

The book cover is made of dark brown leather with a rich, aged patina. It is framed by an elaborate, multi-layered gold-tooled border. The innermost border consists of a repeating pattern of stylized, symmetrical foliate motifs. This is followed by a thin line of small, repeating geometric or floral stamps. The outermost border is a more complex, wider band featuring larger, more intricate foliate and scrollwork designs. The leather surface shows signs of wear, including some staining and a small dark mark near the top center.

# A Little Riak Book

Eric B. Redmond





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# Chapter 1

## Introduction

### What is Riak

Riak is an open-source, distributed key/value database for high availability, fault-tolerance, and near-linear scalability. It's a mouthful, that basically means Riak has remarkably high uptime, and grows with you.

In an increasingly interconnected world, major shifts have occurred in data management. The web and connected devices have spurred an explosion of both data collection and access unseen in the history of the world. The amount of data stored and managed has grown at a staggering rate, and in parallel, more people than ever require fast and reliable access to this data. This trend is known as *Big Data*.

---

### So What is Big Data?

There's a lot of discussion around what constitutes Big Data.

I have a 6 Terabyte RAID in my house to store videos and other backups. Does that count? On the other hand, CERN grabbed about [200 Petabytes](#) looking for the Higgs boson.

It's a hard number to pin down, because Big Data is a personal figure. What's big to one might be small to another. This is why many definitions don't refer to byte count at all, but instead about relative potentials. A reasonable, albeit wordy, [definition of Big Data](#) is given by Gartner.

Big Data are high-volume, high-velocity, and/or high-variety information assets that require new forms of processing to enable enhanced decision making, insight discovery and process optimization.

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The sweet spot of Riak is high-volume (data that's available to read and write when you need it), high-velocity (easily responds to growth), and high-variety information assets (you can store any type of data as a value).

Riak was built as a solution to real Big Data problems, based on the *Amazon Dynamo* design. Dynamo is a highly available design—meaning that it responds to requests quickly at very large scales, even if your application is storing and serving terabytes of data a day. Riak had been used in production prior to being open-sourced in 2009. It's currently used by Github, Comcast, Voxer, Disqus and others, with the larger systems storing hundreds of TBs of data, and handling several GBs per node daily.

Riak was written on the Erlang programming language. Erlang was chosen due to its strong support for concurrency, solid distributed communication, hot code loading, and fault-tolerance. It runs on a virtual machine, so running Riak requires an Erlang installation.

So should you use Riak? A good rule of thumb for potential users is to ask yourself if every moment of downtime will cost you in some way (money, users, etc). Not all systems require such extreme amounts of uptime, and if you don't, Riak may not be for you.

## About This Book

This is not an “install and follow along” guide. This is a “read and comprehend” guide. Don't feel compelled to have Riak, or even have a computer handy, when starting this book. You may feel like installing at some point, and if so, instructions can be found on the [Riak docs](#).

In my opinion, the most important section of this book is the concepts chapter. If you already have a little knowledge it may start slow, but it picks up in a hurry. After laying the theoretical groundwork, we'll move onto helping developers use Riak, by learning how to query it and tinker with some settings. Finally, we'll go over the basic details that operators should know, such as how to set up a Riak cluster, configure some values, use optional tools, and more.

## Chapter 2

# Concepts

Believe me, dear reader, when I suggest that thinking in a distributed fashion is awkward. When I had first encountered Riak, I was not prepared for some of its more preternatural concepts. Our brains just aren't hardwired to think in a distributed, asynchronous manner. Richard Dawkins coined the term *Middle World*—the serial, rote land humans encounter every day, which exists between the extremes of the very small strangeness of quarks and the vastness of outer space. We don't consider these extremes clearly because we don't encounter them on a daily basis, just like distributed computations and storage. So we create models and tools to bring the physical act of scattered parallel resources in line, on our more ordinary synchronous terms. Where Riak takes great pains to simplify the hard parts, it does not pretend that they don't exist. Just like you can never hope to program at an expert level without any knowledge of memory or CPU management, so too can you never safely develop a highly available clusters without a firm grasp of a few underlying concepts.

## The Landscape

The existence of databases like Riak is the culmination of two basic trends: accessible technology spurring different data requirements, and gaps in the data management market.

First, as we've seen steady improvements in technology along with reductions in cost, vast amounts of computing power and storage are now within the grasp of nearly anyone. Along with our increasingly interconnected world caused by the web and shrinking, cheaper computers (like smartphones), this has catalyzed an exponential growth of data, and a demand for more predictability and speed by savvier users. In other words, more data is being created on the front-end, while more data is being managed on the backend.

Second, relational database management systems (RDBMS) have become focused over the years for a standard set of use-cases, like business intelligence. They were also technically tuned for squeezing performance out of single larger servers, like optimizing disk access, even while cheap commodity (and virtualized) servers made horizontal growth increasingly attractive. As cracks in relational implementations became apparent, custom implementations arose in response to specific problems not originally envisioned by the relational DBs.

These new databases are collected under the moniker *NoSQL*, and Riak is of its ilk.

## Database Models

Modern databases can be loosely grouped into the ways they represent data. Although I'm presenting 5 major types (the last 4 are considered NoSQL models), these lines are often blurred—you can use some key/value stores as a document store, you can use a relational database to just store key/value data.

---

### A Quick note on JOINS

Unlike relational databases, but similar to document and columnar stores, objects cannot be joined by Riak. Client code is responsible for accessing values and merging them, or by other code such as mapreduce.

The ability to easily join data across physical servers is a tradeoff that separates single node databases like relational and graph, from *naturally partitionable* systems like document, columnar, and key/value stores.

This limitation changes how you model data. Relational normalization (organizing data to reduce redundancy) exists for systems that can cheaply join data together per request. However, the ability to spread data across multiple nodes requires a denormalized approach, where some data is duplicated, and computed values may be stored for the sake of performance.

---

1. **Relational.** Traditional databases usually use SQL to model and query data. They are useful for data which can be stored in a highly structured schema, yet require flexible querying. Scaling a relational database (RDBMS) traditionally occurs by more powerful hardware (vertical growth).

Examples: *PostgreSQL, MySQL, Oracle*

2. **Graph.** These exist for highly interconnected data. They excel in modeling complex relationships between nodes, and many implementations can handle multiple billions of nodes and relationships (or edges and vertices). I tend to include *triplestores* and *object DBs* to be specialized variants.

Examples: *Neo4j, Graphbase, InfiniteGraph*

3. **Document.** Document datastores model hierarchical values called documents, represented in formats such as JSON or XML, and do not enforce a document schema. They generally support distributing across multiple servers (horizontal growth).

Examples: *CouchDB, MongoDB, Couchbase*

4. **Columnar.** Popularized by [Google's BigTable](#), this form of database exists to scale across multiple servers, and groups like data into column families. Column values can be individually versioned and managed, though families are defined in advance, not unlike RDBMS schemas.

Examples: *HBase, Cassandra, BigTable*

5. **Key/Value.** Key/Value, or KV stores, are conceptually like hashtables, where values are stored and accessed by an immutable key. They range from single-server varieties like *Memcached* used for high-speed caching, to multi-datacenter distributed systems like *Riak Enterprise*.

Examples: *Riak*, *Redis*, *Voldemort*

## The Fallacies of Distributed Computing

One detour in the land of distributed databases is the condition that they are distributed systems replete with their benefits and handicaps. Engineers at Sun Microsystems created this list of [eight fallacies](#) that engineers new to distributed systems often fall victim to. They still apply today, even when operating a database like Riak.

1. The network is reliable.
2. Latency is zero.
3. Bandwidth is infinite.
4. The network is secure.
5. Topology doesn't change.
6. There is one administrator.
7. Transport cost is zero.
8. The network is homogeneous.

I always recommend to initiates to take the time to grock this list. Keeping these points in the back of your mind can save days of pain and expense in the future.

## Riak Components

Riak is a Key/Value (KV) database, built from the ground up to safely distribute data across a cluster of physical servers, called nodes. A Riak cluster is also known as a Ring (we'll cover why later).

For now, we'll only consider the concepts required to be a Riak users, and cover operations later.

Riak functions similarly to a very large hashtable. Depending on your background, you may instead call it a map, or dictionary, or object. But the idea is the same: you store a value with an immutable key, and retrieve it later.

## Key and Value

If Riak were a variable that functioned as a hashtable, you might set the value of your favorite food using the *key* `favorite`.

```
hashtable["favorite"] = "pizza"
```

And retrieve the *value* `pizza` by using the same key as before.

```
food = hashtable["favorite"]  
food == "pizza"
```

One day you burn the roof of your mouth. In anger, you decided your favorite food is now cold pizza.

```
hashtable["favorite"] = "cold pizza"
```

Successive requests for `favorite` will now return `cold pizza`.

For convenience, we call a key/value pair an *object*. Together our `favorite/pizza` pair is referred to as the “favorite object”, rather than the more verbose “favorite key and its value”.

## Buckets

*Buckets* are how Riak allows you to categorizes objects. You can group multiple objects into logical collections, where identical keys will not overlap between buckets.

You can think of buckets as [namespaces](#).

Using our `favorite` example from above, we can specify a favorite food, versus a favorite animal, by using the same key. Unless you’re a Midwest farm kid like me, these categories probably won’t overlap much.

```
food["favorite"] = "pizza"  
animals["favorite"] = "red panda"
```

You could have just named your keys `edible_favorite` and `animal_favorite`, but buckets allow for cleaner key naming, and has other added benefits that I’ll outline later.

Buckets are so useful in Riak that all keys must belong to a bucket. There is no global namespace.

In fact in Riak, the true definition of an object key is actually `bucket/key`.

## Replication and Partitions

Distributing data across several nodes is how Riak is able to remain highly available, while tolerant of outages and network partitions. Riak combines two styles of distribution to achieve this: [replication](#) and [partitions](#).



## Replication

**Replication** is the act of duplicating data across multiple servers. Riak replicates by default.

The obvious benefit of replication is that if one node goes down, nodes that contain replicated data remain available to serve requests. In other words, the system remains *available*.

For example, imagine you have a list of country keys, whose values contain those countries' capitals. If all you do is replicate that data to 2 servers, you would have 2 duplicate databases.

### Node A

```
"Afghanistan": "Kabul"  
"Albania":      "Tirana"  
"Algeria":      "Algiers"  
...  
"Yemen":        "Sanaa"  
"Zambia":       "Lusaka"  
"Zimbabwe":    "Harare"
```

### Node B

```
"Afghanistan": "Kabul"  
"Albania":      "Tirana"  
"Algeria":      "Algiers"  
...  
"Yemen":        "Sanaa"  
"Zambia":       "Lusaka"  
"Zimbabwe":    "Harare"
```

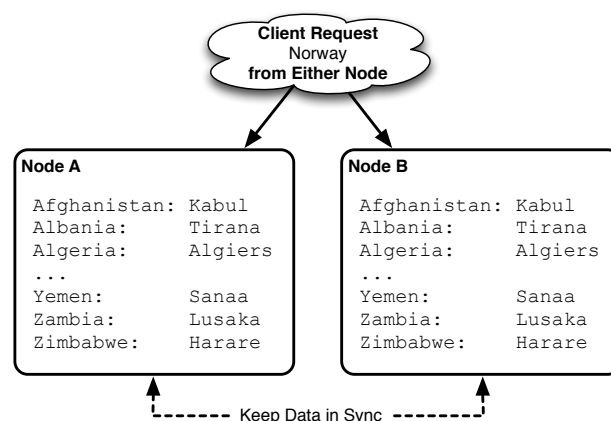


Figure 2.1: Replication

The downside with replication is that you are multiplying the amount of storage required for every duplicate. There is also some network overhead with this approach, since values

must also be routed to all replicated nodes on write. But there is a more insidious problem with this approach, which I will cover shortly.

## Partitions

A **partition** is how we divide a set of keys onto separate physical servers. Rather than duplicate values, we pick one server to exclusively host a range of keys, and the other servers to host remaining non-overlapping ranges.

With partitioning, our total capacity can increase without any big expensive hardware, just lots of cheap commodity servers. If we decided to partition our database into 1000 parts across 1000 nodes, we have (hypothetically) reduced the amount of work any particular server must do to 1/1000th.

For example, if we partition our countries into 2 servers, we might put all countries beginning with letters A-N into Node A, and O-Z into Node B.

### Node A

```
"Afghanistan": "Kabul"  
"Albania":     "Tirana"  
"Algeria":     "Algiers"  
...  
"Norway":      "Oslo"
```

### Node B

```
"Oman":        "Muscat"  
...  
"Yemen":       "Sanaa"  
"Zambia":      "Lusaka"  
"Zimbabwe":    "Harare"
```

There is a bit of overhead the partition approach. Some service must keep track of what range of values live on which node. A requesting application must know that the key `Spain` will be routed to Node B, not Node A.

There's also another downside. Unlike replication, simple partitioning of data actually *decreases* uptime. If one node goes down, that entire partition of data is unavailable. This is why Riak uses both replication and partitioning.

## Replication+Partitions

Since partitions allow us to increase capacity, and replication improves availability, Riak combines them. We partition data across multiple nodes, as well as replicate that data into multiple nodes.

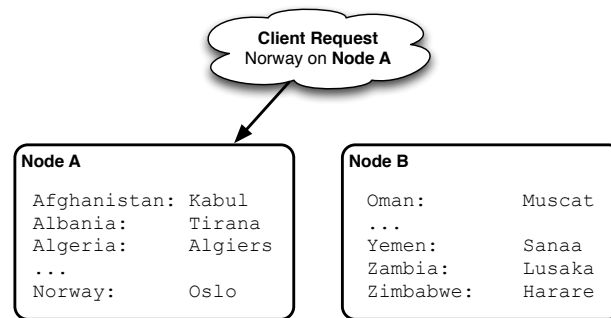


Figure 2.2: Partitions

Where our previous example partitioned data into 2 nodes, we can replicate each of those partitions into 2 more nodes, for a total of 4.

