

# TSKS01 Digital Communication

## Lecture 4

Digital Modulation – Detection and Standard Constellations

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## Last Time – Digital Modulation

Signals:  $s_i(t) = \sum_{j=0}^{N-1} s_{i,j} \phi_j(t)$ ,  $i = 0, 1, \dots, M-1$ ,  $0 \leq t < T$

↑  
ON basis

AWGN:  $W(t) = \sum_{j=0}^{N-1} W_j \phi_j(t) + W'(t)$

↑  
Irrelevant

$W_j = (W, \phi_j)$   
Gaussian with mean 0.

Received:  $X(t) = s_i(t) + W(t) = \sum_{j=0}^{N-1} X_j \phi_j(t) + W'(t)$

$X_j = (X, \phi_j)$   
Gaussian with mean  $s_{i,j}$ .

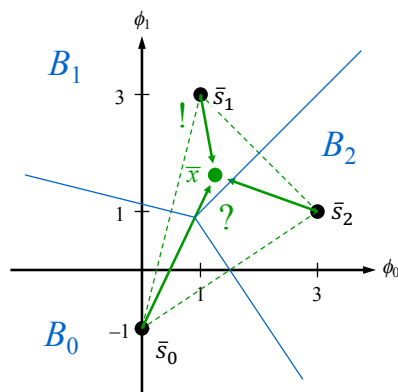
Vectors:  $\bar{X} = \bar{s}_i + \bar{W}$

$$\begin{pmatrix} X_0 \\ \vdots \\ X_{N-1} \end{pmatrix} = \begin{pmatrix} s_{i,0} \\ \vdots \\ s_{i,N-1} \end{pmatrix} + \begin{pmatrix} W_0 \\ \vdots \\ W_{N-1} \end{pmatrix}$$

Orthogonal noise components  
are statistically independent.

$$\sigma_{W_i}^2 = \sigma_{X_i}^2 = R_W(f) = N_0/2$$

## ML Decision Regions



Interpret  $\bar{x}$  as the nearest signal.

Result:

Decision regions consist of all points closest to a signal point.

Notation:

$B_i$  is the decision region of the signal vector  $\bar{s}_i$ .  
Thus also of the signal  $s_i(t)$  and of the message  $a_i$ .

Borders are orthogonal to straight lines between signals:  
In 2 dimensions: Lines.

In 3 dimensions: Planes.

Higher dim: Hyperplanes.

Borders cut the lines mid-way.

## Error Probability

Symbol error probability:

$$P_e = \Pr\{\hat{A} \neq A\} = \sum_{i=0}^{M-1} \Pr\{A = a_i\} \cdot \Pr\{\hat{A} \neq a_i | A = a_i\}$$

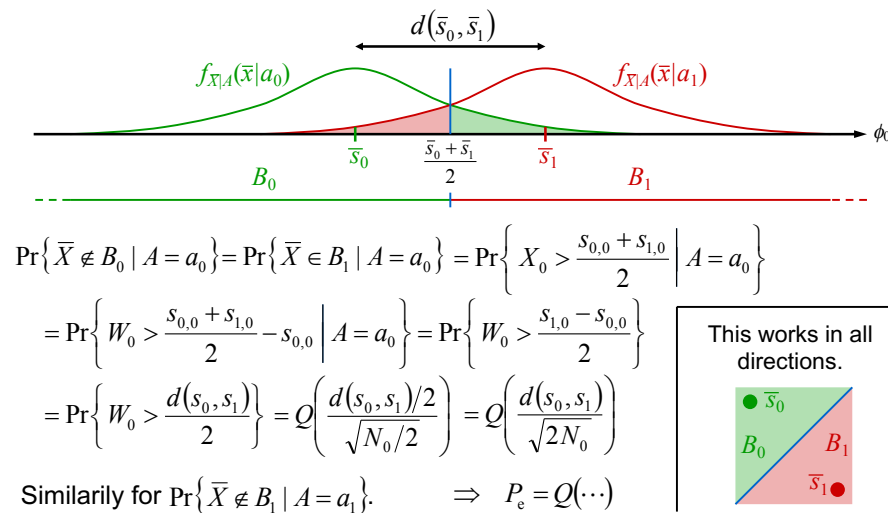
$$= \sum_{i=0}^{M-1} \Pr\{A = a_i\} \cdot \Pr\{\bar{X} \notin B_i | A = a_i\}$$

