



THE UNIVERSITY OF TEXAS
AT AUSTIN

CS383C NUMERICAL ANALYSIS

Final Exam

Edited by L^AT_EX

Department of Computer Science

STUDENT

Jimmy Lin

xl5224

COURSE COORDINATOR

Robert A. van de Geijn

UNIQUE NUMBER

53180

RELEASE DATE

Dec. 2 2014

DUE DATE

Dec. 10 2014

TIME SPENT

10 hours

December 6, 2014

Exercises

I	Cholesky Factorization	2
1	SPD	2
1.1	Show A_{00} is SPD	2
1.2	$l_{10}^T = a_{10}^T L_{00}^{-T}$ is well defined	2
1.3	$\alpha_{11} - l_{10}^T l_{10} > 0$	2
1.4	Show equality	2
2		2
2.1	Another proof of Cholesky Factorization Theorem	2
2.2	Bordered Cholesky Algorithm	2
3	Cost of Bordered Algorithm	2
II	Method of Relatively Robust Representations	3
1	LDL^T Factorization for indefinite matrices	3
2	LDL^T Factorization for tridiagonal matrices	4
3	UDU^T Factorization for indefinite matrices	5
4	UDU^T Factorization for tridiagonal matrices	6
5	Twisted Factorization: ϕ_1	7
6	Twisted Factorization: Eigenvector	9

Part I

Cholesky Factorization

1 SPD

1.1 Show A_{00} is SPD

Proof. Since A is SPD, then

$$A^T = A \quad (1)$$

$$\forall x, x^T A x \geq 0 \quad (2)$$

From (1), we have

$$A^T = \begin{pmatrix} A_{00} & a_{10} \\ a_{10}^T & \alpha_{11} \end{pmatrix}^T = \begin{pmatrix} A_{00}^T & a_{10} \\ a_{10}^T & \alpha_{11} \end{pmatrix} = A \quad (3)$$

Then

$$A_{00}^T = A_{00} \quad (\text{symmetry of } A)$$

Also, we denote arbitrary $x = \begin{pmatrix} x_0 \\ \chi_1 \end{pmatrix}$ and then from (2), we have

$$x^T A x = \begin{pmatrix} x_0 \\ \chi_1 \end{pmatrix}^T \begin{pmatrix} A_{00} & a_{10} \\ a_{10}^T & \alpha_{11} \end{pmatrix} \begin{pmatrix} x_0 \\ \chi_1 \end{pmatrix} = x_0^T A_{00} x_0 + 2\chi_1 a_{10}^T x_0 + \chi_1^2 \alpha_{11} > 0, \forall x_0, \chi_1 \quad (4)$$

Let $\chi_1 = 0$, then we have

$$x_0^T A_{00} x_0 > 0, \forall x_0 \quad (\text{positive definiteness})$$

In terms of (symmetry of A) and (positive definiteness), then it is proved that A_{00} is SPD. \square

1.2 $l_{10}^T = a_{10}^T L_{00}^{-T}$ is well defined

Since L_{00} is non-singular, then it is easy to derive that L_{00}^T is also non-singular. Then L_{00}^{-T} exists. Hence,

$$l_{10}^T = a_{10}^T L_{00}^{-T} \quad (5)$$

is well-defined.

1.3 $\alpha_{11} - l_{10}^T l_{10} > 0$

1.4 Show equality

$$L \cdot L^T = \begin{pmatrix} L_{00} & l_{10} \\ l_{10}^T & \lambda_{11} \end{pmatrix} \begin{pmatrix} L_{00} & l_{10} \\ l_{10}^T & \lambda_{11} \end{pmatrix}^T = \begin{pmatrix} L_{00} L_{00}^T & L_{00} l_{10} \\ l_{10}^T L_{00}^T & l_{10}^T l_{10} + \lambda_{11}^2 \end{pmatrix} \quad (6)$$

