

Introduction to Communication Engineering

SSY121, Lecture # 2

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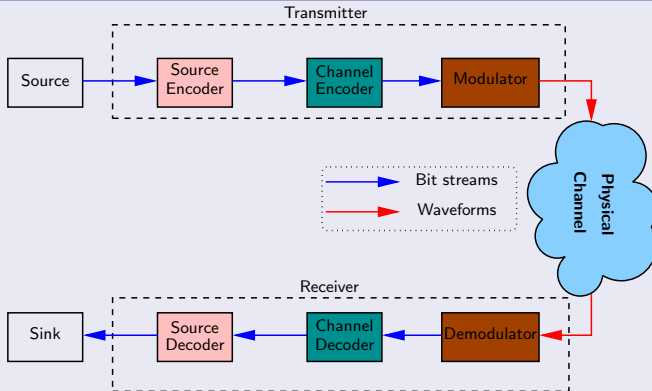
Outline

- 1 The Big Picture
- 2 The Sampling Rx
- 3 A Sequence of M -ary Pulses
- 4 The Nyquist Criterion
- 5 The Sampling Rx (cont.)
- 6 The Linear Rx for one pulse

Part I

The Sampling Rx and Nyquist Pulses

Shannon's Communication Model



The General Diagram

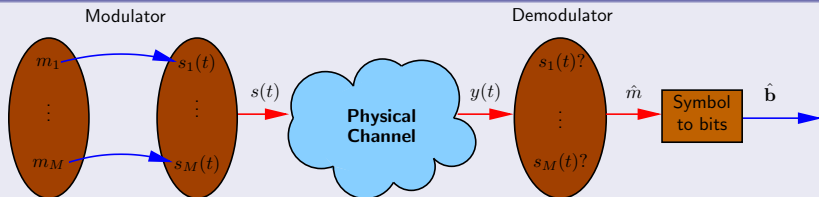


- The length- l binary codeword $\mathbf{b} = [b_1, \dots, b_l]$ has to be transmitted
- Each of the possible $M = 2^l$ codewords is mapped to a message (symbol) $m \in \{m_1, \dots, m_M\}$
- The modulator sends a continuous-time baseband signal $s(t)$ through the physical channel each T_s [s], where $s(t)$ is selected from the set of signal alternatives $\mathcal{S} = \{s_1(t), \dots, s_M(t)\}$.
- The channel introduces some distortion such that $y(t) \neq s(t)$

The General Problem

Using only the channel observation $y(t)$ during one symbol period T_s , *guess* what the transmitted message m was.

The General Problem



- A symbol is what is transmitted in each time slot (every T_s [s])
- T_s is the symbol duration and T_b is the bit duration, where $lT_b = T_s$ ($T_b = T_s$ for **binary** transmission since $l = 1$)
- The symbol rate (baud rate) is the number of symbols per second $R_s = 1/T_s$
- The bit rate is the number of bits per second $R_b = 1/T_b = l/T_s = lR_s$

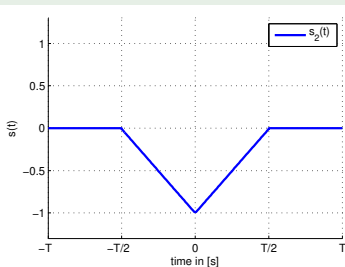
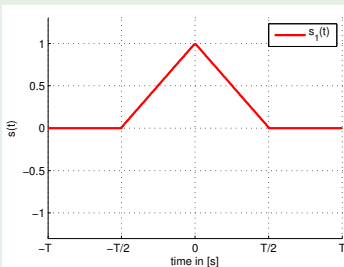
Example (Binary transmission with triangular pulse)

In binary transmission $l = 1$, $M = 2$, $b \in \{0, 1\}$ and therefore $b = 0 \Rightarrow m = m_1 \Rightarrow s_1(t)$, and $b = 1 \Rightarrow m = m_2 \Rightarrow s_2(t)$.

