

Premier University



Assignment

Assignment Title	Applying Dijkstra and Knapsack Algorithm in Real Life Scenario
Course code	CSE 225
Course name	Algorithm Design and Analysis
Date of Submission	13/02/2024

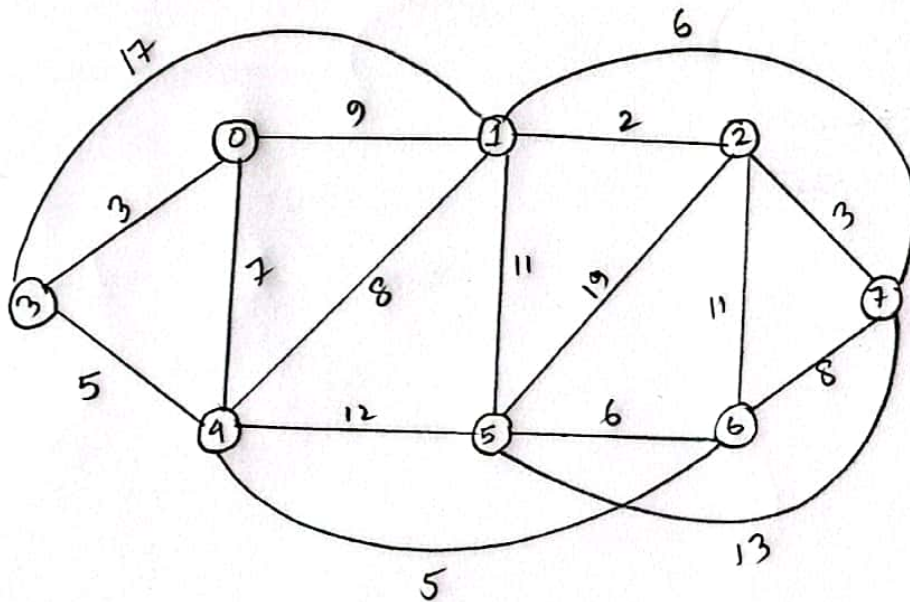
Submitted to
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Semester	4th
Session	Fall 2023

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Objective:

To design a intelligent transportation Optimization System for a smart city using dijkstra and 0/1 Knapsack algorithm. Finding best route for time efficiency to visit from a single city to every city for the inhabitants and maximize sight seeing value for the tourists in a limited budget.

Design best route from single city:

Here, In This graph (0-7) \rightarrow 8 nodes/vertex representing 8 city and each edge representing 17 different Paths with timing cost to visit from one city to another city. In This graph I will apply, Single Source Shortest Path algorithm (Dijkstra) to find the best route. As, there is no negative cost. It's an Undirected graph also.

Algorithm choice for Shortest Path:

I will choose dijkstra algorithm for finding minimum cost for the inhabitants. It's an fundamental Shortest Path from a single Source Vertex to all other vertices in a weighted graph. It employs a greedy strategy, iteratively selecting the vertex and updating the distances to its neighbors accordingly.

Graph representation:

I've chosen adjacency list for the graph representation, because it's a memory efficient approach, rather than adjacency matrix. An adjacency list is a data structure used to represent a graph where each vertex maintains a list of its neighboring vertices.

Data Structure:

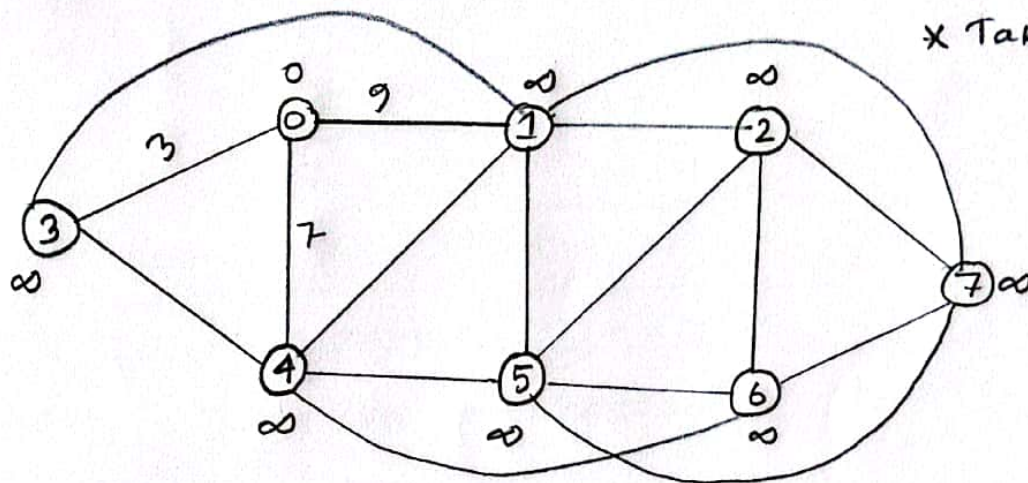
I will use C++ STL vectors and Pairs to store the graph, to keep track of visited vertices, and the distances.

