Digital Communications SSY125, Lecture 14

Coding beyond reliable communication

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December 4, 2019

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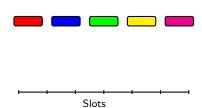
- A very powerful tool for many applications and research fields:
 - Uncoordinated multiple access
 - Post-quantum cryptography
 - DNA storage
 - Caching
 - Distributed computing

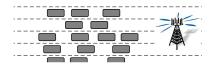


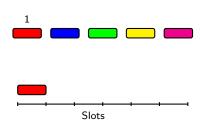


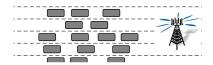


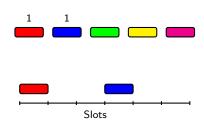


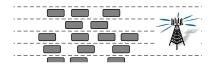


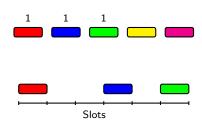


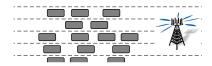


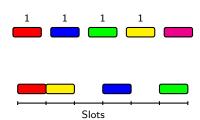


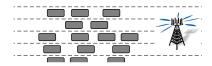


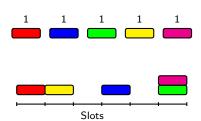


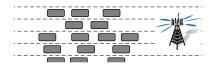














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- Decoding using successive interference cancelation, exploiting the introduced redundancy

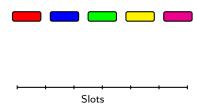


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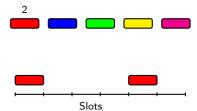


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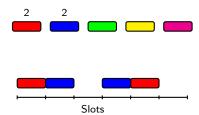


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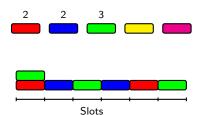


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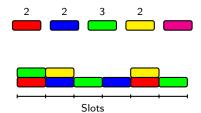


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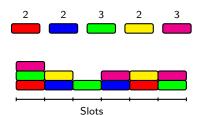


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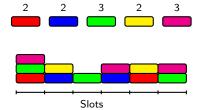


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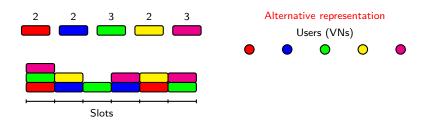
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Alternative representation

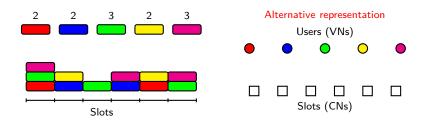


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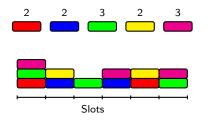


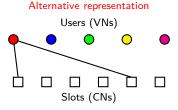
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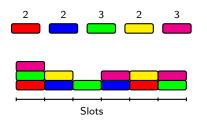
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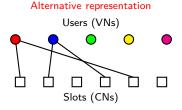






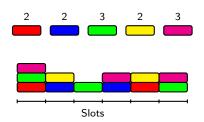
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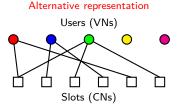






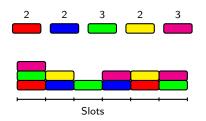
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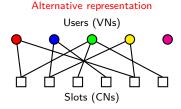






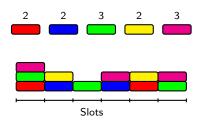
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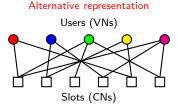






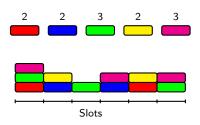
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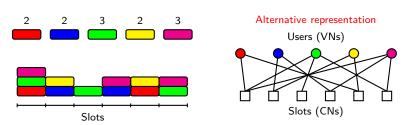
Alternative representation
Users (VNs)

Slots (CNs)

A coding problem!

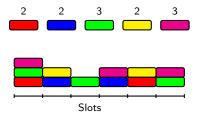


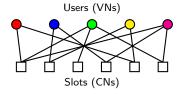
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ullet A coding problem! o Can be analyzed using the theory of codes on graphs

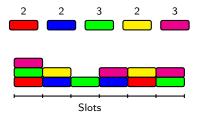
Coded slotted ALOHA: Decoding

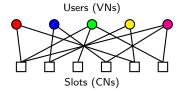




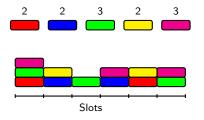
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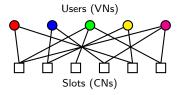
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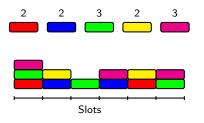


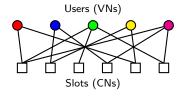
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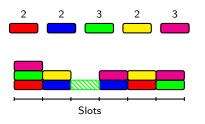


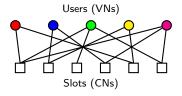
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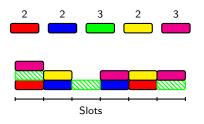


