



# Cohesive Subgraph Detection

## Clique Model

Never Stand Still

Faculty of Engineering

Computer Science and Engineering

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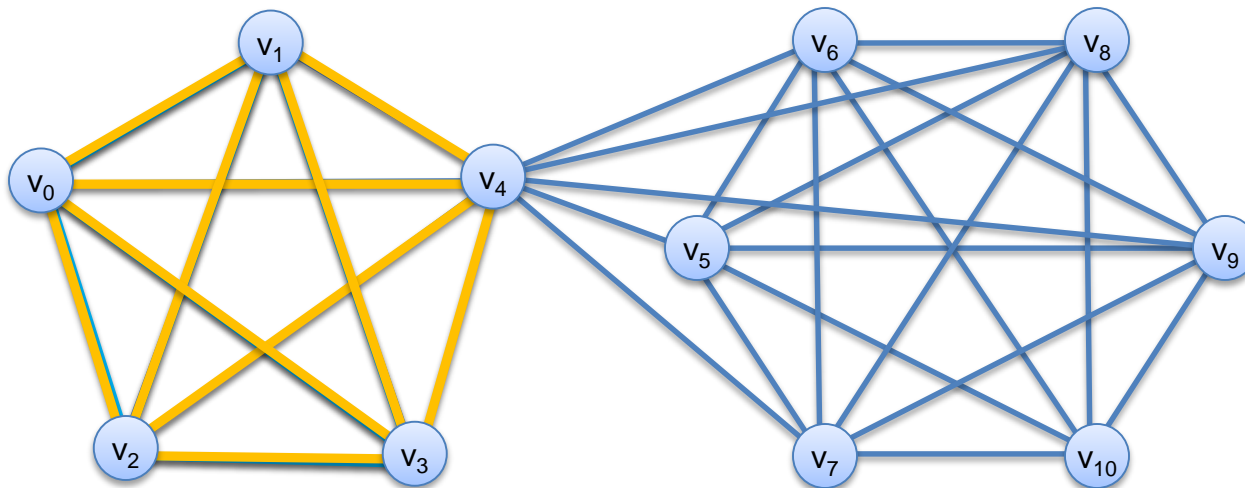
The University of New South Wales, Australia

# Outline

- Clique Model
- K-Core vs K-Truss
- K-Edge Connected
- K-Vertex Connected

# Clique

- Given a graph  $G$ , a clique is a set of nodes such that for any pair of them have an edge
- A clique is called maximal clique if there exist no other bigger cliques that contain it

