

Digital Communications

SSY125, Lecture 7

Analysis of Linear Modulations (Chapter 6)

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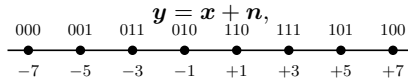


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Analysis of Linear Modulations



AWGN channel,



- $x_i \in \mathcal{X} = \{X_1, X_2, \dots, X_M\} \subset \mathbb{C}$, with **average energy per symbol** E_s .
- $|\mathcal{X}| = M \rightarrow m = \log M$ bits per symbol.
- $N_i \sim \mathcal{CN}(0, N_0/2)$.
- The **energy per information bit** is

$$E_b = \frac{E_s}{\log M}.$$

Optimum Decoding Rule



- Based on y , infer u optimally \rightarrow Based on y infer x optimally.
- Criterion: minimizing the probability of error \rightarrow maximum a posteriori (MAP) decoding rule,

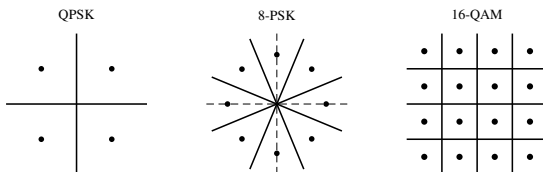
$$\hat{x}_{\text{MAP}} = \arg \max_x p(x|y).$$



- If all sequences x are equiprobable, the MAP decoding rule boils down to the maximum likelihood (ML) decoding rule,

$$\hat{x}_{\text{ML}} = \arg \max_x p(y|x),$$

Optimum Decoding Rule



- If the channel is **memoryless**: **symbol-by-symbol** decoder that makes an independent decision on each received symbol y_i ,

$$\hat{x} = \arg \max_{x \in \mathcal{X}} p(y|x).$$

- For transmission over the Gaussian channel (and equiprobable symbols),

$$\begin{aligned}\hat{x} &= \arg \max_{x \in \mathcal{X}} p(y|x) \\ &= \arg \min_{x \in \mathcal{X}} \|y - x\|^2,\end{aligned}$$

i.e., selecting among all possible constellation symbols $x \in \mathcal{X}$ the one at **minimum Euclidean distance** to the received value y .

Performance of Modulation Schemes

The evaluation of a modulation scheme is based on three parameters:

- The **error probability** (symbol error probability or bit error probability);
- The E_b/N_0 required to achieve such probability of error, and
- The **spectral efficiency** $R = \log M$.

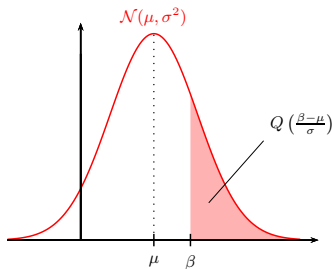
The Union Bound

Given a number of events E_1, \dots, E_N

$$\Pr\left(\bigcup_{i=1}^n E_i\right) \leq \sum_{i=1}^n \Pr(E_i),$$

where $E_i \cup E_j$ stands for the union of the events E_i and E_j .

The Q-function



For $Z \sim \mathcal{N}(\mu, \sigma^2)$,

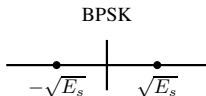
$$\Pr(Z > \beta) = Q\left(\frac{\beta - \mu}{\sigma}\right),$$

where $Q(x) = \frac{1}{\sqrt{2\pi}} \int_x^\infty \exp(-\tau^2/2) d\tau$ is the tail probability of the standard normal distribution.

- $Q(-x) = 1 - Q(x)$.
- For $0 < a \ll b < c$,

$$Q(a) + Q(b) + Q(c) \approx Q(a).$$

Symbol Error Probability of BPSK



- $\mathcal{X} = \{X_1, X_2\} \subset \mathbb{R}$, where $X_1 = -\sqrt{E_s}$ and $X_2 = \sqrt{E_s}$.
- The received signal is

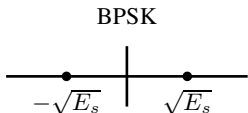
$$y = x + n,$$

where $n \sim \mathcal{N}(0, \sigma^2)$.

