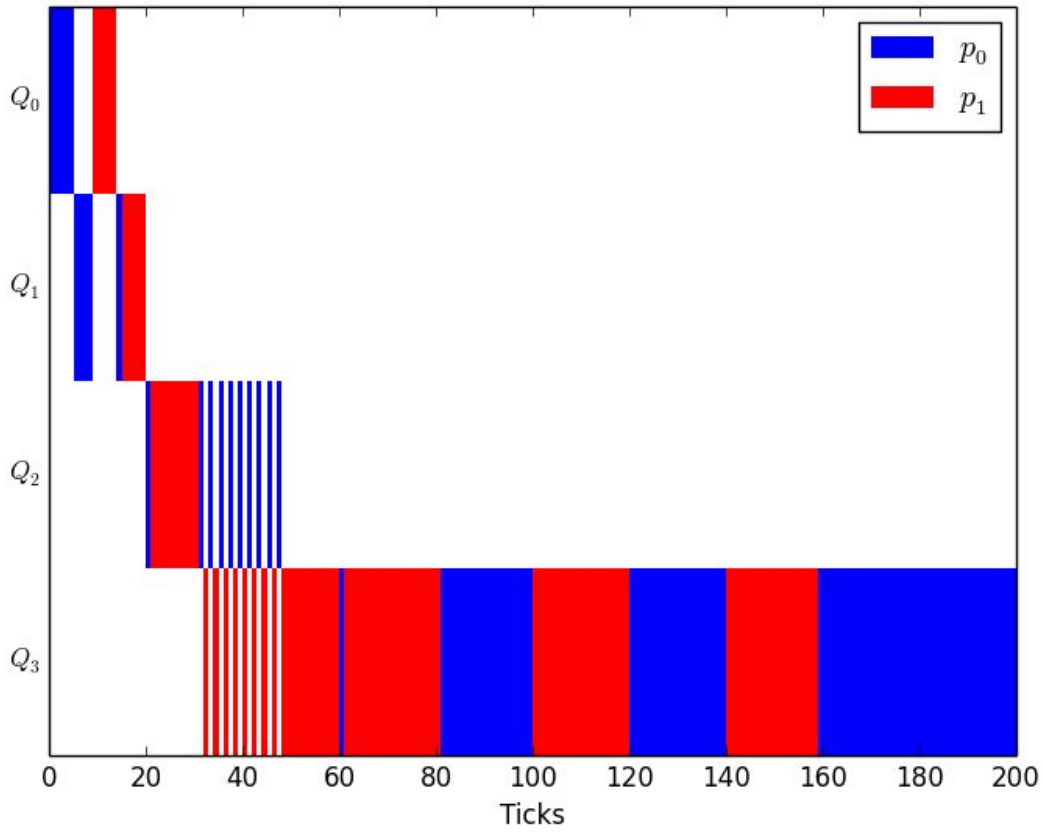

- Workload: Before the time slice of p0 is expired, p0 relinquishes CPU (at time ticks 9 by calling `sleep()`).
- Observations:
    - **Rule 3**: Whenever p1 goes sleep, p2, as the highest priority level process gets scheduled.
    - **Rule 4**: The frequent context switch at around timeticks 30 is due to the mechanism of sleep() syscall, which wakes up the sleeping process every timeticks, and check if it is ready to run. The process is considered to have used up the entire time tick even it voluntarily relinquishes the CPU. In this graph, p1 continues to consume its allowance in each priority level even when it is shortly waked up and goes sleep immediately.