### Nodes A & C

```

"Afghanistan": "Kabul"
"Albania":     "Tirana"
"Algeria":     "Algiers"
...
"Norway":      "Oslo"
  
```

### Nodes B & D

```

"Oman":        "Muscat"
...
"Yemen":       "Sanaa"
"Zambia":      "Lusaka"
"Zimbabwe":    "Harare"
  
```

Our server count has increased, but so has our capacity and reliability. If you're designing a horizontally scalable system by partitioning data, you must deal with replicating those partitions.

The Riak team suggests a minimum of 5 nodes for a Riak cluster, and replicating to 3 nodes (this setting is called `n_val`, for the number of *nodes* on which to replicate each object).

## The Ring

Riak follows the *consistent hashing* technique, that conceptually maps objects to the edge of a circle or ring. It has the benefit of reducing the amount of data that must be rebalanced when a node goes down.

Riak partitions are not mapped alphabetically (as we used in the examples above), but instead, a partition marks a range of key hashes (SHA-1 function applied to a key). The

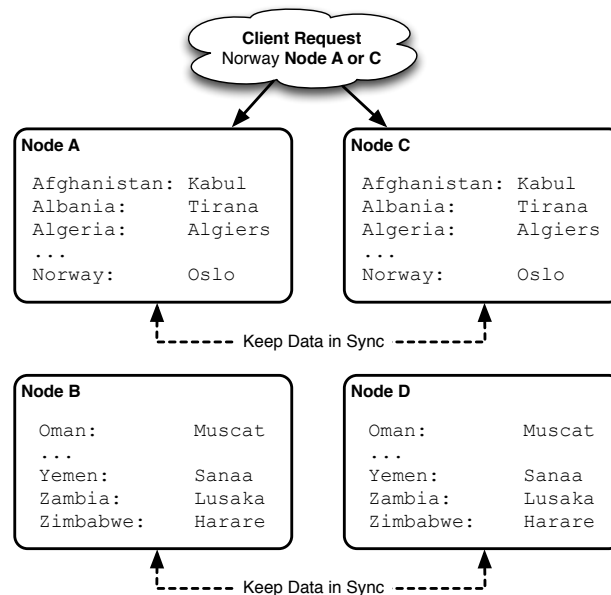


Figure 2.3: Replication Partitions

maximum hash value is  $2^{160}$ , and divided into some number of partitions—64 partitions by default (the Riak config setting is `ring_creation_size`).

Let's walk through what all that means. If you have the key `favorite`, applying the SHA-1 algorithm would return `dc2b 258d 7221 3f8d 05d1 5973 a66d c156 847b 83f4` in hexadecimal. With 64 partitions, each partition has  $1/64$  of the  $2^{160}$  possible values, making the first partition range from 0 to  $2^{154}-1$ , the second range is  $2^{154}$  to  $2^{154}+1$ , and so on, up to the last partition  $63 \cdot 2^{154}-1$  to  $2^{160}-1$ .

We won't do all of the math, but trust me when I say `favorite` falls within the range of partition 3.

If we visualize our 64 partitions as a ring, `favorite` falls here.

You may have wondered, "Didn't he say that Riak suggests a minimum of 5 nodes? How can we put 64 partitions on 5 nodes?" We just give each node more than one partition, which Riak calls a  *vnode* , or  *virtual node* .

We count around the ring of vnodes in order, assigning each node to the next available vnode, until all vnodes are accounted for. So partition/vnode 1 would be owned by Node A, vnode 2 owned by Node B, up to vnode 5 owned by Node E. Then we continue by giving Node A vnode 6, Node B vnode 7, and so on, until our vnodes have been exhausted, leaving us this list.

- A = [1,6,11,16,21,26,31,36,41,46,51,56,61]
- B = [2,7,12,17,22,27,32,37,42,47,52,57,62]
- C = [3,8,13,18,23,28,33,38,43,48,53,58,63]
- D = [4,9,14,19,24,29,34,39,44,49,54,59,64]
- E = [5,10,15,20,25,30,35,40,45,50,55,60]

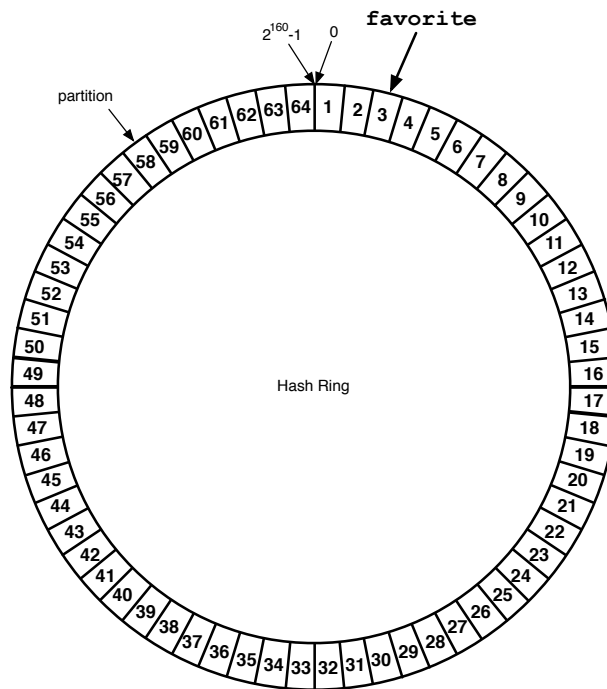


Figure 2.4: Riak Ring

So far we've partitioned the ring, but what about replication? When we write a new value to Riak, it will replicate the result in some number of nodes, defined by a setting called `n_val`. In our 5 node cluster it defaults to 3.

So when we write our `favorite` object to vnode 3, it will be replicated to vnodes 4 and 5. This places the object in physical nodes C, D, and E. Once the write is complete, even if node C crashes, the value is still available on 2 other nodes. This is the secret of Riak's high availability.

We can visualize the Ring with its vnodes, managing nodes, and where `favorite` will go.

The Ring is more than just a circular array of hash partitions. It's also a system of metadata that gets copied to every node. Each node is aware of every other node in the cluster, which nodes own which vnodes, and other system data.

Armed with this information, requests for data can target any node. It will horizontally access data from the proper nodes, and return the result.

## Practical Tradeoffs

So far we've covered the good parts of partitioning and replication: highly available when responding to requests, and inexpensive capacity scaling on commodity hardware. With the clear benefits of horizontal scaling, why is it not more common?

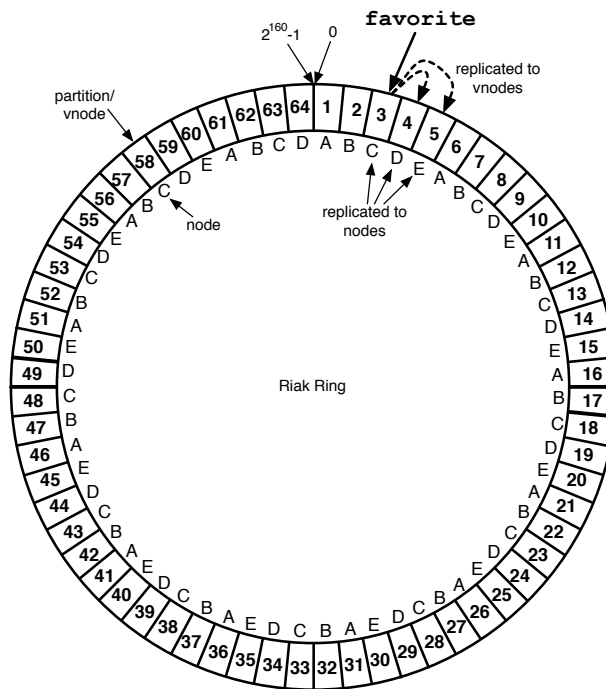


Figure 2.5: Riak Ring

## CAP Theorem

Classic RDBMS databases are *write consistent*. Once a write is confirmed, successive reads are guaranteed to return the newest value. If I save the value `cold pizza` to my key `favorite`, every future read will consistently return `cold pizza` until I change it.

But when values are distributed, *consistency* might not be guaranteed. In the middle of an object's replication, two servers could have different results. When we update `favorite` to `cold pizza` on one node, another node might contain the older value `pizza`, because of a network connectivity problem. If you request the value of `favorite` on either side of a network partition, two different results could possibly be returned—the database is inconsistent.

We do have an alternative choice. Rather than lose consistency, you could choose to lose *availability*. We may, for instance, decide to lock the entire database during a write, and simply refuse to serve requests until that value has been replicated to all relevant nodes. Clients have to wait while their results can be brought into a consistent state (ensuring all replicas will return the same value) or fail if the nodes have trouble communicating. For many high-traffic read/write use-cases, like an online shopping cart where even minor delays will cause people to just shop elsewhere, this is not an acceptable sacrifice.

This tradeoff is known as Brewer's CAP theorem. CAP loosely states that you can have a C (consistent), A (available), or P (partition-tolerant) system, but you can only choose 2. Assuming your system is distributed, you're going to be partition-tolerant, meaning, that your network can tolerate packet loss. If a network partition occurs between nodes, your servers still run.



---

## Not Quite C

Strictly speaking, Riak has a tunable latency-availability tradeoff. The concession is similar to tuning A/C. To decrease read/write latency effectively improves the odds of consistency, by making requests unavailable in certain circumstances in the way a CP system would.

Currently, no setting can make Riak truly CP in the general case, but features for a few strict cases are being researched.

---

## N/R/W

A question the CAP theorem demands you answer with a distributed system is: do I give up strict consistency, or give up ensured availability? If a request comes in, do I lock out requests until I can enforce consistency across the nodes? Or do I serve requests at all costs, with the caveat that the database may become inconsistent?

Riak's solution is based on Amazon Dynamo's novel approach of a *tunable* AP system. It takes advantage of the fact that, though the CAP theorem is true, you can choose what kind of tradeoffs you're willing to make. Riak is highly available to serve requests, with the ability to tune its level of availability (nearing, but never quite reaching, full consistency).

Riak allows you to choose how many nodes you want to replicate an object to, and how many nodes must be written to or read from per request. These values are settings labeled `n_val` (the number of nodes to replicate to), `r` (the number of nodes read from before returning), and `w` (the number of nodes written to before considered successful).

A thought experiment might help clarify.

## N

With our 5 node cluster, having an `n_val=3` means values will eventually replicate to 3 nodes, as we've discussed above. This is the *N value*. You can set the other values to be all `n_val` nodes with the shorthand `all`.

## W

But you may not wish to wait for all nodes to be written to before returning. You can choose to write to all 3 (`w=3` or `w=all`), which means my values are more likely to be consistent, or choose to write only 1 (`w=1`), and allow the remaining 2 nodes to write asynchronously, but return a response quicker. This is the *W value*.

In other words, setting `w=all` would help ensure your system was more likely to be consistent, at the expense of waiting longer, with a chance that your write would fail if fewer than 3 nodes were available (meaning, over half of your total servers are down).

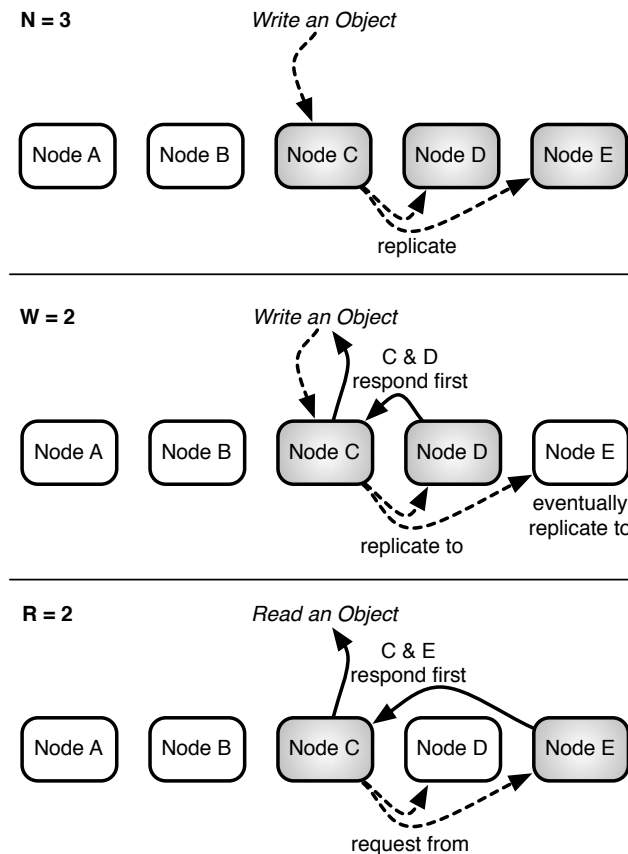


Figure 2.6: NRW

**R**

The same goes for reading. To ensure you have the most recent value, you can read from all 3 nodes containing objects ( $r=all$ ). Even if only 1 of 3 nodes has the most recent value, we can compare all nodes against each other and choose the latest one, thus ensuring some consistency. Remember when I mentioned that RDBMS databases were *write consistent*? This is close to *read consistency*. Just like  $w=all$ , however, the read will fail unless 3 nodes are available to be read. Finally, if you only want to quickly read any value,  $r=1$  has low latency, and is likely consistent if  $w=all$ .

In general terms, the N/R/W values are Riak's way of allowing you to trade less consistency for more availability.

**Vector Clock**

If you've followed thus far, I only have one more conceptual wrench to throw at you. I wrote earlier that with  $r=all$ , we can "compare all nodes against each other and choose the latest one." But how do we know which is the latest value? This is where Vector Clocks come into play.

Vector clocks measure a sequence of events, just like a normal clock. But since we can't reasonably keep dozens, or hundreds, or thousands of servers in sync (without really exotic hardware, like geosynchronized atomic clocks, or quantum entanglement), we instead keep

track of how, and who, modifies an object. It's as easy as keeping a vector (or array) of which clients change an object in which order. That way we can tell if an object is being updated or if a write conflict has occurred.

Let's use our `favorite` example again, but this time we have 3 people trying to come to a consensus on their favorite food: Aaron, Britney, and Carrie. We'll track the value each has chosen, and the relevant vector clock, or *vclock*.