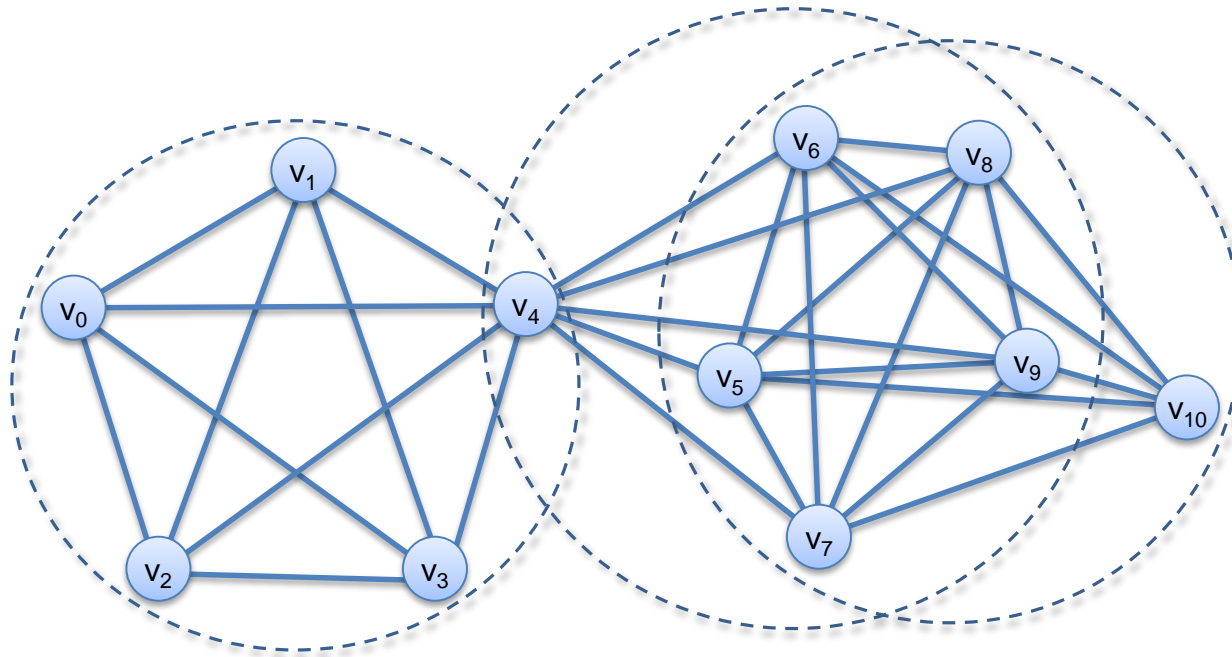


# Application

- Community detection
- Gene expression and motif discovery
- Anomaly detection
- Stock market data visualization
- Signal transmission analysis
- .....

# Maximal Clique Enumeration

- Given a graph  $G$ , find all the maximal cliques in  $G$ .
  - NP-Hard Problem



# In-Memory MCE

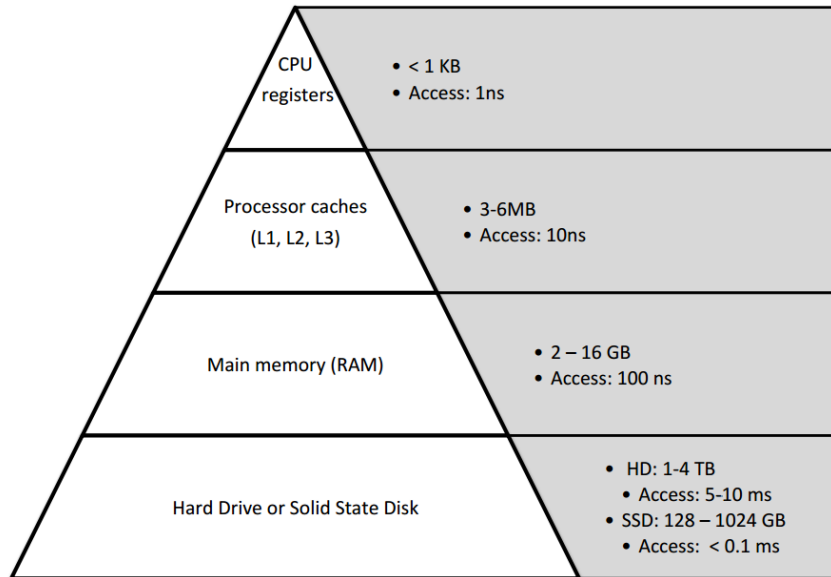
- Bron-Kerbosch Algorithm
  - First Practical In-Memory MCE Algorithm
    - In-Memory means all the input and auxiliary data structure can be loaded in main memory during the computation
    - C. Bron et al., “Algorithm 457: finding all cliques of an undirected graph”, *Communications of the ACM*, **16** (9): 575–577, 1973
  - Based on a recursive backtracking paradigm

# I/O Efficient MCE

- Why I/O Efficient?
  - Real graph is massive
    - Facebook contains 1.32 billion nodes and 140 billion edges
    - EU-2015 (sub-domain of web graph) contains 1.07 billion nodes and 91 billion edges
  - Memory is fast but small while disk is slow but large