And, I will use C++ STL set to find the minimum cost containing nodes/vertices. I've chosen set, because, it's by default can provide me the minimum weight in $(\log n)$ time. I might choose priority-queue, but declaration is a little bit complex (for small \rightarrow greater) values. Overall worst case complexity for dijkstra is $O(n^2)$.

2 2
Step 01:

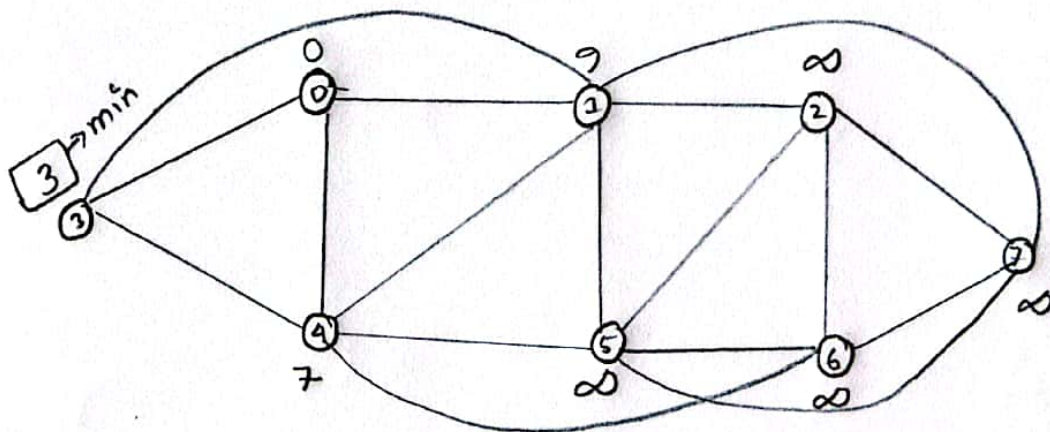
* Here, for the drawing Complexity, I will draw only the edges of current vertex in each step

Source node = 0



Step 02:

Current Vertex = 0

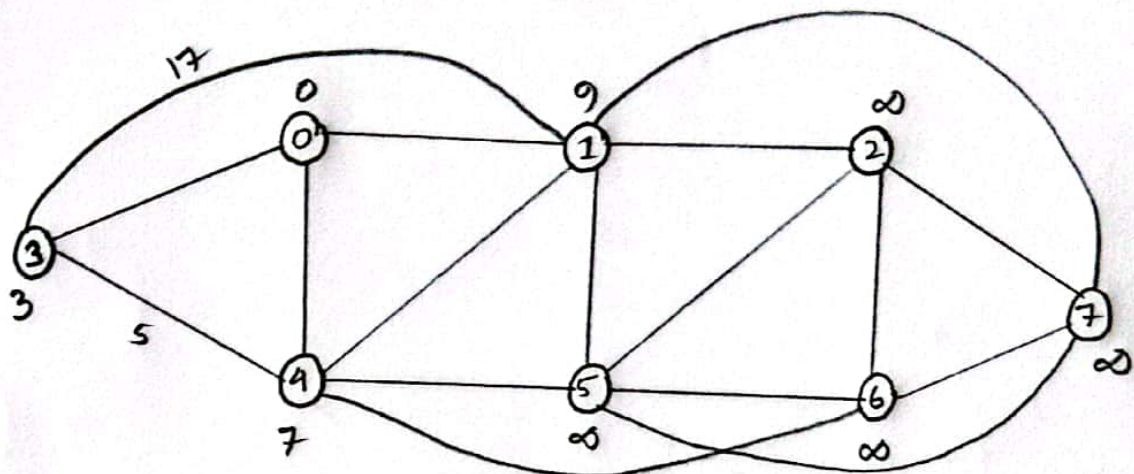


In this step, taking minimum distance from Source Vertex "0", we visit to the new vertex "3", because 3 is the minimum distance here.