ML detection:  $\Pr\{A = a_i\} = \frac{1}{M}$  for  $i = 0, 1, \dots, M-1$ :

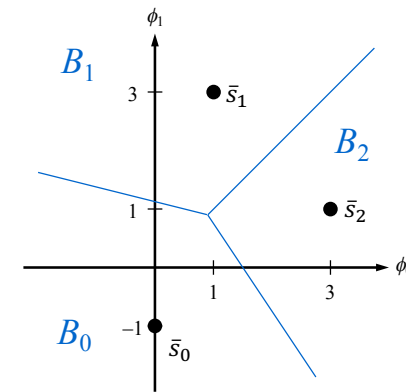
$$P_e = \frac{1}{M} \sum_{i=0}^{M-1} \Pr\{\bar{X} \notin B_i | A = a_i\} = \frac{1}{M} \sum_{i=0}^{M-1} \int \cdots \int_{\bar{x} \notin B_i} f_{\bar{X}|A}(\bar{x} | a_i) dx_0 \cdots dx_{N-1}$$

This is generally hard to calculate!

## Special Case: Two signals in $N = 1$ Dimension



## Back to $M$ Signals in $N$ Dimensions



We had:

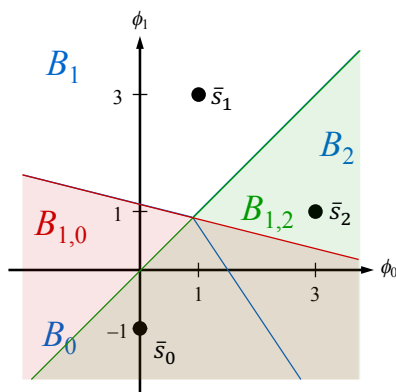
$$P_e = \frac{1}{M} \sum_{i=0}^{M-1} \Pr\{\bar{X} \notin B_i | A = a_i\}$$

$$= \frac{1}{M} \sum_{i=0}^{M-1} \sum_{j \neq i} \Pr\{\bar{X} \in B_j | A = a_i\}$$

**Hard to calculate!**  
Could we take this to the simpler case in one dimension?

Interpret  $\bar{x}$  as the nearest signal.

## The Union Bound



Interpret  $\bar{x}$  as the nearest signal.

**Upper bound** by overestimating the decision regions.

We had:

$$P_e = \frac{1}{M} \sum_{i=0}^{M-1} \sum_{j \neq i} \Pr\{\bar{X} \in B_j | A = a_i\}$$

Define overestimated regions:

$$B_{i,j} = \{\bar{x} : d(\bar{x}, \bar{s}_j) < d(\bar{x}, \bar{s}_i)\}$$

Overestimated error probability:

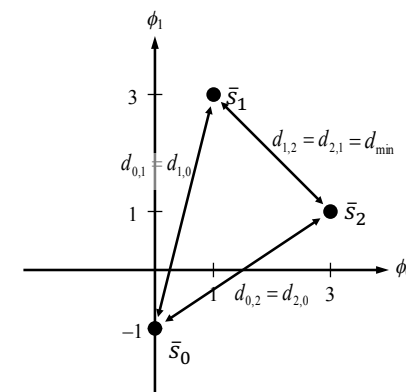
$$P_e \leq \frac{1}{M} \sum_{i=0}^{M-1} \sum_{j \neq i} \Pr\{\bar{X} \in B_{i,j} | A = a_i\}$$

In the one-dimensional case:

$$P_e \leq \frac{1}{M} \sum_{i=0}^{M-1} \sum_{j \neq i} Q\left(\frac{d_{i,j}}{\sqrt{2N_0}}\right)$$

Distances:  $d_{i,j} = d(\bar{s}_i, \bar{s}_j)$

## The Nearest Neighbour Approximation



Interpret  $\bar{x}$  as the nearest signal.