# I/O Efficient MCE

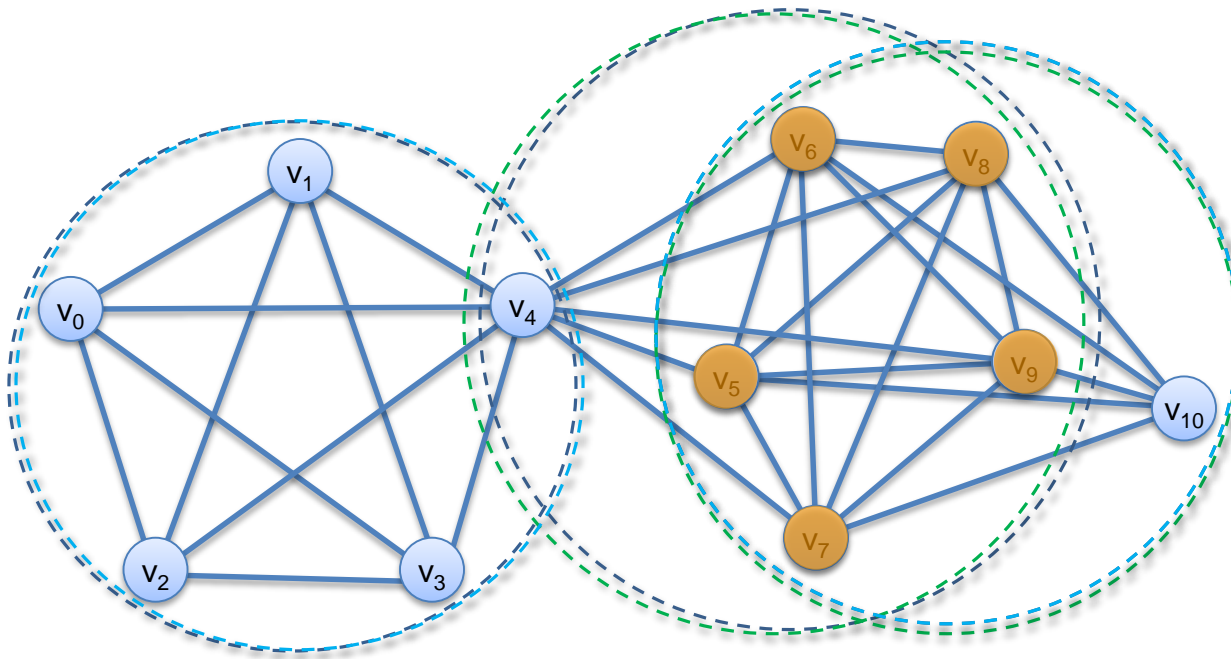
- Why I/O Efficient?
  - Memory Hierarchy





# Diversified Top-K Clique Search

- Traditional models vs diversified top-k clique



# Diversified Top-K Clique Search

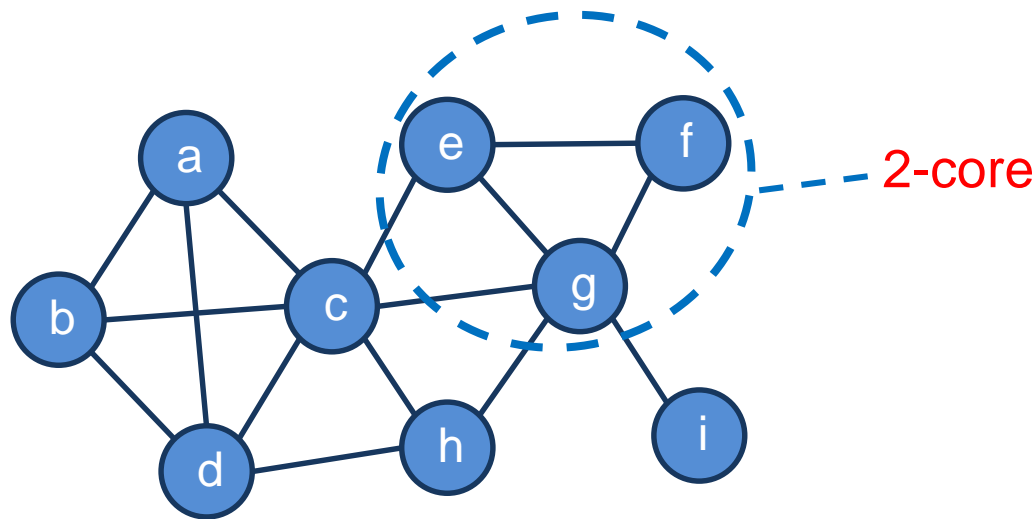
- Our Solution
  - treat it as an online k coverage problem
  - store k maximal cliques in memory
  - update these k candidate maximal cliques while enumerating cliques
    - replace small existing cliques with big new cliques

# Diversified Top-K Clique Search

- PNP-Index
  - An naïve implementation for candidate set maintenance needs  $O(|\mathcal{A}| * k * |C_{max}|)$
  - With the help of PNP-Index, our algorithm can only take  $O(\sum_{C \in \mathcal{A}} |C|)$  time

