DSE 2023 Summer School Lausanne

Lecture 1: Numerical Dynamic Programming

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Topics

- Standard DP algorithms (VFI, HPI, OPI)
- State-dependent discounting

Code in

- Julia
- Python / Numba
- Python / NumPy
- Python / JAX

Based on discussion in Ch. 6 of https://dp.quantecon.org/

Example: Inventory Management

Consider an inventory model with Bellman equation

$$v(y) = \max_{a \in \Gamma(y)} \left\{ r(y, a) + \beta \sum_{y'} v(y') R(y, a, y') \right\}$$

- $y \in \{0, \dots, K\}$ is current inventory
- a is current inventory order (units of stock)
- r(y,a) is current profits
- R(y, a, y') = prob. next state is y' given y, a

In the code, reward (profit) r(y,a) is affected by a fixed cost per order

hence investment is lumpy (S-s dynamics)

One limitation: constant discounting

For this model we would expect

$$\beta = \frac{1}{1 + \text{interest rate}}$$

And interest rates vary over time

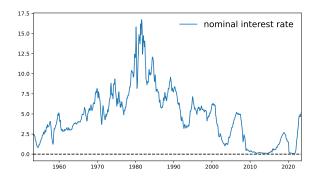


Figure: Nominal US interest rates

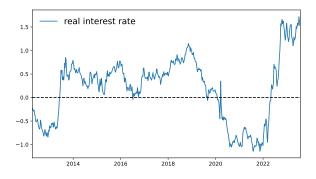


Figure: Real US interest rates

Let's shift to a setting with state-dependent discounting

Example. Discount time t reward r_t via

current value
$$= \mathbb{E} \, \beta_1 \, \beta_2 \, \cdots \, \beta_t \, r_t$$

Strategy:

- 1. study computation of lifetime values under state-dependent discounting
- 2. learn how to optimize these values across feasible policies

Let's start with step 1:

Valuation with non-constant discounting

Let $(X_t)_{t\geqslant 0}$ be a state process

- ullet the discount factor varies with X_t
- ullet profits / payoffs vary with X_t

To construct the state process, we

let X be finite and let P be a Markov matrix on X:

$$P\geqslant 0\quad\text{and}\quad \sum_{x'\in \mathsf{X}}P(x,x')=1$$

• let $(X_t)_{t\geqslant 0}$ be P-Markov on X:

$$\mathbb{P}\{X_{t+1} = x' \,|\, X_t = x\} = P(x, x')$$

Consider the discount factor process

$$\delta_t := \prod_{i=1}^t b(X_{i-1}, X_i)$$

- Example. $b \equiv \beta \in (0,1)$ implies $\delta_t = \beta^t$ (const. discounting)
- Assume b > 0

Let A be the discount operator

$$A(x, x') := b(x, x')P(x, x')$$

Let

$$\rho(A) = \max_{\lambda \in \sigma(A)} |\lambda|, \qquad \sigma(A) := \text{ eigenvalues of } A$$

Let \mathbb{R}^{X} = all real-valued functions on X

Theorem. If $h \in \mathbb{R}^X$ and $\rho(A) < 1$, then

$$v(x) := \mathbb{E}_x \sum_{t=0}^{\infty} \delta_t \, h(X_t)$$
 is finite for all $x \in \mathsf{X}$

Moreover, I - A is bijective and

$$v = \sum_{t \ge 0} A^t h$$
$$= (I - A)^{-1} h$$

Sketch of proof: An inductive argument shows that

$$\mathbb{E}_x \, \delta_t \, h(X_t) = (A^t h)(x)$$

Hence

$$v(x) = \mathbb{E}_x \sum_{t=0}^{\infty} \delta_t h(X_t) = \sum_{t=0}^{\infty} \mathbb{E}_x \delta_t h(X_t) = \sum_{t=0}^{\infty} (A^t h)(x)$$

That is, $v = \sum_{t \ge 0} A^t h$

By $\rho(A) < 1$ and the Neumann series lemma,

$$\sum_{t \ge 0} A^t = (I - A)^{-1}$$

Note that the condition $\rho(A) < 1$ is weak

- necessary as well as sufficient for convergence in many settings
- ullet permits discount factor >1 with positive probability

The last fact means

- can handle negative interest rates
- can handle models with large preference shocks
 - ZLB literature
 - asset pricing literature

Next steps

- 1. Obtain general results on DP with state-dependent discounting
- 2. Apply them to the inventory problem
- 3. Solve with Python / Julia using different algorithms

Our general results will use the theorem on slide 10

a condition to ensure that valuations are finite

MDPs with State-Dependent Discounting

We consider a Markov decision process (MDP) with Bellman equation

$$v(x) = \max_{a \in \Gamma(x)} \left\{ r(x, a) + \sum_{x'} v(x')\beta(x, a, x')P(x, a, x') \right\}$$

for all $x \in X$

- X, A finite state & action spaces
- Γ a nonempty correspondence from $X \to A$
- G := all (x, a) s.t. $a \in \Gamma(x)$
- β maps $G \times X$ to $(0, \infty)$

