

## Linear Dynamical Systems - Summary

*Lecturer:* Dr. Oron Sabag*Name:* Hadar Tal**Contents**

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# 1 Mathematical Tools

## 1.1 Completing The Square

### Lemma 1.1 (Completing the Square (scalars))

To complete the square for a quadratic equation of the form

$$ax^2 + bx + c = 0,$$

we can rewrite it as

$$a(x + d)^2 + e = 0,$$

where

$$d = \frac{b}{2a} \quad \text{and} \quad e = c - \frac{b^2}{4a}.$$

*Proof.*

$$\begin{aligned} ax^2 + bx + c = 0 & \rightarrow x^2 + \frac{b}{a}x + \frac{c}{a} = 0 & \rightarrow x^2 + \frac{b}{a}x = -\frac{c}{a} & \rightarrow \\ x^2 + \frac{b}{a}x + \left(\frac{b}{2a}\right)^2 = -\frac{c}{a} + \left(\frac{b}{2a}\right)^2 & \rightarrow \left(x + \frac{b}{2a}\right)^2 = \frac{b^2}{4a^2} - \frac{c}{a} & \rightarrow \left(x + \frac{b}{2a}\right)^2 = \frac{b^2 - 4ac}{4a^2} & \rightarrow \\ a\left(x + \frac{b}{2a}\right)^2 = a\left(\frac{b^2 - 4ac}{4a^2}\right) & \rightarrow a\left(x + \frac{b}{2a}\right)^2 = \frac{b^2 - 4ac}{4a} & \rightarrow a\left(x + \frac{b}{2a}\right)^2 + c - \frac{b^2}{4a} = 0 \end{aligned}$$

□

### Lemma 1.2 (Completing the Square for Quadratic Forms)

Given a quadratic form  $x^T Ax + b^T x + c$ , where  $A$  is a symmetric positive definite matrix,  $b$  is a vector, and  $c$  is a scalar, the expression can be completed to a perfect square as follows:

$$x^T Ax + b^T x + c = (x + A^{-1}b/2)^T A(x + A^{-1}b/2) + c - \frac{1}{4}b^T A^{-1}b$$

## 2 The Shortest Path Problem

### Definition 2.1 (The Shortest Path Problem)

Given a directed graph  $G = (V, E)$ , each edge in the graph has a cost  $a_{ij}^t$ , where  $i$  is the outgoing node,  $j$  is the node to which the edge is connected, and  $t \in \{0, \dots, N+1\}$  refers to the time. We adopt the convention that no edge implies an infinite cost  $a_{ij}^t = \infty$ .

The objective is to minimize the cumulative cost on a path from the source node  $S_0 = S$  to the terminal node  $S_{N+1} = T$ . Formally, we aim to solve the optimization

$$J^* = \min_{\{n_i \in \mathcal{S}_i\}_{i=0}^{N+1}} \sum_{t=0}^N a_{n_t, n_{t+1}}^t. \quad (2.1)$$

### Definition 2.2 (Cost-to-Go Function)

We define  $J_k(i)$  as the cost-to-go function corresponding to the minimal cost from time  $k$  until the end when starting at node  $i$ . Formally, for  $k = 0, \dots, N$ , define

$$J_k(i) = \min_{\{n_j \in \mathcal{S}_j | j=k+1, \dots, N, n_k=i\}} \sum_{j=k}^N a_{n_j, n_{j+1}}^j, \quad \forall i \in \mathcal{S}_k. \quad (2.2)$$

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#### Algorithm 1: Dynamic Programming Solution for the Shortest Path Problem (Cost-to-Go)

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**Input:** Cost matrix  $a_{ij}^t$  and nodes  $\mathcal{S}_0, \mathcal{S}_1, \dots, \mathcal{S}_{N+1}$

**Output:** Cost-to-go functions  $J_k(i)$

Initialize  $J_N(i) = a_{iT}^N$  ;

**for**  $k = N-1, \dots, 0$  **do**

**for**  $i \in \mathcal{S}_k$  **do**

$J_k(i) = \min_{j \in \mathcal{S}_{k+1}} [a_{ij}^k + J_{k+1}(j)]$

**end**

**end**

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### Definition 2.3 (Cost-to-Arrive Function)