We had the union bound:

$$P_e \leq \frac{1}{M} \sum_{i=0}^{M-1} \sum_{j \neq i} Q\left(\frac{d_{i,j}}{\sqrt{2N_0}}\right)$$

Dominated by the smallest distance:

$$d_{\min} = \min_{i \neq j} d_{i,j}$$

$$n_i = \# j : d_{i,j} = d_{\min}$$

Nearest neighbour approximation:

$$P_e \approx \frac{1}{M} \sum_{i=0}^{M-1} \sum_{j: d_{i,j} = d_{\min}} Q\left(\frac{d_{\min}}{\sqrt{2N_0}}\right)$$

$$= \frac{1}{M} \sum_{i=0}^{M-1} n_i Q\left(\frac{d_{\min}}{\sqrt{2N_0}}\right)$$

## Comments

$$E_{\text{avg}} = \frac{1}{M} \sum_{i=0}^{M-1} E_i, \quad E_i = \int_{-\infty}^{\infty} s_i^2(t) dt < \infty$$

At high SNR  $E_{\text{avg}}/N_0$ :

Both the union bound on (and the nearest neighbour of) the error probability are close to the real error probability.

Alternative upper bound:

As the union bound, but only consider pairs of points whose decision regions share a common border.

Alternative approximation:

As the nearest neighbour approximation, but consider the two or three smallest distances.

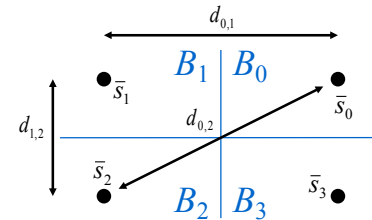
Very simple approximation:

$$P_e \approx Q\left(\frac{d_{\min}}{\sqrt{2N_0}}\right)$$

Very simple upper bound:

$$P_e \leq (M-1) \cdot Q\left(\frac{d_{\min}}{\sqrt{2N_0}}\right)$$

## Special Case: Orthogonal Decision Borders



Define notation:  $q_{i,j} = Q\left(\frac{d_{i,j}}{\sqrt{2N_0}}\right)$

Pythagoras:  $d_{0,2}^2 = d_{0,1}^2 + d_{1,2}^2$

UB:  $P_e \leq q_{0,1} + q_{1,2} + q_{0,2}$

Alternative bound:  $P_e \leq q_{0,1} + q_{1,2}$

NN:  $P_e \approx q_{1,2}$

Alternative approx.:  $P_e \approx q_{0,1} + q_{1,2}$

Exact:  $P_e = q_{0,1} + q_{1,2} - q_{0,1}q_{1,2}$   
(orthogonal noise components are independent)

Lower bound:  $P_e > q_{1,2}$

## Designing a Constellation

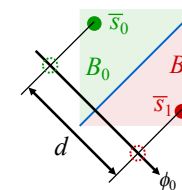
- Practical limitations
  - Energy per symbol – average or maximum
  - Energy per bit – average or maximum

- Recall definitions:

$$E_i = \int_0^T s_i^2(t) dt = \|\bar{s}_i\|^2$$

- Average per symbol:  $E_{\text{avg}} = \frac{1}{M} \sum_{i=0}^{M-1} E_i$
- Maximum per symbol:  $E_{\text{max}} = \max_{0 \leq i \leq M-1} E_i$
- Per bit: Divide  $E$  by  $\log_2(M)$

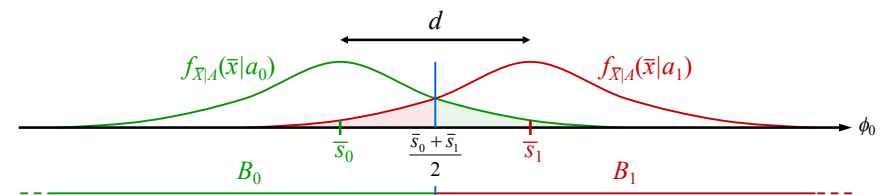
## Binary Constellations ( $M = 2$ )



$$P_e = Q\left(\frac{d}{\sqrt{2N_0}}\right)$$

$$E_{\text{avg}} = \frac{E_0 + E_1}{2} = \frac{\|\bar{s}_0\|^2 + \|\bar{s}_1\|^2}{2}$$

$$E_{\text{max}} = \max(\|\bar{s}_0\|^2, \|\bar{s}_1\|^2)$$



## On-Off Keying (OOK)

**Basis:**

$$\phi(t) = \sqrt{2/T} \cos(2\pi f_c t), \quad 0 \leq t < T$$

**Signals:**

$$s_0(t) = 0, \quad s_1(t) = A \cdot \phi(t)$$

**Energies:**

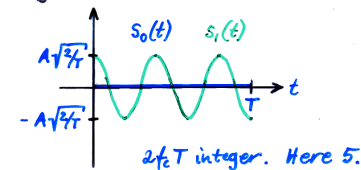
$$E_0 = 0, \quad E_1 = A^2$$

$$E_{\text{avg}} = \frac{0 + A^2}{2} = \frac{A^2}{2}, \quad E_{\text{max}} = E_1 = A^2$$

