



THE UNIVERSITY OF TEXAS  
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CS383C NUMERICAL ANALYSIS  
**HW05 Numerical Stability**

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### Exercise 3.

#### 3.1

Write 1 as floating number.

$$.1 \underbrace{00 \cdots 0}_{t-1} \times 2^1 \quad (1)$$

#### 3.2

Show that  $\mathbf{u} = \frac{1}{2} \cdot 2^{1-t}$

*Proof.* Let  $\chi = .\delta_0\delta_1 \cdots \delta_{t-1}\delta_t \cdots$  and  $\check{\chi}$  to be the value stored in t-digit floating number with rounding mechanism. Then if  $\delta_t = 0$ , then  $\chi = \check{\chi}$  and  $|\delta\chi| = 0 \leq 2^{e-t-1}$ . But if  $\delta_t = 1$ , due to the rounding mechanism, then  $\chi < \check{\chi}$  and

$$|\delta\chi| = |\chi - \check{\chi}| = |.\delta_0\delta_1 \cdots \delta_{t-1}\delta_t \cdots \times 2^e - .\delta_0\delta_1 \cdots \delta'_{t-1} \times 2^e| \leq .\underbrace{00 \cdots 0}_t 1 \times 2^e = 2^{e-t-1} \quad (2)$$

For  $\chi$ , since  $\delta_0 = 1$  (normalized)

$$|\chi| = |.\delta_0\delta_1 \cdots \times 2^e| \geq .1 \times 2^e \geq 2^{e-1} \quad (3)$$

Thus,

$$\frac{|\delta\chi|}{|\chi|} \leq \frac{2^{e-t-1}}{2^{e-1}} = \frac{1}{2} \cdot 2^{1-t} \quad (4)$$

Then,

$$|\delta\chi| \leq \frac{1}{2} \cdot 2^{1-t} |\chi| \quad (5)$$

Now, we have

$$\mathbf{u} = \frac{1}{2} 2^{1-t} \quad (6)$$

□

### Exercise 10.

Show that  $|AB| \leq |A||B|$ .

*Proof.* Let  $C = AB$ . And the  $(i, j)$  entry of  $|C|$  is given by

$$|c_{i,j}| = \left| \sum_{p=0}^k a_{i,p} b_{p,j} \right| \leq \sum_{p=0}^k |a_{i,p} b_{p,j}| \leq \sum_{p=0}^k |a_{i,p}| |b_{p,j}| \quad (7)$$

which equals  $(i, j)$  entry of  $|A||B|$ . Hence, we have

$$|AB| \leq |A||B| \quad (8)$$

□

## Exercise 12.

### 12.1

Show that if  $|A| \leq |B|$ , then  $\|A\|_1 \leq \|B\|_1$ .

*Proof.* Partition  $A_{m \times n} = \left( \begin{array}{c|c|c|c} a_0 & a_1 & \dots & a_{n-1} \end{array} \right)$  where  $a_j$  indicates the  $j$ -th column of matrix  $A$ . Similarly, we partition  $B_{m \times n} = \left( \begin{array}{c|c|c|c} b_0 & b_1 & \dots & b_{n-1} \end{array} \right)$  where  $b_j$  indicates the  $j$ -th column of matrix  $B$ . Then we use  $a_{ij}$  and  $b_{ij}$  to denote  $i$ -th element of  $a_j$  and  $b_j$  respectively.

$$\|A\|_1 = \max_{0 \leq j < n} \|a_j\|_1 = \max_{0 \leq j < n} \sum_{i=0}^{m-1} |a_{ij}| \leq \max_{0 \leq j < n} \sum_{i=0}^{m-1} |b_{ij}| = \max_{0 \leq j < n} \|b_j\|_1 = \|B\|_1 \quad (9)$$

□

**Lemma 1.** For arbitrary matrix  $A = \left( \begin{array}{c|c|c|c} a_0 & a_1 & \dots & a_{n-1} \end{array} \right)$ ,  $\|A\|_1 = \max_{0 \leq j < n} \|a_j\|_1$ .

