

Function Call Convention

Compass Security Schweiz AG Tel +41 55 214 41 60 Werkstrasse 20 Postfach 2038 CH-8645 Jona

Fax +41 55 214 41 61 team@csnc.ch www.csnc.ch

Content



Intel Architecture

Memory Layout

C Arrays

Assembler

Shellcode

Function Calls

Debugging

Buffer Overflow

BoF Exploit

Remote Exploit

Exploit Mitigations

Defeat Exploit Mitigations

Function Call Convention



Function call convention:

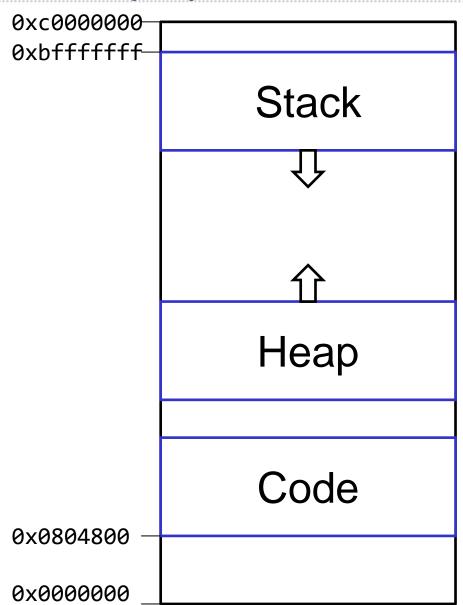
- → How functions work
- → Program-metadata on the stack

Stack based buffer overflow:

→ Overwrite program-metadata on the stack

x32 Memory Layout





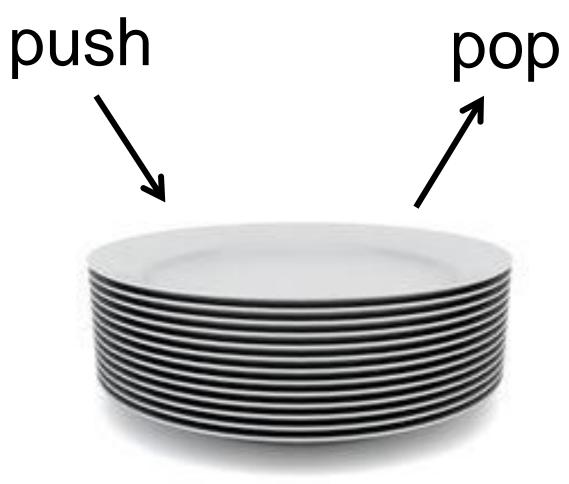




Stacks

How do they work?





Stack



Push 0x1

Push 0x2

Push 0x3

Pop

Push 0x4

0x01	
0x02	
0x04	





Functions and the Stack

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Fax +41 55 214 41 61 team@csnc.ch www.csnc.ch



What is a function?

- **→** Self contained subroutine
- ★ Re-usable
- → Can be called from anywhere
- → After function is finished: Jump to the calling function (calee)



```
void main(void) {
 int blubb = 0;
 foobar (blubb) ;
 return;
void foobar (int arg1) {
 char compass1[];
 char compass2[];
```



What does the function foobar() need?

- → Function Argument:
 - →blubb
- → Local variables
 - **+**Compass1
 - **+**Compass2
- → And: Address of next instruction in main()
 - **→**&return



Saved IP (&__libc_start)
Saved Frame Pointer
Local Variables <main>

SIP SFP blubb

Stack Frame <main>

Argument for <foobar>
Saved IP (&return)

Saved Frame Pointer

Local Variables <foobar>

&blubb
SIP
SFP
compass1

Stack Frame <foobar>

compass2

push por



```
void main(void) {
                       Save ptr
  int blubb = 0; \leftarrow -
  foobar(blubb);
                                 &blubb
  return; <
                                 SIP
                 pointer
                                 SFP
                                 compass1
void foobar (int arg1)
                                 compass2
 char compass1[];
  char compass2[];
                     allocate
```