Let Σ be the set of **feasible policies**

$$\Sigma := \{ \sigma \in \mathsf{A}^\mathsf{X} : \sigma(x) \in \Gamma(x) \text{ for all } x \in \mathsf{X} \}$$

Let

- $r_{\sigma}(x) := r(x, \sigma(x)) = \text{rewards under } \sigma$
- $P_{\sigma}(x,x') := P(x,\sigma(x),x') = \text{transitions under } \sigma$
- $\beta_{\sigma}(x,x') := \beta(x,\sigma(x),x') = \text{discounting under } \sigma$

Note that P_{σ} is Markov dynamics for the state under σ

When it exists, the **lifetime value** v_{σ} of σ obeys

$$v_{\sigma}(x) = \mathbb{E}_x \sum_{t=0}^{\infty} \delta_t^{\sigma} r_{\sigma}(X_t)$$

where

$$\delta_t^{\sigma} := \prod_{i=1}^t \beta_{\sigma}(X_{i-1}, X_i)$$

and

$$(X_t)_{t\geqslant 0}$$
 is $P_\sigma ext{-Markov}$ with $X_0=x$

Let

$$L_{\sigma}(x, x') := \beta_{\sigma}(x, x') P_{\sigma}(x, x')$$

Proposition. If $\rho(L_{\sigma}) < 1$, then

- $I-L_{\sigma}$ is bijective
- v_{σ} is finite and

$$v_{\sigma} = \sum_{t \ge 0} L_{\sigma}^t v = (I - L_{\sigma})^{-1} r_{\sigma}$$

Proof: Follows directly from the result on slide 10

Let T_{σ} be the **policy operator** corresponding to σ :

$$T_{\sigma} v := r_{\sigma} + L_{\sigma} v$$

Note that

$$v_{\sigma} = (I - L_{\sigma})^{-1} r_{\sigma}$$
 $\iff v_{\sigma} \text{ solves eq. } (I - L)v = r_{\sigma}$
 $\iff v_{\sigma} \text{ solves eq. } v = r_{\sigma} + L_{\sigma}v$
 $\iff v_{\sigma} \in \operatorname{fix}(T_{\sigma})$

By the results on slide 10,

$$\rho(L_{\sigma}) < 1 \implies \text{fix}(T_{\sigma}) = \{v_{\sigma}\}$$

Fact. T_{σ} is globally stable on \mathbb{R}^{X} when $\rho(L_{\sigma}) < 1$:

$$\forall v \in \mathbb{R}^{\mathsf{X}}, \qquad T_{\sigma}^{k} v \to v_{\sigma} \quad \text{as} \quad k \to \infty$$
 (1)

• See proof in https://dp.quantecon.org/

Hence we have three ways to compute v_{σ} :

- 1. $v_{\sigma} = \sum_{t \geqslant 0} L_{\sigma}^t r_{\sigma}$
- $2. \ v_{\sigma} = (I L_{\sigma})^{-1} r_{\sigma}$
- 3. via (1)

We need a uniform bound, for all $\sigma \in \Sigma$:

Proposition. If \exists an L such that

- $L_{\sigma} \leqslant L$ for all $\sigma \in \Sigma$
- $\rho(L) < 1$

then, for all $\sigma \in \Sigma$, the unique fixed point of T_{σ} in \mathbb{R}^{X} is

$$v_{\sigma} = (I - L_{\sigma})^{-1} r_{\sigma}$$

Proof: Immediate from

- 1. $0 \leqslant A \leqslant B$ implies $\rho(A) \leqslant \rho(B)$
- 2. $0 \leqslant L_{\sigma} \leqslant L$ and $\rho(L) < 1$

What we have achieved

- found a condition for finite lifetime values across all policies
- found methods to compute these lifetime values

Yet to do

- show how to maximize lifetime values across $\sigma \in \Sigma$
- implement numerically

Defining optimality

Suppose the stated condition holds

• \exists matrix L with $\rho(L) < 1$ and $L_{\sigma} \leqslant L$ for all σ

Then v_{σ} is finite for all σ

Hence we can define the value function

$$v^*(x) := \max_{\sigma \in \Sigma} v_{\sigma}(x) \qquad (x \in \mathsf{X})$$

A policy σ is called **optimal** if $v_{\sigma} = v^*$

The Bellman operator takes the form

$$(Tv)(x) = \max_{a \in \Gamma(x)} \left\{ r(x, a) + \sum_{x'} v(x')\beta(x, a, x')P(x, a, x') \right\}$$