*Proof.* This lemma has been proved in Notes on Norms. □

### 12.2

Show that if  $|A| \leq |B|$ , then  $\|A\|_\infty \leq \|B\|_\infty$ .

*Proof.* Partition  $A_{m \times n} = \left( \begin{array}{c} \frac{a_0}{a_1} \\ \vdots \\ a_{m-1} \end{array} \right)$ , where  $a_i$  indicates the  $i$ -th row of matrix  $A$ . Similarly, we

partition  $B_{m \times n} = \left( \begin{array}{c} \frac{b_0}{b_1} \\ \vdots \\ b_{m-1} \end{array} \right)$ , where  $b_i$  indicates the  $i$ -th row of matrix  $B$ . Then we use  $a_{ij}$  and  $b_{ij}$

to denote  $j$ -th element of  $a_i$  and  $b_i$  respectively.

$$\|A\|_\infty = \max_{0 \leq i < m} \|a_i\|_1 = \max_{0 \leq i < m} \sum_{j=0}^{n-1} |a_{ij}| \leq \max_{0 \leq i < m} \sum_{j=0}^{n-1} |b_{ij}| = \max_{0 \leq i < m} \|b_i\|_1 = \|B\|_\infty \quad (10)$$

□

**Lemma 2.** For arbitrary matrix  $A = \left( \begin{array}{c} \frac{a_0}{a_1} \\ \vdots \\ a_{m-1} \end{array} \right)$ ,  $\|A\|_\infty = \max_{0 \leq i < m} \|a_i\|_1$ .

*Proof.* This lemma has been proved in Notes on Norms. □

### 12.3

Show that if  $|A| \leq |B|$ , then  $\|A\|_F \leq \|B\|_F$ . Let  $A, B \in \mathbb{R}^{m \times n}$  and  $a_{ij}$ ,  $b_{ij}$  be  $(i, j)$  entry of  $A$ ,  $B$  respectively.

*Proof.*

$$\|A\|_F^2 = \sum_{i=0}^{m-1} \sum_{j=0}^{n-1} |a_{ij}|^2 \leq \sum_{i=0}^{m-1} \sum_{j=0}^{n-1} |b_{ij}|^2 = \|B\|_F^2 \quad (11)$$

□

**Exercise 13.**

$$\check{\kappa} = [(\chi_0\psi_0 + \chi_1\psi_1) + \chi_2\psi_2] \quad (12)$$

$$= [[\chi_0\psi_0 + \chi_1\psi_1] + [\chi_2\psi_2]] \quad (13)$$

$$= [[[ \chi_0\psi_0 ] + [ \chi_1\psi_1 ] ] + [ \chi_2\psi_2 ]] \quad (14)$$

$$= [[\chi_0\psi_0(1 + \epsilon_*^{(0)}) + \chi_1\psi_1(1 + \epsilon_*^{(1)})] + \chi_2\psi_2(1 + \epsilon_*^{(2)})] \quad (15)$$

$$= [(\chi_0\psi_0(1 + \epsilon_*^{(0)}) + \chi_1\psi_1(1 + \epsilon_*^{(1)}))(1 + \epsilon_+^{(1)}) + \chi_2\psi_2(1 + \epsilon_*^{(2)})] \quad (16)$$

$$= \left( (\chi_0\psi_0(1 + \epsilon_*^{(0)}) + \chi_1\psi_1(1 + \epsilon_*^{(1)}))(1 + \epsilon_+^{(1)}) + \chi_2\psi_2(1 + \epsilon_*^{(2)}) \right) (1 + \epsilon_+^{(2)}) \quad (17)$$