Obviously, $L \cdot L^T = A$ if only if

$$A_{00} = L_{00} L_{00}^T \quad (7)$$

$$a_{10} = L_{00} l_{10} \quad (8)$$

$$\alpha_{11} = L_{10}^T l_{10} + \lambda_{11}^2 \quad (9)$$

2

2.1 Another proof of Cholesky Factorization Theorem

2.2 Bordered Cholesky Algorithm

3 Cost of Bordered Algorithm

Part II

Method of Relatively Robust Representations

1 LDL^T Factorization for indefinite matrices

```

% Copyright 2014 The University of Texas at Austin
%
% For licensing information see
%     http://www.cs.utexas.edu/users/flame/license.html
%
% Programmed by: Jimmy Xin Lin
%     jimmylin@utexas.edu

function [ A_out ] = LDL_unb( A )

[ ATL, ATR, ...
  ABL, ABR ] = FLA_Part_2x2( A, ...
                             0, 0, 'FLA_TL' );
while ( size( ATL, 1 ) < size( A, 1 ) )

    [ A00,  a01,    A02,  ...
      a10t, alpha11, a12t, ...
      A20,  a21,    A22 ] = FLA_Repart_2x2_to_3x3( ATL, ATR, ...
                                                    ABL, ABR, ...
                                                    1, 1, 'FLA_BR' );

    %-----%

    l21 = a21 / alpha11;
    A22 = A22 - l21 * a12t;
    a21 = l21;

    %-----%

    [ ATL, ATR, ...
      ABL, ABR ] = FLA_Cont_with_3x3_to_2x2( A00,  a01,    A02,  ...
                                              a10t, alpha11, a12t, ...
                                              A20,  a21,    A22, ...
                                              'FLA_TL' );

end

A_out = [ ATL, ATR
          ABL, ABR ];

return

```

2 LDL^T Factorization for tridiagonal matrices

Codes:

```
function [ A_out ] = LDL_TRI_unb( A )

    % here we leave out the partion and repartion code in the notes

    %-----%

    alpha11 = alpha11 % = delta11 (no-op)
    l21 = alpha21 / alpha11; % scalar
    alpha22 = alpha22 - l21 * alpha21;
    alpha21 = l21;

    %-----%

    % here we leave out the continue code in the notes

return
```

Costs:

- Divide: $1 \cdot n = n$ (l21 update)
- Multiply: $1 \cdot n = n$ (alpha22 update)
- Add/Subtract: $1 \cdot n = n$ (alpha21 update)

In terms of above analysis, the approximate cost is $\mathcal{O}(n)$.

Analytics: The way I come up with this algorithm is to instantiate the LDL^T factorization in last question to the case of tridiagonal matrices. That is, treat the α_{21} and l_{21} as vectors with only one non-zero entry.

3 UDU^T Factorization for indefinite matrices

```
% Copyright 2014 The University of Texas at Austin
%
% For licensing information see
%      http://www.cs.utexas.edu/users/flame/license.html
%
% Programmed by: Jimmy Xin Lin
%      jimmylin@utexas.edu

function [ A_out ] = UDU_unb( A )

[ ATL, ATR, ...
  ABL, ABR ] = FLA_Part_2x2( A, ...
                             0, 0, 'FLA_BR' );

while ( size( ABR, 1 ) < size( A, 1 ) )

    [ A00,  a01,      A02,  ...
      a10t, alpha11, a12t, ...
      A20,  a21,      A22 ] = FLA_Repart_2x2_to_3x3( ATL, ATR, ...
                                                       ABL, ABR, ...
                                                       1, 1, 'FLA_TL' );

    %-----%

    % alpha11 = alpha11 = delta11 (no-operation)
    l01 = a01 / alpha11;
    A00 = A00 - l01 * a01;
    a01 = l01;

    %-----%

    [ ATL, ATR, ...
      ABL, ABR ] = FLA_Cont_with_3x3_to_2x2( A00,  a01,      A02,  ...
                                              a10t, alpha11, a12t, ...
                                              A20,  a21,      A22, ...
                                              'FLA_BR' );

end

A_out = [ ATL, ATR
          ABL, ABR ];

return
```

4 UDU^T Factorization for tridiagonal matrices

```
function [ A_out ] = UDU_TRI_unb( A )

    % here we leave out the partion and repartion code in the notes

    %-----%

    alpha33 = alpha33 % = delta11 (no-op)
    l23 = alpha23 / alpha33; % scalar
    alpha22 = alpha22 - l23 * alpha23;
    alpha23 = l23;

    %-----%

    % here we leave out the continue code in the notes

return
```


