a. So that Y knows it is really X.
3. Why is it necessary to have a hash of Y and  $K_y$ ?
  - a. The certifying authority needs both to certify the binding.
4. What would happen if Z had a public key for X, but it was not trustworthy?
  - a. It wouldn't have a hash of X.

### Lecture 55

1. What happens at the root of a chain of trust?
  - a. It should be an unimpeachable authority that holds all chains.
2. Why does an X.509 certificate include a “validity interval”?
  - a. So the user can check the times.
3. What would it mean if the hash and the received value did not match?
  - a. Something or someone has appended or removed something.

### Lecture 56

1. What are some protocols previously discussed?
  - a. HTTP
2. What may happen if one step of a protocol is ignored?
  - a. It doesn't work
3. Why must the ciphers commute in order to accomplish the task in slide 4?
  - a. They both need the same key.
4. Describe how an attacker can extract M from the protocol in slide 6.
  - a. An eavesdropper who stores the three messages can XOR combinations of them to extract any of M,  $K_a$ , and  $K_b$ .

5. Describe how an attacker can extract  $K_a$  from the protocol in slide 6.
  - a. An eavesdropper who stores the three messages can XOR combinations of them to extract any of  $M$ ,  $K_a$ , and  $K_b$ .
6. Describe how an attacker can extract  $K_b$  from the protocol in slide 6.
  - a. An eavesdropper who stores the three messages can XOR combinations of them to extract any of  $M$ ,  $K_a$ , and  $K_b$ .
7. Why are cryptographic protocols difficult to design and easy to get wrong?
  - a. Because you just need one easy point of entry or a few pieces of data to get a key or value.

## Lecture 57

1. Explain the importance of protocols in the context of the internet.
  - a. An internet protocol is a structured dialogue among two or more parties (host to router, host to host, router to router...) from anywhere in the world that can connect to a connected subnetwork.
2. Explain the importance of cryptographic protocols in the context of the internet.
  - a. Cryptographic protocols in the context of the internet allows the secure transfer of data – so that no one can change or eavesdrop a message.
3. What are the assumptions of the protocol in slide 6?
  - a. A can open the message from B by using their secret key and B can open the message by using B's secret key. The assumption is that A can verify it is B and B can verify it is A.
4. What are the goals of the protocol in slide 6?
  - a. Verification
5. Are the goals of the protocol in slide 6 satisfied? Explain.
  - a. No, they are not because  $K_b$  and  $K_a$  are public and any one can use them.
6. How is the protocol in slide 6 flawed?
  - a.  $K_b$  and  $K_a$  are public and any one can use them.

## Lecture 58

1. Why is it important to know if a protocol includes unnecessary steps or messages?
  - a. To know if data is repeated.
2. Why is it important to know if a protocol encrypts items that could be sent in the clear?
  - a. To know the sensitivity of the information.

## Lecture 59

1. Why might it be difficult to answer what constitutes an attack on a cryptographic protocol?
  - a. Attack subtle, there are many different types of attacks, and there can be holes in protocols people aren't aware of.
2. Describe potential dangers of a replay attack.
  - a. There is double data - this can reset or increment values.
3. Are there attacks where an attacker gains no secret information? Explain.
  - a. Yes, attackers can add information without reading.
4. What restrictions are imposed on the attacker?

- a. They have to follow some part of the protocol.
- 5. Why is it important that protocols are asynchronous?
  - a. So that the attacker can't synchronize an attack.

## Lecture 60

1. Would the Needham-Schroeder protocol work without nonces?
  - a. It would work but it wouldn't be secure.
2. For each step of the NS protocol, answer the two questions on slide 5.
  - a. What is the sender trying to say with this message?
    - i. Sends  $N_a$
    - ii.  $K_{ab}$
    - iii.  $K_{ab}$
    - iv.  $N_b$
    - v.  $N_b - 1$
  - b. What is the receiver entitled to believe after receiving the message?
    - i. A's and B's messages are fresh/ $N_b$  is a new nonce.

## Lecture 61

1. As in slide 5, if A's key were later changed, after having  $K_{as}$  compromised, how could A still be impersonated?
  - a. A impersonator requests a new nonce with B.
2. Is it fair to ask the question of a key being broken?
  - a. Yes
3. How might you address these flaws if you were the protocol designer?
  - a. Question the key everytime.

## Lecture 62

1. What guarantees does Otway-Rees seem to provide to A and B?
  - a.
2. Are there guarantees that Needham-Schroeder provides that Otway-Rees does not or vice versa?
3. How could you fix the flawed protocol from slide 4?

## Lecture 63

1. Why is the verification of protocols important?
  - a. Protocols are crucial to the Internet – we should strive to get them right.
2. What is a belief logic?
  - a. Belief logic allows reasoning about what principals within the protocol should be able to infer from the messages they see.
3. A protocol is a program; where do you think beliefs come in?
  - a. Initial assumptions.

## Lecture 64

1. What is a modal logic?
  - a. A logic that express belief of the state/qualify a statement.
2. Explain the intuition behind the message meaning inference rule.
  - a. A is an authority on X and can be trusted on X. If A trusts X, X can be trusted.
3. Explain the intuition behind the nonce verification inference rule.
  - a. If A believes X is fresh and A believes B once said X, then A believes B believes X.
4. Explain the intuition behind the jurisdiction inference rule.
  - a. If A believes B has jurisdiction over X and A believes B believes X, then A believes X.
5. What is idealization and why is it needed?
  - a. Idealization attempts to turn the message sent into its intended semantics. It is needed to omit parts of the message that do not contribute to the beliefs of the recipients.

## Lecture 65

1. Why do you think plaintext is omitted in a BAN idealization?
  - a. It does not contribute to the beliefs of the recipients.
2. Some idealized steps seem to refer to beliefs that will happen later in the protocol. Why would that be?
  - a. It means that if we do this the later state should be true.
3. One benefit of a BAN proof is that it exposes assumptions. Explain that.
  - a. In order to get to state x, one must start from state y.