Saved IP (&__libc_start)
Saved Frame Pointer
Local Variables <main>

SIP SFP blubb

Stack Frame <main>

Argument for <foobar>

Saved IP (&return)

Saved Frame Pointer

Local Variables <foobar>

SIP (&return)

SFP

&blubb

compass1

compass2

push / pop

Stack Frame <foobar>



SIP: Stored Instruction Pointer

- ◆ Copy of EIP
- → Points to the address where control flow continues after end of function
 - ◆(return, ret)
- → Usually points into the code section



Attention! Assembler ahead!

→ AT&T vs Intel syntax

Intel syntax:

mov eax, 1

mov ebx, 0ffh

int 80h

AT&T syntax:

movl \$1, %eax

movl \$0xff, %ebx

int \$0x80

Don't hang me if I messed this up somewhere



In ASM:

call 0x11223344 <&foobar>



<function code> (0x11223344)

ret
pop eip



In ASM:

call 0x11223344 <&foobar>



mov ebp, esp
<function code>
mov esp, ebp
ret





In ASM:

call 0x11223344 <&foobar>

push EIP

jmp 0x11223344

mov ebp, esp

<function code>

mov esp, ebp

ret

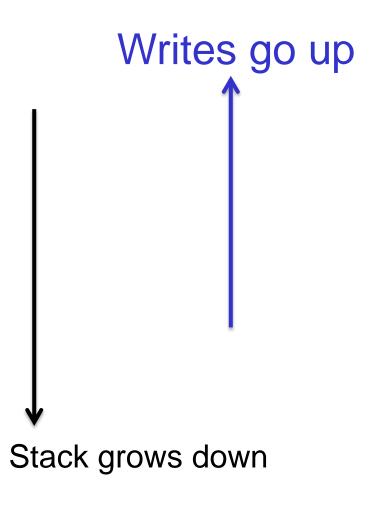
pop eip

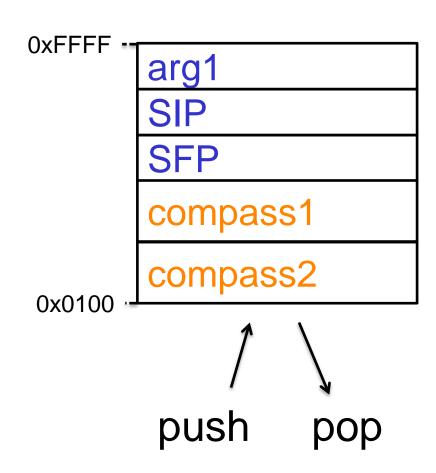
Prolog

Function

Epilog









Recap:

- → User data is on the stack
- → Also: important stuff is on the stack (Instruction Pointer, SIP)
- → Stack grows down
- → Writes go up





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```
int add(int x, int y) {
  int sum;
  sum = x + y;
  return sum;
}
```



$$c = add(3, 4)$$

C

ASM

ASM, detailed



add():

push 4
push 3
push EIP
jmp add

push ebp mov ebp, esp, sub esp, 0x10

mov eax, DWORD PTR [ebp + 0xc] mov edx, DWORD PTR [ebp + 0x8] add eax, edx mov DWORD PTR [ebp - 0x04], eax mov eax, DWORD PTR [ebp - 0x04]

leave ret



add():

```
push 4
push 3
push EIP
jmp add
```

```
push ebp
mov ebp, esp,
sub esp, 0x10
```

```
mov eax, DWORD PTR [ebp + 0xc]
mov edx, DWORD PTR [ebp + 0x8]
add eax, edx
mov DWORD PTR [ebp - 0x04], eax
mov eax, DWORD PTR [ebp - 0x04]
```

```
mov esp, ebp ; leave pop ebp ; leave ret
```



add():