When Aaron sets the `favorite` object to `pizza`, a hypothetical vector clock could contain his name, and the number of updates he's performed.

```
vclock: [Aaron: 1]
value:  pizza
```

Britney now comes along, and reads `favorite`, but decides to update `pizza` to `cold pizza`. When using *vclocks*, she must provide the *vclock* returned from the request she wants to update. This is how Riak can help ensure you're updating a previous value, and not merely overwriting with your own.

```
vclock: [Aaron: 1, Britney: 1]
value:  cold pizza
```

At the same time as Britney, Carrie decides that `pizza` was a terrible choice, and tried to change Aaron's value to `lasagna`.

```
vclock: [Aaron: 1, Carrie: 1]
value:  lasagna
```

This presents a problem, because there are now two vector clocks in play that diverge from `[Aaron: 1]`. So Riak can store both values and both *vclocks*.

Later in the day Britney checks again, but this time she gets the two conflicting values, with two *vclocks*.

```
vclock: [Aaron: 1, Britney: 1]
value:  cold pizza
---
vclock: [Aaron: 1, Carrie: 1]
value:  lasagna
```

It's clear that a decision must be made. Since two people generally agreed on `pizza`, Britney resolves the conflict by deciding on Aaron's original `pizza` value, and updating with her `vclock`.

```
vclock: [Aaron: 1, Britney: 2]
value:  pizza
```

Now we are back to the simple case, where requesting the value of `favorite` will just return the agreed upon `pizza`.

Beyond the ability for vector clocks to provide a reasonable history of updates, is also used when reading values from two conflicting nodes. This is how we can compare the reads of multiple nodes and decide upon the most recent version.

If you're a programmer, you may notice that this is not unlike a version control system, like **git**, where conflicting branches may require manual merging into one.

The Riak mechanism uses internal hashing and system clocks to stop unbounded `vclock` growth. We'll dig into more details of Riak's `vclocks` in the next chapter.

## Riak and ACID

---

### Distributed Relational is Not Exempt

You may have wondered why we don't just distribute a standard relational database. After all, MySQL has the ability to cluster, and it's ACID, right? Yes and no.

A single node in the cluster is ACID, but the entire cluster is not without a loss of availability, and often worse, increased latency. When you write to a primary node, and a secondary node is replicated to, a network partition can occur. To remain available, the secondary will not be in sync (eventually consistent). Have you ever lost data between a failure and a backup? Same idea.

Or, the entire transaction can fail, making the whole cluster unavailable. Even ACID databases cannot escape the scourge of CAP.

---

Unlike single node databases like Neo4j or PostgreSQL, Riak does not support *ACID* transactions. Locking across multiple servers would kill write availability, and equally concerning, increase latency. While ACID transactions promise *Atomicity*, *Consistency*, *Isolation*, and *Durability*—Riak and other NoSQL databases follow *BASE*, or *Basically Available*, *Soft state*, *Eventually consistent*.

The BASE acronym was meant as shorthand for the goals of non-ACID-transactional databases like Riak. It is an acceptance that distribution is never perfect (basically available), all data is in flux (soft state), and that true consistency is generally untenable (eventually consistent).

Be wary if anyone promises highly available distributed ACID transactions—it's usually couched in some diminishing adjective or caveat like *row transactions*, or *per node trans-*

*actions*, which basically mean *not transactional* in terms you would normally use to define it.

As your server count grows—especially as you introduce multiple datacenters—the odds of partitions and node failures drastically increase. My best advice is to design for it upfront.

## Wrapup

Riak is designed to bestow a range of real-world benefits, but equally, to handle the fallout of wielding such power. Consistent hashing and vnodes are an elegant solution to horizontally scaling across servers. N/R/W allows you to dance with the CAP theorem by fine-tuning against its constraints. And vector clocks allow another step closer to true consistency by allowing you to manage conflicts that will occur at high load.

We'll cover other technical concepts as needed, like the gossip protocol, hinted handoff, or read-repair.

Next we'll go through Riak as a user. We'll check out lookups, take advantage of write hooks, and alternative query options like secondary indexing, search, and mapreduce.





# Chapter 3

## Developers

---

### A Note on “Node”

It's worth mentioning that I use the word “node” a lot. Realistically, this means a physical/virtual server, but really, the workhorses of Riak are vnodes.

When you write to multiple vnodes, Riak will attempt to spread values to as many physical servers as possible. However, this isn't guaranteed (for example, if you have 64 vnodes, and only 2 physical ones, setting replication to 5 is perfectly fine, if not a bit silly). You're safe conceptualizing nodes as Riak instances, and it's simpler than qualifying “vnode” all the time. If something applies specifically to a vnode, I'll mention it.

---

*We're going to hold off on the details of installing Riak at the moment. If you'd like to follow along, it's easy enough to get started by following the [install documentation](#) on the website. If not, this is a perfect section to read while you sit on a train without internet connection.*

Developing with a Riak database is quite easy to do, once you understand some of the finer points. It is a key/value store, in the technical sense (you associate values with keys, and retrieve them using the same keys) but it offers so much more. You can embed write hooks to fire before or after a write, or index data for quick retrieval. Riak has a SOLR-based search, and lets you run mapreduce functions to extract and aggregate data across a huge cluster in a relatively short timespan. We'll show some of the bucket-specific settings developers can configure.

## Lookup

---

### Supported Languages

Riak has official drivers for the following languages: Erlang, Java, PHP, Python, Ruby

Including community supplied drivers, supported languages are even more numerous: C/C++, Clojure, Common Lisp, Dart, Go, Groovy, Haskell, Javascript (jquery and nodejs), Lisp Flavored Erlang, .NET, Perl, PHP, Play, Racket, Scala, Smalltalk

There are also dozens of other [project-specific addons](#).

---

Since Riak is a KV database, the most basic commands are setting and getting values. We'll use the HTTP interface, via curl, but we could just as easily use Erlang, Ruby, Java, or any other supported language.

The basic structure of a Riak request is setting a value, reading it, and maybe eventually deleting it. The actions are related to HTTP methods (PUT, GET, POST, DELETE).

```
PUT    /riak/bucket/key
GET    /riak/bucket/key
DELETE /riak/bucket/key
```

## PUT

The simplest write command in Riak is putting a value. It requires a key, value, and a bucket. In curl, all HTTP methods are prefixed with -X. Putting the value `pizza` into the key `favorite` under the `food` bucket is done like this:

```
curl -XPUT http://localhost:8098/riak/food/favorite \
  -H Content-Type:text/plain \
  -d pizza
```

I threw a few curveballs in there. The -d flag denotes the next string will be the value. We've kept things simple with the string `pizza`, declaring it as text with the proceeding line -H 'Content-Type:text/plain'. This defined the HTTP MIME type of this value as plain text. We could have set any value at all, be it XML or JSON—even an image or a video. Any HTTP MIME type is valid content (which is anything, really).

## GET

The next command reads the value `pizza` under the bucket/key `food/favorite`.

```
curl -XGET http://localhost:8098/riak/food/favorite
pizza
```

This is the simplest form of read, responding with only the value. Riak contains much more information, which you can access if you read the entire response, including the HTTP header.

In curl you can access a full response by way of the -i flag. Let's perform the above query again, adding that flag.

```
curl -i -XGET http://localhost:8098/riak/food/favorite
HTTP/1.1 200 OK
X-Riak-Vclock: a85hYGBgzGDKBVIcypz/fgaUHjmdwZTImMfKcN3h1Um+LAA=
```

```
Vary: Accept-Encoding
Server: MochiWeb/1.1 WebMachine/1.9.0 (someone had painted...
Link: </riak/food>; rel="up"
Last-Modified: Wed, 10 Oct 2012 18:56:23 GMT
ETag: "1yHn7L0XMEoMVXRGp4g0om"
Date: Thu, 11 Oct 2012 23:57:29 GMT
Content-Type: text/plain
Content-Length: 5

pizza
```

The anatomy of HTTP is a bit beyond this little book, but let's look at a few parts worth noting.

**Status Codes** The first line gives the HTTP version 1.1 response code 200 OK. You may be familiar with the common website code 404 Not Found. There are many kinds of [HTTP status codes](#), and the Riak HTTP interface stays true to their intent: **1xx Informational, 2xx Success, 3xx Further Action, 4xx Client Error, 5xx Server Error**

Different actions can return different response/error codes. Complete lists can be found in the [official API docs](#).

**Timings** A block of headers represents different timings for the object or the request.

- **Last-Modified** - The last time this object was modified (created or updated).
- **ETag** - An *entity tag* which can be used for cache validation by a client.
- **Date** - The time of the request.
- **X-Riak-Vclock** - A logical clock which we'll cover in more detail later.

**Content** These describe the HTTP body of the message (in Riak's terms, the *value*).

- **Content-Type** - The type of value, such as `text/xml`.
- **Content-Length** - The length, in bytes, of the message body.

Some other headers like `Link` will be covered later in this chapter.

## POST

Similar to PUT, POST will save a value. But with POST a key is optional. All it requires is a bucket name, and it will generate a key for you.

Let's add a JSON value to represent a person under the `people` bucket. The response header is where a POST will return the key it generated for you.

```
curl -i -XPOST http://localhost:8098/riak/people \
  -H Content-Type:application/json \
  -d {"name":"aaron"}
HTTP/1.1 201 Created
Vary: Accept-Encoding
Server: MochiWeb/1.1 WebMachine/1.9.2 (someone had painted...
Location: /riak/people/DNQGJY0KtchMirkidasA066yj5V
Date: Wed, 10 Oct 2012 17:55:22 GMT
Content-Type: application/json
Content-Length: 0
```

You can extract this key from the `Location` value. Other than not being pretty, this key is just as if you defined your own key via PUT.

**Body** You may note that no body was returned with the response. For any kind of write, you can add the `returnbody=true` parameter to force a value to return, along with value-related headers like `X-Riak-Vclock` and `ETag`.

```
curl -i -XPOST http://localhost:8098/riak/people?returnbody=true \
  -H Content-Type:application/json \
  -d {"name":"billy"}
HTTP/1.1 201 Created
X-Riak-Vclock: a85hYGBgzGDKBVIcypz/fgaUHjmdwZTImMfKkD3z10m+LAA=
Vary: Accept-Encoding
Server: MochiWeb/1.1 WebMachine/1.9.0 (someone had painted...
Location: /riak/people/DnetI8GHIBK2yBF0Ecj1EhHprss
Link: </riak/people>; rel="up"
Last-Modified: Tue, 23 Oct 2012 04:30:35 GMT
ETag: "7DsE7SEqAtY12d8T1HMkWZ"
Date: Tue, 23 Oct 2012 04:30:35 GMT
Content-Type: application/json
Content-Length: 16
```

```
{"name":"billy"}
```

This is true for PUTs and POSTs.

## DELETE

The final basic operation is deleting keys, which is similar to getting a value, but sending the DELETE method to the `url/bucket/key`.

```
curl -XDELETE http://localhost:8098/riak/people/DNQGJY0KtchMirkidasA066yj5V
```

A deleted object in Riak is internally marked as deleted, by writing a marker known as a *tombstone*. Later, another process called a *reaper* clears the marked objects from the backend (possibly, the reaper may be turned off).

This detail isn't normally important, except to understand two things:

1. In Riak, a *delete* is actually a *write*, and should be considered as such.
2. Checking for the existence of a key is not enough to know if an object exists. You might be reading a key between a delete and a reap. You must read for tombstones.

## Lists

Riak provides two kinds of lists. The first lists all *buckets* in your cluster, while the second lists all *keys* under a specific bucket. Both of these actions are called in the same way, and come in two varieties.

The following will give us all of our buckets as a JSON object.

```
curl http://localhost:8098/riak?buckets=true
{"buckets":["food"]}
```

And this will give us all of our keys under the `food` bucket.

```
curl http://localhost:8098/riak/food?keys=true
{
  ...
  "keys": [
    "favorite"
  ]
}
```

If we had very many keys, clearly this might take a while. So Riak also provides the ability to stream your list of keys. `keys=stream` will keep the connection open, returning results in chunks of arrays. When it has exhausted its list, it will close the connection. You can see the details through curl in verbose (`-v`) mode (much of that response has been stripped out below).

```
curl -v http://localhost:8098/riak/food?list=stream
...
* Connection #0 to host localhost left intact
...
{"keys":["favorite"]}
{"keys":[]}
* Closing connection #0
```

You should note that list actions should *not* be used in production (they're really expensive operations). But they are useful for development, investigations, or for running occasional analytics at off-peak hours.

## Buckets

Although we've been using buckets as namespaces up to now, they are capable of more.

Different use-cases will dictate whether a bucket is heavily written to, or largely read from. You may use one bucket to store logs, one bucket could store session data, while another may store shopping cart data. Sometimes low latency is important, while other times it's high durability. And sometimes we just want buckets to react differently when a write occurs.

## Quorum

The basis of Riak's availability and tolerance is that it can read from, or write to, multiple nodes. Riak allows you to adjust these N/R/W values (which we covered under Concepts) on a per-bucket basis.

### N/R/W

N is the number of total nodes that a value should be replicated to, defaulting to 3. But we can set this `n_val` to any number fewer than the total number of nodes.

Any bucket property, including `n_val`, can be set by sending a `props` value as a JSON object to the bucket URL. Let's set the `n_val` to 5 nodes, meaning that objects written to `cart` will be replicated to 5 nodes.

```
curl -i -XPUT http://localhost:8098/riak/cart \
  -H "Content-Type: application/json" \
  -d {"props":{"n_val":5}}
```

You can take a peek at the bucket's properties by issuing a GET to the bucket.

*Note: Riak returns unformatted JSON. If you have a command-line tool like `jsonpp` (or `json_pp`) installed, you can pipe the output there for easier reading. The results below are a subset of all the props values.*

```
curl http://localhost:8098/riak/cart | jsonpp
{
  "props": {
    ...
    "dw": "quorum",
    "n_val": 5,
    "name": "cart",
    "postcommit": [],
    "pr": 0,
    "precommit": [],
    "pw": 0,
    "r": "quorum",
    "rw": "quorum",
```



```

    "w": "quorum",
    ...
  }
}

```

As you can see, `n_val` is 5. That's expected. But you may also have noticed that the `props` returned both `r` and `w` as `quorum`, rather than a number. So what is a *quorum*?

**Symbolic Values** A *quorum* is a one more than half of all the total replicated nodes ( $\text{floor}(N/2) + 1$ ). This is an important figure, since if more than half of all nodes are written to, and more than half of all nodes are read from, then you will get the most recent value (under normal circumstances).

