

CprE 281: Digital Logic

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http://www.ece.iastate.edu/~alexs/classes/

Addition of Unsigned Numbers

CprE 281: Digital Logic Iowa State University, Ames, IA Copyright © Alexander Stoytchev

Administrative Stuff

- HW5 is out
- It is due on Monday Oct 3 @ 4pm.
- Please write clearly on the first page (in BLOCK CAPITAL letters) the following three things:
 - Your First and Last Name
 - Your Student ID Number
 - Your Lab Section Letter
- Also, please
 - Staple your pages

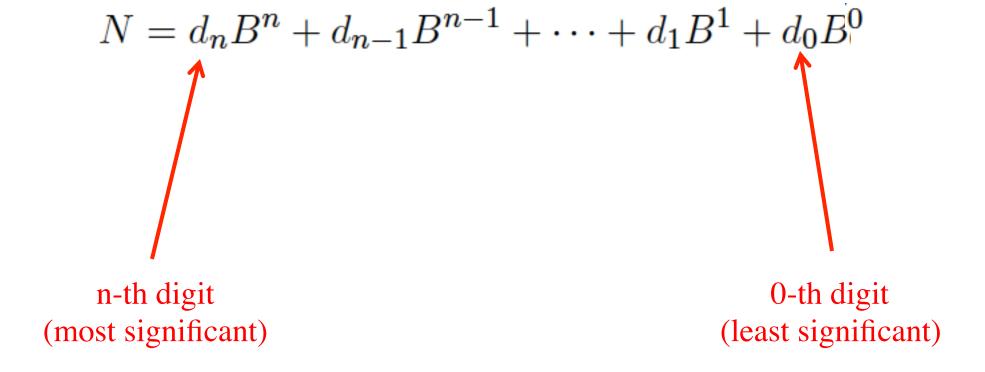
Administrative Stuff

- Labs next week
- Mini-Project
- This is worth 3% of your grade (x2 labs)
- http://www.ece.iastate.edu/~alexs/classes/ 2016_Fall_281/labs/Project-Mini/

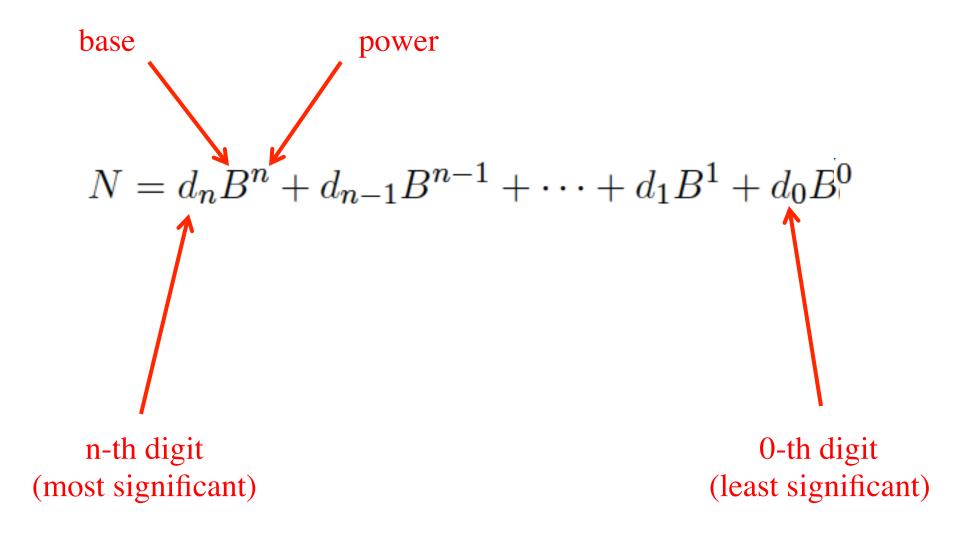
Number Systems

$$N = d_n B^n + d_{n-1} B^{n-1} + \dots + d_1 B^1 + d_0 B^0$$

Number Systems



Number Systems



The Decimal System

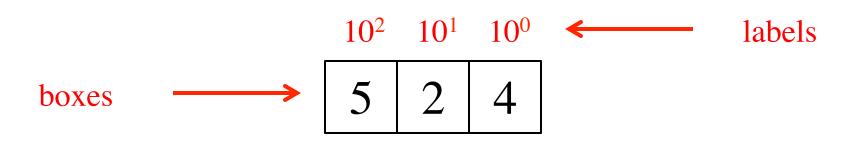
$$524_{10} = 5 \times 10^2 + 2 \times 10^1 + 4 \times 10^0$$