# $k$ -Core

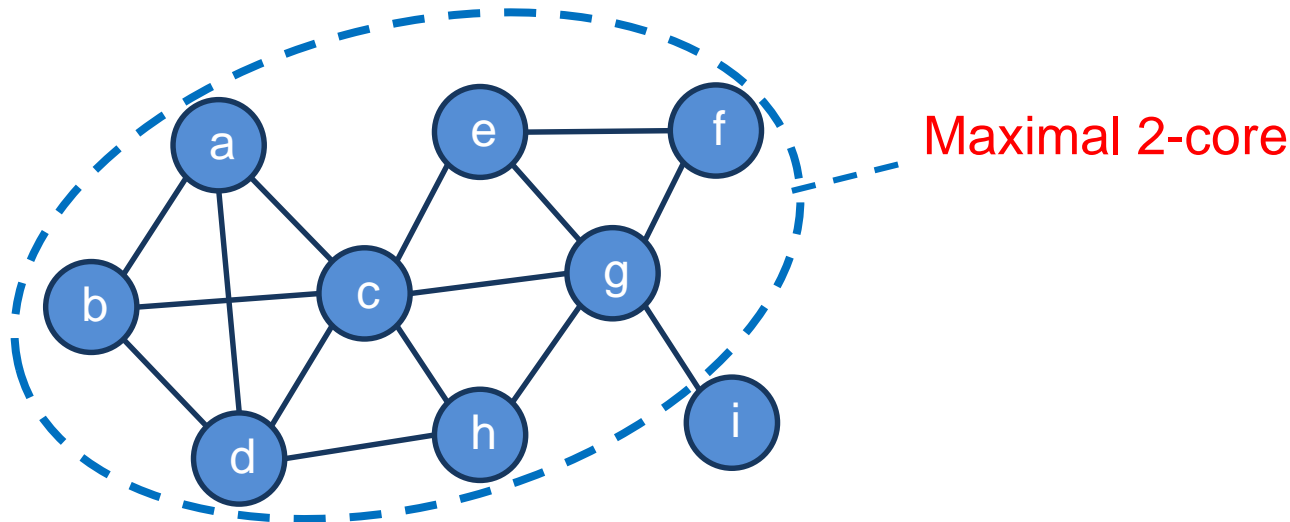
- Given a graph  $G$ , the  $k$ -core of  $G$  is a subgraph where each node has at least  $k$  neighbors (i.e.,  $k$  adjacent nodes, or a degree of  $k$ ).



S. B. Seidman. Network structure and minimum degree. *Social networks*, 5(3):269–287, 1983.

# Maximal $k$ -Core

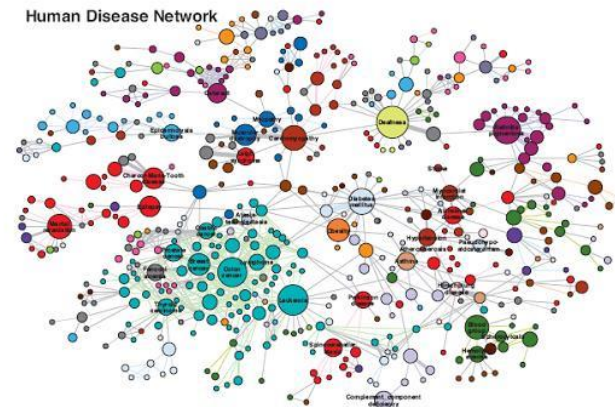
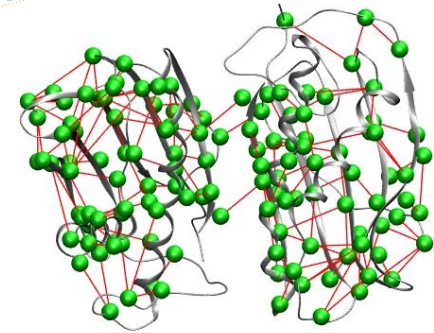
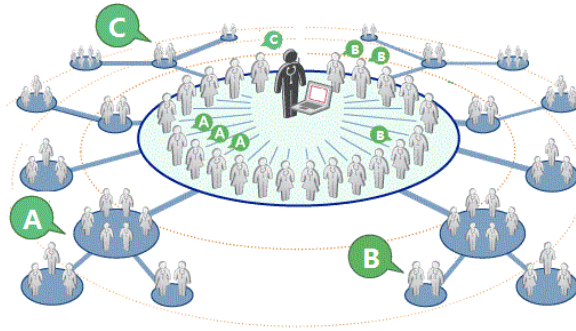
- A  $k$ -core  $C$  is called maximal if any supergraph of  $C$  is not a  $k$ -core (i.e., no another  $k$ -core which contains  $C$ ).