Here's an example with the above `n_val` of 5 (`{A,B,C,D,E}`). Your `w` is a quorum (which is 3, or  $\text{floor}(5/2)+1$ ), so a PUT may respond successfully after writing to `{A,B,C}` (`{D,E}` will eventually be replicated to). Immediately after, a read quorum may GET values from `{C,D,E}`. Even if D and E have older values, you have pulled a value from node C, meaning you will receive the most recent value.

What's important is that your reads and writes *overlap*. As long as  $r+w > n$ , you'll be able to get the newest values. In other words, you'll have consistency.

A quorum is an excellent default, since you're reading and writing from a balance of nodes. But if you have specific requirements, like a log that is often written to, but rarely read, you might find it make more sense to write to a single node, but read from all of them. This affords you an overlap

```

curl -i -XPUT http://localhost:8098/riak/logs \
-H "Content-Type: application/json" \
-d {"props":{"w":"one","r":"all"}}

```

- `all` - All replicas must reply, which is the same as setting `r` or `w` equal to `n_val`
- `one` - Setting `r` or `w` equal to 1
- `quorum` - A majority of the replicas must respond, that is, "half plus one".

**More than R's and W's** Some other values you may have noticed in the bucket's `props` object are `pw`, `pr`, and `dw`.

`pr` and `pw` ensure that many *primary* nodes are available before a read or write. Riak will read or write from backup nodes if one is unavailable, because of network partition or some other server outage. This `p` prefix will ensure that only the primary nodes are used, *primary* meaning the vnode which matches the bucket plus N successive vnodes.

Finally `dw` represents the minimal *durable* writes necessary for success. For a normal `w` write to count a write as successful, a vnode need only promise a write has started, with no guarantee that write has been written to disk, aka, is durable. The `dw` setting means the backend service (for example Bitcask) has agreed to write the value. Although a high `dw`

value is slower than a high *w* value, there are cases where this extra enforcement is good to have, such as dealing with financial data.

**Per Request** It's worth noting that these values (except for *n\_val*) are not actually bound to the bucket, but are instead are just defaults for the bucket. R/W values actually apply *per request*. This is an important distinction, since it allows us to override these values on an as-needed basis: *r*, *pr*, *w*, *pw*, *dw*, *rw*.

Consider you have data that you find very important (say, credit card checkout), and want to help ensure it will be written to every node's disk before success. You could add *?dw=all* to the end of your write.

```
curl -i -XPUT http://localhost:8098/riak/cart/cart1?dw=all \
  -H "Content-Type: application/json" \
  -d {"paid":true}
```

This is a consistent write, since if a single node is not available the write will fail. You wouldn't normally do this since you lose Riak's availability, but it's nice to have the option.

## Hooks

Another utility of buckets are their ability to enforce behaviors on writes by way of hooks. You can attach functions to run either before, or after, a value is committed to a bucket.

Functions that run before a write is called pre-commit, and has the ability to cancel a write altogether if the incoming data is considered bad in some way. A simple pre-commit hook is to check if a value exists at all.

I put my custom files under the riak installation `./custom/my_validators.erl`.

```
-module(my_validators).
-export([value_exists/1]).

%% Object size must be greater than 0 bytes
value_exists(RiakObject) ->
  case erlang:byte_size(riak_object:get_value(RiakObject)) of
    Size when Size == 0 ->
      {fail, "A value sized greater than 0 is required"};
    _ -> RiakObject
  end.
```

Then compile the file.

```
erlc my_validators.erl
```

Install the file by informing the Riak installation of your new code in `app.config` (restart Riak).

```
{riak_kv,
  ...,
  {add_paths, ["/custom"]}}
}
```

All you need to do is set the Erlang module and function as a JSON value to the bucket's pre-commit array {"mod":"my\_validators","fun":"value\_exists"}.

```
curl -i -XPUT http://localhost:8098/riak/cart \
  -H "Content-Type:application/json" \
  -d {"props":{"precommit":[{"mod":"my_validators","fun":"value_exists"}]}}
```

If you try and post to the `cart` bucket without a value, you should expect a failure.

```
curl -XPOST http://localhost:8098/riak/cart \
  -H "Content-Type:application/json"
A value sized greater than 0 is required
```

You can also write pre-commit functions in Javascript, though Erlang code will execute faster.

Post-commits are similar in form and function, but they react after a commit has occurred.

## Entropy

Entropy is a byproduct of eventual consistency. In other words: although eventual consistency says a write will replicate to other nodes in time, there can be a bit of delay while all nodes do not contain the same value.

That difference is *entropy*, and so Riak has created several *anti-entropy* strategies (also called *AE*). We've already talked about how an R/W quorum can deal with differing values when write/read requests overlap at least one node. Riak can repair entropy, or allow you the option to do so yourself.

Riak has a couple strategies related to the case of nodes that do not agree on value.

### Last Write Wins

The most basic, and least reliable, strategy for curing entropy is called *last write wins*. It's the simple idea that the last write based on a node's system clock will overwrite an older one. You can turn this on in the bucket by setting the `last_write_wins` property to `true`.

Realistically, this exists for speed and simplicity, when you really don't care about true order of operations, or the slight possibility of losing data. Since it's impossible to keep server clocks truly in sync (without the proverbial geosynchronized atomic clocks), this is a best guess as to what "last" means, to the nearest millisecond.

## Vector Clock

As we saw under Concepts, *vector clocks* are Riak's way of tracking a true sequence of events of an object. We went over the concept, but let's see how and when Riak vclocks are used.

Every client needs its own id, sent on every write as X-Riak-ClientId.

## Siblings

*Siblings* occur when you have conflicting values, with no clear way for Riak to know which value is correct. Riak will try and resolve these conflicts itself, however, you can instead choose for Riak to create siblings if you set a bucket's `allow_mult` property to `true`.

```
curl -i -XPUT http://localhost:8098/riak/cart \
  -H "Content-Type:application/json" \
  -d {"props":{"allow_mult":true}}
```

Siblings arise in a couple cases.

1. A client writes a value passing a stale vector clock, or missing one altogether.
2. Two clients write at the same time with two different client IDs with the same vector clock value.

We'll use the second case to manufacture our own conflict.

## Creating an Example Conflict

Imagine a shopping cart exists for a single refrigerator, but several people in a household are able to order food for it.

First casey (a vegan) places 10 orders of kale in his cart. To track who is adding to the refrigerator with ID `fridge-97207`, his PUT adds his name as a client id.

Casey writes `[{"item":"kale","count":10}]`.

```
curl -i -XPUT http://localhost:8098/riak/cart/fridge-97207?returnbody=true \
  -H Content-Type:application/json \
  -H X-Riak-ClientId:casey \
  -d [{"item":"kale","count":20}]
```

HTTP/1.1 200 OK

X-Riak-Vclock: a85hYGBgzGDKBVIcypz/fgaUHjmTwZTImMfKsMKK7RRfFgA=

Vary: Accept-Encoding

Server: MochiWeb/1.1 WebMachine/1.9.0 (someone had painted...

Link: </riak/cart>; rel="up"

Last-Modified: Thu, 01 Nov 2012 00:13:28 GMT

```
ETag: "2IGTrV8g1NXEfkPZ45WfAP"
Date: Thu, 01 Nov 2012 00:13:28 GMT
Content-Type: application/json
Content-Length: 28
```

```
[{"item":"kale","count":10}]
```

His roommate mark, reads what's in the order, and updates with his own addition. In order for Riak to know the order of operations, Mark adds the most recent vector clock to his PUT.

Mark writes [{"item":"kale","count":10},{"item":"milk","count":1}].

```
curl -i -XPUT http://localhost:8098/riak/cart/fridge-97207?returnbody=true \
  -H Content-Type:application/json \
  -H X-Riak-ClientId:mark \
  -H X-Riak-Vclock:a85hYGBgzGDKBVlcpz/fgaUHjmTwZTImMfKsMKK7RRfFgA= \
  -d [{"item":"kale","count":20},{"item":"milk","count":1}]
HTTP/1.1 200 OK
X-Riak-Vclock: a85hYGBgzGDKBVlcpz/fgaUHjmTwZTIIMfKcMaK7RRfFgA=
Vary: Accept-Encoding
Server: MochiWeb/1.1 WebMachine/1.9.0 (someone had painted...
Link: </riak/cart>; rel="up"
Last-Modified: Thu, 01 Nov 2012 00:14:04 GMT
ETag: "62NRijQH3mRYPrybFneZaY"
Date: Thu, 01 Nov 2012 00:14:04 GMT
Content-Type: application/json
Content-Length: 54
```

```
[{"item":"kale","count":10},{"item":"milk","count":1}]
```

If you look closely, you'll notice that the vclock sent is not actually identical with the one returned.

- a85hYGBgzGDKBVlcpz/fgaUHjmTwZTImMfKsMKK7RRfFgA=
- a85hYGBgzGDKBVlcpz/fgaUHjmTwZTIIMfKcMaK7RRfFgA=

The vclock was incremented to keep track of the change over time.

Now let's add a third roommate, andy, who loves almonds. Before Mark updates the shared cart with milk, Andy adds his own order to Casey's kale, using the vector clock from Casey's order (the last order Andy was aware of).

Andy writes [{"item":"kale","count":10},{"item":"almonds","count":12}].

```
curl -i -XPUT http://localhost:8098/riak/cart/fridge-97207?returnbody=true \
  -H Content-Type:application/json \
  -H X-Riak-ClientId:andy \
```

```
-H X-Riak-Vclock:a85hYGBgzGDKBVicypz/fgaUHjmTwZTImMfKsMKK7RRfFgA= \
-d [{"item":"kale","count":20},{"item":"almonds","count":12}]
HTTP/1.1 300 Multiple Choices
X-Riak-Vclock: a85hYGBgzGDKBVicypz/fgaUHjmTwZTImMfKoG7LdoovCwA=
Vary: Accept-Encoding
Server: MochiWeb/1.1 WebMachine/1.9.0 (someone had painted...
Last-Modified: Thu, 01 Nov 2012 00:24:07 GMT
ETag: "54Nx22W9M7JUKJnLBrRehj"
Date: Thu, 01 Nov 2012 00:24:07 GMT
Content-Type: multipart/mixed; boundary=Ql300enxVdaMF3YlXF0dm05bvrs
Content-Length: 491
```

```
--Ql300enxVdaMF3YlXF0dm05bvrs
Content-Type: application/json
Link: </riak/cart>; rel="up"
Etag: 62NRijQH3mRYPRybFneZaY
Last-Modified: Thu, 01 Nov 2012 00:14:04 GMT
```

```
[{"item":"kale","count":20},{"item":"milk","count":1}]
--Ql300enxVdaMF3YlXF0dm05bvrs
Content-Type: application/json
Link: </riak/cart>; rel="up"
Etag: 7kfvPXisoVBfC43IiPKYNb
Last-Modified: Thu, 01 Nov 2012 00:24:07 GMT
```

```
[{"item":"kale","count":20},{"item":"almonds","count":12}]
--Ql300enxVdaMF3YlXF0dm05bvrs--
```

Whoa! What's all that?

Since there was a conflict between what Mark and Andy both set the fridge value to be, Riak kept both values.

## V-Tag

Since we're using the HTTP client, Riak returned a 300 Multiple Choices code with a multipart/mixed mime type. It's up to you to separate the results, however, you can request a specific value by it's Etag, also called a Vtag.

Issuing a plain get on the /cart/fridge-97207 key will also return the vtags of all siblings.

```
curl http://localhost:8098/riak/cart/fridge-97207
Siblings:
62NRijQH3mRYPRybFneZaY
```

```
7kfvPXisoVBfC43IiPKYNb
```

What can you do with this tag? Namely, you request the value of a specific sibling by its `vtag`. To get the first sibling in the list (Mark's milk):

```
curl http://localhost:8098/riak/cart/fridge-97207?vtag=62NRijQH3mRYPRybFneZaY
[{"item":"kale","count":20}, {"item":"milk","count":1}]
```

If you want to retrieve all sibling data, tell Riak that you'll accept the multipart message by adding `-H "Accept:multipart/mixed"`.

```
curl -i -XPUT http://localhost:8098/riak/cart/fridge-97207 \
-H Accept:multipart/mixed
```

---

## Use-Case Specific?

When siblings are created, it's up to the application to know how to deal with the conflict. In our example, do we want to accept only one of the orders? Should we remove both milk and almonds and only keep the kale? Should we calculate the cheaper of the two and keep the cheapest option? Should we merge all of the results into a single order? This is why we asked Riak not to resolve this conflict automatically... we want this flexibility.

---

## Resolving Conflicts

With all of our options available, we want to resolve this conflict. Since the way of resolving this conflict is largely *use-case specific*, our application can decide how to proceed.

Let's merge the values into a single result set, taking the larger *count* if the *item* is the same. Pass in the *vclock* of the multipart object, so Riak knows you're resolving the conflict, and you'll get back a new vector clock.

Successive reads will receive a single (merged) result.

```
curl -i -XPUT http://localhost:8098/riak/cart/fridge-97207?returnbody=true \
-H Content-Type:application/json \
-H X-Riak-Vclock:a85hYGBgzGDKBVlcpz/fgaUHjmTwZTInMfKoG7LdoovCwA= \
-d [{"item":"kale","count":20}, {"item":"milk","count":1}, {"item":"almonds","count":12}]
```

It's worth noting that you should never set both `allow_multi` and `last_write_wins` to `true`. It can cause undefined behavior.

## Read Repair

When a successful read happens, but not all replicas agree upon the value, this triggers a *read repair*. This means that Riak will update the replicas with the most recent value. This can happen when, either an object is not found (the vnode has no copy), or a vnode contains an older value (older means that it is an ancestor of the newest vector clock). Unlike `last_write_wins` or manual conflict resolution, read repair is (obviously, I hope, by the name) triggered by a read, rather than a write.

If your nodes get out of sync (for example, if you increase the `n_val` on a bucket), you can force read repair by performing a read operation for all of that bucket's keys. They may return with `not found` the first time, will pull the newest values. Successive reads should now work.

## Querying

So far we've only dealt with key-value lookups. The truth is, key-value is a pretty powerful mechanism that spans a spectrum of use-cases. However, sometimes we need to lookup data by value, rather than key. Sometimes we need to perform some calculations, or aggregations, or search.