Moreover, we could use only one basic pulse such that

$$s(t) = \begin{cases} s_1(t) = +1 \cdot v(t) & \text{if } m = m_1 \\ s_2(t) = -1 \cdot v(t) & \text{if } m = m_2 \end{cases},$$

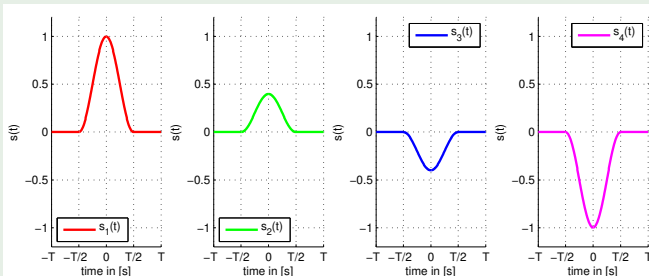
where $v(t)$ is a triangular pulse.



Example (4-ary transmission with arbitrary pulse)

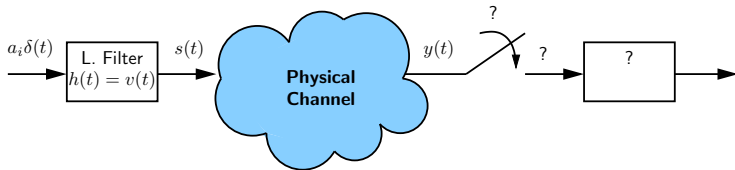
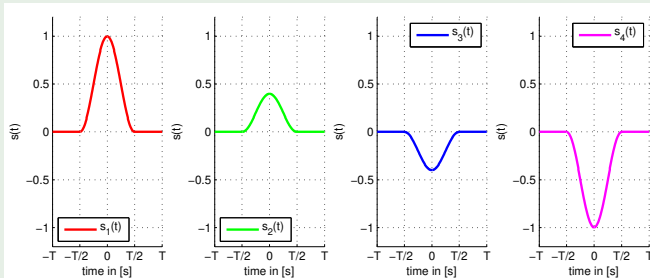
In this case $l = 2$, $M = 4$, $\mathbf{b} \in \{00, 01, 10, 11\}$ and $s_i(t) = a_i \cdot v(t)$ where $a_i \in \{+1, +0.4, -0.4, -1\}$, i.e.,

\mathbf{b}	m	$s(t)$
00	m_1	$s_1(t) = +1.0 \cdot v(t)$
01	m_2	$s_2(t) = +0.4 \cdot v(t)$
10	m_3	$s_3(t) = -0.4 \cdot v(t)$
11	m_4	$s_4(t) = -1.0 \cdot v(t)$



Example (4-ary transmission with arbitrary pulse)

Transmitted signal $s(t) = a_i \cdot v(t)$ where $a_i \in \{+1, +0.4, -0.4, -1\}$



What can we do with in the receiver?

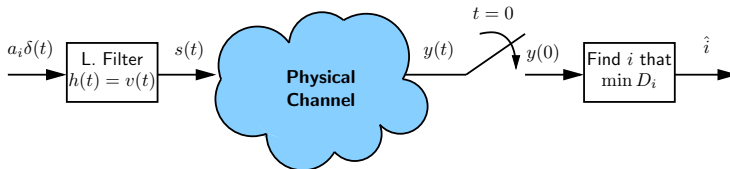
You are given a box which contains:

- a linear filter (with an impulse response you can decide)
- a “bag” with elementary signals (steps, impulses, sinusoids, ...)

