#### Table of contents

```
Problem Statement
```

```
Computable bijections on \{0,1\}^6
4*16 property
Combine
Conclude
```

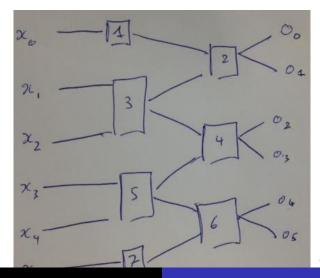
Computable almost-bijections on  $\{0,1\}^6$ 2\*161517 property Conclusion

Conclusion

### Does a 6 bit counter exists?

Can we create a **deterministic** network that will iterate over all 6 bits string – in any order – ?

### Does a 6 bit counter exists?



## Can we bruteforce that problem ?

There are 2<sup>44</sup> differents networks which is approximatively 17000 billions.

We could hope for an answer in a few days.

**But** we can drastically restrict our search space with a **few** observations.

Main idea: we are going to bruteforce on **sub networks** and select those who respect a **necessary** condition.

## How such a network's dynamic would look like?

There are only two possibilities (up to any permutation):

### Implication on the computed function by the network

The function that our counting network coumputes is either:

- ► A bijection
- ▶ An almost-bijection f i.e,  $\exists ! x_0, y_0 \in \{0, 1\}^6$  such that f' is a bijection where:

$$\forall x \neq x_0 \quad f'(x) = f(x)$$
  
 $f'(x_0) = y_0$ 

We are going to check on both cases.

If we have a bijection it will **enumerate** all strings of  $\{0, 1\}^6$ . After reordordering it will look like:

Let's focus on the first two bits, we can organise our sequences this way:

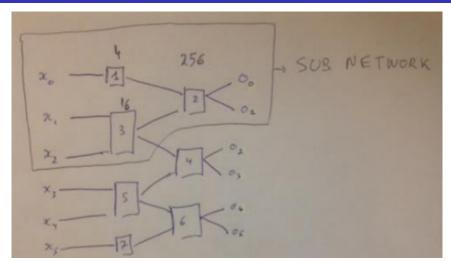
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00000000000000000
                                        11111111111111111
                                                            11111111111111111
00000000000000000
000000000000000000
                    11111111111111111
                                        00000000000000000
                                                            111111111111111111
0000000011111111
                    0000000011111111
                                        0000000011111111
                                                            0000000011111111
0000111100001111
                    0000111100001111
                                        0000111100001111
                                                            0000111100001111
0011001100110011
                    0011001100110011
                                        0011001100110011
                                                            0011001100110011
0101010101010101
                    0101010101010101
                                        0101010101010101
                                                            0101010101010101
```

Let get rid of the last 4 bits:

We see that each of the 2 bits patterns: **00**, **01**, **10**, **11** occurs **16** times in this enumeration.

It's the 4 \* 16 property.

**Hence** the sub network responsible for these 2 bits must have the **4\*16** property to be eligible as a being a sub network of a 6 bits bijective counter network.

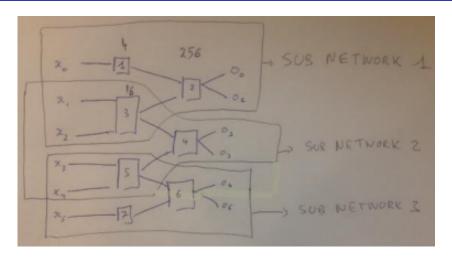


## Compute all the 4 \* 16 sub networks

By exausthive search we find **288** (over 4\*16\*256 = 16384) subnetworks with this property.

By removing equivalent networks we are left with **72 circuits for the first 2 bits**.

# 4\*16 holds for middle and ending bits



# 4\*16 holds for middle and ending bits

- ▶ **72** networks for the first 2 bits
- ▶ 216 networks for the middle 2 bits
- ▶ 72 networks for the last 2 bits

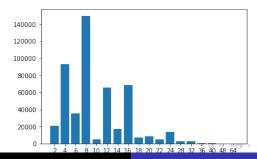
### How to conclude I? Combine it!

**Hence** by combining these we have  $72*216*72 \simeq 10^6$  6 bits networks to test.

#### How to conclude II? Count orbits!

Over all these potential networks we count **497664** bijections. For each of these bijections we have to **count their orbits**, we have a winner **iif it has only 1 orbit**.

We do not find such a network, here there's the histogram of orbits:



#### Conclusion

**Conclusion**: There are no bijective 6bits counters.

# 2\*161517 property

If we have an almost-bijection all 6 bits sequences will be reached by our network but one.

Also one bit string will be reached twice.

It means that at least one of our sub network will see:

- 2 patterns 16 times
- ▶ 1 pattern **15** times
- ▶ 1 pattern **17** times

## 2\*161517 property

We call this the 2\*161517 property.

### 2\*161517 does not occur!

By enumeration, there exist no sub network with the 2\*161517 property.

Hence, we cannot hope for an almost-bijective counter.

# No 6bit counter network :'(

By case distinction, there's no 6bit counter network.

But there are 0 to 62 counters and this kind of methods helps to exhibit at least 24.

### Sources?

All our sources for these computations are available here:

```
https://github.com/cosmo-sterin/ER_MolProg_Project2/tree/master/circuit_sim
```