



8086 Memory Banks

Dept. of Computer Science and Engineering
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CSE 341 Team

Lecture References:

► **Book:**

- *Microprocessors and Interfacing: Programming and Hardware,*

Author: Douglas V. Hall

- *The 8086/8088 Family: Design, Programming, And Interfacing,*

Author: John Uffenbeck.

8086 Memory Organisation

▶ **RECAP:**

- ▶ The 8086 has 20-bit address bus, so it can address 2^{20} or 1,048,576 addresses.
- ▶ Each address can store a byte. Hence, 8086 can store upto 1MB.
- ▶ Each read/write operation takes 1 bus cycle

Address	Contents
0xFFFF	1000 0000
.....
.....
0x00008	0100 1001
0x00007	1100 1100
0x00006	0110 1110
0x00005	0110 1110
0x00004	0000 0000
0x00003	0110 1011
0x00002	0101 0001
0x00001	1100 1001
0x00000	0100 1111

8086 Memory Organisation

- Each read/write operation takes 1 bus cycle for 1 byte data
- To read/write 1 word or 2 bytes of data, the up would need 2 cycles
- To solve this problem by saving processing time, memory is organized into *memory banks*
- Odd and even banks
- With the use of A_0 and \overline{BHE} pins

Address	Contents
0xFFFFF	1000 0000
.....
.....
0x00008	0100 1001
0x00007	1100 1100
0x00006	0110 1110
0x00005	0110 1110
0x00004	0000 0000
0x00003	0110 1011
0x00002	0101 0001
0x00001	1100 1001
0x00000	0100 1111

8086 Memory Organisation

- ▶ Addresses are consecutively numbered
- ▶ All even numbers end with a 0 and all odd numbers end with a 1

Even	Odd
6 = 11 0	5 = 10 1
14 = 111 0	13 = 110 1
20 = 1010 0	27 = 1101 1
42 = 10101 0	35 = 10001 1

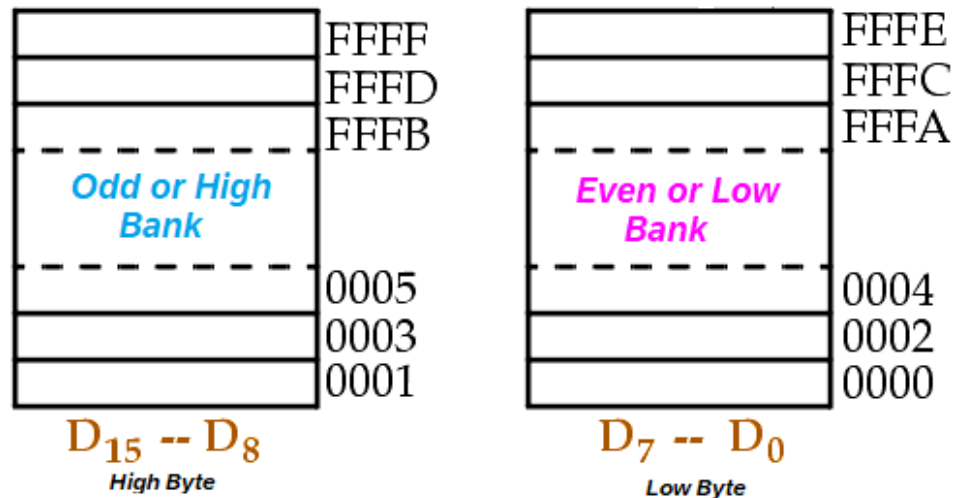
Address	Contents
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0x00004	0000 0000
0x00003	0110 1011
0x00002	0101 0001
0x00001	1100 1001
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even