- For **equiprobable symbols**, $\hat{x} = X_2$ if $y > 0$ and $\hat{x} = X_1$ otherwise.

$$\begin{aligned} P_s^{\text{BPSK}} &= \sum_{x \in \mathcal{X}} \Pr(\hat{x} \neq x | x) P(x) \\ &= \Pr(X_2 | X_1) P(X_1) + \Pr(X_1 | X_2) P(X_2) \\ &= \frac{1}{2} \Pr(X_2 | X_1) + \frac{1}{2} \Pr(X_1 | X_2), \end{aligned}$$

Symbol Error Probability of BPSK



If $X = X_1$, $Y \sim \mathcal{N}(-\sqrt{E_s}, \sigma^2)$, hence

$$\begin{aligned}\Pr(\hat{x} = X_2 | x = X_1) &= \Pr(Y > 0 | X = X_1) \\ &= \Pr(Y > 0) \big|_{Y \sim \mathcal{N}(-\sqrt{E_s}, \sigma^2)} \\ &= Q\left(\frac{\sqrt{E_s}}{\sigma}\right) = Q\left(\sqrt{\frac{2E_s}{N_0}}\right) = Q\left(\sqrt{\frac{2E_b}{N_0}}\right).\end{aligned}$$

Due to symmetry, $\Pr(\hat{x} = X_1 | x = X_2) = p(\hat{x} = X_2 | x = X_1)$, and

$$P_s^{\text{BPSK}} = Q\left(\sqrt{\frac{2E_s}{N_0}}\right) = Q\left(\sqrt{\frac{2E_b}{N_0}}\right).$$

Symbol Error Probability and Euclidean Distance

For BPSK, $d_E(X_1, X_2) = 2\sqrt{E_s}$, hence

$$P_s^{\text{BPSK}} = Q\left(\sqrt{\frac{d_E^2(X_1, X_2)}{2N_0}}\right).$$

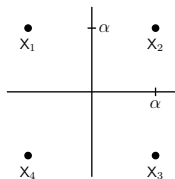
In the general case...

- Consider $X_1 \in \mathbb{C}$ and $X_2 \in \mathbb{C}$ with $d_E(X_1, X_2) = \|X_1 - X_2\|$. Assume $x \in \mathcal{X} = \{X_1, X_2, \dots\}$ is transmitted and we receive $y = x + n$.
- Let \tilde{y} be the projection of y onto the straight line between X_1 and X_2 . Then,

$$\begin{aligned}\Pr(\hat{x} = X_2 | x = X_1) &= \Pr\left(\tilde{Y} > \frac{d_E(X_1, X_2)}{2}\right) \Big|_{\tilde{Y} \sim \mathcal{N}(0, \sigma^2)} \\ &= Q\left(\frac{d_E(X_1, X_2)}{2\sigma}\right) = Q\left(\sqrt{\frac{d_E^2(X_1, X_2)}{2N_0}}\right).\end{aligned}$$

- $\Pr(\hat{x} = X_2 | x = X_1)$ depends on $d_E(X_1, X_2) \Rightarrow$ **Construct constellations with high distance between constellation points!**

Symbol Error Probability of 4-QAM

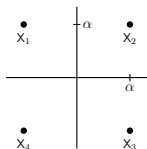


- $\mathcal{X} = \{X_1, X_2, X_3, X_4\}$ with $\alpha = \sqrt{E_s/2}$ such that the average energy per symbol is E_s .
- The ML decision is

$$\hat{x} = \begin{cases} X_1 & \text{if } y_I < 0 \text{ and } y_Q > 0 \\ X_2 & \text{if } y_I > 0 \text{ and } y_Q > 0 \\ X_3 & \text{if } y_I > 0 \text{ and } y_Q < 0 \\ X_4 & \text{if } y_I < 0 \text{ and } y_Q < 0 \end{cases},$$

where y_I and y_Q are the imaginary and quadrature components of y .