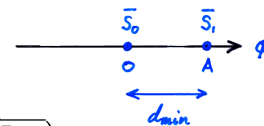
**Error probability (AWGN):**

$$P_e = Q\left(\frac{d_{\min}}{\sqrt{2N_0}}\right) = Q\left(\frac{A}{\sqrt{2N_0}}\right) = Q\left(\sqrt{\frac{E_{\text{avg}}}{N_0}}\right) = Q\left(\sqrt{\frac{E_{\text{max}}}{2N_0}}\right)$$

**Signals:**



**Vectors:**



## Binary Phase-Shift Keying (BPSK)

**Basis:**

$$\phi(t) = \sqrt{2/T} \cos(2\pi f_c t), \quad 0 \leq t < T$$

**Signals:**

$$s_0(t) = +A \cdot \phi(t)$$

$$s_1(t) = -A \cdot \phi(t)$$

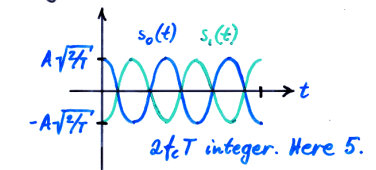
**Energies:**

$$E_0 = E_1 = E_{\text{avg}} = E_{\text{max}} = A^2$$

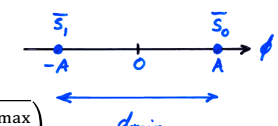
**Error probability (AWGN):**

$$P_e = Q\left(\frac{d_{\min}}{\sqrt{2N_0}}\right) = Q\left(\frac{2A}{\sqrt{2N_0}}\right) = Q\left(\sqrt{\frac{2E_{\text{avg}}}{N_0}}\right) = Q\left(\sqrt{\frac{2E_{\text{max}}}{N_0}}\right)$$

**Signals:**



**Vectors:**



## Binary Frequency-Shift Keying (BFSK)

**Basis:**

$$\phi_0(t) = \sqrt{2/T} \cos(2\pi f_0 t), \quad 0 \leq t < T$$

$$\phi_1(t) = \sqrt{2/T} \cos(2\pi f_1 t), \quad 0 \leq t < T$$

**Signals:**

$$s_0(t) = A \cdot \phi_0(t), \quad s_1(t) = A \cdot \phi_1(t)$$

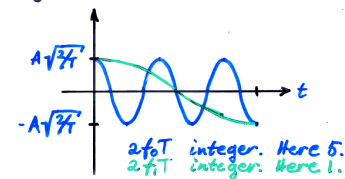
**Energies:**

$$E_0 = E_1 = E_{\text{avg}} = E_{\text{max}} = A^2$$

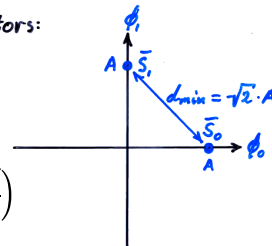
**Error probability (AWGN):**

$$P_e = Q\left(\frac{d_{\min}}{\sqrt{2N_0}}\right) = Q\left(\frac{A}{\sqrt{N_0}}\right) = Q\left(\sqrt{\frac{E_{\text{avg}}}{N_0}}\right) = Q\left(\sqrt{\frac{E_{\text{max}}}{N_0}}\right)$$

**Signals:**



**Vectors:**



## Comparison of Binary Constellations

▪ Fair comparison: Same average energy  $E_{\text{avg}}$ !

▪ Error probability

▪ Q-function is decreasing: Larger argument is better

▪ OOK:  $P_e = Q\left(\sqrt{\frac{E_{\text{avg}}}{N_0}}\right)$

▪ BPSK:  $P_e = Q\left(\sqrt{\frac{2E_{\text{avg}}}{N_0}}\right)$

Lowest error probability  
( $\times 2 = 3 \text{ dB}$  better SNR)

▪ BFSK:  $P_e = Q\left(\sqrt{\frac{E_{\text{avg}}}{N_0}}\right)$

BPSK uses phase difference (+A/-A) to maximize  $d_{\min}$  for given energy!