## Secondary Indexing (2i)

A *secondary index* (2i) is a data structure that lowers the cost of finding non-key values. Like many other databases, Riak has the ability to index data. However, since Riak has no real knowledge of the data it stores (they're just binary values), it uses metadata to index defined by a name pattern to be either integers or binary values.

If your installation is configured to use 2i (shown in the next chapter), simply writing a value to Riak with the header will be indexes, provided it's prefixed by `X-Riak-Index-` and suffixed by `_int` for an integer, or `_bin` for a string.

```
curl -i -XPUT http://localhost:8098/riak/people/casey \
  -H Content-Type:application/json \
  -H X-Riak-Index-age_int:31 \
  -H X-Riak-Index-fridge_bin:fridge-97207 \
  -d {"work":"rodeo clown"}
```

Querying can be done in two forms: exact match and range. Add a couple more people and we'll see what we get: mark is 32, and andy is 35, they both share `fridge-97207`.

What people own `fridge-97207`? It's a quick lookup to receive the keys that have matching index values.

```
curl http://localhost:8098/buckets/people/index/fridge_bin/fridge-97207
{"keys":["mark","casey","andy"]}
```



With those keys it's a simple lookup to get the bodies.

The other query option is an inclusive ranged match. This finds all people under the ages of 32, by searching between 0 and 32.

```
curl http://localhost:8098/buckets/people/index/age_int/0/32
{"keys":["mark","casey"]}
```

That's about it. It's a basic form of 2i, with a decent array of utility.

## MapReduce/Link Walking

MapReduce is a method of aggregating large amounts of data by separating the processing into two phases, map and reduce, that themselves are executed in parts. Map will be executed per object to convert/extract some value, then those mapped values will be reduced into some aggregate result. What do we gain from this structure? It's predicated on the idea that it's cheaper to move the algorithms to where the data lives, than to transferring massive amounts of data to a single server to run a calculation.

This method, popularized by Google, can be seen in a wide array of NoSQL databases. In Riak, you execute a mapreduce job on a single node, which then propagates to the other nodes. The results are mapped and reduced, then further reduced down to the calling node and returned.

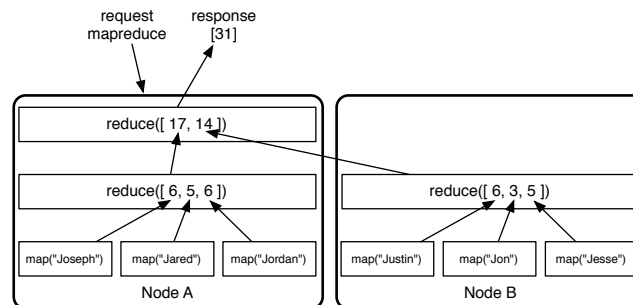


Figure 3.1: MapReduce Returning Name Char Count

Let's assume we have a bucket for log values that stores messages prefixed by either INFO or ERROR. We want to count the number of INFO logs that contain the word "cart".

```
curl -XPOST http://localhost:8098/riak/logs -d INFO: New user added
curl -XPOST http://localhost:8098/riak/logs -d INFO: Kale added to shopping cart
curl -XPOST http://localhost:8098/riak/logs -d INFO: Milk added to shopping cart
curl -XPOST http://localhost:8098/riak/logs -d ERROR: shopping cart cancelled
```

MapReduce jobs can be either Erlang or Javascript code. This time we'll go the easy route and write Javascript. You execute mapreduce by posting JSON to the /mapred path.

```
curl -XPOST http://localhost:8098/mapred \
-H Content-Type: application/json \
```

```

-d @- \
<<EOF
{
  "inputs":"logs",
  "query":[{"
    "map":{
      "language":"javascript",
      "source":"function(riakObject, keydata, arg) {
        var m = riakObject.values[0].data.match(/^INFO.*cart/);
        return [(m ? m.length : 0)];
      }"
    },
    "reduce":{
      "language":"javascript",
      "source":"function(values, arg){
        return [values.reduce(
          function(total, v){ return total + v; }, 0)
        ];
      }"
    }
  ]
}
EOF

```

The result should be `[2]`, as expected. Both map and reduce phases should always return an array. The map phase receives a single riak object, while the reduce phase received an array of values, either the result of multiple map function outputs, or of multiple reduce outputs. I probably cheated a bit by using Javascript's `reduce` function to sum the values, but, well, welcome to the world of thinking in terms of mapreduce!

## Key Filters

Besides executing a map function against every object in a bucket, you can reduce the scope by using *key filters*. Just as it sounds, they are a way of only including those objects that match a pattern... it filters out certain keys.

Rather than passing in a bucket name as a value for `inputs`, instead we pass it a JSON object containing the `bucket` and `key_filters`. The `key_filters` get an array describing how to transform then test each key in the bucket. Any keys that match the predicate will be passed into the map phase, all others are just filtered out.

To get all keys in the `cart` bucket that end with a number greater than 97000, you could tokenize the keys on `-` (remember we used `fridge-97207`) and keep the second half of the string, convert it to an integer, then compare that number to be greater than 97000.

```
"inputs":{
  "bucket":"cart",
  "key_filters":[["tokenize", "-", 2],["string_to_int"],["greater_than",97000]]
}
```

It would look like this to have the mapper just return matching object keys. Pay special attention to the `map` function, and lack of `reduce`.

```
curl -XPOST http://localhost:8098/mapred \
  -H Content-Type: application/json \
  -d @- \
<<EOF
{
  "inputs":{
    "bucket":"cart",
    "key_filters":[
      ["tokenize", "-", 2],
      ["string_to_int"],
      ["greater_than",97000]
    ]
  },
  "query":[{"
    "map":{
      "language":"javascript",
      "source":"function(riakObject, keydata, arg) {
        return [riakObject.key];
      }"
    }
  }]
}
EOF
```

## MR + 2i

Another option when using mapreduce is to combine it with secondary indexes. You can pipe the results of a 2i query into a mapreducer, simply specify the index you wish to use, and either a key for an index lookup, or start and end values for a ranged query.

```
...
"inputs":{
  "bucket":"people",
  "index": "age_int",
  "start": 18,
```

```

    "end": 32
  },
  ...

```

## Link Walking

Conceptually, a link is a one-way relationship from one object to another. *Link walking* is a convenient query option for retrieving data when you start with the object linked from.

Let's add a link to our people, by setting `casey` as the brother of `mark` using the HTTP header `Link`.

```

curl -XPUT http://localhost:8098/riak/people/mark \
  -H Content-Type:application/json \
  -H Link: </riak/people/casey>; riaktag="brother"

```

With a Link in place, now it's time to walk it. Walking is like a normal request, but with the suffix of `/[bucket],[riaktag],[keep]`. In other words, the *bucket* a possible link points to, the value of a *riaktag*, and whether to *keep* the results of this phase (only useful when chaining link walks). Each of these values can be set to a wildcard `_`, meaning you don't care the value.

```

curl http://localhost:8098/riak/people/mark/people,brother,_
--8wuTE7VSpvHlAJ06XovIrGFGalP
Content-Type: multipart/mixed; boundary=991Bi7WVpjYAGUwZlMfJ4nPJR0w

--991Bi7WVpjYAGUwZlMfJ4nPJR0w
X-Riak-Vclock: a85hYGBgzGDKBVIcypz/fgZMzorIYEpkz2NlWCzKcYovCwA=
Location: /riak/people/casey
Content-Type: application/json
Link: </riak/people>; rel="up"
Etag: Wf02eljDiBa5q5nSbTq2s
Last-Modified: Fri, 02 Nov 2012 10:00:03 GMT
x-riak-index-age_int: 31
x-riak-index-fridge_bin: fridge-97207

{"work":"rodeo clown"}
--991Bi7WVpjYAGUwZlMfJ4nPJR0w--

--8wuTE7VSpvHlAJ06XovIrGFGalP--

```

Even without returning the `Content-Type`, this kind of body should look familiar. Link walking always returns a `multipart/mixed`, since a single key can contain any number of links, meaning any number of objects returned.

It gets crazier. You can actually chain together link walks, which will follow the a followed link. If `casey` has links, they can be followed by tacking another link triplet on the end, like so:

```
curl http://localhost:8098/riak/people/mark/people,brother,0/_,_,_
```

Now it may not seem so from what we've seen, but link walking is a specialized case of mapreduce.

There is another phase of a mapreduce query called "link". Rather than executing a function, however, it only requires the same configuration that you pass through the shortcut URL query.

```
...
"query": [{
  "link": {
    "bucket": "people",
    "tag": "brother",
    "keep": false
  }
}]
...
```

As we've seen, mapreduce in Riak is a powerful way of pulling data out of an otherwise straight key/value store. But we have one more method of finding data in Riak.

---

## What Happened to Riak Search?

If you have used Riak before, or have a hold of some older documentation, you may wonder what the difference is between Riak Search and Yokozuna.

In an attempt to make Riak Search user friendly, it was originally developed with a "Solr like" interface. Sadly, due to the complexity of building distributed search engines, it was woefully incomplete. Basho decided that, rather than attempting to maintain parity with Solr, a popular and featureful search engine in it's own right, it made more sense to integrate the two.

---

## Search (Yokozuna)

*Note: This is covering a project still under development. Changes are to be expected, so please refer to the [project page](#) for the most recent information.*

Yokozuna is an extension to Riak that lets you perform searches to find data in a Riak cluster. Unlike the original Riak Search, Yokozuna leverages distributed Solr to perform the inverted indexing and management of retrieving matching values.

Before using Yokozuna, you'll have to have it installed and a bucket set up with an index (these details can be found in the next chapter).

The simplest example is a full-text search. Here we add `ryan` to the `people` table (with a default index).

```
curl -XPUT http://localhost:8091/riak/people/ryan \
  -H Content-Type:text/plain \
  -d Ryan Zezeski
```

To execute a search, request `/search/[bucket]` along with any distributed [Solr parameters](#). Here we query for documents that contain a word starting with `zez`, request the results to be in json format (`wt=json`), only return the Riak key (`fl=_yz_rk`).

```
curl http://localhost:8091/search/people?wt=json&omitHeader=true&fl=_yz_rk&q=zez* | jsonpp
{
  "response": {
    "numFound": 1,
    "start": 0,
    "maxScore": 1.0,
    "docs": [
      {
        "_yz_rk": "ryan"
      }
    ]
  }
}
```

With the matching `_yz_rk` keys, you can retrieve the bodies with a simple Riak lookup.

Yokozuna supports Solr 4.0, which includes filter queries, ranges, page scores, start values and rows (the last two are useful for pagination). You can also receive snippets of matching [highlighted text](#) (`hl,hl.fl`), which is useful for building a search engine (and something we use for [search.basho.com](#)).

## Tagging

Another useful feature of Solr and Yokozuna is the tagging of values. Tagging values give additional context to a Riak value. The current implementation requires all tagged values begin with `X-Riak-Meta`, and be listed under a special header named `X-Riak-Meta-yz-tags`.

```
curl -XPUT http://localhost:8091/riak/people/dave \
  -H Content-Type:text/plain \
  -H X-Riak-Meta-yz-tags: X-Riak-Meta-nickname_s \
  -H X-Riak-Meta-nickname_s:dizzy \
  -d Dave Smith
```

To search by the `nickname_s` tag, just prefix the query string followed by a colon.

```
curl http://localhost:8091/search/people?wt=json&omitHeader=true&q=nickname_s:dizzy | jsonpp
{
  "response": {
```

```
"numFound": 1,
"start": 0,
"maxScore": 1.4054651,
"docs": [
  {
    "nickname_s": "dizzy",
    "id": "dave_25",
    "_yz_ed": "20121102T215100 dave m7psMIomLMu/+dtWx51Kluvvrb8=",
    "_yz_fpn": "23",
    "_yz_node": "dev1@127.0.0.1",
    "_yz_pn": "25",
    "_yz_rk": "dave",
    "_version_": 1417562617478643712
  }
]
}
```

Notice that the docs returned also contain "nickname\_s":"dizzy" as a value. All tagged values will be returned on matching results.

*Expect more features to appear as Yokozuna gets closer to a final release.*

## Wrapup

Riak is a distributed data store with several additions to improve upon the standard key-value lookups, like specifying replication values. Since values in Riak are opaque, many of these methods either: require custom code to extract and give meaning to values, such as *mapreduce*; or allow for header metadata to provide an added descriptive dimension to the object, such as *secondary indexes*, *link walking*, or *search*.

Next we'll peek further under the hood, and see how to set up and manage a cluster of your own, and what you should know.





## Chapter 4

# Operators

In some ways, Riak is downright mundane in its role as the easiest NoSQL database to operate. Want more servers? Add them. A network cable gets cut at 2am? Sleep until morning and fix it. But the fact that it generally hums without a hiccup, does not diminish the importance of understanding this integral part of your application stack.

We've covered the core concepts of Riak, and I've provided a taste of how to go about using Riak, but there is more to Riak than that. There are details you should know if you plan on operating a Riak cluster of your own.

## Clusters

Up to this point you've conceptually read about "clusters" and the "Ring" in nebulous summations. What exactly do we mean, and what are the practical implications of these details for Riak developers and operators?

A *cluster* in Riak is a managed collection of nodes that share a common Ring.

## The Ring

*The Ring* in Riak is actually a two-fold concept.

Firstly, the Ring is a function of the consistent hash space partitions, managed by vnodes. This partition range is treated as circular, from 0 to  $2^{160}-1$  back to 0 again. (If you're wondering, yes this means that we are limited to  $2^{160}$  nodes, which is a limit of a 1.46 quindecillion, or  $1.46 \times 10^{48}$ , node cluster. For comparison, there are only  $1.92 \times 10^{49}$  [silicon atoms on Earth](#).)

When we consider replication, the N value defines how many nodes an object is replicated to. Riak makes a best attempt at spreading that value to as many nodes as it can, so it copies to the next N adjacent nodes, starting with the primary partition and counting around the Ring, if it reaches the last partition, it starts over at the first one.