Analysis of Linear Modulations



$$\begin{aligned}P_s^{4\text{QAM}} &= \Pr(\hat{x} \neq X_1 | x = X_1) = \Pr(Y_I > 0 \cup Y_Q < 0 | X_1) \\&= \Pr(Y_I > 0 | X_1) + \Pr(Y_Q < 0 | X_1) - \Pr(Y_I > 0 \cap Y_Q < 0 | X_1) \\&= \Pr(Y_I > 0 | X_1) + \Pr(Y_Q < 0 | X_1) - \Pr(Y_I > 0 | X_1) \Pr(Y_Q < 0 | X_1),\end{aligned}$$

If X_1 is **transmitted**, $Y_I \sim \mathcal{N}(-\alpha, \sigma^2)$ and $Y_Q \sim \mathcal{N}(\alpha, \sigma^2)$. Thus,

$$\Pr(Y_I > 0 | X_1) = \Pr(Y_I > 0) \Big|_{Y_I \sim \mathcal{N}(-\alpha, \sigma^2)} = Q\left(\frac{\alpha}{\sigma}\right) = Q\left(\sqrt{\frac{2E_b}{N_0}}\right),$$

By symmetry, $\Pr(Y_Q < 0 | X_1) = \Pr(Y_I > 0 | X_1)$, and

$$P_s^{4\text{QAM}} = 2Q\left(\sqrt{\frac{2E_b}{N_0}}\right) - \left(Q\left(\sqrt{\frac{2E_b}{N_0}}\right)\right)^2.$$

Upper Bound on the Symbol Error Probability

For **general constellations**, the exact symbol error probability P_s is **hard to compute** \rightarrow Use union bound to compute an **upper bound**

$$\begin{aligned} P_s^{(M)} &= \sum_{i=1}^M \Pr(\hat{x} \neq X_i | x = X_i) P(X_i) \\ &= \frac{1}{M} \sum_{i=1}^M \Pr(\hat{x} \neq X_i | x = X_i) \\ &= \frac{1}{M} \sum_{i=1}^M \Pr\left(\bigcup_{j \neq i} \hat{x} = X_j | x = X_i\right) \\ &\leq \frac{1}{M} \sum_{i=1}^M \sum_{j \neq i} \Pr(\hat{x} = X_j | x = X_i). \end{aligned}$$

Upper Bound on the Symbol Error Probability

Using

$$\Pr(\hat{x} = X_j | x = X_i) = Q\left(\sqrt{\frac{d_E^2(X_i, X_j)}{2N_0}}\right).$$

the symbol error probability of an M -ary constellation can be upperbounded as

$$P_s^{(M)} \leq \frac{1}{M} \sum_{i=1}^M \sum_{j \neq i} Q\left(\sqrt{\frac{d_E^2(X_i, X_j)}{2N_0}}\right).$$

Upper bound on P_s for 4-QAM

Example: 4-QAM

$$\begin{aligned} P_s^{4\text{QAM}} &\leq \frac{1}{4} \sum_{i=1}^4 \sum_{j \neq i} Q \left(\sqrt{\frac{d_E^2(\mathbf{X}_i, \mathbf{X}_j)}{2N_0}} \right) = \sum_{j=2}^4 Q \left(\sqrt{\frac{d_E^2(\mathbf{X}_1, \mathbf{X}_j)}{2N_0}} \right) \\ &= Q \left(\sqrt{\frac{d_E^2(\mathbf{X}_1, \mathbf{X}_2)}{2N_0}} \right) + Q \left(\sqrt{\frac{d_E^2(\mathbf{X}_1, \mathbf{X}_3)}{2N_0}} \right) + Q \left(\sqrt{\frac{d_E^2(\mathbf{X}_1, \mathbf{X}_4)}{2N_0}} \right). \end{aligned}$$

Using $d_E^2(\mathbf{X}_1, \mathbf{X}_2) = d_E^2(\mathbf{X}_1, \mathbf{X}_4) = \|\alpha(-1+j) - \alpha(1+j)\|^2 = 4\alpha^2 = 2E_s$ and $d_E^2(\mathbf{X}_1, \mathbf{X}_3) = \|\alpha(-1+j) - \alpha(1-j)\|^2 = 8\alpha^2 = 4E_s$,

$$\begin{aligned} P_s^{4\text{QAM}} &\leq 2Q \left(\sqrt{\frac{E_s}{N_0}} \right) + Q \left(\sqrt{\frac{2E_s}{N_0}} \right) \\ &= 2Q \left(\sqrt{\frac{2E_b}{N_0}} \right) + Q \left(\sqrt{\frac{4E_b}{N_0}} \right). \end{aligned}$$

Nearest Neighbor Approximation

- For large M , even the union bound may be **difficult** to compute \rightarrow Derive an **approximation**.

