

# Introduction to Communication Engineering

## SSY121, Lecture # 9

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## Computer lab

The computer lab (room 5225) has updated their hardware, so our headphones can no longer be used. Instead, please use the rooms ES61, ES62, and ES63 on floor 6 in Linsen in the EDIT house. The rooms will not be booked, but they should be available after 17:00.

## Sign up for exam

The exam in this course is as you know on Wed Oct 30 at 08:30-12:30. Please remember to sign up for the exam. Chalmers' rule is that if you have not registered for the exam, you will not be allowed to take the exam. **The last day to register is Oct 10!**

# Outline

## 1 Review and Some Definitions

- The AWGN channel
- Review of the vector channel model for the AWGN channel
- Some Definitions

## 2 Pairwise error probability

## 3 Error probabilities

- The average error probability
- The union bound and a high SNR approximation
- Bit Error probability and Gray code

# Part I

## Review and Some Definitions

## A very nice property of the AWGN

If  $n(t)$  is a white Gaussian random process with PSD  $N_0/2$ , then the random variables  $n_i$ , the projections of  $n(t)$  onto **any** set of orthonormal basis functions, are i.i.d. Gaussian random variables with zero mean and variance  $N_0/2$ . (Theorem 2.5-1 in [Anderson])

## The continuous-time AWGN channel in signal space

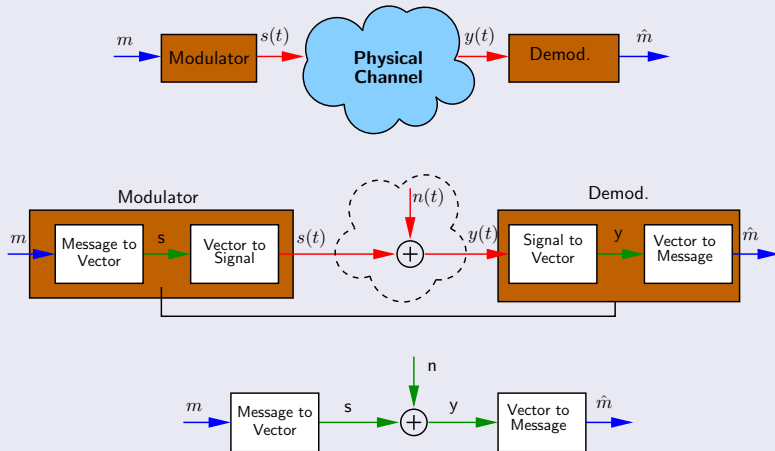
- If  $n(t)$  is AWGN with PSD  $N_0/2$ , the continuous time channel  $y(t) = s(t) + n(t)$ , can be replaced by a vectorial channel model

$$\mathbf{y} = \mathbf{s} + \mathbf{n},$$

where  $\mathbf{s} = [s_1, \dots, s_N]$  is the transmitted symbol,  $\mathbf{n} = [n_1, \dots, n_N]$  is the noise, and  $\mathbf{y} = [y_1, \dots, y_N]$  is the received signal.

- The components of the noise vector  $n_i$  for  $i = 1, \dots, N$  are i.i.d. zero-mean Gaussian random variables with variance  $N_0/2$ .

## The vector channel model for the AWGN channel



## Definitions (see also $L_2$ )

- $M$ : Number of symbols
- $l$ : Number of bits per symbol
- $R_s$ : Symbol rate ( $T_s$ : symbol period)
- $R_b$ : Bit rate ( $T_b$ : bit period).  $R_b = lR_s$
- $E_s$ : Average symbol energy:

$$E_s = \sum_{i=1}^M \mathbb{P}[S = s_i] \|s_i\|^2$$

- $E_b$ : Average bit energy.  $E_b = E_s/l$
- Signal to noise-ratio:  $\text{SNR} = E_s/N_0$
- Decibel: logarithmic measure of the ratio between two power or energies. For example  $\text{SNR}_{\text{dB}} = 10 \log_{10}(E_s/N_0)$

## Example (what can you do after this lecture?)

Consider equally likely 4-PAM transmission where  $s \in \{-3A, -A, A, 3A\}$ .