$$= \chi_0\psi_0(1 + \epsilon_*^{(0)})(1 + \epsilon_+^{(1)})(1 + \epsilon_+^{(2)}) + \chi_1\psi_1(1 + \epsilon_*^{(1)})(1 + \epsilon_+^{(1)})(1 + \epsilon_+^{(2)}) + \chi_2\psi_2(1 + \epsilon_*^{(2)})(1 + \epsilon_+^{(2)}) \quad (18)$$

$$= \begin{pmatrix} \chi_0 \\ \chi_1 \\ \chi_2 \end{pmatrix}^T \begin{pmatrix} \epsilon_0(1 + \epsilon_*^{(0)})(1 + \epsilon_+^{(1)})(1 + \epsilon_+^{(2)}) \\ \epsilon_1(1 + \epsilon_*^{(1)})(1 + \epsilon_+^{(1)})(1 + \epsilon_+^{(2)}) \\ \epsilon_2(1 + \epsilon_*^{(2)})(1 + \epsilon_+^{(2)}) \end{pmatrix} \quad (19)$$

$$= \begin{pmatrix} \chi_0 \\ \chi_1 \\ \chi_2 \end{pmatrix}^T \begin{pmatrix} (1 + \epsilon_*^{(0)})(1 + \epsilon_+^{(1)})(1 + \epsilon_+^{(2)}) & 0 & 0 \\ 0 & (1 + \epsilon_*^{(1)})(1 + \epsilon_+^{(1)})(1 + \epsilon_+^{(2)}) & 0 \\ 0 & 0 & (1 + \epsilon_*^{(2)})(1 + \epsilon_+^{(2)}) \end{pmatrix} \begin{pmatrix} \epsilon_0 \\ \epsilon_1 \\ \epsilon_2 \end{pmatrix} \quad (20)$$

$$= \begin{pmatrix} \chi_0(1 + \epsilon_*^{(0)})(1 + \epsilon_+^{(1)})(1 + \epsilon_+^{(2)}) \\ \chi_1(1 + \epsilon_*^{(1)})(1 + \epsilon_+^{(1)})(1 + \epsilon_+^{(2)}) \\ \chi_2(1 + \epsilon_*^{(2)})(1 + \epsilon_+^{(2)}) \end{pmatrix}^T \begin{pmatrix} \epsilon_0 \\ \epsilon_1 \\ \epsilon_2 \end{pmatrix} \quad (21)$$

**Exercise 15.**

Now we complete the missing part for the Inductive Step Case 1 of Lemma 14.

*Proof.* Case 1:  $\prod_{i=0}^n (1 + \epsilon)^{\pm 1} = \prod_{i=0}^{n-1} (1 + \epsilon)^{\pm 1} (1 + \epsilon_n)$ .

By the inductive hypothesis, there exists a  $\theta_n$  such that

$$(1 + \theta_n) = \prod_{i=0}^{n-1} (1 + \epsilon_i)^{\pm 1} \text{ and } |\theta_n| \leq n\mathbf{u}/(1 - n\mathbf{u}) \quad (22)$$

Then

$$\prod_{i=0}^n (1 + \epsilon)^{\pm 1} = \left( \prod_{i=0}^{n-1} (1 + \epsilon)^{\pm 1} \right) (1 + \epsilon_n) = (1 + \theta_n)(1 + \epsilon_n) = 1 + \underbrace{\theta_n + \epsilon_n + \theta_n \cdot \epsilon_n}_{\theta_{n+1}} \quad (23)$$

which tells us how to pick up  $\theta_{n+1}$ . Then

$$|\theta_{n+1}| = |\theta_n + \epsilon_n + \theta_n \cdot \epsilon_n| \quad (24)$$

$$\leq |\theta_n| + |\epsilon_n| + |\theta_n| \cdot |\epsilon_n| \quad (25)$$

$$= \frac{n\mathbf{u}}{1 - n\mathbf{u}} + \mathbf{u} + \frac{n\mathbf{u}}{1 - n\mathbf{u}} \cdot \mathbf{u} \quad (26)$$

$$= \frac{n\mathbf{u} + \mathbf{u} - n\mathbf{u}^2 + n\mathbf{u}^2}{1 - n\mathbf{u}} \quad (27)$$

$$= \frac{(n+1)\mathbf{u}}{1 - n\mathbf{u}} \quad (28)$$

$$\leq \frac{(n+1)\mathbf{u}}{1 - (n+1)\mathbf{u}} \quad (29)$$

□

## Exercise 18.

### 18.1

Show that if  $n, b \geq 1$ , then  $\gamma_n \leq \gamma_{n+b}$ .