$$P_s^{(M)} \leq \frac{1}{M} \sum_{i=1}^M \sum_{j \neq i} Q \left(\sqrt{\frac{d_E^2(X_i, X_j)}{2N_0}} \right).$$

- The dominant terms are the ones with smaller $d_E(X_i, X_j)$
- Let

$$d_{E,\min}(X_i) = \min_{X_j \neq X_i} d_E(X_i, X_j)$$

and $A_{\min}(X_i)$ the number of constellation points at distance $d_{E,\min}(X_i)$ from X_i (**nearest neighbors** of X_i). Then,

$$P_s^{(M)} \approx \frac{1}{M} \sum_{i=1}^M A_{\min}(X_i) Q \left(\sqrt{\frac{d_{E,\min}^2(X_i)}{2N_0}} \right).$$

Nearest Neighbor Approximation

- Further approximate P_s considering only the **minimum Euclidean distance**,

$$d_{E,\min} = \min_{X_i} d_{E,\min}(X_i).$$

Then,

$$\begin{aligned} P_s^{(M)} &\approx \frac{1}{M} \sum_{i=1}^M A_{\min}(X_i) Q\left(\sqrt{\frac{d_{E,\min}^2(X_i)}{2N_0}}\right) \\ &\leq \frac{1}{M} \sum_{i=1}^M A_{\min}(X_i) Q\left(\sqrt{\frac{d_{E,\min}^2}{2N_0}}\right) \\ &= \left(\frac{1}{M} \sum_{i=1}^M A_{\min}(X_i)\right) Q\left(\sqrt{\frac{d_{E,\min}^2}{2N_0}}\right) = \bar{A}_{\min} Q\left(\sqrt{\frac{d_{E,\min}^2}{2N_0}}\right), \end{aligned}$$

where $\bar{A}_{\min} = \frac{1}{M} \sum_{i=1}^M A_{\min}(X_i)$ is the **average number of nearest neighbors**.

- Requires only **knowledge of \bar{A}_{\min} and $d_{E,\min}$** !

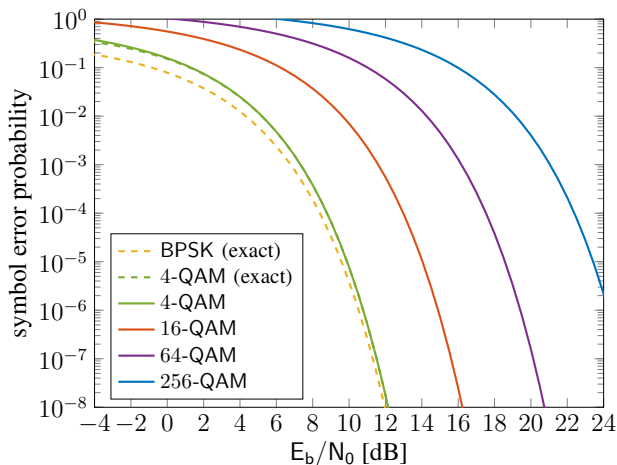
Nearest Neighbor Approximation

Nearest Neighbor Approximation for squared M -QAM

For squared M -QAM constellations, there are 4 constellation points with 2 neighbors, $4(\sqrt{M} - 2)$ points with 3 neighbors and the remaining points have 4 neighbors, hence $\bar{A}_{\min} = 4 - 4/\sqrt{M}$. Furthermore, $d_{E,\min} = \sqrt{\frac{6E_s}{M-1}}$. Hence,

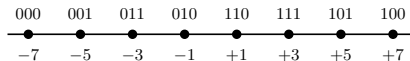
$$\begin{aligned} P_s^{MQAM} &\approx \left(4 - \frac{4}{\sqrt{M}}\right) Q\left(\sqrt{\frac{3E_s}{(M-1)N_0}}\right) \\ &= \left(4 - \frac{4}{\sqrt{M}}\right) Q\left(\sqrt{\frac{3E_b \log M}{(M-1)N_0}}\right). \end{aligned}$$

Nearest Neighbor Approximation of P_s of QAM



- The **power efficiency** decreases with M .
- The **spectral efficiency** increases with M .

