CS 496: Homework Assignment 4 Due: 9 April, 11:55pm

1 Assignment Policies

Collaboration Policy. It is acceptable for students to collaborate in understanding the material but not in solving the problems or programming. Use of the Internet is allowed, but should not include searching for existing solutions.

Under absolutely no circumstances code can be exchanged between students. Excerpts of code presented in class can be used.

Assignments from previous offerings of the course must not be re-used. Violations will be penalized appropriately.

2 Assignment

This assignment has two parts. Part I consist in adding records with mutable fields to EXPLICIT-REFS to obtain EXPLICIT-REFS-MF. Part II consists in writing a series of programs in the language EXPLICIT-REFS-MF that make use of this new language construct.

2.1 Part I

As seen in class in OCaml, a mutable field in a record is one whose contents may be updated through assignment. One can declare a field to be mutable by using "<=" instead of "=". In the next example, a record is defined that holds two items of personal data, a social security number and the age. The ssn field in the example below is immutable but the age field is mutable; age is then updated to 31:

```
let p = {ssn = 10; age <= 30}
in begin
    p.age <= 31; (* updates age to 31 *)
    p.age
    end</pre>
```

Evaluating this expression should produce Ok (NumVal 31). Consider the evaluation of this other expression:

It should produce Ok (RecordVal [("ssn", (false, NumVal 10)); ("age", (true, RefVal 1))]). Notice two things here:

- 1. false indicates the field is immutable and true that it is mutable.
- 2. Mutable fields are always assigned a reference. In this example, the store will hold the expressed value NumVal 31 at location 1.

Updating an immutable field should not be allowed. For example, the following expression should report an error Error "Field not mutable":

```
let p = { ssn = 10; age <= 20}
in begin
    p.ssn <= 11;
    p.age
    end</pre>
```

EXPLICIT-REFS-MF also has a predicate to check whether a run-time value is a number or not. Here is how it should behave:

```
# interp "number?(2)";;
- : exp_val Explicit_refs.Ds.result = Ok (BoolVal true)
# interp "number?(2+1)";;
- : exp_val Explicit_refs.Ds.result = Ok (BoolVal true)

utop # interp "number?(zero?(0))";;
- : exp_val Explicit_refs.Ds.result = Ok (BoolVal false)

utop # interp "number?(proc(x) {x+1})";;
- : exp_val Explicit_refs.Ds.result = Ok (BoolVal false)

utop # interp "number?(proc(x) {x+1})";;
- : exp_val Explicit_refs.Ds.result = Ok (BoolVal false)
```

The updated concrete syntax is:

The updated abstract syntax is as follows:

```
type expr =

2    ...
    Record of (string*(bool*expr)) list

4    Proj of expr*string
    SetField of expr*string*expr

6    IsNumber of expr
```

For example,

You are asked to implement the interpreter extension. The RecordVal constructor has added for you.

As for eval_expr, the case for Record has already been implemented for you. You are asked to implement Proj, SetField and IsNumber:

```
let rec eval_expr : expr -> exp_val ea_result = fun e ->
     match e with
     | IsNumber(e) ->
        failwith "implement"
     Record(fs) ->
       sequence (List.map process_field fs) >>= fun evs \rightarrow
       return (RecordVal (addIds fs evs))
     | Proj(e,id) ->
8
        failwith "implement"
10
      | SetField(e1,id,e2) ->
        failwith "implement"
12
      | IsNumber(e) ->
        failwith "implement"
    process_field (_id,(is_mutable,e)) =
16
     eval_expr e >>= fun ev ->
     if is_mutable
     then return (RefVal (Store.new_ref g_store ev))
18
     else return ev
```

where sequence is define here:

```
let rec sequence : ('a ea_result) list -> ('a list) ea_result =

fun cs ->
  match cs with

| [] -> return []
  | c::t ->

c >>= fun v ->
  sequence t >>= fun vs ->
  return (v::vs)

ds.ml
```

2.2 Part II: Implementing Linked Lists in EXPLICIT-REFS-MF

This part of the assignment asks you to implement linked-list operations in the language EXPLICIT-REFS-MF. Since we do not have a print operation, in order to test your code you will have to insert a debug instruction and inspect the contents of the environment and store. Following the discussion in class using OCaml, a linked-list will be implemented as a record with two fields head and size, both of which are mutable. The head field may either contain o for null or may be a node record. Node records have fields data and next, both of which are mutable. We will start with some examples. EXPLICIT-REFS-MF source code must be placed in text files with extension .exr.