Secondly, the Ring is also used as a shorthand for describing the state of the circular hash ring I just mentioned. This Ring (aka *Ring State*) is a data structure that gets passed around

between nodes, so each knows the state of the entire cluster. Which node manages which vnodes? If a node gets a request for an object managed by other nodes, it consults the Ring and forwards on the request to the proper nodes. It's a local copy of a contract that all of the nodes agree to follow.

Obviously, this contract needs to stay in sync between all of the nodes. If a node is permanently taken offline or a new one added, the other nodes need to readjust, balancing the partitions around the cluster, then updating the Ring with this new structure. This Ring state gets passed between the nodes by means of a *gossip protocol*.

## Gossip

The *gossip protocol* is Riak's method of keeping all nodes current on the state of the Ring. If a node goes up or down, that information is propagated to other nodes. Periodically, nodes will also send their status to a random peer for added consistency.

Propagating changes in Ring is an asynchronous operation, and can take a couple minutes depending on Ring size.

Currently, the ability to change the total number of vnodes of a cluster is not feasible. This means that you *must have an idea of how large you want your cluster to grow in a single datacenter*. Although a basic install starts with 64 vnodes, if you plan any cluster larger than 6 or so servers you should increase vnodes to 256 or 1024.

The number of vnodes must be a power of 2 (eg. 64, 256, 1024, etc).

## How Replication Uses the Ring

Even if you are not a coder, it's worth taking a look at this Ring example. It's also worth remembering that partitions are managed by vnodes, and in conversation are sometimes interchanged, though I'll try and be more precise here.

Let's start with Riak configured to have 8 vnodes/partitions, which are set via `ring_creation_size` in the `etc/app.config` file (we'll dig deeper into this file later).

```
%% Riak Core config
{riak_core, [
    ...
    {ring_creation_size, 8},
```

In this example, I have a total of four Riak nodes running on A@10.0.1.1, B@10.0.1.2, C@10.0.1.3, and D@10.0.1.4.

The following is Erlang, the language Riak was written in. Riak has the amazing, and probably dangerous command `attach`, that attaches an Erlang console to a live Riak node, loaded with all of the Riak modules. So here we're getting a copy of the Ring from the locally running node.

The `riak_core_ring:chash(Ring)` function extracts the total count of partitions (8), with an array of numbers representing the start of the partition, some fraction of the  $2^{160}$  number, and the node name that represents a particular Riak server in the cluster.

```
$ bin/riak attach
(A@10.0.1.1)1> {ok,Ring} = riak_core_ring_manager:get_my_ring().
(A@10.0.1.1)2> riak_core_ring:chash(Ring).
{8,
 [{0,'A@10.0.1.1'},
  {182687704666362864775460604089535377456991567872, 'B@10.0.1.2'},
  {365375409332725729550921208179070754913983135744, 'C@10.0.1.3'},
  {548063113999088594326381812268606132370974703616, 'D@10.0.1.4'},
  {730750818665451459101842416358141509827966271488, 'A@10.0.1.1'},
  {913438523331814323877303020447676887284957839360, 'B@10.0.1.2'},
  {1096126227998177188652763624537212264741949407232, 'C@10.0.1.3'},
  {1278813932664540053428224228626747642198940975104, 'D@10.0.1.4'}]}
```

To discover which partition the bucket/key `food/favorite` object would be stored in, for example, we execute `riak_core_util:chash_key({<<"food">>, <<"favorite">>})` and get a wacky 160 bit Erlang number we named `DocIdx` (document index).

Just to illustrate that Erlang binary value is a real number, the next line makes it a more readable format, similar to the ring partition numbers.

```
(A@10.0.1.1)3> DocIdx = riak_core_util:chash_key({<<"food">>, <<"favorite">>}).
<<80,250,1,193,88,87,95,235,103,144,152,2,21,102,201,9,156,102,128,3>>

(A@10.0.1.1)4> <<I:160/integer>> = DocIdx. I.
462294600869748304160752958594990128818752487427
```

With this `DocIdx` number, we can order the partitions, starting with first number greater than `DocIdx`. The remaining partitions are in numerical order, until we reach zero, then we loop around and continue to exhaust the list.

```
(A@10.0.1.1)5> Preflist = riak_core_ring:preflist(DocIdx, Ring).
[{548063113999088594326381812268606132370974703616, 'D@10.0.1.4'},
 {730750818665451459101842416358141509827966271488, 'A@10.0.1.1'},
 {913438523331814323877303020447676887284957839360, 'B@10.0.1.2'},
 {1096126227998177188652763624537212264741949407232, 'C@10.0.1.3'},
 {1278813932664540053428224228626747642198940975104, 'D@10.0.1.4'},
```

```
{0, 'A@10.0.1.1'},  
{182687704666362864775460604089535377456991567872, 'B@10.0.1.2'},  
{365375409332725729550921208179070754913983135744, 'C@10.0.1.3' }]
```

So what does all this have to do with replication? With the above list, we simply replicate a write down the list *N* times. If we set *N*=3, then the `food/favorite` object will be written to the `D@10.0.1.4` node's partition 5480631... (I truncated the number here), `A@10.0.1.1` partition 7307508..., and `B@10.0.1.2` partition 9134385....

If something has happened to one of those nodes, like a network split (confusingly also called a partition—the “P” in “CAP”), the remaining active nodes in the list become candidates to hold the data.

So if the partition 7307508... write could not connect to node 10.0.1.1, Riak would then attempt to write that partition 7307508... to `C@10.0.1.3` as a fallback (it's the next node in the list prelist after the 3 primaries).

Due to the structure of the Ring, how it distributes partitions, and how it handles failures, it's relatively simple to ensure that data is replicated to as many physical nodes as possible, while being able to remain operational if a node is unavailable, by simply trying the next available node in the list.

## Hinted Handoff

When a node goes down data is replicated to a backup node. But this is not a solution, merely a band aid. So Riak will periodically trigger vnodes to check if they reside on the correct node (according to the Ring). If not, the managing process will attempt to connect with the home node, and if that node responds, will hand off any data it hold back to proper node.

As long as the temporary node cannot connect to the primary, it will continue to access writes and reads on behalf of its incapacitated brethren.

High availability is not the only purpose of hinted handoff. In the case where the ring changes, because a node was added or removed, data must be transferred to its new home. In this case, the same thing will happen: a vnode checks if it's in the correct place, and if not, attempts to transfer its data to its new home node.

## Managing a Cluster

Now that we have a grasp of the general concepts of Riak, how users query it, and how Riak manages replication, it's time to build a cluster. It's so easy to do, in fact, I didn't bother discussing it for most of this book.

## Install

The Riak docs have all of the information you need to [install](#) it per operating system. The general sequence is:

1. Install Erlang
2. Get Riak from a package manager (ala apt-get or Homebrew), or build from source (the results end up under `rel/riak`, with the binaries under `bin`).
3. Run `riak start`

Install Riak on four or five nodes—five being the recommended safe minimum for production.

## Command Line

Most Riak operations can be performed through the command line. We'll concern ourselves with two: `riak` and `riak-admin`.

### **riak**

Simply typing the `riak` command will give a usage list, although not a terribly descriptive one.

```
Usage: riak {start|stop|restart|reboot|ping|console|attach|chkconfig|escript|version}
```

Most of these commands are self explanatory, once you know what they mean. `start` and `stop` are simple enough. `restart` means to stop the running node and restart it inside of the same Erlang VM (virtual machine), while `reboot` will take down the Erlang VM and restart everything.

You can print the current running version. `ping` will return `pong` if the server is in good shape, otherwise you'll get the *just-similar-enough-to-be-annoying* response `pang` (with an *a*), or a simple `Node *X* not responding to pings` if it's not running at all.

`chkconfig` is useful if you want to ensure your `etc/app.config` is not broken (that is to say, it's parsable). I mentioned `attach` briefly above, when we looked into the details of the Ring—it attaches a console to the local running Riak server so you can execute Riak's Erlang code. And finally, `escript` is similar to `console`, except you pass in script file of commands you wish to run.

### **riak-admin**

The `riak-admin` command is the meat operations, the tool you'll use most often. This is where you'll join nodes to the Ring, diagnose issues, check status, and trigger backups.

```
Usage: riak-admin { cluster | join | leave | backup | restore | test |
                  reip | js-reload | erl-reload | wait-for-service |
                  ringready | transfers | force-remove | down |
                  cluster-info | member-status | ring-status | vnode-status |
                  diag | status | transfer-limit |
                  top [-interval N] [-sort reductions|memory|msg_q] [-lines N] }
```

Many of these commands are deprecated, and many don't make sense without a cluster, but a few we can look at now.

`status` outputs a list of information about this cluster. It's mostly the same information you can get from getting `/stats` via HTTP, although the coverage of information is not exact (for example, `riak-admin status` returns `disk`, and `/stats` returns some computed values like `gossip_received`).

```
$ riak-admin status
1-minute stats for 'A@10.0.1.1'
-----
vnode_gets : 0
vnode_gets_total : 2
vnode_puts : 0
vnode_puts_total : 1
vnode_index_reads : 0
vnode_index_reads_total : 0
vnode_index_writes : 0
vnode_index_writes_total : 0
vnode_index_writes_postings : 0
vnode_index_writes_postings_total : 0
vnode_index_deletes : 0
...
```

Adding javascript or erlang files to Riak (as we did in the developers chapter ) are not automatically found by the nodes, but they instead must be informed by either `js-reload` or `erl-reload` command.

`riak-admin` also provides a little `test` command, so you can perform a read/write cycle to a node, which I find useful for testing a client's ability to connect, and the node's ability to write.

Finally, `top` is an analysis command checking the Erlang details of a particular node in real time. Different processes have different process ids (Pids), use varying amounts of memory, queue up so many messages at a time (MsgQ), and so on. This is useful for advanced diagnostics, and is especially useful if you know Erlang, or seek help from other users, the Riak team, or Basho.

Figure 4.1: Top

## Making a Cluster

With several solitary nodes running—assuming they are networked and are able to communicate to each other—launching a cluster is the simplest part.

Executing the `cluster` command will output a descriptive set of commands.

```
$ riak-admin cluster
```

The following commands stage changes to cluster membership. These commands do not take effect immediately. After staging a set of changes, the staged plan must be committed to take effect:

<code>join &lt;node&gt;</code>	Join node to the cluster containing <node>
<code>leave</code>	Have this node leave the cluster and shutdown
<code>leave &lt;node&gt;</code>	Have <node> leave the cluster and shutdown
<code>force-remove &lt;node&gt;</code>	Remove <node> from the cluster without first handing off data. Designed for crashed, unrecoverable nodes
<code>replace &lt;node1&gt; &lt;node2&gt;</code>	Have <node1> transfer all data to <node2>, and then leave the cluster and shutdown
<code>force-replace &lt;node1&gt; &lt;node2&gt;</code>	Reassign all partitions owned by <node1> to <node2> without first handing off data, and remove <node1> from the cluster.

Staging commands:

<code>plan</code>	Display the staged changes to the cluster
<code>commit</code>	Commit the staged changes
<code>clear</code>	Clear the staged changes

To create a new cluster, you must `join` another node (any will do). Taking a node out of the cluster uses `leave` or `force-remove`, while swapping out an old node for a new one uses `replace` or `force-replace`.

I should mention here that using `leave` is the nice way of taking a node out of commission.

However, you don't always get that choice. If a server happens to explode (or simply smoke ominously), you don't need it's approval to remove it from the cluster, but can instead mark it as down.

But before we worry about removing nodes, let's add some first.

```
$ riak-admin cluster join A@10.0.1.1
Success: staged join request for B@10.0.1.2 to A@10.0.1.1
$ riak-admin cluster join A@10.0.1.1
Success: staged join request for C@10.0.1.3 to A@10.0.1.1
```

---

## Don't Wait Too Long

You should always keep in mind the general pattern Riak follows when you make a change to the cluster:

*Join/Leave/Down -> Ring state change -> Gossiped -> Hinted handoff*

Large amounts of data can take time and cause system strain to transfer, so don't wait until it's too late to grow.

Once all changes are staged, you must review the cluster plan. It will give you all of the details of the nodes that are joining the cluster, and what it will look like after each step or *transition*, including the `member-status`, and how the `transfers` plan to handoff partitions.

Below is a simple plan, but there are cases when Riak requires multiple transitions to enact all of your requested actions, such as adding and removing nodes in one stage.

```
$ riak-admin cluster plan
===== Staged Changes =====
Action      Nodes(s)
-----
join        B@10.0.1.2
join        C@10.0.1.3
-----
```

NOTE: Applying these changes will result in 1 cluster transition

```
#####
                After cluster transition 1/1
#####
```

```
===== Membership =====
Status      Ring      Pending      Node
-----
valid       100.0%     34.4%        A@10.0.1.1
valid       0.0%       32.8%        B@10.0.1.2
```



```
valid      0.0%      32.8%      C@10.0.1.3
-----
Valid:3 / Leaving:0 / Exiting:0 / Joining:0 / Down:0
```

WARNING: Not all replicas will be on distinct nodes

```
Transfers resulting from cluster changes: 42
  21 transfers from A@10.0.1.1 to C@10.0.1.3
  21 transfers from A@10.0.1.1 to B@10.0.1.2
```

Making changes to cluster membership can be fairly resource intensive, so Riak defaults to only performing 2 transfers at a time. You can choose to alter this `transfer-limit` from the `riak-admin`, but bear in mind the higher the number, the greater normal operations will be impinged.

At this point, if you find a mistake in the plan, you have the chance to `clear` it and try again. When you are ready, `commit` the cluster to enact the plan.

```
$ dev1/bin/riak-admin cluster commit
Cluster changes committed
```

Without any data, adding a node to a cluster is a quick operation. However, with large amounts of data to be transferred to a new node, it can take quite a while before the service is available.