The Decimal System

$$524_{10} = 5 \times 10^{2} + 2 \times 10^{1} + 4 \times 10^{0}$$
$$= 5 \times 100 + 2 \times 10 + 4 \times 1$$
$$= 500 + 20 + 4$$

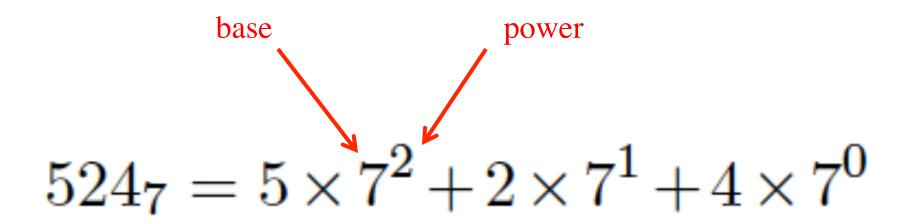
 $=524_{10}$

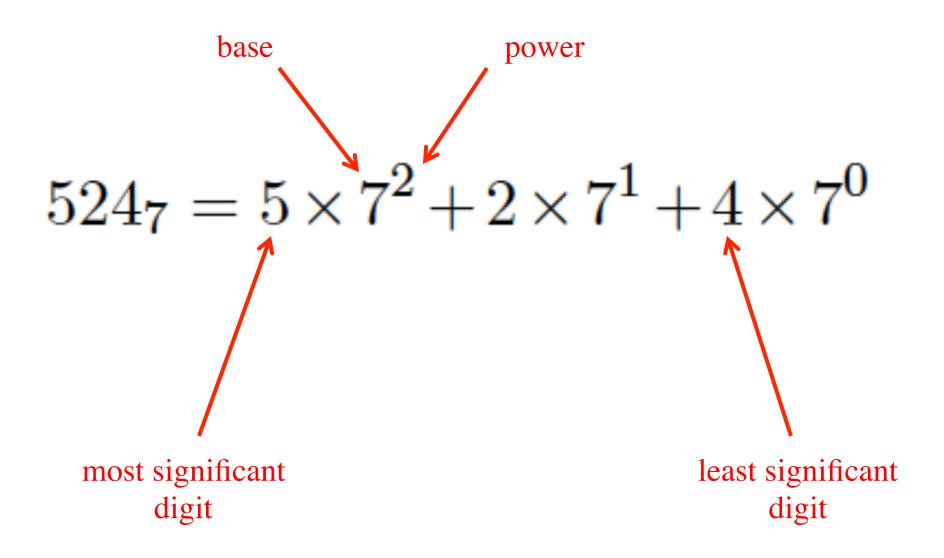
5 2 4



Each box can contain only one digit and has only one label. From right to left, the labels are increasing powers of the base, starting from 0.

$$524_7 = 5 \times 7^2 + 2 \times 7^1 + 4 \times 7^0$$



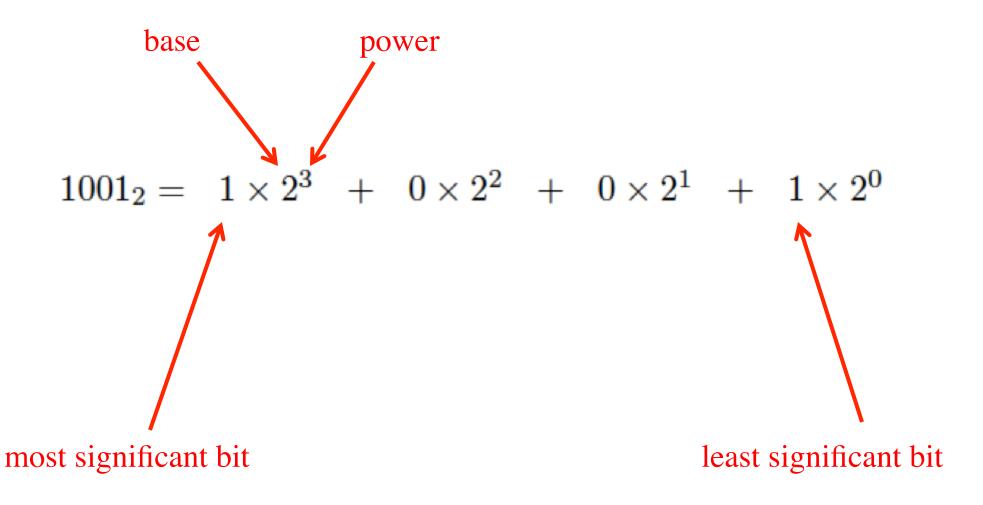


$$524_7 = 5 \times 7^2 + 2 \times 7^1 + 4 \times 7^0$$
$$= 5 \times 49 + 2 \times 7 + 4 \times 1$$
$$= 245 + 14 + 4$$
$$= 263_{10}$$