*Proof.*

$$\gamma_n = \frac{n\mathbf{u}}{1 - n\mathbf{u}} \leq \frac{n\mathbf{u}}{1 - (n+b)\mathbf{u}} \leq \frac{(n+b)\mathbf{u}}{1 - (n+b)\mathbf{u}} = \gamma_{n+b} \quad (30)$$

Note that since  $\mathbf{u}$  is extremely small, then  $1 - n\mathbf{u} > 0$  and  $1 - (n+b)\mathbf{u} > 0$ .  $\square$

### 18.2

Show that if  $n, b \geq 1$ , then  $\gamma_n + \gamma_b + \gamma_n\gamma_b \leq \gamma_{n+b}$ .

*Proof.*

$$\gamma_n + \gamma_b + \gamma_n\gamma_b \quad (31)$$

$$= \frac{n\mathbf{u}}{1 - n\mathbf{u}} + \frac{b\mathbf{u}}{1 - b\mathbf{u}} + \frac{n\mathbf{u}}{1 - n\mathbf{u}} \cdot \frac{b\mathbf{u}}{1 - b\mathbf{u}} \quad (32)$$

$$= \frac{n\mathbf{u} - nb\mathbf{u}^2 + b\mathbf{u} - nb\mathbf{u}^2 + nb\mathbf{u}^2}{(1 - n\mathbf{u})(1 - b\mathbf{u})} \quad (33)$$

$$= \frac{n\mathbf{u} - nb\mathbf{u}^2 + b\mathbf{u}}{1 - n\mathbf{u} - b\mathbf{u} + nb\mathbf{u}^2} \quad (34)$$

$$\leq \frac{n\mathbf{u} + b\mathbf{u}}{1 - n\mathbf{u} - b\mathbf{u} + nb\mathbf{u}^2} \quad (35)$$

$$\leq \frac{n\mathbf{u} + b\mathbf{u}}{1 - n\mathbf{u} - b\mathbf{u}} \quad (36)$$

$$= \frac{(n+b)\mathbf{u}}{1 - (n+b)\mathbf{u}} \quad (37)$$

$$= \gamma_{n+b} \quad (38)$$

Note that since  $\mathbf{u}$  is extremely small, then  $1 - n\mathbf{u} > 0$  and  $1 - (n+b)\mathbf{u} > 0$ . Also, note that  $nb\mathbf{u}^2 \geq 0$ .  $\square$

**Exercise 19.****19.1**  $k = 0$ 

Show that  $\left( \frac{I + \Sigma^{(k)}}{0} \middle| \frac{0}{(1 + \epsilon_1)} \right) (1 + \epsilon_2) = I + \Sigma^{(k+1)}$

*Proof.* Given that if  $k = 0$ , then  $\epsilon_1 = 0$  and  $\Sigma^0$  is  $0 \times 0$  matrix, we have

$$(1 + 0) \cdot (1 + \epsilon_2) = \underbrace{1}_I + \underbrace{\epsilon_2}_{\Sigma^{(1)}} \quad (39)$$

Thus,  $\left( \frac{I + \Sigma^{(k)}}{0} \middle| \frac{0}{(1 + \epsilon_1)} \right) (1 + \epsilon_2) = I + \Sigma^{(k+1)}$  holds for  $k = 0$ .  $\square$

