

Software Verification

Lecture 01. Introduction to Software Analysis

Ji, Yong-Hyeon

24. 07. 11 (Thu)

Coding & Optimization Together (CO2)

Crypto & Security Engineering Lab (CSE)

Department of Information Security, Cryptology, and Mathematics

Table of Contents

1. Motivation
2. Software Analysis
 - 2.1 Software Analysis
 - 2.2 A Hard Limit
 - 2.3 Trade-off
3. Basic Principle
 - 3.1 SW Analysis based on Concrete Execution
 - 3.2 SW Analysis based on Symbolic Execution
 - 3.3 SW Analysis based on Abstract Execution (★)
4. Abstract Interpretation
 - Principles of Abstract Interpretation
5. Summary

Motivation

1. Motivation

1. Software systems are mostly unsafe.

2. Software errors cost the U.S. economy \$60B (≈ 82 조) every year.

(1996) Explosion of the Arian-5 rocket. \$8 billion (≈ 11 조)

(1998) NASA's Mars climate orbiter lost in space. \$125 million (≈ 1700 억)

(2000) Accidents in radiation therapy system. 8 patients died

(2007) Air control system shutdown in LA airport. 6,000 passengers stranded

(2012) Glitch in trading software of Knight Capital. \$440 million (≈ 6000 억)

(2014) Airbag malfunction of Nissan vehicles. \$1 million vehicles recalled

1. Motivation

3. Current Technologies for Safe Software

(a) Code Review

(b) Testing

(c) Debugging

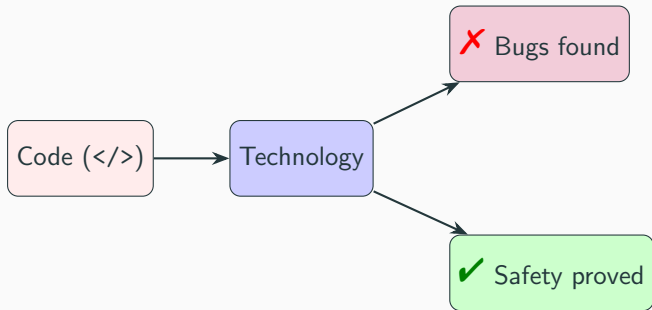
(d) Simulation

(e) ...

Software Analysis

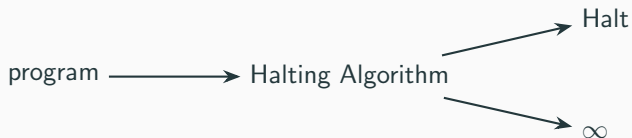
2.1 Software Analysis

- Technology for catching bugs or proving correctness of software.



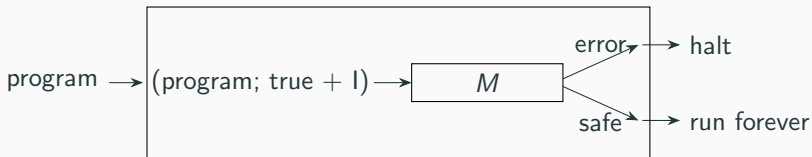
2.2 A Hard Limit

- The Halting problem is not computable. (Alan Turing, 1936)



2.2 A Hard Limit

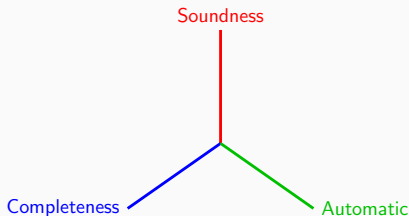
- If exact analysis is possible, we can solve the Halting problem.



- Rice's Theorem (1951): any non-trivial semantic property of a program is undecidable.

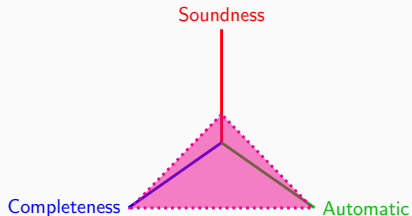
2.3 Trade-off

- Three desirable properties
 - **Soundness**: all program behaviors are captured
 - **Completeness**: only program behaviors are captured
 - **Automation**: without human intervention
- Achieving all of them is generally infeasible

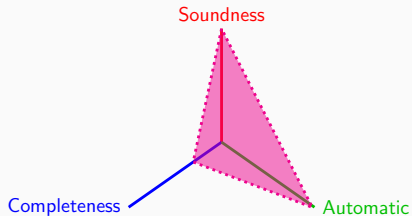


2.3 Trade-off

- Completeness + Automation (Testing)