Binary Numbers (Base 2)

$$1001_2 = 1 \times 2^3 + 0 \times 2^2 + 0 \times 2^1 + 1 \times 2^0$$

Binary Numbers (Base 2)



Binary Numbers (Base 2)

$$1001_2 = 1 \times 2^3 + 0 \times 2^2 + 0 \times 2^1 + 1 \times 2^0 =$$
 $= 1 \times 8 + 0 \times 4 + 0 \times 2 + 1 \times 1 =$
 $= 8 + 0 + 0 + 1 =$
 $= 9_{10}$

Another Example

Powers of 2

$$2^{10} = 1024$$
 $2^9 = 512$
 $2^8 = 256$
 $2^7 = 128$
 $2^6 = 64$
 $2^5 = 32$
 $2^4 = 16$
 $2^3 = 8$
 $2^2 = 4$
 $2^1 = 2$
 $2^0 = 1$

What is the value of this binary number?

• 00101100

· 0 0 1 0 1 1 0 0

• $0*2^7 + 0*2^6 + 1*2^5 + 0*2^4 + 1*2^3 + 1*2^2 + 0*2^1 + 0*2^0$

• 0*128 + 0*64 + 1*32 + 0*16 + 1*8 + 1*4 + 0*2 + 0*1

• 0*128 + 0*64 + 1*32 + 0*16 + 1*8 + 1*4 + 0*2 + 0*1

• 32 + 8 + 4 = 44 (in decimal)

| 27 | 2^6 | 2^5 | 2^4 | 2^3 | 2^2 | 2^1 | 2^0 |
|----|-------|-------|-------|-------|-------|-------|-------|
| 0 | 0 | 1 | 0 | 1 | 1 | 0 | 0 |

Binary numbers

Unsigned numbers

all bits represent the magnitude of a positive integer

Signed numbers

left-most bit represents the sign of a number

| Decimal | Binary | Octal | Hexadecimal |
|---------|--------|-------|---------------|
| 00 | 00000 | 00 | 00 |
| 01 | 00001 | 01 | 01 |
| 02 | 00010 | 02 | 02 |
| 03 | 00011 | 03 | 03 |
| 04 | 00100 | 04 | 04 |
| 05 | 00101 | 05 | 05 |
| 06 | 00110 | 06 | 06 |
| 07 | 00111 | 07 | 07 |
| 08 | 01000 | 10 | 08 |
| 09 | 01001 | 11 | 09 |
| 10 | 01010 | 12 | 0A |
| 11 | 01011 | 13 | 0B |
| 12 | 01100 | 14 | 0C |
| 13 | 01101 | 15 | 0D |
| 14 | 01110 | 16 | $0\mathrm{E}$ |
| 15 | 01111 | 17 | 0F |
| 16 | 10000 | 20 | 10 |
| 17 | 10001 | 21 | 11 |
| 18 | 10010 | 22 | 12 |
| | | | |

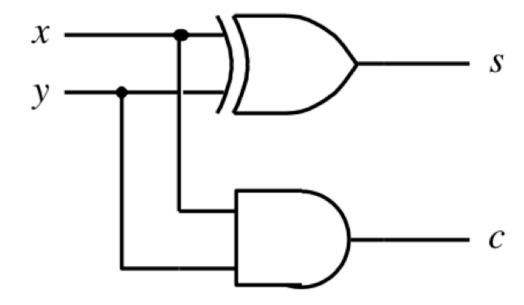
Table 3.1. Numbers in different systems.

Adding two bits (there are four possible cases)

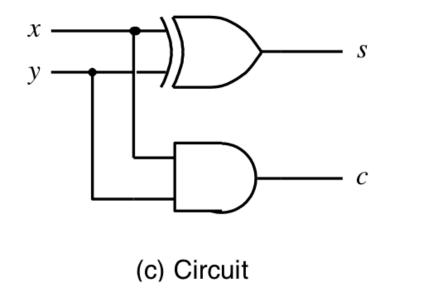
Adding two bits (the truth table)

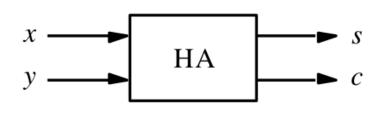
| | Carry | Sum |
|-----|-------|-----|
| x y | С | |
| 0 0 | 0 | 0 |
| 0 1 | 0 | 1 |
| 1 0 | 0 | 1 |
| 1 1 | 1 | 0 |
| | | |

Adding two bits (the logic circuit)



The Half-Adder





(d) Graphical symbol

Addition of multibit numbers

Generated carries
$$\longrightarrow$$
 1 1 1 0 ... c_{i+1} c_i ... $X = x_4 x_3 x_2 x_1 x_0$ 0 1 1 1 1 (15)₁₀ ... x_i ... x_i ... y_i ...