3

Step 03:

Current Vertex = 3



as, $(0 \rightarrow 3 \rightarrow 7) \rightarrow$ Total cost is 8, we apply this-

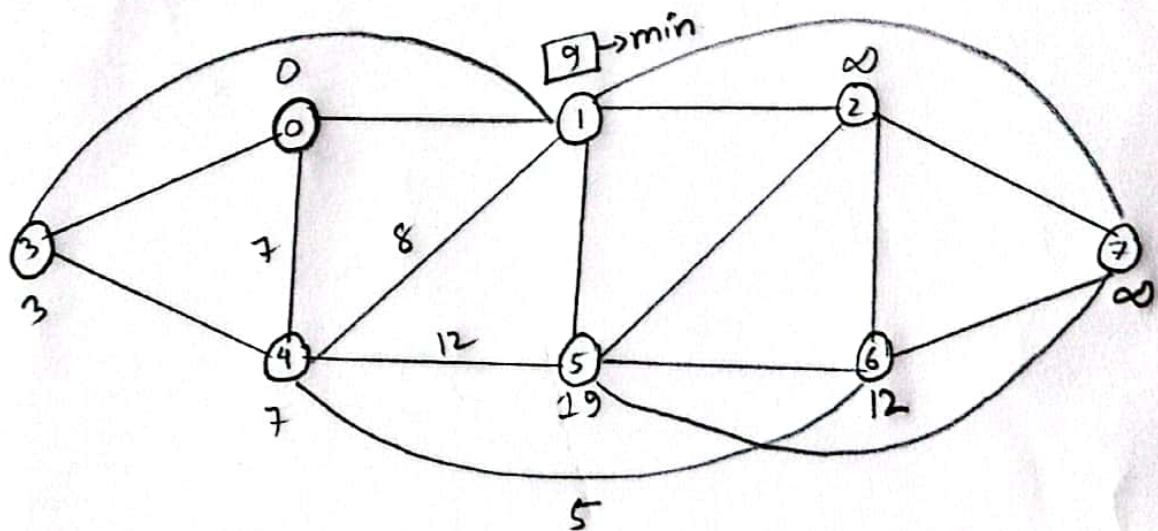
$$d(A,B) = \min(d(A) + d(A,B), d(B)) \rightarrow \text{relaxation formula}$$

-from 0

now, we travel to vertex to 4, because, it's minimal way.

Step 04:

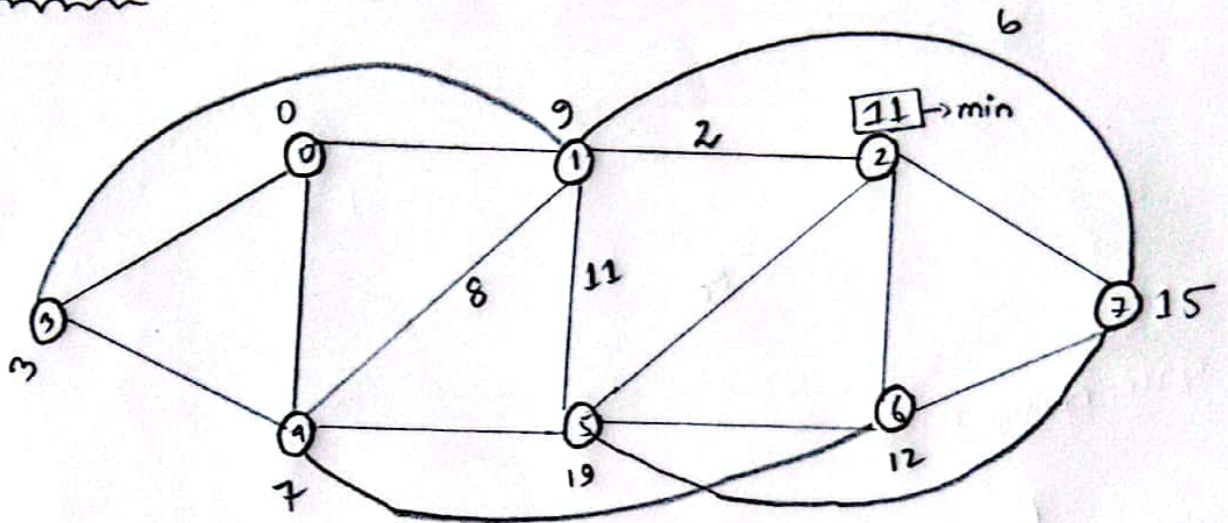
Current Vertex = 4



Here, Vertex 1 has minimum cost. So, we traverse to Vertex 1, for next relaxation.