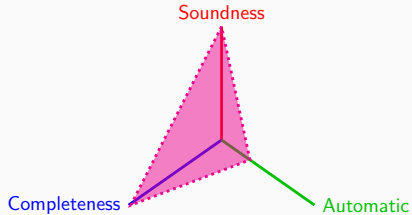


- Sound + Automation (Static Analysis, Compiler)

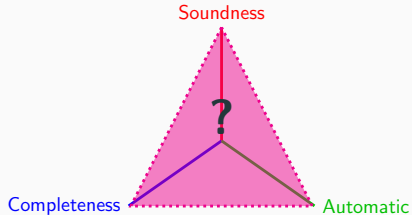


2.3 Trade-off

- Soundness + Completeness (Program Verification)



- And ...



Basic Principle

3. Basic Principle

- Observe the program behavior by “executing” the program
 - Report errors found during the execution
 - When no error is found, report “verified”
- Three types of program execution:
 - **Concrete execution**
 - **Symbolic execution**
 - **Abstract execution**
 - and their combinations, e.g., concolic execution

3. Basic Principle

- Concrete

$$2 * 3 - 6 = 0$$

- Symbolic

$$a * b + (-c) = 0$$

- Abstract

$$(\mathbb{R}, +, *)$$

3.1 SW Analysis based on Concrete Execution

Software Analysis based on Concrete Execution (Testing)

- Execute the program with concrete inputs, analyzing individual program states separately.

3.1 SW Analysis based on Concrete Execution

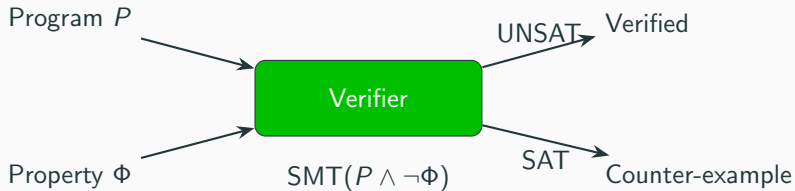
- Error-triggering test? (Test Vector)
- Probability of the error? (assume $0 \leq x, y \leq 10,000$)
- **Types of Fuzzing**
 - Blackbox Fuzzing
 - Greybox Fuzzing
 - Whitebox Fuzzing

3.2 SW Analysis based on Symbolic Execution

Software Analysis based on Symbolic Execution

- Execute the program with symbolic inputs, analyzing each program path only once.

3.2.1 Example: Symbolic Verification



- Represent program behavior and property as a formula in logic.
- Determine the satisfiability of the formula.

3.3 SW Analysis based on Abstract Execution (★)

Software Analysis based on Abstract Execution (Static Analysis)

- Execute the program with abstract inputs, analyzing all program behaviors simultaneously.

Abstract Interpretation

4.1 Principles of Abstract Interpretation

$$30 \times 12 + 11 \times 9 = ?$$

- Dynamic Analysis (Testing): 459
- Static Analysis: a Variety of Answers
 - “integer”, “odd integer”, positive integer”, “ $400 \leq n \leq 500$ ”, etc.

4.1 Principles of Abstract Interpretation

- Static Analysis Process:

1. Choose abstract value (domain), e.g., $\hat{V} = \{\top, e, o, \perp\}$
2. Define the program execution in terms of abstract values:

$\hat{\times}$	\top	e	o	\perp
\top	\top	e	\top	\perp
e	e	e	e	\perp
o	\top	e	o	\perp
\perp	\perp	\perp	\perp	\perp

$\hat{+}$	\top	e	o	\perp
\top	\top	\top	\top	\perp
e	\top	e	o	\perp
o	\top	o	e	\perp
\perp	\perp	\perp	\perp	\perp

3. “Execute” the program:

$$e \hat{\times} e \hat{+} o \hat{\times} o = o$$

Summary

Summary: Software Analysis

- Basically classified based on how programs are interpreted:
 - Techniques based on concrete execution
 - Techniques based on symbolic execution
 - Techniques based on abstract execution
- Each approach has its own strengths and weaknesses: e.g.,

	Sound	Complete	Automatic	When
Testing/Fuzzing	✗	✓	✓	Dynamic (Runtime)
Symbolic Execution	✗	✓	✓	Static/Dynamic
Static Analysis	✓	✗	✓	Static
Program Verification	✓	✓	✗	Static
?	⋮	⋮	⋮	⋮

Questions?