5 Twisted Factorization: ϕ_1

For LDL^T factorization, we have

$$\begin{aligned}
 A &= LDL^T \\
 &= \begin{pmatrix} L_{00} & 0 & 0 \\ \lambda_{10}e_L^T & 1 & 0 \\ 0 & \lambda_{21}e_F & L_{22} \end{pmatrix} \begin{pmatrix} D_{00} & 0 & 0 \\ 0 & \delta_1 & 0 \\ 0 & 0 & D_{22} \end{pmatrix} \begin{pmatrix} L_{00} & 0 & 0 \\ \lambda_{10}e_L^T & 1 & 0 \\ 0 & \lambda_{21}e_F & L_{22} \end{pmatrix}^T \\
 &= \begin{pmatrix} L_{00}D_{00}L_{00}^T & \lambda_{10}L_{00}D_{00}e_L & 0 \\ \lambda_{10}e_L^TD_{00}L_{00}^T & \lambda_{10}^2e_L^TD_{00}e_L + \delta_1 & \lambda_{21}\delta_1e_F^T \\ 0 & \lambda_{21}\delta_1e_F & \lambda_{21}^2\delta_1e_Fe_F^T + L_{22}D_{22}L_{22}^T \end{pmatrix} \quad (10)
 \end{aligned}$$

$$= \begin{pmatrix} A_{00} & \alpha_{10}e_L & 0 \\ \alpha_{10}e_L^T & \alpha_{11} & \alpha_{21}e_F^T \\ 0 & \alpha_{21}e_F & A_{22} \end{pmatrix} \quad (11)$$

And by matching, we have

$$\begin{aligned}
 A_{00} &= L_{00}D_{00}L_{00}^T \\
 \alpha_{10}e_L &= \lambda_{10}L_{00}D_{00}e_L \\
 \alpha_{11} &= \lambda_{10}^2e_L^TD_{00}e_L + \delta_1 \\
 \alpha_{21} &= \lambda_{21}\delta_1 \\
 A_{22} &= \lambda_{21}^2\delta_1e_Fe_F^T + L_{22}D_{22}L_{22}^T
 \end{aligned} \quad (12)$$

Similarly, for UEU^T factorization, we have

$$\begin{aligned}
 A &= UEU^T \\
 &= \begin{pmatrix} U_{00} & v_{01}e_L & 0 \\ 0 & 1 & v_{21}e_F^T \\ 0 & 0 & U_{22} \end{pmatrix} \begin{pmatrix} E_{00} & 0 & 0 \\ 0 & \epsilon_1 & 0 \\ 0 & 0 & E_{22} \end{pmatrix} \begin{pmatrix} U_{00} & v_{01}e_L & 0 \\ 0 & 1 & v_{21}e_F^T \\ 0 & 0 & U_{22} \end{pmatrix}^T \\
 &= \begin{pmatrix} U_{00}E_{00}U_{00}^T + v_{01}\epsilon_1e_Le_L^T & v_{01}\epsilon_1e_L & 0 \\ v_{01}\epsilon_1e_L^T & v_{21}^2e_F^TE_{22}e_F + \epsilon_1 & v_{21}e_F^TE_{22}U_{22}^T \\ 0 & v_{21}U_{22}E_{22}e_F & U_{22}E_{22}U_{22}^T \end{pmatrix} \quad (13)
 \end{aligned}$$

$$= \begin{pmatrix} A_{00} & \alpha_{10}e_L & 0 \\ \alpha_{10}e_L^T & \alpha_{11} & \alpha_{21}e_F^T \\ 0 & \alpha_{21}e_F & A_{22} \end{pmatrix} \quad (14)$$

And by matching, we have

$$\begin{aligned}
 A_{00} &= U_{00}E_{00}U_{00}^T + v_{01}\epsilon_1e_Le_L^T \\
 \alpha_{10} &= v_{01}\epsilon_1 \\
 \alpha_{11} &= v_{21}^2e_F^TE_{22}e_F + \epsilon_1 \\
 \alpha_{21}e_F^T &= v_{21}e_F^TE_{22}U_{22}^T \\
 A_{22} &= U_{22}E_{22}U_{22}^T
 \end{aligned} \quad (15)$$

Now we consider the Twisted Factorization

$$\begin{pmatrix} L_{00} & 0 & 0 \\ \lambda_{10}e_L^T & 1 & v_{21}e_F^T \\ 0 & 0 & U_{22} \end{pmatrix} \begin{pmatrix} D_{00} & 0 & 0 \\ 0 & \phi_1 & 0 \\ 0 & 0 & D_{22} \end{pmatrix} \begin{pmatrix} L_{00} & 0 & 0 \\ \lambda_{10}e_L^T & 1 & v_{21}e_F^T \\ 0 & 0 & U_{22} \end{pmatrix}^T \quad (16)$$

$$= \begin{pmatrix} L_{00}D_{00}L_{00}^T & \lambda_{10}L_{00}D_{00}e_L & 0 \\ \lambda_{10}e_L^TD_{00}L_{00}^T & \phi_1 + \lambda_{10}^2e_L^TD_{00}e_L + v_{21}^2e_F^TE_{22}e_F & v_{21}e_F^TE_{22}U_{22}^T \\ 0 & v_{21}U_{22}E_{22}e_F & U_{22}E_{22}U_{22}^T \end{pmatrix} \quad (17)$$

$$= \begin{pmatrix} A_{00} & \alpha_{10}e_L & 0 \\ \alpha_{10}e_L^T & \alpha_{11} & \alpha_{21}e_F^T \\ 0 & \alpha_{21}e_F & A_{22} \end{pmatrix} \quad (18)$$