The Sampling Rx

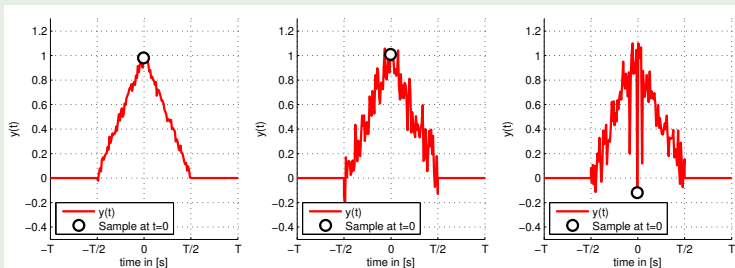
Sample the received signal $y(t)$ at $t = 0$ and choose the signal alternative that is the closest to the sample, i.e.,

$$\hat{i} = \arg \min_{i \in \{1, \dots, M\}} D_i, \text{ where } D_i = |y(0) - s_i(0)|.$$



Example (Sampling Rx for binary transmission with triangular pulse)

Assume that the transmitted message is $m = m_1 \Rightarrow s(t) = s_1(t)$, i.e., a positive triangular pulse. The received signal for three different channels is



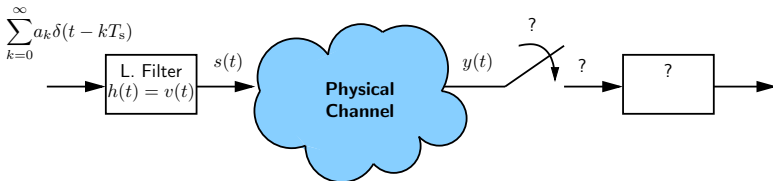
D_1	D_2	\hat{i}
$ +0.98 - 1 = 0.02$	$ +0.98 + 1 = 1.98$	1
$ +1.01 - 1 = 0.01$	$ +1.01 + 1 = 2.01$	1
$ -0.12 - 1 = 1.12$	$ -0.12 + 1 = 0.88$	2

A Sequence of M -ary Pulses

Assume that a basic pulse $v(t)$ is used for transmission. The transmitted signal is a linear superposition of time-shifted versions of $v(t)$, i.e.,

$$s(t) = \sum_{k=0}^{\infty} a_k v(t - kT_s),$$

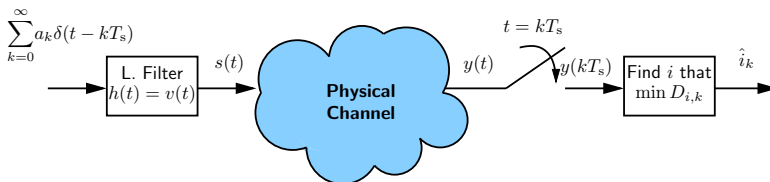
where $a_k \in \mathcal{A}$ where $|\mathcal{A}| = M$.



The Sampling Rx

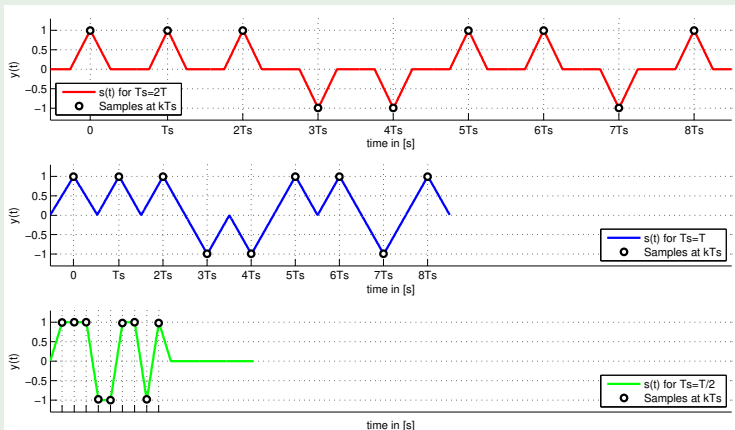
The Sampling Rx simply samples $y(t)$ at time instants $t = kT_s$ for $k = 0, 1, 2, \dots$, each time computing the M distances $D_{i,k}$ and selecting the i that gives the minimum $D_{i,k}$.