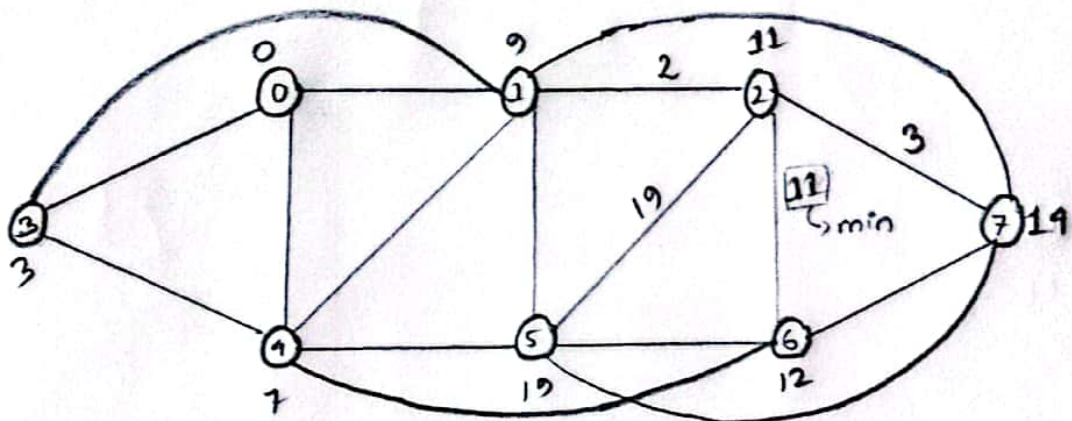
Step 05:

Current vertex = 1



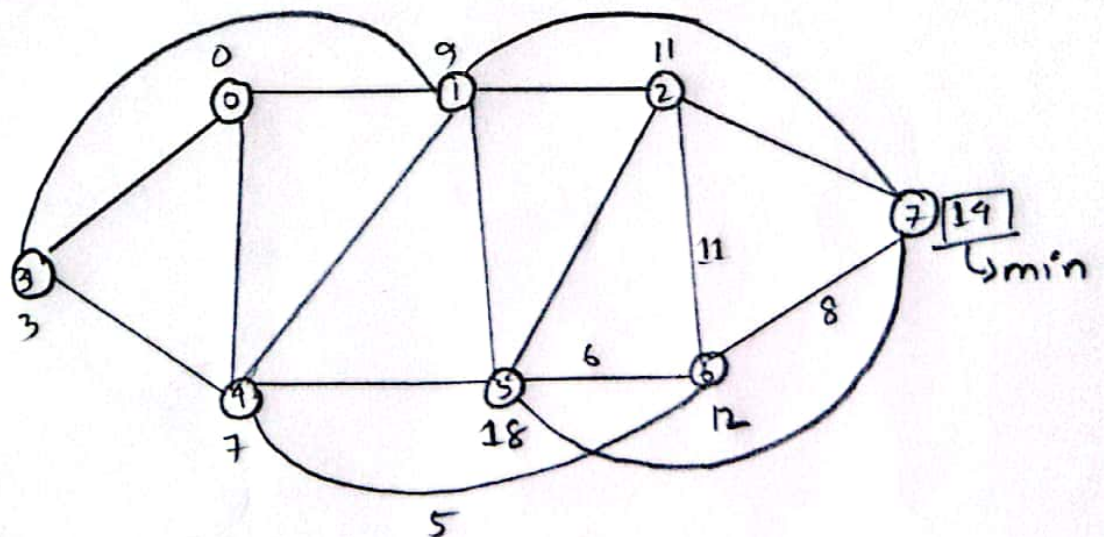
Step 06:

Current Vertex = 2



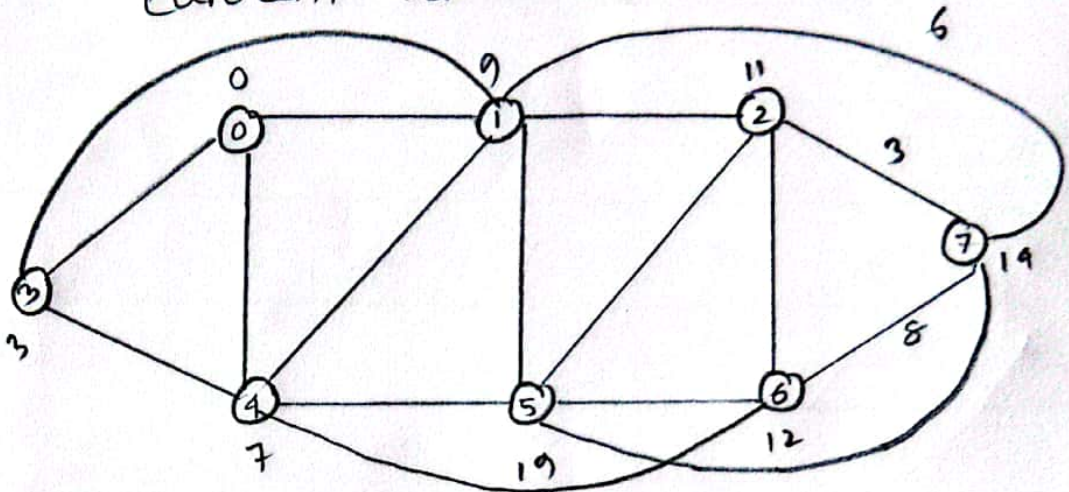
Step 07:

Current vertex = 6.



Step 08:

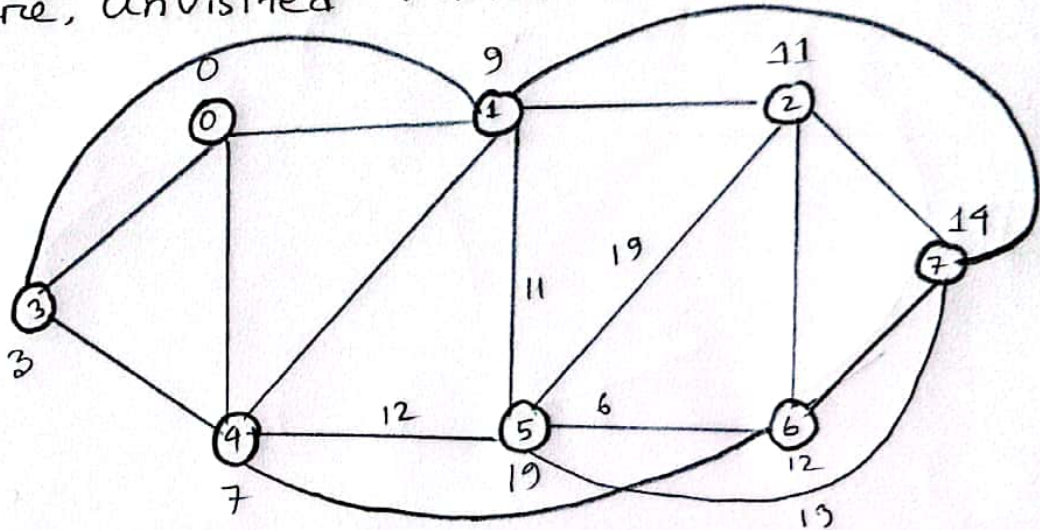
Current vertex = 7



Step 09:

Current vertex = 5.

as there, unvisited vertex is only 5 now.



Final distances: from city 0 to other city —

City 0 : 0
 city 01 : 9
 city 2 : 11
 city 3 : 3
 city 4 : 7
 city 5 : 18
 city 6 : 12
 city 7 : 14

Computation Table:

	0	1	2	3	4	5	6	7
	0	∞	∞	∞	∞	∞	∞	∞
0	<u>0</u>	9	∞	<u>3</u>	7	∞	∞	∞
3	<u>0</u>	9	∞	<u>3</u>	<u>7</u>	∞	∞	∞
4	<u>0</u>	<u>9</u>	∞	<u>3</u>	<u>7</u>	19	12	∞
1	<u>0</u>	<u>9</u>	<u>11</u>	<u>3</u>	<u>7</u>	19	12	15
2	<u>0</u>	<u>9</u>	<u>11</u>	<u>3</u>	<u>7</u>	19	<u>12</u>	14
6	<u>0</u>	<u>9</u>	<u>11</u>	<u>3</u>	<u>7</u>	18	<u>12</u>	<u>19</u>
7	<u>0</u>	<u>9</u>	<u>11</u>	<u>3</u>	<u>7</u>	<u>18</u>	<u>12</u>	<u>19</u>
5	<u>0</u>	<u>9</u>	<u>11</u>	<u>3</u>	<u>7</u>	<u>18</u>	<u>12</u>	<u>19</u>

In the computation table -

Marked with box (\square) are the current minimum cost for the vertex/city and underlined ($\underline{\quad}$)

costs are representing already taken/visited city/vertex, which should not be visit again. to avoid unnecessary computation.

