

南京大学本科生实验报告

课程名称：计算机网络

任课教师：田臣/李文中

助教：

学院	计算机科学与技术系	专业（方向）	计算机科学与技术
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1. 实验名称

Lab 2: Learning Switch

2. 实验目的

使用 Switchyard 框架实现以太网自学习交换机的核心功能。

3. 实验内容

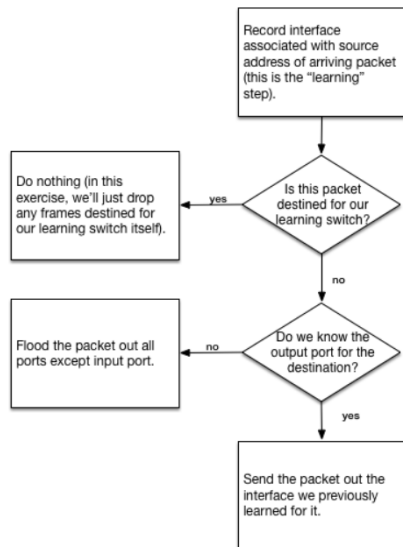
Task 1: Preparation

获取框架代码并将 myswitch.py 复制三份。

Task 2: Basic Switch

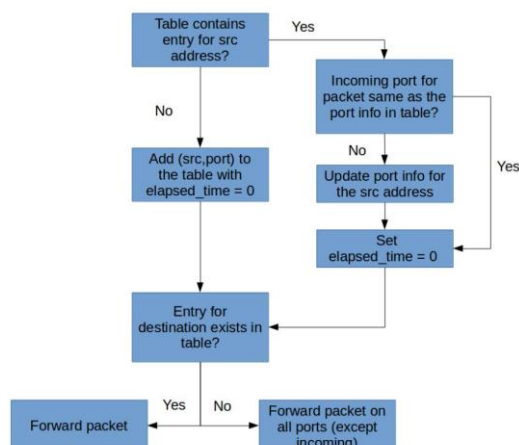
实现以太网交换机的学习、过滤、转发功能，建立交换机的转发表。

逻辑如下图：



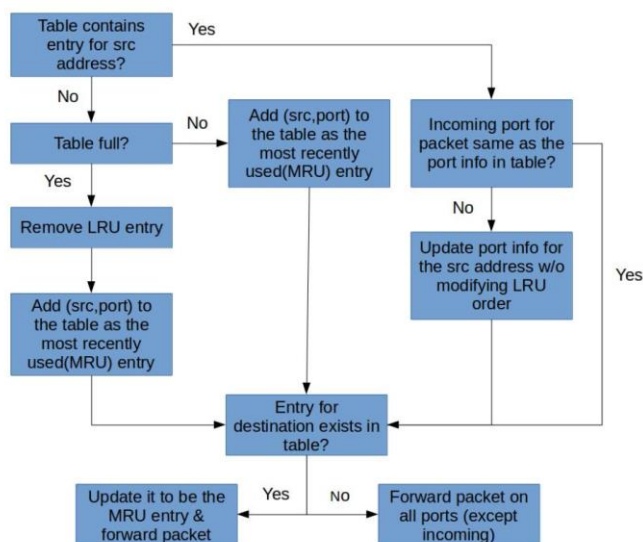
Task 3: Timeouts

在 task2 的基础上，实现转发表的删除机制：超过 10s 后将对应的表项删除。逻辑如下图：



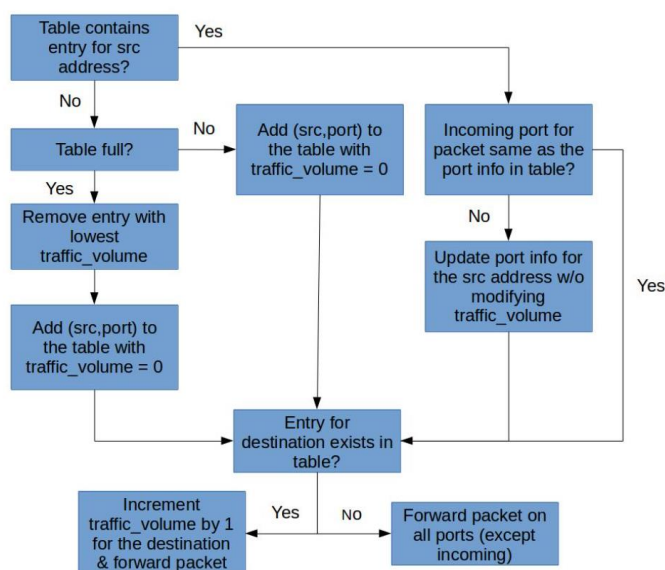
Task 4: Least Recently Used

在 task2 的基础上，实现转发表的删除机制：表项数最大为 5，超过 5 个后删除最近最不常用的表项。逻辑如下图：



Task 5: Least Traffic Volume

在 task2 的基础上，实现转发表的删除机制：表项数最大为 5，超过 5 个后删除流量最小的表项。逻辑如下图：



4. 实验结果

Task 2: Basic Switch

使用命令 ping 让 client 对 server1 ping 两次，使用 wireshark 对 server1 和 server2 进行抓包，可以发现 server1 有两次 ping 的 echo 包和一些 ARP 协议包：