$$\hat{i}_k = \arg \min_{i \in \{1, \dots, M\}} D_{i,k}, \text{ where } D_{i,k} = |y(kT_s) - s_i(0)|.$$



Example (Sampling Rx for binary transmission with triangular pulses)

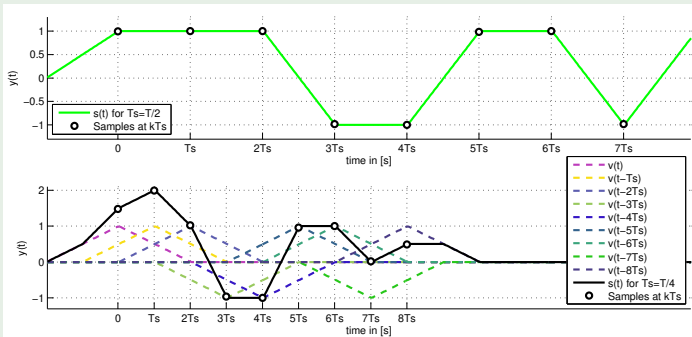
The sequence $\mathbf{b} = [1, 1, 1, 0, 0, 1, 1, 0, 1]$ is transmitted in a perfect channel using a triangular pulse of width T . The symbol rate $R_s = 1/T_s$ is increased...



Increasing and increasing R_s forever?

In the previous example we increased R_s for a fixed BW (fixed T) and the Sampling Rx gave good results. What happens if we keep increasing it?

Example (Previous example with $T_s = T/4$)



Intersymbol Interference

Intersymbol interference (ISI) appears if R_s is too high, i.e., when detecting the k th symbol, interference from other adjacent symbols appear.

Definition (Nyquist pulses)

If the pulse is chosen such that $v(kT_s) = 0$ if $k = \pm 1, \pm 2, \pm 3, \dots$ and $v(0) \neq 0$, the pulse is called “Nyquist pulse”.

Consequences

- $y(kT_s)$ depends only on the k th transmitted symbol (since the shifted versions of $v(t)$ coming from other symbols are zero at that time instant)
- No ISI for the Sampling Rx!
- In a perfect channel $y(kT_s) = a_k v(0) \Rightarrow D_k = 0$

Poisson's Sum Formula

$$\sum_{n=-\infty}^{\infty} V\left(f - \frac{n}{T_s}\right) = T_s \sum_{m=-\infty}^{\infty} v(mT_s) e^{-j2\pi f m T_s}$$

Definition (Nyquist pulses in the frequency domain)

A pulse $v(t)$ is called a Nyquist pulse if and only if its Fourier transform $V(f)$ fulfills the following two conditions

$$\begin{aligned} \sum_{n=-\infty}^{\infty} \Re \left\{ V\left(f - \frac{n}{T_s}\right) \right\} &= T_s v(0) \\ \sum_{n=-\infty}^{\infty} \Im \left\{ V\left(f - \frac{n}{T_s}\right) \right\} &= 0, \end{aligned}$$

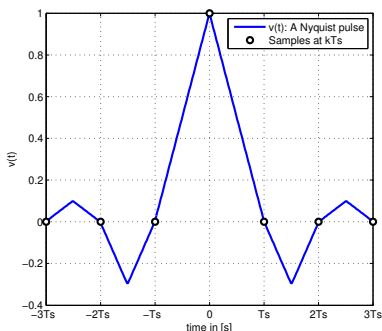
where $T_s v(0)$ is a real constant.