odd

8086 Memory Organisation

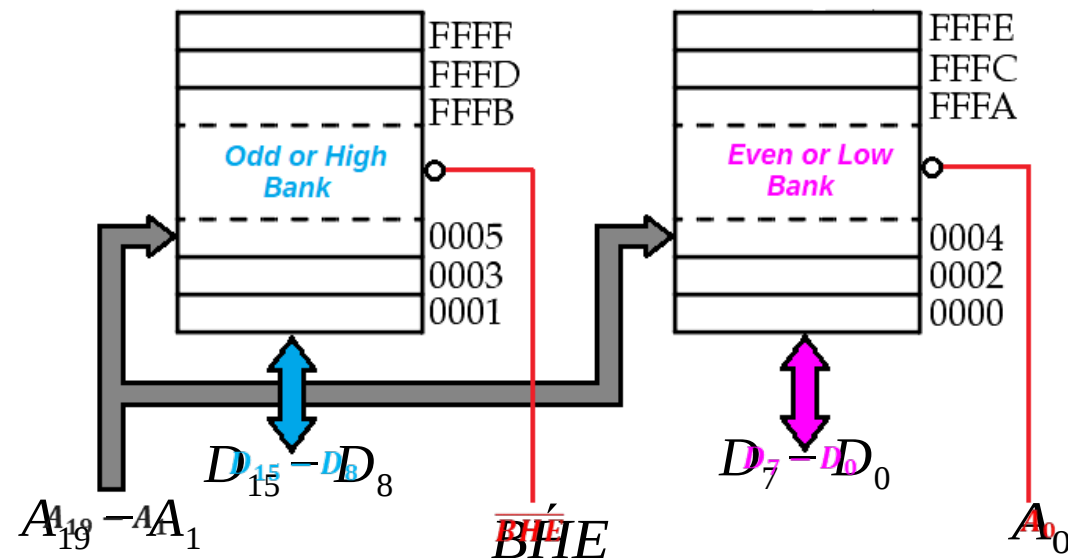
- ▶ Even addresses are considered as the **even/low bank**, which holds the content of the **low byte** while
- ▶ Odd addresses are considered as the **odd/high bank**, which holds the content of the **high byte**



8086 Memory Banks

8086 Pin Recap:

- ▶ Pins for carrying 20-bit address
- ▶ Pins for carrying 16-bit data
- ▶ Pins for carrying low data byte and for high byte – D_{15} for high byte
- ▶ \overline{BHE} when 0 indicates data bus carries high byte of data



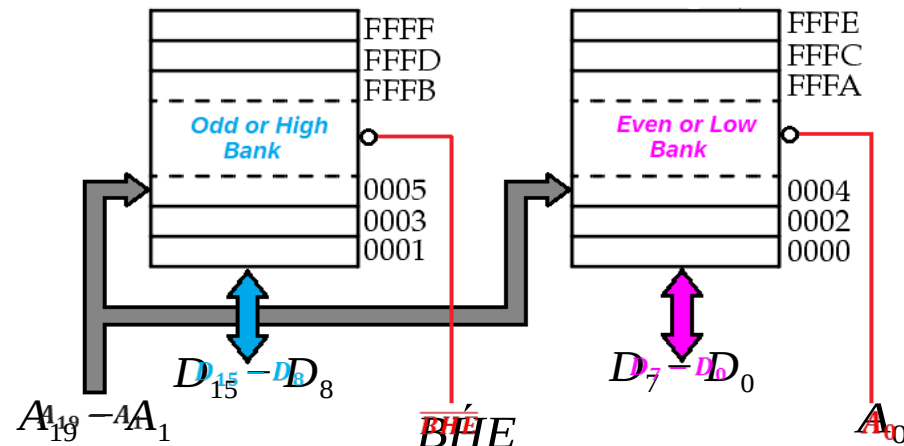
- ▶ The memory for an 8086 is set up in to 2 banks of up to 524,288 bytes or 512kB each
- ▶ This makes it possible to read/write a word with one machine cycle

8086 Memory Addressing

Data can be accessed from the memory in 4 different ways:

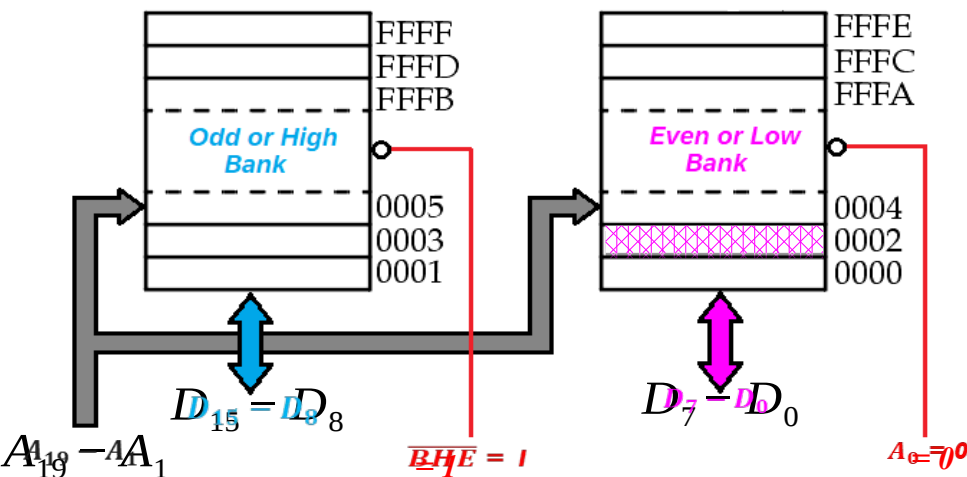
- ▶ 8 - bit data from Even Bank e.g. from address 00002
- ▶ 8 - bit data from Odd Bank e.g. from address 00003
- ▶ 16 - bit data starting from Even Address e.g. from 00002 and 00003
- ▶ 16 - bit data starting from Odd Address. e.g. from 00003 and 00004

		Type of Transfer	Data Lines
0	0	Word i.e. a byte from each bank	
0	1	Byte from odd bank	
1	0	Byte from even bank	
1	1	None	-



A byte from Low/Even bank

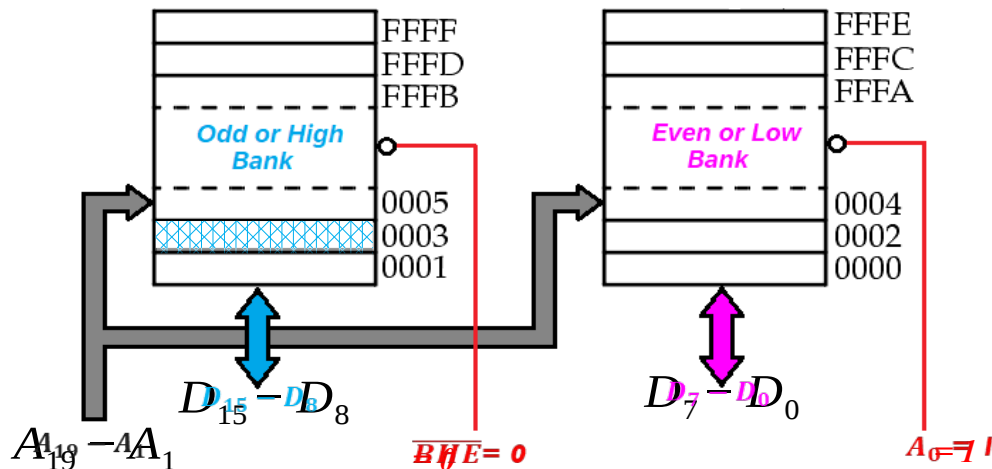
- Requires one bus cycle to read/write a data byte.
- Valid address is provided via A_{19} with A_1 & with $A_0 = 0$ & $\overline{BHE} = 1$
- Byte of data fetched on $D_7 - D_0$
- Low bank enabled, High bank disabled



		Type of Transfer	Data Lines
1	0	Byte from even bank	

A byte from High/Odd bank

- Requires one bus-cycle to read/write a data-byte.
- Valid address is provided via $A_{19} - A_1$ with $A_0 = 1$ & $\overline{BHE} = 0$
- Byte of data fetched on $D_{15} - D_8$
- Low bank disabled, High bank *enabled*