No.	Time	Source	Destination	Protocol	Length	Info
1	0.000000000	30:00:00:00:00:01	Broadcast	ARP	42	Who has 192.168.100.1? Tell 192.168.100.3
2	0.100700951	Private_00:00:01	30:00:00:00:00:01	ARP	42	192.168.100.1 is at 10:00:00:00:00:01
3	0.409730561	192.168.100.3	192.168.100.1	ICMP	98	Echo (ping) request id=0x7e2a, seq=1/256, ttl=
4	0.509874329	192.168.100.1	192.168.100.3	ICMP	98	Echo (ping) reply id=0x7e2a, seq=1/256, ttl=
5	1.033961845	192.168.100.3	192.168.100.1	ICMP	98	Echo (ping) request id=0x7e2a, seq=2/512, ttl=
6	1.134366677	192.168.100.1	192.168.100.3	ICMP	98	Echo (ping) reply id=0x7e2a, seq=2/512, ttl=
7	5.701102035	Private_00:00:01	30:00:00:00:00:01	ARP	42	Who has 192.168.100.3? Tell 192.168.100.1
8	6.034591318	30:00:00:00:00:01	Private_00:00:01	ARP	42	192.168.100.3 is at 30:00:00:00:00:01

而 server2 只有一个 ARP 协议包：

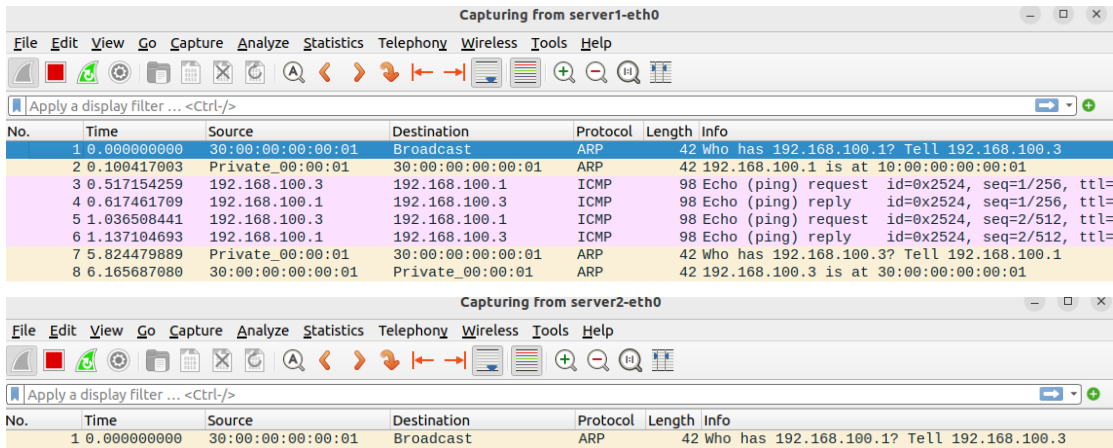
No.	Time	Source	Destination	Protocol	Length	Info
1	0.000000000	30:00:00:00:00:01	Broadcast	ARP	42	Who has 192.168.100.1? Tell 192.168.100.3

说明 client 通过广播 ARP 包来获取 server1 的 MAC 地址，而 server1 回应该广播，该过程使得交换机学习到了他们各自 MAC 地址对应的

端口，从而使得发送 echo 包的时候直接根据转发表进行转发而不再需要进行广播。

Task 3: Timeouts

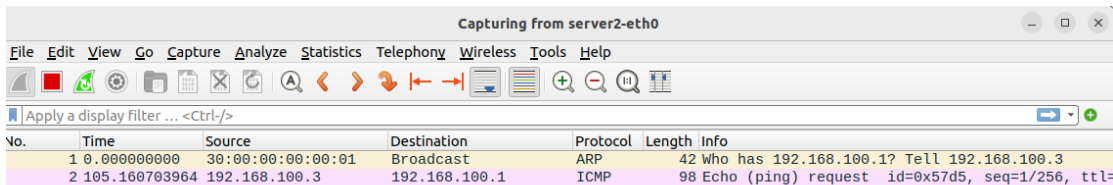
首先使用命令 ping 让 client 对 server1 ping 两次，使用 wireshark 对 server1 和 server2 进行抓包，可以发现结果和 task2 相同：



Wireshark capture from server1-eth0 showing traffic during the first ping test. The capture includes ARP requests and ICMP echo requests and replies.

No.	Time	Source	Destination	Protocol	Length	Info
1	0.000000000	30:00:00:00:00:01	Broadcast	ARP	42	Who has 192.168.100.1? Tell 192.168.100.3
2	0.100417003	Private_00:00:00:01	30:00:00:00:00:01	ARP	42	192.168.100.1 is at 30:00:00:00:00:01
3	0.517154259	192.168.100.3	192.168.100.1	ICMP	98	Echo (ping) request id=0x2524, seq=1/256, ttl=
4	0.617461709	192.168.100.1	192.168.100.3	ICMP	98	Echo (ping) reply id=0x2524, seq=1/256, ttl=
5	1.036508441	192.168.100.3	192.168.100.1	ICMP	98	Echo (ping) request id=0x2524, seq=2/512, ttl=
6	1.137104693	192.168.100.1	192.168.100.3	ICMP	98	Echo (ping) reply id=0x2524, seq=2/512, ttl=
7	5.824479889	Private_00:00:00:01	30:00:00:00:00:01	ARP	42	Who has 192.168.100.3? Tell 192.168.100.1
8	6.165687080	30:00:00:00:00:01	Private_00:00:00:01	ARP	42	192.168.100.3 is at 30:00:00:00:00:01

等待 10s 以上后，再次执行该命令，可以发现 server2 收到了一个 client 发送给 server1 的 echo 包：



Wireshark capture from server2-eth0 showing traffic during the second ping test. The capture includes ARP and ICMP echo requests.