A particular case

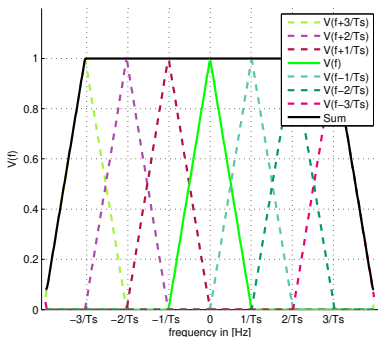
If the pulse is symmetric respect to zero, the previous definition is simplified to

$$\sum_{n=-\infty}^{\infty} V\left(f - \frac{n}{T_s}\right) = T_s v(0)$$

A Nyquist Pulse in t -domain



A Nyquist Pulse in f -domain



Part II

The Correlator/Matched Filter Rx

Symbol Rate BW Tradeoff

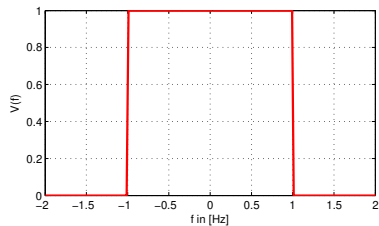
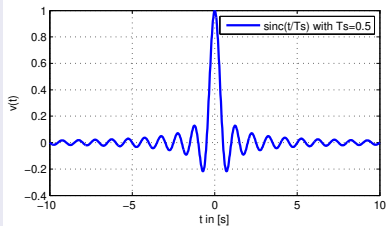
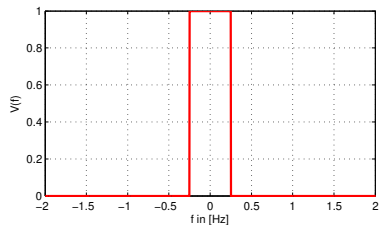
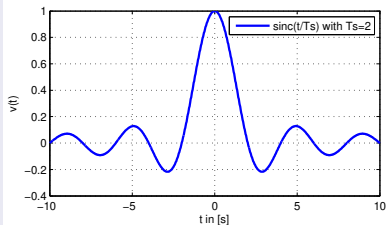
- We know that for a given R_s , ISI-free transmission is indeed possible
- We also know that if R_s increases (T_s decreases), the BW also increases
- We are interested in maximizing R_s using the least possible BW
- The interesting question is then: **Is there a fundamental limit on the BW for ISI-free transmission?**

The answer is Yes, there is a fundamental lower limit

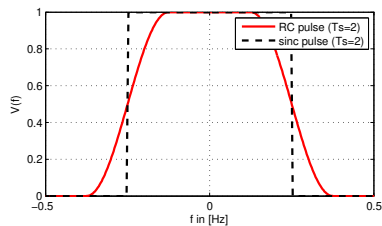
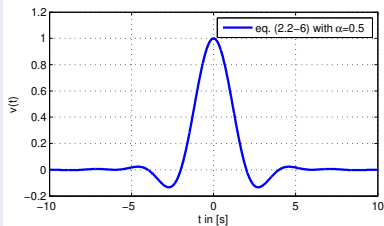
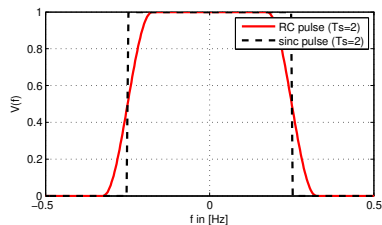
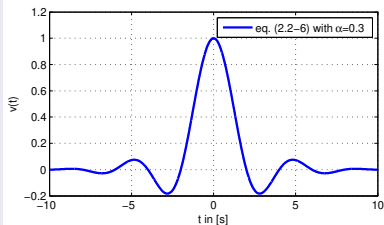
For a given R_s , ISI-free transmission is possible only for $BW \geq \frac{1}{2T_s}$, with equality if and only if $v(t) = \text{sinc}(t/T_s)$. In other words

- Many different pulses offer ISI-free transmission at expenses of $BW > \frac{1}{2T_s}$
- Only one pulse offers ISI-free transmission for $BW = \frac{1}{2T_s}$, the $\text{sinc}(t/T_s)$ pulse
- ISI-free transmission is **not** possible for $BW < \frac{1}{2T_s}$