## Status Options

To check on a launching node's progress, you can run the `wait-for-service` command. It will output the status of the service and stop when it's finally up. You can get a list of available services through the similarly named command.

```
$ riak-admin wait-for-service riak_kv C@10.0.1.3
riak_kv is not up: []
riak_kv is not up: []
riak_kv is up
```

You can also see if the whole ring is ready to go with `ringready`. If the nodes do not agree on the state of the ring, it will output `FALSE`, otherwise `TRUE`.

```
$ riak-admin ringready
TRUE All nodes agree on the ring ['A@10.0.1.1','B@10.0.1.2',
                                   'C@10.0.1.3']
```

For a more complete view of the status of the nodes in the ring, you can check out `member-status`.

```
$ riak-admin member-status
===== Membership =====
Status      Ring      Pending   Node
-----
valid       34.4%     --        A@10.0.1.1
valid       32.8%     --        B@10.0.1.2
valid       32.8%     --        C@10.0.1.3
-----
Valid:3 / Leaving:0 / Exiting:0 / Joining:0 / Down:0
```

And for more details of any current handoffs or unreachable nodes, try `ring-status`. It also lists some information from `ringready` and transfers. Below I turned off a node to show what it might look like.

```
$ riak-admin ring-status
===== Claimant =====
Claimant:   A@10.0.1.1
Status:     up
Ring Ready: true

===== Ownership Handoff =====
Owner:      dev1 at 127.0.0.1
Next Owner: dev2 at 127.0.0.1

Index: 182687704666362864775460604089535377456991567872
Waiting on: []
Complete:   [riak_kv_vnode,riak_pipe_vnode]
...

===== Unreachable Nodes =====
The following nodes are unreachable: [ C@10.0.1.3 ]
```

WARNING: The cluster state will not converge `until` all nodes are up. Once the above nodes come back online, convergence will continue. If the outages are long-term or permanent, you can either mark the nodes `as` down (`riak-admin down NODE`) or forcibly remove the nodes from the cluster (`riak-admin force-remove NODE`) to allow the remaining nodes to settle.

If all of the above information options about your nodes weren't enough, you can list the status of each vnode per node, via `vnode-status`. It'll show each vnode by its partition number, give any status information, and a count of each vnode's keys. Finally, you'll get to see each vnode's backend type—something I'll cover in the next section.

```
$ riak-admin vnode-status
Vnode status information
-----

VNode: 0
Backend: riak_kv_bitcask_backend
Status:
[{"key_count",0},{status,[]}]

VNode: 91343852333181432387730302044767688728495783936
Backend: riak_kv_bitcask_backend
Status:
[{"key_count",0},{status,[]}]

VNode: 182687704666362864775460604089535377456991567872
Backend: riak_kv_bitcask_backend
Status:
[{"key_count",0},{status,[]}]

VNode: 274031556999544297163190906134303066185487351808
Backend: riak_kv_bitcask_backend
Status:
[{"key_count",0},{status,[]}]

VNode: 365375409332725729550921208179070754913983135744
Backend: riak_kv_bitcask_backend
Status:
[{"key_count",0},{status,[]}]
...
```

Some commands we did not cover are either deprecated in favor of their `cluster` equivalents (`join`, `leave`, `force-remove`, `replace`, `force-replace`), or flagged for future removal (`reip` (use `cluster replace`)).

The last command is `diag`, which requires a [Riaknostic](#) installation to give you more diagnostic tools.

I know this was a lot to digest, and probably pretty dry. Walking through command line tools usually is. There are plenty of details behind many of the `riak-admin` commands, too numerous to cover in such a short book. I encourage you to toy around with them on your own installation.

## How Riak is Built

It's difficult to label Riak a single project. It's probably more correct to think of Riak as the center of gravity for a whole system of projects. As we've covered before, Riak is built on Erlang, but that's not the whole story. It's more correct to say Riak is fundamentally Erlang, with some pluggable native C code components (like leveldb), Java (Yokozuna), and even JavaScript (for Mapreduce or commit hooks).

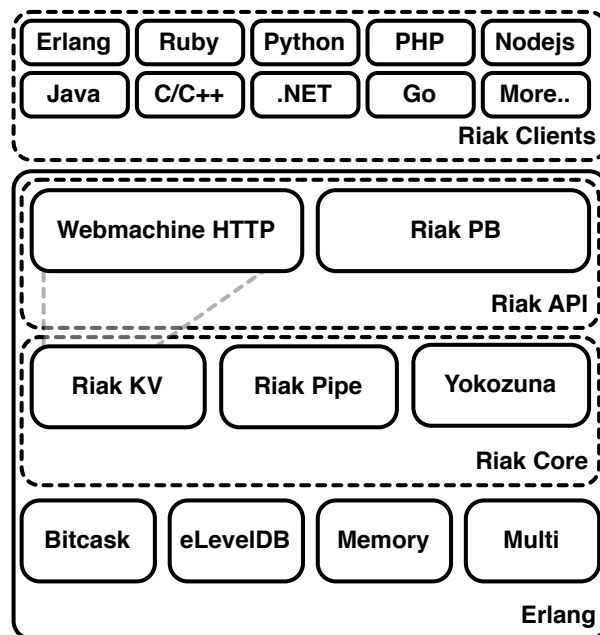


Figure 4.2: Tech Stack

The way Riak stacks technologies is a good thing to keep in mind, in order to make sense of how to configure it properly.

## Erlang

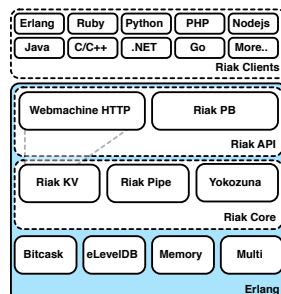


Figure 4.3: Tech Stack Erlang

When you fire up a Riak node, it also starts up an Erlang VM (virtual machine) to run and manage Riak's processes. These include vnodes, process messages, gossips, resource management and more. The Erlang operating system process is found as a `beam.smp` command with many, many arguments.

These arguments are configured through the `etc/vm.args` file. There are a couple setting you should pay special attention to.

```
$ ps -o command | grep beam
/riak/erts-5.9.1/bin/beam.smp \
-K true \
-A 64 \
-W w -- \
-root /riak \
-progname riak -- \
-home /Users/ericredmond -- \
-boot /riak/releases/1.2.1/riak \
-embedded \
-config /riak/etc/app.config \
-pa ./lib/basho-patches \
-name A@10.0.1.1 \
-setcookie testing123 -- \
console
```

The `name` setting is the name of the current Riak node. Every node in your cluster needs a different name. It should have the IP address or dns name of the server this node runs on, and optionally a different prefix—though some people just like to name it *riak* for simplicity (eg: `riak@node15.myhost`).

The `setcookie` parameter is a setting for Erlang to perform inter-process communication (IPC) across nodes. Every node in the cluster must have the same cookie name. I recommend you change the name from `riak` to something a little less likely to accidentally conflict, like `hihohihoitsofftoworkwego`.

My `vm.args` starts with this:

```
## Name of the riak node
-name A@10.0.1.1

## Cookie for distributed erlang. All nodes in the same cluster
## should use the same cookie or they will not be able to communicate.
-setcookie testing123
```

Continuing down the `vm.args` file are more Erlang settings, some environment variables that are set up for the process (prefixed by `-env`), followed by some optional SSL encryption settings.

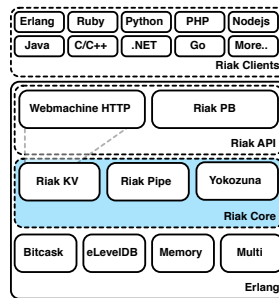


Figure 4.4: Tech Stack Core

## riak\_core

If any single component deserves the title of “Riak proper”, it would be *Riak Core*. Core, and implementations are responsible for managing the partitioned keyspace, launching and supervising vnodes, preference list building, hinted handoff, and things that aren’t related specifically to client interfaces, handling requests, or storage.

Riak Core, like any project, has some hard coded values (for example, how protocol buffer messages are encoded in binary). However, many values can be modified to fit your use-case. The majority of this configuration occurs under `app.config`. This file is Erlang code, so commented lines begin with a `%` character.

The `riak_core` configuration section allows you to change the options in this project. This handles basic settings, like files/directories where values are stored or to be written to, the number of partitions/vnodes in the cluster (`ring_creation_size`), and several port options.

```
%% Riak Core config
{riak_core, [
    %% Default location of ringstate
    {ring_state_dir, "./data/ring"},

    %% Default ring creation size. Make sure it is a power of 2,
    %% e.g. 16, 32, 64, 128, 256, 512 etc
    {ring_creation_size, 64},

    %% http is a list of IP addresses and TCP ports that the Riak
    %% HTTP interface will bind.
    {http, [ {"127.0.0.1", 8098 } ]},

    %% https is a list of IP addresses and TCP ports that the Riak
    %% HTTPS interface will bind.
    {https, [{ "127.0.0.1", 8098 } ]},

    %% Default cert and key locations for https can be overridden
    %% with the ssl config variable, for example:
    {ssl, [
        % {certfile, "./etc/cert.pem"},
    ]}
]}
```

```
%      {keyfile, "./etc/key.pem"}
%      }},

%% riak handoff_port is the TCP port that Riak uses for
%% intra-cluster data handoff.
{handoff_port, 8099 },

%% To encrypt riak_core intra-cluster data handoff traffic,
%% uncomment the following line and edit its path to an
%% appropriate certfile and keyfile. (This example uses a
%% single file with both items concatenated together.)
{handoff_ssl_options, [{certfile, "/tmp/erlserver.pem"}]},

%% Platform-specific installation paths
{platform_bin_dir, "./bin"},
{platform_data_dir, "./data"},
{platform_etc_dir, "./etc"},
{platform_lib_dir, "./lib"},
{platform_log_dir, "./log"}
}],
```

## riak\_kv

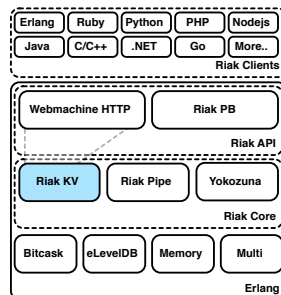


Figure 4.5: Tech Stack KV

Riak KV is the Key/Value implementation of Riak Core. This is where the magic happens, such as handling requests, coordinating them for redundancy and read repair. It's what makes the Riak, as we know it, a KV store rather than something else like a Cassandra-style columnar data store.

HTTP access to KV defaults to the `/riak` path as we've seen in examples throughout the book. This prefix is editable via `raw_name`. Many of the other KV settings are concerned with backward compatibility modes, backend settings, mapreduce, and Javascript integration.

```
%% Riak KV config
{riak_kv, [
  %% raw_name is the first part of all URLs used by the Riak raw HTTP
  %% interface. See riak_web.erl and raw_http_resource.erl for details.
```

```

{raw_name, "riak"},

%% http_url_encoding determines how Riak treats URL encoded
%% buckets, keys, and links over the REST API. When set to on
%% Riak always decodes encoded values sent as URLs and Headers.
%% Otherwise, Riak defaults to compatibility mode where links
%% are decoded, but buckets and keys are not. The compatibility
%% mode will be removed in a future release.
{http_url_encoding, on},

%% Switch to vnode-based vclocks rather than client ids. This
%% significantly reduces the number of vclock entries.
{vnode_vclocks, true},

%% This option toggles compatibility of keylisting with 1.0
%% and earlier versions. Once a rolling upgrade to a version
%% > 1.0 is completed for a cluster, this should be set to true
%% for better control of memory usage during key listing operations
{listkeys_backpressure, true},
...
}],

```

## riak\_pipe

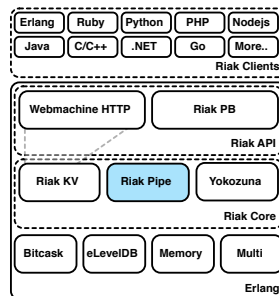


Figure 4.6: Tech Stack Pipe

Riak pipe is an input/output messaging system that forms the basis of Riak's mapreduce. This was not always the case, and MR used to be a dedicated implementation, hence some legacy options. Like the ability to alter the KV path, you can also change HTTP from `/mapred` to a custom path.

```

%% Riak KV config
{riak_kv, [
    %% mapred_name is URL used to submit map/reduce requests to Riak.
    {mapred_name, "mapred"},

    %% mapred_system indicates which version of the MapReduce

```



```

%% system should be used: pipe means riak_pipe will
%% power MapReduce queries, while legacy means that luke
%% will be used
{mapred_system, pipe},

%% mapred_2i_pipe indicates whether secondary-index
%% MapReduce inputs are queued in parallel via their own
%% pipe ( true ), or serially via a helper process
%% ( false or undefined). Set to false or leave
%% undefined during a rolling upgrade from 1.0.
{mapred_2i_pipe, true},

%% directory used to store a transient queue for pending
%% map tasks
%% Only valid when mapred_system == legacy
%% {mapred_queue_dir, "./data/mr_queue" },

%% Number of items the mapper will fetch in one request.
%% Larger values can impact read/write performance for
%% non-MapReduce requests.
%% Only valid when mapred_system == legacy
%% {mapper_batch_size, 5},

%% Number of objects held in the MapReduce cache. These will be
%% ejected when the cache runs out of room or the bucket/key
%% pair for that entry changes
%% Only valid when mapred_system == legacy
%% {map_cache_size, 10000},
...
}]

```

## JavaScript

Though not implemented in pipe, Riak KV's mapreduce implementation is the primary user of the Spidermonkey JavaScript engine—the second use is pre-commit hooks.

```

%% Riak KV config
{riak_kv, [
  ...
  %% Each of the following entries control how many Javascript
  %% virtual machines are available for executing map, reduce,
  %% pre- and post-commit hook functions.
  {map_js_vm_count, 8 },
  {reduce_js_vm_count, 6 },
  {hook_js_vm_count, 2 },

```

```

%% js_max_vm_mem is the maximum amount of memory, in megabytes,
%% allocated to the Javascript VMs. If unset, the default is
%% 8MB.
{js_max_vm_mem, 8},

%% js_thread_stack is the maximum amount of thread stack, in megabytes,
%% allocate to the Javascript VMs. If unset, the default is 16MB.
%% NOTE: This is not the same as the C thread stack.
{js_thread_stack, 16},

%% js_source_dir should point to a directory containing Javascript
%% source files which will be loaded by Riak when it initializes
%% Javascript VMs.
{%js_source_dir, "/tmp/js_source"},
...
}]

```

## yokozuna

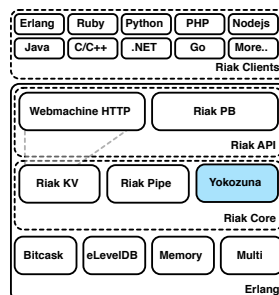


Figure 4.7: Tech Stack Yokozuna

Yokozuna is the newest addition to the Riak ecosystem. It's an integration of the distributed Solr search engine into Riak, and provides some extensions for extracting, indexing, and tagging documents. The Solr server runs its own HTTP interface, and though your Riak users should never have to access it, you can choose which `solr_port` will be used.

```

%% Yokozuna Search
{yokozuna, [
  {solr_port, "8093"},
  {yz_dir, "./data/yz"}
]}

```

## bitcask, eleveldb, memory, multi

Several modern databases have swappable backends, and Riak is no different in that respect. Riak currently supports three different storage engines—*Bitcask*, *eLevelDB*, and

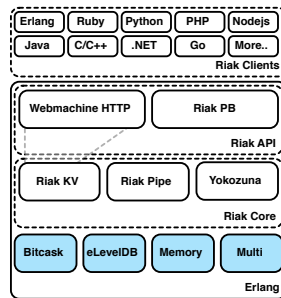


Figure 4.8: Tech Stack Backend

*Memory*—and one hybrid called *Multi*.