S. B. Seidman. Network structure and minimum degree. Social networks, 5(3):269–287, 1983.

# Applications

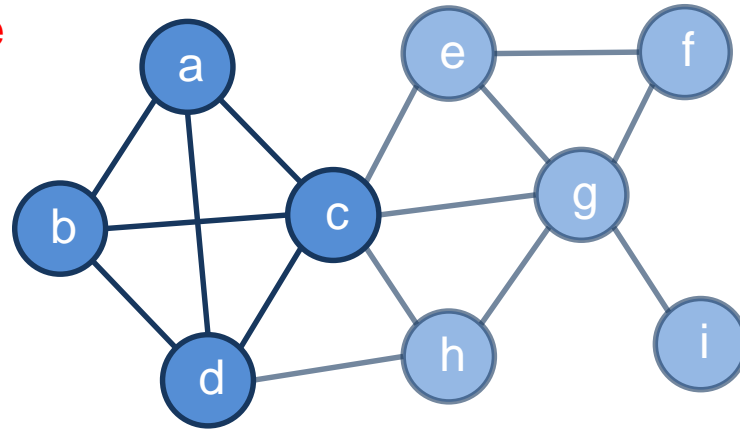
- Community detection
- Social contagion
- User engagement
- Event detection
- Network analysis and visualization
- Influence study
- Graph clustering
- Protein function prediction
- Human Cerebral Cortex
- .....



# Compute Maximal $k$ -Core

- Given a graph  $G$ , the maximal  $k$ -core of  $G$  can be computed by recursively deleting every node and its adjacent edges if its degree is less than  $k$ .