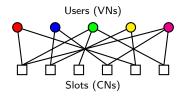
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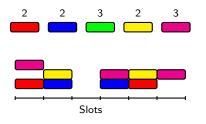


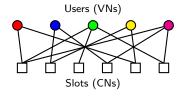
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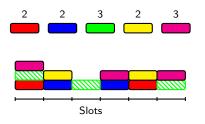


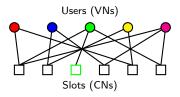
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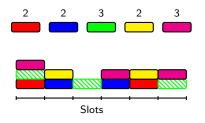


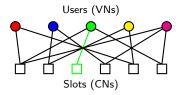
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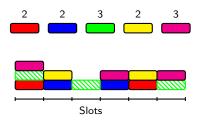


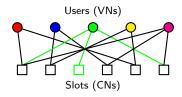
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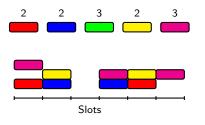


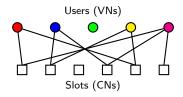
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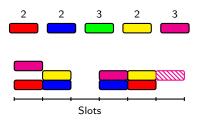


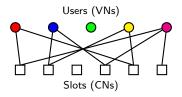
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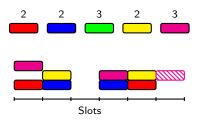


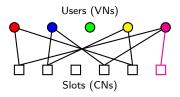
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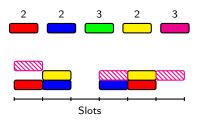


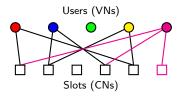
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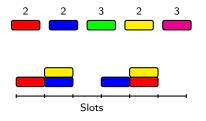


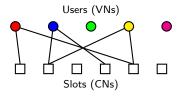
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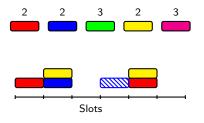


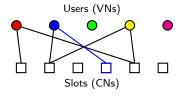
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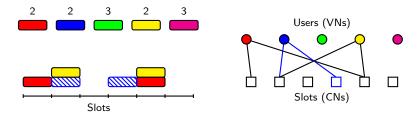


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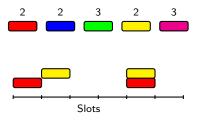


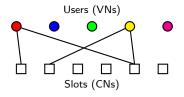


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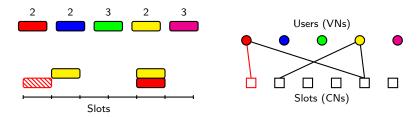


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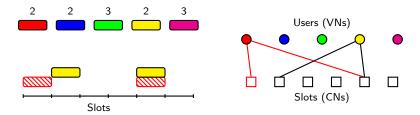




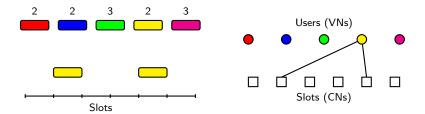
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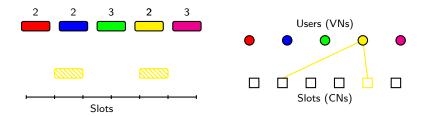
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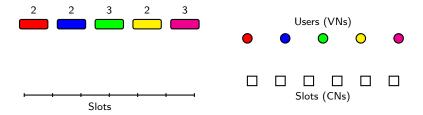
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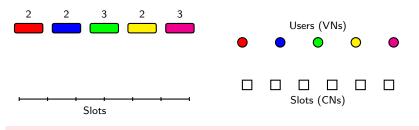
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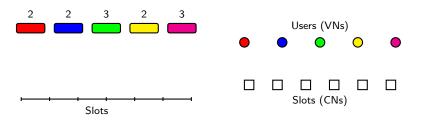


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- Decoding of CSA is equivalent to peeling (BP) decoding on a bipartite graph!
- It can be analyzed using the theory of codes-on-graphs!

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- Compute d that satisfies $de \equiv 1 \pmod{\phi(n)}$.
- (n, e) is the public key, d is the private key.

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- But...a quantum computer can factor n in polynomial time!
- We need to find cryptosystems that are robust against the attack by a quantum computer

Encrypting/Decrypting the message and attack

- Alice turns the message into a number u < n and encrypts u to $c = u^e \pmod n$.
- By using d, Bob can recover the message u, $c^d = (u^e)^d \equiv u \pmod{n}$.
- To recover d, an attacker needs to factor $n \to \text{Very time consuming!}$
- But...a quantum computer can factor n in polynomial time!
- We need to find cryptosystems that are robust against the attack by a quantum computer → Post-quantum cryptography

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- G' is the public key, and (G, S, P) is the secret key.
- The new matrix G' is the generator matrix of another linear code, that is very difficult to decode if the secret key is not known.

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• Bob decodes the codeword c' to obtain uS, and finally computes $u=uS^{-1}$.

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Problem

Very big key sizes (hundreds of thousands of bits, compared to few thousands for RSA)

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 - Efficient decoding.
 - Much shorter key sizes.

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- Archival DNA-based storage already feasible!

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- Once desired information is stored in DNA, it may be rewritten using DNA editing techniques.