Bit position *i*

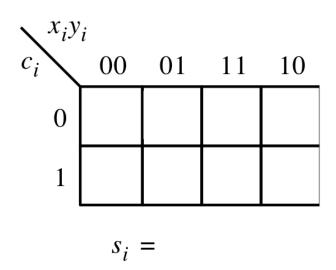
Problem Statement and Truth Table

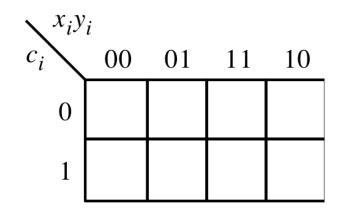
| | c_{i+1} | C_{i} | |
|-----|-----------|---------|-----|
| | | | |
| | | y_i | |
| _ | | | |
| ••• | ••• | s_i | ••• |

| C | \hat{c}_i | x_i | y_i | c_{i+1} | s_i |
|---|-------------|-------|-------|-----------|-------|
| | 0 | 0 | 0 | 0 | 0 |
| | 0 | 0 | | 0 | 1 |
| (| 0 | 1 | 0 | 0 | 1 |
| | 0 | 1 | 1 | 1 | 0 |
| | 1 | 0 | 0 | 0 | 1 |
| | 1 | 0 | 1 | 1 | 0 |
| | 1 | 1 | 0 | 1 | 0 |
| | 1 | 1 | 1 | 1 | 1 |
| | | | | | |

Let's fill-in the two K-maps

| c_i | x_i | y_i | c_{i+1} | s_i |
|-------|-------|-------|-----------|-------|
| 0 | 0 | 0 | 0 | 0 |
| 0 | 0 | 1 | 0 | 1 |
| 0 | 1 | 0 | 0 | 1 |
| 0 | 1 | 1 | 1 | 0 |
| 1 | 0 | 0 | 0 | 1 |
| 1 | 0 | 1 | 1 | 0 |
| 1 | 1 | 0 | 1 | 0 |
| 1 | 1 | 1 | 1 | 1 |

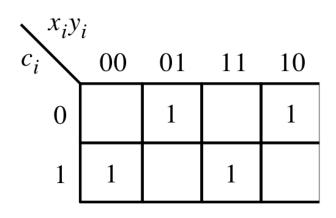




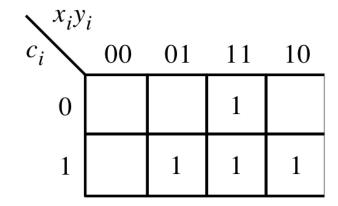
$$c_{i+1} =$$

Let's fill-in the two K-maps

| c_{i} | x_i | y_i | c_{i+1} | s_i |
|---------------------------------|----------------------------|----------------------------|----------------------------|---------------------------------|
| 0 0 0 0 1 1 1 | 0 0 1 1 0 0 | 0 1 0 1 0 1 | 0 0 0 1 0 1 | 0 1 1 0 1 0 0 |
| 1 | 1 | 1 | 1 | 1 |