We define  $J_{N-k}(j)$  as the cost-to-arrive function corresponding to the minimal cost from time 1 until time  $k$  when arriving at node  $j$ . Formally, for  $k = 0, \dots, N$ , define

$$J_{N-k}(j) = \min_{\{n_i \in \mathcal{S}_i | i=1, \dots, k-1, n_k=j\}} \sum_{i=1}^k a_{n_{i-1}, n_i}^i, \quad \forall j \in \mathcal{S}_k. \quad (2.3)$$

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#### Algorithm 2: Forward Algorithm for the Shortest Path Problem (Cost-to-Arrive)

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**Input:** Cost matrix  $a_{ij}^t$  and nodes  $\mathcal{S}_0, \mathcal{S}_1, \dots, \mathcal{S}_{N+1}$

**Output:** Cost-to-arrive functions  $J_{N-k}(j)$

Initialize  $J_N(j) = a_{sj}^0, \forall j \in \mathcal{S}_1$  ;

**for**  $k = 1, \dots, N$  **do**

**for**  $j \in \mathcal{S}_{N-k+1}$  **do**

$J_k(j) = \min_{i \in \mathcal{S}_{N-k}} [a_{ij}^{N-k} + J_{k+1}(i)]$

**end**

**end**

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### 3 Markov Decision Processes (MDPs)

#### Definition 3.1 (MDP)

An MDP is defined by the following elements:

1. The state at time  $k$  is  $x_k$  and takes values in the set  $\mathcal{S}_k$ .
2. The action at time  $k$  is  $u_k$  and takes values from  $\mathcal{U}_k$ .
3. The disturbance at time  $k$  is  $w_k$  and takes values from  $\mathcal{W}_k$ .
4. A dynamical system is given by the function

$$x_{k+1} = f_k(x_k, u_k, w_k), \quad k = 0, \dots, N-1. \quad (3.1)$$

5. The probabilistic law of the disturbance random variable  $w_k$  is characterized by  $P_{W_k}(\cdot|x_k, u_k)$  conditioned on the state  $x_k$  and the action  $u_k$ .
6. A cost function  $g_k : \mathcal{S}_k \times \mathcal{U}_k \rightarrow \mathbb{R}$ .

The cost over a horizon  $N$  is

$$S_\pi(x_0) = \mathbb{E}[g_N(x_N) + \sum_{k=0}^{N-1} g_k(x_k, u_k)], \quad (3.2)$$

where  $g_N(\cdot)$  is the terminal cost.

#### Definition 3.2 (History-dependent policy)

A history-dependent policy is defined by a sequence of functions:

$$\mu_k : \mathcal{S}_1 \times \dots \times \mathcal{S}_k \times \mathcal{U}_1 \times \dots \times \mathcal{U}_{k-1} \rightarrow \mathcal{U}_k, \quad (3.3)$$

such that  $u_k = \mu_k(x_1, x_2, \dots, x_k, u_1, \dots, u_{k-1})$ .

#### Definition 3.3 (Markovian policy)

A Markovian policy is defined by a sequence of functions:

$$\mu_k : \mathcal{S}_k \rightarrow \mathcal{U}_k, \quad (3.4)$$

such that  $u_k = \mu_k(x_k)$ .

#### Remark 3.1 (The Markov property)

We defined the dynamical system using a deterministic function  $f_k(\cdot)$ .

Equivalently, we could describe the evolution with the conditional probability

$$P_k(x_{k+1}|x_k, u_k) = P_s^u(s') \quad (3.5)$$

In particular, we assume that the new state conditioned on the current state and action is not affected by the past. Formally, we assume the Markov chain induced from

$$P(x_{k+1}|x_1, \dots, x_k, u_1, \dots, u_k) = P_k(x_{k+1}|x_k, u_k). \quad (3.6)$$

The MDP described above is called fully observable since the actions depend directly on the state. We will later encounter partially observable MDP where only a noisy version of the state is available to the controller.

## 4 Linear Systems

### Definition 4.1 (Linear System)

A linear system is given by

$$x_{t+1} = Ax_t, \quad t = 0, 1, \dots \quad (4.1)$$

with some initial state  $x_0$ .

- $x_t \in \mathbb{R}^n$  is the state vector.
- $A \in \mathbb{R}^{n \times n}$  is the state-transition matrix.