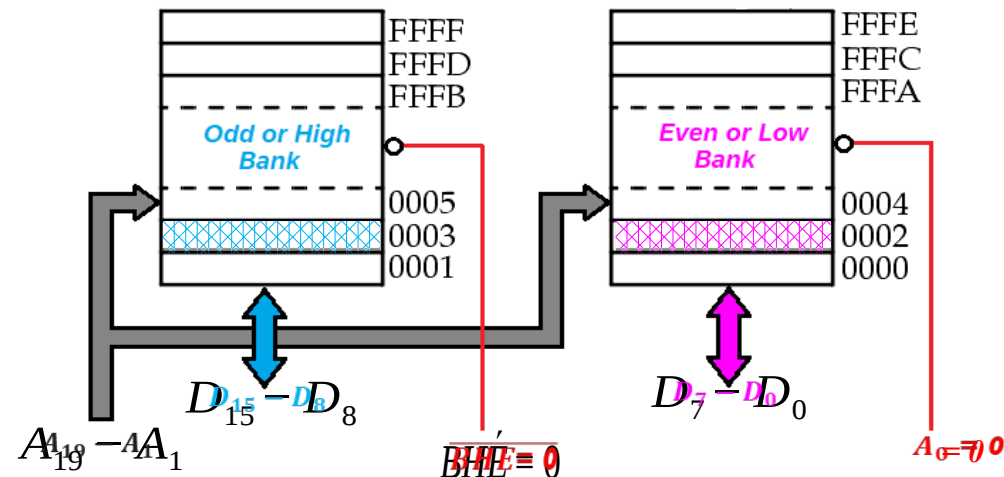


		Type of Transfer	Data Lines
0	1	Byte from odd bank	

An aligned word

starting from an **even** address

- Requires one bus cycle to read/write a data word.
- Valid address is provided via A_{19} with A_1 & with $A_0 = 0$ & $\overline{BHE} = 0$
- Word of data fetched on $D_{15} - D_0$
- Low bank and High bank enabled

