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Question 2:

2.a.i : The node P_i in S will take a color different than the neighbors nodes colors in S form the group of color so we need two or max three colors if the nodes number is odd then look in node in $P \setminus S$, where they do the same and find the node which has the same color to match with.

2.a.ii : The route will be from S to $P \setminus S$

2.a.iii : After matching occurs the P_i will ask for synchrony and get ack, then P_i sends Synchrony ACK and sends the message then P_j will replay

2.b.i : Self-stabilizing Vertex coloring in chapter 2, also Self-stabilizing Maximum Matching in chapter 2.9.

2.b.ii : Yes it solves vertex color matching, Since this algorithm (Maximum Matching) has the ability to find the maximal in every set and the maximum also, It has the algorithm to define the match node and single nodes, free node, chain nodes, and the waiting nodes to match.

- Matching happens when P_i in S pecks to P_j in $P \setminus S$ and P_j in $P \setminus S$ pecks to P_i in S .
- Waiting node when P_i in S pecks to P_j in $P \setminus S$ and P_j in $P \setminus S$ pecks to no one.
- Chain node when P_i in S pecks to P_j in $P \setminus S$ and P_j in $P \setminus S$ pecks k , $k \neq P_i$.
- Free node when P_i in S pecks to no one and P_j in $P \setminus S$ pecks no one.
- Single node when P_i in S pecks to no one and P_j in $P \setminus S$ pecks k , $k \neq P_i$.

2.b.iii :

2.c.i :

2.c.ii :

2.c.iii :