The $\text{sinc}(t/T_s)$ pulse and its spectrum



The Raised-Cosine (RC) pulse and its spectrum



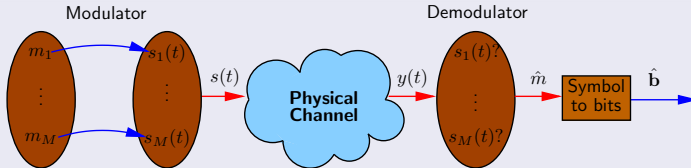
Properties of the Sampling Rx

- Very simple to implement
- Good results for good channel conditions
- Sensitive to bad channel conditions

What do we do if the channel is very noisy?

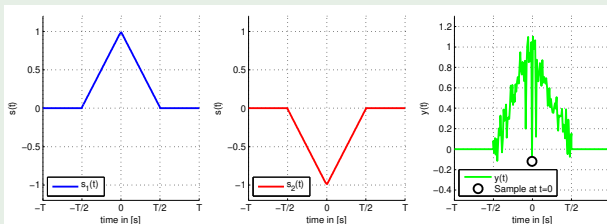
We use a better (but more complex) Rx: **The linear Rx**

The system we are analyzing...



Example (From the previous Example)

A binary message is transmitted using a triangular pulse. At the Rx, and based on $y(t)$, we need to guess which message was transmitted.

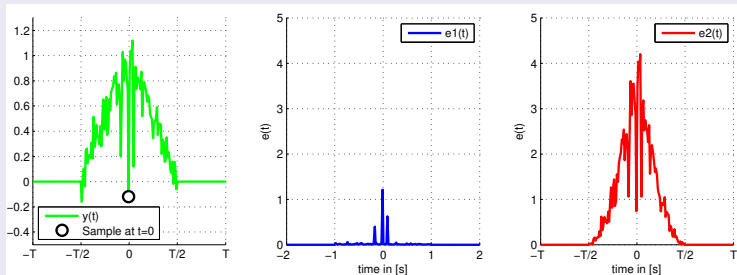


Definition (The error energy)

The error between the received signal $y(t)$ and the i th signal alternative is defined as $e_i(t) = |y(t) - s_i(t)|$, and therefore, one can define the error energy as follows

$$E_{e_i} = \int_{-\infty}^{\infty} [y(t) - s_i(t)]^2 dt.$$

For the previous example



Minimizing the error energy

- The error energies can be used for decision \Rightarrow The problem we need to solve is

$$\hat{i} = \arg \min_i \left\{ \int_{-\infty}^{\infty} [y(t) - s_i(t)]^2 dt \right\}$$

- It can be shown that

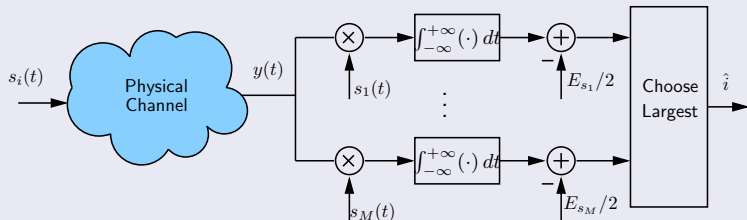
$$\min_i \left\{ \int_{-\infty}^{\infty} [y(t) - s_i(t)]^2 dt \right\} \equiv \max_i \left\{ \int_{-\infty}^{\infty} y(t)s_i(t) dt - \frac{E_{s_i}}{2} \right\},$$

where E_{s_i} is the energy of the i th signal alternative,
 $E_{s_i} = \int_{-\infty}^{\infty} s_i^2(t) dt$.