### Lemma 4.1 (State and Decoupling in Linear Systems)

Given a diagonalizable matrix  $A = TDT^{-1}$ , the state at time  $t$  in a linear system is

$$x_t = TD^tT^{-1}x_0. \quad (4.2)$$

By defining a new state  $z_t = T^{-1}x_t$ , we have

$$z_t = D^tz_0, \quad (4.3)$$

indicating that the states are decoupled, with each entry of  $z_t$  depending only on the corresponding entry of  $z_0$ .

### Remark 4.1

Since the eigenvalues of real matrices may be complex, we have

$$\lambda = a + ib = re^{i\theta} \rightarrow \lambda^t = r^te^{it\theta} (e^{i\theta} = \cos \theta + i \sin \theta).$$

As we increase  $t$ , the magnitude of  $e^{it\theta}$  is clearly unchanged. However, the length of  $r$  determines whether it converges to zero, oscillates, or blows up.

### Definition 4.2 (Stable System)

A system  $A$  is stable if all of its eigenvalues have magnitude smaller than 1, i.e.,  $r < 1$ .

## 5 Linear Systems with Control

### Definition 5.1 (Linear System with Control)

A linear system with control is given by

$$x_{t+1} = Ax_t + Bu_t, \quad t \geq 0, \quad x_0 \in \mathbb{R}^n, \quad (5.1)$$

where we added:

- $u_t \in \mathbb{R}^m$  is the control signal (action).
- $B \in \mathbb{R}^{n \times m}$  is the control matrix.

### Definition 5.2 (State-feedback controller)

A controller (policy) is defined by a sequence of mappings  $\mu_t : \mathbb{R}^n \rightarrow \mathbb{R}^m$  for  $t = 0, 1, \dots, N$  such that  $u_t = \mu_t(x_t)$ .

### Definition 5.3 (State-Feedback, Time-Invariant, Linear Controller)

A state-feedback, time-invariant, linear controller is any mapping of the form

$$u_t = -Kx_t.$$

### Definition 5.4 (Closed-Loop Matrix)

The matrix  $A_K = A - BK$  is called the closed-loop matrix of the system  $A, I.H.T$  -

$$x_{t+1} = A_K x_t = (A - BK)x_t = Ax_t + B(-Kx_t) = Ax_t + Bu_t$$

### Definition 5.5 (Controllability)

The pair  $(A, B)$  is controllable if the system can reach any  $\xi \in \mathbb{R}^n$  from any initial state  $x_0 \in \mathbb{R}^n$  at some finite time.

### Lemma 5.1 (Controllability Matrix)

A pair  $(A, B)$  is controllable if and only if the controllability matrix

$$C \triangleq [B \quad AB \quad \dots \quad A^{n-1}B]$$

has  $\text{rank}(C) = n$ .

### Lemma 5.2 (Poles Placement in Controllable System)

Controllability implies that we can choose the eigenvalues (poles) of  $A - BK$  arbitrarily.

## 6 The Linear Quadratic Regulator (LQR)

### Definition 6.1 (The LQR problem)

For the linear model in (20), find a controller that minimizes

$$J_N(u^N) = \sum_{i=0}^N [x_i^T Q x_i + u_i^T R u_i] + x_{N+1}^T Q_f x_{N+1},$$

where  $u^N \triangleq u_0, u_1, \dots, u_{N-1}$ , and

1.  $Q, Q_f \succeq 0$  are state weights
2.  $R \succ 0$  is the input/action/control weight.

### Lemma 6.1 (LQR matrix formulation)

$$\begin{bmatrix} x_0 \\ \vdots \\ x_N \end{bmatrix} = \begin{bmatrix} 0 & 0 & \cdots & 0 \\ B & 0 & \cdots & 0 \\ AB & B & \cdots & 0 \\ \vdots & \vdots & \ddots & \vdots \\ A^{N-1}B & A^{N-2}B & \cdots & B \end{bmatrix} \begin{bmatrix} u_0 \\ \vdots \\ u_{N-1} \end{bmatrix} + \begin{bmatrix} I \\ A \\ \vdots \\ A^N \end{bmatrix} x_0$$

The matrices above can be written in short as  $x = Gu + Hx_0$ .