```
push 4
push 3
push EIP
jmp add
```

```
push ebp
mov ebp, esp,
sub esp, 0x10
```

```
mov eax, DWORD PTR [ebp + 0xc]
mov edx, DWORD PTR [ebp + 0x8]
add eax, edx
mov DWORD PTR [ebp - 0x04], eax
mov eax, DWORD PTR [ebp - 0x04]
```

mov esp, ebp ; leave

pop ebp ; leave

pop eip ; ret



add():

push 4
push 3
push EIP
jmp add

push ebp mov ebp, esp, sub esp, 0x10

mov esp, ebp ; leave

pop ebp ; leave

pop eip ; ret



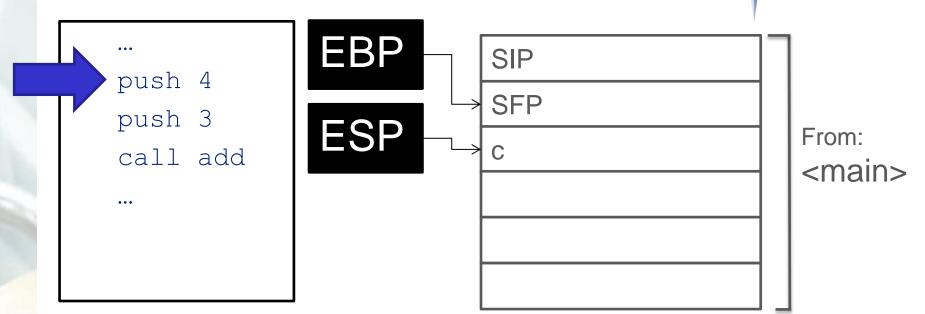


Function Prolog

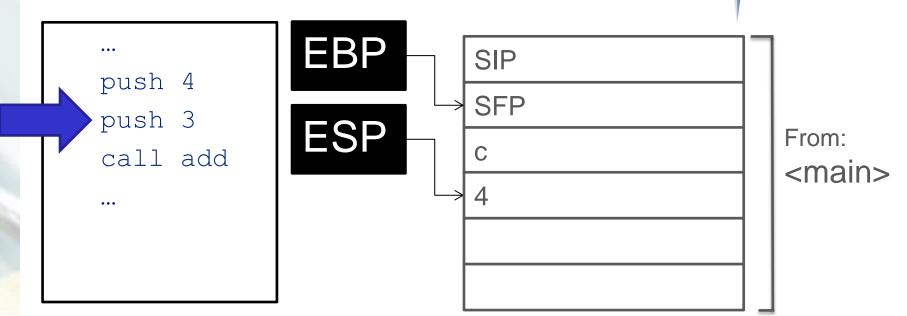
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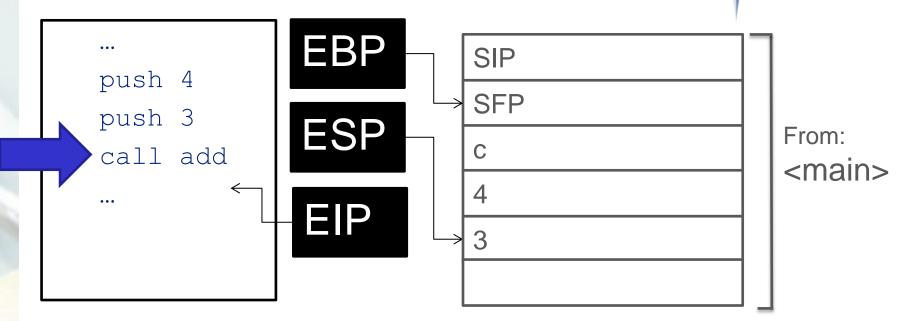