- Compute  $A$  such that  $E_s = 1$
- Compute the relation between  $E_s$  and  $E_b$
- Draw the constellation and the decision regions
- Compute the SNR in dB if  $N_0 = 0.5$
- Compute the conditional PDF  $f_{Y|S}(y|S = -3A)$  for SNR= 3 dB, for SNR= -3 dB, and for SNR= 10 dB
- Compute the probability of mistaking the transmitted symbol  $s = -3A$  by the symbol  $-A$ ,  $A$ , and  $+3A$  for the three SNR values



## Pairwise error probability (PEP) definition

The PEP is simply the probability of mistaking the symbol  $s_i$  by  $s_j$ , i.e.,

$$\text{PEP}^{(i,j)} = \mathbb{P} [\|Y - s_j\| \leq \|Y - s_i\| | S = s_i]$$

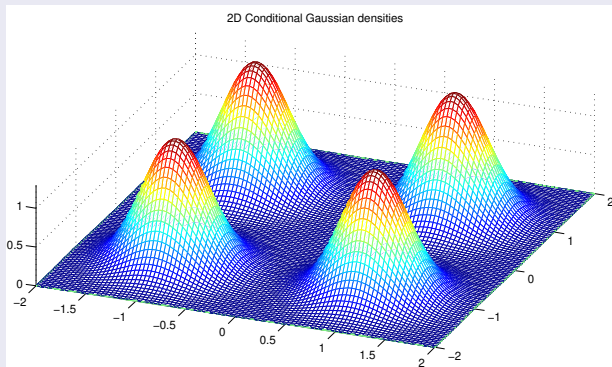
## PEP in 1D for an AWGN channel

$$\begin{aligned} \text{PEP}^{(i,j)} &= \mathbb{P} [\|Y - s_j\| < \|Y - s_i\| | S = s_i] \\ &= \int_{y_0}^{\infty} f_{Y|S}(y|s_i) dy = Q\left(\frac{y_0 - s_i}{\sigma}\right) \\ &= Q\left(\frac{s_i + s_j - 2s_i}{2\sqrt{N_0/2}}\right) = Q\left(\sqrt{\frac{D_{i,j}^2}{2N_0}}\right), \end{aligned}$$

where  $y_0 = (s_i + s_j)/2$  is halfway between the signals,  $\sigma^2 = N_0/2$  the variance of the noise, and  $D_{i,j}^2$  is the squared Euclidean distance between  $s_i$  and  $s_j$ , i.e.,  $D_{i,j}^2 = \|s_j - s_i\|^2$  (see also L7).

## PEP in 2D: Still the same (Why?)

$$\text{PEP}^{(i,j)} = Q\left(\sqrt{\frac{D_{i,j}^2}{2N_0}}\right).$$



## PEP in $N$ dimensions

Still the same! (Why?)

$$\text{PEP}^{(i,j)} = Q \left( \sqrt{\frac{D_{i,j}^2}{2N_0}} \right).$$

## What we need to know?

- The squared Euclidean distance between the two points  $D_{i,j}^2$
- The noise variance  $N_0/2$

## The “border” (decision region)

- In 1D, the “border” is a point
- In 2D, the “border” is a line
- In 3D, the “border” is a plane
- In 4D, ...

## A general error probability formula

The error probability when  $s_i$  is transmitted is

$$\mathbb{P}[\text{error}|S = s_i] = \mathbb{P}[Y \notin \mathcal{R}_i|S = s_i] = \sum_{j \neq i} \mathbb{P}[Y \in \mathcal{R}_j|S = s_i]$$

The **exact** average error probability  $P_e$  is

$$\begin{aligned} P_e &= \sum_{i=1}^M \mathbb{P}[\text{error}|S = s_i] \mathbb{P}[S = s_i] \\ &= \sum_{i=1}^M \mathbb{P}[S = s_i] \sum_{j \neq i} \mathbb{P}[Y \in \mathcal{R}_j|S = s_i] \end{aligned}$$

And what is the problem with this?

In general, we do not know how to compute  $\mathbb{P}[Y \in \mathcal{R}_j|S = s_i]$  😞

## An exception to the rule...

For binary transmission ( $M = 2$ ), we do know how to compute the exact  $P_e$ , since  $\mathbb{P}[Y \in \mathcal{R}_j | S = s_i] = \text{PEP}^{(1,2)}$ .