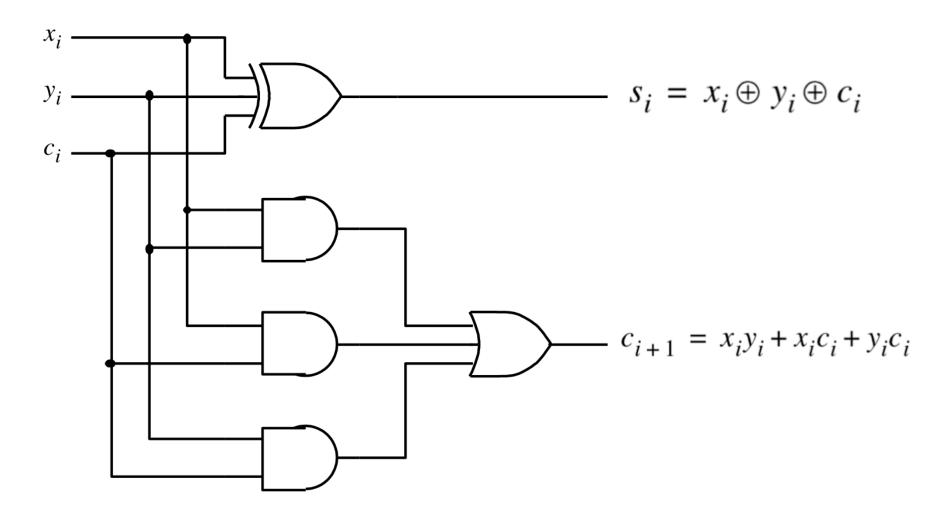


$$s_i = x_i \oplus y_i \oplus c_i$$

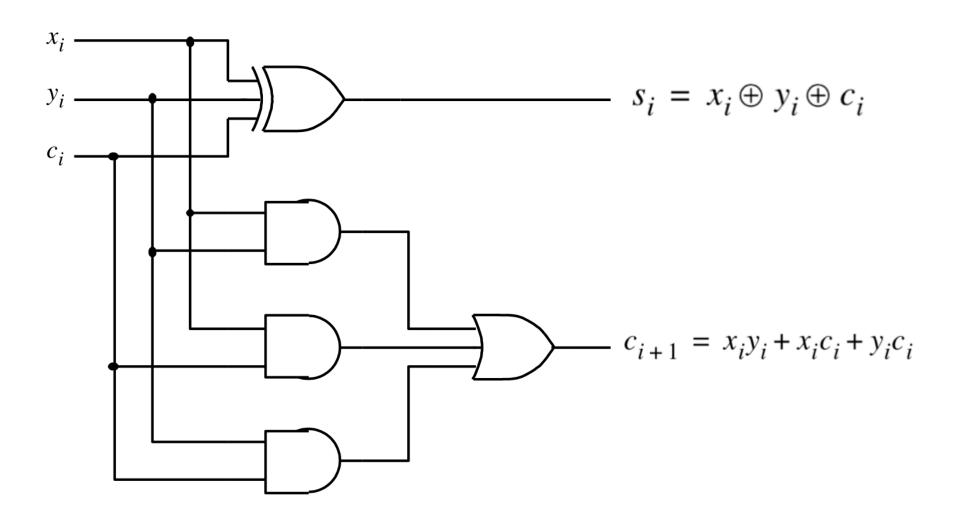


$$c_{i+1} = x_i y_i + x_i c_i + y_i c_i$$

The circuit for the two expressions



This is called the Full-Adder



XOR Magic

$$s_i = \overline{x}_i y_i \overline{c}_i + x_i \overline{y}_i \overline{c}_i + \overline{x}_i \overline{y}_i c_i + x_i y_i c_i$$

XOR Magic

$$s_i = \overline{x}_i y_i \overline{c}_i + x_i \overline{y}_i \overline{c}_i + \overline{x}_i \overline{y}_i c_i + x_i y_i c_i$$

$$s_i = (\overline{x}_i y_i + x_i \overline{y}_i) \overline{c}_i + (\overline{x}_i \overline{y}_i + x_i y_i) c_i$$
$$= (x_i \oplus y_i) \overline{c}_i + (\overline{x}_i \oplus y_i) c_i$$
$$= (x_i \oplus y_i) \oplus c_i$$

XOR Magic

$$s_i = \overline{x}_i y_i \overline{c}_i + x_i \overline{y}_i \overline{c}_i + \overline{x}_i \overline{y}_i c_i + x_i y_i c_i$$

Can you prove this?

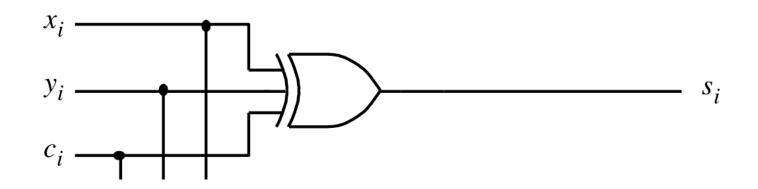
$$s_{i} = (\overline{x}_{i}y_{i} + x_{i}\overline{y}_{i})\overline{c}_{i} + (\overline{x}_{i}\overline{y}_{i} + x_{i}y_{i})c_{i}$$

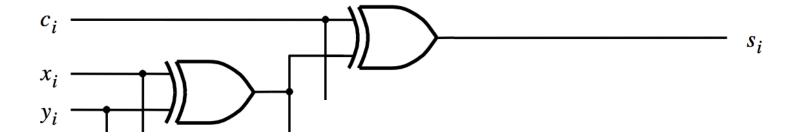
$$= (x_{i} \oplus y_{i})\overline{c}_{i} + (x_{i} \oplus y_{i})e_{i}$$

$$= (x_{i} \oplus y_{i}) \oplus c_{i}$$

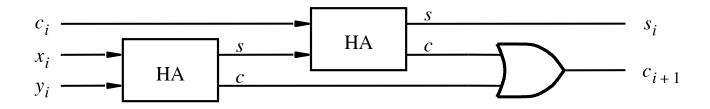
XOR Magic (s_i can be implemented in two different ways)