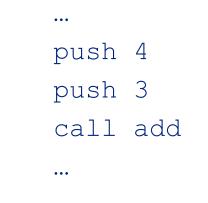


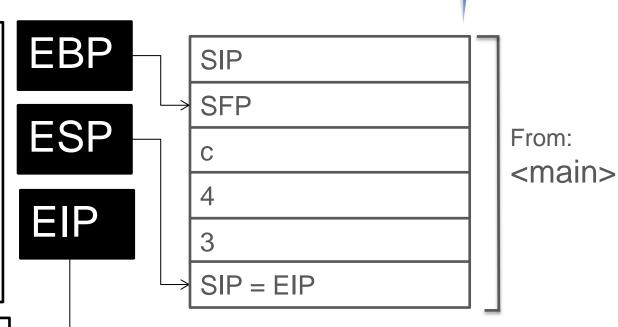






Slide 34



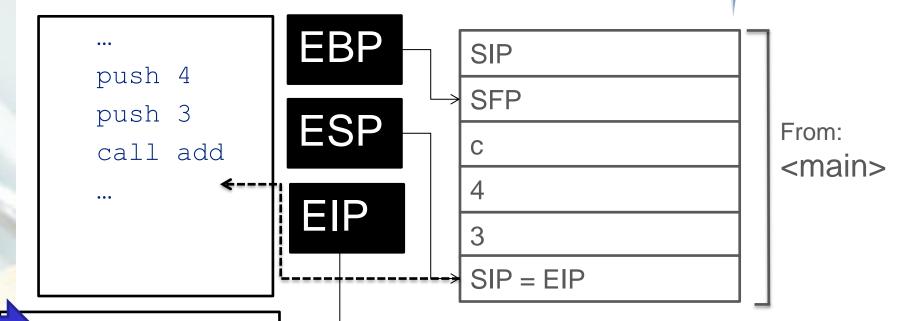


push ebp ←
mov ebp, esp,
sub esp, 0x10

mov esp, ebp
pop ebp
pop eip

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push ebp ←
mov ebp, esp,
sub esp, 0x10

mov esp, ebp
pop ebp

pop eip

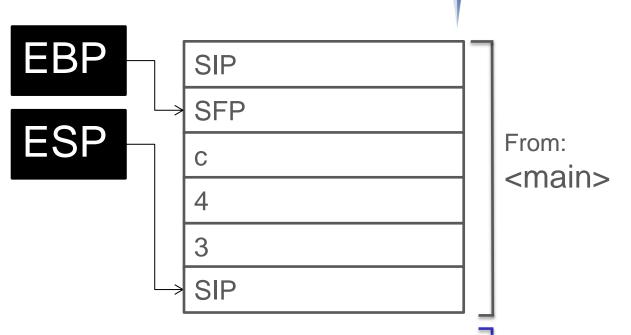
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Slide 35



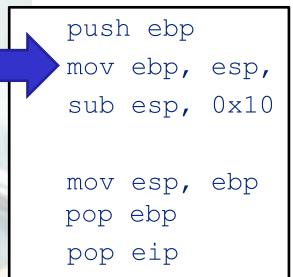
mov ebp, esp,
sub esp, 0x10

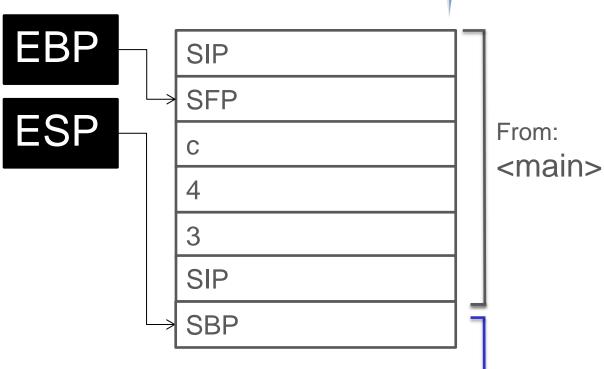
mov esp, ebp
pop ebp