## For the particular case of $M = 2$ and equally likely symbols

For binary modulation ( $M = 2$ ), the exact average error probability  $P_e$  can be written in terms of the PEP

$$\begin{aligned} P_e &= \mathbb{P}[S = s_1] \text{PEP}^{(1,2)} + \mathbb{P}[S = s_2] \text{PEP}^{(2,1)} \\ &= \frac{1}{2} Q \left( \sqrt{\frac{D_{1,2}^2}{2N_0}} \right) + \frac{1}{2} Q \left( \sqrt{\frac{D_{2,1}^2}{2N_0}} \right) \\ &= Q \left( \sqrt{\frac{D_{1,2}^2}{2N_0}} \right) \end{aligned}$$

## The union bound (UB)

A good **approximation** of the average error probability  $P_e$  is

$$\begin{aligned} P_e &= \sum_{i=1}^M \mathbb{P}[S = s_i] \sum_{j \neq i} \mathbb{P}[Y \in \mathcal{R}_j | S = s_i] \\ &\leq \sum_{i=1}^M \mathbb{P}[S = s_i] \sum_{j \neq i} \text{PEP}^{(i,j)} = \sum_{i=1}^M \mathbb{P}[S = s_i] \sum_{j \neq i} Q \left( \sqrt{\frac{D_{i,j}^2}{2N_0}} \right) \end{aligned}$$

## The UB for equally likely symbols

$$P_e \leq \frac{1}{M} \sum_{i=1}^M \sum_{j \neq i} Q \left( \sqrt{\frac{D_{i,j}^2}{2N_0}} \right)$$

## The previous approximation...

- It is actually **not** an approximation (it is exact) for  $M = 2$

## Can we simplify it even more?

- The Q-functions decrease very fast when the argument increases
- For large arguments (high SNR), one of the Q-functions will dominate

## Definition (High-SNR approximation)

For sufficiently high SNR, the average error probability can be approximated by

$$\begin{aligned} P_e &= \sum_{i=1}^M \mathbb{P}[S = s_i] \sum_{j \neq i} \mathbb{P}[Y \in \mathcal{R}_j | S = s_i] \\ &\leq \frac{1}{M} \sum_{i=1}^M \sum_{j \neq i} Q\left(\sqrt{\frac{D_{i,j}^2}{2N_0}}\right) \approx \frac{2K}{M} \cdot Q\left(\sqrt{\frac{D_{\min}^2}{2N_0}}\right), \end{aligned}$$

where  $D_{\min} = \min_{i \neq j} \{D_{i,j}\}$  is the *minimum distance* of the constellation, and  $K$  is the number of signal pairs at minimum distance (section 2.6 in [Andersson])

## Example (Error probability of OOK and BPSK)

Which modulation scheme is better in terms of error probability, OOK or BPSK?

- Normalize both constellations to unit energy
- For a given  $N_0$ , find the PEP
- Plot both of them using Matlab

## Homeworks!

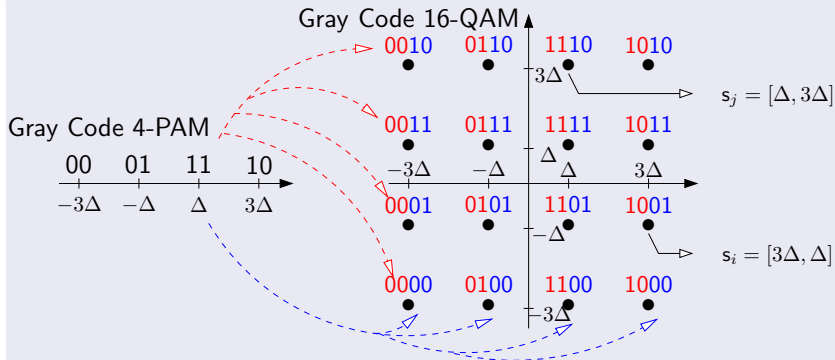
Compute and plot error probabilities (bounds/approximations) for OOK, BPSK, 4-PAM, 8-PAM, QPSK, 16-QAM, 8-PSK, etc.



## Symbol error probability and bit error probability

- The relation between symbol error probability and bit error probability is in general not simple
- It depends on the binary labeling
- The “best” we could do is to have one bit-error for each symbol error
- For the high-SNR approximation, we would like to use a binary labeling such that symbols at  $D_{\min}$  differ in only one bit
- How: **Gray Code**

## Gray code for 4-PAM and 64-QAM



## Today's Summary

- Vector channel model for the AWGN channel.
- Pairwise error probability (PEP), exact average error probability, union bound, and high-SNR approximation.
- Symbol error probability vs. bit error probability and Gray Code.