Bit Error Probability of Linear Modulations



- The bit error probability depends on the **binary labeling**, i.e., **how tuples of $m = \log M$ bits are mapped to the constellation symbols**.
- Let

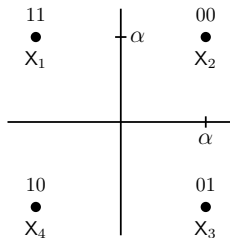
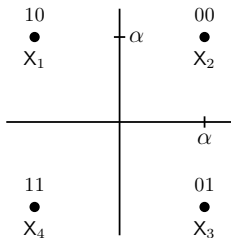
$$\mathbf{L}(x) = (b_1(x), \dots, b_m(x))$$

be the m -bit labeling associated to constellation symbol $x \in \mathcal{X}$.

- Example: $\mathbf{L}(+1) = (1, 1, 0)$, $\mathbf{L}(+5) = (1, 0, 1)$.
- The bit error probability P_b of an M -ary constellation is given by

$$P_b = \frac{1}{m} \sum_{i=1}^m \Pr(i\text{th bit in error}) = \frac{1}{m} \sum_{i=1}^m p_i.$$

Bit Error Probability of 4-QAM with Gray and Lexicographic Labeling



Gray Labeling

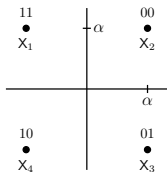
- e_{b_i} : event that bit i is decoded in error.

$$\begin{aligned} p_1 &= \Pr(e_{b_1}) = \sum_{x \in \mathcal{X}} \Pr(e_{b_1}|x)P(x) = \frac{1}{4} \sum_{x \in \mathcal{X}} \Pr(e_{b_1}|x) \\ &= \frac{1}{4}(\Pr(e_{b_1}|X_1) + \Pr(e_{b_1}|X_2) + \Pr(e_{b_1}|X_3) + \Pr(e_{b_1}|X_4)) \\ &= \frac{1}{4}(\Pr(b_1(\hat{x}) = 0|X_1) + \Pr(b_1(\hat{x}) = 1|X_2) \\ &\quad + \Pr(b_1(\hat{x}) = 1|X_3) + \Pr(b_1(\hat{x}) = 0|X_4)), \end{aligned}$$

- All terms are $Q\left(\sqrt{\frac{2E_b}{N_0}}\right)$, thus $p_1 = Q\left(\sqrt{\frac{2E_b}{N_0}}\right)$.
- By symmetry, $p_2 = p_1$, hence,

$$P_b^{4\text{QAM-Gray}} = \frac{1}{2}(p_1 + p_2) = Q\left(\sqrt{\frac{2E_b}{N_0}}\right).$$

Lexicographic Labeling



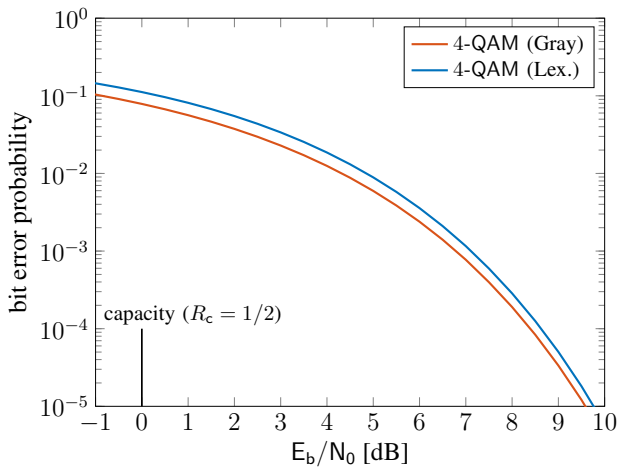
$$p_1 = Q\left(\sqrt{\frac{2E_b}{N_0}}\right)$$
$$p_2 = 2Q\left(\sqrt{\frac{2E_b}{N_0}}\right)\left(1 - Q\left(\sqrt{\frac{2E_b}{N_0}}\right)\right).$$

Thus,

$$P_b^{4\text{QAM-Lex}} = \frac{1}{2}(p_1 + p_2) = \frac{3}{2}Q\left(\sqrt{\frac{2E_b}{N_0}}\right) - \left(Q\left(\sqrt{\frac{2E_b}{N_0}}\right)\right)^2.$$

- Due to the **lack of symmetry**, $p_1 \neq p_2$.

Bit Error Probability for 4-QAM



Bit Error Probability of M -ary Constellations

- For general M -ary constellations and arbitrary labelings, the computation of P_b is cumbersome \rightarrow **Nearest neighbor approximation**.