pop eip



From: <add>







From: <add>



push ebp

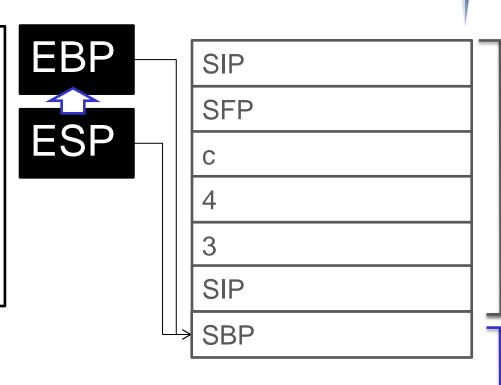
mov ebp, esp,

sub esp, 0x10

mov esp, ebp

pop ebp

pop eip



From: <main>

From: <add>



push ebp

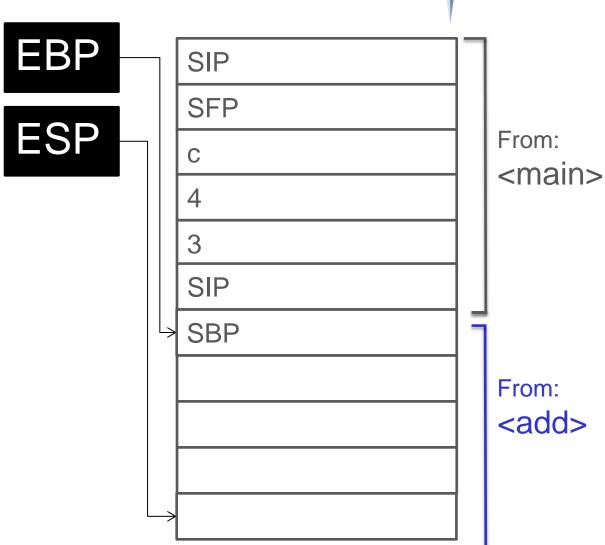
mov ebp, esp,

sub esp, 0x10

mov esp, ebp

pop ebp

pop eip







Execute Function

Compass Security Schweiz AG Tel +41 55 214 41 60 Werkstrasse 20 Postfach 2038 CH-8645 Jona

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x32 Call Convention - Execute Function



EBP

mov eax, DWORD PTR [ebp + 0xc]
mov edx, DWORD PTR [ebp + 0x8]
add eax, edx