**19.2**  $k > 0$ 

*Proof.* For arbitrary  $k > 0$ ,

$$\left( \frac{I + \Sigma^{(k)}}{0} \middle| \frac{0}{(1 + \epsilon_1)} \right) (1 + \epsilon_2) \quad (40)$$

$$= \left( \frac{(I + \Sigma^{(k)})(1 + \epsilon_2)}{0} \middle| \frac{0}{(1 + \epsilon_1)(1 + \epsilon_2)} \right) \quad (41)$$

$$= \left( \frac{I + \epsilon_2 I + \Sigma^{(k)} + \epsilon_2 \Sigma^{(k)}}{0} \middle| \frac{0}{1 + \epsilon_1 + \epsilon_2 + \epsilon_1 \epsilon_2} \right) \quad (42)$$

$$= \underbrace{\left( \frac{I}{0} \middle| \frac{0}{1} \right)}_I + \underbrace{\left( \frac{\epsilon_2 I + \Sigma^{(k)} + \epsilon_2 \Sigma^{(k)}}{0} \middle| \frac{0}{\epsilon_1 + \epsilon_2 + \epsilon_1 \epsilon_2} \right)}_{\Sigma^{(k+1)}} \quad (43)$$

which tells us that  $\left( \frac{I + \Sigma^{(k)}}{0} \middle| \frac{0}{(1 + \epsilon_1)} \right) (1 + \epsilon_2) = I + \Sigma^{(k+1)}$  holds for  $k > 0$ .  $\square$

**Exercise 23.****23.1**

Show that  $\check{\kappa} = (x + \delta x)^T y$ , where  $|\delta x| \leq \gamma_n |x|$ .

*Proof.* Let  $\delta x = \Sigma^{(n)} x$ , where  $\Sigma^{(n)}$  is as in Theorem 20.

$$|\delta x| = |\Sigma^{(n)} x| = \begin{pmatrix} |\theta_n \chi_0| \\ |\theta_n \chi_1| \\ \vdots \\ |\theta_2 \chi_{n-1}| \end{pmatrix} \leq \begin{pmatrix} |\theta_n| |\chi_0| \\ |\theta_n| |\chi_1| \\ \vdots \\ |\theta_2| |\chi_{n-1}| \end{pmatrix} \leq |\theta_n| \begin{pmatrix} |\chi_0| \\ |\chi_1| \\ \vdots \\ |\chi_{n-1}| \end{pmatrix} \leq \gamma_n \begin{pmatrix} |\chi_0| \\ |\chi_1| \\ \vdots \\ |\chi_{n-1}| \end{pmatrix} = \gamma_n |x| \quad (44)$$

Thus, it can be concluded for the backward analysis that

$$|\delta x| \leq \gamma_n |x| \quad (45)$$

$\square$

**23.2**

Show that  $\check{\kappa} = x^T (y + \delta y)$ , where  $|\delta y| \leq \gamma_n |y|$ .

*Proof.* The proof for perturbation on input  $y$  is the same as that of perturbation on input  $x$ .  $\square$

**Exercise 25.**

*Proof.* We partition matrix  $A \in \mathbb{R}^{m \times n}$  and have

$$A = \begin{pmatrix} a_0^T \\ a_1^T \\ \vdots \\ a_{m-1}^T \end{pmatrix} \quad (46)$$

Then in terms of algorithm in Fig. 4 and R1-B,

$$\check{y} = \begin{pmatrix} \psi_0 \\ \psi_1 \\ \vdots \\ \psi_{m-1} \end{pmatrix} = \begin{pmatrix} a_0^T(x + \delta x) \\ a_1^T(x + \delta x) \\ \vdots \\ a_{m-1}^T(x + \delta x) \end{pmatrix} = \begin{pmatrix} a_0^T \\ a_1^T \\ \vdots \\ a_{m-1}^T \end{pmatrix} (x + \delta x) = A(x + \delta x) \quad (47)$$

where  $|\delta x| \leq \gamma_n |x|$  ( $\delta x$  is small). □

**Exercise 27.**