Habib University CS-113 Discrete Mathematics Spring 2018 HW 5 Solutions

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1. Prove, by double counting:

(a)
$$\sum_{k=1}^{n} k \binom{n}{k} = n2^{n-1}$$

Solution:

R.H.S.: The number of pairs (x, subset) where $x \in [n]$ and $subset \subseteq \{[n] \setminus x\}$ is $n2^{n-1}$.

L.H.S.: Alternatively, to create pairs, first choose a subset A of size k from n, and then create pairs (x, subset), where $x \in A$ and $subset = A \setminus x$. Since the first step can be done in $\binom{n}{k}$ ways, the pairs can be created in $k\binom{n}{k}$ ways, where k = 1, 2, ..., n. In total, there are $\sum_{k=1}^{n} k\binom{n}{k}$ pairs.

(b)
$$\sum_{k=-m}^{n} {m+k \choose r} {n-k \choose s} = {m+n+1 \choose r+s+1}$$

Solution:

R.H.S.: The number of ways to choose subsets of size r + 1 + s from [m + n + 1].

L.H.S.: Alternatively, let A be a subset of size r+1+s chosen from [m+n+1] and let x=1,2,...,(m+n+1), be the $r+1^{th}$ element in ordered A. Then the number of such subsets is

$$\sum_{k=0}^{n+m} \binom{k}{r} 1 \binom{n+m-k}{s}$$

where we choose r objects from the first k, and then s objects from the last n + m - k. The 1 in the middle of the summation is symbolic of the fixed $r + 1^{th}$ element. To complete the proof, we observe that

$$\Sigma_{k=0}^{n+m} \binom{k}{r} 1 \binom{n+m-k}{s} = \Sigma_{k=-m}^{n} \binom{m+k}{r} \binom{n-k}{s}$$

2. You and a friend play a game with a pile of N gold coins. On each turn, a player can remove 1, 3, or 6 coins from the pile. The winner is the one who takes the last coin. For 0 < N < 1000, how many starting positions are winning positions for the first player?

Solution: 666 (I will add the complete derivation soon).

3. Six friends, who each bought a gift, distribute them amongst each other such that everyone receives a gift. How many ways are there to distribute the gifts such that no friend gets their own gift?

1

Solution: 265 (I will add the complete derivation soon).

4. 0 points Find the number of non-negative integer solutions of the equation

$$x_1 + x_2 + \dots + x_n = y$$
 where $x_1 < x_2 < \dots < x_n$

5. Your friend's been working hard on developing the covert *Kabootar Messaging System (KMS)*. Their technique transforms each bit string into a unique bit string before sending over an insecure network. Once on the other side, the original bit string is fully reconstructed. Your friend claims that, on average, the *KMS* decreases the length of strings before transmission, and in any case, does not increase the length of a string. Prove that they are wrong.

Solution: For sake of absurdity, suppose that KMS does what it promises, i.e., generates a bijection $A \to B$, where A is the set of all strings over some alphabet Σ and B the set of encrypted strings over Σ , such that on average, strings in A are mapped to smaller strings in B and none to larger strings. Let $A_n \subset A$ be the first n strings in lexicographic order in A. Similarly let $B_n \subset B$ be the first n strings in lexicographic order in B. According to KMS, there exists a bijection from $A_n \to B_n$. But since the average length of strings in B_n is smaller than A_n , and no string in B_n is larger (in lexicographic terms) than any string in A_n , B_n must contain less strings than A_n . Therefore, $|B_n| < n$, and the original claim is absurd.