$$s_i = x_i \oplus y_i \oplus c_i$$

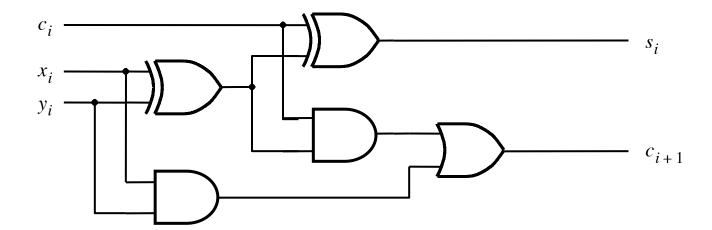




A decomposed implementation of the full-adder circuit

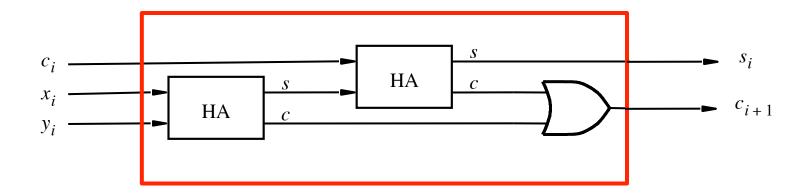


(a) Block diagram

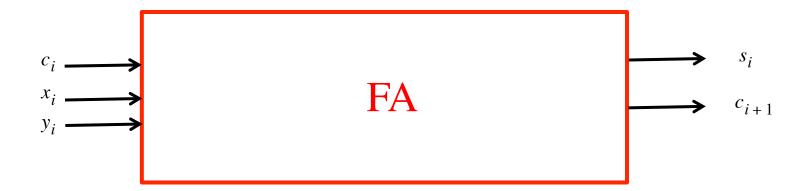


(b) Detailed diagram

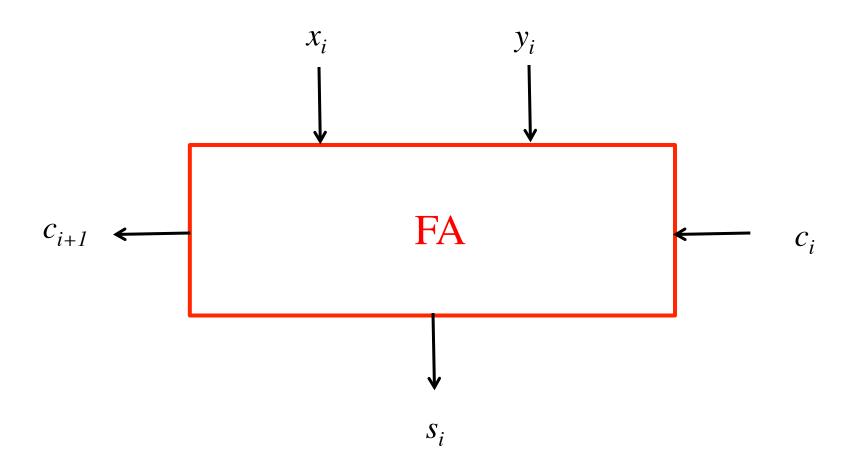
The Full-Adder Abstraction



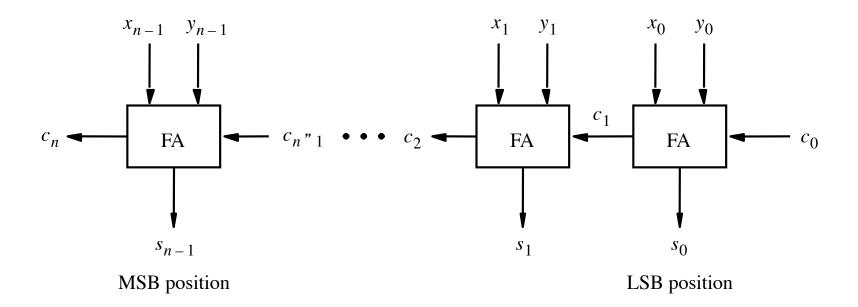
The Full-Adder Abstraction



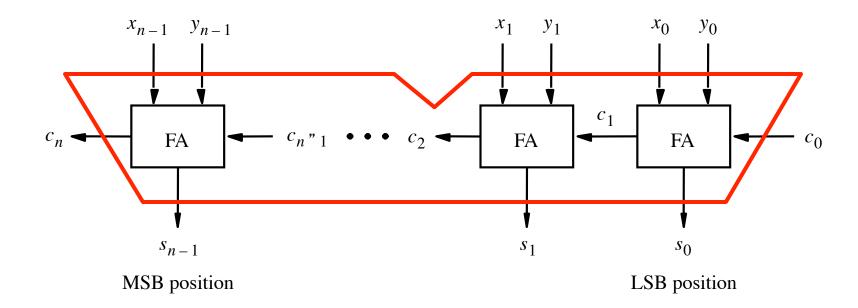
We can place the arrows anywhere



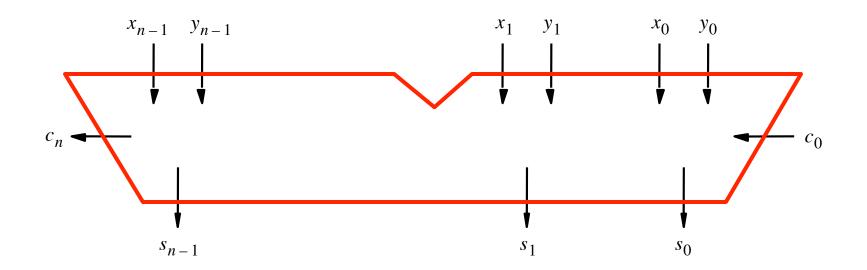
n-bit ripple-carry adder



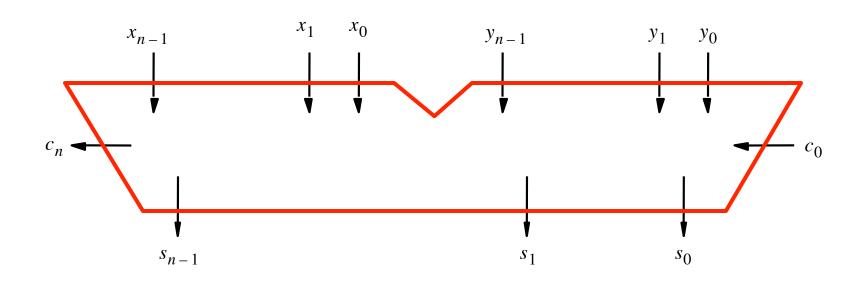
n-bit ripple-carry adder abstraction



n-bit ripple-carry adder abstraction



The x and y lines are typically grouped together for better visualization, but the underlying logic remains the same



Design Example:

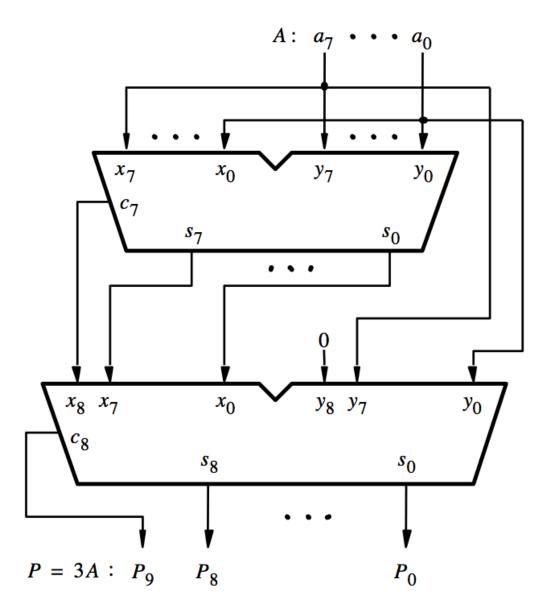
Create a circuit that multiplies a number by 3

How to Get 3A from A?

•
$$3A = A + A + A$$

•
$$3A = (A+A) + A$$

•
$$3A = 2A + A$$



[Figure 3.6a from the textbook]

Decimal Multiplication by 10

What happens when we multiply a number by 10?

$$4 \times 10 = ?$$

$$542 \times 10 = ?$$

$$1245 \times 10 = ?$$

Decimal Multiplication by 10

What happens when we multiply a number by 10?

$$4 \times 10 = 40$$

$$542 \times 10 = 5420$$

$$1245 \times 10 = 12450$$

Decimal Multiplication by 10

What happens when we multiply a number by 10?

$$4 \times 10 = 40$$

$$542 \times 10 = 5420$$

$$1245 \times 10 = 12450$$

You simply add a zero as the rightmost number

Binary Multiplication by 2

What happens when we multiply a number by 2?

011 times 2 = ?

101 times 2 = ?

110011 times 2 = ?

Binary Multiplication by 2

What happens when we multiply a number by 2?