Drawback: Phase coherence (e.g., OOK only needs to compute  $||\vec{x}||$ )

## Non-Binary Constellations

Symbol error probability:  $P_e$

Bit error probability:  $P_b$

Maximum symbol energy:  $E_{\max}$

Average symbol energy:  $E_{\text{avg}}$

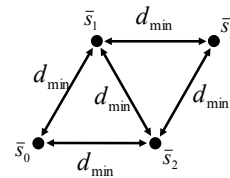
Average bit energy:  $E_b$

Number of bits:  $k$

Number of signals:  $M = 2^k$

Relation:  $E_{\text{avg}} = kE_b$

Example:



Number of nearest neighbours:

$$n_0 = n_3 = 2 \quad n_1 = n_2 = 3$$

Nearest neighbour approximation:

$$P_e \approx \frac{1}{M} \sum_{i=0}^{M-1} n_i Q\left(\frac{d_{\min}}{\sqrt{2N_0}}\right)$$

$$= \frac{1}{4}(2+3+3+2) \cdot Q(\dots) = \frac{5}{2} Q(\dots)$$

## Error Probabilities

Exact:

$$P_e = \frac{1}{M} \sum_{i=0}^{M-1} \sum_{j \neq i} \Pr\{\bar{X} \in B_j \mid A = a_i\}$$

Union bound:

$$P_e \leq \frac{1}{M} \sum_{i=0}^{M-1} \sum_{j \neq i} Q\left(\frac{d_{i,j}}{\sqrt{2N_0}}\right)$$

Nearest neighbour approx:

$$P_e \approx \frac{1}{M} \sum_{i=0}^{M-1} \underbrace{\sum_{j: d_{ij}=d_{\min}} Q\left(\frac{d_{\min}}{\sqrt{2N_0}}\right)}_{n_i Q(\dots)}$$

$$P_b = \frac{1}{M} \sum_{i=0}^{M-1} \sum_{j \neq i} \frac{N_{i,j}}{k} \Pr\{\bar{X} \in B_j \mid A = a_i\}$$

$$P_b \leq \frac{1}{M} \sum_{i=0}^{M-1} \sum_{j \neq i} \frac{N_{i,j}}{k} Q\left(\frac{d_{i,j}}{\sqrt{2N_0}}\right)$$

$$P_b \approx \frac{1}{M} \sum_{i=0}^{M-1} \sum_{j: d_{ij}=d_{\min}} \frac{N_{i,j}}{k} Q\left(\frac{d_{\min}}{\sqrt{2N_0}}\right)$$

Approx. of average fraction of bits that differ between a signal and a neighbour, counting only nearest neighbours.

Number of bit positions where  $s_i$  and  $s_j$  differ.

## Gray Codes

Subsequent patterns differ in only one bit position.

Makes nearest neighbours differ in only one bit position:  $P_b \approx P_e/k$ .

One bit

0  
1

Two bits

0 0 } One bit  
0 1 }  
1 1 } One bit  
1 0 } reflected

Three bits

0 0 0 }  
0 0 1 } Two bits  
0 1 1 }  
0 1 0 }  
1 1 0 }  
1 1 1 } Two bits  
1 0 1 } reflected  
1 0 0 }

In general,  $n$  bits:

0 0 0 ... 0 }  
: : : : n-1 bits  
0 1 0 ... 0 }  
: : : : n-1 bits  
1 1 0 ... 0 } reflected  
: : : : reflected  
1 0 0 ... 0 }

## Amplitude-Shift Keying (ASK)

Basis:

$$\phi(t) = \sqrt{2/T} \cos(2\pi f_c t), \quad 0 \leq t < T$$

Signals ( $M = 4$ ):

$$s_1(t) = -s_2(t) = A \cdot \phi(t)$$

$$s_0(t) = -s_3(t) = 3A \cdot \phi(t)$$

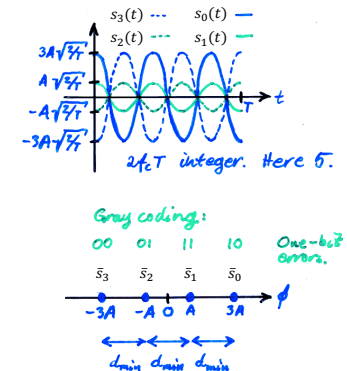
Energies:

$$E_{\text{avg}} = 5A^2, \quad E_{\max} = 9A^2, \quad E_b = 5A^2/2$$

Error probability (AWGN):

$$P_e = \frac{1+2+2+1}{4} Q\left(\frac{d_{\min}}{\sqrt{2N_0}}\right) = \frac{3}{2} Q\left(\frac{2A}{\sqrt{2N_0}}\right) = \frac{3}{2} Q\left(\sqrt{\frac{2E_{\text{avg}}}{5N_0}}\right) = \frac{3}{2} Q\left(\sqrt{\frac{2E_{\max}}{9N_0}}\right)$$

$$P_b \approx \frac{P_e}{2} = \frac{3}{4} Q\left(\sqrt{\frac{2E_{\text{avg}}}{5N_0}}\right) = \frac{3}{4} Q\left(\sqrt{\frac{4E_b}{5N_0}}\right)$$



## Quadruple Phase-Shift Keying (QPSK/4-PSK)

**Basis:**

$$\phi_0(t) = \sqrt{2/T} \cos(2\pi f_c t), \quad 0 \leq t < T$$

$$\phi_1(t) = \sqrt{2/T} \sin(2\pi f_c t), \quad 0 \leq t < T$$

**Signals** ( $i = 0, 1, 2, 3$ ):

$$s_i(t) = A \cos\left(\frac{(2i+1)\pi}{4}\right) \phi_0(t) - A \sin\left(\frac{(2i+1)\pi}{4}\right) \phi_1(t)$$

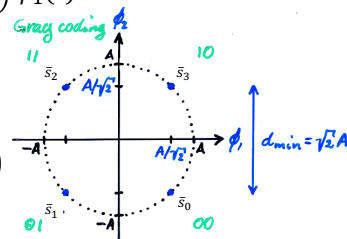
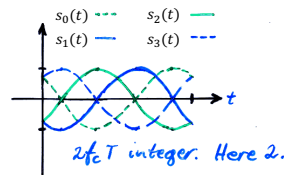
**Energies:**

$$E_i = E_{\text{avg}} = E_{\text{max}} = A^2, \quad E_b = A^2/2$$

**Error probability (AWGN):**

$$P_e \approx 2Q\left(\frac{d_{\min}}{\sqrt{2N_0}}\right) = 2Q\left(\frac{A}{\sqrt{N_0}}\right) = 2Q\left(\sqrt{\frac{E_{\text{avg}}}{N_0}}\right) = 2Q\left(\sqrt{\frac{E_{\text{max}}}{N_0}}\right)$$

$$P_b \approx \frac{P_e}{2} \approx Q\left(\sqrt{\frac{E_{\text{avg}}}{N_0}}\right) = Q\left(\sqrt{\frac{2E_b}{N_0}}\right)$$



## M Phase-Shift Keying (M-PSK)

**Example: 8-PSK**

**Basis:**

$$\phi_0(t) = \sqrt{2/T} \cos(2\pi f_c t), \quad 0 \leq t < T$$

$$\phi_1(t) = \sqrt{2/T} \sin(2\pi f_c t), \quad 0 \leq t < T$$

**Signals** ( $i = 0, 1, 2, 3$ ):

$$s_i(t) = A \cos\left(\frac{(2i+1)\pi}{M}\right) \phi_0(t) - A \sin\left(\frac{(2i+1)\pi}{M}\right) \phi_1(t)$$

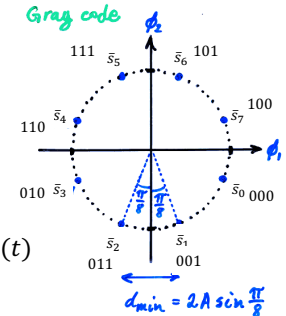
**Energies:**

$$E_i = E_{\text{avg}} = E_{\text{max}} = A^2, \quad E_b = A^2/k$$

**Error probability (AWGN):**

$$P_e \approx 2Q\left(\frac{d_{\min}}{\sqrt{2N_0}}\right) = 2Q\left(\frac{\sqrt{2}A}{\sqrt{N_0}} \sin\left(\frac{\pi}{M}\right)\right) = 2Q\left(\sqrt{\frac{2E_{\text{avg}}}{N_0}} \sin\left(\frac{\pi}{M}\right)\right) = 2Q\left(\sqrt{\frac{2E_{\text{max}}}{N_0}} \sin\left(\frac{\pi}{M}\right)\right)$$

$$P_b \approx \frac{P_e}{k} \approx \frac{2}{k} Q\left(\sqrt{\frac{2E_{\text{avg}}}{N_0}} \sin\left(\frac{\pi}{M}\right)\right) = \frac{2}{k} Q\left(\sqrt{\frac{2kE_b}{N_0}} \sin\left(\frac{\pi}{M}\right)\right)$$



