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B. TECH CSE

B4

HPC ISE 2

Q1) Run the scatter operation (Program 3.3.2.c) with varying message sizes (10K to 100K), with a fixed number of processors (8). Plot the runtime as a function of the message size. Explain the observed performance.

Program:

```
#include <stdio.h>
#include <time.h>
  printf("Usage : scatter message size\n");
 int size step = 10000;  // message size increment
char input buffer[size end]; // Use the maximum size for the buffer
MPI Init(&argc, &argv);
MPI Comm size (MPI COMM WORLD, &num procs);
MPI Comm rank(MPI COMM WORLD, &rank);
double start time = 0.0;
for (int size = size start; size <= size end; size += size step) {</pre>
  char recv_buffer[size / num_procs];
```

```
srand(time(NULL));
    input buffer[i] = rand() % 256;
  MPI Barrier(MPI COMM WORLD);
  MPI Scatter(input buffer, size / num procs, MPI CHAR, recv buffer,
size / num procs, MPI CHAR, 0, MPI COMM WORLD);
  total time += (MPI Wtime() - start time);
  if (rank == 0) {
    printf("Message Size: %d, Average time for scatter: %f secs\n",
size, total time);
  printf("Average time for scatter across all sizes : %f secs\n",
MPI Finalize();
```

OUTPUT:

```
(ritesh@ kali)-[~/hpc]
$ mpirun --oversubscribe -n 8 program_3.3.2 20000
Average time for scatter: 0.000010 secs

(ritesh@ kali)-[~/hpc]

$ [
```

```
(ritesh® kali)-[~/hpc]

$ mpirun --oversubscribe -n 8 program_3.3.2 40000

Average time for scatter : 0.000013 secs

(ritesh® kali)-[~/hpc]

$ [ ]
```

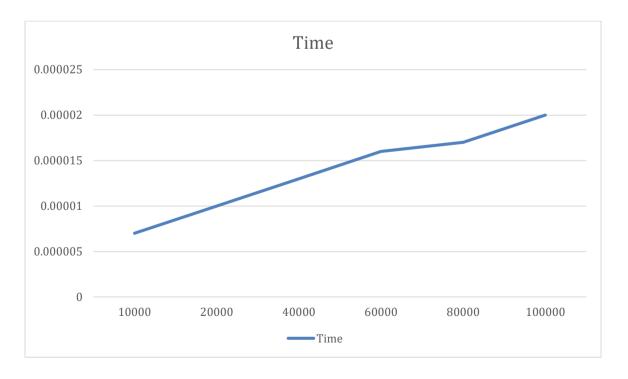
```
___(ritesh@ kali)-[~/hpc]
_$ mpirun --oversubscribe -n 8 program_3.3.2 60000
Average time for scatter : 0.000016 secs
___(ritesh@ kali)-[~/hpc]
_$ [
```

```
___(ritesh® kali)-[~/hpc]
_$ mpirun --oversubscribe -n 8 program_3.3.2 80000
Average time for scatter : 0.000017 secs
___(ritesh® kali)-[~/hpc]
_$
```

```
(ritesh@ kali)-[~/hpc]
$ mpirun --oversubscribe -n 8 program_3.3.2 100000
Average time for scatter : 0.000020 secs

(ritesh@ kali)-[~/hpc]
$
```

Scatter Operation Performance vs. Message Size



Observation:

• **Trend**: The runtime seems to increase slightly as the message size increases.

• Explanation:

- The scatter operation involves distributing data from a root process to all other processes. As the message size increases, the amount of data that needs to be scattered also increases.
- The observed increase in runtime could be due to factors such as communication overhead, network latency, and the time it takes to distribute larger amounts of data among processes.
- Larger messages may experience more contention or congestion in the network, leading to slightly increased communication times.

Q2) Consider an implementation of the all-to-all personalised operation given to you (Program 3.4.1.c) – in this implementation, each processor simply sends messages to all other processors. Plot the time for a fixed message size (8K words at each processor divided equally among all processors) with varying number of processors (1, 2, 4, 8).