Code for the Inhabitants :

```
1 #include <bits/stdc++.h>
2 using namespace std;
3 const int inf = 1e7;
4 const int n = 1e3;
5 vector<pair<int, int>> g[n];
6 // { node , cost}
7 void dijkstra(int source)
8 {
9     vector<int> distance(n, inf);
10    vector<bool> visited(n, 0);
11    set<pair<int, int>> st;
12    // {cost , node} => kept cost on first value to sort based on lowest cost
13    st.insert({0, source});
14    distance[source] = 0;
15    while (!st.empty())
16    {
17        auto node = *st.begin();
18        // will give the minimum weighted pair {cost , node}
19        int parent_node = node.second;
20        int parent_node_cost = node.first;
21        st.erase(st.begin());
22        if (visited[parent_node])
23        {
24            continue;
25        }
26        visited[parent_node] = 1;
27        // Traverse to the child of v, for Relaxation
28        for (auto child : g[parent_node])
29        {
30            int child_node = child.first;
31            int edge_cost = child.second;
32
33            // Relaxation
34            if ((parent_node_cost + edge_cost) < distance[child_node])
35            {
36                distance[child_node] = (parent_node_cost + edge_cost);
37                st.insert({distance[child_node], child_node});
38            }
39        }
40    }
41    cout << "Node\tDistance from " << source << endl;
42    for (int i = 0; i < n; ++i)
43    {
44        if (distance[i] != inf)
45        {
46            cout << i << "\t" << distance[i] << endl;
47        }
48    }
49 }
50 int main()
51 {
52     int node, edge;
53     cin >> node >> edge;
54     for (int i = 0; i < edge; i++)
55     {
56         int u, v, cost;
57         cin >> u >> v >> cost;
58         g[u].push_back({v, cost});
59         g[v].push_back({u, cost});
60         // u/v indexed node connected with v/u node with cost
61     }
62     dijkstra(0);
63
64     return 0;
65 }
66
```

Minimum costs output for 0 number city's Inhabitants to travel into other cities using C++ :

PROBLEMS OUTPUT DEBUG CONSOLE TERMINAL PORTS GITLENS

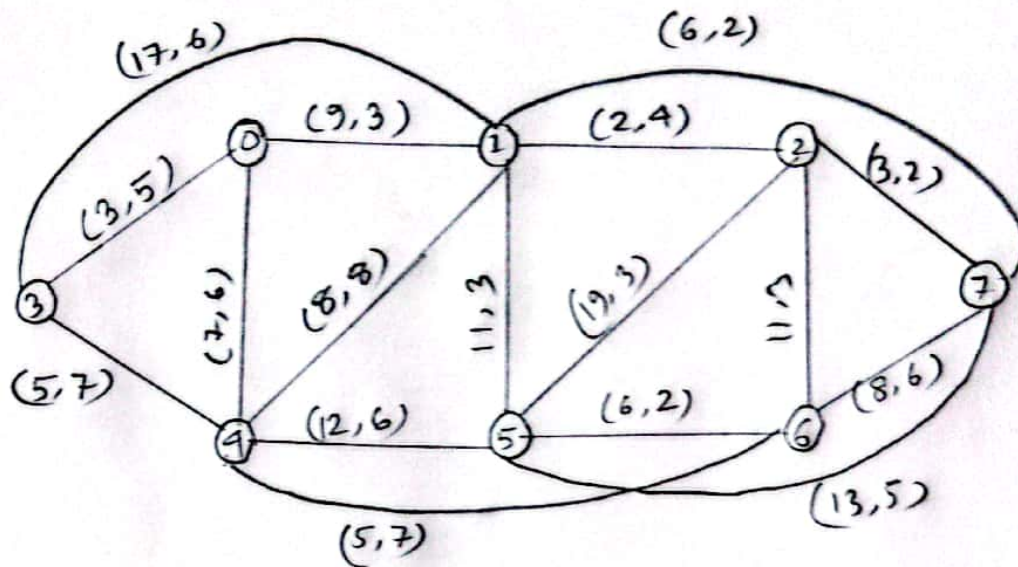
```
8 17
0 1 9
0 4 7
0 3 3
1 4 8
1 5 11
1 7 6
1 3 17
1 2 2
2 7 3
2 6 11
2 5 19
3 4 5
4 6 5
4 5 12
5 7 13
5 6 6
6 7 8
```