## 16-QAM (Quadrature Amplitude Modulation)

**Basis:**

$$\phi_0(t) = \sqrt{2/T} \cos(2\pi f_c t), \quad 0 \leq t < T$$

$$\phi_1(t) = \sqrt{2/T} \sin(2\pi f_c t), \quad 0 \leq t < T$$

**Signals:**

$$s_{i,j}(t) = s_i \phi_0(t) + s_j \phi_1(t), \quad s_i, s_j \in \{\pm A, \pm 3A\}$$

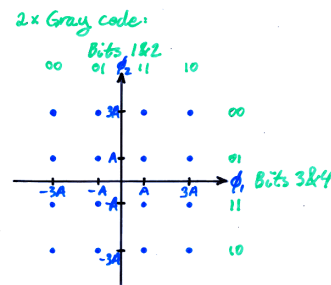
**Energies:**

$$E_i = s_i^2 + s_j^2, \quad E_{\text{avg}} = 10A^2, \quad E_{\text{max}} = 18A^2, \quad E_b = E_{\text{avg}}/4$$

**Error probability (AWGN):**

$$P_e \approx \frac{4 \cdot 4 + 8 \cdot 3 + 4 \cdot 2}{16} Q\left(\frac{2A}{\sqrt{2N_0}}\right) = 3Q\left(\sqrt{\frac{E_{\text{avg}}}{5N_0}}\right) = 3Q\left(\sqrt{\frac{E_{\text{max}}}{9N_0}}\right)$$

$$P_b \approx \frac{P_e}{4} \approx \frac{3}{4} Q\left(\sqrt{\frac{E_{\text{avg}}}{5N_0}}\right) = \frac{3}{4} Q\left(\sqrt{\frac{4E_b}{5N_0}}\right)$$



## Frequency-Shift Keying (FSK)

**Basis:**

$$\phi_i(t) = \sqrt{2/T} \cos\left(2\pi\left(f_c + \frac{i}{2T}\right)t\right), \quad 0 \leq t < T$$

$$i = \{0, 1, \dots, M-1\}$$

**Signals:**

$$s_i(t) = A \cdot \phi_i(t)$$

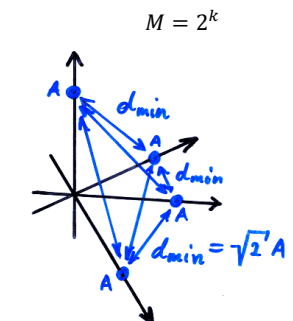
**Energies:**

$$E_i = E_{\text{avg}} = E_{\text{max}} = A^2, \quad E_b = A^2/k$$

**Error probability (AWGN):**

$$P_e \approx (M-1)Q\left(\frac{d_{\min}}{\sqrt{2N_0}}\right) = (M-1)Q\left(\sqrt{\frac{E_{\text{avg}}}{N_0}}\right) = (2^k-1)Q\left(\sqrt{\frac{E_{\text{avg}}}{N_0}}\right)$$

$$P_b \approx 2^{k-1}Q\left(\sqrt{\frac{E_{\text{avg}}}{N_0}}\right) = 2^{k-1}Q\left(\sqrt{\frac{kE_b}{N_0}}\right)$$





# Orthogonal Frequency Division Multiplex (OFDM)

- Principle for  $N$  dimensional signal
  - Use many 2-dimensional modulations (e.g., PSK/QAM)
  - Use  $N/2$  different frequencies:

$$f_k = f_0 + \frac{k}{T}, \quad 0 \leq k < N/2,$$

for base frequency  $f_0$  and  $2f_0T$  being an integer

- OFDM signal generation

$$s(t) = \sum_{k=0}^{N/2-1} (s_{2k} \cos(2\pi f_k t) - s_{2k+1} \sin(2\pi f_k t)) I_{\{0 \leq t < T\}}(t)$$

for  $0 \leq t < T$ .