011 times 2 = 0110

101 times 2 = 1010

110011 times 2 = 1100110

Binary Multiplication by 2

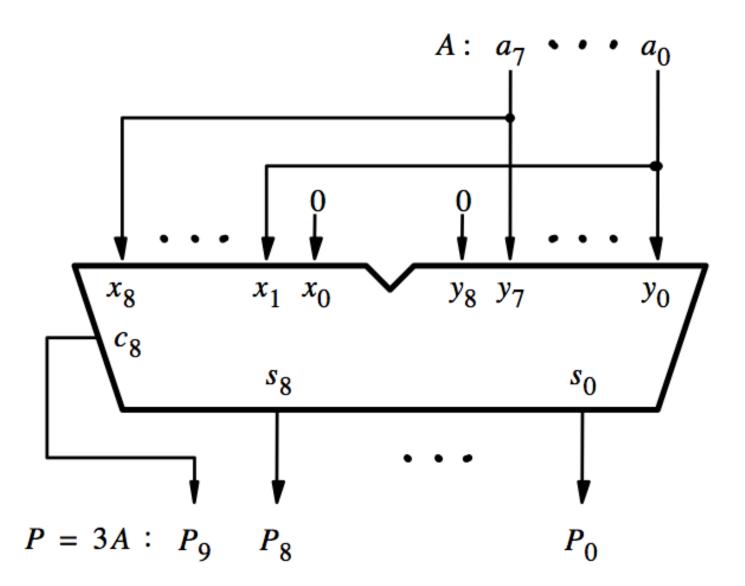
What happens when we multiply a number by 2?

$$011 \text{ times } 2 = 0110$$

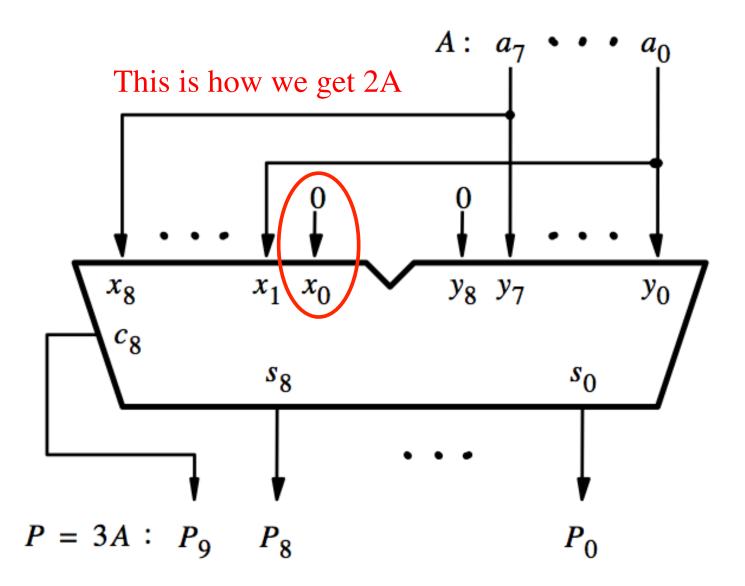
$$101 \text{ times } 2 = 1010$$

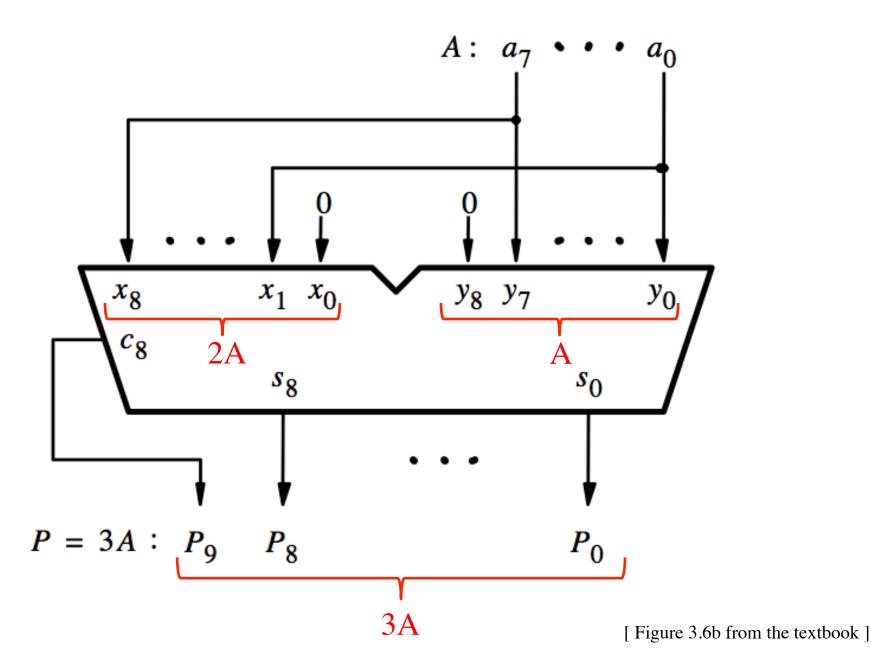
110011 times 2 = 1100110

You simply add a zero as the rightmost number



[Figure 3.6b from the textbook]





Questions?

THE END