### Lemma 6.2 (LQR Closed-Form Solution)

The optimal control input  $\mathbf{u}$  that minimizes the cost function  $\mathbf{u}^T \mathbf{u} + \mathbf{x}^T \mathbf{x}$  can be achieved with

$$\mathbf{u} = -(I + G^T G)^{-1} G^T H x_0.$$

*Proof.*

$$\begin{aligned} \min_{\mathbf{u}} \mathbf{x}^T \mathbf{x} + \mathbf{u}^T \mathbf{u} &= \min_{\mathbf{u}} (G\mathbf{u} + Hx_0)^T (G\mathbf{u} + Hx_0) + \mathbf{u}^T \mathbf{u} \\ &= \min_{\mathbf{u}} [\mathbf{u}^T G^T G \mathbf{u} + 2x_0^T H^T G \mathbf{u} + x_0^T H^T H x_0 + \mathbf{u}^T \mathbf{u}] \\ &= \min_{\mathbf{u}} [\mathbf{u}^T (I + G^T G) \mathbf{u} + 2x_0^T H^T G \mathbf{u} + x_0^T H^T H x_0] \\ &= \min_{\mathbf{u}} [(\mathbf{u} + (I + G^T G)^{-1} G^T H x_0)^T (I + G^T G) (\mathbf{u} + (I + G^T G)^{-1} G^T H x_0) \\ &\quad - x_0^T H^T G (I + G^T G)^{-1} G^T H x_0 + x_0^T H^T H x_0] \\ &= \min_{\mathbf{u}} [(\mathbf{u} + (I + G^T G)^{-1} G^T H x_0)^T (I + G^T G) (\mathbf{u} + (I + G^T G)^{-1} G^T H x_0) \\ &\quad + x_0^T H^T (I - G(I + G^T G)^{-1} G^T) H x_0] \\ &= x_0^T H^T (I - G(I + G^T G)^{-1} G^T) H x_0 \\ &= x_0^T H^T (I + G G^T)^{-1} H x_0, \end{aligned}$$

where the optimal control input is

$$\mathbf{u} = -(I + G^T G)^{-1} G^T H x_0.$$

□

### Remark 6.1

This solution is the optimal solution. However, it is not efficient since we should compute the inverse of  $I + G G^T$  that grows linearly with  $N$ , i.e.,  $O(N^3)$  computations.

### Definition 6.2 (Cost-to-Go Function)

The cost-to-go function (Value-function) is defined as

$$V_t(z) = \min_{u_t, u_{t+1}, \dots, u_{N-1}} \left[ \sum_{i=t}^{N-1} (x_i^T Q x_i + u_i^T R u_i) + x_N^T Q_f x_N \right]$$

for  $x_t = z$ .

**Theorem 6.1 (Properties of the Value Function)**

The value function satisfies the following properties.

1.  $V_t(z)$  is a quadratic function (of the variable  $z$ ).  
That is, we can write  $V_t(z) = z^T P_t z$  with some  $P_t \succeq 0$ .
2. The optimal controller  $u_t$  is given by

$$u_t = -(R + B^T P_{t+1} B)^{-1} B^T P_{t+1} A x_t.$$

3. The sequence  $P_t$  can be computed recursively with

$$P_t = \begin{cases} Q_f & t = N \\ Q + A^T P_{t+1} A - A^T P_{t+1} B (R + B^T P_{t+1} B)^{-1} B^T P_{t+1} A & t < N. \end{cases}$$

Note that we can compute  $P_t$  and  $K_t$  offline prior to the control.  
(Tools used in the proof - completion of the square, induction.)

**Definition 6.3 (Stabilizability)**

The pair  $(A, B)$  is stabilizable if

$$\exists K \in \mathbb{R}^{m \times n} : \rho(A - BK) < 1,$$

where  $\rho(A - BK)$  denotes the spectral radius of  $A - BK$ , i.e., the largest absolute value of its eigenvalues.

An equivalent characterization of stabilizability is

$$\nexists(x, \lambda) \text{ s.t. } xA = \lambda x \wedge |\lambda| \geq 1 \wedge xB = 0.$$

In other words, we can control the unstable modes.

If we want  $J^*(x_0) < \infty$  for all  $x_0$ , a necessary and sufficient condition is that of stabilizability.



