1.

Method	Vector.cs	Mecrosoft .Net framework
Count()	O(1)	O(1)
Capacity()	O(1)	O(1)
Add()	O(1) if count< capacity	O(1) if count< capacity
	O(n) if count>= capacity	O(n) if count>= capacity
IndexOf()	O(n)	O(n)
Insert()	O(n)	O(n)
Clear()	O(n)	O(n)
Contains()	O(n)	O(n)
Remove()	O(n)	O(n)
RemoveAt()	O(n)	O(n)

2

The below function is  $O(n \log(n))$  and  $\Omega(n \log(n))$ , therefore, it's  $\theta(n \log n)$ 

The name of this function is Merge Sort.

3 a. 
$$f=x^1/2$$
  $g=\log(n)$  
$$\lim_{n\to\infty}\frac{n^{\frac{1}{2}}}{\log n}=\infty$$
 f grows faster than g  $f\in\Omega(g)$ 

g = 2

$$\lim_{n\to\infty}\frac{1500}{2}=750$$

f is 750 times bigger than g, they grow at the same rate (f = 750g)

 $f \in \theta(g)$ 

c. 
$$f = 800*2^n$$

 $g = 3^n$ 

$$\lim_{n\to\infty} \frac{800*2^n}{3^n} = 0$$

f grows slower than g

 $f \in O(g)$ 

d. 
$$f = 4^{n+13}$$

 $g = 2^{(2n+2)}$ 

$$\lim_{n \to \infty} \frac{4^{n+13}}{2^{2n+2}} = 16777216$$

f is 16777216 bigger than g, they grows at the same rate (f(n) = 16777216 \* g(n))

 $f \in \theta(g)$ 

e.  $f = 9n \log(n)$ 

 $g = n \log(9n)$ 

$$\lim_{n \to \infty} \frac{9n * log(n)}{n * log(9n)} = 9$$

f is 9 time bigger than g, they grows at the same rate (f(n) = 9 \* g(n))

 $f \in \theta(g)$ 

$$f. f = O(g)$$

f = n!

g = (n+1)!

$$\lim_{n \to \infty} \frac{n!}{(n+1)!} = 0$$

f grows slower than g

 $f \in O(g)$