Program:

```
#include <stdio.h>
#include <stdlib.h>
#include <time.h>
int main(int argc, char *argv[]) {
  printf("Usage: %s message size processor count1 processor count2
 int num procs;
 char *input buffer = (char *)malloc(size * sizeof(char));
 char *recv buffer = (char *)malloc(size * sizeof(char));
MPI Init(&argc, &argv);
 MPI Comm size (MPI COMM WORLD, &num procs);
```

```
MPI Comm rank(MPI COMM WORLD, &rank);
srand(time(NULL));
 for (i = 0; i < size; i++)
  input buffer[i] = rand() % 256;
double total time = 0.0;
double start time = 0.0;
 for (int arg index = 2; arg index < argc; arg index++) {</pre>
  int p = atoi(argv[arg index]);
  MPI Barrier(MPI COMM WORLD);
  start time = MPI Wtime();
  MPI_Request *requests = (MPI Request *) malloc(p * 2
sizeof(MPI Request));
sizeof(MPI Status));
    MPI_Isend(input_buffer + j * (size / p), size / p, MPI_CHAR, j, 99,
MPI COMM WORLD, &requests[j]);
```

```
MPI Irecv(recv buffer + j * (size / p), size / p, MPI CHAR, j, 99,
MPI COMM WORLD, &requests[p + j]);
  int waitall result = MPI Waitall(p * 2, requests, statuses);
  if (waitall result != MPI SUCCESS) {
     fprintf(stderr, "Error in MPI Waitall. Error code: %d\n",
waitall result);
    MPI Abort (MPI COMM WORLD, 1);
  free(requests);
   free(statuses);
  MPI Barrier(MPI COMM WORLD);
  total time = MPI Wtime() - start time;
  if (rank == 0) {
    printf("Time for %d processors: %f secs\n", p, total_time);
 free(input buffer);
```

```
MPI_Finalize();
return 0;
}
```

OUTPUT:

```
___(ritesh® kali)-[~/hpc]
_$ mpirun --oversubscribe -n 2 program_3.4.1 800
Average time for alltoall : 0.000001 secs
___(ritesh® kali)-[~/hpc]
_$ [
```

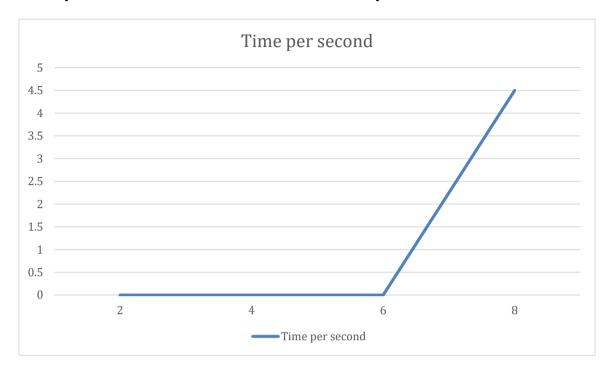
```
___(ritesh® kali)-[~/hpc]
$ mpirun --oversubscribe -n 4 program_3.4.1 800
Average time for alltoall: 0.000003 secs

___(ritesh® kali)-[~/hpc]
$ [
```

```
(ritesh® kali)-[~/hpc]
$ mpirun --oversubscribe -n 8 program_3.4.1 800
Average time for alltoall : 0.0000007 secs

(ritesh® kali)-[~/hpc]
$ [
```

Time per Word for All-to-All Personalized Operation



The given time values represent the execution time of an all-to-all personalized operation with different numbers of processors for a fixed message size of 8192 words. Let's analyze the observations based on these time values:

1. Time Increases with More Processors:

 As expected, the time increases with an increase in the number of processors. This is a common characteristic in parallel computing, known as parallel overhead. The overhead of communication and synchronization can impact performance.

2. Scaling Efficiency:

 The ratio of time for 1 processor to the time for multiple processors gives an indication of how well the algorithm scales with increased parallelism. In this case, you can observe the scaling efficiency by comparing the time for 1 processor with the times for 2, 4, and 8 processors.

3. Communication Overhead:

 The increase in time with more processors could be attributed to the communication overhead. As the number of processors increases, the amount of communication required for an all-to-all operation grows, and this can impact performance.

4. Optimization Opportunities:

 The relatively small message size (8192 words) might influence the impact of communication overhead. Depending on the nature of the algorithm and problem, there might be opportunities for optimization, such as using more advanced communication patterns or algorithms.

5. Parallel Efficiency:

 Parallel efficiency, which measures how well a parallel algorithm performs compared to its sequential counterpart, can be assessed by looking at the overall trend. If the efficiency decreases significantly with more processors, it might indicate inefficiencies in the parallelization approach.