No.	Time	Source	Destination	Protocol	Length	Info
1	0.000000000	30:00:00:00:00:01	Broadcast	ARP	42	Who has 192.168.100.1? Tell 192.168.100.3
2	105.160703964	192.168.100.3	192.168.100.1	ICMP	98	Echo (ping) request id=0x57d5, seq=1/256, ttl=

说明经过 10s 以上后，交换机遗忘了之前的表项需要重新学习，而 server2 只收到一个 request 包说明在第一个包之后交换机有重新学习了 server1 对应的端口，第二次 ping 并不需要再次广播。

测试结果如下：

```

Passed:
1  An Ethernet frame with a broadcast destination address
    should arrive on eth1
2  The Ethernet frame with a broadcast destination address
    should be forwarded out ports eth0 and eth0 and eth2
3  An Ethernet frame from 20:00:00:00:00:01 to
    30:00:00:00:00:02 should arrive on eth0
4  Ethernet frame destined for 30:00:00:00:00:02 should arrive
    on eth1 after self-learning
5  Timeout for 60s
6  An Ethernet frame from 20:00:00:00:00:01 to
    30:00:00:00:00:02 should arrive on eth0
7  Ethernet frame destined for 30:00:00:00:00:02 should be
    flooded out eth1 and eth2
8  An Ethernet frame should arrive on eth2 with destination
    address the same as eth2's MAC address
9  The hub should not do anything in response to a frame
    arriving with a destination address referring to the hub
    itself.

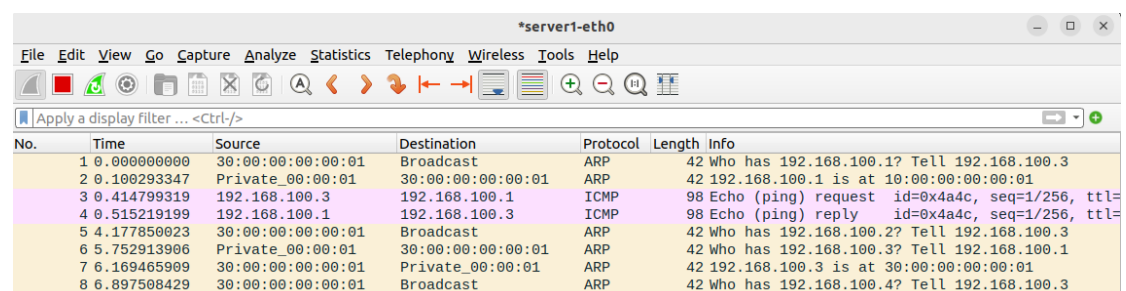
All tests passed!

```

Task 4: Least Recently Used

为方便，此处新增节点 server3，将表项数量上限修改为 3。使用 ping 让 client 分别对三个节点 ping 一次，可以发现通过 ARP 广播获取各节点的 MAC 地址后，每个节点捕获到了发给自己的 echo 包及自己的回应：

server1:



No.	Time	Source	Destination	Protocol	Length	Info
1	0.000000000	30:00:00:00:00:01	Broadcast	ARP	42	Who has 192.168.100.1? Tell 192.168.100.3
2	0.100293347	Private_00:00:01	30:00:00:00:00:01	ARP	42	192.168.100.1 is at 10:00:00:00:00:01
3	0.414799319	192.168.100.3	192.168.100.1	ICMP	98	Echo (ping) request id=0x4a4c, seq=1/256, ttl=
4	0.515219199	192.168.100.1	192.168.100.3	ICMP	98	Echo (ping) reply id=0x4a4c, seq=1/256, ttl=
5	4.177850023	30:00:00:00:00:01	Broadcast	ARP	42	Who has 192.168.100.2? Tell 192.168.100.3
6	5.752913906	Private_00:00:01	30:00:00:00:00:01	ARP	42	Who has 192.168.100.3? Tell 192.168.100.1
7	6.169465909	30:00:00:00:00:01	Private_00:00:01	ARP	42	192.168.100.3 is at 30:00:00:00:00:01
8	6.897508429	30:00:00:00:00:01	Broadcast	ARP	42	Who has 192.168.100.4? Tell 192.168.100.3

server2:

Capturing from server2-eth0						
No.	Time	Source	Destination	Protocol	Length	Info
1	0.000000000	30:00:00:00:00:01	Broadcast	ARP	42	Who has 192.168.100.1? Tell 192.168.100.3
2	4.177849942	30:00:00:00:00:01	Broadcast	ARP	42	Who has 192.168.100.2? Tell 192.168.100.3
3	4.278413066	20:00:00:00:00:01	30:00:00:00:00:01	ARP	42	192.168.100.2 is at 20:00:00:00:00:01
4	4.701289136	192.168.100.3	192.168.100.2	ICMP	98	Echo (ping) request id=0x3717, seq=1/256, ttl=
5	4.801819528	192.168.100.2	192.168.100.3	ICMP	98	Echo (ping) reply id=0x3717, seq=1/256, ttl=
6	6.897507798	30:00:00:00:00:01	Broadcast	ARP	42	Who has 192.168.100.4? Tell 192.168.100.3
7	9.849678044	20:00:00:00:00:01	30:00:00:00:00:01	ARP	42	Who has 192.168.100.3? Tell 192.168.100.2
8	10.237683122	30:00:00:00:00:01	20:00:00:00:00:01	ARP	42	192.168.100.3 is at 30:00:00:00:00:01

server3:

Capturing from server3-eth0						
No.	Time	Source	Destination	Protocol	Length	Info
1	0.000000000	30:00:00:00:00:01	Broadcast	ARP	42	Who has 192.168.100.1? Tell 192.168.100.3
2	4.177849953	30:00:00:00:00:01	Broadcast	ARP	42	Who has 192.168.100.2? Tell 192.168.100.3
3	6.897508510	30:00:00:00:00:01	Broadcast	ARP	42	Who has 192.168.100.4? Tell 192.168.100.3
4	6.997834449	50:00:00:00:00:01	30:00:00:00:00:01	ARP	42	192.168.100.4 is at 50:00:00:00:00:01
5	7.417937499	192.168.100.3	192.168.100.4	ICMP	98	Echo (ping) request id=0xd5bf, seq=1/256, ttl=
6	7.518338064	192.168.100.4	192.168.100.3	ICMP	98	Echo (ping) reply id=0xd5bf, seq=1/256, ttl=
7	12.665386725	50:00:00:00:00:01	30:00:00:00:00:01	ARP	42	Who has 192.168.100.3? Tell 192.168.100.4
8	13.065282779	30:00:00:00:00:01	50:00:00:00:00:01	ARP	42	192.168.100.3 is at 30:00:00:00:00:01

这是与自学习特性相符的。之后我们再次让 client ping 一次 server1
(最早 ping 的节点),可以发现三个节点都收到了相同的 request 包:

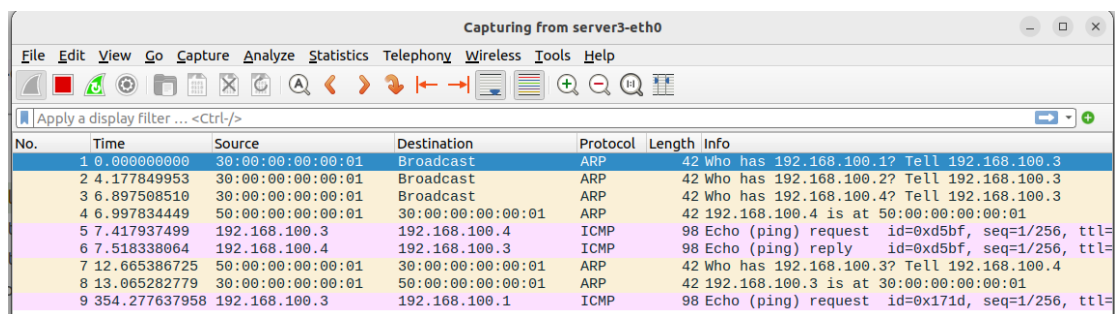
server1:

*server1-eth0						
No.	Time	Source	Destination	Protocol	Length	Info
1	0.000000000	30:00:00:00:00:01	Broadcast	ARP	42	Who has 192.168.100.1? Tell 192.168.100.3
2	0.100293347	Private_00:00:01	30:00:00:00:00:01	ARP	42	192.168.100.1 is at 10:00:00:00:00:01
3	0.414799319	192.168.100.3	192.168.100.1	ICMP	98	Echo (ping) request id=0x4a4c, seq=1/256, ttl=
4	0.515219199	192.168.100.1	192.168.100.3	ICMP	98	Echo (ping) reply id=0x4a4c, seq=1/256, ttl=
5	4.177850023	30:00:00:00:00:01	Broadcast	ARP	42	Who has 192.168.100.2? Tell 192.168.100.3
6	5.752913906	Private_00:00:01	30:00:00:00:00:01	ARP	42	Who has 192.168.100.3? Tell 192.168.100.1
7	6.169465909	30:00:00:00:00:01	Private_00:00:01	ARP	42	192.168.100.3 is at 30:00:00:00:00:01
8	6.897508429	30:00:00:00:00:01	Broadcast	ARP	42	Who has 192.168.100.4? Tell 192.168.100.3
9	354.277637817	192.168.100.3	192.168.100.1	ICMP	98	Echo (ping) request id=0x171d, seq=1/256, ttl=
10	354.378062384	192.168.100.1	192.168.100.3	ICMP	98	Echo (ping) reply id=0x171d, seq=1/256, ttl=
11	359.394480659	30:00:00:00:00:01	Private_00:00:01	ARP	42	Who has 192.168.100.1? Tell 192.168.100.3
12	359.494629647	Private_00:00:01	30:00:00:00:00:01	ARP	42	192.168.100.1 is at 10:00:00:00:00:01
13	359.545657498	Private_00:00:01	30:00:00:00:00:01	ARP	42	Who has 192.168.100.3? Tell 192.168.100.1
14	359.913801840	30:00:00:00:00:01	Private_00:00:01	ARP	42	192.168.100.3 is at 30:00:00:00:00:01

server2:

Capturing from server2-eth0						
No.	Time	Source	Destination	Protocol	Length	Info
1	0.000000000	30:00:00:00:00:01	Broadcast	ARP	42	Who has 192.168.100.1? Tell 192.168.100.3
2	4.177849942	30:00:00:00:00:01	Broadcast	ARP	42	Who has 192.168.100.2? Tell 192.168.100.3
3	4.278413066	20:00:00:00:00:01	30:00:00:00:00:01	ARP	42	192.168.100.2 is at 20:00:00:00:00:01
4	4.701289136	192.168.100.3	192.168.100.2	ICMP	98	Echo (ping) request id=0x3717, seq=1/256, ttl=
5	4.801819528	192.168.100.2	192.168.100.3	ICMP	98	Echo (ping) reply id=0x3717, seq=1/256, ttl=
6	6.897507798	30:00:00:00:00:01	Broadcast	ARP	42	Who has 192.168.100.4? Tell 192.168.100.3
7	9.849678044	20:00:00:00:00:01	30:00:00:00:00:01	ARP	42	Who has 192.168.100.3? Tell 192.168.100.2
8	10.237683122	30:00:00:00:00:01	20:00:00:00:00:01	ARP	42	192.168.100.3 is at 30:00:00:00:00:01
9	354.277637386	192.168.100.3	192.168.100.1	ICMP	98	Echo (ping) request id=0x171d, seq=1/256, ttl=

server3:



No.	Time	Source	Destination	Protocol	Length	Info
1	0.000000000	30:00:00:00:00:01	Broadcast	ARP	42	Who has 192.168.100.1? Tell 192.168.100.3
2	4.177849953	30:00:00:00:00:01	Broadcast	ARP	42	Who has 192.168.100.2? Tell 192.168.100.3
3	6.897508510	30:00:00:00:00:01	Broadcast	ARP	42	Who has 192.168.100.4? Tell 192.168.100.3
4	6.997834449	50:00:00:00:00:01	30:00:00:00:00:01	ARP	42	192.168.100.4 is at 50:00:00:00:00:01
5	7.417937499	192.168.100.3	192.168.100.4	ICMP	98	Echo (ping) request id=0xd5bf, seq=1/256, ttl=
6	7.518338064	192.168.100.4	192.168.100.3	ICMP	98	Echo (ping) reply id=0xd5bf, seq=1/256, ttl=
7	12.665386725	50:00:00:00:00:01	30:00:00:00:00:01	ARP	42	Who has 192.168.100.3? Tell 192.168.100.4
8	13.065282779	30:00:00:00:00:01	50:00:00:00:00:01	ARP	42	192.168.100.3 is at 30:00:00:00:00:01
9	354.277637958	192.168.100.3	192.168.100.1	ICMP	98	Echo (ping) request id=0x171d, seq=1/256, ttl=

这说明表项数达到上限后，在第三次 ping 时，交换机学习到 server3 的同时选择遗忘了最近最不常用的 server1，这与我们的预期相符。

测试结果如下：

```
30:00:00:00:00:02 should arrive on eth2
6 Ethernet frame destined for 30:00:00:00:00:02 should arrive
  on eth1 after self-learning
7 An Ethernet frame from 30:00:00:00:00:04 to
  20:00:00:00:00:01 should arrive on eth3
8 Ethernet frame destined to 20:00:00:00:00:01 should arrive
  on eth0 after self-learning
9 An Ethernet frame from 20:00:00:00:00:01 to
  30:00:00:00:00:04 should arrive on eth0
10 Ethernet frame destined to 20:00:00:00:00:01 should arrive
  on eth3 after self-learning
11 An Ethernet frame from 40:00:00:00:00:05 to
  20:00:00:00:00:01 should arrive on eth4
12 Ethernet frame destined to 20:00:00:00:00:01 should arrive
  on eth0 after self-learning
13 An Ethernet frame from 30:00:00:00:00:05 to
  20:00:00:00:00:01 should arrive on eth4
14 Ethernet frame destined to 20:00:00:00:00:01 should arrive
  on eth0 after self-learning
15 An Ethernet frame from 20:00:00:00:00:05 to
  30:00:00:00:00:02 should arrive on eth4
16 Ethernet frame destined to 30:00:00:00:00:02 should be
  flooded to eth0, eth1, eth2 and eth3
17 An Ethernet frame should arrive on eth2 with destination
  address the same as eth2's MAC address
18 The hub should not do anything in response to a frame
  arriving with a destination address referring to the hub
  itself.
```

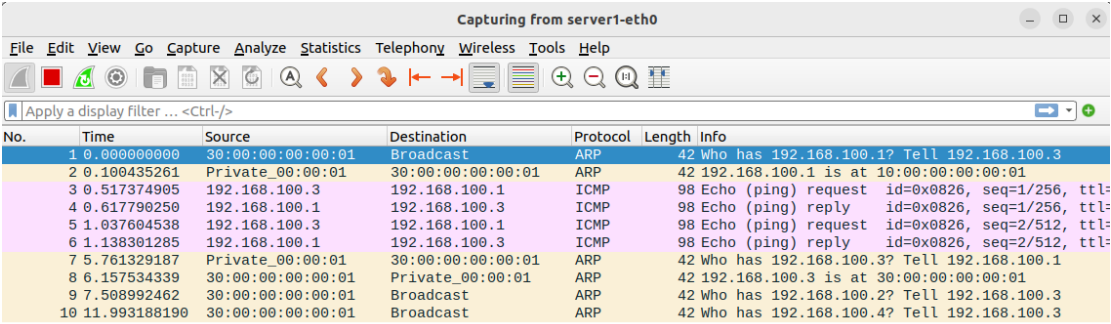
All tests passed!

Task 5: Least Traffic Volume

为方便，此处新增节点 server3，将表项数量上限修改为 3。我们首

先让 client 向 server1 ping 两次，向 server2 ping 一次，向 server3 ping 一次，捕获到的包如下：

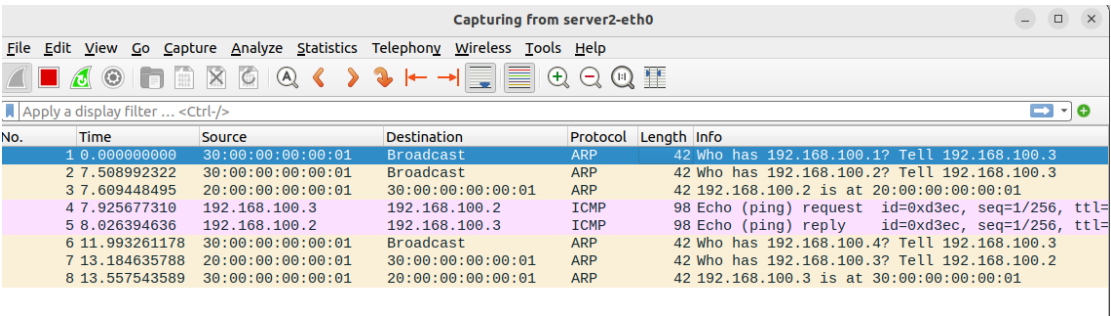
server1:



Capturing from server1-eth0

No.	Time	Source	Destination	Protocol	Length	Info
1	0.000000000	30:00:00:00:00:01	Broadcast	ARP	42	Who has 192.168.100.1? Tell 192.168.100.3
2	0.100435261	Private_00:00:01	30:00:00:00:00:01	ARP	42	192.168.100.1 is at 10:00:00:00:00:01
3	0.517374905	192.168.100.3	192.168.100.1	ICMP	98	Echo (ping) request id=0x0826, seq=1/256, ttl=
4	0.617790250	192.168.100.1	192.168.100.3	ICMP	98	Echo (ping) reply id=0x0826, seq=1/256, ttl=
5	1.037604538	192.168.100.3	192.168.100.1	ICMP	98	Echo (ping) request id=0x0826, seq=2/512, ttl=
6	1.138301285	192.168.100.1	192.168.100.3	ICMP	98	Echo (ping) reply id=0x0826, seq=2/512, ttl=
7	5.761329187	Private_00:00:01	30:00:00:00:00:01	ARP	42	Who has 192.168.100.3? Tell 192.168.100.1
8	6.157534339	30:00:00:00:00:01	Private_00:00:01	ARP	42	192.168.100.3 is at 30:00:00:00:00:01
9	7.508992462	30:00:00:00:00:01	Broadcast	ARP	42	Who has 192.168.100.2? Tell 192.168.100.3
10	11.993188190	30:00:00:00:00:01	Broadcast	ARP	42	Who has 192.168.100.4? Tell 192.168.100.3

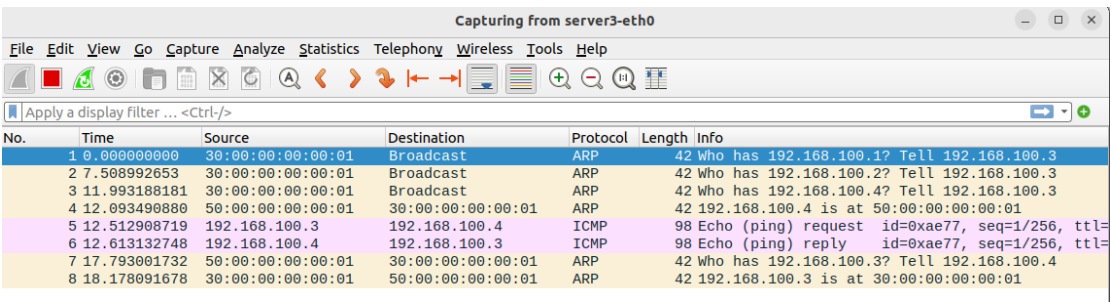
server2:



Capturing from server2-eth0

No.	Time	Source	Destination	Protocol	Length	Info
1	0.000000000	30:00:00:00:00:01	Broadcast	ARP	42	Who has 192.168.100.1? Tell 192.168.100.3
2	7.508992322	30:00:00:00:00:01	Broadcast	ARP	42	Who has 192.168.100.2? Tell 192.168.100.3
3	7.609448495	20:00:00:00:00:01	30:00:00:00:00:01	ARP	42	192.168.100.2 is at 20:00:00:00:00:01
4	7.925677310	192.168.100.3	192.168.100.2	ICMP	98	Echo (ping) request id=0xd3ec, seq=1/256, ttl=
5	8.026394636	192.168.100.2	192.168.100.3	ICMP	98	Echo (ping) reply id=0xd3ec, seq=1/256, ttl=
6	11.993261178	30:00:00:00:00:01	Broadcast	ARP	42	Who has 192.168.100.4? Tell 192.168.100.3
7	13.184635788	20:00:00:00:00:01	30:00:00:00:00:01	ARP	42	Who has 192.168.100.3? Tell 192.168.100.2
8	13.557543589	30:00:00:00:00:01	20:00:00:00:00:01	ARP	42	192.168.100.3 is at 30:00:00:00:00:01

server3:



Capturing from server3-eth0

No.	Time	Source	Destination	Protocol	Length	Info
1	0.000000000	30:00:00:00:00:01	Broadcast	ARP	42	Who has 192.168.100.1? Tell 192.168.100.3
2	7.508992653	30:00:00:00:00:01	Broadcast	ARP	42	Who has 192.168.100.2? Tell 192.168.100.3
3	11.993188181	30:00:00:00:00:01	Broadcast	ARP	42	Who has 192.168.100.4? Tell 192.168.100.3
4	12.093490880	50:00:00:00:00:01	30:00:00:00:00:01	ARP	42	192.168.100.4 is at 50:00:00:00:00:01
5	12.512908719	192.168.100.3	192.168.100.4	ICMP	98	Echo (ping) request id=0xae77, seq=1/256, ttl=
6	12.613132748	192.168.100.4	192.168.100.3	ICMP	98	Echo (ping) reply id=0xae77, seq=1/256, ttl=
7	17.793001732	50:00:00:00:00:01	30:00:00:00:00:01	ARP	42	Who has 192.168.100.3? Tell 192.168.100.4
8	18.178091678	30:00:00:00:00:01	50:00:00:00:00:01	ARP	42	192.168.100.3 is at 30:00:00:00:00:01

可以发现除了广播 ARP 包外，都只收到了发送给自己的 echo 包，符合预期。在让 client ping 一次 server2，三个节点收到的包如下：

server1:

Capturing from server1-eth0						
No.	Time	Source	Destination	Protocol	Length	Info
1	0.000000000	30:00:00:00:00:01	Broadcast	ARP	42	Who has 192.168.100.1? Tell 192.168.100.3
2	0.100435261	Private_00:00:01	30:00:00:00:00:01	ARP	42	192.168.100.1 is at 10:00:00:00:00:01
3	0.517374905	192.168.100.3	192.168.100.1	ICMP	98	Echo (ping) request id=0x0826, seq=1/256, ttl=
4	0.617790250	192.168.100.1	192.168.100.3	ICMP	98	Echo (ping) reply id=0x0826, seq=1/256, ttl=
5	1.037604538	192.168.100.3	192.168.100.1	ICMP	98	Echo (ping) request id=0x0826, seq=2/512, ttl=
6	1.138301285	192.168.100.1	192.168.100.3	ICMP	98	Echo (ping) reply id=0x0826, seq=2/512, ttl=
7	5.761329187	Private_00:00:01	30:00:00:00:00:01	ARP	42	Who has 192.168.100.3? Tell 192.168.100.1
8	6.157534339	30:00:00:00:00:01	Private_00:00:01	ARP	42	192.168.100.3 is at 30:00:00:00:00:01
9	7.508992462	30:00:00:00:00:01	Broadcast	ARP	42	Who has 192.168.100.2? Tell 192.168.100.3
10	11.993188190	30:00:00:00:00:01	Broadcast	ARP	42	Who has 192.168.100.4? Tell 192.168.100.3
11	218.281485593	192.168.100.3	192.168.100.2	ICMP	98	Echo (ping) request id=0xff87, seq=1/256, ttl=

server2:

Capturing from server2-eth0						
No.	Time	Source	Destination	Protocol	Length	Info
1	0.000000000	30:00:00:00:00:01	Broadcast	ARP	42	Who has 192.168.100.1? Tell 192.168.100.3
2	7.508992322	30:00:00:00:00:01	Broadcast	ARP	42	Who has 192.168.100.2? Tell 192.168.100.3
3	7.609448495	20:00:00:00:00:01	30:00:00:00:00:01	ARP	42	192.168.100.2 is at 20:00:00:00:00:01
4	7.925677310	192.168.100.3	192.168.100.2	ICMP	98	Echo (ping) request id=0xd3ec, seq=1/256, ttl=
5	8.026394636	192.168.100.2	192.168.100.3	ICMP	98	Echo (ping) reply id=0xd3ec, seq=1/256, ttl=
6	11.993261178	30:00:00:00:00:01	Broadcast	ARP	42	Who has 192.168.100.4? Tell 192.168.100.3
7	13.184635788	20:00:00:00:00:01	30:00:00:00:00:01	ARP	42	Who has 192.168.100.3? Tell 192.168.100.2
8	13.557543589	30:00:00:00:00:01	20:00:00:00:00:01	ARP	42	192.168.100.3 is at 30:00:00:00:00:01
9	218.281485223	192.168.100.3	192.168.100.2	ICMP	98	Echo (ping) request id=0xff87, seq=1/256, ttl=
10	218.381710344	192.168.100.2	192.168.100.3	ICMP	98	Echo (ping) reply id=0xff87, seq=1/256, ttl=
11	223.509311312	30:00:00:00:00:01	20:00:00:00:00:01	ARP	42	Who has 192.168.100.2? Tell 192.168.100.3
12	223.609598462	20:00:00:00:00:01	30:00:00:00:00:01	ARP	42	192.168.100.2 is at 20:00:00:00:00:01
13	223.616886417	20:00:00:00:00:01	30:00:00:00:00:01	ARP	42	Who has 192.168.100.3? Tell 192.168.100.2
14	224.029537969	30:00:00:00:00:01	20:00:00:00:00:01	ARP	42	192.168.100.3 is at 30:00:00:00:00:01

server3:

Capturing from server3-eth0						
No.	Time	Source	Destination	Protocol	Length	Info
1	0.000000000	30:00:00:00:00:01	Broadcast	ARP	42	Who has 192.168.100.1? Tell 192.168.100.3
2	7.508992653	30:00:00:00:00:01	Broadcast	ARP	42	Who has 192.168.100.2? Tell 192.168.100.3
3	11.993188181	30:00:00:00:00:01	Broadcast	ARP	42	Who has 192.168.100.4? Tell 192.168.100.3
4	12.093490880	50:00:00:00:00:01	30:00:00:00:00:01	ARP	42	192.168.100.4 is at 50:00:00:00:00:01
5	12.512908719	192.168.100.3	192.168.100.4	ICMP	98	Echo (ping) request id=0xae77, seq=1/256, ttl=
6	12.613132748	192.168.100.4	192.168.100.3	ICMP	98	Echo (ping) reply id=0xae77, seq=1/256, ttl=
7	17.793001732	50:00:00:00:00:01	30:00:00:00:00:01	ARP	42	Who has 192.168.100.3? Tell 192.168.100.4
8	18.178091678	30:00:00:00:00:01	50:00:00:00:00:01	ARP	42	192.168.100.3 is at 30:00:00:00:00:01
9	218.281486195	192.168.100.3	192.168.100.2	ICMP	98	Echo (ping) request id=0xff87, seq=1/256, ttl=

可以发现三个节点都收到发送给 server2 的包，这说明在表项数量达到上限后，交换机选择遗忘了流量最小的表项，即 server2，所以在向 server2 发送时会广播该包。