City	Distance from 0
0	0
1	9
2	11
3	3
4	7
5	18
6	12
7	14

 SAKIB  AA Test

(8)

Maximize Site Seeing for tourists:



Here, (cost, ss-value) Pair taken.

Algorithm choice:

I choose 0/1 Knapsack algorithm for maximizing Site-Seeing values. This algorithm involves selecting a subset of items with maximum value while respecting a weight constraint. Each item can either be included in the Knapsack (1) or excluded (0), hence the name. It's a classic optimization problem often solved using dynamic programming techniques to efficiently explore all possible combinations of items and weights.

Data Structures:

I will use C++ STL (vectors) to keep the values of (weight and cost) and a 2D Vector for store the results of subproblems in the dynamic programming solution to the 0/1 Knapsack Problem.

☐ Site-Seeing values and costs from graph:

Budget of Tourist: 15

Nodes	SS-Value	cost
0 → 1	3	9
0 → 4	6	7
0 → 3	5	3
1 → 4	8	8
1 → 5	3	11
1 → 7	2	6
1 → 3	6	17
1 → 2	4	2
2 → 7	2	3
2 → 6	3	11
2 → 5	3	19
3 → 4	7	5
4 → 6	7	5
4 → 5	6	12
5 → 7	5	13
5 → 6	2	6
6 → 7	6	8

Resultant Table:

	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
1	0	0	0	0	0	0	0	0	0	3	3	3	3	3	3	3
2	0	0	0	0	0	0	0	6	6	6	6	6	6	6	6	6
3	0	0	0	5	5	5	5	6	6	6	11	11	11	11	11	11
4	0	0	0	5	5	5	5	6	8	8	11	13	13	13	13	14
5	0	0	0	5	5	5	5	6	8	8	11	13	13	13	13	14
6	0	0	0	5	5	5	5	6	8	8	11	13	13	13	13	14
7	0	0	0	5	5	5	5	6	8	8	11	13	13	13	13	14
8	0	0	4	5	5	5	9	9	9	10	12	13	15	17	17	17
9	0	0	4	5	5	5	9	9	11	11	12	13	15	17	17	17
10	0	0	4	5	5	5	9	9	11	11	12	13	15	17	17	17
11	0	0	4	5	5	5	9	9	11	11	12	13	15	17	17	17
12	0	0	4	5	5	5	9	9	12	12	16	16	16	18	18	19
13	0	0	4	5	5	5	9	9	12	12	16	16	18	19	19	23
14	0	0	4	5	5	5	9	9	12	12	16	16	19	19	19	23
15	0	0	4	5	5	5	9	9	12	12	16	16	19	19	19	23
16	0	0	4	5	5	5	9	9	12	12	16	16	19	19	19	23
17	0	0	4	5	5	5	9	9	12	12	16	16	19	19	19	23

Code for maximizing site-seeing values for tourists:

```
1 #include <bits/stdc++.h>
2 using namespace std;
3
4 vector<int> ss_value = {3, 6, 5, 8, 3, 2, 6, 4, 2, 3, 3, 7, 7, 6, 5, 2, 6};
5 vector<int> cost = {9, 7, 3, 8, 11, 6, 17, 2, 3, 11, 19, 5, 5, 12, 13, 6, 8};
6 int budget = 15; // Tourists Budget
7
8 int knapsack(int W)
9 {
10     int n = cost.size();
11     vector<vector<int>> dp(n + 1, vector<int>(W + 1, 0));
12     for (int i = 1; i <= n; i++)
13     {
14         for (int w = 1; w <= W; w++)
15         {
16             if (cost[i - 1] <= w)
17             {
18                 dp[i][w] = max(dp[i - 1][w], dp[i - 1][w - cost[i - 1]] + ss_value[i - 1]);
19             }
20             else
21             {
22                 dp[i][w] = dp[i - 1][w];
23             }
24         }
25     }
26
27     // Printing the resultant table
28     cout << "Resultant Table (0/1 Knapsack Table):" << endl;
29     cout << "-----" << endl;
30     cout << setw(6) << " ";
31     for (int w = 0; w <= W; w++)
32     {
33         cout << setw(6) << w;
34     }
35     cout << endl;
36     cout << "-----" << endl;
37     for (int i = 0; i <= n; i++)
38     {
39         cout << setw(4) << i << " |";
40         for (int w = 0; w <= W; w++)
41         {
42             cout << setw(6) << dp[i][w];
43         }
44         cout << endl;
45     }
46     cout << "-----" << endl;
47
48     return dp[n][W];
49 }
50
51 int main()
52 {
53     cout << "Maximum Site-Seeing ss_value : " << knapsack(budget) << endl;
54
55     return 0;
56 }
57
```


