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#### Problem Statement

Computable bijections on  $\{0,1\}^6$

Computable bijections on  $\{0,1\}^6$

Conclusion

## Does a 6 bit counter exists ?

Can we create a **deterministic** network that will iterate over all 6 bits string – in any order – ?



# Can we brute force that problem ?

There are  $2^{44}$  different networks which is approximately 17000 billions.

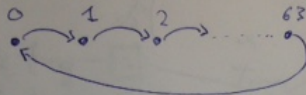
We could hope for an answer in a few days.

**But** we can drastically restrict our search space with a **few** observations.

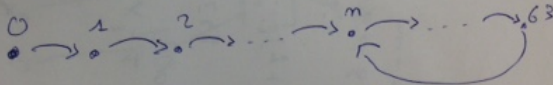
## How such a network's dynamic would look like ?

There are only two possibilities (up to any permutation):

Case 1:



Case 2:



## Implication on the computed function by the network

The function that our counting network computes is either:

- ▶ **A bijection**
- ▶ **An almost bijection**  $f$  i.e,  $\exists! x_0, y_0 \in \{0, 1\}^6$  such that  $f'$  is a bijection where:

$$\begin{aligned}\forall x \neq x_0 \quad f'(x) &= f(x) \\ f'(x_0) &= y_0\end{aligned}$$

We are going to check on both cases.

## The $4 * 16$ property

If we have a bijection it will **enumerate** all strings of  $\{0,1\}^6$ .  
After reordering it will look like:

```
0000000000000000000000000000000011111111111111111111111111111111
000000000000000000001111111111111100000000000000001111111111111111
00000000011111111000000001111111100000000111111110000000011111111
0000111100001111000011110000111100001111000011110000111100001111
0011001100110011001100110011001100110011001100110011001100110011
0101010101010101010101010101010101010101010101010101010101010101
```

## The $4 * 16$ property

Let's focus on the two first bits, we can organise our sequences this way:

0000000000000000	0000000000000000	1111111111111111	1111111111111111
0000000000000000	1111111111111111	0000000000000000	1111111111111111
0000000011111111	0000000011111111	0000000011111111	0000000011111111
0000111100001111	0000111100001111	0000111100001111	0000111100001111
0011001100110011	0011001100110011	0011001100110011	0011001100110011
0101010101010101	0101010101010101	0101010101010101	0101010101010101

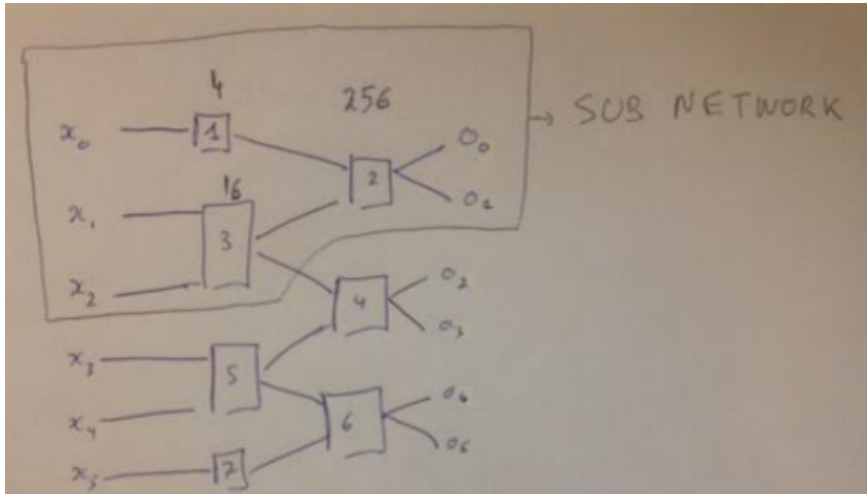
We see that **each of the 2 bits pattern happens 16 times**.  
 It's the  **$4 * 16$  property**.



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## The 4 \* 16 property



## Compute all the $4 * 16$ sub networks

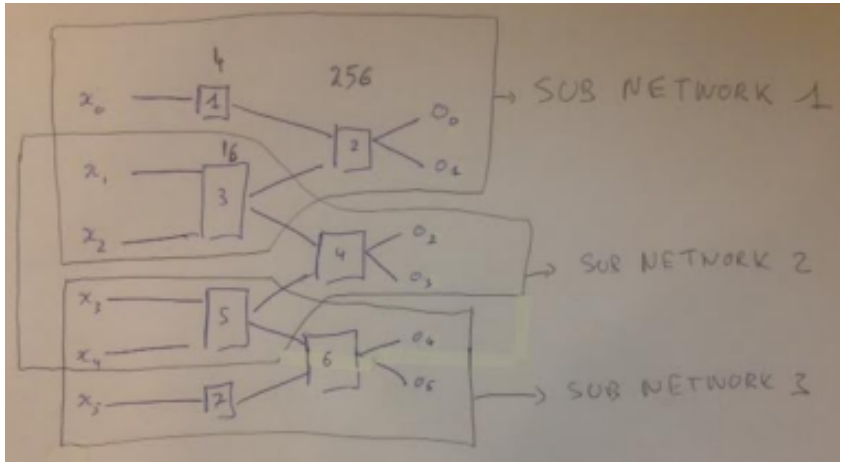
By exhaustive search we find **288** (over 16384) sub networks with this property.

By removing equivalent networks we are left with **72 circuits for the 2 first bits**.

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4\*16 holds for middle and ending bits



## 4\*16 holds for middle and ending bits

- ▶ **72** networks for the 2 first bits
- ▶ **216** networks for the 2 middle bits
- ▶ **72** networks for the 2 last bits

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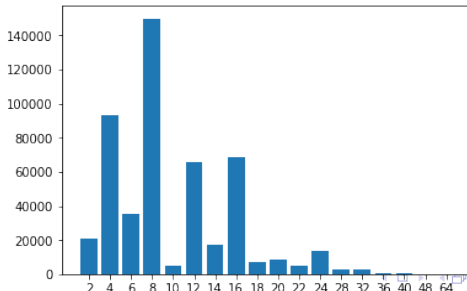
# How to conclude I ? Combine it !

**Hence** by combining these we have  $72 * 216 * 72 \simeq 10^6$   
6 bits networks to test.

## How to conclude II ? Count orbits !

Over all these potential networks we count **497664** bijections.  
For each of these bijections we have to **count their orbits**, we  
have a winner **iif it has only 1 orbit**.

We do not find such a network, here there's the histogram of  
orbits:



# Conclusion

**Conclusion:** There are no bijective 6bits counters.

## 2\*161517 property

If we have an almost-bijection **all 6 bits sequences will be reached by our network but one.**

Also one bit string will be reached **twice.**

It means that **at least one** of our sub network will see:

- ▶ 2 patterns **16** times
- ▶ 1 pattern **15** times
- ▶ 1 pattern **17** times



# 2\*161517 property

We call this the **2\*161517 property**.

$2^{*161517}$  does not occur !

By enumeration, **there exist no sub network with the  $2^{*161517}$  property.**

**Hence**, we cannot hope for an **almost-bijective counter**.

## No 6bit counter network :'(

By case distinction, **there's no 6bit counter network.**

But there are **0 to 62** counters and this kind of methods helps to exhibit **at least 24**.

## Sources ?

All our sources for these computations are available here:

`https://github.com/cosmo-sterin/ER\_MolProg\_Project2/  
tree/master/circuit\_sim`