Then we have

$$\alpha_{11} = \phi_1 + \lambda_{10}^2 e_L^T D_{00} e_L + v_{21}^2 e_F^T E_{22} e_F \quad (19)$$

To satisfy (12), (15) and (19) at the same time, we need to have

$$\phi_1 = \frac{\delta_1 + \epsilon_1 - \lambda_{10}^2 e_L^T D_{00} e_L - v_{21}^2 e_F^T E_{22} e_F}{2} \quad (20)$$

Complexity:

- computation of $e_L^T D_{00} e_L$ or $e_F^T E_{22} e_F$ is $\mathcal{O}(1)$. (constant time)
- computation of the factorized matrix, it requires $\mathcal{O}(n)$ for assembling components of U and L so as to derive the resulted matrix.

6 Twisted Factorization: Eigenvector

Separate terms of the known condition as follows:

$$\underbrace{\begin{pmatrix} L_{00} & 0 & 0 \\ \lambda_{10}e_L^T & 1 & v_{21}e_F^T \\ 0 & 0 & U_{22} \end{pmatrix}}_S \underbrace{\begin{pmatrix} D_{00} & 0 & 0 \\ 0 & 0 & 0 \\ 0 & 0 & E_{22} \end{pmatrix}}_y \begin{pmatrix} L_{00} & 0 & 0 \\ \lambda_{10}e_L^T & 1 & v_{21}e_F^T \\ 0 & 0 & U_{22} \end{pmatrix}^T \begin{pmatrix} x_0 \\ \chi_1 \\ x_2 \end{pmatrix} = \begin{pmatrix} 0 \\ 0 \\ 0 \end{pmatrix} \quad (21)$$

Then we derive the form of S

$$S \triangleq \begin{pmatrix} L_{00} & 0 & 0 \\ \lambda_{10}e_L^T & 1 & v_{21}e_F^T \\ 0 & 0 & U_{22} \end{pmatrix} \cdot \begin{pmatrix} D_{00} & 0 & 0 \\ 0 & 0 & 0 \\ 0 & 0 & E_{22} \end{pmatrix} = \begin{pmatrix} L_{00}D_{00} & 0 & 0 \\ \lambda_{10}e_L^TD_{00} & 0 & v_{21}e_F^TE_{22} \\ 0 & 0 & U_{22}E_{22} \end{pmatrix} \quad (22)$$

Then relate it to y

$$S \cdot y = \begin{pmatrix} L_{00}D_{00} & 0 & 0 \\ \lambda_{10}e_L^TD_{00} & 0 & v_{21}e_F^TE_{22} \\ 0 & 0 & U_{22}E_{22} \end{pmatrix} \begin{pmatrix} y_0 \\ \psi_1 \\ y_2 \end{pmatrix} = \begin{pmatrix} 0 \\ 0 \\ 0 \end{pmatrix} \quad (23)$$

where

$$\begin{aligned} L_{00}^Tx_0 + \lambda_{10}e_L\chi_1 &= y_0 \\ \chi_1 &= \psi_1 \\ v_{21}\chi_1e_F + U_{22}^Tx_2 &= y_2 \end{aligned} \quad (24)$$

Solve (23), we have

$$\begin{aligned} y_0 &= 0 \\ \psi_1 &= c \text{ (constant)} \\ y_2 &= 0 \end{aligned} \quad (25)$$

In terms of (24), for the vector x , we need to solve the following system

$$\begin{aligned} L_{00}^Tx_0 + \lambda_{10}e_L\chi_1 &= 0 \\ \chi_1 &= c \\ v_{21}\chi_1e_F + U_{22}^Tx_2 &= 0 \end{aligned} \quad (26)$$

Note that this equation system has infinity number of solutions unless we set c fixed. Here, we set $c = 1$ for simplicity. Then

$$L_{00}^Tx_0 = -\lambda_{10}e_L \quad (27)$$

$$U_{22}^Tx_2 = -v_{21}e_F \quad (28)$$

which is actually two gaussian elimination problem. In terms of the special structure of L_{00} and U_{22} , the solution takes complexity $\mathcal{O}(n^2)$.