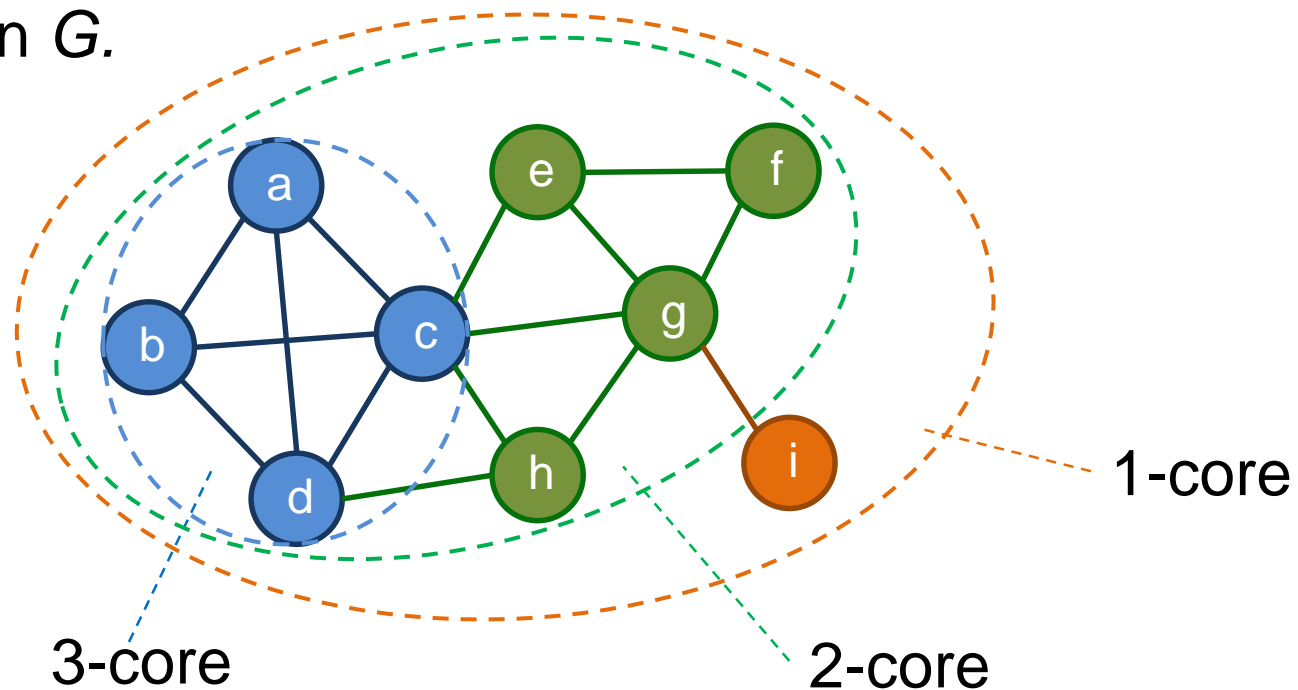
Maximal 3-core



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# $k$ -Core Decomposition

- **Core number** of a node  $v$ : the largest value of  $k$  such that there is a  $k$ -core containing  $v$ .
- Core decomposition: compute the **core number** of each node in  $G$ .





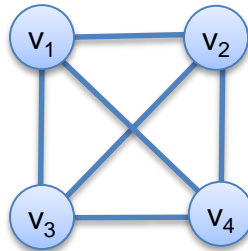
# ***k*-Truss**

- Given a graph  $G$ , the  $k$ -truss of  $G$  is a subgraph where edge is at least involved in  $(k-2)$  triangles.
- $k$ -truss is an enhancement of  $k$ -core; each vertex of  $k$ -truss has a degree at least  $k-1$ .

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# k-edge Connectivity

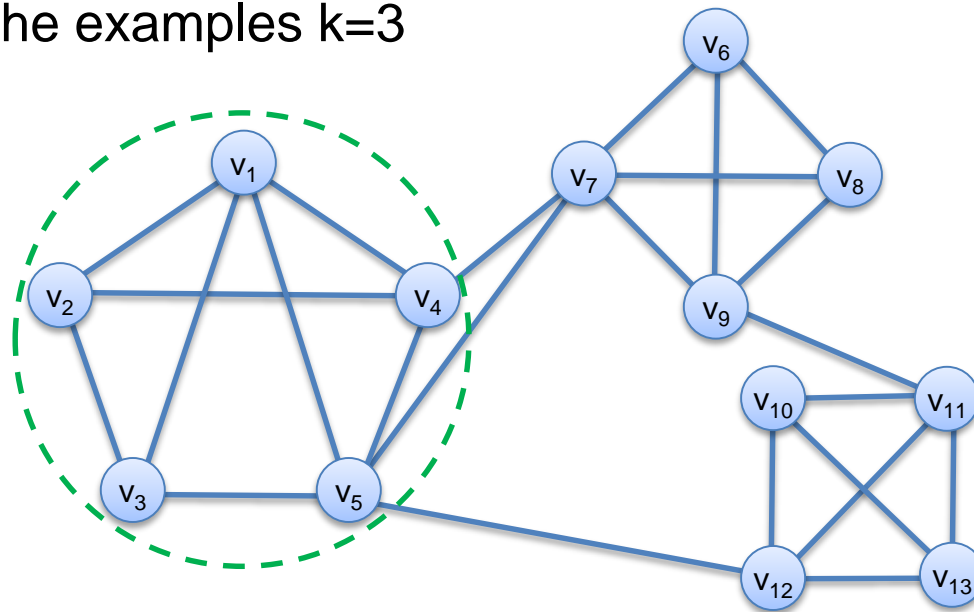
- A graph is  $k$ -edge connected if it is still connected after removing any set of  $(k-1)$  edges from it



# k-edge Connected Component

- A k-edge connected component (k-ECC) of a graph  $G$  is a maximal subgraph  $g$  of  $G$  such that  $g$  is k-edge connected

➤ All the examples  $k=3$



# Application

- Community detection
- Social behaviour mining
- Graph visualization
- Steiner Component Search
- Hierarchy Study in Networks
- .....

# k-Vertex Connectivity

- A graph is  $k$ -vertex connected if it is still connected after removing any set of  $(k-1)$  vertex from it