## 7 Kalman Filter

### Definition 7.1 (Kalman Filter State Space)

$$\begin{aligned} x_{i+1} &= Fx_i + Gw_i \\ y_i &= Hx_i + v_i \end{aligned} \quad (7.1)$$

where:

- $x_{i+1}$  is the state vector at time  $i + 1$ ,
- $x_i$  is the state vector at time  $i$ ,
- $y_i$  is the measurement vector at time  $i$ ,
- $w_i$  is the process noise (zero mean, uncorrelated),
- $v_i$  is the measurement noise (zero mean, uncorrelated).

### Definition 7.2 (Kalman Filter Covariance Matrix)

Formally, the following covariance matrix describes the model:

$$\mathbb{E} \left[ \begin{pmatrix} w_i \\ v_i \\ x_0 \end{pmatrix} \begin{pmatrix} w_j^* & v_j^* & x_0^* & 1 \end{pmatrix} \right] = \begin{pmatrix} \begin{pmatrix} Q & S \\ S^* & R \end{pmatrix} \delta_{ij} & 0 & 0 \\ 0 & \Pi_0 & 0 \end{pmatrix}, \quad (7.2)$$

where  $\begin{pmatrix} Q & S \\ S^* & R \end{pmatrix}$  and  $\Pi_0$  are positive semidefinite matrices and  $\delta_{ij}$  equals 1 if  $i = j$  and is zero otherwise. Note that  $w_i$  is uncorrelated as a process over time but its coordinates at a fixed time can be correlated via  $Q$ .

**Markings:**

- $P_i$  - The error covariance matrix at time  $i$

$$P_i \triangleq (x_i - \hat{x}_i)(x_i - \hat{x}_i)^T \quad (7.3)$$

- $R_{e,i}$  - The covariance of the innovation (or residual) at time  $i$

$$R_{e,i} \triangleq HP_iH^* + R \quad (7.4)$$

- $K_{p,i}$  - The optimal Kalman gain at time  $i$

$$K_{p,i} \triangleq (FP_iH^* + GS)R_{e,i}^{-1} \quad (7.5)$$

### Kalman Filter Optimality

We suggest the following predictor:

$$\hat{x}_{i+1|i} = F\hat{x}_{i|i-1} + K_{p,i}(y_i - H\hat{x}_{i|i-1}) \quad (7.6)$$

#### Lemma 7.1

$$\tilde{x}_{i+1} = (F - K_{p,i}H)\tilde{x}_i + (G - K_{p,i}) \begin{pmatrix} w_i \\ v_i \end{pmatrix}. \quad (7.7)$$

*Proof.*

$$\begin{aligned} \tilde{x}_{i+1} &= x_{i+1} - \hat{x}_{i+1|i} \\ &= (Fx_i + Gw_i) - (F\hat{x}_i + K_{p,i}(y_i - H\hat{x}_i)) \\ &= Fx_i + Gw_i - F\hat{x}_i - K_{p,i}(Hx_i + v_i - H\hat{x}_i) \\ &= Fx_i + Gw_i - F\hat{x}_i - K_{p,i}Hx_i - K_{p,i}v_i + K_{p,i}H\hat{x}_i \\ &= Fx_i - F\hat{x}_i - K_{p,i}Hx_i + K_{p,i}H\hat{x}_i + Gw_i - K_{p,i}v_i \\ &= (F - K_{p,i}H)(x_i - \hat{x}_i) + Gw_i - K_{p,i}v_i \\ &= (F - K_{p,i}H)\tilde{x}_i + Gw_i - K_{p,i}v_i. \end{aligned}$$

□

**Lemma 7.2**

For  $j < i$ , the recursion can be evolved as

$$\begin{aligned}\tilde{x}_i &= (F - K_{p,i-1}H)\tilde{x}_{i-1} + (G - K_{p,i-1}) \begin{pmatrix} w_{i-1} \\ v_{i-1} \end{pmatrix} \\ &= \dots \\ &= \phi_p(i, j)\tilde{x}_j + \xi_i(j),\end{aligned}$$

where

$$\begin{aligned}\phi_p(i, j) &= \prod_{k=j}^{i-1} (F - K_{p,k}H), \\ \xi_i(j) &= \sum_{k=j}^{i-1} \phi_p(i, k+1)(Gw_k - K_{p,k}v_k).\end{aligned}$$