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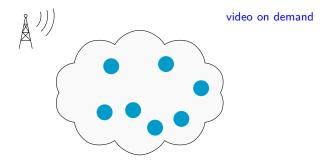
DNA storage

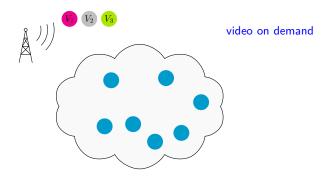
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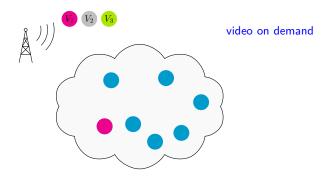
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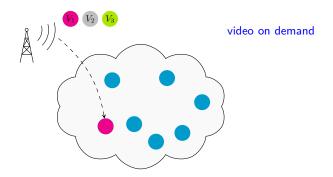
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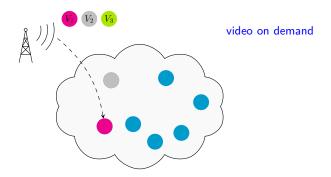
- Coding to correct insertions, deletions and substitutions.
- Constrained coding to avoid some DNA patterns.

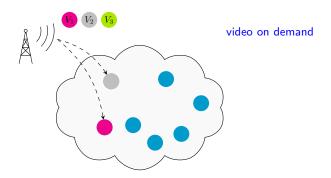


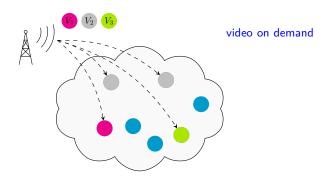


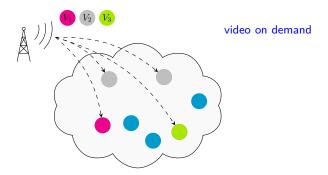




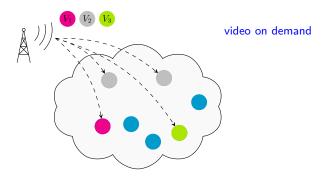




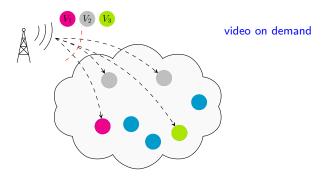




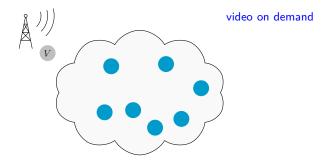
Mobile data traffic is expected to rise at a rate of 45% per year

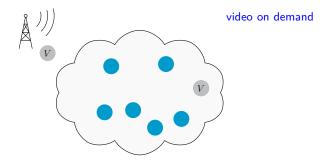


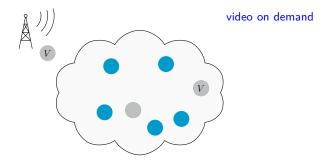
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- It is forecasted to attain 51 exabytes per month by 2021 [Ericsson]

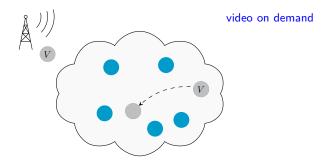


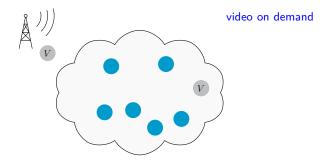
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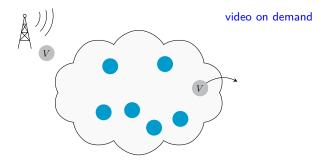


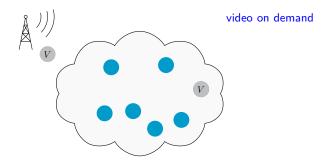


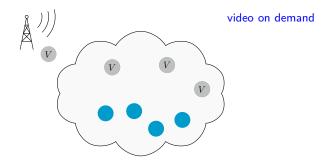


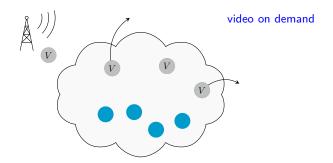


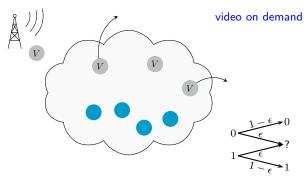




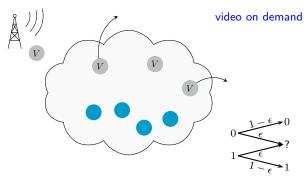




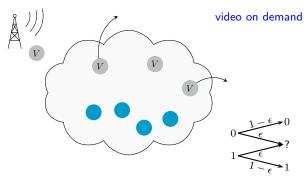




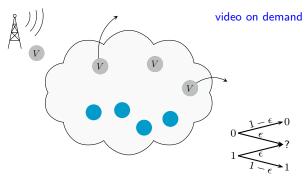
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Idea

Use erasure correcting codes to store content!

And many more applications...

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- Distributed computing
- Distributed storage
- Security and privacy
- Quantum key distribution
- Quantum error correction
- ..