Using a backend is simply a matter of setting the `storage_backend` with one of the following values.

- `riak_kv_bitcask_backend` - The catchall Riak backend. If you don't have a compelling reason to *not* use it, this is my suggestion.
- `riak_kv_eleveldb_backend` - A Riak-friendly backend which uses Google's `leveldb`. This is necessary if you have too many keys to fit into memory, or wish to use 2i.
- `riak_kv_memory_backend` - A main-memory backend, with time-to-live (TTL). Meant for transient data.
- `riak_kv_multi_backend` - Any of the above backends, chosen on a per-bucket basis.

```
%% Riak KV config
{riak_kv, [
  %% Storage_backend specifies the Erlang module defining
  %% the storage mechanism that will be used on this node.
  {storage_backend, riak_kv_memory_backend}
]},
```

Then, with the exception of *Multi*, each memory configuration is under one of the following options.

```
%% Memory Config
{memory_backend, [
  {max_memory, 4096}, %% 4GB in megabytes
  {ttl, 86400} %% 1 Day in seconds
]}

%% Bitcask Config
{bitcask, [
  {data_root, "./data/bitcask"},
  {open_timeout, 4}, %% Wait time to open a keydir (in seconds)
  {sync_strategy, {seconds, 60}} %% Sync every 60 seconds
]}
```

```

}},

%% eLevelDB Config
{eleveldb, [
  {data_root, "./data/leveldb"},
  {write_buffer_size_min, 31457280 }, %% 30 MB in bytes
  {write_buffer_size_max, 62914560}, %% 60 MB in bytes
  {max_open_files, 20}, %% Maximum number of files open at once per partition
  {cache_size, 8388608} %% 8MB default cache size per-partition
]},

```

With the Multi backend, you can even choose different backends for different buckets. This can make sense, as one bucket may hold user information that you wish to index (use `eleveldb`), while another bucket holds volatile session information that you may prefer to simply remain resident (use `memory`).

```

%% Riak KV config
{riak_kv, [
  ...
  %% Storage_backend specifies the Erlang module defining
  %% the storage mechanism that will be used on this node.
  {storage_backend, riak_kv_multi_backend},

  %% Choose one of the names you defined below
  {multi_backend_default, <<"bitcask_multi">>},

  {multi_backend, [
    %% Here s where you set the individual backends
    {<<"bitcask_multi">>, riak_kv_bitcask_backend, [
      %% bitcask configuration
      {config1, ConfigValue1},
      {config2, ConfigValue2}
    ]},
    {<<"memory_multi">>, riak_kv_memory_backend, [
      %% memory configuration
      {max_memory, 8192} %% 8GB
    ]}
  ]},
]},

```

You can put the `memory_multi` configured above to the `session_data` bucket by just setting its backend property.

```

$ curl -XPUT http://riaknode:8098/riak/session_data \
-H "Content-Type: application/json" \
-d {"props":{"backend":"memory_multi"}}

```

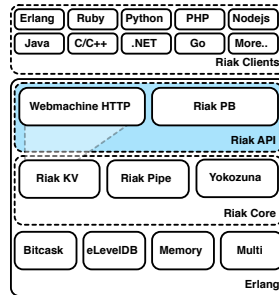
**riak\_api**

Figure 4.9: Tech Stack API

So far, all of the components we've seen have been inside the Riak house. The API is the front door. *In a perfect world*, the API would manage two implementations: Protocol buffers (PB), an efficient binary protocol framework designed by Google; and HTTP. Unfortunately the HTTP client interface is not yet ported, leaving only PB for now—though I like to consider this as a mere implementation detail, to be unraveled from KV soon.

But because they are not yet separated, only PB is configured under `riak_api`, while HTTP still remains under KV.

In any case, Riak API represents the client facing aspect of Riak. Implementations handle how data is encoded and transferred, and this project handles the services for presenting those interfaces, managing connections, providing entry points.

```
%% Riak Client APIs config
{riak_api, [
  %% pb_backlog is the maximum length to which the queue of pending
  %% connections may grow. If set, it must be an integer >= 0.
  %% By default the value is 5. If you anticipate a huge number of
  %% connections being initialised *simultaneously*, set this number
  %% higher.
  %% {pb_backlog, 64},

  %% pb_ip is the IP address that the Riak Protocol Buffers interface
  %% will bind to. If this is undefined, the interface will not run.
  {pb_ip, "127.0.0.1"},

  %% pb_port is the TCP port that the Riak Protocol Buffers interface
  %% will bind to
  {pb_port, 8087 }
]},
```

**Other projects**

Other projects add depth to Riak, though aren't strictly necessary, in a functional sense. Two of these projects are Lager, Riak's chosen logging system; and Sysmon, a useful system

monitor. They both have meaningful defaults, and are also have well documented settings at with their repository docs

- <https://github.com/basho/lager>
- [https://github.com/basho/riak\\_sysmon](https://github.com/basho/riak_sysmon)

```
%% Lager Config
{lager, [
  %% What handlers to install with what arguments
  %% If you wish to disable rotation, you can either set the size to 0
  %% and the rotation time to "", or instead specify a 2-tuple that only
  %% consists of {Logfile, Level}.
  {handlers, [
    {lager_file_backend, [
      {"/log/error.log", error, 10485760, "$D0", 5},
      {"/log/console.log", info, 10485760, "$D0", 5}
    ]}
  ]},

  %% Whether to write a crash log, and where.
  %% Commented/omitted/undefined means no crash logger.
  {crash_log, "/log/crash.log"},

  ...

  %% Whether to redirect error_logger messages into lager - defaults to true
  {error_logger_redirect, true}
]},

%% riak_sysmon config
{riak_sysmon, [
  %% To disable forwarding events of a particular type, set 0
  {process_limit, 30},
  {port_limit, 2},

  %% Finding reasonable limits for a given workload is a matter
  %% of experimentation.
  {gc_ms_limit, 100},
  {heap_word_limit, 40111000},

  %% Configure the following items to false to disable logging
  %% of that event type.
  {busy_port, true},
  {busy_dist_port, true}
]},
```

## Backward Incompatibility

Riak is a project in evolution. And as such, it has a lot of projects that have been created, but over time are being replaced with newer versions. Obviously this baggage can be confounding if you are just learning Riak—especially as you run across deprecated configuration, or documentation.

- InnoDB - The MySQL engine once supported by Riak, but now deprecated.
- Luke - The legacy mapreduce implementation replaced by Riak Pipe.
- Search - The search implementation replaced by Yokozuna.
- Merge Index - The backend created for the legacy Riak Search.
- SASL - A logging engine improved by Lager.

## Tools

### Riaknostic

You may recall that we skipped the `diag` command while looking through `riak-admin`, but it's time to circle back around.

Riaknostic is a diagnostic tool for Riak, meant to run a suite of checks against an installation to discover potential problems. If it finds any, it also recommends potential resolutions.

Riaknostic exists separately from the core project, but is meant to be downloaded and integrated with an installation.

<http://riaknostic.basho.com/>

```
$ wget https://github.com/basho/riaknostic/downloads/riaknostic-1.0.2.tar.gz -P /tmp
$ cd /riak/lib
$ tar xzvf /tmp/riaknostic-1.0.2.tar.gz
```

That's all you need to do to access your buffet of options.

```
$ riak-admin diag --list
```

Available diagnostic checks:

disk	Data directory permissions and atime
dumps	Find crash dumps
memory_use	Measure memory usage
nodes_connected	Cluster node liveness
ring_membership	Cluster membership validity
ring_preflists	Check ring satisfies <code>n_val</code>
ring_size	Ring <code>size</code> valid
search	Check whether search is enabled on all nodes

I'm a bit concerned that my disk might be slow, so I ran the `disk diagnostic`.

```
$ riak-admin diag disk
21:52:47.353 [notice] Data directory /riak/data/bitcask is not mounted with noatime .\
Please remount its disk with the noatime flag to improve performance.
```

Riaknostic returns an analysis and suggestion for improvement. Had my disk configuration been ok, the command would have returned nothing.

## Riak Control

The last tool we'll look at is the aptly named [Riak Control](#). It's a web application for managing Riak clusters, watching, and drilling down into the details of your nodes to get a comprehensive view of the system. That's the idea, anyway. It's forever a work in progress, and it does not yet have parity with all of the command-line tools we've looked at. However, it's great to quick checkups and routing configuration changes.

After [downloading](#) the project, is to alter some `app.config` settings, to both configure users, and to adhere to Control's security requirements (you're opening up your cluster to remote administration, so it's pretty important to get this right).

The first thing is to enable SSL and HTTPS in the `riak_core` section we saw above. You can just uncomment these lines, set the `https` port to 8069, and point the `certfile` and `keyfile` to your SSL certificate. If you have an intermediate authority, add the `cacertfile` too.

```
%% Riak Core config
{riak_core, [
  %% https is a list of IP addresses and TCP ports that the Riak
  %% HTTPS interface will bind.
  {https, [{ "127.0.0.1", 8069 }]},

  %% Default cert and key locations for https can be overridden
  %% with the ssl config variable, for example:
  {ssl, [
    {certfile, "./etc/cert.pem"},
    {keyfile,  "./etc/key.pem"},
    {cacertfile, "./etc/cacert.pem"}
  ]},
]}
```

Then, you'll have to enable Riak Control in your `app.config`, and add a user. Note that the user password is plain text. Yeah it sucks, so be careful to not open your Control web access to the rest of the world, or you risk giving away the keys to the kingdom.

```
%% riak_control config
{riak_control, [
  %% Set to false to disable the admin panel.
  {enabled, true},
]}
```



```

%% Authentication style used for access to the admin
%% panel. Valid styles are userlist <TODO>.
{auth, userlist},

%% If auth is set to userlist then this is the
%% list of usernames and passwords for access to the
%% admin panel.
{userlist, [{"admin", "lovesecretsexgod"}
            ]},

%% The admin panel is broken up into multiple
%% components, each of which is enabled or disabled
%% by one of these settings.
{admin, true}
1}

```

With Control in place, restart your node and connect via a browser (note you're using `https`) `https://localhost:8069/admin`. After you log in using the user you set, you should see a snapshot page, which communicates the health of your cluster.

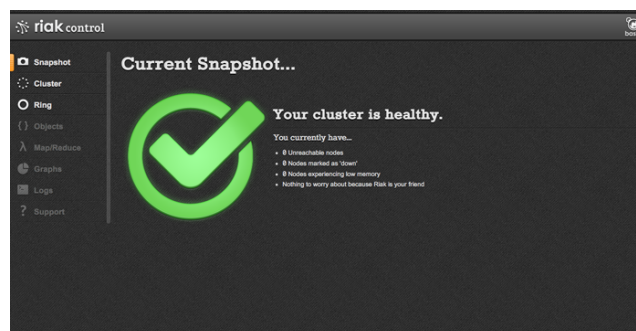


Figure 4.10: Snapshot View

If something is wrong, you'll see a huge red "X" instead of the green check mark, along with a list of what the trouble is.

From here you can drill down into a view the cluster's nodes, with details on memory usage, partition distribution, and other status. You can also add and configure and these nodes.

There is more in line for Riak Control, like performing mapreduce queries, stats views, graphs, and more coming down the pipe. It's not a universal toolkit quite yet, but it has a phenomenal start.

## Wrapup

Once you comprehend the basics of Riak, it's a simple thing to manage. If this seems like a lot to swallow, take it from a long-time relational database guy (me), Riak is a comparatively simple construct, especially when you factor in the complexity of distributed systems in

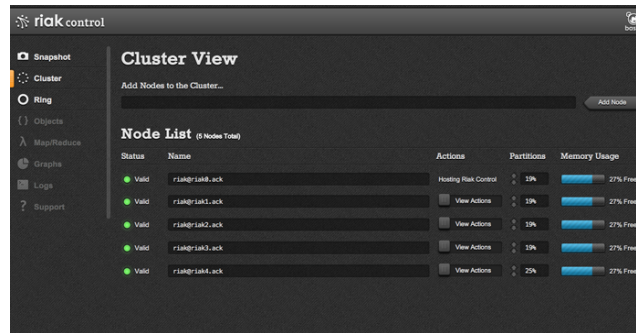


Figure 4.11: Cluster View

general. Riak manages much of the daily tasks an operator might do themselves manually, such as sharding by keys, adding/removing nodes, rebalancing data, supporting multiple backends, and allowing growth with unbalanced nodes. And due to Riak's architecture, the best part of all is when a server goes down at night, you can sleep (do you remember what that was?), and fix it in the morning.

## Chapter 5

# Notes

### A Short Note on MDC

*MDC*, or Multi Data Center, is a commercial extension to Riak provided by Basho. While the documentation is freely available, the source code is not. If you get the scale where keeping multiple Riak cluster in sync across the globe is necessary, I would recommend considering this option.

### A Short Note on RiakCS

*RiakCS* is Basho's commercial extension to Riak to allow your cluster to act as a remote storage mechanism, comparable to (and compatible with) Amazon's S3. There are several reasons you may wish to host your own cloud storage mechanism (security, legal reasons, you already own lots of hardware, cheaper at scale).