Given $v \in \mathbb{R}^{X}$, a policy σ is called v-greedy if

$$\sigma(x) \in \operatorname*{argmax}_{a \in \Gamma(x)} \left\{ r(x,a) + \sum_{x'} v(x') \beta(x,a,x') P(x,a,x') \right\}$$

for all x in X

Proposition. If there exists a linear operator L on \mathbb{R}^{X} such that

$$\rho(L) < 1 \quad \text{and} \quad \beta(x,a,x') P(x,a,x') \leqslant L(x,x')$$

for all $(x, a) \in G$ and $x' \in X$, then

- 1. T_{σ} is globally stable on \mathbb{R}^{X} with unique fixed point v_{σ}
- 2. $v_{\sigma} = (I L_{\sigma})^{-1} r_{\sigma}$
- 3. T is globally stable on \mathbb{R}^{X} with unique fixed point v^*
- 4. a policy $\sigma \in \Sigma$ is optimal if and only if it is v^* -greedy
- 5. at least one optimal policy exists

Algorithms

Now we have

- conditions for optimality
- general DP optimality results

Next we need algorithms

Let's consider

- value function iteration (VFI)
- Howard policy iteration (HPI)
- optimistic policy iteration (OPI)

Algorithm 1: VFI

input $v_0 \in \mathbb{R}^{\mathsf{X}}$

input au, a tolerance level for error

$$\varepsilon \leftarrow +\infty$$
$$k \leftarrow 0$$

while $\varepsilon > \tau$ do

$$\begin{vmatrix} v_{k+1} \leftarrow T v_k \\ \varepsilon \leftarrow \|v_k - v_{k+1}\|_{\infty} \\ k \leftarrow k + 1 \end{vmatrix}$$

end

Compute a v_k -greedy policy σ

return σ

Algorithm 2: HPI

input $\sigma_0 \in \Sigma$, an initial guess of σ^*

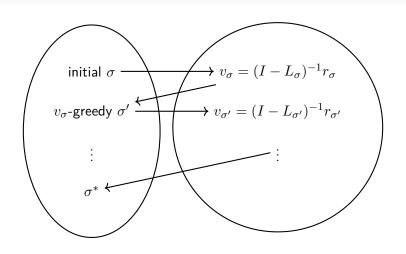
$$k \leftarrow 0$$
$$\varepsilon \leftarrow +\infty$$

while $\varepsilon > 0$ do

$$\begin{split} v_k &\leftarrow (I - L_{\sigma_k})^{-1} r_{\sigma_k} \\ \sigma_{k+1} &\leftarrow \text{a } v_k\text{-greedy policy} \\ \varepsilon &\leftarrow \mathbb{1}\{\sigma_k \neq \sigma_{k+1}\} \\ k &\leftarrow k+1 \end{split}$$

end

return σ_k



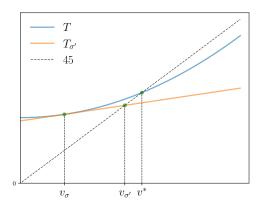


Figure: HPI as a version of Newton's method

Algorithm 3: OPI

```
input v_0 \in \mathbb{R}^X, an initial guess of v^*
input \tau, a tolerance level for error
input m \in \mathbb{N}, a step size
k \leftarrow 0
\varepsilon \leftarrow +\infty
while \varepsilon > \tau do
      \sigma_k \leftarrow \text{a } v_k\text{-greedy policy}
      v_{k+1} \leftarrow T_{\sigma_k}^m v_k
    \varepsilon \leftarrow \|v_k - v_{k+1}\|_{\infty}k \leftarrow k+1
```

end return σ_k

Proposition. Under the stated condition, VFI, HPI and OPI all converge

Moreover, HPI converges to an exact optimal policy in finitely many steps

For details and proofs see Ch. 6 of https://dp.quantecon.org/

Back to the inventory problem

Replace β with $\beta(Z_t)$ where $(Z_t)_{t\geqslant 0}$ is Q-Markov on Z

The Bellman equation becomes

$$v(y,z) = \max_{a \in \Gamma(y)} \left\{ r(y,a) + \beta(z) \sum_{z',y'} v(y',z') Q(z,z') R(y,a,y') \right\}$$

Given $\sigma \in \Sigma$, the **policy operator** is

$$(T_{\sigma} v)(y,z) = r(y,\sigma(y,z)) +$$

$$\beta(z) \sum_{z',y'} v(y',z') Q(z,z') R(y,\sigma(y,z),y')$$

Proposition Let

$$L(z, z') := \beta(z)Q(z, z')$$

If r(L) < 1, then all of the optimality results on slide 24 apply

Ex. Check the details

• See Ch. 6 of https://dp.quantecon.org/ if you get stuck



Code that solves the model can be found at

https://github.com/QuantEcon/dse_2023

Illustrates

- Python vs Julia
- Numba vs NumPy vs JAX
- Power of GPUs if you have one