Resultant table and output for the tourists :

```
PROBLEMS  OUTPUT  DEBUG CONSOLE  TERMINAL  PORTS  GITLENS
• SAKIB  AA Test
  cd "c:\PU Projects\OJ Problem Solve\Codeforces Random Problem Solve\AA Test\" ; if ($?) { g++ knapsack.cpp -o knapsack } ; if (
  $?) { .\knapsack }
Resultant Table (0/1 Knapsack Table):
-----
      0   1   2   3   4   5   6   7   8   9  10  11  12  13  14  15
-----
0 |  0   0   0   0   0   0   0   0   0   0   0   0   0   0   0   0
1 |  0   0   0   0   0   0   0   0   0   3   3   3   3   3   3   3
2 |  0   0   0   0   0   0   0   6   6   6   6   6   6   6   6   6
3 |  0   0   0   5   5   5   5   6   6   6  11  11  11  11  11  11
4 |  0   0   0   5   5   5   5   6   8   8  11  13  13  13  13  14
5 |  0   0   0   5   5   5   5   6   8   8  11  13  13  13  13  14
6 |  0   0   0   5   5   5   5   6   8   8  11  13  13  13  13  14
7 |  0   0   0   5   5   5   5   6   8   8  11  13  13  13  13  14
8 |  0   0   4   5   5   9   9   9   9  10  12  13  15  17  17  17
9 |  0   0   4   5   5   9   9   9   9  11  12  13  15  17  17  17
10 |  0   0   4   5   5   9   9   9   9  11  12  13  15  17  17  17
11 |  0   0   4   5   5   9   9   9   9  11  12  13  15  17  17  17
12 |  0   0   4   5   5   9   9  11  12  12  16  16  16  18  18  19
13 |  0   0   4   5   5   9   9  11  12  12  16  16  18  19  19  23
14 |  0   0   4   5   5   9   9  11  12  12  16  16  18  19  19  23
15 |  0   0   4   5   5   9   9  11  12  12  16  16  18  19  19  23
16 |  0   0   4   5   5   9   9  11  12  12  16  16  18  19  19  23
17 |  0   0   4   5   5   9   9  11  12  12  16  16  18  19  19  23
-----
Maximum Site-Seeing Value : 23
• SAKIB  AA Test
  ↵
git main
```

From the resultant table, maximum site-seeing value for tourist is 23, in the budget of 15.

Complex - Problem Solving Questions:

(a) Does the solution need in-depth engineering knowledge?

⇒ Yes, it requires deep expertise in transportation systems, algorithms, data analysis, and software development.

(b) Does the solution involve wide-ranging or conflicting technical, engineering and other issues?

⇒ Yes, it involves a variety of technical challenges, including optimizing routes, integrating real-time data and balancing user preferences.

(c) Is the solution well-known, or does it require abstract thinking and analysis to formulate?

⇒ While components are known, integrating them into a smart city context requires abstract thinking and analysis.

(d) Does the solution involve infrequently encountered issues?

⇒ Yes, especially in dynamically updating routes and balancing conflicting user preferences.

(e) Does the Solution need adherence to Standards and Codes of Practice?

⇒ Yes, adherence to data Privacy, Security and transportation regulations is necessary.

(f) Does the Solution involve Stakeholders with Conflicting technical requirements?

⇒ Yes, inhabitants and tourists may have conflicting Priorities in route optimization.

(g) Does the Solution involve Interdependence between Sub-Problems or Parts?

⇒ Yes, various Components of the System are Interdependent, requiring careful Coordination