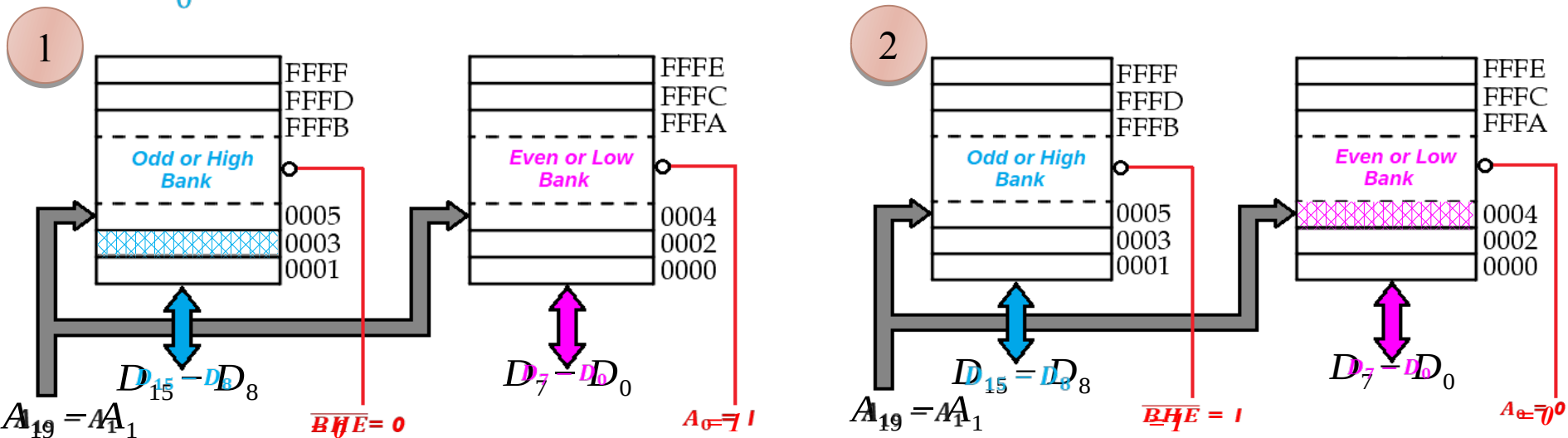


		Type of Transfer	Data Lines
0	0	Word i.e. a byte from each bank	

An unaligned word

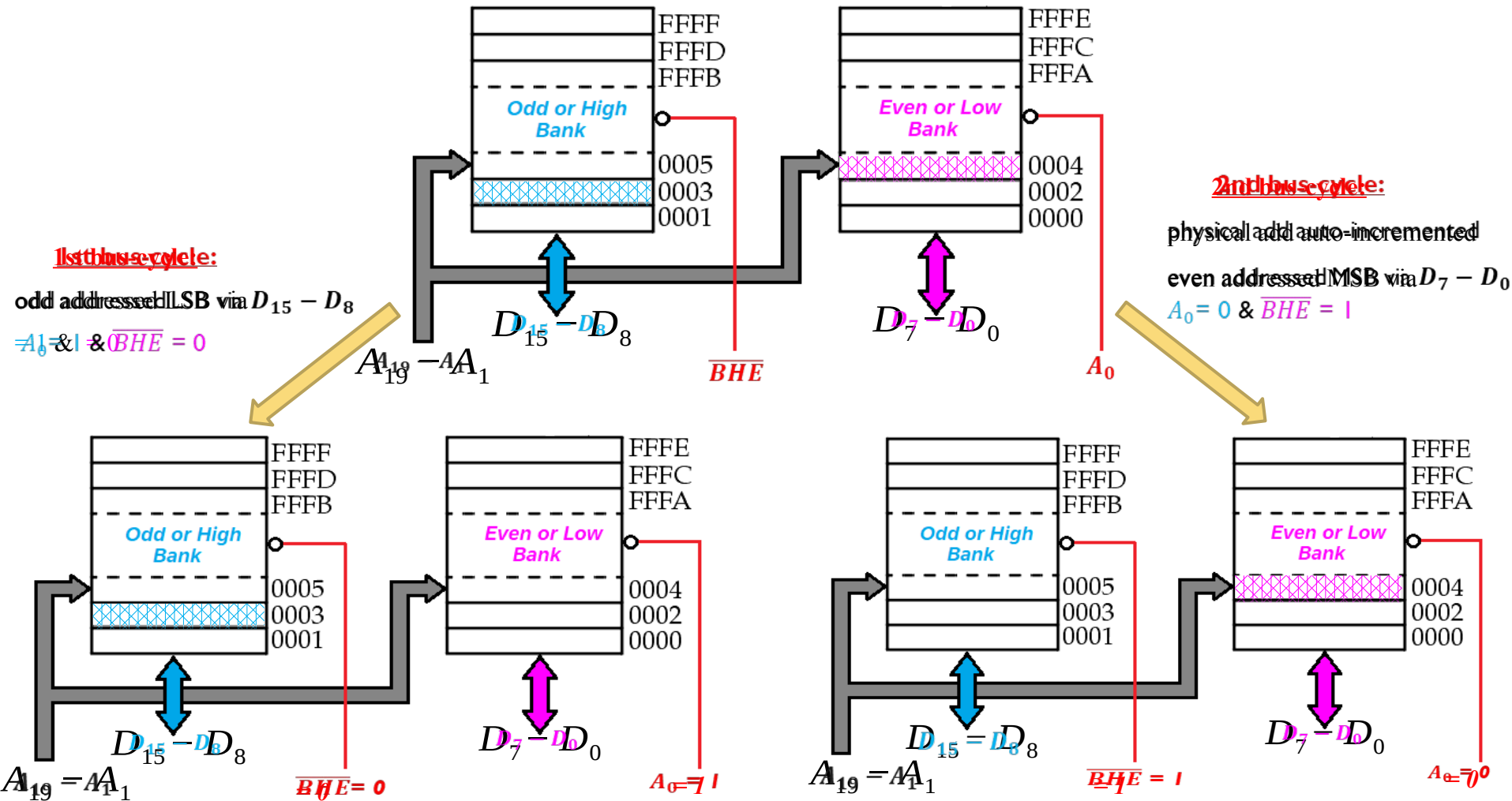
starting from an **odd** address

- Requires **two bus-cycles** to read/write a data-word
- During the 1st bus cycle**, odd addressed LSB of the word is accessed from the high memory bank via $D_{15} - D_8$ of data bus; $A_0 = 1$ & $\overline{BHE} = 0$
- During 2nd bus cycle**, physical address is auto-incremented to access the even address MSB of the word from the Low bank via $D_7 - D_0$; $A_0 = 0$ & $\overline{BHE} = 1$



An unaligned word starting from an **odd** address

2 bus-cycles required



Thank You

Questions are welcome in the discussion class