Example 1 below creates a linked-list with two nodes.

```
(* Example 1 *)
2
   let 11 = { head <= 0; size <= 0}
                                              (* 0 in head signals null *)
   in let add_front = proc (x) { proc (1) {
                 1.head <={ data <=x; next <= 1.head };</pre>
                 1.size <= 1.size+1
              end
              } }
   in begin
         ((add_front 2) 11);
         ((add_front 3) 11);
12
         debug(11) (* required to inspect the list *)
14
                                                                       ll_add_front.exr
```

You may run this in utop as follows:

Example 2 below creates a linked list with two nodes and then adds one to each data element in each node.

```
end } }
   in letrec bump_helper (nd) =
10
       if number?(nd)
       then 0
       else (begin
12
                  nd.data <= nd.data + 1;
                  (bump_helper nd.next)
14
             end)
16
   in let bump = proc (ll) { (bump_helper ll.head) }
   in begin
         ((add_front 2) 11);
         ((add_front 3) 11);
         (bump 11);
20
         debug(11)
       \verb"end"
22
                                                                              11_bump.exr
```

You may run this in utop as follows:

```
utop # interpf "ll_bump";;
>>Environment:
 11->{head <= RefVal (0); size <= RefVal (1)},
 add_front->ProcVal(x,Proc(1,BeginEnd(Var 1.head<={data=Var x;next=Var 1.head};
   Var 1.size <= Add(Var 1.size, Int 1))),[11->{head <= RefVal (0); size <= RefVal (1)}]),
 bump_helper->Rec(nd,IfThenElse(Number?(Var nd),Int 0,BeginEnd(
   Var nd.data<=Add(Var nd.data,Int 1);App(Var bump_helper,Var nd.next)))),</pre>
 bump->ProcVal(11, App(Var bump_helper, Var 11.head), [11->{head <= RefVal (0);
  size <= RefVal (1)},add_front->ProcVal(x,Proc(1,BeginEnd(Var 1.head<={data=Var x;next=Var 1.head}
>>Store:
 0->{data <= RefVal (4); next <= RefVal (5)},
1->NumVal 2,
 2->NumVal 3,
3 \rightarrow NumVal 0,
 4 \rightarrow NumVal 4,
5->{data <= RefVal (2); next <= RefVal (3)}
 - : exp_val Explicit_refs.Ds.result = Error "Reached breakpoint"
                                                                               utop
```

You are asked to implement the following exercises. The stubs are provided for you in the src folder. You may include helper functions.

1. ll_add_last.exr. For example,

```
utop # interpf "ll_add_last";;
>>Environment:

11->{head <= RefVal (0); size <= RefVal (1)},
    add_last_helper->...,

add_last->...,
>>Store:

0->{data <= RefVal (2); next <= RefVal (3)},
    1->NumVal 3,

2->NumVal 2,
    3->{data <= RefVal (4); next <= RefVal (5)},

4->NumVal 3,

5->{data <= RefVal (6); next <= RefVal (7)},

6->NumVal 4,
    7->NumVal 0

15 -: exp_val Explicit_refs.Ds.result = Error "Reached breakpoint"
```

utop

2. ll_remove_first.exr For example,

```
utop # interpf "ll_remove_first";;
 >>Environment:
 11->{head <= RefVal (0); size <= RefVal (1)},</pre>
 add_front->ProcVal(x,Proc(1,BeginEnd(Var 1.head<={data=Var x;next=Var 1.head};
   Var 1.size <= Add(Var 1.size, Int 1))),[11->{head <= RefVal (0); size <= RefVal (1)}]),
 remove_first-> ...,
 >>Store:
 0 \rightarrow \{ data \le RefVal (4); next \le RefVal (5) \},
1->NumVal 2,
 2->NumVal 2,
3->NumVal 0,
 4->NumVal 3,
5->{data <= RefVal (2); next <= RefVal (3)},
 6->NumVal 4,
7->{data <= RefVal (4); next <= RefVal (5)}
 - : exp_val Explicit_refs.Ds.result = Error "Reached breakpoint"
                                                                           utop
```

You may assume the list is not empty.

3. ll_remove_last.exr For example,

```
utop # interpf "ll_remove_last";;
>>Environment:
 11->{head <= RefVal (0); size <= RefVal (1)},
 add_front->ProcVal(x,Proc(1,BeginEnd(Var 1.head<={data=Var x;next=Var 1.head};
   Var 1.size <= Add(Var 1.size, Int 1))),[11->{head <= RefVal (0); size <= RefVal (1)}]),
remove_last_helper->...,
 remove_last->...,
>>Store:
 0->{data <= RefVal (6); next <= RefVal (7)},
1->NumVal 2,
 2->NumVal 2,
3->NumVal 0,
 4->NumVal 3.
5->NumVal 0,
 6 -> NumVal 4,
7->{data <= RefVal (4); next <= RefVal (5)}
 - : exp_val Explicit_refs.Ds.result = Error "Reached breakpoint"
                                                                        utop
```

You may assume the list is not empty.

3 Submission instructions

Submit a file named HW4.zip through Canvas. Include the original stub provided. Please write the names of the members of the team in the as a comment at the top of the interp.ml file.