mov DWORD PTR [ebp - 0x04], eax
mov eax, DWORD PTR [ebp - 0x04]

SIP SFP EBP+0xc EBP+0x8 SIP SBP EBP-0x04 sum

From: <main>

From: <add>





Function Epilog

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push ebp

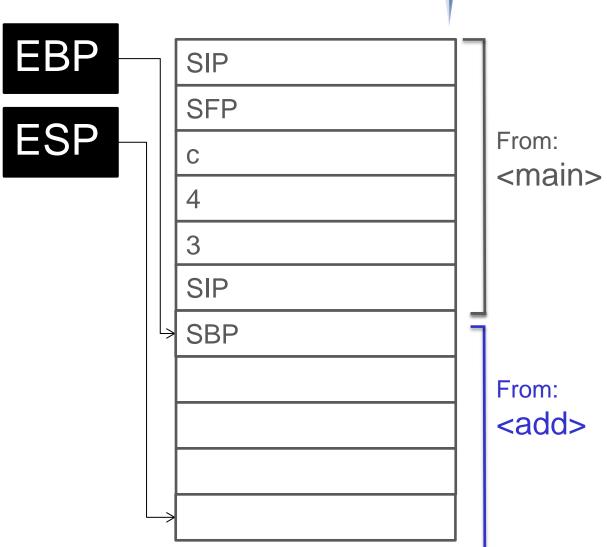
mov ebp, esp,

sub esp, 0x10

mov esp, ebp

pop ebp

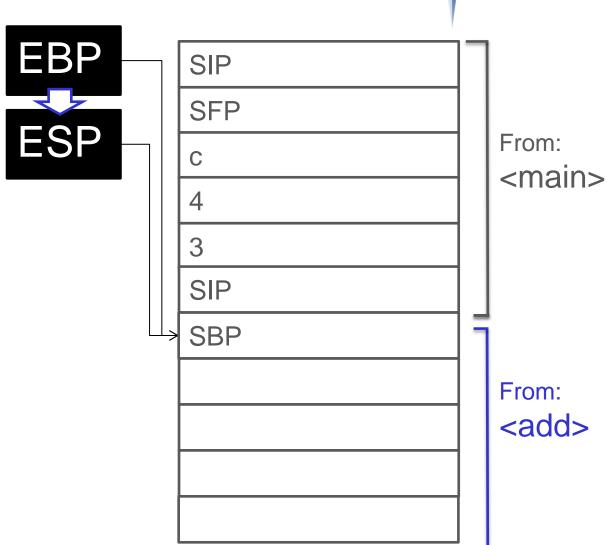
pop eip



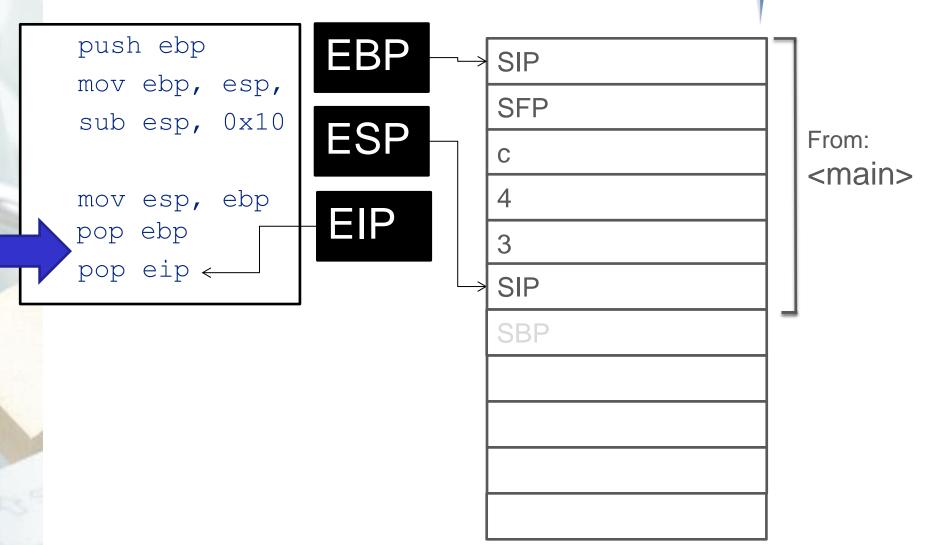


push ebp
mov ebp, esp,
sub esp, 0x10

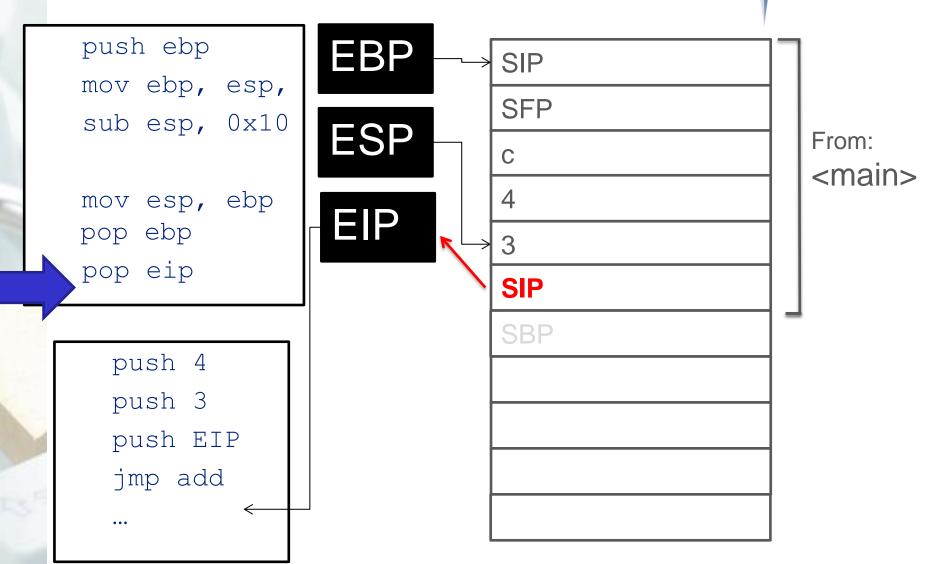
mov esp, ebp
pop ebp
pop eip



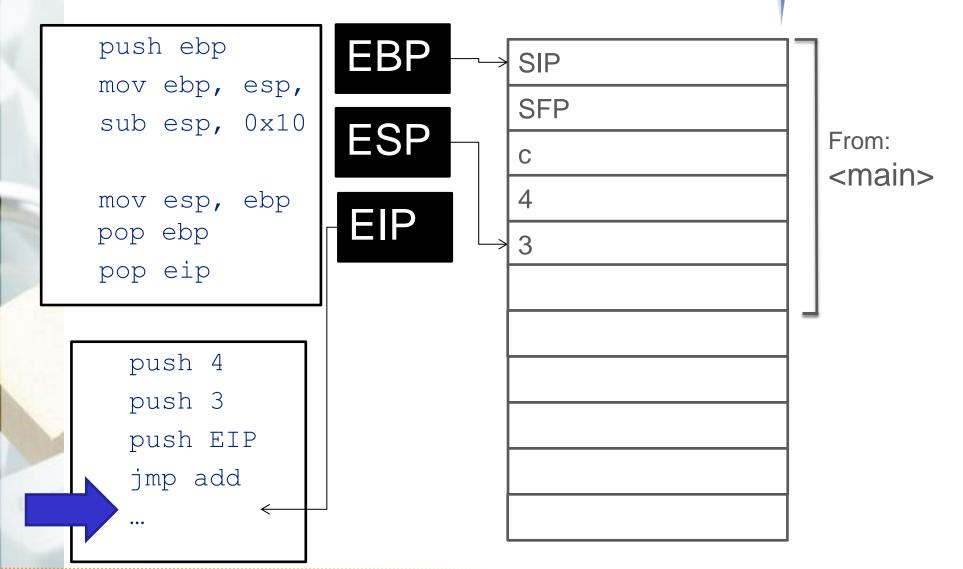












x32 Call Convention - Function Calling



