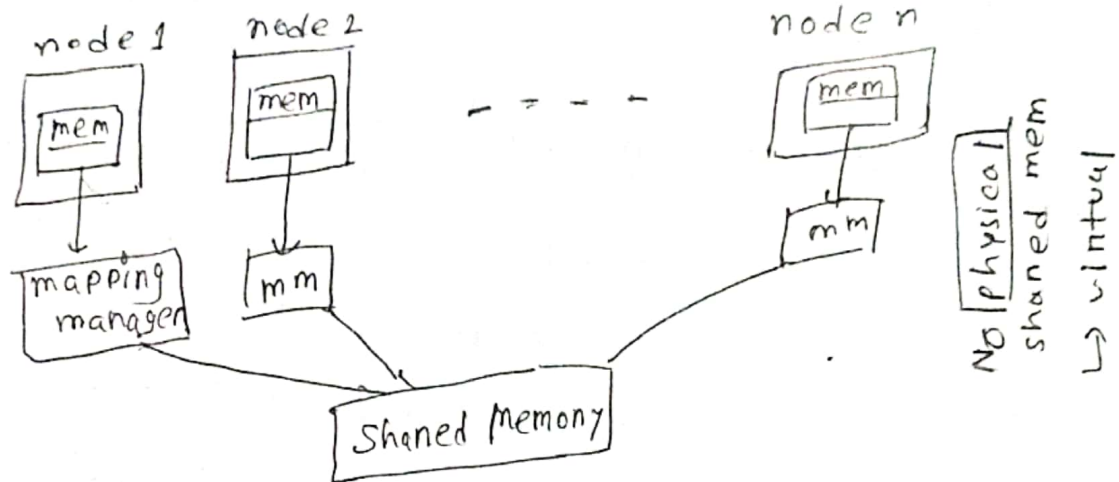


collection of many nodes/devices which are connected through some network and all have local memories.

Distributed Shared Memory

→ provides a virtual address space

shared among all nodes in DSM



What is mapping manager?

→ A layer of s/w, perhaps bundled with the OS or as a runtime library routine.

→ When a process accesses data in the shared address space, the mapping manager maps shared memory address to physical memory (local or remote)

Pros:

- 1) Easy Abstraction: Same address space
→ easy data migration
→ simpler than RPC
- 2) Easier portability → common interface
- 3) Locality of data:
→ data fetched in large blocks
→ need in future
- 4) Larger memory space: 7) Variables direct share
- 5) Better performance 8) Less cost
- 6) Flexible communication environment

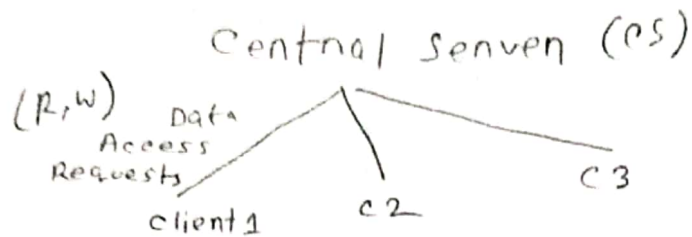
Cons:

- 1) ~~fast~~ ^{slow} Access
- 2) Not efficient
- 3) Process can't run simultaneously
- 4) Common bus → more traffic
- 5) communication delay

Types of algo:

- 1) Central - server
- 2) Data migration
- 3) Read - replication
- 4) Full replication

1) Central Server Algorithm



CS → data stored
C2 → access
data stored in

- simple to implement
- CS maintains all the shared data $\begin{cases} \rightarrow \text{read} \\ \rightarrow \text{write} \end{cases}$
- cons: performance not good
 - not reliable

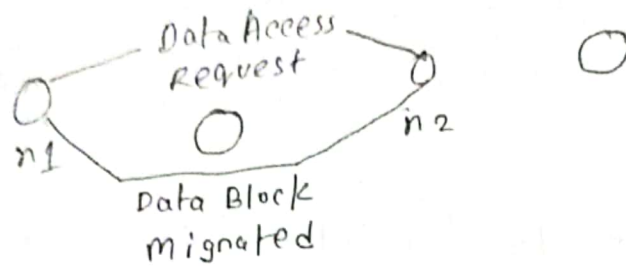
Soln: 1) Use a mapping function to distribute / locate data

2) partition shared data between several servers.

- * if timeout → resend request
- * repeated failure → send it
- * Detect duplicate write requests
 - with associated sequence number

2) Migration Algo:

→ migrate data elements on request



→ on block of data req^d, a block of request
on that data migrate req^d for further
access and processing

→ only 1 node can access data at a time

→ whole block is migrated to that node

3) Read replication Algo:

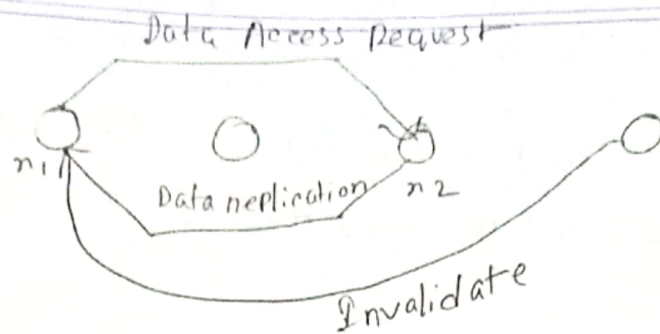
→ replicate data at multiple nodes for read
access

→ write:

1) invalidate all copies of shared data at
various nodes

2) update with modified value

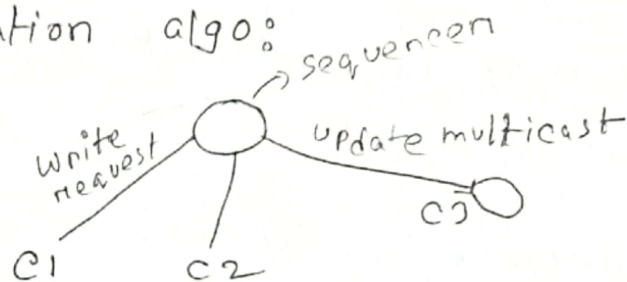
→ low read and high write cost



2 imple:

- 1) INV
- 2) PLUS

4) Full-replication algo:



* Memory coherence:

→ most updated value is returned by read operation (expected)

→ control access

* Sequential consistency:

→ The result of any execution of operations of all processors is the same as if they were executed in sequential order.

* General consistency:

All copies of mem loc contains same data - after write by cpu complete

* Processon consistency:

operations issued by processon are performed in the order they're issued.

* Weak consistency:

memory is consistent only after a synchronization operation.

* Release consistency:

→ further relaxation of weak consistency

→ acquire and release
(lock) (unlock)

* Coherece protocol?

→ why needed:

• how ensure -

→ all replicas have same info

→ nodes not access state data
(old)

Types of protocols: (2)

1) write-invalidate protocol:

→ invalidates data of all copies except 1 before write

→ can't access invalidated data

Ex: IVY, clouds, Dash

Pros: - Locality of reference
- ~~too~~ many updates

good performance
in these

Cons: inefficient

2) Write - update protocol:

→ all copies data update

Cons: more complex

* Cache coherence protocols: 1) PLUS system
2) Munin " +