# OFDM continued

Alternative representation

$$s(t) = \sum_{k=0}^{N/2-1} (s_{2k} \cos(2\pi f_k t) - s_{2k+1} \sin(2\pi f_k t)) I_{\{0 \leq t < T\}}(t)$$

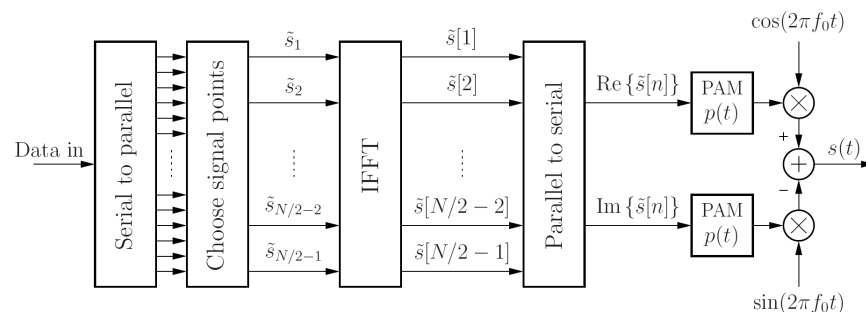
$$= \sum_{k=0}^{N/2-1} \operatorname{Re} \left\{ \tilde{s}_k e^{\frac{j2\pi k}{T} t} e^{j2\pi f_0 t} \right\} I_{\{0 \leq t < T\}}(t) \quad \text{with } \tilde{s}_k = s_{2k} + js_{2k+1}$$

Complex baseband representation:

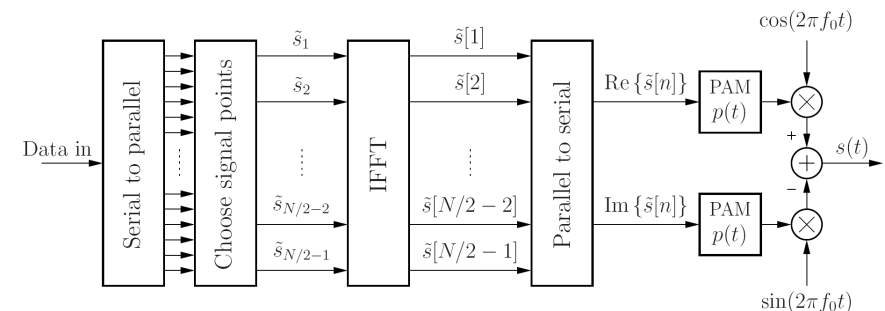
- $\tilde{s}(t) = \sum_{k=0}^{N/2-1} \tilde{s}_k e^{\frac{j2\pi k}{T} t}$  for  $0 \leq t < T$
- Sampled:  $\tilde{s}[n] = \tilde{s}\left(\frac{nT}{N/2}\right)$  for  $0 \leq n < N/2$
- $\tilde{s}[n]$  is the IDFT of  $\tilde{s}_k$  for  $0 \leq k < N/2$

## Generating an OFDM Signal

- Generate  $\tilde{s}(t)$  by PAM of  $\tilde{s}[n]$  (using sinc pulse function)
  - Generate  $s(t)$  from  $\tilde{s}(t)$  as
- $$s(t) = \operatorname{Re}\{\tilde{s}(t)e^{j2\pi f_0 t}\} = \operatorname{Re}\{\tilde{s}(t)\} \cos(2\pi f_0 t) - \operatorname{Im}\{\tilde{s}(t)\} \sin(2\pi f_0 t)$$



## Detection of an OFDM Signal



Detector – Opposite process

- Sample
- FFT
- ML decisions for each frequency separately

## Example: LTE (4G)

- Bandwidth: 10 MHz
  - OFDM with  $N = 1202$
  - $N/2$  subcarriers with 2-dimensional constellations
  - Constellations: BPSK, QPSK, 16-QAM, 64-QAM
  - Effective bandwidth: 9 MHz
- Data rate (no errors)
  - BPSK:  $9 \cdot 10^6 = 9 \text{ Mbit/s}$
  - QPSK:  $2 \cdot 9 \cdot 10^6 = 18 \text{ Mbit/s}$
  - 16-QAM:  $4 \cdot 9 \cdot 10^6 = 36 \text{ Mbit/s}$
  - 64-QAM:  $6 \cdot 9 \cdot 10^6 = 54 \text{ Mbit/s}$