```
call <addr> =
 push EIP
  jmp <addr>
leave =
 mov esp, ebp
 pop ebp
ret =
 pop eip
```

x32 Call Convention - Function Calling



Why "leave"?

→ Opposite of "enter"

"enter":

```
push ebp
mov ebp, esp
sub esp, imm
```

Why no "enter" used?

- enter:
 - ★ 8 cycle latency
 - → 10-20 micro ops
- call <addr>; mov ebp, esp; sub esp, imm:
 - → 3 cycles latency
 - → 4-6 micro ops





Function Call in x64

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x32 Call Convention - Function Call in x64



Differences between x32 and x64 function calls:

Arguments are in registers (not on stack) RDI, RSI, RDX, R8, R9

x32 Call Convention - Function Call in x64



Differences between x32 and x64 function calls

Different ASM commands doing the same thing

```
callq (call)
leaveq (leave)
retq (ret)
```

x32 Call Convention - Function Calling



Recap:

- → When a function is called:
 - → EIP is pushed on the stack (=SIP)
 - → ("call" is doing implicit "push EIP")
- At the end of the function:
 - → SIP is recovered into EIP
 - → ("ret" is doing implicit "pop EIP")

Function Call Convention Cheat Sheet



x32	Parameter	Syscall nr in
x32 userspace	stack	
x32 syscalls	ebx, ecx, edx, esi, edi, ebp	eax

x64	Parameter	Syscall nr in
x64 userspace	rdi, rsi, rdx, rcx, r8, r9	
x64 syscall	rdi, rsi, rdx, r10, r8, r9	rax

http://stackoverflow.com/questions/2535989/what-are-the-calling-conventions-for-unix-linux-system-calls-on-x86-64