测试结果如下：

```
12 Ethernet frame destined to 20:00:00:00:00:01 should arrive
    on eth0 after self-learning
13 An Ethernet frame from 30:00:00:00:00:04 to
    20:00:00:00:00:01 should arrive on eth3
14 Ethernet frame destined to 20:00:00:00:00:01 should arrive
    on eth0 after self-learning
15 An Ethernet frame from 40:00:00:00:00:05 to
    20:00:00:00:00:01 should arrive on eth4
16 Ethernet frame destined to 20:00:00:00:00:01 should arrive
    on eth0 after self-learning
17 An Ethernet frame from 20:00:00:00:00:05 to
    30:00:00:00:00:02 should arrive on eth4
18 Ethernet frame destined to 30:00:00:00:00:02 should arrive
    on eth1 after self-learning
19 An Ethernet frame from 30:00:00:00:00:02 to
    20:00:00:00:00:05 should arrive on eth1
20 Ethernet frame destined to 20:00:00:00:00:05 should arrive
    on eth4 after self-learning
21 An Ethernet frame from 20:00:00:00:00:03 to
    40:00:00:00:00:05 should arrive on eth2
22 Ethernet frame destined to 40:00:00:00:00:05 should be
    flooded to eth0, eth1, eth3 and eth4
23 An Ethernet frame should arrive on eth2 with destination
    address the same as eth2's MAC address
24 The hub should not do anything in response to a frame
    arriving with a destination address referring to the hub
    itself.
```

All tests passed!

5. 核心代码

Task 2: Basic Switch

采用字典来存储表项，key 为 MAC 地址，value 为对应端口。

```
eth = packet.get_header(Ethernet)
if eth is None:
    log_info("Received a non-Ethernet packet?!")
    return
src = eth.src
dst = eth.dst
table[src] = fromIface
if eth.dst in mymacs:
    log_info("Received a packet intended for me")
else:
    if table.get(dst) == None or dst == 'ff:ff:ff:ff:ff:ff':
        for intf in my_interfaces:
            if fromIface != intf.name:
```

```

        log_info (f"Flooding packet {packet} to {intf.name}")
        net.send_packet(intf, packet)
    else:
        log_info (f"Forwarding packet {packet} to {table[dst]}")
        net.send_packet(table[dst], packet)

```

Task 3: Timeouts

采用字典来存储表项，key 为 MAC 地址，value 为由对应端口和产生时间组成的元组。

```

eth = packet.get_header(Ethernet)
if eth is None:
    log_info("Received a non-Ethernet packet?!")
    return
src = eth.src
dst = eth.dst
t = time.time()
table[src] = (fromIface, t)
if eth.dst in mymacs:
    log_info("Received a packet intended for me")
else:
    if dst == 'ff:ff:ff:ff:ff:ff' or table.get(dst) == None or t -
table[dst][1] > T:
        if table.get(dst) != None:
            table.pop(dst)
        for intf in my_interfaces:
            if fromIface != intf.name:
                log_info (f"Flooding packet {packet} to {intf.name}")
                net.send_packet(intf, packet)
    else:
        log_info (f"Forwarding packet {packet} to {table[dst][0]}")
        net.send_packet(table[dst][0], packet)

```

Task 4: Least Recently Used

采用字典来存储表项，key 为 MAC 地址，value 为由对应端口和权值组成的列表组成。每次收到包，若该 src 的 MAC 地址存在，则将端口进行更新，否则新加该表项，权值设为 1；之后将所有表项权值加 1，若表项数量超过 N，删除权值最大的表项；若 dst 的 MAC 地址存在，将其权值设为 1。

```

eth = packet.get_header(Ethernet)
if eth is None:
    log_info("Received a non-Ethernet packet?!")
    return
src = eth.src
dst = eth.dst
if table.get(src) != None:
    table[src][0] = fromIface
else:
    table[src] = [fromIface, 1]
for k in table.keys():
    table[k][1] += 1
if len(table) > N:
    ks, vs = list(table.keys()), list(table.values())
    k = ks[vs.index(max(vs, key=lambda v : v[1])]]
    table.pop(k)
if eth.dst in mymacs:
    log_info("Received a packet intended for me")
else:
    if dst == 'ff:ff:ff:ff:ff:ff' or table.get(dst) == None:
        for intf in my_interfaces:
            if fromIface != intf.name:
                log_info(f"Flooding packet {packet} to {intf.name}")
                net.send_packet(intf, packet)
    else:
        log_info(f'Forwarding packet {packet} to {table[dst][0]}')
        net.send_packet(table[dst][0], packet)
        table[dst][1] = 1

```

Task 5: Least Traffic Volume

采用字典来存储表项，key 为 MAC 地址，value 为由对应端口和流量组成的列表组成。初始流量设为 0，每次 dst 的 MAC 地址若存在，则将其对应流量加 1，表满后删除流量最小的表项。

```

eth = packet.get_header(Ethernet)
if eth is None:
    log_info("Received a non-Ethernet packet?!")
    return
src = eth.src
dst = eth.dst
if table.get(src) == None:
    table[src] = [fromIface, 0]

```

```

else:
    table[src][0] = fromIface
if len(table) > N:
    k = None
    v = 9999
    for key, val in table.items():
        if v > val[1] and key != src:
            k = key
            v = val[1]
    table.pop(k)
if eth.dst in mymacs:
    log_info("Received a packet intended for me")
else:
    if dst == 'ff:ff:ff:ff:ff:ff' or table.get(dst) == None:
        for intf in my_interfaces:
            if fromIface != intf.name:
                log_info(f"Flooding packet {packet} to {intf.name}")
                net.send_packet(intf, packet)
    else:
        log_info(f'Forwarding packet {packet} to {table[dst][0]}')
        net.send_packet(table[dst][0], packet)
        table[dst][1] += 1

```

6. 总结与感想

本次实验对我来说有一定的挑战性，我也花了一些时间才勉强完成。总的来说，收获很大，不仅对 learning switch 有了更深入的理解，也对 python 有了更多的了解。