- The problem we need to solve is then

$$\hat{i} = \arg \max_i \left\{ \int_{-\infty}^{\infty} y(t)s_i(t) dt - \frac{E_{s_i}}{2} \right\}$$

The Correlator Rx



Some Comments

- If all the signal alternatives have the same energy ($E_{s_1} = E_{s_2} = \dots$), the problem is reduced to

$$\hat{i} = \arg \max_i \left\{ \int_{-\infty}^{\infty} y(t) s_i(t) dt \right\}$$

- Since the “heart” of the processing are correlation integrals, this Rx is called the **correlator Rx** (CR)

The Correlation Integral

- It can be proved that

$$\int_{-\infty}^{\infty} y(t)s_i(t) dt = y(t) * h_i(t) \Big|_{t=0},$$

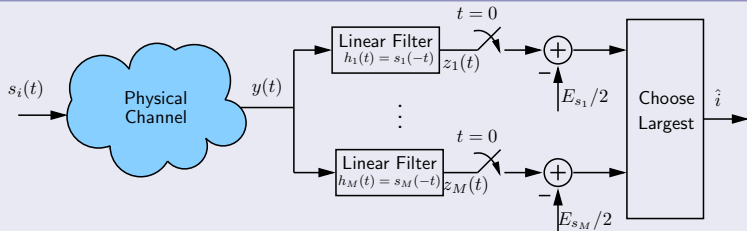
where $h_i(t) = s_i(-t)$, or in frequency domain $H(f) = S^*(f)$.

- **Proof?**
- The previous result implies that the correlation integrals in the previous block diagram can be replaced by linear filters and sampling at $t = 0$.

The Causality Problem

For a linear filter to be realizable, it must be causal, i.e., if $s_i(t) \neq 0$ for $-T_1 \leq t \leq T_2$ ($T_1, T_2 > 0$), the filter with impulse response $h_i(t) = s_i(-t)$ is not causal (not realizable). To have a causal filter the impulse response must be shifted T_2 [s], i.e., $h_i(t) = s_i(-t + T_2)$. The corresponding samples will be delayed in T_2 [s].

The Matched Filter Rx



Some Comments

- Since the filters are *matched* to the signal alternatives, it is called the **matched filter Rx** (MFR)
- Both the **correlator Rx** (CR) and MFR can be implemented using *linear filters* and are therefore called **linear Rx** (LR)
- In terms of performance, the CR and the MFR are **equivalent**
- However, the performance of the CR and MFR is different than the performance of **Sampling Rx**

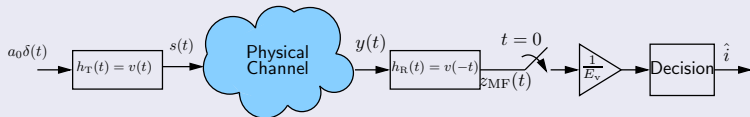
One M -ary Pulse

The transmitted signal is

$$s(t) = a_0 v(t)$$

where $a_0 \in \mathcal{A}$ is the amplitude of the transmitted symbol, and E_v is the energy of the pulse $v(t)$.

The MF Rx for this case?



Example (MF Rx for 4-ary transmission)

Assume $\mathcal{A} = \{-3, -1, 1, 3\}$ and a channel which is perfect.

Why do we need the gain $1/E_v$?

A sequence of M -ary pulses

The transmitted signal is

$$\begin{aligned}s(t) &= \sum_{k=0}^{\infty} a_k v(t - kT_s) \\ &= a_0 v(t) + a_1 v(t - T_s) + a_2 v(t - 2T_s) + \dots,\end{aligned}$$

where $a_k \in \mathcal{A}$ represent the amplitude of the symbol transmitted at the k th time instant.

The MF Rx for this case?

HW for L₃!!

Today's Summary

- Sampling Rx (one pulse and a train of pulses)
- Nyquist pulses. No ISI for Sampling Rx
- Rate and BW tradeoff
- Sinc pulses and RC pulses
- MFR/CR (one pulse and a train of pulses)
- Scaling by $1/E_v$